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### Abstract

Procedure to handle host mobility in a layer 2 Network with EVPN control plane is defined as part of RFC 7432. EVPN has since evolved to find wider applicability across various IRB use cases that include distributing both MAC and IP reachability via a common EVPN control plane. MAC Mobility procedures defined in RFC 7432 are extensible to IRB use cases if a fixed 1:1 mapping between VM IP and MAC is assumed across VM moves. Generic mobility support for IP and MAC that allows these bindings to change across moves is required to support a broader set of EVPN IRB use cases, and requires further consideration. EVPN all-active multi-homing further introduces scenarios that require additional consideration from mobility perspective. This document enumerates a set of design considerations applicable to mobility across these EVPN IRB use cases and defines generic sequence number assignment procedures to address these IRB use cases.

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# 1. Requirements Language and Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

\*EVPN-IRB: A BGP-EVPN distributed control plane based integrated routing and bridging fabric overlay discussed in [EVPN-IRB]

\*Underlay: IP or MPLS fabric core network that provides IP or MPLS routed reachability between EVPN PEs.

\*Overlay: VPN or service layer network consisting of EVPN PEs OR VPN provider-edge (PE) switch-router devices that runs on top of an underlay routed core.

\*EVPN PE: A PE switch-router in a data-center fabric that runs overlay BGP-EVPN control plane and connects to overlay CE host devices. An EVPN PE may also be the first-hop layer-3 gateway for CE/host devices. This document refers to EVPN PE as a logical function in a data-center fabric. This EVPN PE function may be physically hosted on a top-of-rack switching device (ToR) OR at layer(s) above the ToR in the Clos fabric. An EVPN PE is typically also an IP or MPLS tunnel end-point for overlay VPN flow

\*Symmetric EVPN-IRB: An overlay fabric first-hop routing architecture as defined in [EVPN-IRB], wherein, overlay host-tohost routed inter-subnet flows are routed at both ingress and egress EVPN PEs.

\*Asymmetric EVPN-IRB: An overlay fabric first-hop routing architecture as defined in [EVPN-IRB], wherein, overlay host-tohost routed inter-subnet flows are routed and bridged at ingress PE and bridged at egress PEs.

\*ARP: Address Resolution Protocol [RFC 826]. ARP references in this document are equally applicable to ND as well.

\*ND: IPv6 Neighbor Discovery Protocol [RFC 4861].

\*Ethernet-Segment: physical Ethernet or LAG port that connects an access device to an EVPN PE, as defined in [RFC 7432].

\*ESI: Ethernet Segment Identifier as defined in [RFC 7432].

\*LAG: Layer-2 link-aggregation, also known as layer-2 bundle portchannel, or bond interface.

\*EVPN all-active multi-homing: PE-CE all-active multi-homing achieved via a multi-homed layer-2 LAG interface on a CE with member links to multiple PEs and related EVPN procedures on the PEs.

\*RT-2: EVPN route type 2 carrying both MAC and IP reachability.

\*RT-5: EVPN route type 5 carrying IP prefix reachability.

\*MAC-IP: IP association for a MAC, referred to in this document may be IPv4, IPv6 or both.

\*SYNC MAC route: In the context of EVPN multi-homing, this refers to a local MAC route SYNCed from another PE sharing the same ESI.

\*SYNC MAC-IP route: In the context of EVPN multi-homing, this refers to a local MAC-IP route SYNCed from another PE sharing the same ESI.

\*SYNC MAC sequence number: In the context of EVPN multi-homing, this refers to sequencee number received with a SYNC MAC route.

\*SYNC MAC-IP sequence number: In the context of EVPN multi-homing, this refers to sequencee number received with a SYNC MAC-IP route.

# 2. Introduction

EVPN-IRB enables capability to advertise both MAC and IP routes via a single MAC+IP RT-2 advertisement. MAC is imported into local bridge MAC table and enables L2 bridged traffic across the network overlay. IP is imported into the local ARP table in an asymmetric IRB design OR imported into the IP routing table in a symmetric IRB design, and enables routed traffic across the layer 2 network overlay. Please refer to [EVPN-IRB] for more background on EVPN IRB forwarding modes.

To support EVPN mobility procedure, a single sequence number mobility attribute is advertised with the combined MAC+IP route. A single sequence number advertised with the combined MAC+IP route to resolve both MAC and IP reachability implicitly assumes a 1:1 fixed mapping between IP and MAC. While a fixed 1:1 mapping between IP and MAC is a common use case that could be addressed via existing MAC mobility procedure, additional IRB scenarios need to be considered, that don't necessarily adhere to this assumption. Following IRB mobility scenarios are considered:

\*VM move results in VM IP and MAC moving together \*VM move results in VM IP moving to a new MAC association \*VM move results in VM MAC moving to a new IP association

While existing MAC mobility procedure can be leveraged for MAC+IP move in the first scenario, subsequent scenarios result in a new MAC- IP association. As a result, a single sequence number assigned independently per-[MAC, IP] is not sufficient to determine most recent reachability for both MAC and IP, unless the sequence number assignment algorithm is designed to allow for changing MAC-IP bindings across moves.

