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**Requirements for Generalized MPLS (GMPLS) Routing
for Automatically Switched Optical Network (ASON)**

[draft-ietf-ccamp-gmpls-ason-routing-reqts-03.txt](#)

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Abstract

The Generalized MPLS (GMPLS) suite of protocols has been defined to

control different switching technologies as well as different applications. These include support for requesting TDM connections including SONET/SDH and Optical Transport Networks (OTNs).

This document concentrates on the routing requirements on the GMPLS suite of protocols to support the capabilities and functionalities for an Automatically Switched Optical Network (ASON) as defined by ITU-T.

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1. Contributors

This document is the result of the CCAMP Working Group ASON Routing Requirements design team joint effort.

2. Conventions used in this document

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC 2119](#) [[RFC2119](#)].

3. Introduction

The GMPLS suite of protocols provides among other capabilities support for controlling different switching technologies. These include support for requesting TDM connections utilizing SONET/SDH (see ANSI T1.105/ITU-T G.707) as well as Optical Transport Networks (OTN, see ITU-T G.709). However, there are certain capabilities that are needed to support the ITU-T G.8080 control plane architecture for Automatically Switched Optical Network (ASON). Therefore, it is

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desirable to understand the corresponding requirements for the GMPLS protocol suite. The ASON control plane architecture is defined in [[G.8080](#)], ASON routing requirements are identified in [[G.7715](#)], and in [[G.7715.1](#)] for ASON link state protocols. These Recommendations apply to all G.805 layer networks (e.g. SDH and OTN), and provide protocol neutral functional requirements and architecture.

This document focuses on the routing requirements for the GMPLS suite of protocols to support the capabilities and functionality of ASON control planes. This document summarizes the ASON requirements using ASON terminology. This document does not address GMPLS applicability or GMPLS capabilities. Any protocol (in particular, routing) applicability, design or suggested extensions is strictly outside the scope of this document. ASON (Routing) terminology sections are provided in Appendix 1 and 2.

The ASON routing architecture is based on the following assumptions:

- A network is subdivided based on operator decision and criteria (e.g. geography, administration, and/or technology), the network subdivisions are defined in ASON as Routing Areas (RAs).
- The routing architecture and protocols applied after the network is subdivided is an operator's choice. A multi-level hierarchy of RAs, as defined in ITU-T [G.7715] and [G.7715.1], provides for a hierarchical relationship of RAs based on containment, i.e. child RAs are always contained within a parent RA. The hierarchical containment relationship of RAs provides for routing information abstraction, thereby enabling scalable routing information representation. The maximum number of hierarchical RA levels to be supported is NOT specified (outside the scope).
- Within an ASON RA and for each level of the routing hierarchy, multiple routing paradigms (hierarchical, step- by-step, source-based), centralized or distributed path computation, and multiple different routing protocols MAY be supported. The architecture does NOT assume a one-to-one correspondence of a routing protocol and a RA level and allows the routing protocol(s) used within different RAs (including child and parent RAs) to be different. The realization of the routing paradigm(s) to support the hierarchical levels of RAs is NOT specified.
- The routing adjacency topology (i.e. the associated Protocol Controller (PC) connectivity) and transport topology is NOT assumed to be congruent.
- The requirements support architectural evolution, e.g. a change in the number of RA levels, as well as aggregation and segmentation of RAs.

The description of the ASON routing architecture provides for a conceptual reference architecture, with definition of functional components and common information elements to enable end-to-end routing in the case of protocol heterogeneity and facilitate management of ASON networks. This description is only conceptual: no physical partitioning of these functions is implied.

4. ASON Routing Architecture and Requirements

The fundamental architectural concept is the RA and it's related functional components (see Appendix 2 on terminology). The routing services offered by a RA are provided by a Routing Performer (RP). An RP is responsible for a single RA, and it MAY be functionally

realized using distributed Routing Controllers (RC). The RC, itself, MAY be implemented as a cluster of distributed entities (ASON refers to the cluster as a Routing Control Domain (RCD)). The RC components for a RA receive routing topology information from their associated Link Resource Manager(s) (LRMs) and store this information in the Routing Information Database (RDB). The RDB is replicated at each RC bounded to the same Routing Area (RA), and MAY contain information about multiple transport plane network layers. Whenever the routing topology changes, the LRM informs the corresponding RC, which in turn updates its associated RDB. In order to assure RDB synchronization, the RCs co-operate and exchange routing information. Path computation functions MAY exist in each RC, MAY exist on selected RCs within the same RA, or MAY be centralized for the RA.

