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H. Chen  
Futurewei  
M. Toy  
Verizon  
X. Liu  
Volta Networks  
L. Liu  
Fujitsu  
Z. Li  
China Mobile  
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**PCEP Link State Abstraction**  
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Abstract

This document presents extensions to the Path Computation Element Communication Protocol (PCEP) for a child PCE to abstract its domain information to its parent for supporting a hierarchical PCE system.

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## [1.](#) Introduction

A hierarchical PCE architecture is described in [RFC 6805](#), in which a parent PCE maintains an abstract domain topology, which contains its child domains (seen as vertices in the topology) and the connections among them.

For a domain for which a child PCE is responsible, connections attached to the domain may comprise inter-domain links and Area Border Routers (ABRs). For a parent PCE to have the abstract domain topology, each of its child PCEs needs to advertise its connections to the parent PCE.



In addition to the connections attached to the domain, there may be some access points in the domain, which are the addresses in the domain to be accessible outside of the domain. For example, an address of a server in the domain that provides a number of services to users outside of the domain is an access point.

This document presents extensions to the Path Computation Element Communication Protocol (PCEP) for a child PCE to advertise the information about its connections and access points to its parent PCE and for the parent PCE to build and maintain the abstract domain topology based on the information. The extensions may reduce configurations, thus simplify operations on a PCE system.

A child PCE is simply called a child and a parent PCE is called a parent in the following sections.

## **2. Terminology**

ABR: Area Border Router. Router used to connect two IGP areas (Areas in OSPF or levels in IS-IS).

ASBR: Autonomous System (AS) Border Router. Router used to connect together ASes via inter-AS links.

TED: Traffic Engineering Database.

This document uses terminology defined in [[RFC5440](#)].

## **3. Conventions Used in This Document**

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [[RFC2119](#)].

## **4. Connections and Accesses**

A connection is an inter-domain link between two domains in general. An ABR is also a connection, which connects two special domains called areas in a same Autonomous System (AS).

An access point in a domain is an address in the domain to be accessible to the outside of the domain. An access point is simply called an access.



#### **4.1. Information on Inter-domain Link**

An inter-domain link connects two domains in two different ASes. Since there is no IGP running over an inter-domain link, we may not obtain the information about the link generated by an IGP. We may suppose that IP addresses are configured on inter-domain links.

For a point-to-point (P2P) link connecting two ASBRs A and B in two different domains, from A's point of view, the following information about the link may be obtained:

- 1) Link Type: P2P
- 2) Local IP address
- 3) Remote IP address
- 4) Traffic engineering metric
- 5) Maximum bandwidth
- 6) Maximum reservable bandwidth
- 7) Unreserved bandwidth
- 8) Administrative group
- 9) SRLG

We will have a link ID if it is configured; otherwise no link ID (i.e., the Router ID of the neighbor) may be obtained since no IGP adjacency over the link is formed.

For a broadcast link connecting multiple ASBRs in a number of domains, on each of the ASBRs X, the same information about the link as above may be obtained except for the followings:

- a) Link Type: Multi-access,
- b) Local IP address with mask length, and
- c) No Remote IP address.

In other words, the information about the broadcast link obtained by ASBR X comprises a), b), 4) to 9), but does not include any remote IP address or link ID. We will have a link ID if it is configured; otherwise no link ID (i.e., the interface address of the designated router for the link) may be obtained since no IGP selects it.

A parent constructs an abstract AS domain topology after receiving the information about each of the inter-domain links described above from its children.

[RFC 5392](#) and [RFC 5316](#) describe the distributions of inter-domain links in OSPF and IS-IS respectively. For each inter-domain link,



its neighboring AS number and neighboring ASBR Identity (TE Router ID) need to be configured in IGP (OSPF or IS-IS).

In addition, an IGP adjacency between a network node running IGP and a PCE running IGP as a component needs to be configured and fully established if we want the PCE to obtain the inter-domain link information from IGP.

These configurations and IGP adjacency establishment are not needed if the extensions in this draft are used.

