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Domain Name System (DNS) Cookies
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Abstract

DNS cookies are a lightweight DNS transaction security mechanism that provides limited protection to DNS servers and clients against a variety of increasingly common denial-of-service and amplification / forgery or cache poisoning attacks by off-path attackers. DNS Cookies are tolerant of NAT, NAT-PT, and anycast and can be incrementally deployed.

Status of This Document

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DNS Cookies

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1. Introduction

As with many core Internet protocols, the Domain Name System (DNS) was originally designed at a time when the Internet had only a small pool of trusted users. As the Internet has grown exponentially to a global information utility, the DNS has increasingly been subject to abuse.

This document describes DNS cookies, a lightweight DNS transaction security mechanism specified as an OPT [[RFC6891](#)] option. The DNS cookies mechanism provides limited protection to DNS servers and clients against a variety of increasingly common abuses by off-path attackers. It is compatible with and can be used in conjunction with other DNS transaction forgery resistance measures such as those in [[RFC5452](#)].

The protection provided by DNS cookies is similar to that provided by using TCP for DNS transactions. To bypass the weak protection provided by using TCP requires, among other things, that an off-path attacker guessing the 32-bit TCP sequence number in use. To bypass the weak protection provided by DNS Cookies requires such an attacker to guess a 64-bit pseudo-random "cookie" quantity. Where DNS Cookies are not available but TCP is, falling back to using TCP is

reasonable.

If only one party to a DNS transaction supports DNS cookies, the mechanism does not provide a benefit or significantly interfere; but, if both support it, the additional security provided is automatically available.

The DNS cookies mechanism is designed to work in the presence of NAT and NAT-PT boxes and guidance is provided herein on supporting the DNS cookies mechanism in anycast servers.

[1.1](#) Contents of This Document

In [Section 2](#), we discuss the threats against which the DNS cookie mechanism provides some protection.

[Section 3](#) describes existing DNS security mechanisms and why they are not adequate substitutes for DNS cookies.

[Section 4](#) describes the COOKIE OPT option.

[Section 5](#) provides a protocol description.

[Section 6](#) discusses some NAT and anycast related DNS Cookies design considerations.

[Section 7](#) discusses incremental deployment considerations.

Sections [8](#) and [9](#) describe IANA and Security Considerations.

[1.2](#) Definitions

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [[RFC2119](#)].

"Off-path attacker", for a particular DNS client and server, is defined as an attacker who cannot observe the DNS request and response messages between that client and server.

"Soft state" indicates information learned or derived by a host which may be discarded when indicated by the policies of that host but can be later re-instantiated if needed. For example, it could be discarded after a period of time or when storage for caching such data becomes full. If operations requiring that soft state continue after it has been discarded, it will be automatically re-generated, albeit at some cost.

"Silently discarded" indicates that there are no DNS protocol message consequences; however, it is RECOMMENDED that appropriate network management facilities be included in implementations, such as a counter of the occurrences of each such event type.

"IP address" is used herein as a length independent term and includes both IPv4 and IPv6 addresses.

[2.](#) Threats Considered

DNS cookies are intended to provide significant but limited protection against certain attacks by off-path attackers as described below. These attacks include denial-of-service, cache poisoning, and answer forgery.

[2.1](#) Denial-of-Service Attacks

The typical form of the denial-of-service attacks considered herein is to send DNS requests with forged source IP addresses to a server. The intent can be to attack that server or some other selected host as described below.

There are also on-path denial of service attacks that attempt to saturate a server with DNS requests having correct source addresses. Cookies do not protect against such attacks but successful cookie validation improves the probability that the correct source IP address for the requests is known. This facilitates contacting the managers of or taking other actions for the networks from which the requests originate.

[2.1.1](#) DNS Amplification Attacks

A request with a forged IP source address generally causes a response to be sent to that forged IP address. Thus the forging of many such requests with a particular source IP address can result in enough traffic being sent to the forged IP address to interfere with service to the host at the IP address. Furthermore, it is generally easy in the DNS to create short requests that produce much longer responses, thus amplifying the attack.