In this context, communication between RCs within the same RA is realized using a particular routing protocol (or multiple protocols). In ASON, the communication component is represented by the protocol controller (PC) component(s) and the protocol messages are conveyed over the ASON control plane's Signaling Control Network (SCN). The PC MAY convey information for one or more transport network layers (refer to [Section 4.2](#) Note). The RC is protocol independent and RC communications MAY be realized by multiple, different PCs within a RA.

The ASON routing architecture defines a multi-level routing hierarchy of RAs based on a containment model to support routing information abstraction. [\[G.7715.1\]](#) defines the ASON hierarchical link state routing protocol requirements for communication of routing information within an RA (one level) to support hierarchical routing information dissemination (including summarized routing information for other levels). The Communication between any of the other functional component(s) (e.g. SCN, LRM, and between RCDs (RC-RC communication between RAs)), is outside the scope of [\[G.7715.1\]](#) protocol requirements and, thus, is also outside the scope of this document.

ASON Routing components are identified by identifiers that are drawn from different name spaces (see [\[G.7715.1\]](#)). These are control plane identifiers for transport resources, components, and SCN addresses. The formats of those identifiers in a routing protocol realization SHALL be implementation specific and outside the scope of this document.

The failure of a RC, or the failure of communications between RCs, and the subsequent recover from the failure condition MUST NOT

disrupt calls in progress and their associated connections. Calls being set up MAY fail to complete, and the call setup service MAY be unavailable during recovery actions.

4.1 Multiple Hierarchical Levels of ASON Routing Areas (RAs)

[G.8080] introduces the concept of Routing Area (RA) in reference to a network subdivision. RAs provide for routing information abstraction. Except for the single RA case, RAs are hierarchically contained: a higher level (parent) RA contains lower level (child) RAs that in turn MAY also contain RAs, etc. Thus, RAs contain RAs that recursively define successive hierarchical RA levels.

However, the RA containment relationship describes only an architectural hierarchical organization of RAs. It does not restrict a specific routing protocol's realization (e.g. OSPF multi-areas, path computation, etc.). Moreover, the realization of the routing paradigm to support a hierarchical organization of RAs and the number of hierarchical RA levels to be supported is routing protocol specific and outside the scope of this document.

In a multi-level hierarchy of RAs, it is necessary to distinguish among RCs for the different levels of the RA hierarchy. Before any pair of RCs establishes communication, they MUST verify they are bounded to the same parent RA (see [Section 4.2](#)). A RA identifier (RA ID) is required to provide the scope within which the RCs can communicate. To distinguish between RCs bounded to the same RA, an RC identifier (RC ID) is required; the RC ID MUST be unique within its containing RA.

A RA represents a partition of the data plane and its identifier (i.e. RA ID) is used within the control plane as a reference to the data plane partition. Each RA SHALL be uniquely identifiable within a carrier's network. RA IDs MAY be associated with a transport plane name space whereas RC IDs are associated with a control plane name space.

4.2 Hierarchical Routing Information Dissemination

Routing information can be exchanged between RCs bounded to adjacent levels of the RA hierarchy i.e. Level N+1 and N, where Level N represents the RAs contained by Level N+1. The links connecting RAs MAY be viewed as external links (inter-RA links), and the links representing connectivity within a RA MAY be viewed as internal links (intra-RA links). The external links to a RA at one level of the hierarchy may be internal links in the parent RA. Intra-RA links of a child RA MAY be hidden from the parent RA's view.

The physical location of RCs for adjacent RA levels, their relationship and their communication protocol(s) are outside the scope of this document. No assumption is made regarding how RCs communicate between adjacent RA levels. If routing information is

exchanged between a RC, its parent, and its child RCs, it SHOULD include reachability and MAY include (upon policy decision) node and link topology. Only the RCs of the parent RA communicate, RCs of one child's RA never communicate with the RCs of other child RAs. There SHOULD not be any dependencies on the different routing protocols used within a RA or in different RAs.

Multiple RCs bounded to the same RA MAY transform (filter, summarize, etc.) and then forward information to RCs at different levels. However in this case the resulting information at the receiving level must be self-consistent; this MAY be achieved using a number of mechanisms.

