Security Considerations for RFC5011 Publishers
draft-hardaker-rfc5011-security-considerations-03

Abstract

   This document describes the math behind the minimum time-length that
   a DNS zone publisher must wait before using a new DNSKEY to sign
   records when supporting the RFC5011 rollover strategies.

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1. Introduction

   RFC5011 [RFC5011] defines a mechanism by which DNSSEC validators can
   extend their list of trust anchors when they've seen a new key
   published in a zone. However, RFC5011 [intentionally] provides no
   guidance to the publishers of DNSKEYs about how long they must wait
   before switching to the newly published key for signing records.
   Because of this lack of guidance, zone publishers may derive
   incorrect assumptions about safe usage of the RFC5011 DNSKEY
   advertising and rolling process. This document describes the minimum
   security requirements from a publishers point of view and is intended
   to compliment the guidance offered in RFC5011 (which is written to
   provide timing guidance solely to the Validating Resolvers point of
   view).

   To verify this lack of understanding is wide-spread, the authors
   reached out to 5 DNSSEC experts to ask them how long they thought
   they must wait before using a new KSK that was being rolled according
   to the 5011 process. All 5 experts answered with an insecure value,
   and thus we have determined that this lack of operational guidance is
   causing security concerns today. We hope that this document will
   rectify this understanding and provide better guidance to zone
   publishers that wish to make use of the RFC5011 rollover process.

   One important note about ICANN's upcoming 2017 KSK rollover plan for
   the root zone: the timing values chosen for rolling the KSK in the
   root zone appear completely safe, and are not in any way affected by
   the timing concerns introduced by this draft

1.1. Requirements notation

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT",
"SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this
document are to be interpreted as described in [RFC2119].

2. Background

The RFC5011 process describes a process by which a Validating
Resolver may accept a newly published KSK as a trust anchor for
validating future DNSSEC signed records. This document augments that
information with additional constraints, as required from the DNSKEY
publication point of view. Note that it does not define any other
operational guidance or recommendations about the RFC5011 process
from a publication point of view and restricts itself to solely the
security and operational ramifications of switching to a new key too
soon. Failure of a DNSKEY publisher to follow the minimum
recommendations associated with this draft will result in potential
denial-of-service attack opportunities against validating resolvers.

3. Terminology

Trust Anchor Publisher The entity responsible for publishing a
   DNSKEY that can be used as a trust anchor.

4. Timing associated with RFC5011 processing

RFC5011's process of safely publishing a new key and then making use
of that key falls into a number of high-level steps:

1. Publish a new DNSKEY in the zone but continue to sign with the
    old one.

2. Wait a period of time.

3. Begin using the new DNSKEY to sign the appropriate resource
    records.

4. Optionally mark the older DNSKEY as revoked and publish the
    revoked key.

This document discusses step 2 of the above process. Some
interpretations of RFC5011 have erroneously determined that the wait
time is equal to RFC5011's "hold down time".

This document describes an attack based on this (common) erroneous
belief, which results in a denial of service attack against the zone
if that value is used.

5. Denial of Service Attack Considerations

If an attacker is able to provide a RFC5011 validating engine with
past responses, such as when it is in-path or able to otherwise
perform any number of cache poisoning attacks, the attacker may be
able to leave the RFC5011-compliant validator without an appropriate
DNSKEY trust anchor. This scenario will remain until an
administrator manually fixes the situation.

The following timeline illustrates this situation.

5.1. Enumerated Attack Example

The following example settings are used in the example scenario
within this section:

TTL (all records) 1 day

DNSKEY RRSIG Signature Validity 10 days

Zone resigned every 1 day

Given these settings, the following sequence of events depicts how a
Trust Anchor Publisher that waits for only the RFC5011 hold time
timer length of 30 days subjects its users to a potential Denial of
Service attack. The timing schedule listed below is based on a new
trust anchor (a Key Signing Key (KSK)) being published at time T+0.
All numbers in this sequence refer to days before and after such an
event. Thus, T-1 is the day before the introduction of the new key,
and T+15 is the 15th day after the key was introduced into the
fictitious zone being discussed.

In this dialog, we consider two keys being published:

K_old The older KSK being replaced.

K_new The new KSK being transitioned into active use, using the
    RFC5011 process.

In this dialog, the following actors are playing roles in this
situation:

Zone Signer The owner of a zone intending to publish a new Key-
    Signing-Key (KSK) that will become a trust anchor by validators
    following the RFC5011 process.

RFC5011 Validator A DNSSEC validator that is using the RFC5011
    processes to track and update trust anchors.

Attacker An attacker intent on foiling the RFC5011 Validator's
    ability to successfully adopt the Zone Signer's K_new key as a new
    trust anchor.

5.1.1. Attack Timing Breakdown

The following series of steps depicts the timeline in which an attack
occurs that foils the adoption of a new DNSKEY by a Trust Anchor
Publisher that revokes the old key too quickly.

T-1 The last signatures are published by the Zone Signer that signs
    only K_old using K_old. The Attacker queries for, retrieves and
    caches this keyset and corresponding signatures.

T-0 The Zone Signer adds K_new to his zone and signs the zone's key
    set with K_old. The RFC5011 Validator retrieves the new key set
    and corresponding signature set and notices the publication of
    K_new. The RFC5011 Validator starts the (30-day) hold-down timer
    for K_new.

T+5 The RFC5011 Validator queries for the zone's keyset per the
    Active Refresh schedule, discussed in Section 2.3 of RFC5011.
    Instead of receiving the intended published keyset, the Attacker
    successfully replays the keyset and associated signatures that
    they recorded at T-1. Because the signature lifetime is 10 days
    (in this example), the replayed signature and keyset is accepted
    as valid (being only 6 days old) and the RFC5011 Validator cancels
    the hold-down timer for K_new.