Purpose of this draft is to define additional sequence number assignment and handling procedures to adequately address generic mobility support across EVPN-IRB overlay use cases that allow MAC-IP bindings to change across VM moves and can support mobility for both MAC and IP components carried in an EVPN RT-2 for these use cases.

In addition, for hosts on an ESI multi-homed to multiple GW devices, additional procedure is proposed to ensure synchronized sequence number assignments across the multi-homing devices.

Content presented in this draft is independent of data plane encapsulation used in the overlay being MPLS or NVO Tunnels. It is also largely independent of the EVPN IRB solution being based on symmetric OR asymmetric IRB design as defined in [EVPN-INTER-SUBNET].

In addition to symmetric and asymmetric IRB, mobility solution for a routed overlay, where traffic to an end host in the overlay is always IP routed using EVPN RT-5 is also presented in this document.

To summarize, this draft covers mobility mobility for the following independent of the overlay encapsulation being MPLS or an NVO Tunnel:

\*Symmetric EVPN IRB overlay \*Asymmetric EVPN IRB overlay \*Routed EVPN overlay

### 2.1. Document Structure

Following sections of the document should be condidered informnative:

- \*section 4 and 5 provide the necessary background and problem statement being addressed in this document.
- \*section 6 lists the resulting design considerations for the document.

Following sections of the document should be condidered normative:

- \*section 8 describes the mobility and sequence number assigment procedures in an EVPN-IRB overlay required to address the scenarios described in section 4.
- \*section 9 describes the mobility procedures for a routed overlay nextwork as opposed to an IRB overlay.
- \*section 10 describes corresponding duplicate detection procedures for EVPN-IRB and routed overlays.

## 3. Optional MAC only RT-2

In an EVPN IRB scenario, where a single MAC+IP RT-2 advertisement carries both IP and MAC routes, a MAC only RT-2 advertisement is redundant for host MACs that are advertised via MAC+IP RT-2. As a result, a MAC only RT-2 is an optional route that may not be advertised from or received at an EVPN PE. This is an important consideration for mobility scenarios discussed in subsequent sections.

MAC only RT-2 may still be advertised for non-IP host MACs that are not advertised via MAC+IP RT-2.

### 4. Mobility Use Cases

This section describes the IRB mobility use cases considered in this document. Procedures to address them are covered later in section 6 and section 7.

\*Host move results in Host IP and MAC moving together

\*Host move results in Host IP moving to a new MAC association

\*Host move results in Host MAC moving to a new IP association

#### 4.1. Host MAC+IP Move

This is the baseline case, wherein a host move results in both host MAC and IP moving together with no change in MAC-IP binding across a move. Existing MAC mobility defined in RFC 7432 may be leveraged to apply to corresponding MAC+IP route to support this mobility scenario.

#### 4.2. Host IP Move to new MAC

This is the case, where a host move results in VM IP moving to a new MAC binding.

# 4.2.1. VM Reload

A host reload or an orchestrated host move that results in host being re-spawned at a new location may result in host getting a new MAC assignment, while maintaining existing IP address. This results in a host IP move to a new MAC binding:

IP-a, MAC-a ---> IP-a, MAC-b

# 4.2.2. MAC Sharing

This takes into account scenarios, where multiple hosts, each with a unique IP, may share a common MAC binding, and a host move results in a new MAC binding for the host IP.

As an example, hosts running on a single physical server, each with a unique IP, may share the same physical server MAC. In yet another scenario, an L2 access network may be behind a firewall, such that all hosts IPs on the access network are learnt with a common firewall MAC. In all such "shared MAC" use cases, multiple local MAC-IP ARP entries may be learnt with the same MAC. A host IP move, in such scenarios (for e.g., to a new physical server), could result in new MAC association for the host IP.

# 4.2.3. Problem

In both of the above scenarios, a combined MAC+IP EVPN RT-2 advertised with a single sequence number attribute implicitly assumes a fixed IP to MAC mapping. A host IP move to a new MAC breaks this assumption and results in a new MAC+IP route. If this new MAC+IP route is independently assigned a new sequence number, the sequence number can no longer be used to determine most recent host IP reachability in a symmetric EVPN-IRB design OR the most recent IP to MAC binding in an asymmetric EVPN-IRB design.

