

GROW Working Group  
Internet-Draft  
Intended status: Standards Track  
Expires: November 4, 2013

D. McPherson  
Verisign, Inc.  
S. Amante  
Level 3 Communications  
E. Osterweil  
Verisign, Inc.  
L. Blunk  
Merit Network, Inc.  
D. Mitchell  
Twitter, Inc.  
May 3, 2013

IRR & Routing Policy Configuration Considerations  
<[draft-ietf-grow-irr-routing-policy-considerations-01](#)>

## Abstract

The purpose of this document is to catalog past issues influencing the efficacy of Internet Routing Registries (IRR) for inter-domain routing policy specification and application in the global routing system over the past two decades. Additionally, it provides a discussion regarding which of these issues are still problematic in practice, and which are simply artifacts that are no longer applicable but continue to stifle inter-provider policy-based filtering adoption and IRR utility to this day.

## Status of this Memo

This Internet-Draft is submitted in full conformance with the provisions of [BCP 78](#) and [BCP 79](#).

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at <http://datatracker.ietf.org/drafts/current/>.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

This Internet-Draft will expire on November 4, 2013.

## Copyright Notice

Copyright (c) 2013 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to [BCP 78](#) and the IETF Trust's Legal Provisions Relating to IETF Documents (<http://trustee.ietf.org/license-info>) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Simplified BSD License text as described in Section 4.e of the Trust Legal Provisions and are provided without warranty as described in the Simplified BSD License.

## Table of Contents

<a href="#">1.</a>	Introduction . . . . .	<a href="#">3</a>
<a href="#">2.</a>	Background . . . . .	<a href="#">3</a>
<a href="#">3.</a>	Historical Artifacts Influencing IRR Efficacy . . . . .	<a href="#">3</a>
<a href="#">4.</a>	Accuracy and Integrity of Data Contained within the IRR . . . . .	<a href="#">4</a>
<a href="#">4.1.</a>	Lack of Resource Certification . . . . .	<a href="#">4</a>
<a href="#">4.2.</a>	Incentives to Maintain Data within the IRR . . . . .	<a href="#">5</a>
4.3.	Inability for Third Parties to Remove (Stale) Information from the IRRs . . . . .	<a href="#">6</a>
<a href="#">4.4.</a>	Lack of Authoritative IRR for Resources . . . . .	<a href="#">6</a>
<a href="#">4.5.</a>	Conclusions with respect to Data in the IRR . . . . .	<a href="#">7</a>
<a href="#">5.</a>	Operation of the IRR Infrastructure . . . . .	<a href="#">8</a>
<a href="#">5.1.</a>	Replication of Resources Among IRRs . . . . .	<a href="#">8</a>
<a href="#">5.2.</a>	Updating Routing Policies from Updated IRR Resources . . . . .	<a href="#">9</a>
<a href="#">6.</a>	Historical BGP Protocol Limitations . . . . .	<a href="#">11</a>
<a href="#">7.</a>	Historical Limitations of Routers . . . . .	<a href="#">12</a>
<a href="#">7.1.</a>	Incremental Updates to Policy on Routers . . . . .	<a href="#">12</a>
<a href="#">7.2.</a>	Storage Requirements for Policy on Routers . . . . .	<a href="#">12</a>
<a href="#">7.3.</a>	Updating Configuration on Routers . . . . .	<a href="#">13</a>
<a href="#">8.</a>	Security Considerations . . . . .	<a href="#">14</a>
<a href="#">9.</a>	IANA Considerations . . . . .	<a href="#">14</a>
<a href="#">10.</a>	Acknowledgements . . . . .	<a href="#">14</a>

[11. Informative References](#) . . . . . [14](#)  
[Authors' Addresses](#) . . . . . [15](#)

[1. Introduction](#)

The purpose of this document is to catalog past issues influencing the efficacy of Internet Routing Registries (IRR) for inter-domain routing policy specification and application in the global routing system over the past two decades. Additionally, it provides a discussion regarding which of these issues still pose problems in practice, and which are no longer obstacles, but whose perceived drawbacks continue to stifle inter-provider policy-based filtering support and IRR utility to this day.

