

6LoWPAN Generic Compression of Headers and Header-like Payloads  
[draft-bormann-6lowpan-ghc-06](#)

## Abstract

This short I-D provides a simple addition to 6LoWPAN Header Compression that enables the compression of generic headers and header-like payloads, without a need to define a new header compression scheme for each new such header or header-like payload.

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## [1.](#) Introduction

### [1.1.](#) The Header Compression Coupling Problem

6LoWPAN-HC [[RFC6282](#)] defines a scheme for header compression in 6LoWPAN [[RFC4944](#)] packets. As with most header compression schemes, a new specification is needed for every new kind of header that needs to be compressed. In addition, [[RFC6282](#)] does not define an extensibility scheme like the ROHC profiles defined in ROHC [[RFC3095](#)] [[RFC5795](#)]. This leads to the difficult situation that 6LoWPAN-HC tended to be reopened and reexamined each time a new header receives consideration (or an old header is changed and reconsidered) in the 6LoWPAN/roll/CoRE cluster of IETF working groups. While [[RFC6282](#)] finally got completed, the underlying problem remains unsolved.

The purpose of the present contribution is to plug into [[RFC6282](#)] as is, using its NHC (next header compression) concept. We add a slightly less efficient, but vastly more general form of compression for headers of any kind and even for header-like payloads such as those exhibited by routing protocols, DHCP, etc. The objective is an extremely simple specification that can be defined on a single page and implemented in a small number of lines of code, as opposed to a general data compression scheme such as that defined in [[RFC1951](#)].

### [1.2.](#) Terminology



The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC 2119](#) [[RFC2119](#)].

The term "byte" is used in its now customary sense as a synonym for "octet".

### [1.3.](#) Notation

This specification uses a trivial notation for code bytes and the bitfields in them the meaning of which should be mostly obvious. More formally speaking, the meaning of the notation is:

Potential values for the code bytes themselves are expressed by templates that represent 8-bit most-significant-bit-first binary numbers (without any special prefix), where 0 stands for 0, 1 for 1, and variable segments in these code byte templates are indicated by sequences of the same letter such as kkkkkkk or ssss, the length of which indicates the length of the variable segment in bits.

In the notation of values derived from the code bytes, 0b is used as a prefix for expressing binary numbers in most-significant-bit first notation (akin to the use of 0x for most-significant-digit-first hexadecimal numbers in the C programming language). Where the abovementioned sequences of letters are then referenced in such a binary number in the text, the intention is that the value from these bitfields in the actual code byte be inserted.

Example: The code byte template

101nssss

stands for a byte that starts (most-significant-bit-first) with the bits 1, 0, and 1, and continues with five variable bits, the first of which is referenced as "n" and the next four are referenced as "ssss". Based on this code byte template, a reference to

0b0ssss000

means a binary number composed from a zero bit, the four bits that are in the "ssss" field (for 101nssss, the four least significant bits) in the actual byte encountered, kept in the same order, and three more zero bits.



## 2. 6LoWPAN-GHC

The format of a GHC-compressed header or payload is a simple bytecode. A compressed header consists of a sequence of pieces, each of which begins with a code byte, which may be followed by zero or more bytes as its argument. Some code bytes cause bytes to be laid out in the destination buffer, some simply modify some decompression variables.

At the start of decompressing a header or payload within a L2 packet (= fragment), variables "sa" and "na" are initialized as zero.

The code bytes are defined as follows:

code byte	Action	Argument
0kkkkkkk	Append k = 0b0kkkkkkk bytes of data in the bytecode argument (k < 96)	k bytes of data
1000nnnn	Append 0b0000nnnn+2 bytes of zeroes	
10010000	STOP code (end of compressed data, see <a href="#">Section 3.2</a> )	
101nssss	Set up extended arguments for a backreference: sa += 0b0ssss000, na += 0b0000n000	
11nnnkkk	Backreference: n = na+0b00000nnn+2; s = 0b00000kkk+sa+n; append n bytes from previously output bytes, starting s bytes to the left of the current output pointer; set sa = 0, na = 0	

Note that the following bit combinations are reserved at this time: 011xxxxx, and 1001nnnn (where 0b0000nnnn > 0).

