

Network Working Group

Internet-Draft

[<draft-crowcroft-rmfp-00.txt>](#)

J Crowcroft (UCL)

Z Wang (UCL)

A Ghosh (UTS)

C Diot (INRIA)

Nov 1996

## **RMFP: A Reliable Multicast Framing Protocol**

Status of this Memo

This document is an Internet-Draft. Internet-Drafts are working documents of the Internet Engineering Task Force (IETF), its areas, and its working groups. Note that other groups may also distribute working documents as Internet-Drafts.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as ``work in progress.''

To learn the current status of any Internet-Draft, please check the ``[id-abstracts.txt](#)'' listing contained in the Internet-Drafts Shadow Directories on [ds.internic.net](#) (US East Coast), [nic.nordu.net](#) (Europe), [ftp.isi.edu](#) (US West Coast), or [munnari.oz.au](#) (Pacific Rim).

The distribution of this memo is unlimited. It is filed as [<draft-crowcroft-rmfp-01.txt>](#), and expires 15 May, 1997. Please send comments to the authors.

### **1. Introduction**

There has been considerable interest in reliable multicast, and a number of reliable multicast transport systems have been proposed in the past years.

Reliable multicast transport is considerably more complex than reliable unicast. It is difficult to build a generic reliable transport protocol for multicast, much as TCP is a generic transport protocol for unicast, since different applications often have very different reliability requirements and modes of operation.

In this document we propose a framing protocol for reliable multicast transport - Reliable Multicast Framing Protocol (RMFP). RMFP runs over multicast UDP and itself does not provide any reliability (or functionality in a larger extend). Reliability and other protocol functionalities will be defined in specific profiles. The purpose of RMFP is to provide a common framework upon which a set of reliable multicast systems can be built and share similar functionalities where exist.

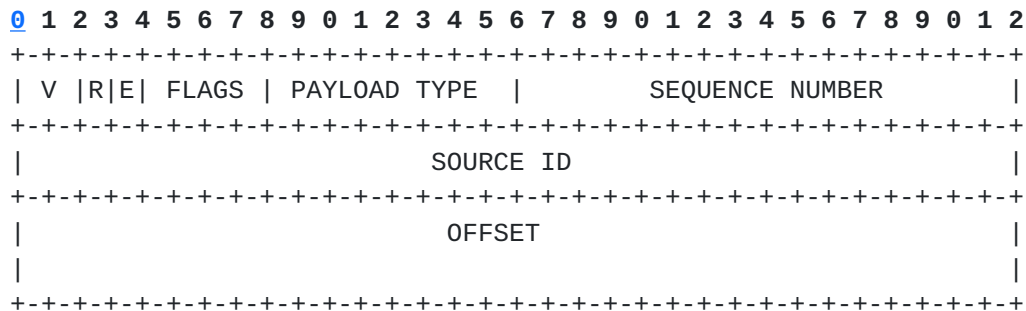
The philosophy of RMFP is in many respects similar to the one

of RTP. However, RMFP is different from RTP, as we believe that using RTP for reliable multicast is not a right approach and will not lead to a clean application design.

This draft is intended to stimulate more discussion on the one issue of a generic framing protocol for reliable multicast.

## 2. RMFP Packet Format

RMFP packet header includes common per-packet related fields. An application may include application-specific fields in a preamble header.



Version(V): 2 bits

This field identifies the version of RMFP.

Retransmission (R): 1 bit

This bit, when set, indicates that it is a retransmitted information.

Forward Error Correction (E): 1 bit

This bit, when set, indicates that FEC is used. The exact format of FEC is determined by Payload Type and its profile.

Flags: 4 bits

The flags are used for indicating significant features such as object (or ADUs) boundaries. Object boundaries can be used for multiplexing multiple objects within a single session. For example, one can multicast several files within one session.

0000: reserved

0001: start mark - the start of an object

0010: end mark - the end of an object

other: reserved

Payload Type: 8 bits

This field identifies the format of the payload and determines its interpretation by the application. Profiles will be defined for each payload type.