The DNS Cookies mechanism can severely limit the traffic amplification obtained by attacker requests that are off the path between the server and the request's source address. Enforced DNS cookies would make it hard for an off path attacker to cause any more than rate-limited short error responses to be sent to a forged IP address so the attack would be attenuated rather than amplified. DNS cookies make it more effective to implement a rate limiting scheme for error responses from the server. Such a scheme would further restrict selected host denial-of-service traffic from that server.

[2.1.2](#) DNS Server Denial-of-Service

DNS requests that are accepted cause work on the part of DNS servers. This is particularly true for recursive servers that may issue one or more requests and process the responses thereto, in order to determine their response to the initial request. And the situation can be even worse for recursive servers implementing DNSSEC ([\[RFC4033\]](#) [\[RFC4034\]](#) [\[RFC4035\]](#)) because they may be induced to perform burdensome cryptographic computations in attempts to verify the authenticity of data they retrieve in trying to answer the request.

The computational or communications burden caused by such requests may not depend on a forged IP source address, but the use of such addresses makes

- + the source of the requests causing the denial-of-service attack harder to find and
- + restriction of the IP addresses from which such requests should be honored hard or impossible to specify or verify.

Use of DNS cookies should enable a server to reject forged requests from an off path attacker with relative ease and before any recursive queries or public key cryptographic operations are performed.

[2.2](#) Cache Poisoning and Answer Forgery Attacks

The form of the cache poisoning attacks considered is to send forged replies to a resolver. Modern network speeds for well-connected hosts are such that, by forging replies from the IP addresses of a DNS server to a resolver for names that resolver has been induced to resolve or for common names whose resource records have short time-to-live values, there can be an unacceptably high probability of randomly coming up with a reply that will be accepted and cause false DNS information to be cached by that resolver (the Dan Kaminsky attack [\[Kaminsky\]](#)). This can be used to facilitate phishing attacks and other diversion of legitimate traffic to a compromised or malicious host such as a web server.

With the use of DNS cookies, a resolver can generally reject such forged replies.

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[3.](#) Comments on Existing DNS Security

Two forms of security have been added to DNS, data security and message/transaction security.

[3.1](#) Existing DNS Data Security

DNS data security is one part of DNSSEC and is described in [\[RFC4033\]](#), [\[RFC4034\]](#), [\[RFC4035\]](#), and updates thereto. It provides data origin authentication and authenticated denial of existence. DNSSEC is being deployed and can provide strong protection against forged data and cache poisoning; however, it has the unintended effect of making some denial-of-service attacks worse because of the cryptographic computational load it can require and the increased size in DNS response packets that it tends to produce.

[3.2](#) DNS Message/Transaction Security

The second form of security that has been added to DNS provides "transaction" security through TSIG [\[RFC2845\]](#) or SIG(0) [\[RFC2931\]](#). TSIG could provide strong protection against the attacks for which the DNS Cookies mechanism provides weak protection; however, TSIG is non-trivial to deploy in the general Internet because of the burdens it imposes. Among these burdens are pre-agreement and key distribution between client and server, keeping track of server side key state, and required time synchronization between client and server.

TKEY [\[RFC2930\]](#) can solve the problem of key distribution for TSIG but some modes of TKEY impose a substantial cryptographic computation load and can be dependent on the deployment of DNS data security (see [Section 3.1](#)).

SIG(0) [\[RFC2931\]](#) provides less denial of service protection than TSIG or, in one way, even DNS cookies, because it does not authenticate requests, only complete transactions. In any case, it also depends on the deployment of DNS data security and requires computationally burdensome public key cryptographic operations.

[3.3](#) Conclusions on Existing DNS Security

The existing DNS security mechanisms do not provide the services provided by the DNS Cookies mechanism: lightweight message authentication of DNS requests and responses with no requirement for

pre-configuration or per client server side state.

4. DNS Cookie Option

The DNS Cookie Option is an OPT RR [[RFC6891](#)] option that can be included in the RDATA portion of an OPT RR in DNS requests and responses. The option length varies depending on the circumstances in which it is being used. There are two cases as described below. Both use the same OPTION-CODE; they are distinguished by their length.

