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**Authoritative DNS-over-TLS Operational Considerations**  
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**Abstract**

DNS over TLS (DoT) has been gaining attention, primarily as a means of communication between stub resolvers and recursive resolvers. There have also been discussions and experiments involving the use of DoT to communicate with authoritative nameservers (Authoritative DNS over TLS or "ADoT"), including communication between recursive and authoritative resolvers. However, we have identified a number of operational concerns with ADoT. These operational concerns need to be addressed prior to ADoT's deployment at scale by DNS operators in order to maintain the stability and resilience of the global DNS. The document also provides some suggested next steps to advance the operator community's understanding of ADoT's operational impact.

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## **1. Introduction**

This is an operational considerations document that focuses on the factors operators need to consider when implementing Authoritative DNS over TLS. An evaluation of the merits of DNS over TLS are beyond the scope and intent of this document.

Typically, DNS communication between stub resolvers, recursive resolvers, and authoritative servers is not encrypted. Some argue that this can pose a privacy challenge for Internet users, because their access to named network resources can potentially be tracked through their DNS communication. In principle, any network element along the path between the user and resolving or authoritative nameservers could observe this unencrypted traffic. Many of these concerns are addressed in [[RFC7626](#)].

[RFC8310] proposes using DNS over TLS (DoT) in order to encrypt DNS traffic.

Historically, much of the work on DNS encryption has focused on the stub-to-recursive path as the recursive-to-authoritative server path does not leak user specific information. However, with the increased deployment of EDNS0 Client Subnet [[RFC7871](#)], recursive-to-authoritative encryption is becoming an area of interest. Therefore, this document focuses on the recursive-to-authoritative aspect of DoT and we use the term Authoritative DNS over TLS (ADoT) in order to differentiate it from the stub-to-recursive path.

The addition of ADoT, while providing encryption for DNS communication, also introduces other factors that might impact the stability and resiliency of authoritative nameserver operations which may have been optimized for unencrypted DNS, often focusing on UDP transport.

The objective of this document is to try to describe the problem space, make suggestions about solutions, and propose next steps that can help inform both recursive and authoritative operators on how to assess and address the challenges posed by ADoT deployment.

## **1.1. Background and Motivation**

### **1.1.1. Why operational considerations are so important for ADoT**

The main concerns for most authoritative operators are the stability, resiliency, scalability, and performance of their platforms. These concerns need to be weighed against the benefits, if any, offered to the end user by encrypting DNS queries to the authoritative servers and the associated benefit of protection from modifications in transit.

The communication between the recursive-to-authoritative server is less able to be associated with a particular user than information communicated along the stub-to-recursive path given that the originating IP address is that of the recursive resolver, not the user's, and the QNAME is potentially minimized. Therefore, initial deployments of ADoT may offer an immediate expansion of the attack surface (additional port, transport protocol, and computationally expensive crypto operations for an attacker to exploit) while potentially providing limited benefit to end users.

### **1.1.2. Other considerations related to ADoT**

As resolvers add encryption on the client-to-recursive path, they may also change the way they handle data on the recursive-to-authoritative path. This is expressed in Mozilla Trusted Recursive Resolver (TRR) requirements [[1](#)], for example, which require participating resolvers to perform QNAME minimization [[RFC7816](#)], and

TRR requirement #6, which forbids the EDNS0 Client subnet (ECS) from being propagated unless the recursive-to-authoritative path is encrypted.

The latter requirement may have the possible unintended consequence of reducing the authoritative name servers' ability to provide a best response to DNS queries, until such time as they deploy DNS encryption.

Given that recursive resolvers should be configured to prevent ECS transmission to root, top-level, and effective top-level domain (TLD) servers [\[RFC7871\] section 12.1 \[2\]](#) - the ECS encryption requirement motivates consideration of authoritative DNS encryption below these levels.

At the higher levels, techniques such as QNAME minimization and Aggressive Use of DNSSEC-Validated Cache [\[RFC8198\]](#) arguably provide an alternate path toward mitigating the risk of disclosure of sensitive information without the operational risk of DNS encryption.

Resolver requirements may change as the understanding of DNS encryption options evolve, but in the meantime, they provide motivation for authoritative name server operators to weigh the risks and benefits of DNS encryption, hence the importance of understanding these operational considerations.

## **[2. Terminology](#)**

### **[2.1. Requirements Language](#)**

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [BCP 14 \[RFC2119\] \[RFC8174\]](#) when, and only when, they appear in all capitals, as shown here.

### **[2.2. Definitions](#)**

ADoT: Authoritative DNS over TLS - the use of DNS over TLS (DoT) when communicating with an authoritative DNS server. eg: between the recursive resolver and the authoritative DNS server (recursive-to-authoritative). We use DoT when discussing generic DNS over TLS or when referring to encrypted communication between stub and recursive resolver.

Attack Surface: The sum of attack vectors where an unauthorized user (attacker) can try to enter or extract data from the environment or compromise a service via resource starvation.

Authoritative Operator: An operator of an authoritative DNS server.

CDN: Content Delivery Network - distributed network of servers which

proxy traffic between content providers and end users in order to provide high availability and high performance.

