

Network Working Group  
Internet-Draft  
Intended status: Standards Track  
Expires: February 7, 2020

C. Filsfils, Ed.  
D. Dukes, Ed.  
Cisco Systems, Inc.  
S. Previdi  
Huawei  
J. Leddy  
Individual  
S. Matsushima  
Softbank  
D. Voyer  
Bell Canada  
August 6, 2019

**IPv6 Segment Routing Header (SRH)  
draft-ietf-6man-segment-routing-header-22**

Abstract

Segment Routing can be applied to the IPv6 data plane using a new type of Routing Extension Header called the Segment Routing Header. This document describes the Segment Routing Header and how it is used by Segment Routing capable nodes.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of [BCP 78](#) and [BCP 79](#).

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at <https://datatracker.ietf.org/drafts/current/>.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

This Internet-Draft will expire on February 7, 2020.

Copyright Notice

Copyright (c) 2019 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to [BCP 78](#) and the IETF Trust's Legal Provisions Relating to IETF Documents

(<https://trustee.ietf.org/license-info>) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Simplified BSD License text as described in Section 4.e of the Trust Legal Provisions and are provided without warranty as described in the Simplified BSD License.

Table of Contents

- [1. Introduction . . . . .](#) [3](#)
- [1.1. Requirements Language . . . . .](#) [3](#)
- [2. Segment Routing Header . . . . .](#) [4](#)
- [2.1. SRH TLVs . . . . .](#) [6](#)
- [2.1.1. Padding TLVs . . . . .](#) [8](#)
- [2.1.2. HMAC TLV . . . . .](#) [9](#)
- [3. SR Nodes . . . . .](#) [12](#)
- [3.1. Source SR Node . . . . .](#) [12](#)
- [3.2. Transit Node . . . . .](#) [12](#)
- [3.3. SR Segment Endpoint Node . . . . .](#) [12](#)
- [4. Packet Processing . . . . .](#) [12](#)
- [4.1. Source SR Node . . . . .](#) [12](#)
- [4.1.1. Reduced SRH . . . . .](#) [13](#)
- [4.2. Transit Node . . . . .](#) [13](#)
- [4.3. SR Segment Endpoint Node . . . . .](#) [14](#)
- [4.3.1. FIB Entry Is Locally Instantiated SRv6 SID . . . . .](#) [14](#)
- [4.3.2. FIB Entry is a Local Interface . . . . .](#) [16](#)
- [4.3.3. FIB Entry Is A Non-Local Route . . . . .](#) [17](#)
- [4.3.4. FIB Entry Is A No Match . . . . .](#) [17](#)
- [5. Intra SR Domain Deployment Model . . . . .](#) [17](#)
- [5.1. Securing the SR Domain . . . . .](#) [17](#)
- 5.2. SR Domain as a single system with delegation among components . . . . . [18](#)
- [5.3. MTU Considerations . . . . .](#) [19](#)
- [5.4. ICMP Error Processing . . . . .](#) [19](#)
- [5.5. Load Balancing and ECMP . . . . .](#) [19](#)
- [5.6. Other Deployments . . . . .](#) [20](#)
- [6. Illustrations . . . . .](#) [20](#)
- [6.1. Abstract Representation of an SRH . . . . .](#) [20](#)
- [6.2. Example Topology . . . . .](#) [21](#)
- [6.3. Source SR Node . . . . .](#) [22](#)
- [6.3.1. Intra SR Domain Packet . . . . .](#) [22](#)
- [6.3.2. Inter SR Domain Packet - Transit . . . . .](#) [22](#)
- [6.3.3. Inter SR Domain Packet - Internal to External . . . . .](#) [22](#)
- [6.4. Transit Node . . . . .](#) [23](#)
- [6.5. SR Segment Endpoint Node . . . . .](#) [23](#)
- [6.6. Delegation of Function with HMAC Verification . . . . .](#) [23](#)
- [6.6.1. SID List Verification . . . . .](#) [23](#)



<a href="#">7.</a>	<a href="#">Security Considerations</a>	<a href="#">24</a>
<a href="#">7.1.</a>	<a href="#">Source Routing Attacks</a>	<a href="#">24</a>
<a href="#">7.2.</a>	<a href="#">Service Theft</a>	<a href="#">25</a>
<a href="#">7.3.</a>	<a href="#">Topology Disclosure</a>	<a href="#">25</a>
<a href="#">7.4.</a>	<a href="#">ICMP Generation</a>	<a href="#">25</a>
<a href="#">7.5.</a>	<a href="#">Applicability of AH</a>	<a href="#">26</a>
<a href="#">8.</a>	<a href="#">IANA Considerations</a>	<a href="#">26</a>
<a href="#">8.1.</a>	<a href="#">Segment Routing Header Flags Register</a>	<a href="#">27</a>
<a href="#">8.2.</a>	<a href="#">Segment Routing Header TLVs Register</a>	<a href="#">27</a>
<a href="#">9.</a>	<a href="#">Implementation Status</a>	<a href="#">28</a>
<a href="#">9.1.</a>	<a href="#">Linux</a>	<a href="#">28</a>
<a href="#">9.2.</a>	<a href="#">Cisco Systems</a>	<a href="#">28</a>
<a href="#">9.3.</a>	<a href="#">FD.io</a>	<a href="#">28</a>
<a href="#">9.4.</a>	<a href="#">Barefoot</a>	<a href="#">28</a>
<a href="#">9.5.</a>	<a href="#">Juniper</a>	<a href="#">29</a>
<a href="#">9.6.</a>	<a href="#">Huawei</a>	<a href="#">29</a>
<a href="#">10.</a>	<a href="#">Contributors</a>	<a href="#">29</a>
<a href="#">11.</a>	<a href="#">Acknowledgements</a>	<a href="#">29</a>
<a href="#">12.</a>	<a href="#">References</a>	<a href="#">29</a>
<a href="#">12.1.</a>	<a href="#">Normative References</a>	<a href="#">29</a>
<a href="#">12.2.</a>	<a href="#">Informative References</a>	<a href="#">30</a>
	<a href="#">Authors' Addresses</a>	<a href="#">32</a>

## [1.](#) Introduction

Segment Routing can be applied to the IPv6 data plane using a new type of Routing Header called the Segment Routing Header. This document describes the Segment Routing Header and how it is used by Segment Routing capable nodes.

The Segment Routing Architecture [[RFC8402](#)] describes Segment Routing and its instantiation in two data planes MPLS and IPv6.

The encoding of IPv6 segments in the Segment Routing Header is defined in this document.

This document uses the terms Segment Routing, SR Domain, SRv6, Segment ID (SID), SRv6 SID, Active Segment, and SR Policy as defined in [[RFC8402](#)].

