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**Extensions to Path Computation Element Communication Protocol (PCEP)  
for Backup Ingress of a Traffic Engineering Label Switched Path  
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**Abstract**

This document presents extensions to the Path Computation Element Communication Protocol (PCEP) for a PCC to send a request for computing a backup ingress for an MPLS TE P2MP LSP or an MPLS TE P2P LSP to a PCE and for a PCE to compute the backup ingress and reply to the PCC with a computation result for the backup ingress.

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## **1. Introduction**

[RFC4090](#) "Fast Reroute Extensions to RSVP-TE for LSP Tunnels" describes two methods to protect P2P LSP tunnels or paths at local repair points. The local repair points may comprise a number of intermediate nodes between an ingress node and an egress node along the path. The first method is a one-to-one backup method, where a detour backup P2P LSP for each protected P2P LSP is created at each potential point of local repair. The second method is a facility bypass backup protection method, where a bypass backup P2P LSP tunnel is created using MPLS label stacking to protect a potential failure point for a set of P2P LSP tunnels. The bypass backup tunnel can protect a set of P2P LSPs that have similar backup constraints.

[RFC4875](#) "Extensions to RSVP-TE for P2MP TE LSPs" describes how to use the one-to-one backup method and facility bypass backup method to protect a link or intermediate node failure on the path of a P2MP LSP.

However, there is no mention of locally protecting an ingress node failure in a protected P2MP LSP or P2P LSP.

The methods for protecting an ingress node of a P2MP LSP or P2P LSP may be classified into two categories.

A first category uses a backup P2MP LSP that is from a backup ingress node to the number of destination nodes for the P2MP LSP, and a backup P2P LSP that is from a backup ingress node to the destination node for the P2P LSP. The disadvantages of this class of methods include more network resource such as computer power and link bandwidth consumption since the backup P2MP LSP or P2P LSP is from the backup ingress node to the number of destination nodes or the destination respectively.

A second category uses a local P2MP LSP or P2P LSP for protecting the ingress of a P2MP LSP or P2P LSP locally. The local P2MP LSP is from a backup ingress node to the next hop nodes of the ingress of the P2MP LSP. The local P2P LSP is from a backup ingress node to the next hop node of the ingress of the P2P LSP. It is desirable to let PCE compute these backup ingress nodes.

This document defines extensions to the Path Computation Element Communication Protocol (PCEP) for a PCC to send a request for computing a backup ingress node for an MPLS TE P2MP LSP or an MPLS TE P2P LSP to a PCE and for a PCE to compute the backup ingress node and reply to the PCC with a computation result for the backup ingress node.



## **2. Terminology**

This document uses terminologies defined in [RFC5440](#), [RFC4090](#), and [RFC4875](#).

## **3. Conventions Used in This Document**

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119](#).

## **4. Extensions to PCEP**

This section describes the extensions to PCEP for computing a backup ingress of an MPLS TE P2MP LSP and an MPLS TE P2P LSP.

### **4.1. Backup Ingress Capability Advertisement**

#### **4.1.1. Capability TLV in Existing PCE Discovery Protocol**

There are a couple of options for advertising a PCE capability for computing a backup ingress for an MPLS TE P2MP LSP or an MPLS TE P2P LSP.

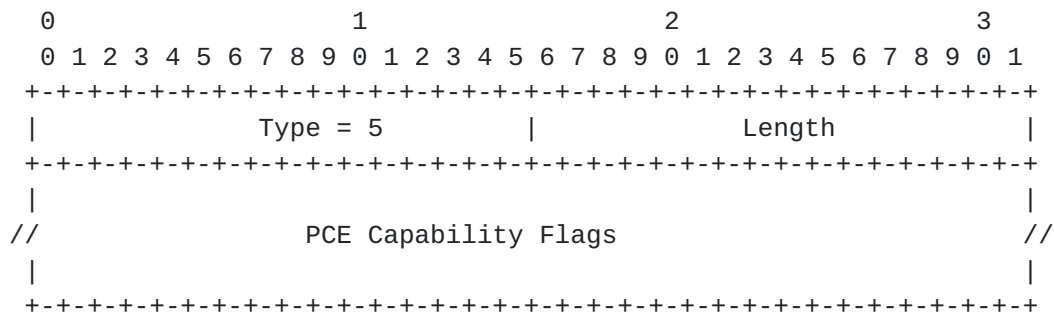
The first option is to define a new flag in the OSPF and ISIS PCE Capability Flags to indicate the capability that a PCE is capable to compute both a backup ingress for an MPLS TE P2MP LSP and a backup ingress for an MPLS TE P2P LSP.