In a request sent by a client to a server when the client does not know the server's cookie, its length is 8, consisting of an 8 byte Client Cookie as shown in Figure 1.

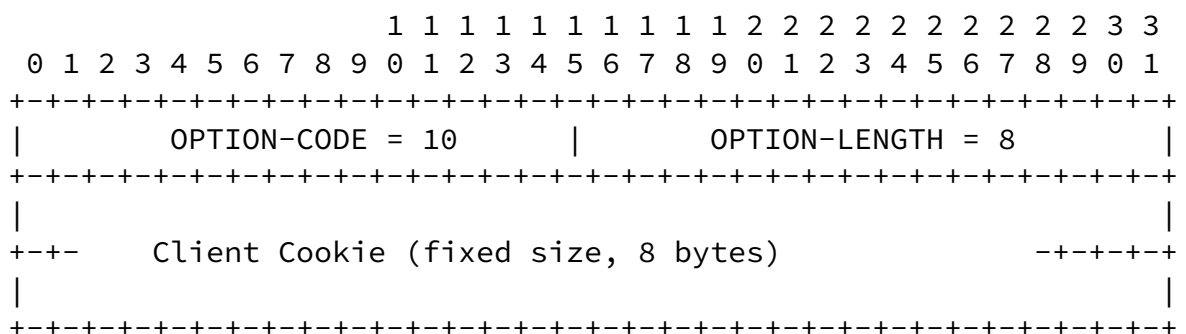


Figure 1. COOKIE Option, Unknown Server Cookie

In a request sent by a client when a server cookie is known and in

all responses, the length is variable from 16 to 40 bytes, consisting of an 8 bytes Client Cookie followed by the variable 8 to 32 bytes Server Cookie as shown in Figure 2. The variability of the option length stems from the variable length Server Cookie. The Server Cookie is an integer number of bytes with a minimum size of 8 bytes for security and a maximum size of 32 bytes for implementation convenience.

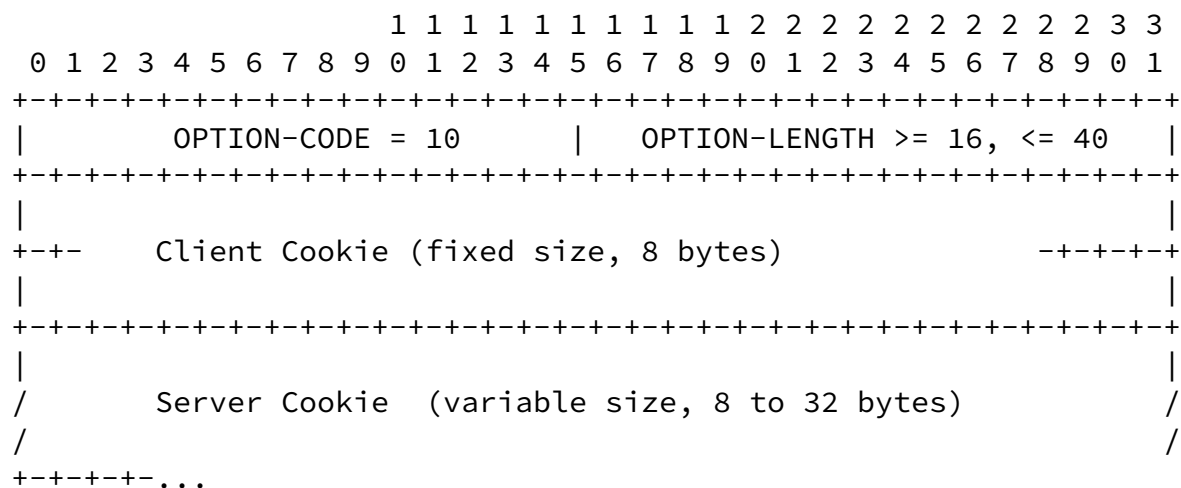


Figure 2. COOKIE Option, Known Server Cookie

[4.1](#) Client Cookie

The Client Cookie SHOULD be a pseudo-random function of the server IP address and a secret quantity known only to the client. This client secret SHOULD have at least 64 bits of entropy [[RFC4086](#)] and be changed periodically (see [Section 5.5](#)). The selection of the pseudo-random function is a matter private to the client as only the client needs to recognize its own DNS cookies.

