Network Working Group Internet-Draft Intended status: Informational Expires: July 27, 2016 M. Chen, Ed. L. Zheng, Ed. Huawei Technologies G. Mirsky, Ed. Ericsson G. Fioccola, Ed. Telecom Italia January 24, 2016

# IP Flow Performance Measurement Framework draft-chen-ippm-coloring-based-ipfpm-framework-05

#### Abstract

This document specifies a measurement method, the IP flow performance measurement (IPFPM). With IPFPM, data packets are marked into different blocks of markers by changing one or more bits of packets. No additional delimiting packet is needed and the performance is measured in-service and in-band without the insertion of additional traffic.

#### Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in <u>RFC 2119</u> [<u>RFC2119</u>].

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This Internet-Draft will expire on July 27, 2016.

Chen, et al.

Expires July 27, 2016

[Page 1]

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### Table of Contents

$\underline{1}$ . Introduction	2
<u>2</u> . Terminology	<u>4</u>
<u>3</u> . Overview and Concept	<u>4</u>
$\underline{4}$ . Consideration on Marking Bits	<u>6</u>
5. Reference Model and Functional Components	<u>6</u>
<u>5.1</u> . Reference Model	<u>6</u>
<u>5.2</u> . Measurement Control Point	7
<u>5.3</u> . Measurement Agent	7
<u>6</u> . Period Number	<u>8</u>
$\underline{7}$ . Re-ordering Tolerance	<u>8</u>
<u>8</u> . Packet Loss Measurement	<u>9</u>
9. Packet Delay Measurement	<u>10</u>
<u>10</u> . IANA Considerations	<u>11</u>
<u>11</u> . Security Considerations	<u>11</u>
<u>12</u> . Acknowledgements	<u>12</u>
<u>13</u> . Contributing Authors	<u>12</u>
<u>14</u> . References	<u>13</u>
<u>14.1</u> . Normative References	<u>13</u>
<u>14.2</u> . Informative References	<u>14</u>
Authors' Addresses	15

# **1**. Introduction

Performance Measurement (PM) is an important tool for service provider for Service Level Agreement (SLA) verification, troubleshooting (e.g., fault localization or fault delimitation) and network visualization. Measurement methods could be roughly put into two categories - active measurement method and passive measurement method. Active method measures performance or reliability parameters by the examination of traffic (IP Packets) injected into the network, expressly for the purpose of measurement by intended measurement

point . On the contrary, passive method measures some performance or reliability parameter associated with the existing traffic (packets) on the network. Both passive and active methods have their strengths and should be regarded as complementary. There are certain scenarios where active measurements alone is not enough or applicable and passive measurements are

desirable[I-D.deng-ippm-passive-wireless-usecase].

With active measurement method, the rate, numbers and interval between the injected packets may affect the accuracy of the results. And for injected test packets it may not be guaranteed to always be in-band with the data traffic in the pure IP network due to Equal Cost Multi-Path (ECMP).

The Multiprotocol Label Switching (MPLS) PM protocol [RFC6374] for packet loss could be considered an example of passive performance measurement method. By periodically inserting auxiliary Operations, Administration and Maintenance (OAM) packets, the traffic is delimited by the OAM packets into consecutive blocks, and the receivers count the packets and calculate the packets loss each block. However, solutions like [RFC6374] depend on the fixed positions of the delimiting OAM packets for packets counting, and thus are vulnerable to out-of-order arrival of packets. This could happen particularly with out-of-band OAM channels, but might also happen with in-band OAM because of the presence of multipath forwarding within the network. Out of order delivery of data and the delimiting OAM can give rise to inaccuracies in the performance measurement figures. The scale of these inaccuracies will depend on data speeds and the variation in delivery, but with out-of-band OAM, this could result in significant differences between real and reported performance.