Sequence Number: 16 bits

The sequence number increments by one for each data packet sent. Sequence number can be used to determine packet losses (including both data packet and retransmitted packets)

Source ID: 32 bits

This field identifies the source. It is generated randomly similar to the SSRC field in RTP. It can be used to detect packet losses.

Offset: 64 bits

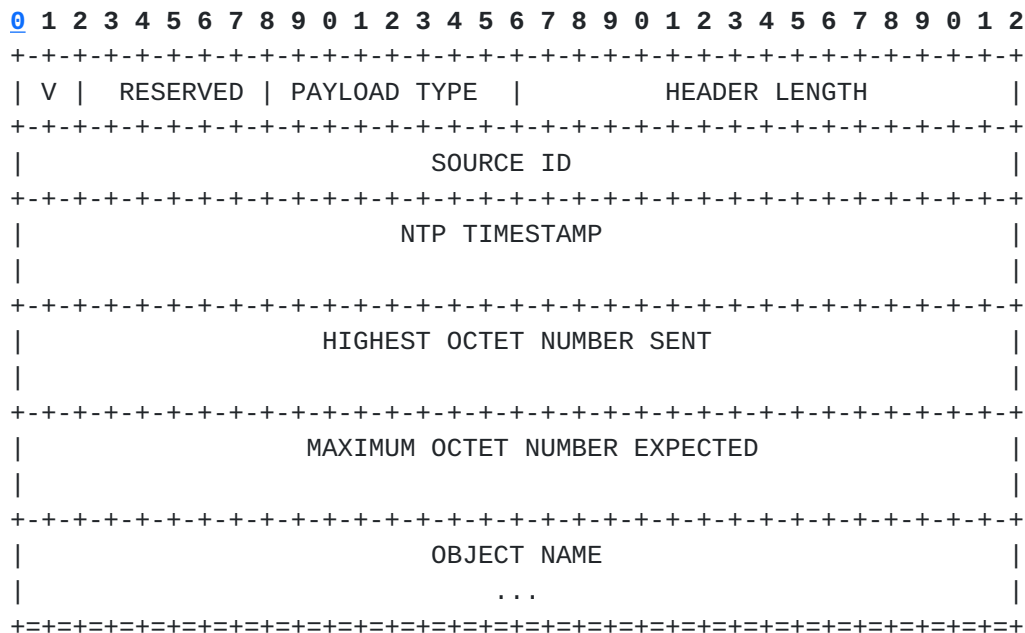
This field identifies the position of the data relative to the beginning of the session.

### 3. RMFP Control Packet Format

RMFP control packets include sender's report packets and receiver's report packets.

Sender's Report Packet

Sender's report is sent periodically by the sender about the data transmitted in the session.



Version(V): 2 bits

This field identifies the version number.

Payload Type: 8 bits

This field is set to xxx for Sender's Report Packets

Header Length: 16 bits

This field specifies the length of the header.

Source ID: 32 bits  
This field identifies the source of the sender

NTP Timestamp: 64 bits  
The NTP timestamp when the report is sent.