[RFC 7752](#) (BGP-LS) describes the distributions of TE link state information including inter-domain link state. A BGP peer between a network node running BGP and a PCE running BGP as a component needs to be configured and the peer relation must be established before the PCE can obtain the inter-domain link information from BGP. However, some networks may not run BGP.

#### **[4.2.](#) Information on ABR**

For an AS running IGP and containing multiple areas, an ABR connects two or more areas. For each area connected to the ABR, the PCE as a child responsible for the area sends its parent the information about the ABR, which indicates the identifier (ID) of the ABR.

A parent has the information about each of its children, which includes the domain such as the area for which the child is responsible. The parent knows all the areas to which each ABR connects after receiving the information on the ABR from each of its children.

#### **[4.3.](#) Information on Access Point**

For an IP address in a domain to be accessible outside of the domain, the PCE as a child responsible for the domain sends its parent the information about the address.

The parent has all the access points (i.e., IP addresses) to be accessible outside of all its children's domains after receiving the information on the access points from each of its children.

### **[5.](#) Extensions to PCEP**

This section focuses on procedures for abstracting domain information after briefing messages containing the abstract information.





### **5.1. Messages for Abstract Information**

A child abstracts its domain to its parent through sending its parent a message containing the abstract information on the domain. After the relation between the child and the parent is determined, the parent has some information on the child, which includes the child's ID and domain. The message does not need to contain this information. It comprises the followings:

- o For new or updated Connections and Accesses,
  - \* Indication of Update Connections and Accesses
  - \* Detail Information about Connections and Accesses
- o For Connections and Accesses down,
  - \* Indication of Withdraw Connections and Accesses
  - \* ID Information about Connections and Accesses

For a P2P link from ASBR A to B and a broadcast link connecting to A, the detail information on the links includes A's ID, the information on the P2P link and the information on the broadcast link described in [Section 4](#). The ID information on the links includes A's ID, 1) to 3) for the P2P link and a) to b) for the broadcast link described in [Section 4](#). A link ID for a link is included if it is configured.

For an ABR X, the information on X includes X's ID and a flag indicating that X is ABR.

For an Access X (address), the detail information on X includes X and a cost associated with it. The ID information on X is X itself.

There are a few ways to encode the information above into a message. For example, one way is to extend an existing Notification message for including the information. Another way is to use a new message. These are put in [Appendix A](#) for your reference.

### **5.2. Procedures**

#### **5.2.1. Child Procedures**

##### **5.2.1.1. New or Changed Connections and Accesses**

After a child determines its parent, it sends the parent a message containing the information about the connections (i.e., inter-domain



links and ABRs) from its domain to its adjacent domains and the access points in its domain.

For any new or changed inter-domain links, ABRs and access points in the domain for which a child is responsible, the child sends its parent a message containing the information about these links, ABRs and access points with indication of Update Connections and Accesses.

For example, for a new inter-domain P2P link from ASBR A in a child's domain to ASBR B in another domain, the child sends its parent a message containing an indication of Update Connections and Accesses, A's ID, and the detail information on the link described in [section 4.1](#).

For multiple new or changed inter-domain links from ASBR A, the child sends its parent a message having an indication of Update Connections and Accesses, and A's ID followed by the detail information about each of the links.

In another example, for a new or changed inter-domain broadcast link connected to ASBR X, an ABR Y and an access point 10.10.10.1/32 with cost 10 in a child's domain, the child sends its parent a message containing an indication of Update Connections and Accesses, and X's ID followed by the detail information about the link attached to X and the detail information about ABR Y, and the information on access 10.10.10.1/32 with cost 10.

For changes on the attributes (such as bandwidth) of an inter-domain link, a threshold may be used to control the frequency of updates that are sent from a child to its parent. At one extreme, the threshold is set to let a child send its parent a update message for any change on the attributes of an inter-domain link. At another extreme, the threshold is set to make a child not to send its parent any update message for any change on the attributes of an inter-domain link. Typically, the threshold is set to allow a child to send its parent a update message for a significant change on the attributes of an inter-domain link.

#### **[5.2.1.2](#). Connections and Accesses Down**

For any inter-domain links, ABRs and access points down in the domain for which a child is responsible, the child sends its parent a message containing the information about these links, ABRs and access points with indication of Withdraw Connections and Accesses.