For further discussion of the Client Cookie field, see [Section 5.1](#). For example methods of determining a Client Cookie, see [Appendix A](#).

In order to provide minimal authentication, a client MUST send client COOKIES that will usually be different for any two servers at different IP addresses.

[4.2](#) Server Cookie

The Server Cookie SHOULD consist of or include a 64-bit or larger pseudo-random function of the request source IP address, the request Client Cookie, and a secret quantity known only to the server. (See [Section 6](#) for a discussion of why the Client Cookie is used as input to the Server Cookie but the Server Cookie is not used as an input to the Client Cookie.) This server secret SHOULD have at least 64 bits of entropy [[RFC4086](#)] and be changed periodically (see [Section 5.5](#)). The selection of the pseudo-random function is a matter private to the server as only the server needs to recognize its own DNS cookies.

For further discussion of the Server Cookie field see [Section 5.2](#). For example methods of determining a Server Cookie, see [Appendix B](#).

In order to provide minimal authentication, a server MUST send server COOKIES that will usually be different for clients at any two different IP addresses or with different client COOKIES.

[5](#). DNS Cookies Protocol Specification

This section discusses using DNS Cookies in the DNS Protocol. The cycle of originating a request, responding to that request, and processing the response are covered in Sections [5.1](#), [5.2](#), and [5.3](#). A de facto extension to QUERY to allow pre-fetching a Server Cookie is specified in [Section 5.4](#). Rollover of the client and server secrets and transient retention of the old cookie or secret is covered in [Section 5.5](#).

DNS clients and servers SHOULD implement DNS cookies to decrease their vulnerability to the threats discussed in [Section 2](#).

[5.1](#) Originating Requests

A DNS client that implements DNS Cookies includes one DNS COOKIE OPT option containing a Client Cookie in every DNS request it sends unless DNS cookies are disabled.

If the client has a cached Server Cookie for the server against its IP address it uses the longer cookie form and includes that Server Cookie in the option along with the Client Cookie (Figure 2). Otherwise it just sends the shorter form option with a Client Cookie (Figure 1).

[5.2](#) Responding to Request

The Server Cookie, when it occurs in a COOKIE OPT option in a request, is intended to weakly assure the server that the request came from a client that is both at the source IP address of the request and using the Client Cookie included in the option. This weak assurance is provided by the Server Cookie that server sent to that client in an earlier response appearing as the Server Cookie field in the request.

At a server where DNS Cookies are not implemented and enabled, presence of a COOKIE OPT option is ignored and the server responds as if no COOKIE OPT option had been included in the request.

When DNS Cookies are implemented and enabled, there are five possibilities: (1) there is no OPT RR at all in the request or there is a OPT RR but the the COOKIE OPT option is absent from the OPT RR; (2) a COOKIE OPT is present but is not a legal length or otherwise malformed; (3) there is a valid length cookie option in the request with no Server Cookie; (4) there is a valid length COOKIE OPT in the request with a Server Cookie but that Server Cookie is invalid; or

(5) there is a valid length COOKIE OPT in the request with a correct Server Cookie.

The five possibilities are discussed in the subsections below.

In all cases of multiple COOKIE OPT options in a request, only the first (the one closest to the DNS header) is considered. All others are ignored.

[5.2.1](#) No Opt RR or No COOKIE OPT option

If there is no OPT record or no COOKIE OPT option present in the request then the server responds to the request as if the server doesn't implement the COOKIE OPT.