ECS: EDNS0 Client Subnet [[RFC7871](#)] - an extension to EDNS0 where the client's subnet is included in the DNS query, intended to provide a hint to authoritative servers who may wish to provide different answers in an attempt to provide higher performance for end users based on their network location.

EPSK: External Pre-Shared Key - TLS 1.3 [[RFC8446](#)] uses the same PSK extension for keys established both during handshake (resumption PSK) and keys established externally. The EPSK acronym was introduced in draft [I-D.[draft-wood-tls-external-psk-importer-00](#)] in order to disambiguate External vs Resumption PSKs.

Performance:

- o QRTT: Query Round-Trip Time - the time it takes between sending a query and receiving a response.
- o Best Response - whether or not the authoritative server, if dynamic responses are used as they are in CDNs, are able to determine or infer location and provide the most local response. It is a key part of the end-to-end performance for end users to get not just \_an\_ answer quickly but to get the best and most local answer.
- o System Performance - the cost in system resources such as CPU/IO/Memory.
- o Queries Per Second (QPS) - the maximum number of simultaneous queries that a DNS server can handle, on a per second basis.

Authoritative Server - see [[RFC8499](#)]

Recursive Resolver - see [[RFC8499](#)]

Stub Resolver - see [[RFC8499](#)]

TLS: - see [[RFC8446](#)]

SUDN: Special-Use Domain Names - see [[RFC6761](#)]

### **3. Key Issues and Questions**

#### **3.1. Signaling Support for ADoT**

[RFC8310] does not define a method for a nameserver to advertise its support of DoT other than to have the client make a connection attempt to the default port of 853. The extra round-trip to check for ADoT support imposes a penalty for clients and resolvers that either do not remember the nameserver or have not communicated with

that nameserver before. The extra round-trip required may lead some implementers down a similar path to happy eyeballs [[RFC8305](#)] which, in the case of DNS, would send the same query over both encrypted and un-encrypted channels at the same time. A happy eyeballs type approach, which we'll call "leaky resolvers", would defeat the purpose of the encryption protection for the testing query, but may enable subsequent queries to be sent over a private channel with the first query being subject to on-path adversaries. An implementation could use some constant query string as a test query. However any query included in the set of queries comprising the iterative resolution for a QNAME first sent over an encrypted channel that leaks the original stub QNAME, SHOULD NOT be used.

### **[3.2.](#) Port number**

[RFC7858] [section 3](#) [3] indicates that port 853 MUST be used for session establishment unless otherwise negotiated and configured by both the client and server. In the stub-to-recursive connection, changing the port is something that can be done at stub configuration time however, managing this negotiation between the recursive-to-authoritative server is not scalable or standardized. The scalability problem is due to the fact that recursive resolvers communicate with thousands of authoritative servers, therefore port/service discovery for each of these authoritatives becomes difficult.

Static use of a pre-defined port provides on-path adversaries the ability to more easily drop or manipulate traffic intended for that port, possibly triggering resolvers to downgrade a connection back to a traditional DNS query, eliminating the encryption protections. This attack is more likely to happen on the stub-to-recursive connection but is also a possible threat for recursive-to-authoritative connections.

### **[3.3.](#) TLS version**

Implementers of ADoT should read, understand, and follow the guidance provided in [BCP195](#) [4], also known as [[RFC7525](#)], when deploying DoT on their platforms. At the time of writing, [[RFC7525](#)] did not include coverage for TLS 1.3. However, TLS 1.3 should be included in the document that obsoletes this BCP. Until this happens, TLS 1.3 SHOULD be preferred over TLS 1.2, as 1.3 offers both security and performance enhancements. Additionally, operators should monitor TLS version issues and cipher suite vulnerabilities for the version of TLS that their platforms offer.

In the absence of any widespread ADoT deployments, it is easier to limit TLS version 1.3 or greater. The absence of widespread adoption also allows the IETF to create and enforce standards/policies that ensure TLS versions are kept current going forward.

### **[3.4.](#) Resumptions**

TLS resumption allows clients and servers to use information from a previously established session in order to bootstrap the cryptographic state while avoiding a full handshake. The resumption mechanism is redesigned in TLS 1.3 [\[RFC8446\] section 2.2 \[5\]](#) and [section 2.3 \[6\]](#), eliminating both [\[RFC5077\]](#) session tickets and session ID resumption.

Resumption improves both connection and resource (socket and CPU) efficiency, therefore operators SHOULD allow for TLS resumption. However, special consideration should be given to 0-RTT resumption as it is vulnerable to replay attacks [\[RFC3552\]](#) see [Section 3.3.1 \[7\]](#). The replay attack may not be as important for DNS, as DNS queries are generally idempotent. However consideration should be given to possible side-channel attacks [\[8\]](#).

### **[3.5.](#) Operational Monitoring**

Many operators use external passive monitors in order to understand the health and performance of their infrastructure. Infrastructure monitoring is also often done to retain a copy of traffic for forensic purposes - such as the BIND "packet of death" [\[9\]](#) scenario. These legacy monitoring systems may break with the use of TLS 1.3. Therefore alternatives may need to be deployed/developed in order to maintain effective operational performance and security monitoring functionality.