For the purposes of the backreferences, the expansion buffer is initialized with a predefined dictionary, at the end of which the target buffer begins. This dictionary is composed of the pseudo-header for the current packet as defined in [\[RFC2460\]](#), followed by a 16-byte static dictionary (Figure 1). These dictionary bytes are therefore available for backreferencing, but not copied into the final result.



```
16 fe fd 17 fe fd 00 01 00 00 00 00 00 01 00 00
```

Figure 1: The 16 bytes of static dictionary (in hex)

### 3. Integrating 6LoWPAN-GHC into 6LoWPAN-HC

6LoWPAN-GHC is intended to plug in as an NHC format for 6LoWPAN-HC [RFC6282]. This section shows how this can be done (without supplying the detailed normative text yet, although it could be implemented from this page).

#### 3.1. Compressing payloads (UDP and ICMPv6)

GHC is by definition generic and can be applied to different kinds of packets. All the examples given in [Appendix A](#) are for ICMPv6 packets; a single NHC value suffices to define an NHC format for ICMPv6 based on GHC (see below).

In addition it may be useful to include an NHC format for UDP, as many headerlike payloads (e.g., DHCPv6) are carried in UDP. [RFC6282] already defines an NHC format for UDP (11110CPP). What remains to be done is to define an analogous NHC byte formatted, e.g. as shown in Figure 2, and simply reference the existing specification, indicating that for 0b11010cpp NHC bytes, the UDP payload is not supplied literally but compressed by 6LoWPAN-GHC.

```

      0   1   2   3   4   5   6   7
+---+---+---+---+---+---+---+---+
| 1 | 1 | 0 | 1 | 0 | C |   P   |
+---+---+---+---+---+---+---+

```

Figure 2: Proposed NHC byte for UDP GHC (actual value to be allocated by IANA)

To stay in the same general numbering space, we propose 0b11011111 as the NHC byte for ICMPv6 GHC (Figure 3).

```

      0   1   2   3   4   5   6   7
+---+---+---+---+---+---+---+---+
| 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
+---+---+---+---+---+---+---+

```

Figure 3: Proposed NHC byte for ICMPv6 GHC (actual value to be allocated by IANA)

#### 3.2. Compressing extension headers





If the compression of specific extension headers is considered desirable, this can be added in a similar way, e.g. as in Figure 4 (however, probably only EID 0 to 3 need to be assigned). As there is no easy way to extract the length field from the GHC-encoded header before decoding, this would make detecting the end of the extension header somewhat complex. The easiest (and most efficient) approach is to completely elide the length field (in the same way NHC already elides the next header field in certain cases) and reconstruct it only on decompression. To serve as a terminator for the extension header, the reserved bytecode 0b10010000 has been assigned as a stop marker. Note that the stop marker is only needed for extension headers, not for final payloads, the decompression of which is automatically stopped by the end of the packet.

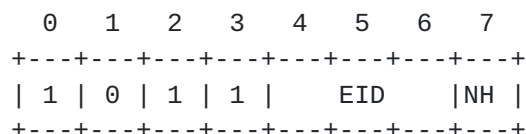


Figure 4: Proposed NHC byte for extension header GHC

### 3.3. Indicating GHC capability

The 6LoWPAN baseline includes just [\[RFC4944\]](#), [\[RFC6282\]](#), [\[RFC6775\]](#) (see [\[I-D.bormann-6lowpan-roadmap\]](#)). To enable the use of GHC, 6LoWPAN nodes need to know that their neighbors implement it. While this can also simply be administratively required, a transition strategy as well as a way to support mixed networks is required.

One way to know a neighbor does implement GHC is receiving a packet from that neighbor with GHC in it ("implicit capability detection"). However, there needs to be a way to bootstrap this, as nobody ever would start sending packets with GHC otherwise.

To minimize the impact on [RFC6775], we propose adding an ND option 6LoWPAN Capability Indication (6CIO), as illustrated in Figure 5.