### [1.1.](#) Requirements Language

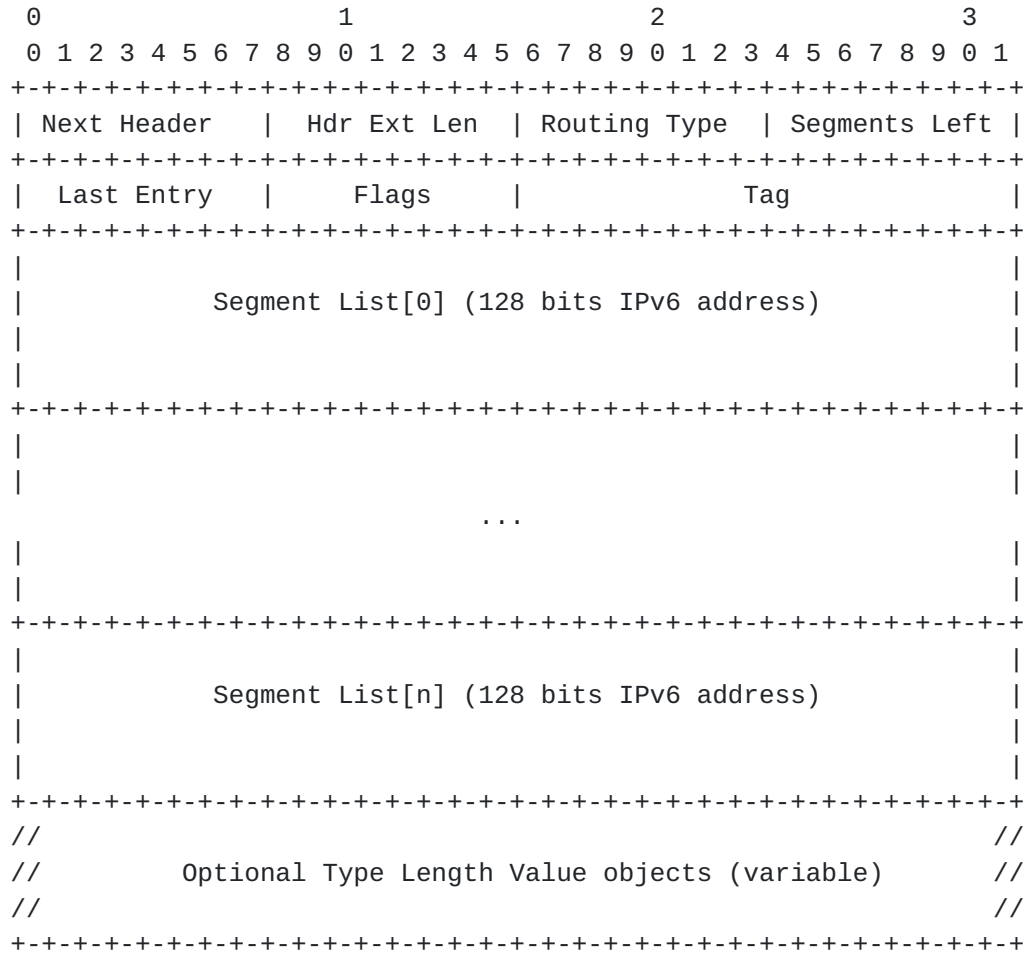
The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [BCP 14](#) [[RFC2119](#)] [[RFC8174](#)] when, and only when, they appear in all capitals, as shown here.



## 2. Segment Routing Header

Routing Headers are defined in [\[RFC8200\]](#). The Segment Routing Header has a new Routing Type (suggested value 4) to be assigned by IANA.

The Segment Routing Header (SRH) is defined as follows:



where:

- o Next Header: Defined in [\[RFC8200\] Section 4.4](#)
- o Hdr Ext Len: Defined in [\[RFC8200\] Section 4.4](#)
- o Routing Type: TBD, to be assigned by IANA (suggested value: 4).
- o Segments Left: Defined in [\[RFC8200\] Section 4.4](#)
- o Last Entry: contains the index (zero based), in the Segment List, of the last element of the Segment List.



- o Flags: 8 bits of flags. [Section 8.1](#) creates an IANA registry for new flags to be defined. The following flags are defined:

```

  0 1 2 3 4 5 6 7
+-+--+--+--+--+--+
|U U U U U U U U|
+-+--+--+--+--+--+

```

U: Unused and for future use. MUST be 0 on transmission and ignored on receipt.

- o Tag: tag a packet as part of a class or group of packets, e.g., packets sharing the same set of properties. When tag is not used at source it MUST be set to zero on transmission. When tag is not used during SRH Processing it SHOULD be ignored. Tag is not used when processing the SID defined in [Section 4.3.1](#). It may be used when processing other SIDs which are not defined in this document. The allocation and use of tag is outside the scope of this document.
- o Segment List[n]: 128 bit IPv6 addresses representing the nth segment in the Segment List. The Segment List is encoded starting from the last segment of the SR Policy. I.e., the first element of the segment list (Segment List [0]) contains the last segment of the SR Policy, the second element contains the penultimate segment of the SR Policy and so on.
- o Type Length Value (TLV) are described in [Section 2.1](#).

In the SRH, the Next Header, Hdr Ext Len, Routing Type, and Segments Left fields are defined in [Section 4.4 of \[RFC8200\]](#). Based on the constraints in that section Next Header, Header Ext Len, and Routing Type are not mutable while Segments Left is mutable.

The mutability of the TLV value is defined by the most significant bit in the type, as specified in [Section 2.1](#).

[Section 4.3](#) defines the mutability of the remaining fields in the SRH (Flags, Tag, Segment List) in the context of the SID defined in this document.

New SIDs defined in the future MUST specify the mutability properties of the Flags, Tag, and Segment List and indicate how the HMAC TLV ([Section 2.1.2](#)) verification works. Note, that in effect these fields are mutable.





Consistent with the source routing model, the source of the SRH always knows how to set the segment list, Flags, Tag and TLVs of the SRH for use within the SR Domain. How it achieves this is outside the scope of this document, but may be based on topology, available SIDs and their mutability properties, the SRH mutability requirements of the destination, or any other information.

## 2.1. SRH TLVs

This section defines TLVs of the Segment Routing Header.

A TLV provides meta-data for segment processing. The only TLVs defined in this document are the HMAC ([Section 2.1.2](#)) and PAD ([Section 2.1.1](#)) TLVs. While processing the SID defined in [Section 4.3.1](#), all TLVs are ignored unless local configuration indicates otherwise ([Section 4.3.1.1.1](#)). Thus, TLV and HMAC support is optional for any implementation, however an implementation adding or parsing TLVs MUST support PAD TLVs. Other documents may define additional TLVs and processing rules for them.

TLVs are present when the Hdr Ext Len is greater than  $(\text{Last Entry} + 1) * 2$ .

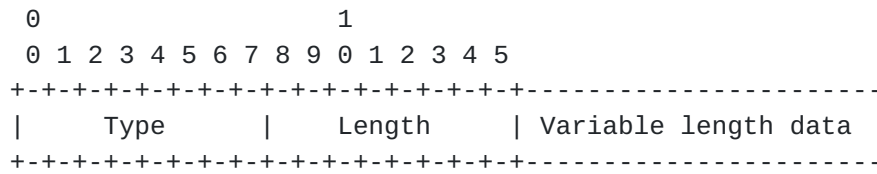
While processing TLVs at a segment endpoint, TLVs MUST be fully contained within the SRH as determined by the Hdr Ext Len. Detection of TLVs exceeding the boundary of the SRH Hdr Ext Len results in an ICMP Parameter Problem, Code 0, message to the Source Address, pointing to the Hdr Ext Len field of the SRH, and the packet being discarded.

