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                          RIFT Auto-EVPN

Abstract

   This document specifies procedures that allow an EVPN overlay to be
   fully and automatically provisioned when using RIFT as underlay and
   leveraging its no touch ZTP architecture.

Status of This Memo

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1. Introduction

RIFT is a protocol that focuses heavily on operational simplicity.
[RIFT] natively supports Zero Touch Provisioning (ZTP) functionality
that allows each node in an underlay network to automatically derive
its place in the topology and configure itself accordingly when
properly cabled. RIFT can also disseminate Key-Value information
contained in Key-Value Topology Information Elements (KV-TIEs).
These KV-TIEs can contain any information and therefore be used for
any purpose. Leveraging RIFT to provision EVPN overlays without any
need for configuration and leveraging KV capabilities to easily
validate correct operation of such overlay without a single point of
failure would provide significant benefit to operators in terms of
simplicity and robustness of such a solution.

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT",
"SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this
document are to be interpreted as described in RFC 2119 [RFC2119].

2. Design Considerations

EVPN supports various service models, this document defines a method
for the VLAN-Aware service model defined in [RFC7432]. Other service
models may be considered in future revisions of this document.

Each model has its own set of requirements for deployment. For example, a functional BGP overlay is necessary to exchange EVPN NLRI regardless of the service model. Furthermore, the requirements are made up of individual variables, such as each node's loopback address and AS number for the BGP session. Some of these variables may be coordinated across each node in a network, but are ultimately locally significant (e.g. route distinguishers). Similarly, calculation of some variables will be local only to each device. RIFT contains currently enough topology information in each node to calculate all those necessary variables automatically.

Once the EVPN overlay is configured and becomes operational KV TIEs can be used to distribute state information to allow for validation of basic operational correctness without need for further tooling.

3. System ID

The 64-bit RIFT System ID that uniquely identifies a node as defined in [RIFT](#) [[RIFT](#)].

4. Fabric ID

RIFT operates on variants of Clos substrate which are commonly called an IP Fabric. Since EVPN VLANs can be either contained within one fabric or span them, Auto-EVPN introduces the concept of a Fabric ID into RIFT.

This section describes an optional extension to LIE packet schema in the form of a 16-bit Fabric ID that identifies a nodes membership within a particular fabric. Auto-EVPN capable nodes MUST support this extension but MAY not advertise it when not participating in Auto-EVPN. A non-present Fabric ID and value of 0 is reserved as ANY_FABRIC and MUST NOT be used for any other purpose.

Fabric ID MUST be considered in existing adjacency FSM rules so nodes that support Auto-EVPN can interoperate with nodes that do not. The LIE validation is extended with following clause and if it is not met, miscabling should be declared:

```
(if fabric_id is not advertised by either node OR
if fabric_id is identical on both nodes)
AND
(if auto_evpn_version is not advertised by either node OR
if auto_evpn_version is identical on both nodes)
```

The appendix details [LIE](#) ([Appendix A.1.2](#)) and [Node-TIE](#) ([Appendix A. 2.2](#)) schema changes.

5. Auto-EVPN Device Roles

Auto-EVPN requires that each node understand its given role within
the scope of the EVPN implementation so each node derives the
necessary variables and provides the necessary overlay
configuration. For example, a leaf node performing VXLAN gateway
functions does not need to derive its own Cluster ID or learn one
from the route reflector that it peers with.

5.1. All Participating Nodes

Not all nodes have to participate in Auto-EVPN but when they do they
do assume EVPN roles and MUST derive according variables:

IPv6 Loopback Address
Unique IPv6 loopback address used in BGP sessions.

Router ID
The BGP Router ID.

Autonomous System Number
The ASN for IBGP sessions.

Cluster ID
The Cluster ID for Top-of-Fabric IBGP route reflection.