[5.2.2](#) Malformed COOKIE OPT option

If the COOKIE OPT is too short to contain a Client Cookie then FORMERR is generated. If the COOKIE OPT is longer than that required to hold a COOKIE OPT with just a Client Cookie (8) but is shorter than the minimum COOKIE OPT with both a Client and Server Cookie (16) then FORMERR is generated. If the COOKIE OPT is longer than the maximum valid COOKIE OPT (40) then a FORMERR is generated.

In summary, valid cookie lengths are 8 and 16 to 40 inclusive.

[5.2.3](#) Only a Client Cookie

Based on server policy, including rate limiting, the server chooses one of the following:

- (1) Silently discard the request.
- (2) Send a BADCOOKIE error response.
- (3) Process the request and provide a normal response. The RCODE is NOERROR unless some non-cookie error occurs in processing the request.

If the server responds, choosing 2 or 3 above, it SHALL generate its own COOKIE OPT containing both the Client Cookie copied from the request and a Server Cookie it has generated and adds this COOKIE OPT to the response's OPT record. Servers MUST, at least occasionally, respond to such requests to inform the client of the correct Server

Cookie. This is necessary so that such a client can bootstrap to the weakly secure state where requests and responses have recognized Server Cookies and Client Cookies. A server is not expected to maintain per client state to achieve this. For example, it could respond to every Nth request across all clients.

If the request was received over TCP, the server SHOULD take the weak authentication provided by the use of TCP into account and SHOULD choose 3. In this case, if the server is not willing to accept the weak security provided by TCP as a substitute for the weak security provided by DNS Cookies but instead chooses 2, there is some danger of an indefinite loop of retries (see [Section 5.3](#)).

[5.2.4](#) A Client Cookie and an Invalid Server Cookie

The server examines the Server Cookie to determine if it is a valid Server Cookie it has generated. This examination will result in a determination of whether the Server Cookie is valid or not. If the cookie is invalid, it can be because of a stale Server Cookie, or a client's IP address or Client Cookie changing without the DNS server being aware, or an anycast server cluster that is not consistently configured, or an attempt to spoof the client.

The server SHALL process the request as if the invalid Server Cookie was not present as described in [Section 5.2.3](#).

[5.2.5](#) A Client Cookie and a Valid Server Cookie

When a valid Server Cookie is present in the request the server can assume that the request is from a client that it has talked to before and defensive measures for spoofed UDP requests, if any, are no longer required.

The server SHALL process the request and include a COOKIE OPT in the response by (a) copying the complete COOKIE OPT from the request or (b) generating a new COOKIE OPT containing both the Client Cookie copied from the request and a valid Server Cookie it has generated.

[5.3](#) Processing Responses

The Client Cookie, when it occurs in a COOKIE OPT option in a DNS reply, is intended to weakly assure the client that the reply came from a server at the source IP address used in the response packet because the Client Cookie value is the value that client would send

to that server in a request. In a DNS reply with multiple COOKIE OPT options, all but the first (the one closest to the DNS Header) are ignored.

A DNS client where DNS cookies are implemented and enabled examines the response for DNS cookies and MUST discard the response if it contains an illegal COOKIE OPT option length or an incorrect Client Cookie value. If the COOKIE OPT option Client Cookie is correct, the client caches the Server Cookie provided even if the response is an error response (RCODE non-zero).

If the reply extended RCODE is BADCOOKIE and the Client Cookie matches what was sent, it means that the server was unwilling to process the request because it did not have the correct Server Cookie in it. The client SHOULD retry the request using the new Server Cookie from the response. Repeated BADCOOKIE responses to requests that use the Server Cookie provided in the previous response may be an indication that the shared secrets / secret generation method in an anycast cluster of servers are inconsistent. If the reply to a retried request with a fresh Server Cookie is BADCOOKIE, the client SHOULD retry using TCP as the transport since the server will likely process the request normally based on the weak security provided by TCP (see [Section 5.2.3](#)).

If the RCODE is some value other than BADCOOKIE, including zero, the further processing of the response proceeds normally.