# +----+ | Underlay Network Fabric| +----+ +----+ +----+ +----+ +----+ +----+ +----+ | PE1 | | PE2 | | PE3 | | PE4 | | PE5 | | PE6 | +----+ +----+ +----+ +----+ +----+ +----+ $\land$ / / $\backslash$ / \ ESI-1 / \ ESI-2 / \ / \ ESI-3 / ESI-3 / +\---/+ +\---/+ +\---/+ | \ / | | \ / | | \ / | | \ / | | \ / | | \ / | +--+-+ +--+-+ +--+-+ | | | Server-MAC1 Server-MAC2 Server-MAC3 | | | [VM-IP1, VM-IP2] [VM-IP3, VM-IP4] [VM-IP5, VM-IP6]

#### Figure 1

As an example, consider a topology shown in Figure 1, with host VMs sharing the physical server MAC. In steady state, [IP1, MAC1] route is learnt at [PE1, PE2] and advertised to remote PEs with a sequence number N. Now, VM-IP1 is moved to Server-MAC2. ARP or ND based local learning at [PE3, PE4] would now result in a new [IP1, MAC2] route being learnt. If route [IP1, MAC2] is learnt as a new MAC+IP route and assigned a new sequence number of say 0, mobility procedure for VM-IP1 will not trigger across the overlay network.

A sequence number assignment procedure needs to be defined to unambiguously determine the most recent IP reachability, IP to MAC binding, and MAC reachability for such a MAC sharing scenario.

### 4.3. Host MAC move to new IP

This is a scenario where host move or re-provisioning behind a new gateway location may result in host getting a new IP address assigned, while keeping the same MAC.

# 4.3.1. Problem

Complication with this scenario is that MAC reachability could be carried via a combined MAC+IP route while a MAC only route may not be advertised at all. A single sequence number association with the MAC+IP route again implicitly assumes a fixed mapping between MAC and IP. A MAC move resulting in a new IP association for the host MAC breaks this assumption and results in a new MAC+IP route. If this new MAC+IP route independently assumes a new sequence number, this mobility attribute can no longer be used to determine most recent host MAC reachability.

		++				
	Underlay Network Fabric					
		+	+			
++	++	++	++	++	++	
PE1	PE2	PE3	PE4	PE5	PE6	
++	++	++	++	++	++	
\	/	λ.	/	Υ	/	
\ ESI-1 /		\ ESI	\ ESI-2 /		\ ESI-3 /	
$\land$ /		λ.	\ /		$\land$ /	
+\/+		+\	+\/+		+\/+	
\ /		\	\ /		$  \setminus /  $	
++		+	++		++	
1					1	
Server1		Serv	Server2		Server3	
1						
[VM-IP1-M1, VM-IP2-M2]		-M2] [VM-IP3-	-M3, VM-IP4-	M4] [VM-IP5-	M5, VM-IP6-M6]	

#### Figure 2

As an example, consider a host VM IP1-M1 that is learnt locally at [PE1, PE2] and advertised to remote hosts with a sequence number N. Consider a scenario where this VM with MAC M1 is re-provisioned at server 2, however, as part of this re-provisioning, assigned a different IP address say IP7. [IP7, M1] is learnt as a new route at [PE3, PE4] and advertised to remote PEs with a sequence number of 0. As a result, L3 reachability to IP7 would be established across the overlay, however, MAC mobility procedure for MAC1 will not trigger as a result of this MAC-IP route advertisement. If an optional MAC only route is also advertised, sequence number associated with the MAC only route would trigger MAC mobility as per [RFC7432]. However, in the absence of an additional MAC only route advertisement, a single sequence number advertised with a combined MAC+IP route may not be sufficient to update MAC reachability across the overlay.

A MAC-IP sequence number assignment procedure needs to be defined to unambiguously determine the most recent MAC reachability in such a scenario without a MAC only route being advertised.

Further, PE1/PE2, on learning new reachability for [IP7, M1] via PE3/PE4 MUST probe and delete any local IPs associated with MAC M1, such as [IP1, M1] in the above example.

Arguably, MAC mobility sequence number defined in [RFC7432], could be interpreted to apply only to the MAC part of MAC-IP route, and would hence cover this scenario. It could hence be interpreted as a clarification to [RFC7432] and one of the considerations for a common sequence number assignment procedure across all MAC-IP mobility scenarios detailed in this document.

# 5. EVPN All Active multi-homed ES

+-----+ | Underlay Network Fabric| +----+ +----+ | PE1 | | PE2 | +---+ +----+  $\backslash \backslash$ 11 \\ ESI-1 //  $\backslash \backslash$ /X +\\---//+ | \\ // | +---+ CE1

Figure 3

Consider an EVPN-IRB overlay network shown in Figure 2, with hosts multi-homed to two or more PE devices via an all-active multi-homed ES. MAC and ARP entries learnt on a local ES may also be synchronized across the multi-homing PE devices sharing this ES. This MAC and ARP SYNC enables local switching of intra and inter subnet ECMP traffic flows from remote hosts. In other words, local MAC and ARP entries on a given ES may be learnt via local learning and / or via sync from another PE device sharing the same ES.

For a host that is multi-homed to multiple PE devices via an allactive ES interface, local learning of host MAC and MAC-IP at each PE device is an independent asynchronous event, that is dependent on traffic flow and or ARP / ND response from the host hashing to a directly connected PE on the MC-LAG interface. As a result, sequence number mobility attribute value assigned to a locally learnt MAC or MAC-IP route at each device may not always be the same, depending on transient states on the device at the time of local learning.