For example, for the inter-domain P2P link from ASBR A down, the child sends its parent a message containing an indication of Withdraw



Connections and Accesses, and A's ID, which is followed by the ID information about the link.

For multiple inter-domain links from ASBR A down, the child sends its parent a message having an indication of Withdraw Connections and Accesses, and A's ID, which is followed by the ID information about each of the links.

#### **5.2.1.3. Child and Parent in Same Organization**

If a child and its parent are in a same organization, the child may send its parent the information inside its domain. For a parent, after all its children in its organization send their parent the information in their domains, connections and access points, it has in its TED the detail information inside each of its children's domains and the connections among these domains. The parent can compute a path crossing these domains directly and efficiently without sending any path computation request to its children.

#### **5.2.1.4. Child as a Parent**

There are a few ways in which a child as a parent abstracts its domain information to its parent.

One way is that the child sends its parent all its domain information if the child and the parent are in a same organization. The information includes the detail network topology inside each of the child's domains, the inter-domain links connecting the domains that the child's children are responsible and the inter-domain links connecting these domains to other adjacent domains.

In another way, the child abstracts each of the domains that its children are responsible as a cloud (or say abstract node) and these clouds are connected by the inter-domain links attached to the domains. The child sends its parent all the inter-domain links attached to any of the domains.

In a third way, the child abstracts all its domains including the domains for which its children are responsible as a cloud. This abstraction is described below in details.

If a parent P1 is also a child of another parent P2, P1 as a child sends its parent P2 a message containing the information about the connections and access points. P1 as a parent has the connections among its children's domains. But these connections are hidden from its parent P2. P1 may have connections from its children's domains to other domains. P1 as a child sends its parent P2 these connections.



P1 as a parent has the access points in its children's domains to be accessible outside of the domains. P1 as child may not send all of these to its parent P2. It sends its parent some of these access points according to some local policies.

From P2's point of view, its child P1 is responsible for one domain, which has some connections to its adjacent domains and some access points to be accessible.

### **5.2.2. Parent Procedures**

#### **5.2.2.1. Process Connections and Accesses**

A parent stores into its TED the connections and accesses for each of its children according to the messages containing connections and accesses received. For a message containing Update Connections and Accesses, it updates the connections and accesses in the TED accordingly. For a message containing Withdraw Connections and Accesses, it removes the connections and accesses from the TED.

After receiving the messages for connections and accesses from its children, the parent builds and maintains the TED for the topology of its children's domains, in which each of the domains is seen as a cloud or an abstract node. The information inside each of the domains is hidden from the parent. There are connections among the domains and the access points in the domains to be accessible in the topology.

For a new P2P link from node A to B with no link ID configured, when receiving a message containing the link from a child, the parent stores the link from A into its TED, where A is attached to the child's domain as a cloud. It finds the link's remote end B using the remote IP address of the link. After finding B, it associates the link attached to A with B and the link attached to B with A. This creates a bidirectional connection between A and B.

For a new P2P link from node A to B with link ID configured, when receiving a message containing the link, the parent stores the link from A into its TED. It finds the link's remote end B using the link ID (i.e., B's ID).

For a new broadcast link connecting multiple nodes with no link ID configured, when the parent receives a message containing the link attached to node X, it stores the link from X into its TED. It finds the link's remote end P using the link's local IP address with network mask. P is a Pseudo node identified by the local IP address of the designated node selected from the nodes connected to the link. After finding P, it associates the link attached to X with P and the





link connected to P with X. If P is not found, a new Pseudo node P is created. The parent associates the link attached to X with P and the link attached to P with X. This creates a bidirectional connection between X and P.

The first node and second node from which the parent receives a message containing the link is selected as the designed node and backup designed node respectively. After the designed node is down, the backup designed node becomes the designed node and the node other than the designed node with the largest local IP address connecting to the link is selected as the backup designed node.

When the old designed node is down and the backup designed node becomes the new designed node, the parent updates its TED through removing the link between each of nodes X and old P (the Pseudo node corresponding to the old designed node) and adding a link between each of nodes X (still connecting to the broadcast link) and new P (the Pseudo node corresponding to the new designed node).

