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**A PMIPv6-based solution for Distributed Mobility Management
draft-bernardos-mext-dmm-pmip-01**

Abstract

The number of mobile users and their traffic demand is expected to be ever-increasing in future years, and this growth can represent a limitation for deploying current mobility management schemes that are intrinsically centralized, e.g., Mobile IPv6 and Proxy MIPv6. For this reason it has been waded a need for distributed and dynamic mobility management approaches, with the objective of reducing operators' burdens, evolving to a cheaper and more efficient architecture.

This draft describes a solution to distribute the data forwarding plane on Proxy Mobile IPv6 domains, thus trying to overcome the suboptimal data path introduced when the LMA is traversed.

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1. Introduction

Current IP mobility solutions, standardized with the names of Mobile IPv6 [[RFC3775](#)], or Proxy Mobile IPv6 [[RFC5213](#)], just to cite the two most relevant examples, offer mobility support at the cost of handling operations at a cardinal point, the mobility anchor, and burdening it with data forwarding and control mechanisms for a great amount of users. As stated in [[I-D.chan-distributed-mobility-ps](#)], centralized mobility solutions are prone to several problems and limitations: longer (sub-optimal) routing paths, scalability problems, signaling overhead (and most likely a longer associated handover latency), more complex network deployment, higher vulnerability due to the existence of a potential single point of failure, and lack of granularity on the mobility management service (i.e., mobility is offered on a per-node basis, not being possible to define finer granularity policies, as for example per-application).

There are basically two main approaches being researched now: one aimed at making Mobile IPv6 work in a distributed way (a complete solution can be found in [[I-D.bernardos-mext-dmm-cmip](#)]), and another one doing the same exercise for Proxy Mobile IPv6. In this draft we describe a solution to achieve a DMM behavior for network-based mobility support (i.e., inspired by PMIPv6). This document is based on a research paper of the same authors, currently under submission, called "A Network-based Localized Mobility Solution for Distributed Mobility Management" [[Net-basedDMM](#)].

2. Terminology

The following terms used in this document are defined in the Proxy Mobile IPv6 specification [[RFC5213](#)]:

Local Mobility Anchor (LMA)

Mobile Access Gateway (MAG)

Mobile Node (MN)

Binding Cache Entry (BCE)

Proxy Care-of Address (P-CoA)

Proxy Binding Update (PBU)

Proxy Binding Acknowledgement (PBA)

The following terms are defined and used in this document:

MAAR (Mobility Anchor and Access Router). First hop routers where the mobile nodes attach to. They also play the role of mobility managers for the IPv6 prefixes they anchor.

CMD (Central Mobility Database). Node that stores the BCEs for the MNs in the mobility domain.

A-MAAR (Anchor MAAR). MAAR which was previously visited by the MN and is still involved in an active flow using an IPv6 prefix it has advertised to the MN (i.e., MAAR where that IPv6 prefix is anchored).

S-MAAR (Serving MAAR). MAAR which the MN is currently attached to.

3. Description of the solution

The purpose of Distributed Mobility Management approaches is to overcome the limitations of the traditional centralized mobility management by bringing the mobility anchor closer to the MN. Following this idea, in our proposal, the central anchor is moved to the edge of the network, being deployed in the default gateway of the mobile node. That is, the first elements that provide IP connectivity to a set of MNs are also the mobility managers for those MNs. In the following, we will call these Mobility Anchor and Access Routers (MAARs).

