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**Making BGP filtering a habit: Impact on policies
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Abstract

This document describes how unexpected traffic flows can emerge across an autonomous system, as the result of other autonomous systems filtering, or restricting the propagation of overlapping prefixes. We provide a review of the techniques to detect the occurrence of this issue and defend against it.

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[1.](#) Introduction

It is common practice for network operators to propagate a more specific (overlapping) prefix in the BGP routing system, along with the covering prefix that they originate. It is also possible for some Autonomous Systems (ASes) to apply different policies to the overlapping and the covering prefix.

While BGP makes independent, policy driven decisions for the selection of the best path to be used for a given IP prefix, routers must forward packets using the longest-prefix-match rule, which "precedes" any BGP policy ([RFC1812](#) [1]). The existence of a prefix p that is more specific than a prefix p' in the Forwarding Information Base (FIB) will let packets whose destination matches p be forwarded

according to the next hop selected as best for p (the overlapping prefix). This process takes place by disregarding the policies applied in the control plane for the selection of the best next-hop for p'. When an Autonomous System filters overlapping prefixes and forwards packets according to the covering prefix, the discrepancy in the routing policies applied to covering and overlapping prefixes can create unexpected traffic flows that infringe the policies of other ASes, still holding a path towards the overlapping prefix.

The objective of this draft is to shed light on possible side effects associated with overlapping prefix filtering. This document presents examples of such side effects and discusses approaches towards solutions to the problem.

The rest of the document is organized as follows: In [Section 2](#) we provide some scenarios in which the filtering of overlapping prefixes leads to the creation of unexpected traffic flows. [Section 3](#) and [Section 4](#) discuss some techniques that ASes can use for, respectively, detect and react to unexpected traffic flows. We conclude in [Section 5](#).

[1.1](#). Terminology

Overlapping prefix: A prefix in the routing table with an address range that is covered by another prefix present in the routing table.

Covering prefix: A prefix in the routing table with an address range partially covered by other prefixes.

We re-use the definitions of customer-transit peering and settlement-free peering of [RFC4384](#) [2].

Selective advertisement: The behavior of only advertising a self originated BGP path for a prefix over a strict subset of the eBGP sessions of the AS.

Selective propagation: The behavior of only propagating a BGP path for a prefix over a strict subset of the eBGP sessions of an AS.

Local filtering: The behavior of explicitly ignoring a BGP path received over an eBGP session.

Remote filtering: The behavior of triggering selective propagation of a BGP path at a distant AS. Note that this is typically achieved by tagging a self-originated path with BGP communities defined by the distant AS.

Unexpected traffic flow: Traffic flowing between two neighboring ASes of an AS, although the transit policy of that AS is to not provide connectivity between these two neighbors. A traffic flow across an AS, between two of its transit providers, or between a transit provider and one of its settlement-free peers, are classical examples of unexpected traffic flows.

2. Unexpected Traffic Flows

In this section, we describe how overlapping prefix filtering can lead to unexpected traffic flows in other, remote, ASes. We differentiate cases in which the filtering is performed locally from those where the filtering is triggered remotely.

2.1. Local filtering

Local filtering can be motivated by different reasons, such as: (1) Traffic engineering, where an AS wants to control their local outbound traffic distribution using only the policy applied to the covering prefix. (2) Enforcing contract compliance, where, for instance, an AS avoids a settlement-free peer to attract traffic to one link by using selective advertisement, when this is not allowed by their peering agreement.

Figure 1 illustrates a scenario in which one AS is motivated to perform local filtering due to outbound traffic engineering. The figure depicts AS64504, and two of its neighboring ASes, AS64502 and AS64505. AS 64504 has a settlement-free peering with AS64502 and is a customer of AS64505. AS64504 receives from AS64505 prefixes 2001:DB8::/32 and 2001:DB8::/34, covering and overlapping prefixes, respectively. AS64504 receives only the covering prefix 2001:DB8::/32 from AS64502.



Figure 1: Local Filtering

Due to economical reasons, AS64504 might prefer to send traffic to AS64502 instead of AS64505. However, even if paths received from AS64502 are given a large local preference, routers in AS64504 will still send traffic to prefix 2001:DB8::/34 via neighbor AS64505. This situation may push AS64504 to apply an inbound filter for the overlapping prefix, 2001:DB8::/34, on the session with AS64505. After the filter is applied, traffic to the overlapping prefix will be sent to AS64502.