[2. Background](#)

IRRs can be used to express a multitude of Internet number bindings and policy objectives, to include bindings between an origin AS and a given prefix, AS and community import and export policies for a given AS, as well as AS macros (as-sets in RPSL-speak) that convey the set of ASes for which a given AS intends to include in some common group.

As quoted from [Section 7 of \[RFC1787\]](#), Routing in a Multi-Provider Internet:

While ensuring Internet-wide coordination may be more and more difficult, as the Internet continues to grow, stability and consistency of the Internet-wide routing could significantly benefit if the information about routing requirements of various organizations could be shared across organizational boundaries. Such information could be used in a wide variety of situations ranging from troubleshooting to detecting and eliminating conflicting routing requirements. The scale of the Internet implies that the information should be distributed. Work is currently underway to establish depositories of this information (Routing Registries), as well as to develop tools that analyze, as well as utilize this information.

### 3. Historical Artifacts Influencing IRR Efficacy

The term IRR is often used, incorrectly, as a broad catch-all term to categorize issues related to the accuracy of data in the IRR, the Routing Policy Specification Language (RPSL) and the operational deployment of policy (derived from RPSL contained within the IRR) to routers. It is important to classify these issues into distinct categories so that the reader can understand which categories of issues are historical artifacts that are no longer applicable and which categories of issues still exist and might be addressed by the IETF.

McPherson, et al.

Expires November 4, 2013

[Page 3]

---

Internet-Draft

IRR & Routing Policy Considerations

May 2013

The following sections will separate out challenges related to the IRR into the following categories. First, accuracy and integrity of data contained within the IRR. Second, the Resource Policy Specification Language (RPSL) used to represent various types of data in the IRR. Third, operation of the IRR infrastructure, i.e.: synchronization of resources from one IRR to other IRRs. Finally, the methods related to extraction of policy from the IRR and the input plus activation of that policy within routers.

### 4. Accuracy and Integrity of Data Contained within the IRR

The following section will examine issues related to accuracy and integrity of data contained within the IRR.

#### 4.1. Lack of Resource Certification

One of the largest weaknesses often cited with the IRR system is that the data contained within the IRRs is out of date or lacks integrity. This is largely attributable to the fact that historically there has existed no method for developing cryptographically verifiable bindings of an IP prefix to the originating Autonomous System (AS). That is, there has never existed a resource certification infrastructure that enables a resource holder to authorize a particular autonomous system to originate network layer reachability advertisements for a given IPv4 or IPv6 prefix. It should be noted that this is not a weakness of the underlying Routing Policy Specification Language (RPSL) [[RFC2622](#)], but rather, was largely the result of no clear demand by the operator community for Internet

Number Resource Registries to provide sufficient resource certification infrastructure that would enable a resource holder to develop a cryptographic binding between, for example, a given AS number and an IP prefix.

Another noteworthy (but slightly different) deficiency in the IRR system is the absence of a tangible tie between the resource and the resource holder. If a resource holder's authorization cannot be certified, then consumers cannot verify attestations made. In effect, without resource certification, consumers are basically only certifying the assertions that the creator / maintainer of the resource object has made (not if that assertion is valid).

The RIPE community addressed this last issue by putting in a foundation policy [[RIPE2007-01](#)], which requires a contractual link between the RIPE NCC and the end user in direct assignment + ASN assignment cases, which weren't previously covered by LIR contracts for address allocations. There were a couple of intentions with this policy:

1. There was an encumbrance placed in the policy for the LIR to charge the end-user for provider independent resources. This action created a collection mechanism for PI address resources (IPv4 / IPv6 space, ASNs).
2. It guaranteed that all RIPE NCC allocated/assigned space would be subject to a contractual link, and that this contractual chain might end up actually meaning something when it came to the issue of who made what claim about what number resource.
3. It tied into the RIPE NCC's object grandfathering policy which ties the registration details of the end-user to the object registered in the IRR database.

One of the central observations of this policy was that without a chain-of-ownership functionality in IRR databases, the discussion of certifying their contents becomes moot.