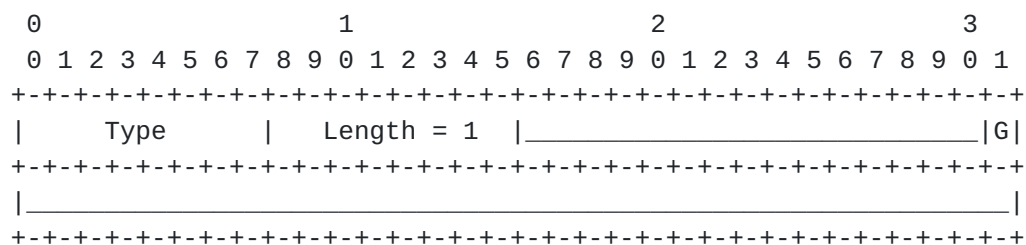


Figure 5: 6LoWPAN Capability Indication Option (6CIO)



The G bit indicates whether the node sending the option is GHC capable.

Once a node receives either an explicit or an implicit indication of GHC capability from another node, it may send GHC-compressed packets to that node. Where that capability has not been recently confirmed, similar to the way PLPMTUD [[RFC4821](#)] finds out about changes in the network, a node SHOULD make use of NUD (neighbor unreachability detection) failures to switch back to basic 6LoWPAN header compression [[RFC6282](#)].

### **3.4. Using the 6CIO Option**

The 6CIO option will typically only be ever sent in 6LoWPAN-ND RS packets (it then cannot itself be GHC compressed unless the host desires to limit itself to talking to GHC capable routers); the resulting 6LoWPAN-ND RA can already make use of GHC and thus indicate GHC capability implicitly, which in turn allows the nodes to use GHC in the 6LoWPAN-ND NS/NA exchange.

6CIO can also be used for future options that need to be negotiated between 6LoWPAN peers; an IANA registry will administrate the flags. (Bits marked by underscores in Figure 5 are reserved for future allocation, i.e., they MUST be sent as zero and MUST be ignored on reception until allocated. Length values larger than 1 MUST be accepted by implementations in order to enable future extensions; the additional bits in the option are then reserved in the same way. For the purposes of the IANA registry, the bits are numbered in msb-first order from the 16th bit of the option onward, i.e., the G bit is flag number 15.) (Additional bits may also be used by a follow-on version of this document if some bit combinations that have been left reserved here are then used in an upward compatible manner.)

Where the use of this option by other specifications is envisioned, the following items have to be kept in mind:

- o The option can be used in any ND packet.
- o Specific bits are set in the option to indicate that a capability is present in the sender. (There may be other ways to infer this information, as is the case in this specification.) Bit combinations may be used as desired. The absence of the capability `_indication_` is signaled by setting these bits to zero; this does not necessarily mean that the capability is absent.
- o The intention is not to modify the semantics of the specific ND packet carrying the option, but to provide the general capability indication described above.



- o Specifications have to be designed such that receivers that do not receive or do not process such a capability indication can still interoperate (presumably without exploiting the indicated capability).
- o The option is meant to be used sparsely, i.e. once a sender has reason to believe the capability indication has been received, there no longer is a need to continue sending it.

#### 4. IANA considerations

[This section to be removed/replaced by the RFC Editor.]

In the IANA registry for the "LOWPAN\_NHC Header Type" (in the "IPv6 Low Power Personal Area Network Parameters"), IANA needs to add the assignments in Figure 6.

10110IIN: Extension header	GHC	[RFCthis]
11010CPP: UDP	GHC	[RFCthis]
11011111: ICMPv6	GHC	[RFCthis]

Figure 6: IANA assignments for the NHC byte

IANA needs to allocate an ND option number for the 6CIO ND option format in the Registry "IPv6 Neighbor Discovery Option Formats" [[RFC4861](#)].

IANA needs to create a registry for "6LoWPAN capability bits" within the "Internet Control Message Protocol version 6 (ICMPv6) Parameters". The bits are allocated by giving their numbers as small non-negative integers as defined in section [Section 3.4](#), preferably in the range 0..47. The policy is "RFC Required" [[RFC5226](#)]. The initial content of the registry is as in Figure 7:

0..14: unassigned	
15: GHC capable bit (G bit)	[RFCthis]
16..47: unassigned	

Figure 7



## 5. Security considerations

The security considerations of [\[RFC4944\]](#) and [\[RFC6282\]](#) apply. As usual in protocols with packet parsing/construction, care must be taken in implementations to avoid buffer overflows and in particular (with respect to the back-referencing) out-of-area references during decompression.