2.1.1. Unexpected traffic flows caused by local filtering of overlapping prefixes

In this section, we show how the decision of AS64504 to perform local filtering creates unexpected traffic flows in AS64502. Figure 2 shows the whole picture of the scenario; where AS64501 is a customer of AS64503 and AS64502. AS64503 is a settlement-free peer with AS64502. AS64503 and AS64502 are customers of AS64505. The AS originating the two prefixes, AS64501, performs selective advertisement with the overlapping prefix and only announces it to AS64503.

After AS64504 locally filters the overlapping prefix, traffic from AS64504 to prefix 2001:DB8::/34 is forwarded towards AS64502. Because AS64502 receives the more specific prefix from AS64503, traffic from AS64504 to 2001:DB8::/34 follows the path AS64504-AS64502-AS64503-AS64501. AS64502's BGP policies are implemented to avoid transporting traffic between AS64504 and AS64503. However, due to the discrepancies of routes from the overlapping and covering prefixes, unexpected traffic flows between AS64504 and AS64503 exist in AS64502's network.

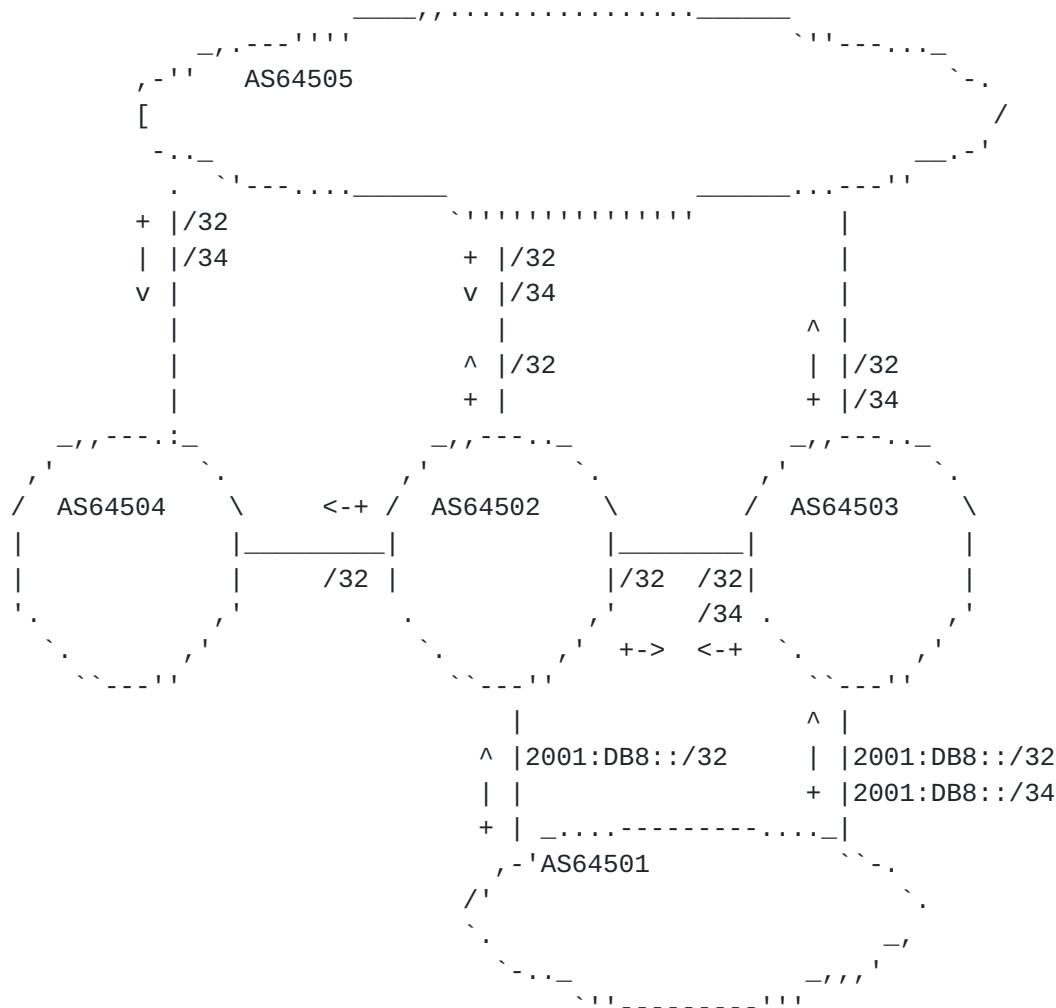


Figure 2: Unexpected traffic flows due to local filtering

2.2. Remote filtering

ISPs can tag the BGP paths that they propagate to neighboring ASes with communities, in order to tweak the propagation behavior of the ASes that handle these paths [1]. Some ISPs allow their customers to use such communities to let the receiving AS not export the path to some selected neighboring ASes. By combining communities, the prefix could be advertised only to a given peer of the AS providing this feature. Remote filtering can be leveraged by an AS to, for instance, limit the scope of prefixes and hence perform a more granular inbound traffic engineering.