However, MAARs leverage on the Central Mobility Database (CMD) to access and update information related to the MNs, stored as mobility sessions; hence, a centralized node maintains a global view on the status of the network. The CMD is queried whenever a MN is detected to join the mobility domain. It might be a fresh attachment or a handover, but as MAARs do not store any mobility session, they contact the CMD to retrieve the data of interest and eventually take the appropriate action. The procedure adopted for the query and the messages exchange sequence might vary to optimize the update latency and/or the signaling overhead. Here is presented one method for the initial registration, and three different approaches to update the mobility sessions using PBUs and PBAs. Each approach assigns a different role to the CMD:

- o The CMD is a PBU/PBA relay
- o The CMD is only a MAAR locator
- o PBU/PBA is a PBU/PBA proxy

3.1. Initial registration

Upon the MN's attachment to a MAAR, say MAAR1, an IPv6 global prefix belonging to the MAAR's prefix pool is reserved for it (Pref1). The prefix is sent in a PBU with the MN's Identifier (MN-ID) to the CMD, which, since the session is new, stores a Binding Cache Entry containing as main fields the MN-ID, the MN's prefix and MAAR1's address as Proxy-CoA. The CMD replies to MAAR1 with a PBA indicating that the MN's registration is fresh and no past status is available. MAAR1 sends a Router Advertisement (RA) in unicast to the MN including the prefix reserved before, that can be used by the MN to configure an IPv6 address (e.g., with stateless auto-configuration). The address is routable at the MAAR, in the sense that it is on the path of packets addressed to the MN. Moreover, the MAAR acts as plain router for those packets, as no encapsulation nor special handling takes place. Figure 1 illustrates this scenario.

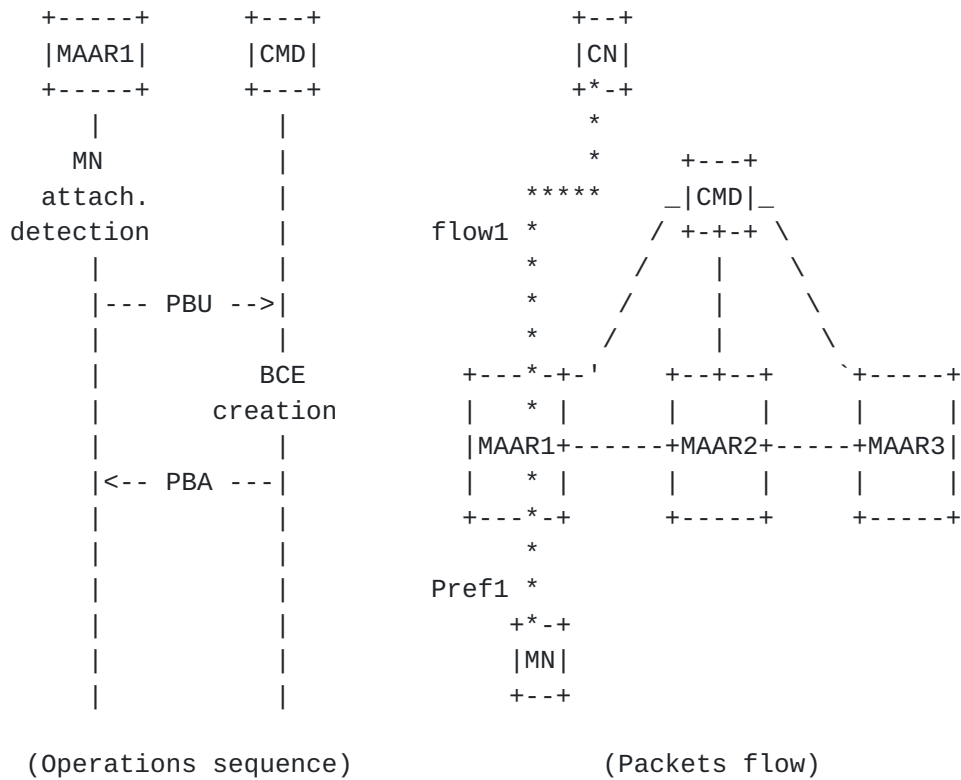


Figure 1: First attachment to the network

3.2. The CMD as PBU/PBA relay

When the MN moves from its current access, it associates to MAAR2 (now the S-MAAR), which delegates another IPv6 prefix (Pref2) and sends it to the CMD for registration. The CMD has already an entry for the MN, binding the MN-ID to its former location; thus, it