This document specifies a different measurement method, the IP flow performance measurement (IPFPM). With IPFPM, data packets are marked into different blocks of markers by changing one or more bits of packets without altering normal processing in the network. No additional delimiting packet is needed and the performance can be measured in-service without the insertion of additional traffic. Furthermore, because marking based IP performance measurement does not require extra OAM packets for traffic delimitation, it can be used in situations where there is packets re-ordering. IP Flow Information eXport (IPFIX) [<u>RFC7011</u>] is used for reporting the measurement data of IPFPM to a central calculation element for performance metrics calculation. Several new Information Elements of IPFIX are defined for IPFPM. These are described in the companion document [I-D.chen-ippm-ipfpm-report].

# 2. Terminology

The acronyms used in this document will be listed here.

#### **3**. Overview and Concept

The concept of marking IP packets for performance measurement is described in [<u>I-D.tempia-opsawg-p3m</u>]. Marking of packets in a specific IP flow to different markings divides the flows into different consecutive blocks. Packets in a block have same marking and consecutive blocks will have different markings. This enables the measuring node to count and calculate packet loss and/or delay based on each block of markers without any additional auxiliary OAM packets. The following figure (Figure 1) is an example that illustrates the different markings in a single IP flow in alternate 0 and 1 blocks.

Figure 1: Packet Marking

For packet loss measurement, there are two ways to mark packets: fixed packet numbers or fixed time period for each block of markers. This document considers only fixed time period method. The sender and receiver nodes count the transmitted and received packets/octets based on each block of markers. By counting and comparing the transmitted and received packets/octets, the packet loss can be detected.

For packet delay measurement, there are three solutions. One is similar to the packet loss, it still marks the IP flows to different blocks of markers and uses the time of the marking change as the reference time for delay calculations. This solution requires that there must not be any out-of-order packets; otherwise, the result will not be accurate. Because it uses the first packet of each block of markers for delay measurement, if there is packet reordering, the first packet of each block at the sender will be probably different from the first packet of the block at the receiver. An alternate way is to periodically mark a single packet in the IP flow. Within a given time period, there is only one packet that can be marked. The sender records the timestamp when the marked packet is transmitted, the receiver records the timestamp when received the marked packet. With the two timestamps, the packet delay can be computed. An additional method consists in taking into account the average arrival time of the packets within a single block (i.e. the same block of markers used for packet loss measurement). The network device

locally sums all the timestamps and divides by the total number of packets received, so the average arrival time for that block of packets can be calculated. By subtracting the average arrival times of two adjacent devices it is possible to calculate the average delay between those nodes. This method is robust to out of order packets and also to packet loss (only an error is introduced dependent from the number of lost packets).

A centralized calculation element Measurement Control Point (MCP) is introduced in <u>Section 5.2</u> of this document, to collect the packet counts and timestamps from the senders and receivers for metrics calculation. The IP Flow Information eXport (IPFIX) [RFC7011] protocol is used for collecting the performance measurement statistic information [I-D.chen-ippm-ipfpm-report]. For the statistic information collected, the MCP has to know exactly what packet pair counts (one from the sender and the other is from the receiver) are based on the same block of markers and a pair of timestamps (one from the sender and the other is from the receiver) are based on the same marked packet. In case of average delay calculation the MCP has to know in addition to the packet pair counters also the pair of average timestamps for the same block of markers. The "Period Number" based solution <u>Section 6</u> is introduced to achieve this.

For a specific IP flow to be measured, there may be one or more upstream and downstream Measurement Agents (MAs)(<u>Section 5.3</u>). An IP flow can be identified by the Source IP (SIP) and Destination IP (DIP) addresses, and it may combine the SIP and DIP with any or all of the Protocol number, the Source port, the Destination port, and the Type of Service (TOS) to identify an IP flow. For each flow, there will be a flow identifier that is unique within a certain administrative domain. To simplify the process description, the flows discussed in this document are all unidirectional. A bidirectional flow can be seen as two unidirectional flows.