#### [4.2.](#) Incentives to Maintain Data within the IRR

A second problem with data contained in the IRRs is that the incentives for resource holders to maintain both accurate and up-to-

date information in one, or more, IRRs; including deletion of out-of-date or stale data from the IRRs can diminish rapidly when changing their peering policies (such as switching transit providers). Specifically, there is a very strong incentive for an ISP's customers to register new routing information in the IRR, because some ISPs enforce a strict policy that they will only build or update a customer's prefix-lists applied to the customer's inbound eBGP sessions based off information found within the IRRs. Unfortunately, there is little incentive for an ISP's customers to remove out-of-date information from an IRR, most likely attributed to the fact that some ISPs do not use, or enforce use of, data contained within the IRRs to automatically build incoming policy applied to customer's eBGP sessions. For example, if a customer is terminating service from one ISP that requires use of IRR data to build incoming policy and, at the same time, enabling service with another ISP that does not require use of IRR data, then that customer will likely let the data in the IRR become stale or inaccurate.

Further, policy filters are almost exclusively generated based on the origin AS information contained within IRR route objects and used by providers to filter downstream transit customers. Since providers only look for route objects containing the origin AS of their downstream customer(s), stale route objects with historical and, possibly, incorrect origin AS information are ignored. Assuming that the downstream customer(s) do not continue to announce those routes with historical, or incorrect, origin AS information in BGP to the

upstream provider, there is no significant incentive to remove them as they do not impact offline policy filter generation nor routing on the Internet. On the other hand, the main incentive that causes the Service Provider to work with downstream customer(s) is when the resultant filter list becomes too large that it is difficult for it to be stored on-board PE routers; however, this is more practically an operational issue with very old, legacy PE routers, not more modern PE router hardware with more robust Control Plane Engines.

#### 4.3. Inability for Third Parties to Remove (Stale) Information from the IRRs

A third challenge with data contained in IRRs is that it is not possible for IRR operators, and ISPs who use them, to proactively remove (perceived) out-of-date or "stale" resources in an IRR, on

behalf of resource holders who may not be diligent in maintaining this information themselves. The reason is that, according to the Resource Policy Specification Language (RPSL) [[RFC2622](#)], only the resource holder ('mntner') specified in a 'mnt-by' value field of an RPSL resource is authorized to add, modify or delete their own resources within the IRR. To address this issue, the 'auth-override' mechanism [[RFC2725](#)] was later developed that would have enabled a third party to update and/or remove "stale" resources from the IRR. While it is unclear if this was ever implemented or deployed, it does provide language semantics needed to overcome this obstacle.

Nevertheless, with such a mechanism in place, there is still a risk that the original RPSL resource holder would not receive notifications (via the 'notify' attribute in various RPSL resources) about the pending or actual removal of a resource from the IRR in time to halt that action if those resources were still being used. In this case, if the removal of a resource was not suspended, it could potentially result in an unintentional denial of service for the RPSL resource holder when, for example, ISPs automatically generated and deployed a new policy based on the newly removed resources from the IRR, thus dropping any reachability announcement for the removed resource in eBGP.

#### [4.4.](#) Lack of Authoritative IRR for Resources

According to [[RFC2622](#)], within an RPSL resource "the source attribute specifies the registry where the object is registered." Note that this source attribute only exists within the RPSL resource itself. Unfortunately, given a specific resource, (e.g.: a specific IPv4 or IPv6 prefix), most of the time it is impossible to tell a priori the authoritative IRR where to query and retrieve an authoritative copy of that resource.

This makes it difficult for consumers of data from the IRR to automatically know the authoritative IRR of a resource holder, which will contain their most up-to-date set of resources. This is typically not a problem for an ISP that has a direct (customer) relationship with the resource holder, because the ISP will ask the resource holder which (authoritative) IRR to pull their resources from on, for example, a "Customer BGP Order Form". However, third parties that do not have a direct relationship with the resource

holder, have a difficult time attempting to infer the authoritative IRR, preferred by the resource holder, that likely contains the most up-to-date set of resources. As a result, it would be helpful for third parties if there was a robust referral mechanism so that a query to one IRR would be automatically redirected toward the authoritative IRR for the most up-to-date and authoritative copy of that particular resource. This problem is worked around by individual IRR operators storing a local copy of other IRR's resources, through periodic mirroring, which allows the individual IRR to respond to a client's query with all registered instances of a particular IRR resource that exist in both the local IRR and all other IRRs. Of course, the problem with this approach is that an individual IRR operator is under no obligation to mirror all other IRRs and, in practice, some IRRs do not mirror the resources from all other IRRs. This could lead to the false impression that a particular resource is not registered or maintained at a particular IRR. Furthermore, the authentication process of accepting updates by any given IRR may, or may not, be robust enough to overcome impersonation attacks. As a result, there is no rigorous assurance that a mirrored RPSL statement was actually made by the authorized resource holder.