The second option is to define two new flags. One new flag in the OSPF and ISIS PCE Capability Flags indicates the capability that a PCE is capable to compute a backup ingress for an MPLS TE P2MP LSP; and another new flag in the OSPF and ISIS PCE Capability Flags indicates the capability that a PCE is capable to compute a backup ingress for an MPLS TE P2P LSP.

This second option is preferred now.

The format of the PCE-CAP-FLAGS sub-TLV is as follows:





Type: 5  
 Length: Multiple of 4 octets  
 Value: This contains an array of units of 32-bit flags  
 numbered from the most significant as bit zero, where  
 each bit represents one PCE capability.

The following capability bits have been assigned by IANA:

Bit	Capabilities
0	Path computation with GMPLS link constraints
1	Bidirectional path computation
2	Diverse path computation
3	Load-balanced path computation
4	Synchronized path computation
5	Support for multiple objective functions
6	Support for additive path constraints (max hop count, etc.)
7	Support for request prioritization
8	Support for multiple requests per message
9	Global Concurrent Optimization (GCO)
10	P2MP path computation
11-31	Reserved for future assignments by IANA.

Reserved bits SHOULD be set to zero on transmission and MUST be ignored on receipt.

For the second option, one bit such as bit 11 may be assigned to indicate that a PCE is capable to compute a backup ingress for an MPLS TE P2MP LSP and another bit such as bit 12 may be assigned to indicate that a PCE is capable to compute a backup ingress for an MPLS TE P2P LSP.





Bit	Capabilities
11	Backup ingress computation for P2MP LSP
12	Backup ingress computation for P2P LSP
13-31	Reserved for future assignments by IANA.

#### **4.1.2. Open Message Extension**

If a PCE does not advertise its backup ingress computation capability during discovery, PCEP should be used to allow a PCC to discover, during the Open Message Exchange, which PCEs are capable of supporting backup ingress computation.

To achieve this, we extend the PCEP OPEN object by defining a new optional TLV to indicate the PCE's capability to perform backup ingress computation for an MPLS TE P2MP LSP and an MPLS TE P2P LSP.

We request IANA to allocate a value such as 8 from the "PCEP TLV Type Indicators" subregistry, as documented in Section below ("Backup Ingress Capability TLV"). The description is "backup ingress capable", and the length value is 2 bytes. The value field is set to indicate the capability of a PCE for backup ingress computation for an MPLS TE LSP in details.

We can use flag bits in the value field in the same way as the PCE Capability Flags described in the previous section.

The inclusion of this TLV in an OPEN object indicates that the sender can perform backup ingress computation for an MPLS TE P2MP LSP or an MPLS TE P2P LSP.

The capability TLV is meaningful only for a PCE, so it will typically appear only in one of the two Open messages during PCE session establishment. However, in case of PCE cooperation (e.g., inter-domain), when a PCE behaving as a PCC initiates a PCE session it SHOULD also indicate its path computation capabilities.

#### **4.2. Request and Reply Message Extension**

This section describes extensions to the existing RP (Request Parameters) object to allow a PCC to request a PCE for computing a backup ingress of an MPLS TE P2MP LSP or an MPLS TE P2P LSP when the PCE receives the PCEP request.