forwards the PBU to the MAAR indicated as Proxy CoA, in this case MAAR1 (now the A-MAAR), and it updates the P-CoA field with the S-MAAR's address. Upon PBU reception, MAAR1 replies to the CMD with a PBA to ensure that the new location has successfully changed, containing the prefix anchored at MAAR1. The CMD updates the BCE adding the P-MAAR address in the list of old P-CoAs and forwards the PBA to the new S-MAAR, containing the previous Proxy-CoA and the prefix anchored to it, so that a tunnel can be established between the two MAARs and new routes are set appropriately to recover the flow(s).

Now packets destined to Pref1 are first received by MAAR1, encapsulated into the tunnel and forwarded to MAAR2, which finally delivers them to their destination. In uplink, when the MN transmits packets using Pref1 for the source address, they are sent to MAAR2, as it is MN's new default gateway, then tunneled to MAAR1 which routes them towards the next hop to destination. Conversely, packets carrying Pref2 are routed by MAAR2 without any special packet handling both for uplink and downlink. The procedure is depicted in Figure 2.

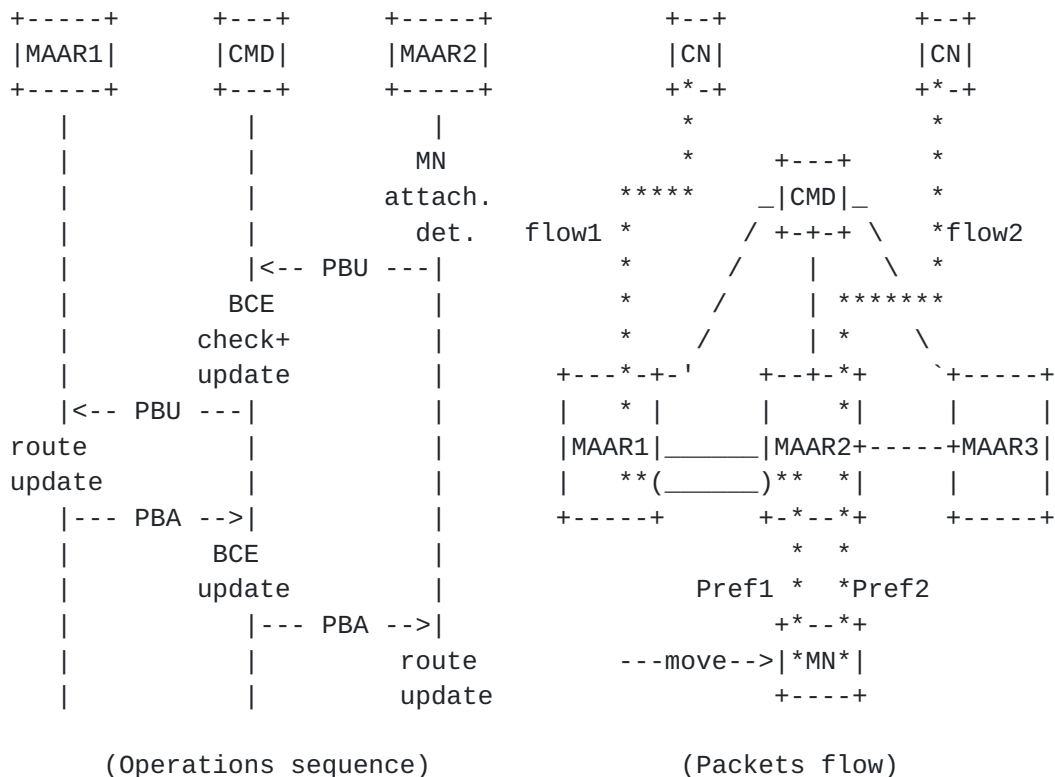


Figure 2: Scenario after a handover, CMD as relay

For next MN's movements the process is repeated except for the number of P-MAARs involved, that rises accordingly to the number of prefixes