One additional consideration is that an attacker may send a forged packet that makes a second node believe a third victim node is GHC-capable. If it is not, this may prevent packets sent by the second node from reaching the third node (at least until robustness features such as those discussed in [Section 3.3](#) kick in).

No mitigation is proposed (or known) for this attack, except that a victim node that does implement GHC is not vulnerable. However, with unsecured ND, a number of attacks with similar outcomes are already possible, so there is little incentive to make use of this additional attack. With secured ND, 6CIO is also secured; nodes relying on secured ND therefore should use 6CIO bidirectionally (and limit the implicit capability detection to secured ND packets carrying GHC) instead of basing their neighbor capability assumptions on receiving any kind of unprotected packet.

## 6. Acknowledgements

Colin O'Flynn has repeatedly insisted that some form of compression for ICMPv6 and ND packets might be beneficial. He actually wrote his own draft, [\[I-D.oflynn-6lowpan-icmp6hc\]](#), which compresses better, but addresses basic ICMPv6/ND only and needs a much longer spec (around 17 pages of detailed spec, as compared to the single page of core spec here). This motivated the author to try something simple, yet general. Special thanks go to Colin for indicating that he indeed considers his draft superseded by the present one.

The examples given are based on pcap files that Colin O'Flynn, Owen Kirby, Olaf Bergmann and others provided.

The static dictionary was developed, and the bit allocations validated, based on research by Sebastian Dominik.

Erik Nordmark provided input that helped shaping the 6CIO option.

Yoshihiro Ohba insisted on clarifying the notation used for the definition of the bytetimes and their bitfields.





## [7. References](#)

### [7.1. Normative References](#)

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### [7.2. Informative References](#)

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[RFC4821] Mathis, M. and J. Heffner, "Packetization Layer Path MTU Discovery", [RFC 4821](#), March 2007.

[RFC5795] Sandlund, K., Pelletier, G., and L-E. Jonsson, "The RObust Header Compression (ROHC) Framework", [RFC 5795](#), March 2010.

## [Appendix A](#). Examples

This section demonstrates some relatively realistic examples derived from actual PCAP dumps taken at previous interops. (TBD: Add a couple DHCP examples.)

Figure 8 shows an RPL DODAG Information Solicitation, a quite short RPL message that obviously cannot be improved much.

IP header:

```
60 00 00 00 00 08 3a ff fe 80 00 00 00 00 00 00
02 1c da ff fe 00 20 24 ff 02 00 00 00 00 00 00
00 00 00 00 00 00 00 00 1a
```

Payload:

```
9b 00 6b de 00 00 00 00
```

Dictionary:

```
fe 80 00 00 00 00 00 00 02 1c da ff fe 00 20 24
ff 02 00 00 00 00 00 00 00 00 00 00 00 00 00 1a
00 00 00 08 00 00 00 3a 16 fe fd 17 fe fd 00 01
00 00 00 00 00 01 00 00
```

copy: 04 9b 00 6b de

4 nulls: 82

Compressed:

```
04 9b 00 6b de 82
```

Was 8 bytes; compressed to 6 bytes, compression factor 1.33

Figure 8: A simple RPL example

Figure 9 shows an RPL DODAG Information Object, a longer RPL control message that is improved a bit more. Note that the compressed output exposes an inefficiency in the simple-minded compressor used to generate it; this does not devalue the example since constrained nodes are quite likely to make use of simple-minded compressors.