#### **5.2.2.2. Detail Topology in a Domain**

If a parent is in a same organization as its child, it stores into its TED the detail information inside the child's domain when receiving a message containing the information from the child; otherwise, it discards the information and issues a warning indicating that the information is sent to a wrong place.

### **6. Security Considerations**

The mechanism described in this document does not raise any new security issues for the PCEP protocols.

### **7. IANA Considerations**

This section specifies requests for IANA allocation.

### **8. Acknowledgement**

The authors would like to thank Jescia Chen, Adrian Farrel, and Eric Wu for their valuable comments on this draft.

### **9. References**

#### **9.1. Normative References**



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- [RFC5440] Vasseur, JP., Ed. and JL. Le Roux, Ed., "Path Computation Element (PCE) Communication Protocol (PCEP)", [RFC 5440](#), DOI 10.17487/RFC5440, March 2009, <<https://www.rfc-editor.org/info/rfc5440>>.
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## **9.2. Informative References**

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## Appendix A. Message Encoding

### A.1. Extension to Existing Message

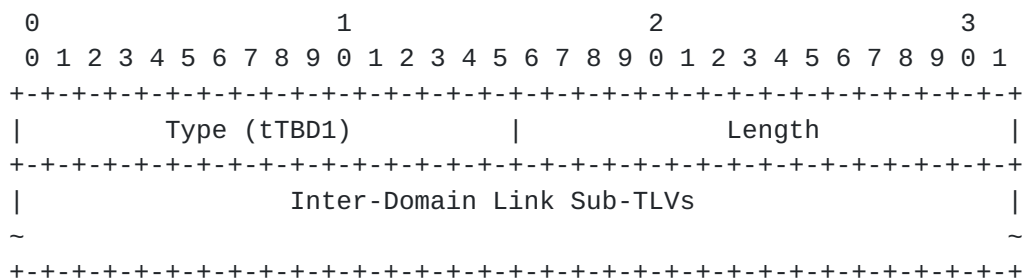
An existing Notification message may be extended to advertise the information about connections and access points. The following new Notification-type (NT) and Notification-value (NV) of a NOTIFICATION object in the message are defined:

- o NT=8 (TBD): Connections and Accesses
  - \* NV=1: Update Connections and Accesses. A NT=8 and NV=1 indicates that the child sends its parent updates on the information about Connections and Accesses, and TLVs containing the information are in the object.
  - \* NV=2: Withdraw Connections and Accesses. A NT=8 and NV=2 indicates that the child asks its parent to remove Connections and Accesses indicated by TLVs in the object.

### A.1.1. TLVs

Four TLVs are defined for connections and accesses. They are Inter-Domain link TLV, Router-ID TLV, Access IPv4/IPv6 Prefix TLV.

The format of the Inter-Domain link TLV is illustrated below.

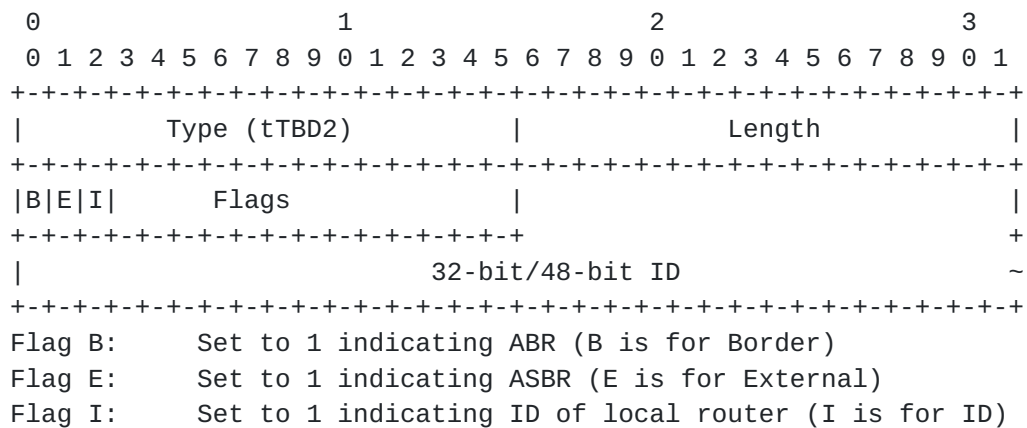


An Inter-Domain link TLV describes a single inter-domain link. It comprises a number of inter-domain link sub-TLVs for the information described in [section 4](#), which are the sub-TLVs defined in [RFC 3630](#) or their equivalents except for the local IP address with mask length defined below.