that the MN wishes to maintain. Indeed, once the CMD receives the first PBU from the new S-MAAR, it forwards copies of the PBU to all the A-MAARs indicated in the BCE as current P-CoA (i.e., the MAAR prior to handover) and old P-CoAs. They reply with a PBA to the CMD, which aggregates them into a single one to notify the S-MAAR, that finally can establish the tunnels with the A-MAARs. It should be noted that this design separates the mobility management at the prefix granularity, and it can be tuned in order to erase old mobility sessions when not required, while the MN is reachable through the latest prefix acquired. Moreover, the latency associated to the mobility update is bound to the PBA sent by the furthest A-MAAR, that takes the longest time to reach the CMD. The drawback can be mitigated introducing a timeout at the CMD, by which, after its expiration, all the PBAs so far collected are transmitted, and the remaining are sent later upon their arrival.

3.3. The CMD as MAAR locator

The latency experienced in the approach shown before can be mitigated if the A-MAARs are allowed to signal directly their information to the new S-MAAR. This procedure reflect what was described in [Section 3.2](#) up to the moment the A-MAAR receives the PBU. At that point an A-MAAR is aware of the new MN's location (i.e., S-MAAR) and, besides sending a PBA to the CMD, it also sends a PBA to the S-MAAR including the prefix it is anchoring. The CMD is relieved from forwarding the PBA to the S-MAAR, as it receives a copy directly from the A-MAAR with the necessary information to build the tunnel and set the appropriate routes. In Figure 3 is illustrated the new messages sequence, while the data forwarding is unaltered.

