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## A Vocabulary of Path Properties draft-enghardt-panrg-path-properties-02

#### Abstract

This document defines and categorizes information about Internet paths that an entity, such as a host, might have or want to have. This information is expressed as properties of paths between two hosts.

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## **<u>1</u>**. Introduction

Because the current Internet provides an IP-based best-effort bit pipe, hosts have little information about paths to other hosts. A Path Aware Network exposes information about one or multiple paths through the network to hosts or the network infrastructure.

It is impossible to provide an exhaustive list of path properties, as with every new technology and protocol, novel properties might become relevant. In this document, we specify a set of path properties which might be useful in the following use cases: Traffic policies, network monitoring, and path selection.

- o Traffic policies: Entities such as network operators or end users may want to define traffic policies leveraging path awareness. Such policies can allow or disallow sending traffic over specific networks or nodes, select an appropriate protocol depending on the capabilities of the on-path devices, or adjust protocol parameters to an existing path. An example of a traffic policy is a video streaming application choosing an (initial) video quality based on the achievable data rate, or the monetary cost of the link using a volume-based or flat-rate cost model. Another example is an enterprise network where all traffic has to go through a firewall, in which case the host needs to be aware of on-path firewalls.
- Network monitoring: Network operators can use path properties (e.g., measured by on-path devices), to observe Quality of Service (QoS) characteristics of recent end-user traffic, and identify potential problems with their network early on, before the enduser complains.
- Path selection: In some cases, entities can choose to use a certain path (or subset of paths) from a set of paths to achieve a specific goal. As the possible benefits of a well chosen path varies based on the goal, as a baseline, a path selection

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algorithm should aim to not perform worse than the default case most of the time. Depending on the goal, an entity may prefer paths with different properties, e.g., retrieving a small webpage as quickly as possible requires low latency paths, or retrieving a large file in a peer-to-peer network requires paths with high achievable data rate. Additionally, there may be trade-offs between path properties (e.g., latency and data rate), and entities may influence these trade-offs with their choices. A network (e.g., an AS) can adjust its path selection for internal or external routing based on the path properties. In BGP, the Multi Exit Discriminator (MED) attribute decides which path to choose if other attributes are equal: in a path aware network

network (e.g., an AS) can adjust its path selection for internal or external routing based on the path properties. In BGP, the Multi Exit Discriminator (MED) attribute decides which path to choose if other attributes are equal; in a path aware network, instead of using this single MED value, other properties such as maximum or available/expected data rate could additionally be used to improve load balancing. A host might be able to select between a set of paths, either if there are several paths to the same destination (e.g., if the host is a mobile device with two wireless interfaces, both providing a path), or if there are several destinations, and thus several paths, providing the same service (e.g., Application-Layer Traffic Optimization (ALTO) [RFC5693], an application layer peer-to-peer protocol allowing hosts a better-than-random peer selection). Care needs to be taken when selecting paths based on path properties, as path properties that were previously measured may have become outdated and, thus, useless to predict the path properties of packets sent now.

Such path properties may be relatively dynamic, e.g. current Round Trip Time, close to the origin, e.g. nature of the access technology on the first hop, or far from the origin, e.g. list of ASes traversed.

Usefulness over time is fundamentally different for dynamic and nondynamic properties. The merit of a momentary measurement of a dynamic path property diminishes greatly as time goes on, e.g. the merit of an RTT measurement from a few seconds ago is quite small, while a non-dynamic path property might stay relevant, e.g. a NAT can be assumed to stay on a path during the lifetime of a connection, as the removal of the NAT would break the connection.

Non-dynamic properties are further separated into (local) domain properties related to the first few hops of the connection, and backbone properties related to the remaining hops. Domain properties expose a high amount of information to hosts and strongly influence the connection behavior while there is little influence and information about backbone properties.

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Path Properties

Dynamic properties are not separated into domain and backbone properties, since most of these properties are defined for a complete path and it is difficult and seldom useful to define them on part of the path. There are exceptions such as dynamic wireless access properties, but these do not justify separation into different categories.