#### [4.2.1.](#) RP Object Extension

The following flags are added into the RP Object:

The I bit is added in the flag bits field of the RP object to tell the receiver of the message that the request/reply is for computing a backup ingress of an MPLS TE P2MP LSP and an MPLS TE P2P LSP.

o I ( Backup Ingress bit - 1 bit):

0: This indicates that this is not PCReq/PCRep for backup ingress.

1: This indicates that this is PCReq or PCRep message for backup ingress.

The IANA request is referenced in Section below (Request Parameter Bit Flags) of this document.

This I bit with the N bit defined in [RFC6006](#) can indicate whether the request/reply is for a backup ingress of an MPLS TE P2MP LSP or an MPLS TE P2P LSP.

o I = 1 and N = 1: This indicates that this is a PCReq/PCRep message for backup ingress of an MPLS TE P2MP LSP.

o I = 1 and N = 0: This indicates that this is a PCReq/PCRep message for backup ingress of an MPLS TE P2P LSP.

#### [4.2.2.](#) External Source Node

In addition to the information about the path that an MPLS TE P2MP LSP or an MPLS TE P2P LSP traverses, a request message may comprise other information that may be used for computing the backup ingress for the P2MP LSP or P2P LSP. For example, the information about an external source node, from which data traffic is delivered to the ingress node of the P2MP LSP or P2P LSP and transported to the egress node(s) via the P2MP LSP or P2P LSP, is useful for computing a backup ingress node.

The PCC can specify an external source node (ESN) Object. The ESN Object has the same format as the IRO object defined in [\[RFC5440\]](#) except that it only supports IPv4 and IPv6 prefix sub-objects.



The object can only be carried in a PCReq message. A Path Request may carry at most one external source node Object.

The Object-Class and Object-types will need to be allocated by IANA. The IANA request is documented in Section below. (PCEP Objects).

Alternatively, we may use END-POINTS to represent an external source node in a request message for computing a backup ingress node of MPLS LSP.

To represent an external source node efficiently, we define a new type of END-POINTS objects for computing a backup ingress node of MPLS LSP. The format of the new END-POINTS object body for IPv4 (Object-Type 3) is as follows:

```

      0               1               2               3
      0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|                               External Source Type (11)                               |
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|                               External Source IPv4 address                               |
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+

```

The new type of END-POINTS is External Source Node Type (11). The final value for the type will be assigned by IANA. This new type of END-POINTS object contains an external source node IPv4 address.

#### **4.2.3. Constraints between Ingress and Backup Ingress**

A request message sent to a PCE from a PCC for computing a backup ingress of an MPLS TE P2MP LSP or an MPLS TE P2P LSP may comprise a constraint indicating that there must be a path from the backup ingress node to be computed to the ingress node of the P2MP LSP or P2P LSP and that the length of the path is within a given hop limit such as one hop.

This constraint can be considered as default by a PCE or explicitly sent to the PCE by a PCC [TBD].

#### **4.2.4. Constraints for Backup Path**

A request message sent to a PCE from a PCC for computing a backup ingress of a P2MP LSP or P2P LSP may comprise a constraint indicating that the backup ingress node to be computed may not be a node on the P2MP LSP or P2P LSP. In addition, the request message may comprise a list of nodes, each of which is a candidate for the backup ingress node.



A request message sent to a PCE from a PCC for computing a backup ingress of a P2MP LSP or P2P LSP may comprise a constraint indicating that there must be a path from the backup ingress node to be computed to the next-hop nodes of the ingress node of the P2MP LSP or P2P LSP and that there is not an internal node of the path from the backup ingress to the next-hop nodes on the P2MP LSP or P2P LSP .

Most of these constraints for the backup path can be considered as default by a PCE. The constraints for the backup path may be explicitly sent to the PCE by a PCC [TBD].

#### **4.2.5. Backup Ingress Node**

The PCE may send a reply message to the PCC in return to the request message for computing a new backup ingress node. The reply message may comprise information about the computed backup ingress node, which is contained in the path from the backup ingress computed to the next-hop node(s) of the ingress node of the P2MP LSP or P2P LSP.