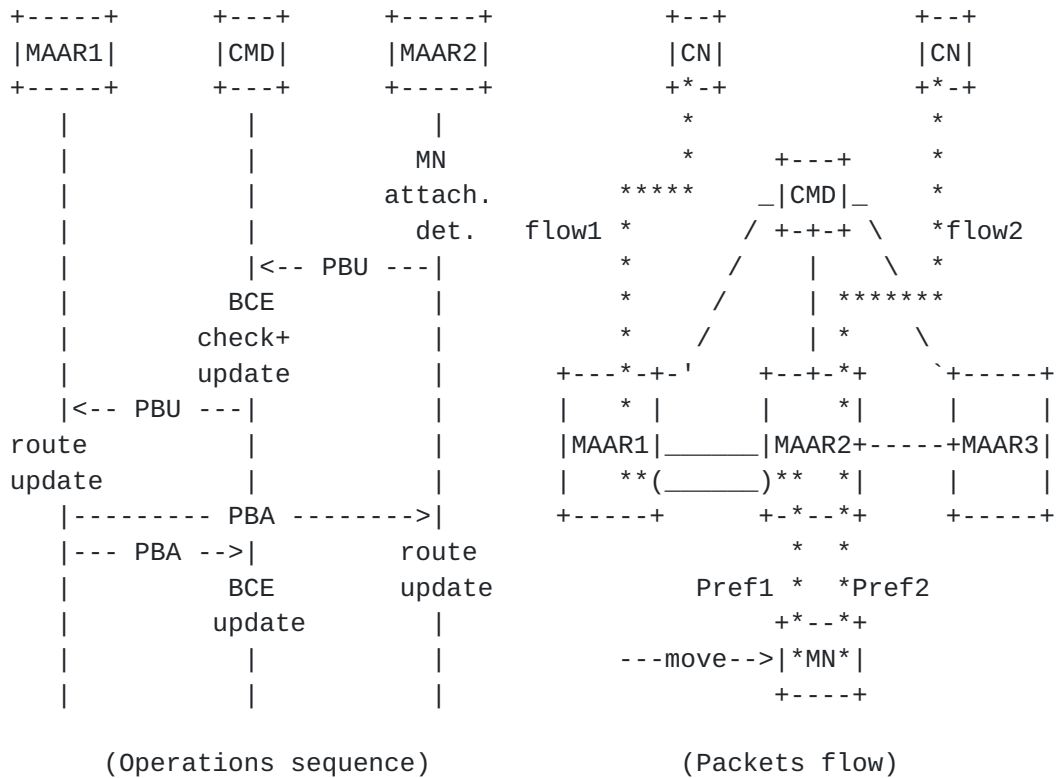


Figure 3: Scenario after a handover, CMD as locator

3.4. The CMD as PBU/PBA proxy

A further enhancement of previous solutions can be achieved when the CMD sends the PBA to the new S-MAAR before notifying the A-MAARs of the location change. Indeed, when the CMD receives the PBU for the new registration, it is already in possess of all the information that the new S-MAAR requires to set up the tunnel and the routes. Thus the PBA is sent to the S-MAAR immediately after a PBU is received. In parallel, a PBU is sent by the CMD to the A-MAARs to notify them about the new MN's location, so they receive the information to establish the tunnel and routes on their side. When A-MAARs complete the update, they send a PBA to the CMD to indicate that the operation is concluded and the information are updated in all network nodes. This scheme is depicted in Figure 4, where, again, the data forwarding is kept untouched.

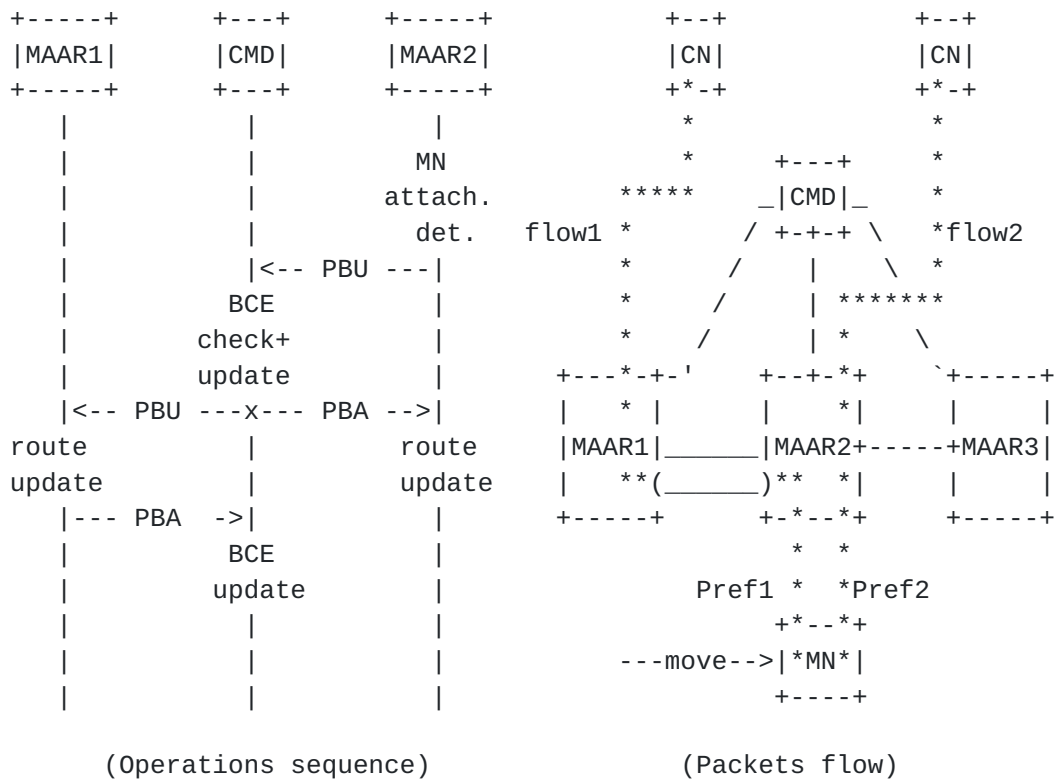


Figure 4: Scenario after a handover, CMD as proxy

4. IANA Considerations

TBD.

5. Security Considerations

TBD.

6. Acknowledgments

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