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LISP EID Anonymity
draft-farinacci-lisp-eid-anonymity-02

Abstract

This specification will describe how ephemeral LISP EIDs can be used to create source anonymity. The idea makes use of frequently changing EIDs much like how a credit-card system uses a different credit-card numbers for each transaction.

Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [[RFC2119](#)].

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LISP EID Anonymity

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Table of Contents

1.	Introduction	2
2.	Definition of Terms	3
3.	Overview	3
4.	Design Details	4
5.	Other Types of Ephemeral-EIDs	4
6.	Interworking Considerations	5
7.	Multicast Considerations	5
8.	Performance Improvements	5
9.	Security Considerations	6
10.	IANA Considerations	6
11.	References	6
11.1.	Normative References	6
11.2.	Informative References	7
Appendix A.	Acknowledgments	8
Appendix B.	Document Change Log	8
B.1.	Changes to draft-farinacci-lisp-eid-anonymity-02	8
B.2.	Changes to draft-farinacci-lisp-eid-anonymity-01	8
B.3.	Changes to draft-farinacci-lisp-eid-anonymity-00	8
	Authors' Addresses	8

[1.](#) Introduction

The LISP architecture [[RFC6830](#)] specifies two namespaces, End-Point IDs (EIDs) and Routing Locators (RLOCs). An EID identifies a node in the network and the RLOC indicates the EID's topological location. Typically EIDs are globally unique so a end-node system can connect to any other end-node system on the Internet. Privately used EIDs are allowed when scoped within a VPN but must always be unique within that scope. Therefore, address allocation is required by network administration to avoid address collisions or duplicate address use. In a multiple namespace architecture like LISP, typically the EID

will stay fixed while the RLOC can change. This occurs when the EID is mobile or when the LISP site the EID resides in changes its connection to the Internet.

LISP creates the opportunity where EIDs are fixed and won't change. This can create a privacy problem more so than what we have on the Internet today. This draft will examine a technique to allow an end-node system to use a temporary address. The lifetime of a temporary address can be the same as a lifetime of an address in use today on the Internet or can have traditionally shorter lifetimes, possibly on the order of a day or even change as frequent as new connection attempts.

[2.](#) Definition of Terms

Ephemeral-EID - is an IP address that is created randomly for use for a temporary period of time. An Ephemeral-EID has all the properties of an EID as defined in [[RFC6830](#)]. Ephemeral-EIDs are not stored in the Domain Name System (DNS) and should not be used in long-term address referrals.

Client End-Node - is a network node that originates and consumes packets. It is a system that originates packets or initiates the establishment of transport-layer connections. It does not offer services as a server system would. It accesses servers and attempts to do it anonymously.

[3.](#) Overview

A client end-node can assign its own ephemeral EID and use it to talk to any system on the Internet. The system is acting as a client where it initiates communication and desires to be an inaccessible resource from any other system. The ephemeral EID is used as a destination address solely to return packets to resources the ephemeral EID connects to.

Here is the procedure a client end-node would use:

1. Client end-node desires to talk on the network. It creates and assigns an ephemeral-EID on any interface.

2. If the client end-node is a LISP xTR, it will register the ephemeral-EID with a globally routable RLOC. If the client end-node is not a LISP xTR, it can send packets on the network where a LISP router xTR will register the ephemeral-EID with its RLOC.
3. The client end-node originates packets with a source address equal to the ephemeral-EID and will receive packets addressed to the ephemeral-EID.
4. When the client end-node decides to stop using the ephemeral-EID, it will deregister it from the mapping system and create and

assign a new ephemeral-EID, or decide to configure a static global address, or participate in DHCP to get assigned a leased address.

Note that the ephemeral-EID can be mobile just like any other EID so if it is initially registered to the mapping system with one or more RLOCs, later the RLOC-set can change as the ephemeral-EID roams.

[4.](#) Design Details

This specification proposes the use of the experimental LISP EID-block 2001:5::/32 when IPv6 is used. See IANA Considerations section for a specific sub-block allocation request. When IPv4 is used, the Class E block 240.0.0.0/4 is being proposed.

The client end-node system will use the rest of the host bits to allocate a random number to be used as the ephemeral-EID. The EID can be created manually or via a programmatic interface. When the EID address is going to change frequently, it is suggested to use a programmatic interface. The probability of address collision is unlikely for IPv6 EIDs but could occur for IPv4 EIDs. A client end-node can create a ephemeral-EID and then look it up in the mapping system to see if it exists. If the EID exists in the mapping system, the client end-node can attempt creation of a new random number for the ephemeral-EID. See [Section 8](#) where ephemeral-EIDs can be preallocated and registered to the mapping system before use.