[5.4](#) QUERYing for a Server Cookie

In many cases a client will learn the Server Cookie for a server as the side effect of another transaction; however, there may be times when this is not desirable. Therefore a means is provided for obtaining a Server Cookie through an extension to the QUERY opcode for which opcode most existing implementations require that QDCOUNT be one (see [Section 4.1.2 of \[RFC1035\]](#)).

For servers with DNS Cookies enabled, the QUERY opcode behavior is extended to support queries with a empty question section (QDCOUNT zero) provided that an OPT record is present with a COOKIE option. Such servers will reply with an empty answer section and a COOKIE option giving the Client Cookie provided in the query and a valid Server Cookie.

If such a query provided just a Client Cookie and no Server Cookie, the response SHALL have the RCODE NOERROR.

This mechanism can also be used to confirm/re-establish a existing Server Cookie by sending a cached Server Cookie with the Client

Cookie. In this case the response SHALL have the RCODE BADCOOKIE if the Server Cookie sent with the query was invalid and the RCODE NOERROR if it was valid.

Servers which don't support the COOKIE option will normally send FORMERR in response to such a query, though REFUSED, NOTIMP, and NOERROR without a COOKIE option are also possible in such responses.

[5.5](#) Client and Server Secret Rollover

The longer a secret is used, the higher the probability it has been compromised. Thus clients and servers MUST NOT continue to use the same secret in new requests and responses for more than 36 days and SHOULD NOT continue to do so for more than 26 hours. These values are chosen to assure that a secret will not be used for longer than about a month and normally no longer than one day. The odd values are to allow for long holiday weekends and daylight savings time shifts and the like while still staying within the limits.

Many clients rolling over their secret at the same time could briefly increase server traffic and exactly predictable rollover times for clients or servers might facilitate guessing attacks. For example, an attacker might increase the priority of attacking secrets they believe will be in effect for an extended period of time. To avoid rollover synchronization and predictability, it is RECOMMENDED that pseudorandom jitter in the range of plus zero to minus at least 40% be applied to the time until a scheduled rollover of a DNS cookie secret.

It is RECOMMENDED that a client keep the Client Cookie it is expecting in a reply associated with the outstanding request to avoid rejection of replies due to a bad Client Cookie right after a change in the client secret. It is RECOMMENDED that a server retain its previous secret for a period of time not less than 1 second or more than 5 minutes, after a change in its secret, and consider requests with Server Cookies based on its previous secret to have a correct Server Cookie during that time.

When a server or client starts receiving an increased level of requests with bad server cookies or replies with bad client cookies, it would be reasonable for it to believe it is likely under attack and it should consider a more frequent rollover of its secret. More rapid rollover decreases the benefit to a cookie guessing attacker if they succeed in guessing a cookie.

6. NAT Considerations and AnyCast Server Considerations

In the Classic Internet, DNS Cookies could simply be a pseudo-random function of the client IP address and a server secret or the server IP address and a client secret. You would want to compute the Server Cookie that way, so a client could cache its Server Cookie for a particular server for an indefinitely amount of time and the server could easily regenerate and check it. You could consider the Client Cookie to be a weak client signature over the server IP address that the client checks in replies and you could extend this weak signature to cover the request ID, for example, or any other information that is returned unchanged in the reply.

But we have this reality called NAT [[RFC3022](#)], Network Address Translation (including, for the purposes of this document, NAT-PT, Network Address and Protocol Translation, which has been declared Historic [[RFC4966](#)]). There is no problem with DNS transactions between clients and servers behind a NAT box using local IP addresses. Nor is there a problem with NAT translation of internal addresses to external addresses or translations between IPv4 and IPv6 addresses, as long as the address mapping is relatively stable. Should the external IP address an internal client is being mapped to

change occasionally, the disruption is little more than when a client rolls-over its DNS COOKIE secret. And normally external access to a DNS server behind a NAT box is handled by a fixed mapping which forwards externally received DNS requests to a specific host.