The backup ingress node is the root or head node of the backup path computed.

#### **4.2.6. Backup Ingress PCEP Error Objects and Types**

In some cases, the PCE may not complete the backup ingress computation as requested, for example based on a set of constraints. As such, the PCE may send a reply message to the PCC that indicates an unsuccessful backup ingress computation attempt. The reply message may comprise a PCEP-error object, which may comprise an error-value, error-type and some detail information.

#### **4.2.7. Request Message Format**

The PCReq message is encoded as follows using RBNF as defined in [[RFC5511](#)].

Below is the message format for a request message:





```
<PCReq Message> ::= <Common Header>
                    [<svec-list>]
                    <request>
<request> ::= <RP>
              <end-point-rro-pair-list>
              [<OF>]
              [<LSPA>]
              [<BANDWIDTH>]
              [<metric-list>]
              [<ESNO>]
              [<IRO>]
              [<LOAD-BALANCING>]
```

where:

<ESNO> is an external source node object.

Figure 1: The Format for a Request Message

The definitions for svec-list, RP, end-point-rro-pair-list, OF, LSPA, BANDWIDTH, metric-list, IRO, and LOAD-BALANCING are described in [RFC5440](#) and [RFC6006](#).

#### **4.2.8. Reply Message Format**

The PCRep message is encoded as follows using RBNF as defined in [\[RFC5511\]](#).

Below is the message format for a reply message:



```
<PCRep Message> ::= <Common Header>
                        <response>
<response> ::= <RP>
                <end-point-path-pair-list>
                [<NO-PATH>]
                [<attribute-list>]
where:

<end-point-path-pair-list> ::=
    [<END-POINTS>]<path>[<end-point-path-pair-list>]

<path> ::= (<ERO>|<SERO>) [<path>]

<attribute-list> ::= [<OF>]
                    [<LSPA>]
                    [<BANDWIDTH>]
                    [<metric-list>]
                    [<IRO>]
```

Figure 2: The Format for a Reply Message

The definitions for RP, NO-PATH, END-POINTS, OF, LSPA, BANDWIDTH, metric-list, IRO, and SERO are described in [RFC5440](#), [RFC6006](#) and [RFC4875](#).

## 5. Security Considerations

The mechanism described in this document does not raise any new security issues for the PCEP, OSPF and IS-IS protocols.

## 6. IANA Considerations

This section specifies requests for IANA allocation.

### 6.1. Backup Ingress Capability Flag

Two new OSPF Capability Flags are defined in this document to indicate the capabilities for computing a backup ingress for an MPLS TE P2MP LSP and an MPLS TE P2P LSP. IANA is requested to make the assignment from the "OSPF Parameters Path Computation Element (PCE) Capability Flags" registry:



Bit	Description	Reference
11	Backup ingress for P2MP LSP	This I-D
12	Backup ingress for P2P LSP	This I-D

## **6.2. Backup Ingress Capability TLV**

A new backup ingress capability TLV is defined in this document to allow a PCE to advertize its backup ingress computation capability. IANA is requested to make the following allocation from the "PCEP TLV Type Indicators" sub-registry.

Value	Description	Reference
8	Backup ingress capable	This I-D

## **6.3. Request Parameter Bit Flags**

A new RP Object Flag has been defined in this document. IANA is requested to make the following allocation from the "PCEP RP Object Flag Field" Sub-Registry:

Bit	Description	Reference
16	Backup ingress (I-bit)	This I-D

## **6.4. PCEP Objects**

An External Source Node Object-Type is defined in this document. IANA is requested to make the following Object-Type allocation from the "PCEP Objects" sub-registry:

Object-Class Value	33
Name	External Source Node
Object-Type	1: IPv4 2: IPv6 3-15: Unassigned
Reference	This I-D



## **7. Acknowledgement**

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