Figure 3 illustrates a scenario in which an AS uses BGP communities to command its provider to selectively propagate an overlapping prefix. Let AS64501 be a customer of AS64502 and AS64503. AS64501

originates prefix 2001:DB8::/32, which it advertises through AS64502 and AS64503. AS64502 and AS64503 are settlement-free peers. Let AS64501 do selective advertisement and only propagate 2001:DB8::/34 over AS64503. AS64503 would normally propagate this prefix to its customers, providers, and peers, including AS64502.

Let us consider that AS64501 decides to limit the scope of the overlapping prefix. AS64501 can make this decision based on its traffic engineering strategy. To achieve this, AS64501 can tag the overlapping prefix with a set of communities that leads AS64503 to only propagate the path to AS64502.

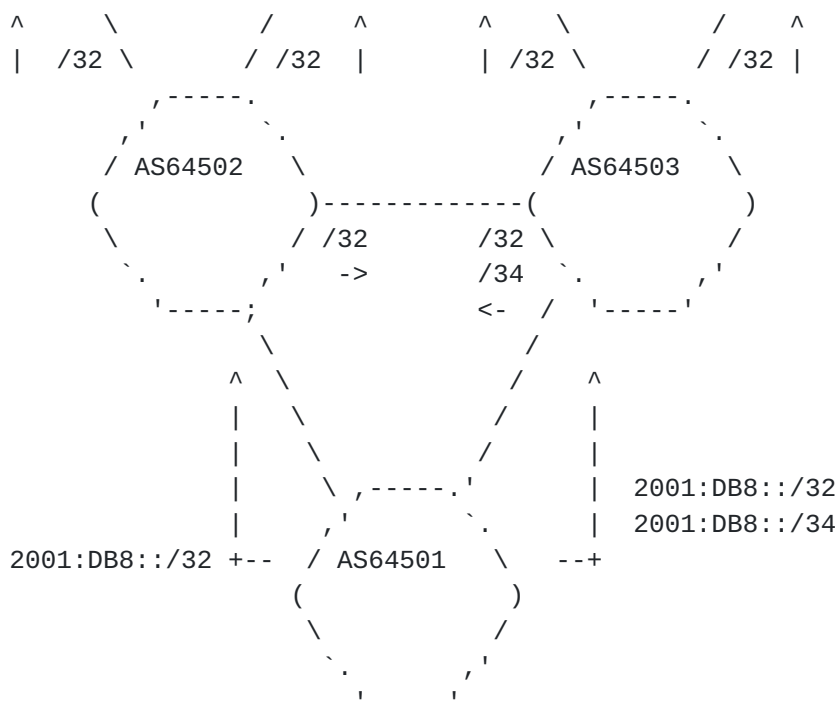


Figure 3: Remote triggered filtering

2.2.1. Unexpected traffic flows caused by remotely triggered filtering of overlapping prefixes

Figure 4 expands the scenario from Figure 3 and includes other AS peering with ASes 64502 and 64503. Due to the limitation on the scope performed on the overlapping prefix, ASes that are not customers of AS64502 will not receive a path for 2001:DB8::/34. These ASes will forward packets destined to 2001:DB8::/34 according to their routing state for 2001:DB8::/32. Let us assume that AS64505 is such an AS, and that its best path towards 2001:DB8::/32 is through AS64502. Packets sent towards 2001:DB8::1 by AS64505 will reach AS64502. However, in the data-plane of the nodes of AS64502, the longest prefix match for 2001:DB8::1 is 2001:DB8::/34, which is

reached through AS64503, a settlement-free peer of AS64502. Since AS64505 is not in the customer branch of AS64502, we are in a situation in which traffic flows between non-customer ASes take place in AS64502.