#### [4.5.](#) Conclusions with respect to Data in the IRR

All of the aforementioned issues related to integrity and accuracy of data within the IRR stem from a distinct lack of resource certification for resources contained within the IRR. Only now is an experimental test bed that reports to provide this function (under the auspices of the Resource PKI [[I-D.ietf-sidr-arch](#)]) being formally discussed, which could also aid in certification of resources within the IRR. It should be noted that the RPKI is only currently able to support signing of Route Origin Authorization (ROA) resources that are the equivalent of 'route' resources in the IRR. It is unclear if, in the future, the RPKI will be extended to support additional resources that already exist in the IRR, e.g.: aut-num, as-net, route-set, etc. Finally, a seemingly equivalent resource certification specification for all resources in the IRR has already been developed [[RFC2725](#)], however it is unclear how widely it was ever implemented or deployed.

#### [5.](#) Operation of the IRR Infrastructure

## 5.1. Replication of Resources Among IRRs

Currently, several IRRs [IRR\_LIST] use a Near-Real-Time Mirroring (NRTM) protocol to replicate each others contents. However, this protocol has a several weaknesses. Namely, there is no way to validate that the copy of mirror is correct and synchronization issues have often resulted. Furthermore, the NRTM protocol does not employ any security mechanisms. The NRTM protocol relies on a pull mechanism and is generally is configured with a poll interval of 5 to 10 minutes. There is currently no mechanism to notify an IRR when an update has occurred in a mirrored IRR so that an immediate update can be made.

Some providers employ a process of mirroring an instance of an IRR that involves downloading a flat text file copy of the IRR that is made available via FTP. These FTP files are exported at regular intervals of typically anywhere between 2 and 24 hours by the IRRs. When a provider fetches those text files, it will process them to identify any updates and reflect changes within its own internally maintained database. The use of an internally maintained database is out-of-scope of this document, but is generally used to assist with more straightforward access to or modification of data by the IRR operator. Providers typically employ a 24 hour cycle to pull updated resources from IRRs. Thus, depending on when resource holders submitted their changes to an IRR, it may take up to 24 hours for those changes to be reflected in their policy configurations. In practice, it appears that the RPKI will also employ a 24 hour cycle whereby changes in resources are pushed out to other RPKI caches [REF?].

IRRs originated from [Section 7 of \[RFC1787\]](#), specifically: "The scale of the Internet implies that the [routing requirements] information should be distributed." Regardless, the practical effect of an organization maintaining its own local cache of IRR resources is an increase in resource resiliency (due to multiple copies of the same resource being geographically distributed), a reduction in query time for resources and, likely, a reduction in bandwidth consumption and associated costs. This is particularly beneficial when, for example, an ISP is evaluating resources in an IRR just to determine if there are any modifications to those resources that will ultimately be reflected in a new routing policy that will get propagated to (Edge) routers in the ISP's network. Cache locality results in reduced bandwidth utilization for each round trip.

On the other hand, it is unclear from where the current provider replication interval of 24 hours originated or even whether it still

---

provides enough freshness in the face of available resources, particularly in today's environment where a typical IRR system consists of a (multi-core) multi-Ghz CPU connected to a network via a physical connection of 100 Mbps or, more likely, higher bandwidth. In addition, it can be observed that the most common circuit size used by ISPs is now 10 Gbps, thus eliminating bandwidth as a significant factor in the transfer of data between IRRs. Furthermore, it should be noted that Merit's Internet Routing Registry Daemon (IRRd) [[MERIT-IRRd](#)], uses 10 minutes as its default for "irr\_mirror\_interval".