Note: there is no implied relationship between multi-layer transport networks and multi-level routing. Implementations may support a hierarchical routing topology (multi-level) with a single routing protocol instance for multiple transport switching layers or a hierarchical routing topology for one transport switching layer.

1. Type of Information Exchanged

The type of information flowing upward (i.e. Level N to Level N+1) and the information flowing downward (i.e. Level N+1 to Level N) are used for similar purposes, namely, the exchange of reachability information and summarized topology information to allow routing across multiple RAs. The summarization of topology

information may impact the accuracy of routing and MAY require additional path calculation.

The following information exchange are expected:

- Level N+1 visibility to Level N reachability and topology (or upward information communication) allowing RC(s) at Level N+1 to determine the reachable endpoints from Level N.
- Level N visibility to Level N+1 reachability and topology (or downward information communication) allowing RC(s) bounded to a RA at Level N to develop paths to reachable endpoints outside of the RA.

2. Interactions between Upward and Downward Communication

When both upward and downward information exchanges contain endpoint reachability information, a feedback loop could potentially be created. Consequently, the routing protocol MUST include a method to:

- prevent information propagated from a Level N+1 RA's RC into the Level N RA's RC to be re-introduced into the Level N+1 RA's RC, and
- prevent information propagated from a Level N-1 RA's RC into the Level N RA's RC to be re-introduced into the Level N-1 RA's

RC.

The routing protocol SHALL differentiate the routing information originated at a given level RA from derived routing information (received from external RAs), even when this information is forwarded by another RC at the same level. This is a necessary condition to be fulfilled by routing protocols to be loop free.

3. Method of Communication

Two approaches exist for communication between Level N and N+1.

- The first approach places an instance of a Level N routing function and an instance of a Level N+1 routing function in the same system. The communications interface is within a single system and is thus not an open interface subject to standardization.
- The second approach places the Level N routing function on a separate system from the Level N+1 routing function. In this case, a communication interface must be used between the systems containing the routing functions for different levels. This communication interface and mechanisms are outside the scope of this document.

4.3 Configuration

4.3.1 Configuring the Multi-Level Hierarchy

The RC MUST support static (i.e. operator assisted) and MAY support automated configuration of the information describing its relationship to parent and its child within the hierarchical structure (including RA ID and RC ID). When applied recursively, the whole hierarchy is thus configured.

4.3.2 Configuring RC Adjacencies

The RC MUST support static (i.e. operator assisted) and MAY support automated configuration of the information describing its associated PC adjacencies to other RCs bounded to the same parent RA. The routing protocol SHOULD support all the types of RC adjacencies described in Section 9 of [[G.7715](#)]. The latter includes congruent topology (with distributed RC) and hubbed topology (e.g. note that the latter does not automatically imply designated RC).

4.4 Evolution

The containment relationships of RAs MAY change, motivated by events such as mergers, acquisitions, and divestitures.

The routing protocol SHOULD be capable of supporting architectural evolution in terms of number of hierarchical levels of RAs, as well

as aggregation and segmentation of RAs. RA IDs uniqueness within an administrative domain MAY facilitate these operations. The routing protocol is not expected to automatically initiate and/or execute these operations. Reconfiguration of the RA hierarchy MAY not disrupt calls in progress, though calls being set up may fail to complete, and the call setup service may be unavailable during reconfiguration actions.

4.5 Routing Attributes

Routing for transport networks is performed on a per layer basis, where the routing paradigms MAY differ among layers and within a layer. Not all equipment support the same set of transport layers or the same degree of connection flexibility at any given layer. A server layer trail may support various clients, involving different adaptation functions. Additionally, equipment may support variable adaptation functionality, whereby a single server layer trail dynamically supports different multiplexing structures. As a result, routing information MAY include layer specific, layer independent, and client/server adaptation information.

4.5.1 Taxonomy of Routing Attributes

Attributes can be organized according to the following categories:

- Node related or link related
- Provisioned, negotiated or automatically configured
- Inherited or layer specific (client layers can inherit some attributes from the server layer while other attributes like Link Capacity are specified by layer).

(Component) link attributes MAY be statically or automatically configured for each transport network layer. This may lead to unnecessary repetition. Hence, the inheritance property of attributes MAY also be used to optimize the configuration process.

ASON uses the term, SNP, for the control plane representation of a transport plane resource. The control plane representation and transport plane topology is NOT assumed to be congruent, the control plane representation SHALL not be restricted by the physical topology. The relational grouping of SNPs for routing is termed a SNPP. The routing function understands topology in terms of SNPP links. Grouping MAY be based on different link attributes (e.g., SRLG information, link weight, etc).