An implementation MAY limit the number and/or length of TLVs it processes based on local configuration. It MAY:

- o Limit the number of consecutive Pad1 ([Section 2.1.1.1](#)) options to 1, if padding of more than one byte is required then PadN ([Section 2.1.1.2](#)) should be used.
- o Limit the length in PadN to 5.
- o Limit the maximum number of non-Pad TLVs to be processed.
- o Limit the maximum length of all TLVs to be processed.

The implementation MAY stop processing additional TLVs in the SRH when these configured limits are exceeded.





Type: An 8 bit value. Unrecognized Types MUST be ignored on receipt.

Length: The length of the Variable length data.

Variable length data: Length bytes of data that is specific to the Type.

Type Length Value (TLV) contain OPTIONAL information that may be used by the node identified in the Destination Address (DA) of the packet.

Each TLV has its own length, format and semantic. The code-point allocated (by IANA) to each TLV Type defines both the format and the semantic of the information carried in the TLV. Multiple TLVs may be encoded in the same SRH.

The highest-order bit of the TLV type specifies whether or not the TLV data of that type can change en route to the packet's final destination:

0: TLV data does not change en route

1: TLV data does change en route

All TLVs specify their alignment requirements using an xn+y format. The xn+y format is defined as per [RFC8200]. The SR Source nodes use the xn+y alignment requirements of TLVs and padding TLVs when constructing an SRH.

The "Length" field of the TLV is used to skip the TLV while inspecting the SRH in case the node doesn't support or recognize the Type. The "Length" defines the TLV length in octets, not including the "Type" and "Length" fields.

The following TLVs are defined in this document:

Padding TLVs

HMAC TLV

Additional TLVs may be defined in the future.



**2.1.1. Padding TLVs**

There are two types of padding TLVs, pad1 and padN, the following applies to both:

Padding TLVs are used for meeting the alignment requirement of the subsequent TLVs.

Padding TLVs are used to pad the SRH to a multiple of 8 octets.

Padding TLVs are used for alignment.

Padding TLVs are ignored by a node processing the SRH TLV.

Multiple Padding TLVs MAY be used in one SRH

**2.1.1.1. PAD1**

Alignment requirement: none

```

0 1 2 3 4 5 6 7
+--+--+--+--+--+
|   Type   |
+--+--+--+--+--+

```

Type: to be assigned by IANA (Suggested value 0)

A single Pad1 TLV MUST be used when a single byte of padding is required. If more than one byte of padding is required a Pad1 TLV MUST NOT be used, the PadN TLV MUST be used.

**2.1.1.2. PADN**

Alignment requirement: none

```

0                               1                               2                               3
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
|   Type   |   Length   |   Padding (variable)   |
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
//                               Padding (variable)                               //
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+

```

Type: to be assigned by IANA (suggested value 4).

Length: 0 to 5



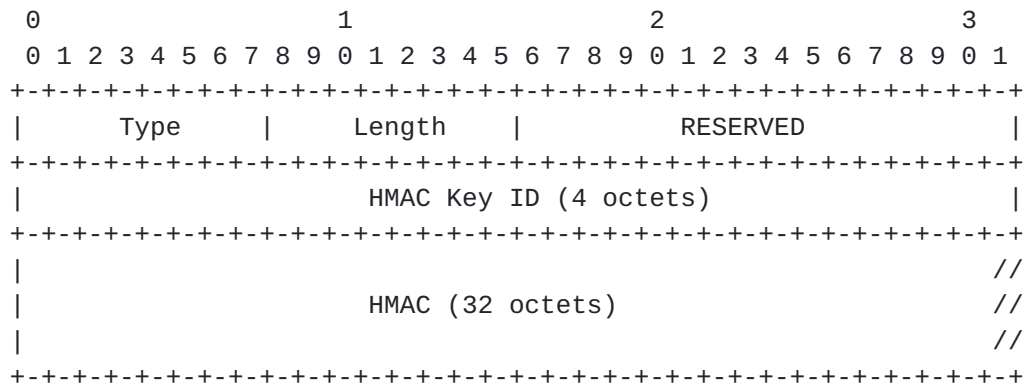
Padding: Length octets of padding. Padding bits have no semantics. They MUST be set to 0 on transmission and ignored on receipt.

The PadN TLV MUST be used when more than one byte of padding is required.

2.1.2. HMAC TLV

Alignment requirement: 8n

The keyed Hashed Message Authentication Code (HMAC) TLV is OPTIONAL and has the following format:



where:

- o Type: to be assigned by IANA (suggested value 5).
- o Length: 38.
- o RESERVED: 2 octets. MUST be 0 on transmission and ignored on receipt.
- o HMAC Key ID: A 4 octet opaque number which uniquely identifies the pre-shared key and algorithm used to generate the HMAC.
- o HMAC: 32 octets of keyed HMAC.

The HMAC TLV is used to verify the source of a packet is permitted to use the current segment in the destination address of the packet, and ensure the segment list is not modified in transit.





### **2.1.2.1. HMAC Generation and Verification**

Local configuration determines when to check for an HMAC and potentially indicates what the HMAC protects, and a requirement on where the HMAC TLV must appear (e.g. first TLV), and whether or not to verify the destination address is equal to the current segment. This local configuration is outside the scope of this document. It may be based on the active segment at an SR Segment endpoint node, the result of an ACL that considers incoming interface, HMAC Key ID, or other packet fields.

An implementation that supports the generation and verification of the HMAC SHOULD support the following default behavior as defined in the remainder of this section.

The HMAC verification begins by checking the current segment is equal to the destination address of the IPv6 header, i.e. destination address is equal to Segment List [Segments Left] and Segments Left is less than or equal to Last Segment+1.

The HMAC field is the output of the HMAC computation as defined in [[RFC2104](#)], using:

- o key: the pre-shared key identified by HMAC Key ID
- o HMAC algorithm: identified by the HMAC Key ID
- o Text: a concatenation of the following fields from the IPv6 header and the SRH, as it would be received at the node verifying the HMAC:
  - \* IPv6 header: source address (16 octets)
  - \* SRH: Last Entry (1 octet)
  - \* SRH: Flags (1 octet)
  - \* SRH: HMAC Key-id (4 octets)
  - \* SRH: all addresses in the Segment List (variable octets)

The HMAC digest is truncated to 32 octets and placed in the HMAC field of the HMAC TLV.

For HMAC algorithms producing digests less than 32 octets, the digest is placed in the lowest order octets of the HMAC field. Remaining octets MUST be set to zero.



If HMAC verification is successful, the packet is forwarded to the next segment.

If HMAC verification fails, an ICMP error message (parameter problem, error code 0, pointing to the HMAC TLV) SHOULD be generated (but rate limited) and SHOULD be logged.

#### **[2.1.2.2.](#) HMAC Pre-Shared Key Algorithm**

The HMAC Key ID field allows for the simultaneous existence of several hash algorithms (SHA-256, SHA3-256 ... or future ones) as well as pre-shared keys.

The HMAC Key ID field is opaque, i.e., it has neither syntax nor semantic except as an identifier of the right combination of pre-shared key and hash algorithm, and except that a value of 0 means that there is no HMAC field.