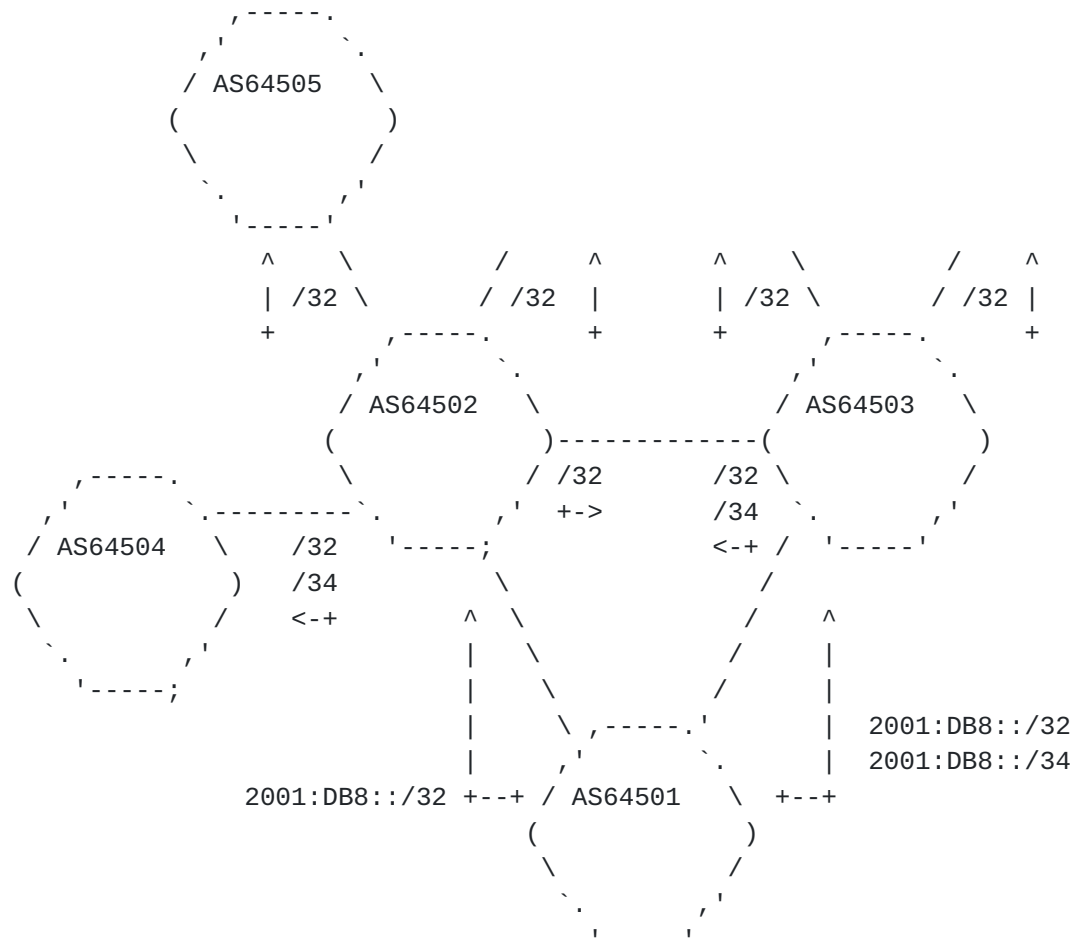


Figure 4: Unexpected traffic flows due to remote triggered filtering

3. Techniques to detect unexpected traffic flows caused by filtering of overlapping prefixes

3.1. Existence of unexpected traffic flows within an AS

To detect if unexpected traffic flows are taking place in its network, an ISP can monitor its traffic data to check if it is providing transit between two of its peers, although his policy is configured to not provide such transit. IPFIX ([RFC7011](#) [3]) is an example of a technology that can be used to export information regarding traffic flows across the network. Traffic information must be analyzed under the perspective of the business relationships with

neighboring ASes. Open source tools such as [\[4\]](#) can be used to this end.

