Jerry Ash Bur Goode Jim Hand AT&T

Raymond Zhang Infonet Services Corporation

April, 2004

Requirements for Header Compression over MPLS

Status of this Memo:

By submitting this Internet-Draft, we certify that any applicable patent or other IPR claims of which we are aware have been disclosed, and any of which we become aware will be disclosed, in accordance with <u>RFC 3668</u> (<u>BCP 79</u>).

By submitting this Internet-Draft, we accept the provisions of <u>Section 4 of RFC 3667</u> (<u>BCP 78</u>).

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF), its areas, and its working groups. Note that other groups may also distribute working documents as Internet-Drafts.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or cite them other than as "work in progress".

The list of current Internet-Drafts can be accessed at http://www.ietf.org/ietf/lid-abstracts.txt.

The list of Internet-Draft Shadow Directories can be accessed at http://www.ietf.org/shadow.html.

This document is a submission of the IETF AVT WG. Comments should be directed to the AVT WG mailing list, avt@ietf.org.

Copyright Notice:

Copyright (C) The Internet Society (2004). All Rights Reserved.

Ash

<<u>draft-ietf-avt-hc-mpls-reqs-00.txt</u>>

[Page 1]

Abstract:

VoIP typically uses the encapsulation voice/RTP/UDP/IP. When MPLS labels are added, this becomes voice/RTP/UDP/IP/MPLS-labels, where, for example, the packet header is at least 48 bytes, while the voice payload is often no more than 30 bytes. Header compression can significantly reduce the overhead through various compression mechanisms, such as enhanced compressed RTP (ECRTP) and robust header compression (ROHC). We consider using MPLS to route compressed packets over an MPLS LSP without compression/decompression cycles at each router. This approach can increase the bandwidth efficiency as well as processing scalability of the maximum number of simultaneous flows that use header compression at each router. In the draft we give a problem statement, goals and requirements, and an example scenario.

Table of Contents:

<u>1</u> . Introduction	
<u>2</u> . Problem Statement	<u>3</u>
$\underline{3}$. Goals & Requirements	<u>4</u>
$\underline{4}$. Candidate Solution Methods & Needs	<u>5</u>
<u>5</u> . Example Scenario	<u>6</u>
<u>6</u> . Security Considerations	<u>7</u>
$\underline{7}$. IANA Considerations	<u>7</u>
<u>8</u> . Normative References	<u>7</u>
$\underline{9}$. Informative References	<u>7</u>
<u>10</u> . Authors' Addresses	<u>8</u>
<u>11</u> . Full Copyright Statement	<u>8</u>

Specification of Requirements

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [<u>RFC 2119</u>].

1. Introduction

Voice over IP (VoIP) typically uses the encapsulation voice/RTP/UDP/IP. When MPLS labels [MPLS-ARCH] are added, this becomes voice/RTP/UDP/IP/MPLS-labels. For an MPLS VPN (e.g., [MPLS-VPN], the packet header is at least 48 bytes, while the voice payload is often no more than 30 bytes, for example. The interest in header compression (HC) is to exploit the possibility of significantly reducing the overhead through various compression mechanisms, such as with enhanced compressed RTP [ECRTP] or robust header compression [ROHC], and also to increase scalability of HC. We consider using MPLS to route compressed packets over an MPLS LSP (label switched path) without compression/decompression cycles at each router. Such an HC over MPLS capability can increase bandwidth efficiency as well as the processing scalability of the maximum number of simultaneous flows which use HC at each router.

To implement HC over MPLS, the ingress router/gateway would have to apply the HC algorithm to the IP packet, the compressed packet routed on an MPLS LSP using MPLS labels, and the compressed header would be decompressed at the egress router/gateway where the HC session

<<u>draft-ietf-avt-hc-mpls-reqs-00.txt</u>>

Ash

[Page 2]

terminates. Figure 1 illustrates an HC over MPLS session established on an LSP that crosses several routers, from R1/HC --> R2 --> R3 --> R4/HD, where R1/HC is the ingress router where HC is performed, and R4/HD is the egress router where header decompression (HD) is done. HC of the RTP/UDP/IP header is performed at R1/HC, and the compressed packets are routed using MPLS labels from R1/HC to R2, to R3, and finally to R4/HD, without further decompression/recompression cycles. The RTP/UDP/IP header is decompressed at R4/HD and can be forwarded to other routers, as needed.