At the HMAC TLV verification node the Key ID uniquely identifies the pre-shared key and HMAC algorithm.

At the HMAC TLV generating node the Key ID and destination address uniquely identify the pre-shared key and HMAC algorithm. Utilizing the destination address with the Key ID allows for overlapping key IDs amongst different HMAC verification nodes. The Text for the HMAC computation is set to the IPv6 header fields and SRH fields as they would appear at the verification node, not necessarily the same as the source node sending a packet with the HMAC TLV.

Pre-shared key roll-over is supported by having two key IDs in use while the HMAC TLV generating node and verifying node converge to a new key.

An implementation supporting HMAC can support multiple hash functions. An implementation supporting HMAC MUST implement SHA-2 [[FIPS180-4](#)] in its SHA-256 variant.

The selection of pre-shared key and algorithm, and their distribution is outside the scope of this document, some options may include:

- o in the configuration of the HMAC generating or verifying nodes, either by static configuration or any SDN oriented approach
- o dynamically using a trusted key distribution protocol such as [[RFC6407](#)]



### **3. SR Nodes**

There are different types of nodes that may be involved in segment routing networks: source SR nodes originate packets with a segment in the destination address of the IPv6 header, transit nodes that forward packets destined to a remote segment, and SR segment endpoint nodes that process a local segment in the destination address of an IPv6 header.

#### **3.1. Source SR Node**

A Source SR Node is any node that originates an IPv6 packet with a segment (i.e. SRv6 SID) in the destination address of the IPv6 header. The packet leaving the source SR Node may or may not contain an SRH. This includes either:

A host originating an IPv6 packet.

An SR domain ingress router encapsulating a received packet in an outer IPv6 header, followed by an optional SRH.

The mechanism through which a segment in the destination address of the IPv6 header and the Segment List in the SRH, is derived is outside the scope of this document.

#### **3.2. Transit Node**

A transit node is any node forwarding an IPv6 packet where the destination address of that packet is not locally configured as a segment nor a local interface. A transit node is not required to be capable of processing a segment nor SRH.

#### **3.3. SR Segment Endpoint Node**

A SR segment endpoint node is any node receiving an IPv6 packet where the destination address of that packet is locally configured as a segment or local interface.

### **4. Packet Processing**

This section describes SRv6 packet processing at the SR source, Transit and SR segment endpoint nodes.

#### **4.1. Source SR Node**

A Source node steers a packet into an SR Policy. If the SR Policy results in a segment list containing a single segment, and there is



no need to add information to SRH flag or TLV, the DA is set to the single segment list entry and the SRH MAY be omitted.

When needed, the SRH is created as follows:

Next Header and Hdr Ext Len fields are set as specified in [\[RFC8200\]](#).

Routing Type field is set as TBD (to be allocated by IANA, suggested value 4).

The DA of the packet is set with the value of the first segment.

The first element of the SRH Segment List is the ultimate segment. The second element is the penultimate segment and so on.

The Segments Left field is set to  $n-1$  where  $n$  is the number of elements in the SR Policy.

The Last Entry field is set to  $n-1$  where  $n$  is the number of elements in the SR Policy.

HMAC TLV may be set according to [Section 7](#).

The packet is forwarded toward the packet's Destination Address (the first segment).

#### [4.1.1](#). Reduced SRH

When a source does not require the entire SID list to be preserved in the SRH, a reduced SRH may be used.

A reduced SRH does not contain the first segment of the related SR Policy (the first segment is the one already in the DA of the IPv6 header), and the Last Entry field is set to  $n-2$  where  $n$  is the number of elements in the SR Policy.

#### [4.2](#). Transit Node

As specified in [\[RFC8200\]](#), the only node allowed to inspect the Routing Extension Header (and therefore the SRH), is the node corresponding to the DA of the packet. Any other transit node MUST NOT inspect the underneath routing header and MUST forward the packet toward the DA according to its IPv6 routing table.

When a SID is in the destination address of an IPv6 header of a packet, it's routed through an IPv6 network as an IPv6 address. SIDs, or the prefix(es) covering SIDs, and their reachability may be





distributed by means outside the scope of this document. For example, [[RFC5308](#)] or [[RFC5340](#)] may be used to advertise a prefix covering the SIDs on a node.

### **4.3. SR Segment Endpoint Node**

Without constraining the details of an implementation, the SR segment endpoint node creates Forwarding Information Base (FIB) entries for its local SIDs.

When an SRv6-capable node receives an IPv6 packet, it performs a longest-prefix-match lookup on the packets destination address. This lookup can return any of the following:

- A FIB entry that represents a locally instantiated SRv6 SID
- A FIB entry that represents a local interface, not locally instantiated as an SRv6 SID
- A FIB entry that represents a non-local route
- No Match

#### **4.3.1. FIB Entry Is Locally Instantiated SRv6 SID**

This document, and section, defines a single SRv6 SID. Future documents may define additional SRv6 SIDs. In which case, the entire content of this section will be defined in that document.

If the FIB entry represents a locally instantiated SRv6 SID, process the next header chain of the IPv6 header as defined in [section 4 of \[RFC8200\]](#). [Section 4.3.1.1](#) describes how to process an SRH, [Section 4.3.1.2](#) describes how to process an upper layer header or no next header.

Processing this SID modifies the Segments Left and, if configured to process TLVs, it may modify the "variable length data" of TLV types that change en route. Therefore Segments Left is mutable and TLVs that change en route are mutable. The remainder of the SRH (Flags, Tag, Segment List, and TLVs that do not change en route) are immutable while processing this SID.

##### **4.3.1.1. SRH Processing**



```
S01. When an SRH is processed {
S02.   If Segments Left is equal to zero {
S03.     Proceed to process the next header in the packet,
        whose type is identified by the Next Header field in
        the Routing header.
S04.   }
S05.   Else {
S06.     If local configuration requires TLV processing {
S07.       Perform TLV processing (see TLV Processing)
S08.     }
S09.     max_last_entry = ( Hdr Ext Len / 2 ) - 1
S10.     If ((Last Entry > max_last_entry) or
S11.       (Segments Left is greater than (Last Entry+1)) {
S12.       Send an ICMP Parameter Problem, Code 0, message to
        the Source Address, pointing to the Segments Left
        field, and discard the packet.
S13.     }
S14.     Else {
S15.       Decrement Segments Left by 1.
S16.       Copy Segment List[Segments Left] from the SRH to the
        destination address of the IPv6 header.
S17.       If the IPv6 Hop Limit is less than or equal to 1 {
S18.         Send an ICMP Time Exceeded -- Hop Limit Exceeded in
        Transit message to the Source Address and discard
        the packet.
S19.       }
S20.       Else {
S21.         Decrement the Hop Limit by 1
S22.         Resubmit the packet to the IPv6 module for transmission
        to the new destination.
S23.       }
S24.     }
S25.   }
S26. }
```

#### **4.3.1.1.1. TLV Processing**

Local configuration determines how TLVs are to be processed when the Active Segment is a local SID defined in this document. The definition of local configuration is outside the scope of this document.

For illustration purpose only, two example local configurations that may be associated with a SID are provided below.