This document addresses the first of the questions in Path-Aware Networking [I-D.irtf-panrg-questions], which is a product of the PANRG in the IRTF.

### **2**. Terminology

- Node: An entity which processes packets, e.g., sends, receives, forwards, or modifies them.
- Host: A node that processes packets that are explicitly addressed to itself.
- Router: A node that processes packets that are not explicitly addressed to itself.
- Link: A medium or communication facility that connects two or more nodes with each other and enables them to exchange packets. A link can be physical, e.g., a WiFi network which connects an Access Point to stations, or virtual, e.g., a virtual switch which connects two virtual machines hosted on the same physical machine.

Path element: Either a node or a link.

- Path: A sequence of adjacent path elements, alternating between nodes and links, starting and ending with a host. A path can be viewed as an abstraction on a specific layer, omitting lower layer path elements. For example, a router implementing IPv6 may be a path element on a path when considering the network layer. If this router does not implement transport layer functionality, it is hidden when a higher layer, such as the transport or application layer, is considered. In the case of multicast or broadcast, a single packet may be sent over multiple paths at once - one path for each combination of sending and receiving host.
- Subpath: Given a path, a subpath is a sequence of adjacent path elements of this path, starting and ending with a node.
- Flow: One or multiple packets which are traversing the same subpath or path. For example, a flow can consist of all packets sent within a TCP session with the same five-tuple between two hosts, or it can consist of all packets sent on the same physical link.

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- Property: A trait of one or a sequence of path elements, or a trait of a flow with respect to one or a sequence of path elements. An example of a link property is the maximum data rate that can be sent over the link. An example of a node property is the administrative domain that the node belongs to. An example of a property of a flow with respect to a subpath is the aggregated one-way delay of the flow being sent from one node to another node over a subpath. A property is thus described by a tuple containing the sequence of path elements, the flow or an empty set if no packets are relevant for the property, the name of the property (e.g., maximum data rate), and the value of the property (e.g., 1Gbps).
- Aggregated property: A collection of multiple values of a property into a single value, according to a function. A property can be aggregated over multiple path elements (i.e., a path), e.g., the MTU of a path as the minimum MTU of all links on the path, over multiple packets (i.e., a flow), e.g., the median one-way latency of all packets between two nodes, or over both, e.g., the mean of the queueing delays of a flow on all nodes along a path. The aggregation function can be numerical, e.g., median, sum, minimum, it can be logical, e.g., "true if all are true", "true if at least 50\% of values are true", or an arbitrary function which maps multiple input values to an output value.
- Measured property: A property that is observed for a specific path element or path, e.g., using measurements. For example, the oneway delay of a specific packet can be measured.
- Estimated property: An approximate calculation or judgment of the value of a property. For example, an estimated property may describe the expected median one-way latency of packets sent on a path within the next second. An estimated property includes the reliability of the estimate. The notion of reliability depends on the property. For example, it may be the confidence level and interval for numerical properties or the likelihood that a property holds for non-numerical properties.

#### <u>3</u>. Domain Properties

Domain path properties relate to path elements within the first hop or the first few hops, which are usually in the same administrative domain as a host considering them.

Due to the potential physical proximity and pre-existing trust or contractual relationships between hosts and path elements within the same domain, domain properties may be more accessible to the host than other properties.

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Furthermore, hosts may be able to influence both which domain they are in and which path elements in this domain to connect to, and they may be able to influence the properties of path elements within this domain. For example, a user might select between multiple potential adjacent links by selecting between multiple available WiFi Access Points. Or when connected to an Access Point, the user may move closer to enable their device to use a different access technology, potentially increasing the data rate available to the device. Another example is a user changing their data plan to reduce the Monetary Cost to transmit a given amount of data across a network.