Two RAs may be linked by one or more SNPP links. Multiple SNPP links MAY be required when component links are not equivalent for routing purposes with respect to the RAs they are attached to, or to the containing RA, or when smaller groupings are required.

4.5.2 Commonly Advertised Information

Advertisements MAY contain the following common set of information regardless of whether they are link or node related:

- RA ID of which the advertisement is bounded
- RC ID of the entity generating the advertisement
- Information to uniquely identify advertisements
- Information to determine whether an advertisement has been updated
- Information to indicate when an advertisement has been derived from a different level RA.

4.5.3 Node Attributes

All nodes belong to a RA, hence, the RA ID can be considered an attribute of all nodes. Given that no distinction is made between abstract nodes and those that cannot be decomposed any further, the same attributes MAY be used for their advertisement. In the following tables, Capability refers to level of support required in the realization of a link state routing protocol, whereas Usage refers to degree of operational and implementation flexibility.

The following Node Attributes are defined:

| Attribute | Capability | Usage |
|--------------|------------|----------|
| ----- | ----- | ----- |
| Node ID | REQUIRED | REQUIRED |
| Reachability | REQUIRED | OPTIONAL |

Table 1. Node Attributes

Reachability information describes the set of endpoints that are reachable by the associated node. It MAY be advertised as a set of associated external (e.g. UNI) address/address prefixes or a set of associated SNPP link IDs/SNPP ID prefixes, the selection of which MUST be consistent within the applicable scope. These are control plane identifiers, the formats of these identifiers in a protocol realization is implementation specific and outside the scope of this document.

Note: no distinction is made between nodes that may have further internal details (i.e., abstract nodes) and those that cannot be decomposed any further. Hence the attributes of a node are not be considered only as single switch attributes but MAY apply to a node at a higher level of the hierarchy that represents a sub-network.

[4.5.4](#) Link Attributes

The following Link Attributes are defined:

| Link Attribute | Capability | Usage |
|--------------------|------------|----------|
| ----- | ----- | ----- |
| Local SNPP link ID | REQUIRED | REQUIRED |

| | | |
|--------------------------------|-------------|----------|
| Remote SNPP link ID | REQUIRED | REQUIRED |
| Layer Specific Characteristics | see Table 3 | |

Table 2. Link Attributes

The SNPP link ID name MUST be sufficient to uniquely identify the corresponding transport plane resource taking into account separation of data and control planes (see [Section 4.5.1](#), the

control plane representation and transport plane topology is not assumed to be congruent). The SNPP link ID format is routing protocol specific.

Note: when the remote end of a SNPP link is located outside of the RA, the remote SNPP link ID is OPTIONAL.

The following link characteristic attributes are defined:

- Signal Type: This identifies the characteristic information of the layer network.
- Link Weight: The metric indicating the relative desirability of a particular link over another e.g. during path computation.
- Resource Class: This corresponds to the set of administrative groups assigned by the operator to this link. A link MAY belong to zero, one or more administrative groups.
- Connection Types: This attribute identifies whether the local SNP represents a TCP, CP, or can be flexibly configured as a TCP.
- Link Capacity: This provides the sum of the available and potential bandwidth capacity for a particular network transport layer. Other capacity measures MAY be further considered.
- Link Availability: This represents the survivability capability such as the protection type associated with the link.
- Diversity Support: This represents diversity information such as the SRLG information associated with the link.
- Local Adaptation Support: This indicates the set of client layer adaptations supported by the TCP associated with the Local SNPP. This is only applicable when the local SNP represents a TCP or can be flexibly configured as a TCP.

| Link Characteristics | Capability | Usage |
|----------------------|------------|----------|
| ----- | ----- | ----- |
| Signal Type | REQUIRED | OPTIONAL |
| Link Weight | REQUIRED | OPTIONAL |
| Resource Class | REQUIRED | OPTIONAL |

| | | |
|------------------------|----------|----------|
| Local Connection Types | REQUIRED | OPTIONAL |
| Link Capacity | REQUIRED | OPTIONAL |

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| | | |
|--------------------------|----------|----------|
| Link Availability | OPTIONAL | OPTIONAL |
| Diversity Support | OPTIONAL | OPTIONAL |
| Local Adaptation support | OPTIONAL | OPTIONAL |

Table 3. Link Characteristics

Note: separate advertisements of layer specific attributes MAY be chosen. However this may lead to unnecessary duplication. This can be avoided using the inheritance property, so that the attributes derivable from the local adaptation information do not need to be advertised. Thus, an optimization MAY be used when several layers are present by indicating when an attribute is inheritable from a server layer.