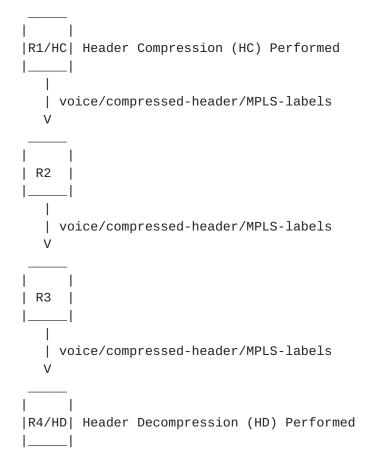


Figure 1. Example of Header Compression over MPLS over Routers R1-->R4

In the example scenario, HC therefore takes place between R1 and R4, and the MPLS path transports voice/compressed-header/MPLS-labels instead of voice/RTP/UDP/IP/MPLS-labels, typically saving 30 octets or more per packet. The MPLS label stack and link-layer headers are not compressed. A signaling method is needed to set up a correspondence between the ingress and egress routers of the HC over MPLS session. In <u>Section 2</u> we give a problem statement, in <u>Section 3</u> we give goals and requirements, and in <u>Section 4</u> we give an example scenario.

2. Problem Statement

As described in the introduction, HC over MPLS can significantly reduce the header overhead through HC mechanisms. The need for HC may be important on low-speed links where bandwidth is more scarce, but it could also be important on backbone facilities, especially where costs are high (e.g., some global cross-sections). VoIP typically will use voice compression mechanisms (e.g., G.729) on low-speed and international routes, in order to conserve bandwidth. With HC,

Ash <<u>draft-ietf-avt-hc-mpls-reqs-00.txt</u>> [Page 3]

significantly more bandwidth could be saved. For example, carrying uncompressed headers for the entire voice load of a large domestic network with 300 million or more calls per day could consume on the order of about 20-40 gigabits-per-second on the backbone network for headers alone. This overhead could translate into considerable bandwidth capacity.

The claim is often made that once fiber is in place, increasing the bandwidth capacity is inexpensive, nearly 'free'. This may be true in some cases, however, on some international cross-sections, especially, facility/transport costs are very high and saving bandwidth on such backbone links is very worthwhile. Decreasing the backbone bandwidth is needed in some areas of the world where bandwidth is very expensive. It is also important in almost all locations to decrease the bandwidth consumption on low-speed links. So although bandwidth is getting cheaper, the value of compression does not go away. It should be further noted that IPv6 will increase the size of headers, and therefore increase the importance of HC for RTP flows.

While hop-by-hop HC could be applied to decrease bandwidth requirements, that implies a processing requirement for compression-decompression cycles at every router hop, which does not scale well for large voice traffic loads. The maximum number of cRTP flows is about 30-50 for a typical customer premise router, depending upon its uplink speed and processing power, while the need may exceed 300-500 for a high-end case. Therefore, HC over MPLS seems to be a viable alternative to get the compression benefits without introducing costly processing demands on the intermediate nodes. By using HC over MPLS, routers merely forward compressed packets without doing a decompression/recompression cycle, thereby increasing the maximum number of simultaneous compressed flows that routers can handle.

Therefore the proposal is to use existing HC techniques, together with MPLS labels, to make the transport of the RTP/UDP/IP headers more efficient over an MPLS network. However, at this time, there are no standards for HC over MPLS, and vendors have not implemented such techniques.

<u>3</u>. Goals & Requirements

The goals of HC over MPLS are as follows:

a. provide more efficient voice transport over MPLS networks,b. increase the scalability of HC to a large number of flows,c. not significantly increase packet delay, delay variation, or loss probability, andd. leverage existing work through use of standard protocols as much as possible.

Therefore the requirements for HC over MPLS are as follows:

a. MUST use existing protocols (e.g., [ECRTP], [ROHC]) to compress
RTP/UDP/IP headers, in order to provide for efficient transport,
tolerance to packet loss, and resistance to loss of session context.
b. MUST allow HC over an MPLS LSP, and thereby avoid hop-by-hop
compression/decompression cycles [e.g., ECRTP-MPLS-PROTO].

Ash <<u>draft-ietf-avt-hc-mpls-reqs-00.txt</u>> [Page 4]

c. MUST minimize incremental performance degradation due to increased delay, packet loss, and jitter.d. MUST use standard protocols to signal context identification and control information (e.g., [RSVP], [RSVP-TE], [LDP]).

It is necessary that the HC method be able to handle out-of-sequence packets. MPLS [MPLS-ARCH] enables 4-byte labels to be appended to IP packets to allow switching from the ingress label switched router (LSR) to the egress LSP on an LSP through an MPLS network. However, MPLS does not guarantee that packets will arrive in order at the egress LSR, since a number of things could cause packets to be delivered out of sequence. For example, a link failure could cause the LSP routing to change, due perhaps to an MPLS fast reroute taking place, or to the interior gateway protocol (IGP) and label distribution protocol (LDP) converging to another route, among other possible reasons. Other causes could include IGP reroutes due to 'loose hops' in the LSP, or BGP route changes reflecting back into IGP reroutes. HC algorithms may be able to handle reordering magnitudes on the order of about 10 packets, which may make the time required for IGP reconvergence (typically on the order of seconds) untenable for the HC algorithm. On the other hand, MPLS fast reroute may be fast enough (on the order of 50 ms. or less) for the HC algorithm to handle packet reordering. The issue of reordering needs to be further considered in the development of the HC over MPLS solution.