- Access Technology: The physical or link layer technology used for transmitting or receiving a flow on one or multiple path elements in the same domain. The Access Technology may be given in an abstract way, e.g., as a WiFi, Wired Ethernet, or Cellular link. It may also be given as a specific technology, e.g., as a 2G, 3G, 4G, or 5G cellular link, or an 802.11a, b, g, n, or ac WiFi link. Other path elements relevant to the access technology may include on-path devices, such as elements of a cellular backbone network. Note that there is no common registry of possible values for this property.
- Monetary Cost: The price to be paid to transmit a specific flow across a subpath.

#### **<u>4</u>**. Backbone Properties

Backbone path properties relate to path elements not within the same domain as a host considering them, thus, in the backbone from the host's point of view.

Typically, backbone properties are less accessible to a host than domain properties, due to the potential increased distance and the lack of pre-existing trust or contractual relationship.

Additionally, hosts are less likely to be able to influence which path elements form their path in the backbone, as well as their properties.

Some path properties relate to the entire path or to subpaths, part of which often lies outside of a host's domain. Thus, such properties are listed as Backbone Properties.

Presence of a certain network function on the path: Indicates that a node performs a certain network function on a flow, e.g., whether the node acts as a proxy, as a firewall, or performs Network Address Translation (NAT). This node may be either in the same domain as the host or in a different domain, i.e., the backbone.

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- Administrative Entity: The administrative entity, e.g., the AS, to which a path element or subpath belongs.
- Disjointness: For a set of two paths, the number of shared path elements can be a measure of intersection (e.g., Jaccard coefficient, which is the number of shared elements divided by the total number of elements). Conversely, the number of non-shared path elements can be a measure of disjointness (e.g., 1 - Jaccard coefficient). A multipath protocol might use disjointness of paths as a metric to reduce the number of single points of failure.
- Path MTU: The maximum size, in octets, of an IP packet that can be transmitted without fragmentation on a subpath.
- Transport Protocols available: Whether a specific transport protocol can be used to establish a connection over a path or subpath. A host may cache its knowledge about recent successfully established connections using specific protocols, e.g., a QUIC connection, or an MPTCP subflow.
- Protocol Features available: Whether a specific protocol feature is available over a path or subpath, e.g., Explicit Congestion Notification (ECN), or TCP Fast Open.

#### 5. Dynamic Properties

Dynamic path properties relate to the transmission of an individual packet or of a flow over a subpath. Properties related to a path element which constitutes a single layer 2 domain are abstracted from the used physical and link layer technology, similar to [<u>RFC8175</u>].

Typically, Dynamic Properties can be measured or approximated, and might be made available in an aggregated form, such as averages or minimums. Dynamic Path Properties can be measured by the host itself or by a different entity. See [<u>ANRW18-Metrics</u>] for a discussion of how to measure some dynamic path properties at the host.

Some dynamic properties are defined in different directions for the same path element, e.g., for transmitting and receiving packets.

- Maximum Data Rate (Transmit/Receive): The theoretical maximum data rate, in bits per second, that can be achieved on a link, subpath, or path, for receiving or transmitting traffic.
- Current Data Rate (Transmit/Receive): The data rate, in bits per second, at which a link is currently receiving or transmitting traffic.

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Latency: The time delay between a node sending a packet and a different node on the same path receiving the same packet.

Latency variation: The variation of the latency within a flow.

- Packet Loss: The percentage of packets within a flow which are sent by one node, but not received by a different node.
- Congestion: Whether a protocol feature such as ECN has provided information that there currently is congestion on a path.

### <u>6</u>. Security Considerations

If nodes are basing policy or path selection decisions on path properties, they need to rely on the accuracy of path properties that other devices communicate to them. In order to be able to trust such path properties, nodes may need to establish a trust relationship or be able to verify the authenticity, integrity, and correctness of path properties received from another node.

### 7. IANA Considerations

This document has no IANA actions.

### 8. Informative References

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Authors' Addresses

Theresa Enghardt TU Berlin

Email: theresa@inet.tu-berlin.de

Cyrill Kraehenbuehl ETH Zuerich

Email: cyrill.kraehenbuehl@inf.ethz.ch