T+10 The RFC5011 Validator queries for the zone's keyset and
    discovers K_new again, signed by K_old (the attacker is unable to
    replay the records cached at T-1, because they have now expired).
    The RFC5011 Validator starts (anew) the hold-timer for K_new.

T+15,T+20, and T+25 The RFC5011 Validator continues checking the
    zone's key set and lets the hold-down timer keep running without
    resetting it.

T+30 The Zone Signer knows that this is the first time at which some
    validators might accept K_new as a new trust anchor, since the
    hold-down timer of a RFC5011 Validator not under attack that had
    queried and retrieved K_new at T+0 would now have reached 30 days.
    However, the hold-down timer of our attacked RFC5011 Validator is
    only at 20 days.

T+35 The Zone Signer (mistakenly) believes that all validators
    following the Active Refresh schedule (Section 2.3 of RFC5011)
    should have accepted K_new as a the new trust anchor (since the

hold down time of 30 days + 1/2 the signature validity period
would have passed). However, the hold-down timer of our attacked
RFC5011 Validator is only at 25 days; The replay attack at T+5
means its new hold-time timer actually started at T+10, and thus
at this time it's real hold-down timer is at T+35 - T+10 = 25
days, which is less than the RFC5011 required 30 days.

T+36 The Zone Signer, believing K_new is safe to use, switches their
active signing KSK to K_new and publishes a new DNSKEY set
signature signed with K_new. Non-attacked RFC5011 validators,
with a hold-down timer of at least 30 days, would have accepted
K_new into their set of trusted keys. But, because our attacked
RFC5011 Validator still has a hold-down timer for K_new at 26
days, it will fail to accept K_new as a trust anchor and since
K_old is no longer being used, all the KSK records from the zone
signed by K_new will be treated as invalid. Subsequently, all
keys in the key set are now unusable, invalidating all records in
the zone of any type and name.

6. Minimum RFC5011 Timing Requirements

Given the attack description in Section 5, the correct minimum length
of time required for the Zone Signer to wait before using K_new is:

```
  waitTime = addHoldDownTime
            + MIN(3 * (DNSKEY RRSIG Signature Validity) / 2,
                    3 * MAX(original TTL of K_old DNSKEY RRSet) / 2,
                    15 days)
            + 2 * MAX(TTL of all records)
```

The most confusing element of the above equation comes from the "3 *
(DNSKEY RRSIG Signature Validity) / 2" element, but is the most
critical to understand and get right.

The RFC5011 "Active Refresh" requirements state that:

A resolver that has been configured for an automatic update
of keys from a particular trust point MUST query that trust
point (e.g., do a lookup for the DNSKEY RRSet and related
RRSIG records) no less often than the lesser of 15 days, half
the original TTL for the DNSKEY RRSet, or half the RRSIG
expiration interval and no more often than once per hour.

The important timing constraint that must be considered is the last
point at which a validating resolver may have received a replayed the
original DNSKEY set (K_old) without the new key. It's the next query
of the RFC5011 validator that the assured K_new will be seen. For
the following paragraph, we use the more common case where (DNSKEY
RRSIG Signature Validity) is greater than (original TTL for the
DNSKEY RRSet), although the equation above takes both cases into
account.

Thus, the worst case scenario of this attack is when the attacker can
replay K_old at just before (DNSKEY RRSIG Signature Validity). If a
RFC5011 validator picks up K_old at this this point, it will not have
a hold down timer started at all. It's not until the next "Active
Refresh" time that they'll pick up K_new with assurance, and thus
start their hold down timer. Thus, this is not at (DNSKEY RRSIG
Signature Validity) time past publication, but rather 3 * (DNSKEY
RRSIG Signature Validity) / 2.

The extra 2 * MAX(TTL of all records) is the standard added safety
margin when dealing with DNSSEC due to caching that can take place.
Because the 5011 steps require direct validation using the signature
validity, the authors aren't yet convinced it is needed in this
particular case.

For the parameters listed in Section 5.1, our example:

```
  waitTime = 30
          + 3 * (10) / 2
          + 2 * (1)        (days)

  waitTime = 47            (days)
```

This hold-down time of 47 days is 12 days longer than the frequently
perceived 35 days in T+35 above.

7. IANA Considerations

This document contains no IANA considerations.

8. Operational Considerations

A companion document to RFC5011 was expected to be published that
describes the best operational practice considerations from the
perspective of a zone publisher and Trust Anchor Publisher. However,
this companion document was never written. The authors of this
document hope that it will at some point in the future, as RFC5011
timing can be tricky as we have shown. This document is intended

only to fill a single operational void that results in security
ramifications (specifically a denial of service attack against an
RFC5011 Validator). This document does not attempt to document any
other missing operational guidance for zone publishers.

9. Security Considerations

This document, is solely about the security considerations with
respect to the Trust Anchor Publisher of RFC5011 trust anchors /
keys. Thus the entire document is a discussion of Security
Considerations

10. Acknowledgements

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help and advice. We would also like to thank Bob Harold, Shane Kerr,
Matthijs Mekking, Duane Wessels, and everyone else who assisted with
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Appendix A. Changes / Author Notes.

From -00 to -02:

Additional background and clarifications in abstract.

Better separation in attack description between attacked and non-
attacked resolvers.

Some language cleanup.

Clarified that this is maths (and math is hard, let's go
shopping!)

Changed to " <?rfc include='reference....'?> " style references.

From -02 to -03:

      Minor changes from Bob Harold

      Clarified why 3/2 signature validity is needed

      Changed min wait time math to include TTL value as well

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