5. Security Considerations

ASON routing protocol MUST deliver the operational security objectives where required. These objectives do not necessarily imply requirements on the routing protocol itself, and MAY be met by other established means.

6. Conclusions

The description of the ASON routing architecture and components is provided in terms of routing functionality. This description is only conceptual: no physical partitioning of these functions is implied.

In summary, the ASON routing architecture assumes:

- A network is subdivided into ASON RAs, which MAY support multiple routing protocols, no one-to-one relationship SHALL be assumed
- Routing Controllers (RC) provide for the exchange of routing information (primitives) for the RA. The RC is protocol independent and MAY be realized by multiple, different protocol controllers within a RA. The routing information exchanged between RCs SHALL be subject to policy constraints imposed at reference

points (External- and Internal-NNI)

- A multi-level RA hierarchy based on containment, only the RCs of the parent RA communicate. RCs of child RAs never communicate with the RCs of other child RAs. There SHOULD not be any dependencies on the different routing protocols used within a child RA and that of its parent. The routing information exchanged within the parent RA SHALL be independent of both the routing protocol operating within a child RA, and any control distribution choice(s), e.g. centralized, fully distributed.
- For a RA, the set of RCs is referred to as an ASON routing (control) domain. The routing information exchanged between routing domains (inter-RA, i.e. inter-domain) SHALL be independent of both the intra-domain routing protocol(s), and the intra-domain control distribution choice(s), e.g. centralized, fully distributed. RCs bounded to different RA levels MAY be co-located within the same physical element or physically distributed.
- The routing adjacency topology (i.e. the associated PC

connectivity topology) and the transport network topology SHALL NOT be assumed to be congruent

- The routing topology SHALL support multiple links between nodes and RAs

In summary, the following functionality is expected from GMPLS routing to instantiate the ASON hierarchical routing architecture realization (see [G.7715] and [G.7715.1]):

- RAs SHALL be uniquely identifiable within a carrier's network, each having a unique RA ID within the carrier's network.
- Within a RA (one level), the routing protocol SHALL support dissemination of hierarchical routing information (including summarized routing information for other levels) in support of an architecture of multiple hierarchical levels of RAs; the number of hierarchical RA levels to be supported by a routing protocol is implementation specific.
- The routing protocol SHALL support routing information based on a common set of information elements as defined in [G.7715] and [G.7715.1], divided between attributes pertaining to links and abstract nodes (each representing either a sub-network or simply a node). [G.7715] recognizes that the manner in which the routing information is represented and exchanged will vary with the routing protocol used.
- The routing protocol SHALL converge such that the distributed RDBs become synchronized after a period of time.

To support hierarchical routing information dissemination within an RA, the routing protocol MUST deliver:

- processing of routing information exchanged between adjacent levels of the hierarchy (i.e. Level N+1 and N) including reachability and upon policy decision summarized topology information
- when multiple RCs bound to a RA transform (filter, summarize, etc.) and then forward information to RC(s) at different levels that the resulting information at the receiving level is self-consistent
- a mechanism to prevent re-introduction of information propagated into the Level N RA's RC back to the adjacent level RA's RC from which this information has been initially received.

In order to support operator assisted changes in the containment relationships of RAs, the routing protocol SHALL support evolution in terms of number of hierarchical levels of RAs. Example: support of non-disruptive operations such as adding and removing RAs at the top/bottom of the hierarchy, adding or removing a hierarchical level of RAs in or from the middle of the hierarchy, as well as aggregation and segmentation of RAs. The number of hierarchical levels to be supported is routing protocol specific, and reflects a containment relationship e.g. a RA insertion involves supporting a different routing protocol domain in a portion of the network.

Reachability information (see [Section 4.5.3](#)) of the set of endpoints reachable by a node may be advertised either as a set of UNI Transport Resource addresses/ address prefixes, or a set of associated SNPP link IDs/SNPP link ID prefixes, assigned and selected consistently in their applicability scope. The formats of the control plane identifiers in a protocol realization are implementation specific. Use of a routing protocol within a RA should not restrict the choice of routing protocols for use in other RAs (child or parent).