Resynchronization and performance also needs to be considered, since HC over MPLS can sometimes have multiple routers in the LSP. Tunneling a HC session over an MPLS LSP with multiple routers in the path will increase the round trip delay and the chance of packet loss, and HC contexts are invalidated due to packet loss. The HC error recovery mechanism can compound the problem when long round trip delays are involved.

4. Candidate Solution Methods & Needs

[cRTP] performs best with very low packet error rates on all hops of the path. When the cRTP decompressor context state gets out of synch with the compressor, it will drop packets associated with the context until the two states are resynchronized. To resynchronize context state at the two ends, the decompressor transmits the CONTEXT_STATE packet to the compressor, and the compressor transmits a FULL_HEADER packet to the decompressor.

[ECRTP] minimizes feedback based error recovery using CONTEXT_STATE packets to make cRTP more tolerant of packet loss, and minimize the need

to use the cRTP error recovery mechanism. ECRTP is therefore a candidate method to make HC over MPLS more tolerant of packet loss and to guard against frequent resynchronizations. ECRTP also is able to accommodate out-of-sequence packets. ECRTP protocol extensions may be required to identify FULL_HEADER, CONTEXT_STATE, and compressed packet types. [cRTP-ENCAP] specifies a separate link-layer packet type defined for HC. Using a separate link-layer packet type avoids the need to add extra bits to the compression header to identify the packet type. However, this approach does not extend well to MPLS encapsulation conventions [MPLS-ENCAP], in which a separate link-layer packet type translates into a separate LSP for each packet type. In order to extend ECRTP to HC over MPLS, each packet type defined in [ECRTP] would need to be identified in an appended packet type field in the ECRTP header.

Ash <<u>draft-ietf-avt-hc-mpls-reqs-00.txt</u>> [Page 5]

[ROHC] is also very tolerant of packet loss, and therefore is a candidate method to guard against frequent resynchronizations. ROHC also achieves a somewhat better level of compression as compared to ECRTP. However, ROHC currently does not accommodate packet reordering needed to protect against out-of-sequence packets that can occur on MPLS LSPs, and would need to be extended to do so. On the other hand, ROHC already has the capability to identify the packet type in the compression header, so no further extension is needed to identify packet type.

Extensions to MPLS signaling are needed to identify the LSP from HC to HD egress point, negotiate the HC algorithm used and protocol parameters, and negotiate the session context IDs (SCIDs) space between the ingress and egress routers on the MPLS LSP. For example, new objects need to be defined for [<u>RSVP-TE</u>] to signal the SCID spaces between the ingress and egress routers, and the HC algorithm used to determine the context; these HC packets then contain the SCID identified by using the RSVP-TE objects. It is also desirable to signal HC over MPLS tunnels with the label distribution protocol [LDP], since many <u>RFC2547</u> VPN [MPLS-VPN] implementations use LDP as the underlying LSP signaling mechanism, and LDP is very scalable. However, extensions to LDP would be needed to signal SCIDs between ingress and egress routers on HC over MPLS LSPs. For example, 'targeted LDP sessions' might be established for signaling SCIDs, or perhaps methods described in [LDP-PWE3] and [GVPLS] to signal pseudo-wires and multipoint-to-point LSPs might be extended to support signaling of SCIDs for HC over MPLS LSPs. These MPLS signaling protocol extensions need coordination with other working groups (e.g., MPLS).

<u>5</u>. Example Scenario

As illustrated in Figure 2, many VoIP flows are originated from customer sites, which are served by routers R1, R2 and R3, and terminated at several large customer call centers, which are served by R5, R6 and R7. R4 is a service-provider router, and all VoIP flows traverse R4. It is essential that the R4-R5, R4-R6, and R4-R7 low-speed links all use HC to allow a maximum number of simultaneous VoIP flows. To allow processing at R4 to handle the volume of simultaneous VoIP flows, it is desired to use HC over MPLS for these flows. With HC over MPLS, R4 does not need to do HC/HD for the flows to the call centers, enabling more scalability of the number of simultaneous VoIP flows with HC at R4.

voice/C-HDR/MPLS-labels	voice/C-HDR/MPLS-labels
R1/HC>	> R5/HD

1		I
voice/C-HDR/MPLS-labels		voice/C-HDR/MPLS-labels
R2/HC>	R4	> R6/HD
voice/C-HDR/MPLS-labels		voice/C-HDR/MPLS-labels
R3/HC>		> R7/HD