5.2. ToF Nodes as Route Reflectors

This section defines an Auto-EVPN role whereby some Top-of-Fabric
nodes act as EVPN route reflectors. It is expected that route
reflectors would establish IBGP sessions with leaf nodes in the same
fabric. The typical route reflector requirements do not change,
however determining which specific values to use requires further
consideration. ToF nodes performing route reflector functionality
MUST derive the following variables:

IPv6 RR Loopback Address
The source address for IBGP sessions with leaf nodes in case
ToF won election for one of the route reflectors in the
fabric.

IPv6 RR Acceptable Prefix Range
Range of addresses acceptable by the route reflector to form a
IBGP session. This range covers ALL possible IPv6 Loopback
Addresses derived by other Auto EVPN nodes in the current
fabric and other Auto-EVPN RRs addresses.

5.3. Leaf Nodes

Leaf nodes derive their role from realizing they are at the bottom
of the fabric, i.e. not having any southbound adjacencies.
Alternately, a node can assume a leaf node if it has only southbound
adjacencies to nodes with explicit LEAF_LEVEL to allow for scenarios
where RIFT leaves do NOT participate in Auto-EVPN.

Leaf nodes MUST derive the following variables:

IPv6 RR Loopback Adresses
Addresses of the RRs present in the fabric. Those addresses
are used to build BGP sessions to the RR.

EVIs
Leaf node derives all the necessary variables to instantiate
EVIs with layer-2 and optionally layer-3 functionality.

If a leaf node is required to perform layer-2 VXLAN gateway
functions, it MUST be capable of deriving the following types of
variables:

Route Distinguisher
The route distinguisher corresponding to a MAC-VRF that
uniquely identifies each node.

Route Target
The route target that corresponds to a MAC-VRF.

MAC VRF name
This is an optional variable to provide a common MAC VRF name
across all leaves.

Set of VLANs
Those are VLANs provisioned either within the fabric or
allowing to stretch across fabrics.

For each VLAN derived in an EVI the following variables MUST be
derived:

VLAN
The VLAN ID.

name
This is an optional variable to provide a common VLAN name
across all leaves.

VNI
The VNI that corresponds to the VLAN ID. This will contribute
to the EVPN Type-2 route.

IRB
> Optional variables of the IRB for the VLAN if the leaf performs layer-3 gateway function.

If a leaf node is required to perform layer-3 VXLAN gateway functions, it MUST additionally be capable of deriving the following types of variables:

IP Gateway MAC Address
> The MAC address associated with IP gateway.

IP Gateway Subnetted Address
> The IPv4 and/or IPv6 gateway address including its subnet length.

Type-5 EVPN IP Prefix with ToFs performing gateway functionality can also be derived and will be described in a future version of this document.

6. Auto-EVPN Variable Derivation

As previously mentioned, not all nodes are required to derive all variables in a given network (e.g. a transit spine node may not need to derive any or participate in Auto-EVPN). Additionally, all derived variables are derived from RIFT's FSM or ZTP mechanism so no additional flooding beside RIFT flooding is necessary for the functionality.

It is also important to mention that all variable derivation is in some way based on combinations of System ID, MAC-VRF ID, Fabric ID, EVI and VLAN and MUST comply precisely with calculation methods specified in the Appendix section to allow interoperability between different implementations.

6.1. Auto-EVPN Version

This section describes extensions to both the RIFT LIE packet and Node-TIE schemas in the form of a 16-bit value that identifies the Auto-EVPN Version. Auto-EVPN capable nodes MUST support this extension, but MAY choose not to advertise it in LIEs and Node-TIEs when Auto-EVPN is not being utilized. The appendix describes LIE (Appendix A.1.1) and Node-TIE (Appendix A.2.1) schema changes in detail.

6.2. MAC-VRF ID

This section describes a variable MAC-VRF ID that uniquely identifies an instance of EVPN instance (EVI) and is used in variable derivation procedures. Each EVPN EVI MUST be associated with a unique MAC-VRF ID, this document does not specify a method

for making that association or ensuring that they are coordinated
properly across fabric(s).