Note that the AS detecting the unexpected traffic flow may simply realize that his policy configuration is broken. The first recommended action upon detection of an unexpected traffic flow is to verify the correctness of the BGP configuration.

Once it has been assessed that the local configuration is correct, the operator should check if the problem detected in the data-plane arose due to filtering of BGP paths by neighboring ASes. The operator should check if the destination address of the unexpected traffic flow is locally routed as per an overlapping prefix only received from non-customer peers. The operator should also check if there are paths to a covering prefix received from a customer, and hence propagated to peers. If these two situations happen at the same time, the neighboring AS at the entry point of the unexpected flow is routing the traffic based on the covering prefix, although the ISP is actually forwarding the traffic via non-customer links.

To detect the origin of the problem, human interaction or looking glasses can be used in order to find out whether local filtering, remote filtering, or selective propagation was performed on the overlapping prefix. Due to the distributed nature and restricted visibility of the steering of BGP policies, such analysis is deemed to not identify the origin of the problem with guaranteed accuracy. We are not aware, at the time of this writing, of any openly available tool that can automatically perform this operation.

[3.2.](#) Contribution to the existence of unexpected traffic flows in another AS

It can be considered problematic to be causing unexpected traffic flows in other ASes. This situation may appear as an abuse to the network resources of other ISPs.

There may be justifiable reasons for one ISP to perform filtering; either to enforce established policies or to provide prefix advertisement scoping features to its customers. These can vary from trouble-shooting purposes to business relationship implementations. Restricting the use of these features for the sake of avoiding the creation of unexpected traffic flows is not a practical option.

It is advisable for an AS to assess the risks of filtering overlapping prefixes before implementing them by obtaining as much data information as possible about its surrounding routing environment. The AS would need information of the routing policies and the relationships among external ASes to detect if its actions

could trigger the appearance of unexpected traffic flows. With this information, the operator could detect other ASes receiving the overlapping prefix from non-customer ASes, while announcing the covering prefix to other non-customer ASes. If the filtering of the overlapping prefix leads other ASes to send traffic for the overlapping prefix to these ASes, an unexpected traffic flow can arise. However, the information required for this operation is difficult to obtain, due to the distributed nature of BGP policies. We are not aware, at the time of this writing, of any openly available tool that can automatically perform this procedure.

4. Techniques to counter unexpected traffic flows

Network Operators can adopt different approaches with respect to unexpected traffic flows. Note that due the complexity of inter-domain routing policies, there is not a single solution that can be applied to all situations. We provide potential solutions that ISPs must evaluate against each particular case. We classify these actions according to whether they are anticipant or reactive.

Reactive approaches are those in which the operator tries to detect the situations via monitoring and solve unexpected traffic flows, manually, on a case-by-case basis.

Anticipant or preventive approaches are those in which the routing system will not let the unexpected traffic flows actually take place when the scenario arises.

We use the scenario depicted in Figure 5 to describe these two kinds of approaches. Based on our analysis, we observe that anticipant approaches can be complex to implement and can lead to undesired effects. Therefore, we conclude that the reactive approach is the more reasonable recommendation to deal with unexpected flows.