A number of solutions have been suggested:

- o TLS Security and Data Center Monitoring: Searching for a Path Forward [\[10\]](#)
- o TLS 1.3: Will Your Network Monitoring Go Blind? [\[11\]](#)

### **[3.6.](#) Architecture**

Operators often reconfigure their architectural designs to best deliver a new product offering or service. Operators should consider the following design alternatives for the new ADoT service:

- o Operators should consider segregating ADoT addresses from traditional DNS over UDP/TCP to enable better attack mitigation, better service monitoring, less service interference, and more stability.
- o Operators should weigh the pros/cons of using a TLS proxy vs direct client-to-host connection. In case of ADoT, the client is most likely a recursive resolver and the host is the authoritative host server.

### **[3.7.](#) Socket efficiency/tuning considerations**

Operators can realize substantial gains in client session establishment and improve overall RTT by tuning sockets setting for best use-case efficiency.

For the ADoT use case, operators should consult [\[RFC7766\] section 6.2 \[12\]](#) and minimally consider the following:

- o Optimal number of persistent connections - consideration should be given to the number of persistent connections maintained for both the recursive resolvers and authoritative servers
- o Optimal read/write buffer size
- o Optimal session timeout
- o Optimal close wait state time
- o Optimal connection time/timeout

### **3.8. Post-Quantum Security**

Given that ADoT deployments will likely have a long lifetime and are being introduced in an era where post-quantum security is now an important design consideration, it is prudent to consider how protections against quantum computers might be integrated into the deployments.

[I-D.[draft-hoffman-c2pq-05](#)] outlines the threat quantum computing presents to classical cryptographic algorithms.

External Pre-Shared Keys (EPSKs) may be less vulnerable to quantum attacks. A proposed approach to combining EPSKs and certificates in TLS is described in [I-D.[draft-housley-tls-tls13-cert-with-extern-psk-03](#)].

## **4. Suggestions for further research and development**

### **4.1. Required studies and analysis**

Unlike stub-to-recursive DNS communication, authoritative nameservers affect users in ways that end users cannot avoid or work around. In the event that all authoritative servers for a zone are unreachable, the zone becomes globally unavailable. Hence, in order to preserve stability and resiliency of authoritative nameservers when deploying ADoT, more empirical studies and analysis MUST be conducted. The following list is a minimal set of studies and considerations that need to be conducted/addressed in order to maintain authoritative stability and resilience.

- o Attack vectors and mitigation: consider the new adversarial powers enabled by ADoT - types of attacks and denial of service, or other security challenges that are created with the addition of ADoT to



authoritative nameservers.

- o Traffic: consider how traffic patterns to authoritative nameservers change with the introduction of ADoT and how these traffic patterns change when the parameters of the service are changed; e.g. persistent connection lifetime, TLS connection parameters, use of TLS session tickets [[RFC5077](#)] or Pre-Shared Key extension in TLS 1.3 [[RFC8446](#)] [section 2.2](#) [[13](#)]. Consider how these traffic pattern changes will affect the architecture and infrastructure for authoritative operators.
- o ADoT capacity and footprint expansion: consider how common scaling techniques impact authoritative operators; e.g. anycast, load balancing, custom hardware.
- o DTLS/UDP - consider if there is any reason to implement DTLS given that we lose the benefit of pipelining requests and must drop back to TLS/TCP in the case of fragmentation.

It is critical to conduct large-scale measurements of DNS infrastructure in order to quantify some of the scalability issues. While these tests may be performed initially in a controlled lab environment, the public Internet is fundamentally more variable. Therefore, global testing at scale on the Internet MUST also be conducted in order to understand and measure potential issues which must be overcome before full global deployment can occur.

#### **[4.2.](#) Authoritative DNS over TLS (ADoT) Profile**

Profiles can be used as a mechanism to help mitigate operational concerns over increased attack surface by restricting features such as computationally expensive processes, insecure ciphers, general starvation vectors, or other features that may limit operational performance.

Therefore, an ADoT application profile draft, taking into account the conclusions of required studies and analysis, may help assuage some of the concerns raised in this document.

### **[5.](#) Security Considerations**

In addition to the applicable security considerations described in RFCs [[RFC7626](#)] and [[RFC8310](#)], considerations focused on future deployment of quantum computers are described in Post-Quantum Security ([Section 3.8](#)). Additional considerations associated with ADoT are TBD based on working group discussions.

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## **6.2. URIs**

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- [2] <https://tools.ietf.org/html/rfc7871#section-12.1>
- [3] <https://tools.ietf.org/html/rfc7858#section-3>
- [4] <https://tools.ietf.org/html/bcp195>
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- [13] <https://tools.ietf.org/html/rfc8446#section-2.2>

## **Appendix A. Acknowledgements**

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## **Appendix B. Change Log**

RFC EDITOR: PLEASE REMOVE THE THIS SECTION PRIOR TO PUBLICATION.

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pre-00

- o Initial draft.

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