RSA/MD5 KEYs and SIGs in the Domain Name System (DNS)

Donald E. Eastlake 3rd

Status of This Document

This draft, file name draft-ietf-dnssec-rsa-00.txt, is intended to be become a Proposed Standard RFC. Distribution of this document is unlimited. Comments should be sent to the DNS security mailing list <dns-security@tis.com> or to the author.

Abstract

A standard method for storing RSA keys and and RSA/MD5 based signatures in the Domain Name System is described which utilizes DNS KEY and SIG resource records.

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1. Introduction

The Domain Name System (DNS) is the global hierarchical replicated distributed database system for Internet addressing, mail proxy, and other information. The DNS has been extended to include digital signatures and cryptographic keys as described in [draft-ietf-dnssec-secext2-*]. Thus the DNS can now be secured and used for secure key distribution.

This document describes how to store RSA keys and and RSA/MD5 based signatures in the DNS. Familiarity with the RSA algorithm is assumed [Schneier]. Implementation of the RSA algorithm in DNS is recommended.

2. RSA Public KEY Resource Records

RSA public keys are stored in the DNS as KEY RRs using algorithm number 1 [draft-ietf-dnssec-secext2-*]. The structure of the algorithm specific portion of the RDATA part of such RRs is as shown below.

Field	Size
exponent length	1 or 3 octets (see text)
exponent	as specified by length field
modulus	remaining space

For interoperability, the exponent and modulus are each currently limited to 4096 bits in length. The public key exponent is a variable length unsigned integer. Its length in octets is represented as one octet if it is in the range of 1 to 255 and by a zero octet followed by a two octet unsigned length if it is longer than 255 bytes. The public key modulus field is a multiprecision unsigned integer. The length of the modulus can be determined from the RDLENGTH and the preceding RDATA fields including the exponent. Leading zero octets are prohibited in the exponent and modulus.

3. RSA/MD5 SIG Resource Records

The signature portion of the SIG RR RDATA area, when using the RSA/MD5 algorithm, is calculated as shown below. The data signed is determined as specified in [draft-ietf-dnssec-secext2-*]. See [draft-ietf-dnssec-secext2-*] for fields in the SIG RR RDATA which precede the signature itself.

```
hash = MD5 ( data )

signature = ( 01 | FF* | 00 | prefix | hash ) ** e (mod n)
```

where MD5 is the message digest algorithm documented in [RFC 1321], "|" is concatenation, "e" is the private key exponent of the signer, and "n" is the modulus of the signer's public key. 01, FF, and 00 are fixed octets of the corresponding hexadecimal value. "prefix" is the ASN.1 BER MD5 algorithm designator prefix specified in PKCS1, that is,

```
hex 3020300c06082a864886f70d020505000410 [NETSEC].
```

This prefix is included to make it easier to use RSAREF (or similar packages such as EuroRef). The FF octet MUST be repeated the maximum number of times such that the value of the quantity being exponentiated is one octet shorter than the value of n.

(The above specifications are identical to the corresponding part of Public Key Cryptographic Standard #1 [PKCS1].)

The size of n, including most and least significant bits (which will be 1) MUST be not less than 512 bits and not more than 4096 bits. n and e SHOULD be chosen such that the public exponent is small.

Leading zero bytes are permitted in the RSA/MD5 algorithm signature.

A public exponent of 3 minimizes the effort needed to verify a signature. Use of 3 as the public exponent is weak for confidentiality uses since, if the same data can be collected encrypted under three different keys with an exponent of 3 then, using the Chinese Remainder Theorem [NETSEC], the original plain text can be easily recovered. This weakness is not significant for DNS security because we seek only authentication, not confidentiality.

4. Performance Considerations

General signature generation speeds are roughly the same for RSA and DSA [RFC xDSA]. With sufficient pre-computation, signature generation with DSA is faster than RSA. Key generation is also faster for DSA. However, signature verification is an order of magnitude slower with DSA when the RSA public exponent is chosen to be small as is recommended for KEY RRs used in domain name system (DNS) data authentication.

Current DNS implementations are optimized for small transfers, typically less than 512 bytes including overhead. While larger transfers will perform correctly and work is underway to make larger transfers more efficient, it is still advisable at this time to make reasonable efforts to minimize the size of KEY RR sets stored within the DNS consistent with adequate security. Keep in mind that in a secure zone, at least one authenticating SIG RR will also be returned.

5. Security Considerations

Many of the general security consideration in [draft-ietf-dnssec-secext2-*] apply. Keys retrieved from the DNS should not be trusted unless (1) they have been securely obtained from a secure resolver or independently verified by the user and (2) this secure resolver and secure obtainment or independent verification conform to security policies acceptable to the user. As with all cryptographic algorithms, evaluating the necessary strength of the key is essential and dependent on local policy.

For interoperability, the RSA key size is limited to 4096 bits. For particularly critical applications, implementors are encouraged to consider the range of available algorithms and key sizes.

References

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Author's Address

Donald E. Eastlake 3rd
CyberCash, Inc.
318 Acton Street
Carlisle, MA 01741 USA

Telephone: +1 978 287 4877
+1 703 620-4200 (main office, Reston, Virginia)
FAX: +1 978 371 7148
EMail: dee@cybercash.com

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