6.3. Loopback Address

First and foremost, RIFT does not advertise anything more specific
than the fabric default route in the southbound direction by
default. However, Auto-EVPN nodes MUST advertise specific loopback
addresses southbound to all other Auto-EVPN nodes so to establish
MP-BGP reachability correctly in all scenarios.

Auto-EVPN nodes MUST derive a ULA-scoped IPv6 loopback address to be
used as both the IBGP source address, as well as the VTEP source
when VXLAN gateways are required. Calculation is done using the 6-
bytes of reserved ULA space, the 2-byte Fabric ID, and the node's 8-
byte System ID. Derivation of the System ID varies slightly
depending upon the node's location/role in the fabric and will be
described in subsequent sections.

IPv4 addresses MAY be supported, but it should be noted that they
have a higher likelihood of collision.

The required algorithm can be found in the [appendix]([Appendix A.
3.3]).

6.3.1. Leaf Nodes as Gateways

Calculation is done using the 6-bytes of reserved ULA space, the 2-
byte Fabric ID, and the node's 8-byte System ID.

6.3.2. ToF Nodes as Route Reflectors

ToF nodes acting as route reflectors MUST derive their loopback
address according to the specific section describing the algorithm.
Calculation is done using the 6-bytes of reserved ULA space, the 2-
byte Fabric ID, and the 8-byte System ID of each elected route
reflector.

6.3.2.1. Route Reflector Election Procedures

Four Top-of-Fabric nodes MUST be elected as an IBGP route reflector.
Each ToF performs the election independently based on system IDs of
other ToFs in the fabric obtained via southbound reflection. The
route reflector election procedures are defined as follows:

  1. ToF node with the highest System ID.

  2. ToF node with the lowest System ID.

  3. ToF node with the 2nd highest System ID.

4. ToF node with the 2nd lowest System ID.

This ordering is necessary to prevent a single node with either the
highest or lowest System ID from triggering changes to route
reflector loopback addresses as it would result in all BGP sessions
dropping.

For example, if two nodes, ToF01 and ToF02 with System IDs
002c6af5a281c000 and 002c6bf5788fc000 respectively, ToF02 would be
elected due to it having the highest System ID of the ToFs
(002c6bf5788fc000). If a ToF determines that it is elected as route
reflector, it uses the knowledge of its position in the list to
derive route reflector v6 loopback address.

Considerations for multiplane route reflector elections will be
included in future revisions.

6.4. Autonomous System Number

Nodes in each fabric MUST derive a private autonomous system number
based on its Fabric ID so that it is unique across the fabric.

The required algorithm for 2-byte ASNs can be found in the appendix
(Appendix A.3.4).

6.5. Cluster ID

Route reflector nodes in each fabric MUST derive a cluster ID that
is based on its Fabric ID so that it is unique across the fabric.
Implementations MAY choose to simply use the AS number as the
cluster ID.

The required algorithm can be found in the appendix (Appendix A.
3.5).

6.6. Router ID

Nodes MUST drive a Router ID that is based on both its System ID and
Fabric ID so that it is unique to both.

The required algorithm can be found in the appendix (Appendix A.
3.6).

6.7. Route Target

Nodes hosting EVPN EVIs MUST derive a route target extended
community based on the MAC-VRF ID for each EVI so that it is unique
across the network. Route targets MUST be of type 0 as per RFC4360.

For example, if given a MAC-VRF ID of 1, the derived route target would be "target:1"

The required algorithm can be found in the [appendix](#) ([Appendix A. 3.7](#)).

6.8. Route Distinguisher

Nodes hosting EVPN EVIs MUST derive a type-0 route distinguisher based on its System ID and Fabric ID so that it is unique per MAC-VRF and per node.

The required algorithm can be found in the [appendix](#) ([Appendix A. 3.8](#)).