IP header:

```
60 00 00 00 00 5c 3a ff fe 80 00 00 00 00 00 00
```



```

02 1c da ff fe 00 30 23 ff 02 00 00 00 00 00 00
00 00 00 00 00 00 00 00 1a
Payload:
9b 01 7a 5f 00 f0 01 00 88 00 00 00 20 02 0d b8
00 00 00 00 00 00 00 00 ff fe 00 fa ce 04 0e 00 14
09 ff 00 00 01 00 00 00 00 00 00 00 08 1e 80 20
ff ff ff ff ff ff ff ff 00 00 00 00 20 02 0d b8
00 00 00 00 00 00 00 ff fe 00 fa ce 03 0e 40 00
ff ff ff ff 20 02 0d b8 00 00 00 00
Dictionary:
fe 80 00 00 00 00 00 00 02 1c da ff fe 00 30 23
ff 02 00 00 00 00 00 00 00 00 00 00 00 00 00 1a
00 00 00 5c 00 00 00 3a 16 fe fd 17 fe fd 00 01
00 00 00 00 00 01 00 00
copy: 06 9b 01 7a 5f 00 f0
ref(9): 01 00 -> ref 11nnnk 0 7: c7
copy: 01 88
3 nulls: 81
copy: 04 20 02 0d b8
7 nulls: 85
ref(68): ff fe 00 -> ref 101nssss 0 8/11nnnk 1 1: a8 c9
copy: 08 fa ce 04 0e 00 14 09 ff
ref(39): 00 00 01 00 00 -> ref 101nssss 0 4/11nnnk 3 2: a4 da
5 nulls: 83
copy: 06 08 1e 80 20 ff ff
ref(2): ff ff -> ref 11nnnk 0 0: c0
ref(4): ff ff ff ff -> ref 11nnnk 2 0: d0
4 nulls: 82
ref(48): 20 02 0d b8 00 00 00 00 00 00 00 ff fe 00 fa ce
-> ref 101nssss 1 4/11nnnk 6 0: b4 f0
copy: 03 03 0e 40
ref(9): 00 ff -> ref 11nnnk 0 7: c7
ref(28): ff ff ff -> ref 101nssss 0 3/11nnnk 1 1: a3 c9
ref(24): 20 02 0d b8 00 00 00 00
-> ref 101nssss 0 2/11nnnk 6 0: a2 f0
Compressed:
06 9b 01 7a 5f 00 f0 c7 01 88 81 04 20 02 0d b8
85 a8 c9 08 fa ce 04 0e 00 14 09 ff a4 da 83 06
08 1e 80 20 ff ff c0 d0 82 b4 f0 03 03 0e 40 c7
a3 c9 a2 f0
Was 92 bytes; compressed to 52 bytes, compression factor 1.77

```

Figure 9: A longer RPL example

Similarly, Figure 10 shows an RPL DAO message. One of the embedded addresses is copied right out of the pseudo-header, the other one is effectively converted from global to local by providing the prefix FE80 literally, inserting a number of nulls, and copying (some of)



the IID part again out of the pseudo-header. Note that a simple implementation would probably emit fewer nulls and copy the entire IID; there are multiple ways to encode this 50-byte payload into 27 bytes.

IP header:

```
60 00 00 00 00 32 3a ff 20 02 0d b8 00 00 00 00
00 00 00 ff fe 00 33 44 20 02 0d b8 00 00 00 00
00 00 00 ff fe 00 11 22
```

Payload:

```
9b 02 58 7d 01 80 00 f1 05 12 00 80 20 02 0d b8
00 00 00 00 00 00 00 ff fe 00 33 44 06 14 00 80
f1 00 fe 80 00 00 00 00 00 00 00 00 ff fe 00
11 22
```

Dictionary:

```
20 02 0d b8 00 00 00 00 00 00 00 ff fe 00 33 44
20 02 0d b8 00 00 00 00 00 00 00 ff fe 00 11 22
00 00 00 32 00 00 00 3a 16 fe fd 17 fe fd 00 01
00 00 00 00 00 01 00 00
```

copy: 0c 9b 02 58 7d 01 80 00 f1 05 12 00 80

ref(68): 20 02 0d b8 00 00 00 00 00 00 00 ff fe 00 33 44

-> ref 101nssss 1 6/11nnnkkk 6 4: b6 f4

copy: 08 06 14 00 80 f1 00 fe 80

9 nulls: 87

ref(74): ff fe 00 11 22 -> ref 101nssss 0 8/11nnnkkk 3 5: a8 dd

Compressed:

```
0c 9b 02 58 7d 01 80 00 f1 05 12 00 80 b6 f4 08
06 14 00 80 f1 00 fe 80 87 a8 dd
```

Was 50 bytes; compressed to 27 bytes, compression factor 1.85

Figure 10: An RPL DAO message

Figure 11 shows the effect of compressing a simple ND neighbor solicitation.