As ASON does not restrict the control plane architecture choice used, either a co-located architecture or a physically separated architecture may be used. A collection of links and nodes such as a sub-network or RA MUST be able to represent itself to the wider

network as a single logical entity with only its external links visible to the topology database.

7. Acknowledgements

The authors would like to thank Kireeti Kompella for having initiated the proposal of an ASON Routing Requirement Design Team.

8. Intellectual Property Considerations

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8.1 IPR Disclosure Acknowledgement

By submitting this Internet-Draft, I certify that any applicable patent or other IPR claims of which I am aware have been disclosed, and any of which I become aware will be disclosed, in accordance with [RFC 3668](#).

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Appendix 1: ASON Terminology

This document makes use of the following terms:

Administrative domain: (Recommendation G.805 For the purposes of [G7715.1] an administrative domain represents the extent of resources which belong to a single player such as a network operator, a service provider, or an end-user. Administrative domains of different players do not overlap amongst themselves.

Control plane: performs the call control and connection control functions. Through signaling, the control plane sets up and releases connections, and may restore a connection in case of a failure.

(Control) Domain: represents a collection of (control) entities that are grouped for a particular purpose. The control plane is subdivided into domains matching administrative domains. Within an administrative domain, further subdivisions of the control plane are recursively applied. A routing control domain is an abstract entity that hides the details of the RC distribution.

External NNI (E-NNI): interfaces are located between protocol controllers between control domains.

Internal NNI (I-NNI): interfaces are located between protocol controllers within control domains.

Link: [See Recommendation G.805] a "topological component" which describes a fixed relationship between a "subnetwork" or "access group" and another "subnetwork" or "access group". Links are not limited to being provided by a single server trail.

Management plane: performs management functions for the Transport Plane, the control plane and the system as a whole. It also provides coordination between all the planes. The following management functional areas are performed in the management plane: performance, fault, configuration, accounting and security management

Management domain: [See Recommendation G.805] A management domain defines a collection of managed objects which are grouped to meet organizational requirements according to geography, technology, policy or other structure, and for a number of functional areas such as configuration, security, (FCAPS), for the purpose of providing control in a consistent manner. Management domains can be disjoint, contained or overlapping. As such the resources within an administrative domain can be distributed into several possible overlapping management domains. The same resource can therefore belong to several management domains simultaneously, but a management domain shall not cross the border of an administrative domain.

SNP: The SNP is a control plane abstraction that represents an actual or potential transport plane resource. SNPs (in different subnetwork partitions) may represent the same transport resource. A one-to-one correspondence should not be assumed.

Transport plane: provides bi-directional or unidirectional transfer of user information, from one location to another. It can also provide transfer of some control and network management information.

The Transport Plane is layered; it is equivalent to the Transport Network defined in G.805.

User Network Interface (UNI): interfaces are located between protocol controllers between a user and a control domain. Note: there is no routing function associated with a UNI reference point.

This document makes use of the following terms:

Routing Area (RA): a RA represents a partition of the data plane and its identifier is used within the control plane as the representation of this partition. Per [G.8080] a RA is defined by a set of sub-networks, the TE links that interconnect them, and the interfaces representing the ends of the TE links exiting that RA. A RA may contain smaller RAs inter-connected by TE links. The limit of subdivision results in a RA that contains two sub-networks and a TE link with a single component link.

Routing Database (RDB): repository for the local topology, network topology, reachability, and other routing information that is updated as part of the routing information exchange and may additionally contain information that is configured. The RDB may contain routing information for more than one Routing Area (RA).

Routing Components: ASON routing architecture functions. These functions can be classified as protocol independent (Link Resource Manager or LRM, Routing Controller or RC) and protocol specific (Protocol Controller or PC).

Routing Controller (RC): handles (abstract) information needed for routing and the routing information exchange with peering RCs by operating on the RDB. The RC has access to a view of the RDB. The RC is protocol independent.

Note: Since the RDB may contain routing information pertaining to multiple RAs (and possibly to multiple layer networks), the RCs accessing the RDB may share the routing information.

Link Resource Manager (LRM): supplies all the relevant component and TE link information to the RC. It informs the RC about any state changes of the link resources it controls.

Protocol Controller (PC): handles protocol specific message exchanges according to the reference point over which the information is exchanged (e.g. E-NNI, I-NNI), and internal exchanges with the RC. The PC function is protocol dependent.

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