However, NAT devices sometimes also map ports. This can cause multiple DNS requests and responses from multiple internal hosts to be mapped to a smaller number of external IP addresses, such as one address. Thus there could be many clients behind a NAT box that appear to come from the same source IP address to a server outside that NAT box. If one of these were an attacker (think Zombie or Botnet), that behind-NAT attacker could get the Server Cookie for some server for the outgoing IP address by just making some random request to that server. It could then include that Server Cookie in the COOKIE OPT of requests to the server with the forged local IP address of some other host and/or client behind the NAT box. (Attacker possession of this Server Cookie will not help in forging responses to cause cache poisoning as such responses are protected by the required Client Cookie.)

To fix this potential defect, it is necessary to distinguish different clients behind a NAT box from the point of view of the server. It is for this reason that the Server Cookie is specified as a pseudo-random function of both the request source IP address and the Client Cookie. From this inclusion of the Client Cookie in the calculation of the Server Cookie, it follows that a stable Client Cookie, for any particular server, is needed. If, for example, the request ID was included in the calculation of the Client Cookie, it

would normally change with each request to a particular server. This would mean that each request would have to be sent twice: first to learn the new Server Cookie based on this new Client Cookie based on the new ID and then again using this new Client Cookie to actually get an answer. Thus the input to the Client Cookie computation must be limited to the server IP address and one or more things that change slowly such as the client secret.

In principle, there could be a similar problem for servers, not due to NAT but due to mechanisms like anycast which may cause requests to a DNS server at an IP address to be delivered to any one of several machines. (External requests to a DNS server behind a NAT box usually occur via port forwarding such that all such requests go to one host.) However, it is impossible to solve this the way the similar

problem was solved for NATed clients; if the Server Cookie was included in the calculation of the Client Cookie the same way the Client Cookie is included in the Server Cookie, you would just get an almost infinite series of errors as a request was repeatedly retried.

For servers accessed via anycast to successfully support DNS COOKIES, the server clones must either all use the same server secret or the mechanism that distributes requests to them must cause the requests from a particular client to go to a particular server for a sufficiently long period of time that extra requests due to changes in Server Cookie resulting from accessing different server machines are not unduly burdensome. (When such anycast-accessed servers act as recursive servers or otherwise act as clients they normally use a different unique address to source their requests to avoid confusion in the delivery of responses.)

For simplicity, it is RECOMMENDED that the same server secret be used by each DNS server in a set of anycast servers. If there is limited time skew in updating this secret in different anycast servers, this can be handled by a server accepting requests containing a Server Cookie based on either its old or new secret for the maximum likely time period of such time skew (see also [Section 5.5](#)).

[7](#). Deployment

The DNS cookies mechanism is designed for incremental deployment and to complement the orthogonal techniques in [[RFC5452](#)]. Either or both techniques can be deployed independently at each DNS server and client.

In particular, a DNS server or client that implements the DNS COOKIE mechanism can interoperate successfully with a DNS client or server that does not implement this mechanism although, of course, in this case it will not get the benefit of the mechanism and the server involved might choose to severely rate limit responses. When such a server or client interoperates with a client or server which also implements the DNS cookies mechanism, they get the weak security benefits of the DNS Cookies mechanism.

[8.](#) IANA Considerations

IANA has assigned the following OPT option value:

Value	Name	Status	Reference
-----	-----	-----	-----
10	COOKIE	Standard	[this document]

IANA has assigned the following DNS error code as an early allocation:

RCODE	Name	Description	Reference
-----	-----	-----	-----
23	BADCOKIE	Bad/missing server cookie	[this document]

[9.](#) Security Considerations

DNS Cookies provide a weak form of authentication of DNS requests and responses. In particular, they provide no protection against "on-path" adversaries; that is, they provide no protection against any adversary that can observe the plain text DNS traffic, such as an on-path router, bridge, or any device on an on-path shared link (unless the DNS traffic in question on that path is encrypted).

For example, if a host is connected via an unsecured IEEE Std 802.11 link (Wi-Fi), any device in the vicinity that could receive and decode the 802.11 transmissions must be considered "on-path". On the other hand, in a similar situation but one where 802.11 Robust Security (WPAv2) is appropriately deployed on the Wi-Fi network nodes, only the Access Point via which the host is connecting is "on-path" as far as the 802.11 link is concerned.

