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## Standard Representation Of Domain Sequence draft-dhody-pce-pcep-domain-sequence-02

## Abstract

The ability to compute shortest constrained Traffic Engineering Label Switched Paths (TE LSPs) in Multiprotocol Label Switching (MPLS) and Generalized MPLS (GMPLS) networks across multiple domains has been identified as a key requirement for P2P and P2MP scenarios. In this context, a domain is a collection of network elements within a common sphere of address management or path computational responsibility such as an IGP area or an Autonomous Systems. This document specifies a standard representation and encoding of a domain sequence, which is defined as an ordered sequence of domains traversed to reach the destination domain. This document also defines new sub-objects to be used to encode domain identifiers.

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## 1. Introduction

A PCE may be used to compute end-to-end paths across multi-domain environments using a per-domain path computation technique [RFC5152]. The so called backward recursive path computation (BRPC) mechanism [RFC5441] defines a PCE-based path computation procedure to compute inter-domain constrained (G)MPLS TE LSPs. However, both per-domain and BRPC techniques assume that the sequence of domains to be crossed from source to destination is known, either fixed by the network operator or obtained by other means. For inter-domain point-to-multi-point (P2MP) tree, [PCE-P2MP-PROCEDURES] assumes the domaintree is known.

The list of domains in a point-to-point (P2P) path or a point-to-multi-point (P2MP) tree is usually a constraint in the path computation request. The PCE decouples the domain to determine the next PCE to forward the request.

According to BRPC mechanism the PCC MAY indicate the sequence of domains to be traversed using the Include Route Object (IRO) defined in [RFC5440].

This document proposes a standard way to represent and encode a domain sequence using IRO in various deployment scenarios including P2P, P2MP and Hierarchical PCE (HPCE) [PCE-HIERARCHY-FWK].

The domain sequence (the set of domains traversed to reach the destination domain) is either administratively predetermined or discovered by some means (H-PCE) that is outside of the scope of this document. Here the focus is only on a standard representation of the domain sequence in all possible scenarios.

### 1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

## 2. Terminology

The following terminology is used in this document.

ABR: OSPF Area Border Router. Routers used to connect two IGP areas.

AS: Autonomous System.

ASBR: Autonomous System Boundary Router.

BN: Boundary Node, Can be an ABR or ASBR.

BRPC: Backward Recursive Path Computation

Domain: Any collection of network elements within a common sphere of address management or path computational responsibility. Examples of domains include Interior Gateway Protocol (IGP) areas and Autonomous Systems (ASs).

Domain-Seq: An ordered sequence of domains traversed to reach the destination domain.

ERO: Explicit Route Object

H-PCE: Hierarchical PCE

IGP: Interior Gateway Protocol. Either of the two routing protocols, Open Shortest Path First (OSPF) or Intermediate System to Intermediate System (IS-IS).

IRO: Include Route Object
IS-IS: Intermediate System to Intermediate System.
OSPF: Open Shortest Path First.
PCC: Path Computation Client: any client application requesting a path computation to be performed by a Path Computation Element.

PCE: Path Computation Element. An entity (component, application, or network node) that is capable of computing a network path or route based on a network graph and applying computational constraints.

P2MP: Point-to-Multipoint
P2P: Point-to-Point
TE LSP: Traffic Engineering Label Switched Path.

## 3. Detail Description

### 3.1. Domains

A domain can be defined as a separate administrative or geographic environment within the network. A domain may be further defined as a zone of routing or computational ability. Under these definitions a domain might be categorized as an Antonymous System (AS) or an Interior Gateway Protocol (IGP) area ( as per [RFC4726] and [RFC4655]). To uniquely identify a domain in the domain sequence both AS and Area-id is important.

### 3.2. Domain-Sequence

A domain-sequence is an ordered sequence of domains traversed to reach the destination domain. In this context a Domain could be an Autonomous System (AS) or an IGP Area. Note that an AS can be further made of multiple Area.

Domain Sequence can be applied as a constraint and carried in path computation request to PCE(s). In case of HPCE [PCE-HIERARCHY-FWK] Parent PCE MAY send the domain sequence as a result in path computation reply.

In this context, ordered sequence is important, in a P2P path, the domains listed appear in the order that they are crossed. In a P2MP path, the domain tree is represented as list of domain sequences.

One main goal of the Domain-Sequence is to enable a PCE to select the next PCE to forward the path computation request based on the domain information.

A PCC or PCE MAY add an additional constraints covering which Boundary Nodes (ABR or ASBR) or Border links (Inter-AS-link) MUST be traversed while defining a domain sequence.

Thus a Domain-Sequence MAY be made up of one or more of -
o AS Number
o Area ID
o Boundary Node ID
o Inter-AS-Link Address

### 3.3. Standard Representation

The IRO (Include Route Object) [RFC5440] is an optional object used to specify a set of specified network elements that the computed path MUST traverse. [RFC5440] in its description of IRO does not constrain the sub-objects to be in a given particular order. When considering a domain sequence, the domain relative ordering is a basic criterion and, as such, this document specifies a new IRO object type.