6.9. EVPN MAC-VRF Services

It's obvious that applications utilizing Auto-EVPN overlay services may require a variety of layer-2 and/or layer-3 traffic considerations. Variables supporting these services are also derived based on some combination of MAC-VRF ID, Fabric ID, and other constant values. Integrated Routing and Bridging (IRB) gateway address derivation also leverages a set of constant "random seed" values to provide additional entropy.

The required derivation procedures can be found in the [appendix](#) ([Appendix A.3](#)).

6.9.1. Untagged Traffic in Multiple Fabrics

This section defines a methods to derive unique VLAN, VNI, MAC, and gateway address values for deployments where untagged traffic is stretched across multiple fabrics.

6.9.1.1. VLAN

Untagged traffic stretched across multiple fabrics MUST derive VLAN tags based on MAC-VRF ID in conjunction with a constant value of 1 (i.e. MAC-VRF ID + 1).

6.9.1.2. VNI

Untagged traffic stretched across multiple fabrics MUST derive VNIs based on MAC-VRF ID and Fabric ID in conjunction with a constant value. These VNIs MUST correspond to EVPN Type-2 routes.

6.9.1.3. MAC Address

The MAC address MUST be a unicast address and also MUST be identical for any IRB gateways that belong to an individual bridge-domain

across fabrics. The last 5-bytes MUST be a hash of the MAC-VRF ID
and a constant value of 1 that is calculated using the previously
mentioned random seed values.

6.9.1.4. IPv6 IRB Gateway Address

The derived IPv6 gateway address MUST be from a ULA-scoped range
that will account for the first 6-bytes. The next 5-bytes MUST be
the last bytes of the derived MAC address. Finally, the remaining 7-
bytes MUST be ::0001.

6.9.1.5. IPv4 IRB Gateway Address

The derived IPv4 gateway address MUST be from a RFC1918 range, which
accounts for the first octet. The next octet MUST a hash of the MAC-
VRF ID and a constant value of 1 that is calculated using the
previously mentioned random seed values. Finally, the remaining 2
octets MUST be 0 and 1 respectively.

6.9.2. Tagged Traffic in Multiple Fabrics

This section defines a methods to derive unique VLAN, VNI, MAC, and
gateway address values for deployments where tagged traffic is
stretched across multiple fabrics.

6.9.2.1. VLAN

Tagged traffic stretched across multiple fabrics MUST derive VLAN
tags based on MAC-VRF ID in conjunction with a constant value of 16
(i.e. MAC-VRF ID + 16).

6.9.2.2. VNI

Tagged traffic stretched across multiple fabrics MUST derive VNIs
based on MAC-VRF ID and Fabric ID in conjunction with a constant
value. These VNIs MUST correspond to EVPN Type-2 routes.

6.9.2.3. MAC Address

The MAC address MUST be a unicast address and also MUST be identical
for any IRB gateways that belong to an individual bridge-domain
across fabrics. The last 5-bytes MUST be a hash of the MAC-VRF ID
and a constant value of 1 that is calculated using the previously
mentioned random seed values.

6.9.2.4. IPv6 IRB Gateway Address

The derived IPv6 gateway address MUST be from a ULA-scoped range
that will account for the first 6-bytes. The next 5-bytes MUST be

the last bytes of the derived MAC address. Finally, the remaining 7-bytes MUST be ::0001.

6.9.2.5. IPv4 IRB Gateway Address

The derived IPv4 gateway address MUST be from a RFC1918 range, which accounts for the first octet. The next octet MUST a hash of the MAC-VRF ID and a constant value of 16 that is calculated using the previously mentioned random seed values. Finally, the remaining 2 octets MUST be 0 and 1 respectively.

6.9.3. Tagged Traffic in a Single Fabric

This section defines a methods to derive unique VLAN, VNI, MAC, and gateway address values for deployments where untagged traffic is contained within a single fabric.