IP header:

```
60 00 00 00 00 30 3a ff 20 02 0d b8 00 00 00 00
00 00 00 ff fe 00 3b d3 fe 80 00 00 00 00 00 00
02 1c da ff fe 00 30 23
```

Payload:

```
87 00 a7 68 00 00 00 00 fe 80 00 00 00 00 00 00
02 1c da ff fe 00 30 23 01 01 3b d3 00 00 00 00
1f 02 00 00 00 00 00 06 00 1c da ff fe 00 20 24
```

Dictionary:

```
20 02 0d b8 00 00 00 00 00 00 00 ff fe 00 3b d3
fe 80 00 00 00 00 00 02 1c da ff fe 00 30 23
00 00 00 30 00 00 00 3a 16 fe fd 17 fe fd 00 01
00 00 00 00 00 01 00 00
```





```

copy: 04 87 00 a7 68
4 nulls: 82
ref(48): fe 80 00 00 00 00 00 00 02 1c da ff fe 00 30 23
  -> ref 101nssss 1 4/11nnnk 6 0: b4 f0
copy: 04 01 01 3b d3
4 nulls: 82
copy: 02 1f 02
5 nulls: 83
copy: 02 06 00
ref(24): 1c da ff fe 00 -> ref 101nssss 0 2/11nnnk 3 3: a2 db
copy: 02 20 24
Compressed:
  04 87 00 a7 68 82 b4 f0 04 01 01 3b d3 82 02 1f
  02 83 02 06 00 a2 db 02 20 24
Was 48 bytes; compressed to 26 bytes, compression factor 1.85

```

Figure 11: An ND neighbor solicitation

Figure 12 shows the compression of an ND neighbor advertisement.

```

IP header:
  60 00 00 00 00 30 3a fe fe 80 00 00 00 00 00 00
  02 1c da ff fe 00 30 23 20 02 0d b8 00 00 00 00
  00 00 00 ff fe 00 3b d3
Payload:
  88 00 26 6c c0 00 00 00 fe 80 00 00 00 00 00 00
  02 1c da ff fe 00 30 23 02 01 fa ce 00 00 00 00
  1f 02 00 00 00 00 00 06 00 1c da ff fe 00 20 24
Dictionary:
  fe 80 00 00 00 00 00 00 02 1c da ff fe 00 30 23
  20 02 0d b8 00 00 00 00 00 00 00 ff fe 00 3b d3
  00 00 00 30 00 00 00 3a 16 fe fd 17 fe fd 00 01
  00 00 00 00 00 01 00 00
copy: 05 88 00 26 6c c0
3 nulls: 81
ref(64): fe 80 00 00 00 00 00 00 02 1c da ff fe 00 30 23
  -> ref 101nssss 1 6/11nnnk 6 0: b6 f0
copy: 04 02 01 fa ce
4 nulls: 82
copy: 02 1f 02
5 nulls: 83
copy: 02 06 00
ref(24): 1c da ff fe 00 -> ref 101nssss 0 2/11nnnk 3 3: a2 db
copy: 02 20 24
Compressed:
  05 88 00 26 6c c0 81 b6 f0 04 02 01 fa ce 82 02
  1f 02 83 02 06 00 a2 db 02 20 24
Was 48 bytes; compressed to 27 bytes, compression factor 1.78

```



Figure 12: An ND neighbor advertisement

Figure 13 shows the compression of an ND router solicitation. Note that the relatively good compression is not caused by the many zero bytes in the link-layer address of this particular capture (which are unlikely to occur in practice): 7 of these 8 bytes are copied from the pseudo-header (the 8th byte cannot be copied as the universal/local bit needs to be inverted).