Lastly, it should be noted that [[RFC2769](#)], Routing Policy System Replication, attempted to offer a more methodical solution for distributed replication of resources between IRRs. It is unclear why this RFC failed to gain traction, but it is suspected that this was due to that specification's reliance on [[RFC2725](#)], Routing Policy System Security, that addressed "the need to assure integrity of the data by providing an authentication and authorization model."

## [5.2](#). Updating Routing Policies from Updated IRR Resources

Ultimately, the length of time it takes to replicate resources among IRRs is, generally, the dominant factor in reflecting changes to resources in policy that is eventually applied within the Control Plane of routers. The length of time to update network elements will vary considerably depending on the size of the ISP and the number of IRR resources that were updated during any given interval. However, there are a variety of common techniques, that are outside the scope of this document, that allow for this automated process to be optimized to considerably reduce the length of time it takes to update policies in the ISP's network.

An ISP will begin the process of updating the policy in their network, by first fetching IRR resources associated with, for example, a customer ASN attached to their network. Next, the ISP constructs a new policy associated to that customer and then evaluates if that new policy is different from existing policy associated with that same customer. If there are no changes between the new and existing policy associated with that customer, then the ISP does not make any changes to the policy in their routers specific to that customer. On the other hand, if the new policy does reflect changes from the existing policy for that customer, then the ISP begins a process of uploading the new policy to the routers attached to that customer.

The process of constructing a new policy involves use of a set of

programs, e.g.: IRRtoolset, that performs recursive expansion of an RPSL aut-num resource, that comprises an arbitrary combination of

other RPSL aut-num, as-set, route and route-set resources, according to procedures defined by RPSL. The end result of this process is, traditionally, a router vendor dependent configuration snippet that defines the routing policy for that customer. This routing policy may consist of the set of IPv4 or IPv6 prefixes, associated prefix lengths and AS\_PATH's that are supposed to be accepted from a customer's eBGP session. However, if indicated in the appropriate RPSL resource, the policy may also set certain BGP Attributes, such as MED, AS\_PATH prepend value, LOCAL\_PREF, etc., at either the incoming eBGP session from the customer or on static routes that are originated by the resource holder.

An ISP's customers may not adequately plan for pre-planned maintenance or, more likely, need to rapidly begin announcing a new IP prefix as a result of, for example, an emergency turn-up of the ISP customer's new downstream customer. Unfortunately, the routine, automated process employed by the ISP means that they may not begin updating their routing policy on their network for up to 24 hours, because the ISP or the IRRs the ISP uses might only mirror changes to IRR resources once every 24 hours. The time interval for the routine/automated process is not responsive to the needs of directly paying customer(s) who need rapid changes in their policy in rare situations. In these situations, when a customer has an urgent need for updates to take affect immediately, they will call up the Network Operations Center (NOC) of their ISP and request that the ISP immediately fetch new IRR objects and push those changes out to their network. This is often accomplished in as little as five minutes from the time a customer contacts their ISP's NOC to the time a new filtering policy is pushed out to the PE routers that are attached to that customer's Attachment Circuits (AC's). It is conceivable that some ISPs automate this using out-of-band mechanisms as well, although the authors are unaware of any existing mechanisms that support this.

Ultimately, the aforementioned latency with respect to "emergency changes" in IRR resources that need to be reflected in near-real-time in the network is compounded if the IRR resources were being used by third party ISPs to perform filtering on their peering circuits, where typically no such policies are employed today for this very

reason. It is likely that the length of time at which IRRs mirror changes will have to be dramatically reduced with a corresponding reduction in time at the ISPs that, in turn, need to evaluate whether changes in a set of IRR resources need to get reflected in updated router policies that will get pushed out to ASBR's, connected to peering circuits, on their network.

## [6.](#) Historical BGP Protocol Limitations

As mentioned previously, after a resource holder made changes to their resources in an IRR, those changes would automatically be distributed to other IRRs, ISPs would then learn of those changes, generate a new BGP policies, and then apply those to the appropriate subset of routers in their ASes. However, in the past, one additional step is necessary in order to have any of those new BGP policies take effect in the Control Plane and to allow/deny the updated resource from a customer to their ISP and from their immediately upstream ISP to the ISP's peers. It was necessary (often manually) to actually induce BGP on each router to apply the new policy within the Control Plane, thus learning of a newly announced/changed IP prefix (or, dropping a deleted IP prefix). Unfortunately, most of these methods were not only highly operationally impactful, but also affected traffic forwarding to IP destinations that were unrelated to the most recent changes to the BGP policy.