IFPFM supports the measurement of Multipoint-to-Multipoint (MP2MP) model, which satisfy all the scenarios that include Point-to-Point (P2P), Point-to-Multipoint (P2MP), Multipoint-to-Point (MP2P), and MP2MP. P2P scenario is obvious and can be used anywhere. P2MP and MP2P are very common in mobile backhaul networks. For example, a Cell Site Gateway (CSG) multi-homing to two Radio Network Controller (RNC) Site Gateways (RSGs) is a typical network design. When there is a failure, there is a requirement to monitor the flows between the CSG and the two RSGs hence to determine whether the fault is in the transport network or in the wireless network (typically called "fault delimitation"). This is especially useful in the situation where the transport network belongs to one service provider and wireless network belongs to other service providers.

# <u>4</u>. Consideration on Marking Bits

The marking bits selection are encapsulation related, different bits for marking should be allocated by different encapsulations. This document does not define any marking bits, the marking bits selection for specific encapsulation will be defined in the relevant documents. Also, in general, in order to support packet loss and delay measurement test simultaneously, at least two marking bits are needed. One bit is for packet loss measurement, the other one is for packet delay measurement.

In theory, so long as there are unused bits could be allocated for marking purpose, marking based measurement mechanism can be applied to any encapsulation. It's relatively easier for new defined encapsulations to allocate marking bits. Example of such case is Bit Indexed Explicit Replication (BIER). Two marking bits for passive performance measurementand has been allocated in BIER encapsulation [I-D.ietf-bier-mpls-encapsulation] (Section 3.). But for sophisticated encapsulations, it's harder or even impossible to allocate bits for marking purpose. IPv4 encapsulation is one of the examples. IPv6 encapsulation is in the same situation, but for IPv6, an alternative solution is to leverage the IPv6 extension header for marking.

Since marking will directly change some bits (of the header) of the real traffic packets, it MUST make sure that the marking operations will not affect the forwarding and process of the packets. Specifically, the marking bits MUST NOT be used for ECMP hashing. In addition, to increase the accuracy of measurement, hardware based implementation is desired, so the marking bits SHOULD be easy for hardware implementation. For example, the marking bits should be better at the fixed positions in a packet header.

### 5. Reference Model and Functional Components

## **<u>5.1</u>**. Reference Model

The outline of the measurement system of large-scale measurement platforms (LMAP) network is introduced in [<u>I-D.ietf-lmap-framework</u>]. It describes the main functional components of the LMAP measurement system, and the interactions between the components. The Measurement Agent (MA) of IPFPM could be considered equivalent to the MA of LMAP. The Measurement Control Point (MCP) of IPFPM could be considered as the combined function of Controller and Collector. IP Flow Information eXport (IPFIX) [<u>RFC7011</u>] protocol is used for collecting the performance measurement data on the MAs and reported to MCP. The details are specified in the companion document

Internet-Draft

[<u>I-D.chen-ippm-ipfpm-report</u>]. The control between MCP and MAs are left for future study. Figure 1 gives the reference model of IPFPM.



## 5.2. Measurement Control Point

The Measurement Control Point (MCP) is responsible for collecting the measurement data from the Measurement Agents (MAs) and calculating the performance metrics according to the collected measurement data. For packet loss, based on each block of markers, the difference between the total counts received from all upstream MAs and the total counts received from all downstream MAs are the lost packet numbers. The MCP must make sure that the counts from the upstream MAs and downstream MAs are related to the same marking/packets block. For packet delay (e.g., one way delay), the difference between the timestamps from the downstream MA and upstream MA is the packet delay. Similarly to packet loss, the MCP must make sure the two timestamps are based on the same marked packet. This document introduces a Period Number (PN) based synchronization mechanism which is discussed in details in Section 6.

### 5.3. Measurement Agent

The Measurement Agent (MA) executes the measurement actions (e.g., marks the packets, counts the packets, records the timestamps, etc.), and reports the data to the Measurement Control Point (MCP). Each MA maintains two timers, one (C-timer, used at upstream MA) is for marking change, the other (R-timer, used at downstream MA) is for reading the packet counts and timestamps. The two timers have the same time interval but are started at different time. A MA can be either an upstream or a downstream MA: the role is specific to an IP flow to be measured. For a specific IP flow, the upstream MA will change the marking and read the packet counts and timestamps when the C-timer expires, the downstream MA just reads the packets counts and timestamps when the R-timer expires. The MA may delay the reading for certain time when R-timer expires, in order to tolerant certain degree of packet re-ordering. Section 7 describes this in details.