Example 1:

```
For any packet received from interface I2
  Skip TLV processing
```

Example 2:

```
For any packet received from interface I1
  If first TLV is HMAC {
    Process the HMAC TLV
  }
  Else {
    Discard the packet
  }
```

#### **4.3.1.2. Upper-layer Header or No Next Header**

When processing the Upper-layer header of a packet matching a FIB entry locally instantiated as an SRv6 SID defined in this document.

```
IF (Upper-layer Header is IPv4 or IPv6) and
  local configuration permits {
  Perform IPv6 decapsulation
  Resubmit the decapsulated packet to the IPv4 or IPv6 module
}
ELSE {
  Send an ICMP parameter problem message to the Source Address and
  discard the packet. Error code (TBD by IANA) "SR Upper-layer
  Header Error", pointer set to the offset of the upper-layer
  header.
}
```

A unique error code allows an SR Source node to recognize an error in SID processing at an endpoint.

#### **4.3.2. FIB Entry is a Local Interface**

If the FIB entry represents a local interface, not locally instantiated as an SRv6 SID, the SRH is processed as follows:

If Segments Left is zero, the node must ignore the Routing header and proceed to process the next header in the packet, whose type is identified by the Next Header field in the Routing Header.

If Segments Left is non-zero, the node must discard the packet and send an ICMP Parameter Problem, Code 0, message to the packet's Source Address, pointing to the unrecognized Routing Type.



#### **4.3.3. FIB Entry Is A Non-Local Route**

Processing is not changed by this document.

#### **4.3.4. FIB Entry Is A No Match**

Processing is not changed by this document.

### **5. Intra SR Domain Deployment Model**

The use of the SIDs exclusively within the SR Domain and solely for packets of the SR Domain is an important deployment model.

This enables the SR Domain to act as a single routing system.

This section covers:

- o securing the SR Domain from external attempt to use its SIDs
- o SR Domain as a single system with delegation between components
- o handling packets of the SR Domain

#### **5.1. Securing the SR Domain**

Nodes outside the SR Domain are not trusted: they cannot directly use the SID's of the domain. This is enforced by two levels of access control lists:

1. Any packet entering the SR Domain and destined to a SID within the SR Domain is dropped. This may be realized with the following logic, other methods with equivalent outcome are considered compliant:
  - \* allocate all the SID's from a block S/s
  - \* configure each external interface of each edge node of the domain with an inbound infrastructure access list (IACL) which drops any incoming packet with a destination address in S/s
  - \* Failure to implement this method of ingress filtering exposes the SR Domain to source routing attacks as described and referenced in [[RFC5095](#)]
2. The distributed protection in #1 is complemented with per node protection, dropping packets to SIDs from source addresses outside the SR Domain. This may be realized with the following





logic, other methods with equivalent outcome are considered compliant:

- \* assign all interface addresses from prefix A/a
- \* at node k, all SIDs local to k are assigned from prefix Sk/sk
- \* configure each internal interface of each SR node k in the SR Domain with an inbound IACL which drops any incoming packet with a destination address in Sk/sk if the source address is not in A/a.

## **5.2. SR Domain as a single system with delegation among components**

All intra SR Domain packets are of the SR Domain. The IPv6 header is originated by a node of the SR Domain, and is destined to a node of the SR Domain.

All inter domain packets are encapsulated for the part of the packet journey that is within the SR Domain. The outer IPv6 header is originated by a node of the SR Domain, and is destined to a node of the SR Domain.

As a consequence, any packet within the SR Domain is of the SR Domain.

The SR Domain is a system in which the operator may want to distribute or delegate different operations of the outer most header to different nodes within the system.

An operator of an SR domain may choose to delegate SRH addition to a host node within the SR domain, and validation of the contents of any SRH to a more trusted router or switch attached to the host. Consider a top of rack switch (T) connected to host (H) via interface (I). H receives an SRH (SRH1) with a computed HMAC via some SDN method outside the scope of this document. H classifies traffic it sources and adds SRH1 to traffic requiring a specific SLA. T is configured with an IACL on I requiring verification of the SRH for any packet destined to the SID block of the SR Domain (S/s). T checks and verifies that SRH1 is valid, contains an HMAC TLV and verifies the HMAC.

An operator of the SR Domain may choose to have all segments in the SR Domain verify the HMAC. This mechanism would verify that the SRH segment list is not modified while traversing the SR Domain.



### **5.3. MTU Considerations**

An SR Domain ingress edge node encapsulates packets traversing the SR Domain, and needs to consider the MTU of the SR Domain. Within the SR Domain, well known mitigation techniques are RECOMMENDED, such as deploying a greater MTU value within the SR Domain than at the ingress edges.

### **5.4. ICMP Error Processing**

ICMP error packets generated within the SR Domain are sent to source nodes within the SR Domain. The invoking packet in the ICMP error message may contain an SRH. Since the destination address of a packet with an SRH changes as each segment is processed, it may not be the destination used by the socket or application that generated the invoking packet.

For the source of an invoking packet to process the ICMP error message, the correct destination address must be determined. The following logic is used to determine the destination address for use by protocol error handlers.

- o Walk all extension headers of the invoking IPv6 packet to the routing extension header preceding the upper layer header.
  - \* If routing header is type 4 (SRH)
    - + Use the SID at Segment List[0] as the destination address of the invoking packet.

ICMP errors are then processed by upper layer transports as defined in [[RFC4443](#)].

For IP packets encapsulated in an outer IPv6 header, ICMP error handling is as defined in [[RFC2473](#)].

### **5.5. Load Balancing and ECMP**

For any inter domain packet, the SR Source node MUST impose a flow label computed based on the inner packet. The computation of the flow label is as recommended in [[RFC6438](#)] for the sending Tunnel End Point.

For any intra domain packet, the SR Source node SHOULD impose a flow label computed as described in [[RFC6437](#)] to assist ECMP load balancing at transit nodes incapable of computing a 5-tuple beyond the SRH.



At any transit node within an SR domain, the flow label MUST be used as defined in [RFC6438] to calculate the ECMP hash toward the destination address. If flow label is not used, the transit node would likely hash all packets between a pair of SR Edge nodes to the same link.

At an SR segment endpoint node, the flow label MUST be used as defined in [RFC6438] to calculate any ECMP hash used to forward the processed packet to the next segment.

## 5.6. Other Deployments

Other deployment models and their implications on security, MTU, HMAC, ICMP error processing and interaction with other extension headers are outside the scope of this document.

## 6. Illustrations

This section provides illustrations of SRv6 packet processing at SR source, transit and SR segment endpoint nodes.

### 6.1. Abstract Representation of an SRH

For a node  $k$ , its IPv6 address is represented as  $A_k$ , its SRv6 SID is represented as  $S_k$ .

IPv6 headers are represented as the tuple of (source, destination). For example, a packet with source address  $A_1$  and destination address  $A_2$  is represented as  $(A_1, A_2)$ . The payload of the packet is omitted.

An SR Policy is a list of segments. A list of segments is represented as  $\langle S_1, S_2, S_3 \rangle$  where  $S_1$  is the first SID to visit,  $S_2$  is the second SID to visit and  $S_3$  is the last SID to visit.