When the client end-node system is co-located with the RLOC and acts as an xTR, it should register the binding before sending packets.

This eliminates a race condition for returning packets not knowing where to encapsulate packets to the ephemeral-EID's RLOCs. See [Section 8](#) for alternatives for fixing this race condition problem. When the client end-node system is not acting as an xTR, it should send some packets so its ephemeral-EID can be discovered by an xTR which supports EID-mobility [[I-D.portoles-lisp-eid-mobility](#)] so mapping system registration can occur before the destination returns packets. When the end-node system is acting as an xTR, the EID and RLOC-set is co-located in the same node. So when the EID is created, the xTR can register the mapping versus waiting for packet transmission.

[5.](#) Other Types of Ephemeral-EIDs

When IPv6 Ephemeral-EIDs are used, an alternative to a random number can be used. For example, the low-order bits of the IPv6 address could be a cryptographic hash of a public-key. Mechanisms from [[RFC3972](#)] could be used for EIDs. Using this approach allows the sender with a hashed EID to be authenticated. So packet signatures

can be verified by the corresponding public-key. When hashed EIDs are used, the EID can change frequently as rekeying may be required for enhanced security.

[6.](#) Interworking Considerations

If a client end-node is communicating with a system that is not in a LISP site, the procedures from [[RFC6832](#)] should be followed. The PITR will be required to originate route advertisements for the ephemeral-EID sub-block [[I-D.draft-ietf-lisp-eid-block](#)] so it can attract packets sourced by non-LISP sites destined to ephemeral-EIDs. However, in the general case, the coarse block from [[I-D.draft-ietf-lisp-eid-block](#)] will be advertised which would cover the sub-block. For IPv4, the 240.0.0.0/4 must be advertised into the IPv4 routing system.

[7.](#) Multicast Considerations

A client end-node system can be a member of a multicast group fairly easily since its address is not used for multicast communication as a receiver. This is due to the design characteristics of IGMP [[RFC3376](#)] [[RFC2236](#)] [[RFC1112](#)] and MLD [[RFC2710](#)] [[RFC3810](#)].

When a client end-node system is a multicast source, there is ephemeral (S,G) state that is created and maintained in the network via multicast routing protocols such as PIM [[RFC4602](#)] and when PIM is used with LISP [[RFC6802](#)]. In addition, when [[I-D.draft-ietf-lisp-signal-free-multicast](#)] is used, ephemeral-EID state is created in the mapping database. This doesn't present any problems other than the amount of state that may exist in the network if not timed out and removed promptly.

However, there exists a multicast source discovery problem when PIM-SSM [[RFC4607](#)] is used. Members that join (S,G) channels via out of band mechanisms. These mechanisms need to support ephemeral-EIDs. Otherwise, PIM-ASM [[RFC4602](#)] or PIM-Bidir [[RFC5015](#)] will need to be used.

8. Performance Improvements

An optimization to reduce the race condition between registering ephemeral-EIDs and returning packets as well as reducing the probability of ephemeral-EID address collision is to preload the mapping database with a list of ephemeral-EIDs before using them. It comes at a expense of rebinding all of registered ephemeral-EIDs when there is an RLOC change. There is work in progress to consider adding a level of indirection here so a single entry gets the RLOC update and the list of ephemeral-EIDs point to the single entry.

9. Security Considerations

When LISP-crypto [[I-D.draft-ietf-lisp-crypto](#)] is used the EID payload is more secure through encryption providing EID obfuscation of the ephemeral-EID as well as the global-EID it is communicating with. But the obfuscation only occurs between xTRs. So the randomness of a ephemeral-EID inside of LISP sites provide a new level of privacy.

10. IANA Considerations

This specification is requesting the sub-block 2001:5:ffff::/48 for ephemeral-EID usage.

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[Appendix A](#). Acknowledgments

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[Appendix B](#). Document Change Log

[RFC Editor: Please delete this section on publication as RFC.]

[B.1](#). Changes to [draft-farinacci-lisp-eid-anonymity-02](#)

- o Posted April 2017.
- o Added section describing how ephemeral-EIDs can use a public key hash as an alternative to a random number.
- o Indciate when an EID/RLOC co-located, that the xTR can register the EID when it is configured or changed versus waiting for a packet to be sent as in the EID/RLOC separated case.

[B.2](#). Changes to [draft-farinacci-lisp-eid-anonymity-01](#)

- o Posted October 2016.
- o Update document timer.

[B.3](#). Changes to [draft-farinacci-lisp-eid-anonymity-00](#)

- o Posted April 2016.
- o Initial posting.

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