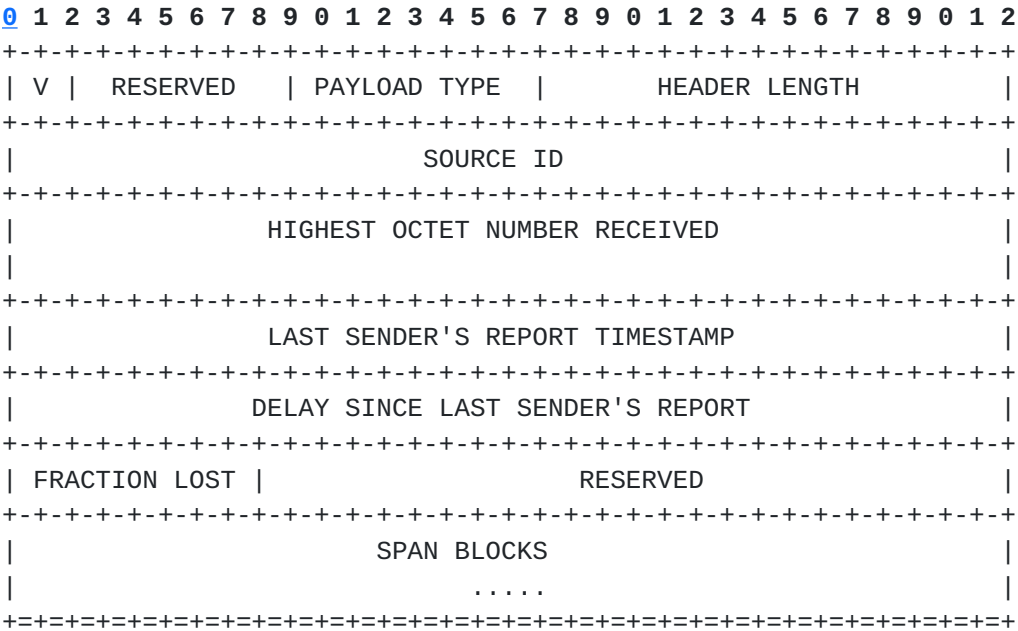
Highest Octet Number Sent: 64 bits  
This field indicates the data sent at the time the report is sent.

Maximum Octet Number Expected: 64 bits  
This field indicates the total size of the object. An application may use the information to allocate space for the session. Set to zero if the size is unknown.

Object Name  
This is a variable length field identifying the name of the object. It may be a filename, a URL, a message name etc.

Receiver's Report Packet

Receiver's report is periodically sent by the receivers to give feedback on congestion and packet losses.



Version(V): 2 bits  
This field identifies the version number.

Payload Type: 8 bits  
This field is set to xxx for Receiver's Report Packets

Header Length: 16 bits

This field specifies the length of the header.

Source ID: 32 bits

This field identifies the source of the report

Highest Octet Number Received: 64 bits

This field indicates the highest octet of the data received so far.

Last Sender's Report Timestamp: 32 bits

The middle 32 bits of the NTP Timestamp of the most recent Sender's Report

Delay Since Last Sender's Report: 32 bits

The delay, expressed in units of 1/65536 seconds, between receiving last Sender's report and sending of this report

Fraction Lost: 8 bits

The fraction of packets lost since last Sender's report, expressed as a fixed point number with the binary point at the left edge of the field. Fraction lost is the loss rate seen by the receiver. The information may be used for congestion control, error recovery purpose by the sender.

SPAN Blocks: 64 bits + 32 bits each block

Each block specifies the offset number and the length of a missing data block. The information is used for retransmission of lost packets.

#### **4. Open Issues**

Profiles for applications

Various and numerous mechanisms can be used to control reliability. Consequently, control information specific to each mechanism cannot be provided in the RMFP protocol. We propose a profile to be defined for each mechanism. The SRM profile could be based on Parnes' work on reliable RTP, for example. Other profiles could be defined for the various FEC types.

Explicit join/leave

Some reliable applications may need an explicit Join and Leave mechanism. It is not clear to us today how this facility should be provided, or if it has to be provided in RMFP (using reports or a new packet type).

#### **5. Authors's Addresses**

J Crowcroft, Zheng Wang  
{j.crowcroft, z.wang}@cs.ucl.ac.uk  
Department of Computer Science  
University College London  
Gower Street  
London  
WC1E 6BT

Atanu Ghosh  
atanu@socs.uts.EDU.AU  
School of Computing Sciences  
University of Technology  
Sydney  
PO Box 123 , Broadway  
NSW 2007  
Australia

Christophe Diot  
Christophe.Diot@sophia.inria.fr  
INRIA  
Sophia Antipolis, 2004  
route des Lucioles  
BP93 06902  
France