We define a new type of IRO Object to define Domain Sequence.

IRO Object-Class is 10.
IRO Object-Type is TBD. (2 suggested value to IANA)


Sub-objects: The IRO is made of sub-objects identical to the ones defined in [RFC3209], [RFC3473], and [RFC3477], where the IRO subobject type is identical to the sub-object type defined in the related documents. Some new sub-objects related to Domain-Sequence are also added in this document.

The following sub-object types are used.

| Type | Sub-object |
| :--- | :--- |
| 1 | IPv4 prefix |
| 2 | IPv6 prefix |
| 4 | Unnumbered Interface ID |
| 32 | Autonomous system number (2 Byte) |
| TBD | Autonomous system number (4 Byte) |
| TBD | OSPF Area id |
| TBD | ISIS Area id |

[RFC3209] defines sub-objects for IPv4, IPv6 and unnumbered Interface ID, which in the context of domain-sequence is used to specify Boundary Node (ABR/ASBR) and Inter-AS-Links.
[RFC3209] also defines 2 octet AS number.

To support 4 octet AS number [RFC4893] following subobject is
defined:

```
0 1 2 3
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-++-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|L| Type | Length | Reserved |
| AS Id (4 bytes)
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
```

Since the length of Area-id is different for OSPF and ISIS, we propose different sub-objects.

For OSPF, the area-id is a 32 bit number. The Subobject looks 0

1
2
3
012345678901234567890122345678901 +-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
| L Type | Length | Reserved |
Area Id (4 bytes)
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+

The length if fixed.

For ISIS, the area-id is of variable length and thus the length of the Subobject is variable. The Area-id is as described in ISIS by ISO standard [ISO 10589]. The Length MUST be at least 4, and MUST be a multiple of 4.


The above sub-objects in various combinations can be used to encode the domain-sequence. When the domain-sequence is used as a constraint in path computation request it is carried in IRO Domain Sequence Object Type. The same sub-objects and their encoding can be used in ERO and path reply message when the domain sequence is computed from Parent PCE.

All other rules of PCEP objects and message processing is as per
[RFC5440].

### 3.4. Mode of Operation

A domain sequence IRO object constraints or defines the domains involved in a muti-domain path computation, typically involving two or more collaborative PCEs.

Consequently, a Domain-Sequence can be used:

1. by a PCE in order to discover or select the next PCE in a collaborative path computation, such as in BRPC [RFC5441];
2. by the Parent PCE to return the domain sequence when unknown, this can further be an input to BRPC procedure;
3. By a PCC (or PCE) to constraint the domains used in a H-PCE path computation, explicitly specifying which domains to be expanded;

A domain sequence can have varying degrees on granularity; it is possible to have a domain sequence composed of, uniquely, AS identifiers. It is also possible to list the involved areas for a given AS.

In any case, the mapping between domains and responsible PCEs is not defined in this document. It is assumed that a PCE that needs to obtain a "next PCE" from a domain sequence is able to do so (e.g. via administrative configuration, or discovery).

The following algorithm can be applied to select the next domain and, if need be, the PCE responsible for that domain. Note the PCC select the PCE(1) based on its own domain information.