$(SA, DA) (S_3, S_2, S_1; SL)$  represents an IPv6 packet with:

- o Source Address is SA, Destination Addresses is DA, and next-header is SRH.
- o SRH with SID list  $\langle S_1, S_2, S_3 \rangle$  with SegmentsLeft = SL.
- o Note the difference between the  $\langle \rangle$  and  $()$  symbols.  $\langle S_1, S_2, S_3 \rangle$  represents a SID list where the leftmost segment is the first segment. Whereas,  $(S_3, S_2, S_1; SL)$  represents the same SID list but encoded in the SRH Segment List format where the leftmost segment is the last segment. When referring to an SR policy in a high-level use-case, it is simpler to use the  $\langle S_1, S_2, S_3 \rangle$



notation. When referring to an illustration of detailed behavior, the (S3, S2, S1; SL) notation is more convenient.

At its SR Policy headend, the Segment List <S1,S2,S3> results in SRH (S3,S2,S1; SL=2) represented fully as:

Segments Left=2  
Last Entry=2  
Flags=0  
Tag=0  
Segment List[0]=S3  
Segment List[1]=S2  
Segment List[2]=S1

**6.2. Example Topology**

The following topology is used in examples below:

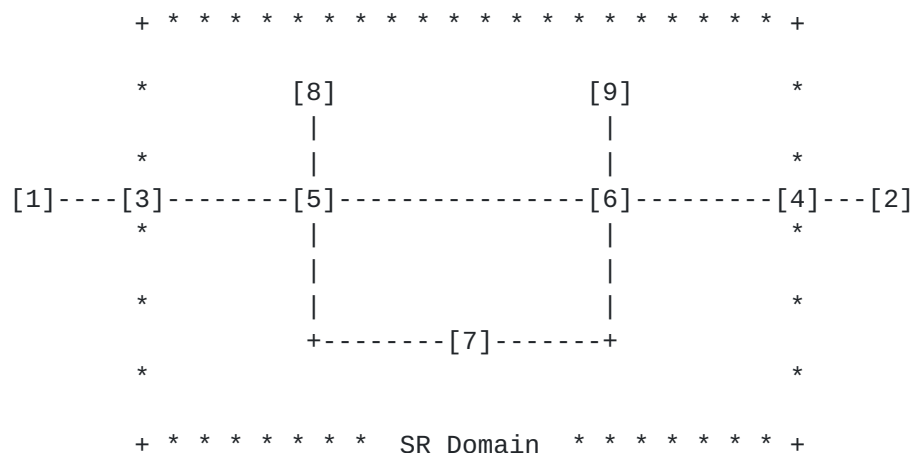


Figure 3

- o 3 and 4 are SR Domain edge routers
- o 5, 6, and 7 are all SR Domain routers
- o 8 and 9 are hosts within the SR Domain
- o 1 and 2 are hosts outside the SR Domain
- o The SR domain is secured as per [Section 5.1](#) and no external packet can enter the domain with a destination address equal to a segment of the domain.





### **6.3. Source SR Node**

#### **6.3.1. Intra SR Domain Packet**

When host 8 sends a packet to host 9 via an SR Policy <S7,A9> the packet is

P1: (A8,S7)(A9,S7; SL=1)

##### **6.3.1.1. Reduced Variant**

When host 8 sends a packet to host 9 via an SR Policy <S7,A9> and it wants to use a reduced SRH, the packet is

P2: (A8,S7)(A9; SL=1)

#### **6.3.2. Inter SR Domain Packet - Transit**

When host 1 sends a packet to host 2, the packet is

P3: (A1,A2)

The SR Domain ingress router 3 receives P3 and steers it to SR Domain egress router 4 via an SR Policy <S7, S4>. Router 3 encapsulates the received packet P3 in an outer header with an SRH. The packet is

P4: (A3, S7)(S4, S7; SL=1)(A1, A2)

If the SR Policy contains only one segment (the egress router 4), the ingress Router 3 encapsulates P3 into an outer header (A3, S4). The packet is

P5: (A3, S4)(A1, A2)

##### **6.3.2.1. Reduced Variant**

The SR Domain ingress router 3 receives P3 and steers it to SR Domain egress router 4 via an SR Policy <S7, S4>. If router 3 wants to use a reduced SRH, Router 3 encapsulates the received packet P3 in an outer header with a reduced SRH. The packet is

P6: (A3, S7)(S4; SL=1)(A1, A2)

#### **6.3.3. Inter SR Domain Packet - Internal to External**

When host 8 sends a packet to host 1, the packet is encapsulated for the portion of its journey within the SR Domain. From 8 to 3 the packet is



P7: (A8,S3)(A8,A1)

In the opposite direction, the packet generated from 1 to 8 is

P8: (A1,A8)

At node 3 P8 is encapsulated for the portion of its journey within the SR domain, with the outer header destined to segment S8.

Resulting in

P9: (A3,S8)(A1,A8)

At node 8 the outer IPv6 header is removed by S8 processing, then processed again when received by A8.

#### **6.4. Transit Node**

Nodes 5 acts as transit nodes for packet P1, and sends packet

P1: (A8,S7)(A9,S7;SL=1)

on the interface toward node 7.

#### **6.5. SR Segment Endpoint Node**

Node 7 receives packet P1 and, using the logic in [Section 4.3.1](#), sends packet

P7: (A8,A9)(A9,S7; SL=0)

on the interface toward router 6.

#### **6.6. Delegation of Function with HMAC Verification**

This section describes how a function may be delegated within the SR Domain to non SR source nodes. In the following sections consider a host 8 connected to a top of rack 5.

##### **6.6.1. SID List Verification**

An operator may prefer to add the SRH at source 8, while 5 verifies the SID list is valid.

For illustration purpose, an SDN controller provides 8 an SRH terminating at node 9, with segment list <S5,S7,S6,A9>, and HMAC TLV computed for the SRH. The HMAC key is shared with 5, node 8 does not know the key. Node 5 is configured with an IACL applied to the



interface connected to 8, requiring HMAC verification for any packet destined to S/s.

Node 8 originates packets with the received SRH with HMAC TLV.

P15:(A8,S5)(A9,S6,S7,S5;SL=3;HMAC)

Node 5 receives and verifies the HMAC for the SRH, then forwards the packet to the next segment

P16:(A8,S7)(A9,S6,S7,S5;SL=2;HMAC)

Node 6 receives

P17:(A8,S6)(A9,S6,S7,S5;SL=1;HMAC)

Node 9 receives

P18:(A8,A9)(A9,S6,S7,S5;SL=0;HMAC)

This use of an HMAC is particularly valuable within an enterprise based SR Domain [[SRN](#)].

## 7. Security Considerations

This section reviews security considerations related to the SRH, given the SRH processing and deployment models discussed in this document.

As described in [Section 5](#), it is necessary to filter packets ingress to the SR Domain, destined to SIDs within the SR Domain (i.e., bearing a SID in the destination address). This ingress filtering is via an IACL at SR Domain ingress border nodes. Additional protection is applied via an IACL at each SR Segment Endpoint node, filtering packets not from within the SR Domain, destined to SIDs in the SR Domain. ACLs are easily supported for small numbers of prefixes, making summarization important, and when the prefixes requiring filtering is kept to a seldom changing set.