As an example, consider a host VM that is deleted from ESI-2 and moved to ESI-1. It is possible for host to be learnt on say, PE1 following deletion of the remote route from [PE3, PE4], while being learnt on PE2 prior to deletion of remote route from [PE3, PE4]. If so, PE1 would process local host route learning as a new route and assign a sequence number of 0, while PE2 would process local host route learning as a remote to local move and assign a sequence number of N+1, N being the existing sequence number assigned at [PE3, PE4].

Inconsistent sequence numbers advertised from multi-homing devices introduces:

\*Ambiguity with respect to how the remote PEs should handle paths with same ESI and different sequence numbers. A remote PE may not program ECMP paths if it receives routes with different sequence numbers from a set of multi-homing PEs sharing the same ESI.

\*Breaks consistent route versioning across the network overlay that is needed for EVPN mobility procedures to work.

As an example, in this inconsistent state, PE2 would drop a remote route received for the same host with sequence number N (as its local sequence number is N+1), while PE1 would install it as the best route (as its local sequence number is 0).

There is need for a mechanism to ensure consistency of sequence numbers advertised from a set of multi-homing devices for EVPN mobility to work reliably.

In order to support mobility for multi-homed hosts using the sequence number mobility attribute, local MAC and MAC-IP routes learnt on a multi-homed ES MUST be advertised with the same sequence number by all PE devices that the ES is multi-homed to. There is need for a mechanism to ensure consistency of sequence numbers assigned across these PEs.

### 6. Design Considerations

To summarize, sequence number assignment scheme and implementation must take following considerations into account:

\*MAC+IP may be learnt on an ES multi-homed to multiple PE devices, hence requires sequence numbers to be synchronized across multihoming PE devices.

\*MAC only RT-2 is optional in an IRB scenario and may not necessarily be advertised in addition to MAC+IP RT-2.

\*Single MAC may be associated with multiple IPs, i.e., multiple host IPs may share a common MAC.

\*Host IP move could result in host moving to a new MAC, resulting in a new IP to MAC association and a new MAC+IP route. \*Host MAC move to a new location could result in host MAC being associated with a different IP address, resulting in a new MAC to IP association and a new MAC+IP route.

\*LOCAL MAC-IP learn via ARP would always accompanied by a LOCAL MAC learn event resulting from the ARP packet. MAC and MAC-IP learning, however, could happen in any order.

\*Use cases discussed earlier that do not maintain a constant 1:1 MAC-IP mapping across moves could potentially be addressed by using separate sequence numbers associated with MAC and IP components of MAC+IP route. Maintaining two separate sequence numbers however adds significant overhead with respect to complexity, debugability, and backward compatibility. Hence, this document addresses these requirements via a single sequence number attribute.

#### 7. Solution Components

This section goes over main components of the EVPN IRB mobility solution proposed in this draft. Later sections will go over exact sequence number assignment procedures resulting from concepts described in this section.

### 7.1. Sequence Number Inheritance

Main idea presented here is to view a LOCAL MAC-IP route as a child of the corresponding LOCAL MAC only route that inherits the sequence number attribute from the parent LOCAL MAC only route:

 $Mx-IPx \dots Mx (seq\# = N)$ 

As a result, both parent MAC and child MAC-IP routes share one common sequence number associated with the parent MAC route. Doing so ensures that a single sequence number attribute carried in a combined MAC+IP route represents sequence number for both a MAC only route as well as a MAC+IP route, and hence makes the MAC only route truly optional. As a result, optional MAC only route with its own sequence number is not required to establish most recent reachability for a MAC in the overlay network. Specifically, this enables a MAC to assume a different IP address on a move, and still be able to establish most recent reachability to the MAC across the overlay network via mobility attribute associated with the MAC+IP route advertisement. As an example, when Mx moves to a new location, it would result in LOCAL Mx being assigned a higher sequence number at its new location as per RFC 7432. If this move results in Mx assuming a different IP address, IPz, LOCAL Mx+IPz route would inherit the new sequence number from Mx.

LOCAL MAC and LOCAL MAC-IP routes would typically be sourced from data plane learning and ARP learning respectively, and could get learnt in control plane in any order. Implementation could either replicate inherited sequence number in each MAC-IP entry OR maintain a single attribute in the parent MAC by creating a forward reference LOCAL MAC object for cases where a LOCAL MAC-IP is learnt before the LOCAL MAC.

Arguably, this inheritance may be assumed from RFC 7432, in which case, the above may be interpreted as a clarification with respect to interpretation of a MAC sequence number in a MAC-IP route.