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```
START
Get the first Sub-Object S1 from the Domain-Sequence
IF S1's Type is Area (OSPF or ISIS)
    IF S1's Domain is same as current PCE's Area
        Remove S1 from Domain-Sequence and Goto START
    ELSE
            Find the next PCE based on S1's Area within the AS
        ENDIF
ELSEIF S1's Type is AS (2 or 4 Byte)
    IF S1's Domain is same as current PCE's AS
            Remove S1 from Domain-Sequence and Goto START
        ELSE
            Get the next Sub-Object S2 from the Domain-Sequence
            IF the S2 is NULL or S2's type is AS
                Find the next PCE based on S1's Domain (AS) only
            ELSEIF S1's Type is Area
                    Find the next PCE based on S1's Domain (AS)
                    and S2's Domain (Area)
        ELSE
        ENDIF
    ENDIF
ENDIF
IF Domain-Sequence is empty or next PCE is not found
    Send PCRep with NO-Path
ENDIF
```

If the Sub-Object is of other type representing Boundary Node or Inter-As-Link, it is not used to select the next PCE, but used only while applying BRPC or any other inter-domain procedure.

### 3.5. Examples

### 3.5.1. Inter-Area Path Computation

In an inter-area path computation where ingress and egress belong to different IGP area, the domain sequence MAYBE represented using a ordered list of AREA sub-objects. AS number MAYBE skipped, as area information is enough to select the next PCE.


AS Number is 100.

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Figure 1: Inter-Area Path Computation

This could be represented as <IRO> as:


AS is optional and it MAY be skipped. PCE should be able to understand both notations.

### 3.5.2. Inter-AS Path Computation

In inter-AS path computation, where ingress and egress belong to different AS, the domain sequence is represented using an ordered list of AS sub-objects. The domain sequence MAY further include decomposed area information in AREA sub-objects.

### 3.5.2.1. Example 1

As shown in Figure 2, where AS to be made of a single area, the area subobject MAY be skipped in the domain sequence as AS is enough to uniquely identify the next domain and PCE.

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Both AS are made of Area 0.

Figure 2: Inter-AS Path Computation

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This could be represented as <IRO> as:


Area is optional and it MAY be skipped. PCE should be able to understand both notations.

### 3.5.2.2. Example 2

As shown in Figure 3, where AS 200 is made up of multiple areas and multiple domain-sequence exist, PCE MAY include both AS and AREA subobject to uniquely identify the next domain and PCE.

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Figure 3: Inter-AS Path Computation

The domain sequence can be carried in IRO as shown below:

| \| IRO | \|Sub | \|Sub | \| Sub | \| Sub | \| Sub | \|Sub |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| \|Object | \|Object | \|Object | \|Object | \|Object | \|Object | \|Object |
| \| Header | \|As 100 | \|Area 1 | \|AS 200 | \|Area 3 | \| Area 0 | \|Area 4 |
| \| | \| | \| | 1 | 1 |  |  |

Combination of both AS and Area uniquely identify a domain in the domain sequence.

Note that an Area domain identifier always belongs to the previous AS that appear before it or, if no AS sub-objects are present, it is assumed to be the current AS.

If the area information cannot be provided, PCE MAY forward the path computation request to the next PCE based on AS only. If multiple PCEs of different area domain exist, PCE MAY apply local policy to select the next PCE. Furthermore the domain sequence (list of areas within AS) in the next PCE MAYBE pre-administered or MAYBE discovered via some mechanism (ex. HPCE).

### 3.5.3. Boundary Node and Inter-AS-Link

A PCC or PCE MAY add additional constraints covering which Boundary Nodes (ABR or ASBR) or Border links (Inter-AS-link) MUST be traversed while defining a domain sequence. In which case the Boundary Node or Link MAY be encoded as a part of the domain-sequence using the existing sub-objects.

Boundary Node (ABR / ASBR) can be encoded using the IPv4 or IPv6 prefix sub-objects. The Inter-AS link can be encoded using the IPv4 or IPv6 prefix or unnumbered interface sub-objects.

For Figure 1, an $A B R$ to be traversed can be specified as:


For Figure 2, an inter-AS-link to be traversed can be specified as:


### 3.5.4. PCE serving multiple domains

A single PCE MAYBE responsible for multiple domains; for example PCE function deployed on an ABR. Domain sequence should have no impact on this. PCE which can support 2 adjacent domains can internally handle this situation without any impact on the neighboring domains.

### 3.5.5. P2MP

In case of P2MP the path domain tree is nothing but a series of Domain Sequences, as shown in the below figure:

```
D1-D3-D6, D1-D3-D5 and D1-D2-D4.
                D1
            / \
            D2 D3
            / / \
            D4 D5 D6
```


### 3.5.6. HPCE

As per [PCE-HIERARCHY-FWK], consider a case as shown in Figure 4 consisting of multiple child PCEs and a parent PCE.

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Figure 4: Hierarchical PCE

In HPCE implementation the initiator PCE - PCE(1) can request the

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parent PCE to determine the domain sequence and return in the path computation reply message (PCRep), using the ERO Object. The ERO can contain an ordered sequence of sub-object such as AS and Area (OSPF/ ISIS). In this case, the PCRep would carry the domain sequence result as:


Note that, in the case of ERO objects, no new PCEP object type is required since the ordering constraint is assumed.

### 3.5.7. Relationship to PCE Sequence

[RFC5886] and [PCE-P2MP-PROCEDURES] along with Domain Sequence introduces the concept of PCE-Sequence, where a sequence of PCEs, based on the domain sequence, should be decided and attached in the PCReq at the very beginning of path computation. An alternative would be to use domain sequences, which simplifies as explained below:

## Advantages

o All PCE must be aware of all other PCEs in all domain for PCESequence. There is no clear method for this. In domain-sequence PCE should be aware of the domains and not all the PCEs serving the domain. PCE needs to be aware of the neighboring PCEs as done by discovery protocols.
o There maybe multiple PCE in a domain, the selection of PCE should not be made at the PCC/PCE(1). This decision is made only at the neighboring PCE which is aware of state of PCEs via notification
messages.
o Domain sequence would be compatible to P2P inter-domain BRPC method as described in [RFC5441].

## 4. IANA Considerations

### 4.1. New IRO Object Type

IANA has defined a registry for Domain-Sequence.

IRO Object-Class 10
IRO Object-Type 2
4.2. Sub-Objects

IANA has defined a registry for following sub-objects.

Type Sub-object
TBD AS Number (4 Byte)
TBD OSPF Area id
TBD ISIS Area id
5. Security Considerations

This document specifies a standard representation of domain sequence, which is used in all inter-domain PCE scenarios as explained in other RFC and drafts. It does not introduce any new security considerations.
6. Manageability Considerations

TBD

## 7. Acknowledgments

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