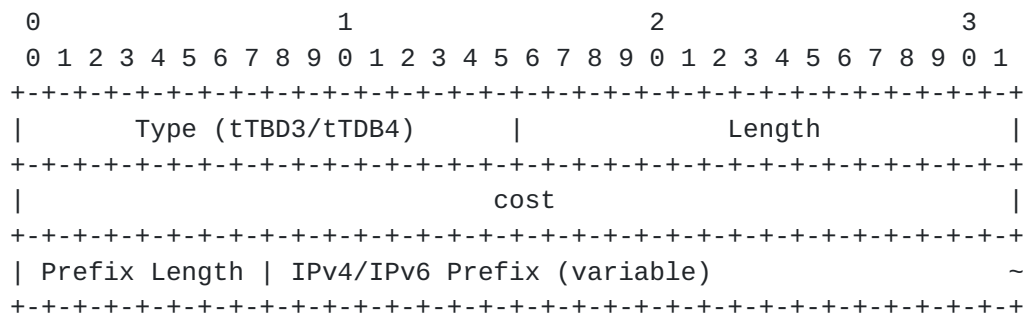
The format of the Router-ID TLV is shown below. Undefined flags MUST be set to zero. The ID indicates the ID of a router. For a router running OSPF, the ID may be the 32-bit OSPF router ID of the router. For a router running IS-IS, the ID may be the 48-bit IS-IS router ID



of the router. For a router not running OSPF or IS-IS, the ID may be the 32-bit ID of the router configured.

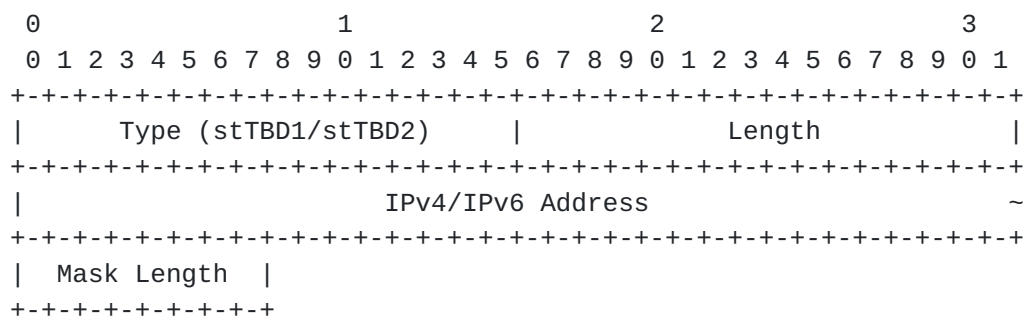


The format of the Access IPv4/IPv6 Prefix TLV is shown as follows. The cost is the metric to the prefix. The Prefix Length indicates the length of the prefix. The IPv4/IPv6 Prefix indicates an access IPv4/IPv6 address prefix.



### A.1.2. Sub-TLVs

The format of the Sub-TLV for a local IPv4/IPv6 address with mask length is shown as follows.







The IPv4/IPv6 Address indicates the local IPv4/IPv6 address of a link. The Mask Length indicates the length of the IPv4/IPv6 address mask.

## **A.2. New Message**

A new message may be defined to advertise the connections and accesses from a child to its parent. The format of the message containing Connections and Access (AC for short) is as follows:

```
<AC Message> ::= <Common Header> <NRP>
                  <Connection-List> [<Access-List>]
where:
  <Connection-List> ::= <Connection> [<Connection-List>]
  <Connection> ::= [<Inter-Domain-Link> | <ABR>]
  <Access-List> ::= <Access-Address> [<Access-List>]
```

Where the value of the Message-Type in the Common Header indicates the new message type. The exact value is to be assigned by IANA. A new RP (NRP) object will be defined, which follows the Common Header.