# 7.2. MAC Sharing

Further, for the shared MAC scenario, this would result in multiple LOCAL MAC-IP siblings inheriting sequence number attribute from a common parent MAC route:

```
Mx-IP1 -----

| | |

Mx-IP2 -----

. |

. +---> Mx (seq# = N)

. |

Mx-IPw -----

| | |

Mx-IPx -----
```

### Figure 4

In such a case, a host-IP move to a different physical server would result in IP moving to a new MAC binding. A new MAC-IP route resulting from this move must now be advertised with a sequence number that is higher than the previous MAC-IP route for this IP, advertised from the prior location. As an example, consider a route MX-IPx that is currently advertised with sequence number N from PE1. IPx moving to a new physical server behind PE2 results in IPx being associated with MAC Mz. A new local Mz-IPx route resulting from this move at PE2 must now be advertised with a sequence number higher than N. This is so that PE devices, including PE1, PE2, and other remote PE devices that are part of the overlay can clearly determine and program the most recent MAC binding and reachability for the IP. PE1, on receiving this new Mz-IPx route with sequence number say, N+1, for symmetric IRB case, would update IPx reachability via PE2 in forwarding, for asymmetric IRB case, would update IPx's ARP binding to Mz. In addition, PE1 would clear and withdraw the stale Mx-IPx route with the lower sequence number.

This also implies that sequence number associated with local MAC Mz and all local MAC-IP children of Mz at PE2 must now be incremented to N+1, and re-advertised across the overlay. While this readvertisement of all local MAC-IP children routes affected by the parent MAC route is an overhead, it avoids the need for two separate sequence number attributes to be maintained and advertised for IP and MAC components of MAC+IP RT-2. Implementation would need to be able to lookup MAC-IP routes for a given IP and update sequence number for it's parent MAC and its MAC-IP children.

# 7.3. Multi-homing Mobility Synchronization

In order to support mobility for multi-homed hosts, local MAC and MAC-IP routes learnt on a shared ES MUST be advertised with the same sequence number by all PE devices that the ES is multi-homed to. This also applies to local MAC only routes. LOCAL MAC and MAC-IP may be learnt natively via data plane and ARP/ND respectively as well as via SYNC from another multi-homing PE to achieve local switching. Local and SYNC route learning can happen in any order. Local MAC-IP routes advertised by all multi-homing PE devices sharing the ES must carry the same sequence number, independent of the order in which they are learnt. This implies:

\*On local or SYNC MAC-IP route learning, sequence number for the local MAC-IP route MUST be compared and updated to the higher value.

\*On local or SYNC MAC route learning, sequence number for the local MAC route MUST be compared and updated to the higher value.

If an update to local MAC-IP sequence number is required as a result of above comparison with SYNC MAC-IP route, it would essentially amount to a sequence number update on the parent local MAC, resulting in inherited sequence number update on the MAC-IP route.

# 8. Requirements for Sequence Number Assignment

Following sections summarize sequence number assignment procedure needed on local and SYNC MAC and MAC-IP route learning events in order to accomplish the above.

#### 8.1. LOCAL MAC-IP learning

A local Mx-IPx learning via ARP or ND should result in computation OR re-computation of parent MAC Mx's sequence number, following which the MAC-IP route Mx-IPx would simply inherit parent MAC's sequence number. Parent MAC Mx Sequence number should be computed as follows:

\*MUST be higher than any existing remote MAC route for Mx, as per RFC 7432.

\*MUST be at least equal to corresponding SYNC MAC sequence number if one is present.

\*If the IP is also associated with a different remote MAC "Mz", MUST be higher than "Mz" sequence number.

Once new sequence number for MAC route Mx is computed as per above, all LOCAL MAC-IPs associated with MAC Mx MUST inherit the updated sequence number.

#### 8.2. LOCAL MAC learning

Local MAC Mx Sequence number should be computed as follows:

\*MUST be higher than any existing remote MAC route for Mx, as per RFC 7432.

\*MUST be at least equal to corresponding SYNC MAC sequence number if one is present.

\*Once new sequence number for MAC route Mx is computed as per above, all LOCAL MAC-IPs associated with MAC Mx MUST inherit the updated sequence number.

Note that the local MAC sequence number might already be present if there was a local MAC-IP learnt prior to the local MAC, in which case the above may not result in any change in local MAC's sequence number.

### 8.3. Remote MAC OR MAC-IP Update

On receiving a remote MAC OR MAC-IP route update associated with a MAC Mx with a sequence number that is

\*either higher than the sequence number assigned to a LOCAL route
for MAC Mx,

\*or equal to the sequence number assigned to a LOCAL route for MAC Mx, but the remote route is selected as best because of lower VTEP IP as per [RFC 7432],

following handling is required on the receiving PE:

\*PE MUST trigger probe and deletion procedure for all LOCAL IPs associated with MAC Mx.

\*PE MUST trigger deletion procedure for LOCAL MAC route for Mx.

### 8.4. REMOTE (SYNC) MAC update

Corresponding local MAC Mx (if present) sequence number should be re- computed as follows:

\*If the current sequence number is less than the received SYNC MAC sequence number, it MUST be increased to be equal to received SYNC MAC sequence number.