6.9.3.1. VLAN

Tagged traffic contained to a single fabric MUST derive VLAN tags based on MAC-VRF ID and Fabric ID in conjunction with a constant value of 17 (i.e. MAC-VRF ID + Fabric ID + 17).

6.9.3.2. VNI

Tagged traffic contained to a single fabric MUST derive VNIs based on MAC-VRF ID and Fabric ID in conjunction with a constant value. These VNIs MUST correspond to EVPN Type-2 routes.

6.9.3.3. MAC Address

The MAC address MUST be a unicast address and also MUST be identical for any IRB gateways that belong to an individual bridge-domain across fabrics. The last 5-bytes MUST be a hash of the MAC-VRF ID and a constant value of 1 that is calculated using the previously mentioned random seed values.

6.9.3.4. IPv6 IRB Gateway Address

The derived IPv6 gateway address MUST be from a ULA-scoped range, which accounts for the first 6-bytes. The next 5-bytes MUST be the last bytes of the derived MAC address. Finally, the remaining 7-bytes MUST be ::0001.

6.9.3.5. IPv4 IRB Gateway Address

The derived IPv4 gateway address MUST be from a RFC1918 range, which accounts for the first octet. The next octet MUST a hash of the MAC-VRF ID and a constant value of 17 that is calculated using the

previously mentioned random seed values. Finally, the remaining 2
octets MUST be 0 and 1 respectively.

6.9.4. Traffic Routed to External Destinations

6.9.4.1. Route Distinguisher

Nodes hosting IP Prefix routes MUST derive a type-0 route
distinguisher based on its System ID and Fabric ID so that it is
unique per IP-VRF and per node.

The required algorithm can be found in the appendix (Appendix A.
3.8).

6.9.4.2. Route Target

Nodes hosting IP prefix routes MUST derive a route target extended
community based on the MAC-VRF ID for each IP-VRF so that it is
unique across the network. Route targets MUST be of type 0.

The required algorithm can be found in the appendix (Appendix A.
3.7).

7. Acknowledgements

TBD

8. Security Considerations

This document introduces no new security concerns to RIFT or other
specifications referenced in this document.

9. References

9.1. Normative References

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Appendix A. Appendix

A.1. RIFT LIE Schema

A.1.1. Auto-EVPN Version

```
struct LIEPacket {
 ...
   /** It provides the optional ID of the configured fabric  */
   25: optional common.FabricIDType        fabric_id;
...
```

A.1.2. Fabric ID

```
...
struct LIEPacket {
 ...
   /** It provides optional version of EVPN ZTP as 256 * MAJOR + MINOR */
   26: optional i16                        auto_evpn_version;
...
```

A.2. RIFT Node-TIE Schema

A.2.1. Auto-EVPN Version

```
struct NodeTIEElement {
...
   /** It provides optional version of EVPN ZTP as 256 * MAJOR + MINOR */
   13: optional i16                        auto_evpn_version;
```

A.2.2. Fabric ID

```
  struct NodeTIEElement {
 ...
    /** It provides the optional ID of the Fabric configured */
   12: optional common.FabricIDType        fabric_id;
```

A.3. Variable Derivation

A.3.1. Random Seed Values

   To be provided in future version of this document.

A.3.2. Fabric ID

To be provided in future version of this document.

A.3.3. Loopback Address

To be provided in future version of this document.

A.3.4. Autonomous System Number

To be provided in future version of this document.

A.3.5. Cluster ID

To be provided in future version of this document.

A.3.6. Router ID

To be provided in future version of this document.

A.3.7. Route Target

To be provided in future version of this document.

A.3.8. Route Distinguisher

To be provided in future version of this document.

A.3.9. VLAN

To be provided in future version of this document.

A.3.10. VNI

To be provided in future version of this document.

A.3.11. Gateway (MAC)

To be provided in future version of this document.

A.3.12. Gateway (IPv6)

To be provided in future version of this document.

A.3.13. Gateway (IPv4)

To be provided in future version of this document.

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