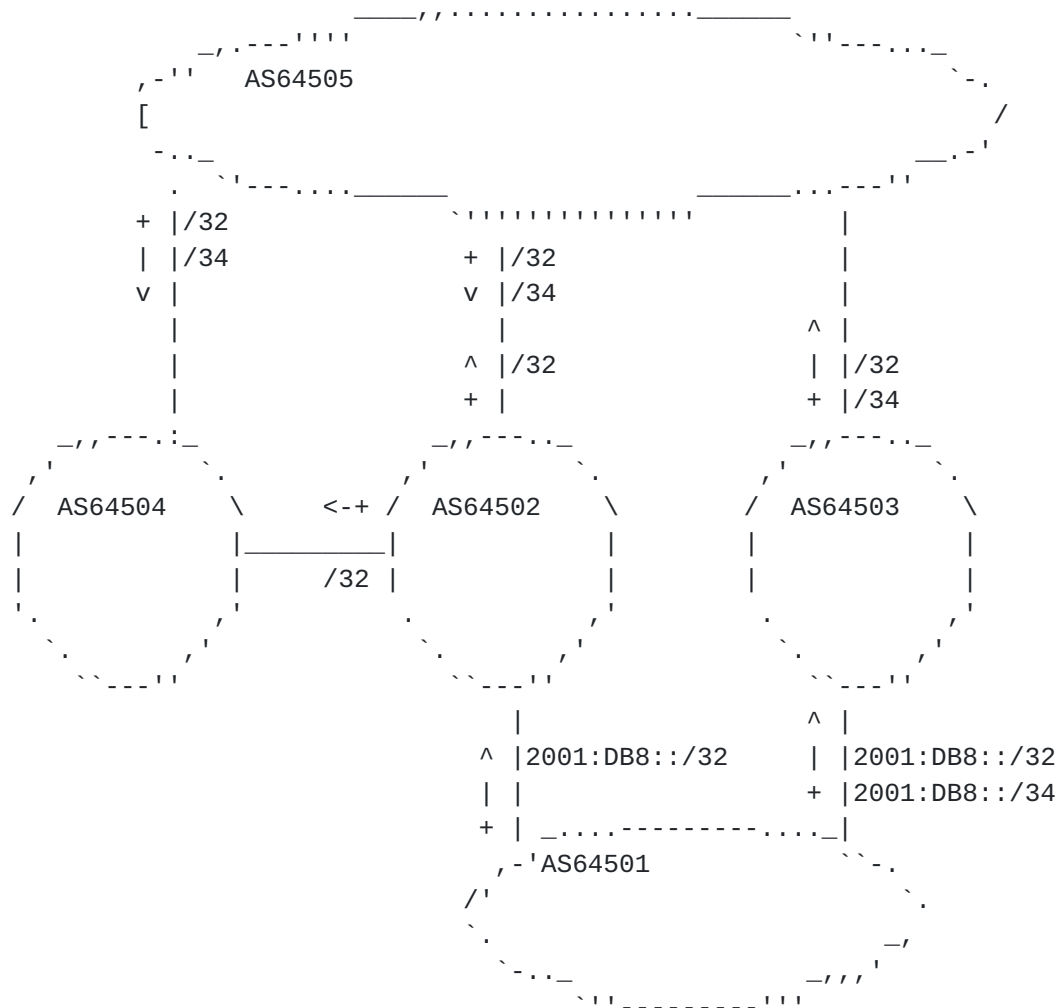


Figure 5: Counter-measures for unexpected traffic flows - Base example

4.1. Reactive counter-measures

An operator who detects unexpected traffic flows originated by any of the cases described in [Section 2](#) can contact the ASes that are likely to have performed the propagation tweaks, inform them of the situation, and persuade them to change their behavior.

If the situation remains, the operator can implement prefix filtering in order to stop the unexpected flows. The operator can decide to perform this action over the session with the operator announcing the overlapping prefix or over the session with the neighboring AS from which it is receiving the traffic. Each of these options carry a different repercussion for the affected AS. We describe briefly the two alternatives.

- o An operator can decide to stop announcing the covering prefix at the peering session with the neighboring AS from which it is receiving traffic to the overlapping prefix. In the example of Figure 5, AS64502 would filter out the prefix 2001:DB8::/32 from the eBGP session with AS64504. In this case, not all traffic heading to the prefix 2001:DB8::/32 from AS64501 would no longer traverse AS64502. AS64502 should evaluate if solving the issues originated by the unexpected traffic flows are worth the loss of this traffic share.
- o An operator can decide to filter out the overlapping prefix at the peering session over which it was received. In the example of Figure 5, AS64502 would filter out the incoming prefix 2001:DB8::/34 from the eBGP session with AS64505. As a result, the traffic destined to that /32 would be forwarded by AS64502 along its link with AS64501, despite the actions performed by AS64501 to have this traffic coming in through its link with AS64503. However, as AS64502 will no longer know a route to the overlapping prefix, it risks losing the traffic share from customers different from AS64501 to that prefix. Furthermore, this action can generate conflicts between AS64502 and AS64501, since AS64502 does not follow the routing information expressed by AS64501 in its BGP announcements.