\*If a LOCAL MAC sequence number is updated as a result of the above, all LOCAL MAC-IPs associated with MAC Mx MUST inherit the updated sequence number.

# 8.5. REMOTE (SYNC) MAC-IP update

If this is a SYNCed MAC-IP on a local ES, it would also result in a derived SYNC MAC Mx route entry, as MAC only RT-2 advertisement is optional. Corresponding local MAC Mx (if present) sequence number should be re-computed as follows:

\*If the current sequence number is less than the received SYNC MAC sequence number, it MUST be increased to be equal to received SYNC MAC sequence number.

\*If a LOCAL MAC sequence number is updated as a result of the above, all LOCAL MAC-IPs associated with MAC Mx MUST inherit the updated sequence number.

### 8.6. Inter-op

In general, if all PE nodes in the overlay network follow the above sequence number assignment procedure, and the PE is advertising both MAC+IP and MAC routes, sequence number advertised with the MAC and MAC+IP routes with the same MAC would always be the same. However, an inter-op scenario with a different implementation could arise, where a PE implementation non-compliant with this document or with RFC 7432 assigns and advertises independent sequence numbers to MAC and MAC+IP routes. To handle this case, if different sequence numbers are received for remote MAC+IP and corresponding remote MAC routes from a remote PE, sequence number associated with the remote MAC route should be computed as:

\*Highest of the all received sequence numbers with remote MAC+IP and MAC routes with the same MAC.

\*MAC sequence number would be re-computed on a MAC or MAC+IP route withdraw as per above.

A MAC and / or IP move to the local PE would now result in the MAC (and hence all MAC-IP) sequence numbers incremented from the above computed remote MAC sequence number.

If MAC only routes are not advertised at all, and different sequence numbers are received with multiple MAC+IP routes for a given MAC, sequence number associated with the derived remote MAC route should still be computed as the highest of the all received MAC+IP sequence numbers with the same MAC.

### 8.7. MAC Sharing Race Condition

In a MAC sharing use case described in section 6.2, a race condition is possible with simultaneous host moves between a pair of PEs. As an example, consider PE1 with local host IPs I1 and I2 sharing MAC M1, and PE2 with local host IPs I3 and I4 sharing MAC M2. A simultaneous move of I1 from PE1 to PE2 and of I3 from PE2 to PE1, such that I3 is learnt on PE1 before I1's local entry has been probed out on PE1 and/or I1 is learnt on PE2 before I3's local entry has been probed out on PE2 may trigger a race condition. This race condition together with MAC sequence number assignment rules defined in section 7.1 can cause new mac-ip routes [I1, M2] and [I3, M1] to bounce a couple of times with an incremented sequence number until stale entries [I1, M1] and [I3, M2] have been probed out from PE1 and PE2 respectively. An implementation MUST ensure proper probing procedures to remove stale ARP, ND, and local MAC entries, following a move, on learning remote routes as defined in section 7.3 (and as per [EVPN-IRB]) to minimize exposure to this race condition.

### 8.8. Mobility Convergence

This sections is to be treated as optional and details ARP and ND probing procedures that MAY be implemented to achieve faster host re- learning and convergence on mobility events.

\*Following a host move from PE1 to PE2, the host's MAC is discovered at PE2 as a local MAC via a data frames received from the host. If PE2 has a prior REMOTE MAC-IP host route for this MAC from PE1, an ARP/ND probe MAY be triggered at PE2 to learn the MAC-IP as a local adjacency and trigger EVPN RT-2 advertisement for this MAC-IP across the overlay with new reachability via PE2. This results in a reliable "event based" host IP learning triggered by a "MAC learning event" across the overlay, and hence faster convergence of overlay routed flows to the host.

\*Following a host move from PE1 to PE2, once PE1 receives a MAC or MAC-IP route from PE2 with a higher sequence number, an ARP/ND probe MAY be triggered at PE1 to clear the stale local MAC-IP neighbor adjacency OR re-learn the local MAC-IP in case the host has moved back or is duplicate.

\*Following a local MAC age-out, if there is a local IP adjacency with this MAC, an ARP/ND probe MAY be triggered for this IP to either re-learn the local MAC and maintain local 13 and 12 reachability to this host OR to clear the ARP/ND entry in case the host is indeed no longer local. Note that this accomplishes clearing of stale ARP entries, triggered by a MAC age-out event even when the ARP refresh timer was longer than the MAC age-out timer. Clearing of stale IP neighbor entries in-turn facilitates traffic convergence in the event that the host was silent and not discovered at its new location. Once stale neighbor entry for the host is cleared, routed traffic flow destined for the host can re-trigger ARP/ND discovery for this host at the new location.

### 8.8.1. Generalized Probing Logic

Above probing logic may be generalized as probing for an IP neighbor anytime a resolving parent MAC route is "inconsistent" with the MAC-IP neighbor route, where being inconsistent is defined as being not present OR conflicting in terms of the route source being local OR remote. MAC-IP to MAC parent relationship described earlier in this document in section 6.1 MAY be used to achieve this logic.