IP header:

```
60 00 00 00 00 18 3a ff fe 80 00 00 00 00 00 00
ae de 48 00 00 00 00 01 ff 02 00 00 00 00 00 00
00 00 00 00 00 00 00 02
```

Payload:

```
85 00 90 65 00 00 00 00 01 02 ac de 48 00 00 00
00 01 00 00 00 00 00 00
```

Dictionary:

```
fe 80 00 00 00 00 00 00 ae de 48 00 00 00 00 01
ff 02 00 00 00 00 00 00 00 00 00 00 00 00 02
00 00 00 18 00 00 00 3a 16 fe fd 17 fe fd 00 01
00 00 00 00 00 01 00 00
```

copy: 04 85 00 90 65

ref(11): 00 00 00 00 01 -> ref 11nnnkkk 3 6: de

copy: 02 02 ac

ref(58): de 48 00 00 00 00 01

-> ref 101nssss 0 6/11nnnkkk 5 3: a6 eb

6 nulls: 84

Compressed:

```
04 85 00 90 65 de 02 02 ac a6 eb 84
```

Was 24 bytes; compressed to 12 bytes, compression factor 2.00

Figure 13: An ND router solicitation

Figure 14 shows the compression of an ND router advertisement. The indefinite lifetime is compressed to four bytes by backreferencing; this could be improved (at the cost of minor additional decompressor complexity) by including some simple runlength mechanism.

IP header:

```
60 00 00 00 00 60 3a ff fe 80 00 00 00 00 00 00
10 34 00 ff fe 00 11 22 fe 80 00 00 00 00 00 00
ae de 48 00 00 00 00 01
```

Payload:

```
86 00 55 c9 40 00 0f a0 1c 5a 38 17 00 00 07 d0
01 01 11 22 00 00 00 00 03 04 40 40 ff ff ff ff
ff ff ff ff 00 00 00 00 20 02 0d b8 00 00 00 00
00 00 00 00 00 00 00 00 20 02 40 10 00 00 03 e8
20 02 0d b8 00 00 00 00 21 03 00 01 00 00 00 00
```



```

    20 02 0d b8 00 00 00 00 00 00 00 ff fe 00 11 22
Dictionary:
    fe 80 00 00 00 00 00 00 10 34 00 ff fe 00 11 22
    fe 80 00 00 00 00 00 00 ae de 48 00 00 00 00 01
    00 00 00 60 00 00 00 3a 16 fe fd 17 fe fd 00 01
    00 00 00 00 00 01 00 00
copy: 0c 86 00 55 c9 40 00 0f a0 1c 5a 38 17
2 nulls: 80
copy: 06 07 d0 01 01 11 22
4 nulls: 82
copy: 06 03 04 40 40 ff ff
ref(2): ff ff -> ref 11nnnk 0 0: c0
ref(4): ff ff ff ff -> ref 11nnnk 2 0: d0
4 nulls: 82
copy: 04 20 02 0d b8
12 nulls: 8a
copy: 04 20 02 40 10
ref(38): 00 00 03 -> ref 101nss 0 4/11nnnk 1 3: a4 cb
copy: 01 e8
ref(24): 20 02 0d b8 00 00 00 00
    -> ref 101nss 0 2/11nnnk 6 0: a2 f0
copy: 02 21 03
ref(84): 00 01 00 00 00 00
    -> ref 101nss 0 9/11nnnk 4 6: a9 e6
ref(40): 20 02 0d b8 00 00 00 00 00 00
    -> ref 101nss 1 3/11nnnk 1 5: b3 cd
ref(136): ff fe 00 11 22
    -> ref 101nss 0 15/101nss 0 1/11nnnk 3 3: af a1 db
Compressed:
    0c 86 00 55 c9 40 00 0f a0 1c 5a 38 17 80 06 07
    d0 01 01 11 22 82 06 03 04 40 40 ff ff c0 d0 82
    04 20 02 0d b8 8a 04 20 02 40 10 a4 cb 01 e8 a2
    f0 02 21 03 a9 e6 b3 cd af a1 db
Was 96 bytes; compressed to 59 bytes, compression factor 1.63

```

Figure 14: An ND router advertisement

Figure 15 shows the compression of a DTLS application data packet with a net payload of 13 bytes of cleartext, and 8 bytes of authenticator (note that the IP header is not relevant for this example and has been set to 0). This makes good use of the static dictionary, and is quite effective crunching out the redundancy in the TLS\_PSK\_WITH\_AES\_128\_CCM\_8 header, leading to a net reduction by 15 bytes.