For each Measurement Task (corresponding to an IP flow) [<u>I-D.ietf-lmap-framework</u>], a MA maintains a pairs of packet counters and a timestamp counter for each block of markers. As for the pair of packet counters, one is for counting packets and the other is for counting octets.

### 6. Period Number

When data are collected on the upstream MA and downstream MA, e.g. packet counts or timestamps, and periodically reported to the MCP, certain synchronization mechanism is required to ensure the data collected are correlated. This document introduces the Period Number (PN) to help the MCP to determine whether any two or more packet counts (from distributed MAs) are related to the same block of markers or any two timestamps are related to the same marked packet.

Period Number assures the data correlation by literally split the packets into different measurement period. The PN is generated each time a MA reads the packet counts or timestamps, and is associated with Each packet count and timestamp reported to the MCP. When the MCP see e.g. two PN associated with two packet counts from a upstream and downstream MA, it consider this two packet counts are for the same measurement period by the same PN, i.e. this two packet counts are related to the same block of markers. The assumption is that the upstream and downstream MAs are time synchronized. This requires that the upstream and downstream MAs are time synchronized. This Protocol (NTP) [RFC5905], or the IEEE 1588 Precision Time Protocol (PTP) [IEEE1588].) The PN is calculated as the modulo of the local time (when the counts or timestamps are read) and the interval of the marking time period.

### 7. Re-ordering Tolerance

In order to allow for a certain degree of packets re-ordering, the R-timer on downstream MA should be started delta-t (Dt) later after the C-timer is started. Dt is a defined period of time and should satisfies the following conditions:

(Time-L - Time-MRO ) < Dt < (Time-L + Time-MRO )</pre>

Where

Time-L: the link delay time between the sender and receiver;

Time-MRO: the maximum re-ordering time difference; if a packet is expected to arrive at t1 but actually arrives at t2, then the Time-MRO = | t2 - t1|.

So, the R-timer should be started at "t + Dt" (where the t is the time when C-timer started).

For simplicity, the C-timer should be started at the beginning of each time period. This document recommends the implementation to support at least these time periods (1s, 10s, 1min, 10min and 1h). So, if the time period is 10s, the C-timer should be started at the time of any multiples of 10 in seconds (e.g., 0s, 10s, 20s, etc.), then the R-timer should be started (e.g., 0s+Dt, 10s+Dt, 20s+Dt, etc.). With this method, each MA can independently start its C-timer and R-timer given that the clocks have been synchronized.

#### 8. Packet Loss Measurement

To simplify the process description, the flows discussed in this document are all unidirectional. A bidirectional flow can be seen as two unidirectional flows. For a specific flow, there will be upstream and downstream MAs and upstream and downstream packet counts/timestamp accordingly.

For packet loss measurement, this document defines the following counters and quantities:

U-CountP[n][m]: U-CountP is a two-dimensional array that stores the number of packets transmitted by each upstream MA in each marking time period. Specifically, parameter "n" is the "period number" of measured blocks of markers while parameter "m" refers to the m-th MA of the upstream MAs.

D-CountP[n][m]: D-CountP is a two-dimensional array that stores the number of packets received by each downstream MA in each marking time period. Specifically, parameter "n" is the "period number" of measured blocks of markers while parameter "m" refers to the m-th MA of the downstream MAs.

U-CountO[n][m]: U-CountO is a two-dimensional array that stores the number of octets transmitted by each upstream MA in each marking time period. Specifically, parameter "n" is the "period number" of measured blocks of markers while parameter "m" refers to the m-th MA of the upstream MAs.

D-CountO[n][m]: D-CountO is a two-dimensional array that stores the number of octets received by each downstream MA in each marking time period. Specifically, parameter "n" is the "period number" of measured blocks of markers while parameter "m" refers to the m-th MA of the downstream MAs.