### 9. Routed Overlay

An additional use case is possible, such that traffic to an end host in the overlay is always IP routed. In a purely routed overlay such as this:

\*A host MAC is never advertised in EVPN overlay control plane.

\*Host /32 or /128 IP reachability is distributed across the overlay via EVPN route type 5 (RT-5) along with a zero or non-zero ESI.

\*An overlay IP subnet may still be stretched across the underlay fabric, however, intra-subnet traffic across the stretched overlay is never bridged.

\*Both inter-subnet and intra-subnet traffic, in the overlay is IP routed at the EVPN PE.

Please refer to [RFC 7814] for more details.

Host mobility within the stretched subnet would still need to be supported for this use. In the absence of any host MAC routes, sequence number mobility EXT-COMM specified in [RFC7432], section 7.7 may be associated with a /32 OR /128 host IP prefix advertised via EVPN route type 5. MAC mobility procedures defined in RFC 7432 can now be applied as is to host IP prefixes:

\*On LOCAL learning of a host IP, on a new ESI, host IP MUST be advertised with a sequence number attribute that is higher than what is currently advertised with the old ESI.

\*On receiving a host IP route advertisement with a higher sequence number, a PE MUST trigger ARP/ND probe and deletion procedure on any LOCAL route for that IP with a lower sequence number. A PE would essentially move the forwarding entry to point to the remote route with a higher sequence number and send an ARP/ND PROBE for the local IP route. If the IP has indeed moved, PROBE would timeout and the local IP host route would be deleted.

Note that there is still only one sequence number associated with a host route at any time. For earlier use cases where a host MAC is advertised along with the host IP, a sequence number is only associated with a MAC. Only if the MAC is not advertised at all, as in this use case, is a sequence number associated with a host IP.

Note that this mobility procedure would not apply to "anycast IPv6" hosts advertised via NA messages with 0-bit=0. Please refer to [EVPN-PROXY-ARP].

# **10.** Duplicate Host Detection

Duplicate host detection scenarios across EVPN IRB can be classified as follows:

\*Scenario A: where two hosts have the same MAC (host IPs may or may not be duplicate).

\*Scenario B: where two hosts have the same IP but different MACs.

\*Scenario C: where two hosts have the same IP and host MAC is not advertised at all.

Duplicate detection procedures for scenario B and C would not apply to "anycast IPv6" hosts advertised via NA messages with 0-bit=0. Please refer to [EVPN-PROXY-ARP].

### 10.1. Scenario A

For all use cases where duplicate hosts have the same MAC, MAC is detected as duplicate via duplicate MAC detection procedure described in RFC 7432. Corresponding MAC-IP routes with the same MAC do not require duplicate detection and MUST simply inherit the DUPLICATE property from the corresponding MAC route. In other words, if a MAC route is in DUPLICATE state, all corresponding MAC-IP routes MUST also be treated as DUPLICATE. Duplicate detection procedure need only be applied to MAC routes.

### 10.2. Scenario B

Due to misconfiguration, a situation may arise where hosts with different MACs are configured with the same IP. This scenario would not be detected by existing duplicate MAC detection procedure and would result in incorrect forwarding of routed traffic destined to this IP.

Such a situation, on LOCAL MAC-IP learning, would be detected as a move scenario via the following local MAC sequence number computation procedure described earlier in section 6.1:

\*If the IP is also associated with a different remote MAC "Mz", MUST be higher than "Mz" sequence number.

Such a move that results in sequence number increment on local MAC because of a remote MAC-IP route associated with a different MAC MUST be counted as an "IP move" against the "IP" independent of MAC. Duplicate detection procedure described in RFC 7432 can now be applied to an "IP" entity independent of MAC. Once an IP is detected as DUPLICATE, corresponding MAC-IP route should be treated as DUPLICATE. Associated MAC routes and any other MAC-IP routes associated with this MAC should not be affected.

### 10.2.1. Duplicate IP Detection Procedure for Scenario B

Duplicate IP detection procedure for such a scenario is specified in [EVPN-PROXY-ARP]. What counts as an "IP move" in this scenario is further clarified as follows:

\*On learning a LOCAL MAC-IP route Mx-IPx, check if there is an existing REMOTE OR LOCAL route for IPx with a different MAC association, say, Mz-IPx. If so, count this as an "IP move" count for IPx, independent of the MAC.

\*On learning a REMOTE MAC-IP route Mz-IPx, check if there is an existing LOCAL route for IPx with a different MAC association, say, Mx-IPx. If so, count this as an "IP move" count for IPx, independent of the MAC. A MAC-IP route SHOULD be treated as DUPLICATE if either of the following two conditions are met:

\*Corresponding MAC route is marked as DUPLICATE via existing duplicate detection procedure.

\*Corresponding IP is marked as DUPLICATE via extended procedure described above.