IP header:

```

00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00

```



```

00 00 00 00 00 00 00 00
Payload:
17 fe fd 00 01 00 00 00 00 00 01 00 1d 00 01 00
00 00 00 00 01 09 b2 0e 82 c1 6e b6 96 c5 1f 36
8d 17 61 e2 b5 d4 22 d4 ed 2b
Dictionary:
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00 00 00 2a 00 00 00 00 16 fe fd 17 fe fd 00 01
00 00 00 00 00 01 00 00
ref(13): 17 fe fd 00 01 00 00 00 00 00 01 00
-> ref 101nssss 1 0/11nnnkkk 2 1: b0 d1
copy: 01 1d
ref(10): 00 01 00 00 00 00 00 01 -> ref 11nnnkkk 6 2: f2
copy: 15 09 b2 0e 82 c1 6e b6 96 c5 1f 36 8d 17 61 e2
copy: b5 d4 22 d4 ed 2b
Compressed:
b0 d1 01 1d f2 15 09 b2 0e 82 c1 6e b6 96 c5 1f
36 8d 17 61 e2 b5 d4 22 d4 ed 2b
Was 42 bytes; compressed to 27 bytes, compression factor 1.56

```

Figure 15: A DTLS application data packet

Figure 16 shows that the compression is slightly worse in a subsequent packet (containing 6 bytes of cleartext and 8 bytes of authenticator, yielding a net compression of 13 bytes). The total overhead does stay at a quite acceptable 8 bytes.

```

IP header:
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00
Payload:
17 fe fd 00 01 00 00 00 00 00 05 00 16 00 01 00
00 00 00 00 05 ae a0 15 56 67 92 4d ff 8a 24 e4
cb 35 b9
Dictionary:
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00 00 00 23 00 00 00 00 16 fe fd 17 fe fd 00 01
00 00 00 00 00 01 00 00
ref(13): 17 fe fd 00 01 00 00 00 00 00 00
-> ref 101nssss 1 0/11nnnkkk 0 3: b0 c3
copy: 03 05 00 16
ref(10): 00 01 00 00 00 00 00 05 -> ref 11nnnkkk 6 2: f2
copy: 0e ae a0 15 56 67 92 4d ff 8a 24 e4 cb 35 b9
Compressed:
b0 c3 03 05 00 16 f2 0e ae a0 15 56 67 92 4d ff

```





8a 24 e4 cb 35 b9  
 Was 35 bytes; compressed to 22 bytes, compression factor 1.59

Figure 16: Another DTLS application data packet

Figure 17 shows the compression of a DTLS handshake message, here a client hello. There is little that can be compressed about the 32 bytes of randomness. Still, the net reduction is by 14 bytes.

IP header:

```
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00
```

Payload:

```
16 fe fd 00 00 00 00 00 00 00 00 00 00 36 01 00 00
2a 00 00 00 00 00 00 00 2a fe fd 51 52 ed 79 a4
20 c9 62 56 11 47 c9 39 ee 6c c0 a4 fe c6 89 2f
32 26 9a 16 4e 31 7e 9f 20 92 92 00 00 00 02 c0
a8 01 00
```

Dictionary:

```
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
00 00 00 43 00 00 00 00 16 fe fd 17 fe fd 00 01
00 00 00 00 00 01 00 00
```

ref(16): 16 fe fd -> ref 101nssss 0 1/11nnnkkk 1 5: a1 cd  
 9 nulls: 87

copy: 01 36

ref(16): 01 00 00 -> ref 101nssss 0 1/11nnnkkk 1 5: a1 cd

copy: 01 2a

7 nulls: 85

copy: 23 2a fe fd 51 52 ed 79 a4 20 c9 62 56 11 47 c9

copy: 39 ee 6c c0 a4 fe c6 89 2f 32 26 9a 16 4e 31 7e

copy: 9f 20 92 92

3 nulls: 81

copy: 05 02 c0 a8 01 00

Compressed:

```
a1 cd 87 01 36 a1 cd 01 2a 85 23 2a fe fd 51 52
ed 79 a4 20 c9 62 56 11 47 c9 39 ee 6c c0 a4 fe
c6 89 2f 32 26 9a 16 4e 31 7e 9f 20 92 92 81 05
02 c0 a8 01 00
```

Was 67 bytes; compressed to 53 bytes, compression factor 1.26

Figure 17: A DTLS handshake packet (client hello)



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