It is possible that the behavior of the neighboring AS causing the unexpected traffic flows opposes the peering agreement. In this case, an operator could account the amount of traffic that has been subject to the unexpected flows, using traffic measurement protocols such as IPFIX, and charge the peer for that traffic. That is, the operator can claim that it has been a provider of that peer for the traffic that transited between the two ASes.

4.2. Anticipant counter-measures

4.2.1. Access lists

An operator can configure its routers to install dynamically an access-list made of the prefixes towards which the forwarding of traffic from that interface would lead to unexpected traffic flows. In the example of Figure 5, AS64502 would install an access-list denying packets matching 2001:DB8::/34 associated with the interface connecting to AS64504. As a result, traffic destined to that prefix would be dropped, despite the existence of a valid route towards 2001:DB8::/32.

This technique actually lets packets destined to a valid prefix be dropped while they are sent from a neighboring AS that may not know about the policy conflict and hence had no means to avoid the

creation of unexpected traffic flows. For this reason, this technique can be considered harmful and is thus not recommended for implementation.

4.2.2. Neighbor-specific forwarding

An operator can technically ensure that traffic destined to a given prefix will be forwarded from an entry point of the network based only on the set of paths that have been advertised over that entry point.

As an example, let us analyze the scenario of Figure 5 from the point of view of AS64502. The edge router connecting to the AS64504 forwards packets destined to prefix 2001:DB8::/34 towards AS64505. Likewise, it forwards packets destined to prefix 2001:DB8::/32 towards AS64501. The router, however, only propagates the path of the covering prefix (2001:DB8::/32) to AS64504. An operator could implement the necessary techniques to force the edge router to forward packets coming from AS64504 based only on the paths propagated to AS64504. Thus, the edge router would forward packets destined to 2001:DB8::/34 towards AS64501 in which case no unexpected traffic flow would occur.

Different techniques could provide this functionality; however, their technical implementation can be complex to design and operate. An operator could, for instance, employ Virtual Routing Forwarding (VRF) tables [RFC4364](#) [4] to store the routes announced to a neighbor and forward traffic exclusively based on those routes. [2] describes this solution and provides a description of its limitations. In the future, new network protocols and architectures could provide this functionality with less overhead for management and device resources.

Note that similarly to the solution described in [Section 4.1](#), this approach could create conflicts between AS64502 and AS64501, since the traffic forwarding performed by AS64502 goes against the policy of AS64501.

5. Conclusions

In this document, we described how the filtering of overlapping prefixes can potentially create unexpected traffic flows in remote ASes. We provided examples of scenarios in which unexpected traffic flows are caused by these practices and introduce some techniques for their detection and prevention. Analyzing the different options for dealing with this kind of problems, we recommend ASes affected by unexpected traffic flows to implement monitoring systems that can detect them and react to them according to the specific situation. Although we observe that there are reasonable situations in which

ASes could filter overlapping prefixes, we encourage network operators to implement this type of filters considering the cases described in this document.

6. Security Considerations

It is possible for an AS to use any of the methods described in this document to deliberately reroute traffic flowing through another AS. The objective of this document is to inform on this potential routing security issue.

7. IANA Considerations

This document has no IANA actions.

8. Acknowledgments

The authors would like to thank Wes George, Jon Mitchell, and Bruno Decraene for their useful suggestions and comments.

9. References

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- [2] Vanbever, L., Francois, P., Bonaventure, O., and J. Rexford, "Customized BGP Route Selection Using BGP/MPLS VPNs", Cisco Systems, Routing Symposium <http://www.cs.princeton.edu/~jrex/talks/cisconag09.pdf>, October 2009.
- [3] "INIT7-RIPE63", <<http://ripe63.ripe.net/presentations/48-How-more-specifics-increase-your-transit-bill-v0.2.pdf>>.
- [4] "pmacct project: IP accounting iconoclasm", <<http://www.pmacct.net>>.

9.2. URIs

- [1] <http://www.ietf.org/rfc/rfc1812.txt>
- [2] <http://www.ietf.org/rfc/rfc4384.txt>
- [3] <http://www.ietf.org/rfc/rfc7011.txt>

[4] <http://www.ietf.org/rfc/rfc4364.txt>

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