Additionally, ingress filtering of IPv6 source addresses as recommended in [BCP38](#) SHOULD be used.

### 7.1. Source Routing Attacks

[RFC5095] deprecates the Type 0 Routing header due to a number of significant attacks that are referenced in that document. Such attacks include bypassing filtering devices, reaching otherwise



unreachable Internet systems, network topology discovery, bandwidth exhaustion, and defeating anycast.

Because this document specifies that the SRH is for use within an SR domain protected by ingress filtering via IACLs; such attacks cannot be mounted from outside an SR Domain. As specified in this document, SR Domain ingress edge nodes drop packets entering the SR Domain destined to segments within the SR Domain.

Additionally, this document specifies the use of IACL on SR Segment Endpoint nodes within the SR Domain to limit the source addresses permitted to send packets to a SID in the SR Domain.

Such attacks may, however, be mounted from within the SR Domain, from nodes permitted to source traffic to SIDs in the domain. As such, these attacks and other known attacks on an IP network (e.g. DOS/DDOS, topology discovery, man-in-the-middle, traffic interception/siphoning), can occur from compromised nodes within an SR Domain.

## **7.2. Service Theft**

Service theft is defined as the use of a service offered by the SR Domain by a node not authorized to use the service.

Service theft is not a concern within the SR Domain as all SR Source nodes and SR segment endpoint nodes within the domain are able to utilize the services of the Domain. If a node outside the SR Domain learns of segments or a topological service within the SR domain, IACL filtering denies access to those segments.

## **7.3. Topology Disclosure**

The SRH is unencrypted and may contain SIDs of some intermediate SR-nodes in the path towards the destination within the SR Domain. If packets can be snooped within the SR Domain, the SRH may reveal topology, traffic flows, and service usage.

This is applicable within an SR Domain but the disclosure is less relevant as an attacker has other means of learning topology, flows, and service usage.

## **7.4. ICMP Generation**

The generation of ICMPv6 error messages may be used to attempt denial-of-service attacks by sending an error-causing destination address or SRH in back-to-back packets. An implementation that correctly follows [Section 2.4 of \[RFC4443\]](#) would be protected by the ICMPv6 rate-limiting mechanism.





**7.5. Applicability of AH**

The SR Domain is a trusted domain, as defined in [\[RFC8402\] Section 2](#) and [Section 8.2](#). The SR Source is trusted to add an SRH (optionally verified via the HMAC TLV in this document), and segments advertised within the domain are trusted to be accurate and advertised by trusted sources via a secure control plane. As such the SR Domain does not rely on the Authentication Header (AH) as defined in [\[RFC4302\]](#) to secure the SRH.

The use of SRH with AH by an SR source node, and processing at a SR segment endpoint node, is not defined in this document. Future documents may define use of SRH with AH and its processing.

**8. IANA Considerations**

This document makes the following registrations in the Internet Protocol Version 6 (IPv6) Parameters "Routing Type" registry maintained by IANA:

Suggested Value	Description	Reference
4	Segment Routing Header (SRH)	This document

This document makes the following registrations in "Type 4 - Parameter Problem" message of the "Internet Control Message Protocol version 6 (ICMPv6) Parameters" registry maintained by IANA:

CODE	NAME/DESCRIPTION
TBD IANA	SR Upper-layer Header Error

This section provides guidance to the Internet Assigned Numbers Authority (IANA) regarding registration of values related to the SRH, in accordance with [BCP 26](#), [\[RFC8126\]](#).

The following terms are used here with the meanings defined in [BCP 26](#): "namespace", "assigned value", "registration".

The following policies are used here with the meanings defined in [BCP 26](#): "Private Use", "First Come First Served", "Expert Review", "Specification Required", "IETF Consensus", "Standards Action".

For registration requests where a Designated Expert should be consulted, the responsible IESG area director should appoint the Designated Expert. The intention is that any allocation will be accompanied by a published RFC. In order to allow for the allocation



of values prior to the RFC being approved for publication, the Designated Expert can approve allocations once it seems clear that an RFC will be published. The Designated expert will post a request to the 6man WG mailing list (or a successor designated by the Area Director) for comment and review, including an Internet-Draft. Before a period of 30 days has passed, the Designated Expert will either approve or deny the registration request and publish a notice of the decision to the 6man WG mailing list or its successor, as well as informing IANA. A denial notice must be justified by an explanation, and in the cases where it is possible, concrete suggestions on how the request can be modified so as to become acceptable should be provided.

### **8.1. Segment Routing Header Flags Register**

This document requests the creation of a new IANA managed registry to identify SRH Flags Bits. The registration procedure is "Expert Review" as defined in [[RFC8126](#)]. Suggested registry name is "Segment Routing Header Flags". Flags is 8 bits.

### **8.2. Segment Routing Header TLVs Register**

This document requests the creation of a new IANA managed registry to identify SRH TLVs. The registration procedure is "Expert Review" as defined in [[RFC8126](#)]. Suggested registry name is "Segment Routing Header TLVs". A TLV is identified through an unsigned 8 bit codepoint value, with assigned values 0-127 for TLVs that do not change en route, and 128-255 for TLVs that may change en route. The following codepoints are defined in this document:

Assigned Value	Description	Reference
0	Pad1 TLV	This document
1	Reserved	This document
2	Reserved	This document
3	Reserved	This document
4	PadN TLV	This document
5	HMAC TLV	This document
6	Reserved	This document
124-126	Experimentation and Test	This document
127	Reserved	This document
252-254	Experimentation and Test	This document
255	Reserved	This document

Values 1,2,3,6 were defined in draft versions of this specification and are Reserved for backwards compatibility with early implementations and should not be reassigned. Values 127 and 255 are



Reserved to allow for expansion of the Type field in future specifications if needed.

## **9. Implementation Status**

This section is to be removed prior to publishing as an RFC.

See [[I-D.matsushima-spring-srv6-deployment-status](#)] for updated deployment and interoperability reports.

### **9.1. Linux**

Name: Linux Kernel v4.14

Status: Production

Implementation: adds SRH, performs END processing, supports HMAC TLV

Details: <https://irtf.org/anrw/2017/anrw17-final3.pdf> and [[I-D.filsfils-spring-srv6-interop](#)]

### **9.2. Cisco Systems**

Name: IOS XR and IOS XE

Status: Production (IOS XR), Pre-production (IOS XE)

Implementation: adds SRH, performs END processing, no TLV processing

Details: [[I-D.filsfils-spring-srv6-interop](#)]

### **9.3. FD.io**

Name: VPP/Segment Routing for IPv6

Status: Production

Implementation: adds SRH, performs END processing, no TLV processing

Details: [https://wiki.fd.io/view/VPP/Segment\\_Routing\\_for\\_IPv6](https://wiki.fd.io/view/VPP/Segment_Routing_for_IPv6) and [[I-D.filsfils-spring-srv6-interop](#)]

### **9.4. Barefoot**

Name: Barefoot Networks Tofino NPU

Status: Prototype



Implementation: performs END processing, no TLV processing

Details: [[I-D.filsfils-spring-srv6-interop](#)]

### **9.5. Juniper**

Name: Juniper Networks Trio and vTrio NPU's

Status: Prototype & Experimental

Implementation: SRH insertion mode, Process SID where SID is an interface address, no TLV processing

### **9.6. Huawei**

Name: Huawei Systems VRP Platform

Status: Production

Implementation: adds SRH, performs END processing, no TLV processing

## **10. Contributors**

Kamran Raza, Zafar Ali, Brian Field, Daniel Bernier, Ida Leung, Jen Linkova, Ebben Aries, Tomoya Kosugi, Eric Vyncke, David Lebrun, Dirk Steinberg, Robert Raszuk, Dave Barach, John Brzozowski, Pierre Francois, Nagendra Kumar, Mark Townsley, Christian Martin, Roberta Maglione, James Connolly, Aloys Augustin contributed to the content of this document.