Historically, a customer would have to (re-)announce the new IP prefix toward their ISP, but only after the ISP had placed the new BGP policies into affect. Alternatively, the ISP would have to reset the entire PE-CE eBGP session either by: a) bouncing the entire interface toward the customer, (e.g.: shutdown/no shutdown); or, b) clearing the eBGP session toward the customer, (e.g.: clear ip bgp neighbor >IP address of CE router<). The latter two cases were, of course, the most highly impactful and thus could generally only be performed off-hours during a maintenance window.

Once the new IP prefix has been successfully announced by the customer and permitted by the newly updated policy at the ISP's PE's (attached to that customer), it would be propagated to that ISP's ASBR's, attached to peers, at the perimeter of the ISP network.

Unfortunately, if those peers had either not yet learned of the changes to resources in the IRR or pushed out new BGP policies (and, reset their BGP sessions immediately afterward) incorporating those changes, then the newly announced route would also get dropped at the ingress ASBR's of the peers.

Ultimately, either of the two scenarios above further lengthen the effective time it would take for changes in IRR resources to take affect within BGP in the network. Fortunately, BGP has been enhanced over the last several years in order that changes within BGP policy will take affect without requiring a service impacting reset of BGP sessions. Specifically, BGP soft-reconfiguration [REF?] and, later, Route Refresh Capability for BGP-4 [[RFC2918](#)] were developed so that ISPs, or their customers, could induce BGP to apply a new policy while leaving both the existing eBGP session active as well as (unaffected) routes active in both the Loc-RIB and, more importantly,

FIB of the router. Thus, using either of these mechanisms, an ISP or its peers currently will deploy a newly generated BGP policy, based on changes to resources within the IRR, and immediately trigger a non-service impacting notification to the BGP process to have those changes take affect in their Control Plane, either allowing a new IP prefix to be announced or an old IP prefix to be dropped. This dramatically reduces the length of time after changes are propagated throughout the IRRs to when those changes are actually operational within BGP policy in an ISP's network.

## [7.](#) Historical Limitations of Routers

### [7.1.](#) Incremental Updates to Policy on Routers

Routers in the mid 1990's rarely supported incrementally updatable prefix filters for BGP, and therefore, if new information was received from a policy or internal configuration database that would impact a policy applied to a given eBGP peer, the entire prefix list or access list would need to be deleted and rewritten, compiled, and installed. This was very tedious and commonly led to leaked routes during the time when the policy was being rewritten, compiled, and applied on a router. Furthermore, application of a new policy would not automatically result in new ingress or egress reachability advertisements, from that new policy, because routers at the time

would require a reset of the eBGP sessions for routing information to be evaluated by the new policy. Of course, resetting of an eBGP session had implications on traffic forwarding during the time the eBGP session was reestablished and new routing information was learned.

Routers now support the ability to perform incremental, and in situ, updates to filter lists consisting of IP prefixes and/or AS\_PATH's that are used within an ingress or egress BGP policy. In addition, routers also can apply those incremental updates to policy, with no traffic disruption, using BGP soft-reconfiguration or BGP Route Refresh, as discussed in the previous section.

## 7.2. Storage Requirements for Policy on Routers

Historically, routers had very limited storage capacity and would have difficulty in storing an extremely large BGP policy on-board. This was typically the result of router hardware vendors using an extremely limited amount of NVRAM for storage of router configurations.

Another challenge with historical router hardware was that writing to NVRAM was extremely slow. Thus, when the router configuration had

changed as a result of, for example, updating a BGP policy to accommodate changes in IRR resources, this would result in extremely long times to write out these configuration changes and, sometimes, due to bugs would result in loss of protocol keep-alives, thus causing an impact to routing or forwarding of packets through the platform.

The above limitations have largely been resolved with more modern equipment from the last few years shipping with increasing amounts of non-volatile storage such as: PCMCIA or USB Flash Cards, Hard Disk Drives or Solid State Disk Drives.