### 10.3. Scenario C

For a purely routed overlay scenario described in section 8, where only a host IP is advertised via EVPN RT-5, together with a sequence number mobility attribute, duplicate MAC detection procedures specified in RFC 7432 can be intuitively applied to IP only host routes for the purpose of duplicate IP detection.

\*On learning a LOCAL host IP route IPx, check if there is an existing REMOTE OR LOCAL route for IPx with a different ESI association. If so, count this as an "IP move" count for IPx.

\*On learning a REMOTE host IP route IPx, check if there is an existing LOCAL route for IPx with a different ESI association. If so, count this as an "IP move" count for IPx.

\*With configurable parameters "N" and "M", If "N" IP moves are detected within "M" seconds for IPx, treat IPx as DUPLICATE.

### 10.4. Duplicate Host Recovery

Once a MAC or IP is marked as DUPLICATE and FROZEN, corrective action must be taken to un-provision one of the duplicate MAC or IP. Un- provisioning a duplicate MAC or IP in this context refers to a corrective action taken on the host side. Once one of the duplicate MAC or IP is un-provisioned, normal operation would not resume until the duplicate MAC or IP ages out, following this correction, unless additional action is taken to speed up recovery.

This section lists possible additional corrective actions that could be taken to achieve faster recovery to normal operation.

# 10.4.1. Route Un-freezing Configuration

Unfreezing the DUPLICATE OR FROZEN MAC or IP via a CLI can be leveraged to recover from DUPLICATE and FROZEN state following corrective un-provisioning of the duplicate MAC or IP.

Unfreezing the frozen MAC or IP via a CLI at a PE should result in that MAC OR IP being advertised with a sequence number that is

higher than the sequence number advertised from the other location of that MAC or IP.

Two possible corrective un-provisioning scenarios exist:

\*Scenario A: A duplicate MAC or IP may have been un-provisioned at the location where it was NOT marked as DUPLICATE and FROZEN.

\*Scenario B: A duplicate MAC or IP may have been un-provisioned at the location where it was marked as DUPLICATE and FROZEN.

Unfreezing the DUPLICATE and FROZEN MAC or IP, following the above corrective un-provisioning scenarios would result in recovery to steady state as follows:

\*Scenario A: If the duplicate MAC or IP was un-provisioned at the location where it was NOT marked as DUPLICATE, unfreezing the route at the FROZEN location will result in the route being advertised with a higher sequence number. This would in-turn result in automatic clearing of local route at the PE location, where the host was un-provisioned via ARP/ND PROBE and DELETE procedure specified earlier in section 8 and in [RFC 7432].

\*Scenario B: If the duplicate host is un-provisioned at the location where it was marked as DUPLICATE, unfreezing the route will trigger an advertisement with a higher sequence number to the other location. This would in-turn trigger re-learning of local route at the remote location, resulting in another advertisement with a higher sequence number from the remote location. Route at the local location would now be cleared on receiving this remote route advertisement, following the ARP/ND PROBE.

Note that the probes referred to in the above scenarios are event driven probes resulting from receiving a route with a higher sequence number. Periodic probes resulting from refresh timers may also occur in addition as completely independent probes.

#### 10.4.2. Route Clearing Configuration

In addition to the above, route clearing CLIs may also be leveraged to clear the local MAC or IP route, to be executed AFTER the duplicate host is un-provisioned:

\*clear mac CLI: A clear MAC CLI can be leveraged to clear a DUPLICATE MAC route, to recover from a duplicate MAC scenario.

\*clear ARP/ND: A clear ARP/ND CLI may be leveraged to clear a DUPLICATE IP route to recover from a duplicate IP scenario.

Note that the route unfreeze CLI may still need to be run if the route was un-provisioned and cleared from the NON-DUPLICATE / NON-FROZEN location. Given that unfreezing of the route via the un-freeze CLI would any ways result in auto-clearing of the route from the "un- provisioned" location, as explained in the prior section, need for a route clearing CLI for recovery from DUPLICATE / FROZEN state is truly optional.

### 11. Security Considerations

This document raises no new security issues for EVPN.

### **12. IANA Considerations**

None.

### 13. Acknowledgements

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# Authors' Addresses

Neeraj Malhotra (editor) Cisco Systems 170 W. Tasman Drive San Jose, CA 95134 United States of America

Email: <u>nmalhotr@cisco.com</u>

Ali Sajassi Cisco Systems 170 W. Tasman Drive San Jose, CA 95134 United States of America

Email: <u>sajassi@cisco.com</u>

Aparna Pattekar Cisco Systems 170 W. Tasman Drive San Jose, CA 95134 United States of America

Email: apjoshi@cisco.com

Jorge Rabadan Nokia 777 E. Middlefield Road Mountain View, CA 94043 United States of America

Email: jorge.rabadan@nokia.com

Avinash Lingala ATT 200 S. Laurel Avenue Middletown, CA 07748 United States of America

Email: <u>ar977m@att.com</u>

John Drake Juniper Networks

# Email: jdrake@juniper.net