## **11. Acknowledgements**

The authors would like to thank Ole Troan, Bob Hinden, Ron Bonica, Fred Baker, Brian Carpenter, Alexandru Petrescu, Punit Kumar Jaiswal, and David Lebrun for their comments to this document.

## **12. References**

### **12.1. Normative References**

[FIPS180-4]

National Institute of Standards and Technology, "FIPS 180-4 Secure Hash Standard (SHS)", March 2012, <<http://csrc.nist.gov/publications/fips/fips180-4/fips-180-4.pdf>>.





- [RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", [BCP 14](#), [RFC 2119](#), DOI 10.17487/RFC2119, March 1997, <<https://www.rfc-editor.org/info/rfc2119>>.
- [RFC2473] Conta, A. and S. Deering, "Generic Packet Tunneling in IPv6 Specification", [RFC 2473](#), DOI 10.17487/RFC2473, December 1998, <<https://www.rfc-editor.org/info/rfc2473>>.
- [RFC4302] Kent, S., "IP Authentication Header", [RFC 4302](#), DOI 10.17487/RFC4302, December 2005, <<https://www.rfc-editor.org/info/rfc4302>>.
- [RFC5095] Abley, J., Savola, P., and G. Neville-Neil, "Deprecation of Type 0 Routing Headers in IPv6", [RFC 5095](#), DOI 10.17487/RFC5095, December 2007, <<https://www.rfc-editor.org/info/rfc5095>>.
- [RFC6407] Weis, B., Rowles, S., and T. Hardjono, "The Group Domain of Interpretation", [RFC 6407](#), DOI 10.17487/RFC6407, October 2011, <<https://www.rfc-editor.org/info/rfc6407>>.
- [RFC8174] Leiba, B., "Ambiguity of Uppercase vs Lowercase in [RFC 2119](#) Key Words", [BCP 14](#), [RFC 8174](#), DOI 10.17487/RFC8174, May 2017, <<https://www.rfc-editor.org/info/rfc8174>>.
- [RFC8200] Deering, S. and R. Hinden, "Internet Protocol, Version 6 (IPv6) Specification", STD 86, [RFC 8200](#), DOI 10.17487/RFC8200, July 2017, <<https://www.rfc-editor.org/info/rfc8200>>.
- [RFC8402] Filsfils, C., Ed., Previdi, S., Ed., Ginsberg, L., Decraene, B., Litkowski, S., and R. Shakir, "Segment Routing Architecture", [RFC 8402](#), DOI 10.17487/RFC8402, July 2018, <<https://www.rfc-editor.org/info/rfc8402>>.

## **12.2. Informative References**

- [I-D.filsfils-spring-srv6-interop]  
Filsfils, C., Clad, F., Camarillo, P., Abdelsalam, A., Salsano, S., Bonaventure, O., Horn, J., and J. Liste, "SRV6 interoperability report", [draft-filsfils-spring-srv6-interop-02](#) (work in progress), March 2019.



- [I-D.matsushima-spring-srv6-deployment-status]  
Matsushima, S., Filsfils, C., Ali, Z., and Z. Li, "SRV6 Implementation and Deployment Status", [draft-matsushima-spring-srv6-deployment-status-01](#) (work in progress), May 2019.
- [RFC2104] Krawczyk, H., Bellare, M., and R. Canetti, "HMAC: Keyed-Hashing for Message Authentication", [RFC 2104](#), DOI 10.17487/RFC2104, February 1997, <<https://www.rfc-editor.org/info/rfc2104>>.
- [RFC4443] Conta, A., Deering, S., and M. Gupta, Ed., "Internet Control Message Protocol (ICMPv6) for the Internet Protocol Version 6 (IPv6) Specification", STD 89, [RFC 4443](#), DOI 10.17487/RFC4443, March 2006, <<https://www.rfc-editor.org/info/rfc4443>>.
- [RFC5308] Hopps, C., "Routing IPv6 with IS-IS", [RFC 5308](#), DOI 10.17487/RFC5308, October 2008, <<https://www.rfc-editor.org/info/rfc5308>>.
- [RFC5340] Coltun, R., Ferguson, D., Moy, J., and A. Lindem, "OSPF for IPv6", [RFC 5340](#), DOI 10.17487/RFC5340, July 2008, <<https://www.rfc-editor.org/info/rfc5340>>.
- [RFC6437] Amante, S., Carpenter, B., Jiang, S., and J. Rajahalme, "IPv6 Flow Label Specification", [RFC 6437](#), DOI 10.17487/RFC6437, November 2011, <<https://www.rfc-editor.org/info/rfc6437>>.
- [RFC6438] Carpenter, B. and S. Amante, "Using the IPv6 Flow Label for Equal Cost Multipath Routing and Link Aggregation in Tunnels", [RFC 6438](#), DOI 10.17487/RFC6438, November 2011, <<https://www.rfc-editor.org/info/rfc6438>>.
- [RFC8126] Cotton, M., Leiba, B., and T. Narten, "Guidelines for Writing an IANA Considerations Section in RFCs", [BCP 26](#), [RFC 8126](#), DOI 10.17487/RFC8126, June 2017, <<https://www.rfc-editor.org/info/rfc8126>>.
- [SRN] Lebrun, D., Jadin, M., Clad, F., Filsfils, C., and O. Bonaventure, "Software Resolved Networks: Rethinking Enterprise Networks with IPv6 Segment Routing", 2018, <<https://inl.info.ucl.ac.be/system/files/sosr18-final15-embedfonts.pdf>>.



Authors' Addresses

Clarence Filsfils (editor)  
Cisco Systems, Inc.  
Brussels  
BE

Email: [cfilsfil@cisco.com](mailto:cfilsfil@cisco.com)

Darren Dukes (editor)  
Cisco Systems, Inc.  
Ottawa  
CA

Email: [ddukes@cisco.com](mailto:ddukes@cisco.com)

Stefano Previdi  
Huawei  
Italy

Email: [stefano@previdi.net](mailto:stefano@previdi.net)

John Leddy  
Individual  
US

Email: [john@leddy.net](mailto:john@leddy.net)

Satoru Matsushima  
Softbank

Email: [satoru.matsushima@g.softbank.co.jp](mailto:satoru.matsushima@g.softbank.co.jp)

Daniel Voyer  
Bell Canada

Email: [daniel.voyer@bell.ca](mailto:daniel.voyer@bell.ca)