## 7.3. Updating Configuration on Routers

Historically, there has not been a standardized network configuration modeling language or an associated method to update router configurations. When an ISP detected a change in resources within the IRR, it would fashion a router vendor dependent BGP policy and

upload that to the router usually via the following method.

First, an updated BGP policy configuration 'snippet' is generated via processes running on an offline server. Next, the operator uses either telnet or SSH to login to the CLI of a target router and issue router vendor dependent CLI commands that will trigger the target router to fetch the new configuration snippet via TFTP, FTP or Secure Copy (SCP) stored on the offline server. The target router will then perform syntax checking on that configuration snippet and, if that passes, merge that configuration snippet into the running configuration of the router's Control Software. At this point, the new BGP policy configuration 'snippet' is active within the Control Plane of the router. One last step remains, which is the operator will issue a CLI command to induce the router to perform a "soft reset", via BGP soft-reconfiguration or BGP Route Refresh, of the associated BGP session in order to trigger the router to apply the new policy to routes learned from that BGP session without disrupting traffic forwarding.

More recently, operators have the ability to use NETCONF/SSH (or, TLS) from an offline server to push a BGP configuration 'snippet' from an offline server toward a target router that has that capability. However, this activity is still dependent on generating, via the offline server, a router vendor dependent XML configuration snippet that would get uploaded via SSH or TLS to the target router.

In the future, the ability to upload new Route Origin Authorization (ROA) information may be provided from the RPKI to routers via the RPKI-RTR [[I-D.ietf-sidr-rpki-rtr](#)] protocol. However, this will not allow operators the ability to upload other configuration information

such as BGP policy information, such as AS\_PATH's, BGP communities, etc. that might be associated with that ROA information, like they can from IRR generated BGP policies.

## [8.](#) Security Considerations

## [9.](#) IANA Considerations

## 10. Acknowledgements

TBD.

## 11. Informative References

[I-D.ietf-sidr-arch]

Lepinski, M. and S. Kent, "An Infrastructure to Support Secure Internet Routing", [draft-ietf-sidr-arch-13](#) (work in progress), May 2011.

[I-D.ietf-sidr-rpki-rtr]

Bush, R. and R. Austein, "The RPKI/Router Protocol", [draft-ietf-sidr-rpki-rtr-20](#) (work in progress), November 2011.

[MERIT-IRRd]

Merit, "IRRd - Internet Routing Registry Daemon", <http://www.irrd.net>.

[RFC1787] Rekhter, Y., "Routing in a Multi-provider Internet", [RFC 1787](#), April 1995.

[RFC2622] Alaettinoglu, C., Villamizar, C., Gerich, E., Kessens, D., Meyer, D., Bates, T., Karrenberg, D., and M. Terpstra, "Routing Policy Specification Language (RPSL)", [RFC 2622](#), June 1999.

[RFC2725] Villamizar, C., Alaettinoglu, C., Meyer, D., and S. Murphy, "Routing Policy System Security", [RFC 2725](#), December 1999.

[RFC2769] Villamizar, C., Alaettinoglu, C., Govindan, R., and D. Meyer, "Routing Policy System Replication", [RFC 2769](#), February 2000.

[RFC2918] Chen, E., "Route Refresh Capability for BGP-4", [RFC 2918](#), September 2000.

[RIPE2007-01]

RIPE NCC, "DRAFT: Autonomous System (AS) Number Assignment Policies and Procedures", Foundation Policy <http://www.ripe.net/ripe/docs/2007-01-asn>.

#### Authors' Addresses

Danny McPherson  
Verisign, Inc.

Email: [dmcpherson@verisign.com](mailto:dmcpherson@verisign.com)

Shane Amante  
Level 3 Communications  
1025 Eldorado Blvd  
Broomfield, CO 80021  
USA

Email: [shane@level3.net](mailto:shane@level3.net)

Eric Osterweil  
Verisign, Inc.

Email: [eosterweil@verisign.com](mailto:eosterweil@verisign.com)

Larry J. Blunk  
Merit Network, Inc.

Email: [ljb@merit.edu](mailto:ljb@merit.edu)

Dave Mitchell  
Twitter, Inc.

Email: [dave@twitter.com](mailto:dave@twitter.com)