Despite these limitations, deployment of DNS Cookies on the global Internet is expected to provide a significant reduction in the available launch points for the traffic amplification and denial of service forgery attacks described in [Section 2](#) above.

Should stronger message/transaction security be desired, it is suggested that TSIG or SIG(0) security be used (see [Section 3.2](#)); however, it may be useful to use DNS Cookies in conjunction with these features. In particular, DNS Cookies could screen out many DNS messages before the cryptographic computations of TSIG or SIG(0) are required and, if SIG(0) is in use, DNS Cookies could usefully screen out many requests given that SIG(0) does not screen requests but only authenticates the response of complete transactions.

[9.1](#) Cookie Algorithm Considerations

The cookie computation algorithm for use in DNS Cookies SHOULD be based on a pseudo-random function at least as strong as 64-bit FNV (Fowler-Noll-Vo [[FNV](#)]) because an excessively weak or trivial algorithm could enable adversaries to guess cookies. However, in light of the weak plain-text token security provided by DNS Cookies, a strong cryptography hash algorithm may not be warranted in many

cases, and would cause an increased computational burden. Nevertheless there is nothing wrong with using something stronger, for example, HMAC-SHA256-64 [[RFC6234](#)], assuming a DNS processor has adequate computational resources available. DNS processors that feel the need for somewhat stronger security without a significant increase in computational load should consider more frequent changes in their client and/or server secret; however, this does require more frequent generation of a cryptographically strong random number [[RFC4086](#)]. See Appendices A and B for specific examples of cookie

computation algorithms.

[10](#). Implementation Considerations

The DNS Cookie Option specified herein is implemented in BIND 9.10 using a experimental option code.

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Normative References

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Appendix A: Example Client Cookie Algorithms

[A.1](#) A Simple Algorithm

An simple example method to compute Client Cookies is the FNV-64 [[FNV](#)] of the server IP address and the client secret. That is

$$\text{Client Cookie} = \text{FNV-64} (\text{Client Secret} \mid \text{Server IP Address})$$

where " \mid " indicates concatenation.

[A.2](#) A More Complex Algorithm

A more complex algorithm to calculate Client Cookies is given below. It uses more computational resources than the simpler algorithm shown in A.1.

$$\text{Client Cookie} = \text{HMAC-SHA256-64} (\text{Client Secret}, \\ \text{Server IP Address})$$

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Appendix B: Example Server Cookie Algorithms

[B.1](#) A Simple Algorithm

An example of a simple method producing a 64-bit Server Cookie is the FNV-64 [[FNV](#)] of the request IP address, the Client Cookie, and the server secret. That is

$$\text{Server Cookie} = \text{FNV-64} (\text{Server Secret} \mid \text{Request IP Address} \mid \text{Client Cookie})$$

where " \mid " represents concatenation.

[B.2](#) A More Complex Algorithm

Since the Server Cookie has a variable size, the server can store various information in that field as long as it is hard for an adversary to guess the entire quantity used for weak authentication. There should be 64 bits of entropy in the Server Cookie; for example it could have a sub-field of 64-bits computed pseudo-randomly with the server secret as one of the inputs to the pseudo-random function. Types of additional information that could be stored include a time stamp and/or a nonce.

The example below is one variation for the Server Cookie that has been implemented in a beta release of BIND where the Server Cookie is 128 bits composed as follows:

Sub-field	Size
-----	-----
Nonce	32 bits

Time	32 bits
Hash	64 bits

With this algorithm, the server sends a new 128-bit cookie back with every request. The Nonce field assures a low probability that there would be a duplicate.

The Time field gives the server time and makes it easy to reject old cookies.

The Hash part of the Server Cookie is the hard-to-guess part. In the beta release of BIND, its computation can be configured to use AES, HMAC-SHA1, or, as shown below, HMAC-SHA256:

```
hash =  
    HMAC-SHA256-64 ( Server Secret,  
                    (Client Cookie | nonce | time | client IP Address) )
```

where "|" represents concatenation.

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