[Note: HC = header compression; C-HDR = compressed header; HD = header decompression]

Figure 2. Example Scenario for Application of HC over MPLS

Ash <<u>draft-ietf-avt-hc-mpls-reqs-00.txt</u>> [Page 6]

<u>6</u>. Security Considerations

The high processing load of HC makes HC a target for denial-of-service attacks. For example, an attacker could send a high bandwidth data stream through a network, with the headers in the data stream marked appropriately to cause HC to be applied. This would use large amounts of processing resources on the routers performing compression and decompression, and these processing resources might then be unavailable for other important functions on the router. This threat is not a new threat for HC, but is addressed and mitigated by HC over MPLS. That is, by reducing the need for performing compression and decompression cycles, as proposed in this draft, the risk of this type of denial-of-service attack is reduced.

7. IANA Considerations

No IANA actions are required.

8. Normative References

[cRTP] Casner, S., Jacobsen, V., "Compressing IP/UDP/RTP Headers for Low-Speed Serial Links", <u>RFC 2508</u>, February 1999.

[cRTP-ENCAP] Engan, M., Casner, S., Bormann, C., "IP Header Compression over PPP", <u>RFC 2509</u>, February 1999.

[ECRTP] Koren, T., et. al., "Compressing IP/UDP/RTP Headers on Links with High Delay, Packet Loss, and Reordering," <u>RFC 3545</u>, July 2003.

[KEY] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", <u>RFC 2119</u>, March 1997.

[LDP] Andersson, L., et. al., "LDP Specification", <u>RFC 3036</u>, January 2001.

[MPLS-ARCH] Rosen, E., et. al., "Multiprotocol Label Switching Architecture," <u>RFC 3031</u>, January 2001.

[ROHC] Bormann, C., et. al., "Robust Header Compression (ROHC)," RFC 3091, July 2001.

[RSVP] Braden, R. et al., "Resource ReSerVation Protocol (RSVP) --Version 1, Functional Specification", <u>RFC 2205</u>, September 1997.

[RSVP-TE] Awduche, D., et. al., "RSVP-TE: Extensions to RSVP for LSP Tunnels", <u>RFC 3209</u>, December 2001.

<u>9</u>. Informative References

[ECRTP-MPLS-PROTO] Ash, G., Goode, B., Hand, J., "Protocol Extensions for Header Compression over MPLS", work in progress.

[GVPLS] Radoaca, V., et. al., "GVPLS/LPE - Generalized VPLS Solution based on LPE Framework," work in progress.

Ash <<u>draft-ietf-avt-hc-mpls-reqs-00.txt</u>> [Page 7]

[LDP-PWE3] Martini, L., et. al., "Pseudowire Setup and Maintenance using LDP", work in progress.

[MPLS-ENCAP] Rosen, E., et. al., "MPLS Label Stack Encoding", <u>RFC 3032</u>, January 2001.

[MPLS-VPN] Rosen, E., Rekhter, Y., "BGP/MPLS VPNs", <u>RFC 2547</u>, March 1999.

10. Authors' Addresses

Jerry Ash AT&T Room MT D5-2A01 200 Laurel Avenue Middletown, NJ 07748, USA Phone: +1 732-420-4578 Email: gash@att.com

Bur Goode AT&T Phone: + 1 203-341-8705 E-mail: bgoode@att.com

Jim Hand Consultant E-mail: hand17@earthlink.net

Raymond Zhang Infonet Services Corporation 2160 E. Grand Ave. El Segundo, CA 90025 USA Email: raymond_zhangr@info.net

<u>11</u>. Full Copyright Statement

Copyright (C) The Internet Society (2004). This document is subject to the rights, licenses and restrictions contained in $\frac{\text{BCP}}{78}$ and except as set forth therein, the authors retain all their rights.

This document and the information contained herein are provided on an "AS IS" basis and THE CONTRIBUTOR, THE ORGANIZATION HE/SHE REPRESENTS OR IS SPONSORED BY (IF ANY), THE INTERNET SOCIETY AND THE INTERNET ENGINEERING TASK FORCE DISCLAIM ALL WARRANTIES, EXPRESS OR IMPLIED, INCLUDING BUT NOT LIMITED TO ANY WARRANTY THAT THE USE OF THE INFORMATION HEREIN WILL NOT INFRINGE ANY RIGHTS OR ANY IMPLIED WARRANTIES OF MERCHANTABILITY OR FITNESS FOR A PARTICULAR PURPOSE.

Ash <<u>draft-ietf-avt-hc-mpls-reqs-00.txt</u>> [Page 8]