LossP: the number of packets transmitted by the upstream MAs but not received at the downstream MAs.

LossO: the total octets transmitted by the upstream MAs but not received at the downstream MAs.

The total packet loss of a flow can be computed as follows:

LossP = U-CountP[1][1] + U-CountP[1][2] + .... + U-CountP[n][m] -D-CountP[1][1] - D-CountP[1][2] - .... - D-CountP[n][m'].

Loss0 = U-Count0[1][1] + U-Count0[1][2] + .... + U-Count0[n][m] -D-Count0[1][1] - D-Count0[1][2] - .... - D-Count0[n][m'].

Where the m and m' are the number of upstream MAs and downstream MAs, respectively.

#### 9. Packet Delay Measurement

For packet delay measurement, there will be only one upstream MA and may be one or more (P2MP) downstream MAs. Although the marking based IPFPM supports P2MP model, this document only discusses P2P model, the P2MP model is left for future study. This document defines the following timestamps and quantities:

U-Time[n]: U-Time is a one-dimension array that stores the time when marked packets are sent; in case the "average delay" method is being used, U-Time stores the average of the time when the packets of the same block are sent; parameter "n" is the "period number" of marked packets.

D-Time[n]: D-Time is a one-dimension array that stores the time when marked packets are received; in case the "average delay" method is being used, D-Time stores the average of the time when the packets of the same block are received; parameter "n" is the "period number" of marked packets. This is only for P2P model.

D-Time[n][m]: D-Time a two-dimension array that stores the time when the marked packet is received by downstream MAs at each marking time period; in case the "average delay" method is being used, D-Time stores the average of the times when the packets of the same block are received by downstream MAs at each marking time period. Here, parameter "n" is the "period number" of marked packets while parameter "m" refers to the m-th MA of the downstream MAs. This is for P2MP model which is left for future study.

```
One-way Delay[n]: The one-way delay metric for packet networks is
described in [RFC2679]. The "n" identifies the "period number" of
the marked packet.
One-way Delay[1] = D-Time[1] - U-Time[1].
One-way Delay[2] = D-Time[2] - U-Time[2].
 . . .
One-way Delay[n] = D-Time[n] - U-Time[n].
In the case of two-way delay, the delay is the sum of the two one-way
delays of the two flows that have the same MAs but have opposite
directions.
Two-way Delay[1] = (D-Time[1] - U-Time[1]) + (D-Time'[1] - U-Time[1]) + (D-Time'[1] - U-Time[1]) + (D-Time'[1] - U-Time[1]) + (D-Time'[1] - U-Time[1]) + (D-Time'[1]) + (
U-Time'[1]).
Two-way Delay[2] = (D-Time[2] - U-Time[2]) + (D-Time'[2] -
U-Time'[2]).
 . . .
Two-way Delay[n] = (D-Time[n] - U-Time[n]) + (D-Time'[n] -
U-Time'[n]).
```

```
Where the D-Time and U-Time are for one forward flow, the D-Time' and U-Time' are for reverse flow.
```

# **10**. IANA Considerations

This document makes no request to IANA.

#### **<u>11</u>**. Security Considerations

This document specifies a passive mechanism for measuring packet loss and delay within a Service Provider's network where the IP packets are marked with the unused bits in IP head field, and then inserting additional OAM packets during the measurement is avoided. Obviously, such mechanism does not directly affect other applications running on the Internet but may lead to potential affects to the measurement itself.

First, the measurement itself may be affected by routers (or other network devices) along the path of IP packets intentionally altering the value of marking bits of packets. Just as mentioned before, the mechanism specified in this document is just in the context of one

Service Provider's network, so the routers (or other network devices) are controllable and thus this kind of attack can be omitted.

Second, the measurement can be harmed by attackers injecting artificial traffic. Then authentication techniques, like digital signatures, may be used to guard against such kind of attack.

## **<u>12</u>**. Acknowledgements

The authors would like to thank Adrian Farrel for his review, suggestion and comments to this document.

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# Internet-Draft

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Chen, et al. Expires July 27, 2016 [Page 13]

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