A new flag W (Withdraw) in the NRP object is defined to indicate whether the connections and access are withdrawn. When flag W is set to one, the parent removes the connections and accesses contained in the message after receiving it. When flag W is set to zero, the parent adds/updates the connections and accesses in the message after receiving it.

An alternative to flag W in the NRP object is a similar flag in each CONNECTION and ACCESS object such as using one bit in Res flags for flag W. For example, when the flag is set to one in the object, the parent removes the connections and accesses in the object after receiving it. When the flag is set to zero in the object, the parent adds/updates the connections and accesses in the object after receiving it.

In another option, one byte in a CONNECTION and ACCESS Object is defined as flags field and one bit is used as flag W. The other undefined bits in the flags field MUST be set to zero.

The objects in the new message are defined below.

### **A.2.1. CONNECTION and ACCESS Object**

A new object, called CONNECTION and ACCESS Object (CA for short), is defined. It has Object-Class octBD1. Four Object-Types are defined under CA object:



- o CA Inter-Domain Link: CA Object-Type is 1.
- o CA ABR: CA Object-Type is 2.
- o CA Access IPv4 Prefix: CA Object-Type is 3.
- o CA Access IPv6 Prefix: CA Object-Type is 4.

Each of these objects are described below.

The format of Inter-Domain Link object body is as follows:

```

Object-Class = ocTBD1 (Connection and Access)
Object-Type = 1 (CA Inter-Domain Link)
  0                               1                               2                               3
  0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+-+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
|W|  Flags  |                               Router-ID TLV          |
+-+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
~
+-+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
|                               Inter-Domain Link TLVs              |
~
+-+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+

```

The Router-ID TLV indicates an ASBR in the domain, which is a local end of inter-domain links. Each of the Inter-Domain Link TLVs describes an inter-domain link and comprises a number of inter-domain link Sub-TLVs. Flag W=1 indicates withdraw the links. W=0 indicates new or changed links.

The format of ABR object body is illustrated below:

```

Object-Class = ocTBD1 (Connection and Access)
Object-Type = 2 (CA ABR)
  0                               1                               2                               3
  0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+-+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
|W|  Flags  |                               Router-ID TLVs          |
+-+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
~
+-+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+

```

Each of the Router-ID TLVs indicates an ABR in the domain. Flag W=1 indicates withdraw the ABRs. W=0 indicates new ABRs.



The format of Access IPv4/IPv6 Prefix object body is as follows:

```

Object-Class = ocTBD1 (Connection and Access)
Object-Type = 3/4 (CA Access IPv4/IPv6 Prefix)
0              1              2              3
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|W|   Flags   |           Access IPv4/IPv6 Prefix TLVs           |
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
~                                                         ~
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+

```

Each of the Access IPv4/IPv6 Prefix TLVs describes an access IPv4/IPv6 address prefix in the domain, which is accessible to outside of the domain. Flag W=1 indicates withdraw the address prefixes. W=0 indicates new address prefixes.

The TLVs in the objects are the same as those described above.

#### Authors' Addresses

Huaimo Chen  
Futurewei  
Boston, MA  
USA

E-Mail: [Huaimo.chen@futurewei.com](mailto:Huaimo.chen@futurewei.com)

Mehmet Toy  
Verizon  
USA

E-Mail: [mehmet.toy@verizon.com](mailto:mehmet.toy@verizon.com)

Xufeng Liu  
Volta Networks  
McLean, VA  
USA

E-Mail: [xufeng.liu.ietf@gmail.com](mailto:xufeng.liu.ietf@gmail.com)



Lei Liu  
Fujitsu  
USA

E-Mail: [liulei.kddi@gmail.com](mailto:liulei.kddi@gmail.com)

Zhenqiang Li  
China Mobile  
No.32 Xuanwumenxi Ave., Xicheng District  
Beijing 100032  
P.R. China

E-Mail: [li\\_zhenqiang@hotmail.com](mailto:li_zhenqiang@hotmail.com)



