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The Accumulated IGP Metric Attribute for BGP

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Abstract

Routing protocols that have been designed to run within a single administrative domain ("IGPs") generally do so by assigning a metric to each link, and then choosing as the installed path between two nodes the path for which the total distance (sum of the metric of each link along the path) is minimized. BGP, designed to provide routing over a large number of independent administrative domains ("autonomous systems"), does not make its path selection decisions through the use of a metric. It is generally recognized that any attempt to do so would incur significant scalability problems, as well as inter-administration coordination problems. However, there are deployments in which a single administration runs several contiguous BGP networks. In such cases, it can be desirable, within that single administrative domain, for BGP to select paths based on a metric, just as an IGP would do. The purpose of this document is to provide a specification for doing so.

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[1. Specification of requirements](#)

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [[RFC2119](#)].

[2. Introduction](#)

There are many routing protocols that have been designed to run within a single administrative domain. These are known collectively as "Interior Gateway Protocols" (IGPs). Typically, each link is assigned a particular "metric" value. The path between two nodes can then be assigned a "distance", which is the sum of the metrics of all the links that belong to that path. An IGP selects the "shortest" (minimal distance) path between any two nodes, perhaps subject to the constraint that if the IGP provides multiple "areas", it may prefer

the shortest path within an area to a path that traverses more than one area. Typically the administration of the network has some routing policy which can be approximated by selecting shortest paths in this way.

BGP, as distinguished from the IGPs, was designed to run over an arbitrarily large number of administrative domains ("autonomous systems", or "ASes") with limited coordination among the various administrations. BGP does not make its path selection decisions based on a metric; there is no such thing as an "inter-AS metric". There are two fundamental reasons for this:

- The distance between two nodes in a common administrative domain may change at any time due to events occurring in that domain. These changes are not propagated around the Internet unless they actually cause the border routers of the domain to select routes with different BGP attributes for some set of address prefixes. This accords with a fundamental principle of scaling, viz., that changes with only local significance must not have global effects. If local changes in distance were always propagated around the Internet, this principle would be violated.
- A basic principle of inter-domain routing is that the different administrative domains may have their own policies, which do not have to be revealed to other domains, and which certainly do not have to be agreed to by other domains. Yet the use of inter-AS metric in the Internet would have exactly these effects.

There are, however, deployments in which a single administration runs a network which has been sub-divided into multiple, contiguous ASes, each running BGP. There are several reasons why a single administrative domain may be broken into several ASes (which, in this case, are not really "autonomous".) It may be that the existing IGPs do not scale well in the particular environment; it may be that a more generalized topology is desired than could be obtained by use of a single IGP domain; it may be that a more finely grained routing policy is desired than can be supported by an IGP. In such deployments, it can be useful to allow BGP to make its routing decisions based on the IGP metric, so that BGP chooses the "shortest" path between two nodes, even if the nodes are in two different ASes within that same administrative domain.

There are in fact some implementations which already do something like this, using the MULTI_EXIT_DISC (MED) attribute to carry the IGP metric. However, that doesn't really provide IGP-like "shortest path" routing, as the BGP decision process gives priority to other factors, such as LOCAL_PREF and AS_PATH length. Also, the standard procedures for use of the MED do not ensure that the IGP metric is

properly accumulated so that it covers all the links along the path.

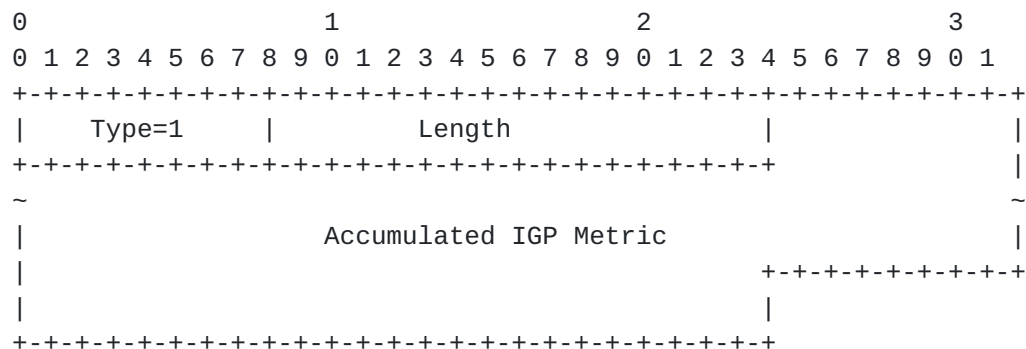
In this document, we define a new optional, non-transitive BGP attribute, called the "Accumulated IGP Metric Attribute", or "AIGP attribute", and specify the procedures for using it.

The specified procedures prevent the AIGP attribute from "leaking out" past the administrative domain boundaries into the Internet.

The specified procedures also ensure that the value in the AIGP attribute has been accumulated all along the path from the destination, i.e., that the AIGP attribute does not appear when there are "gaps" along the path where the IGP metric is unknown.

3. AIGP Attribute

The AIGP Attribute is an optional non-transitive BGP Path Attribute. The attribute type code for the AIGP Attribute is to be assigned by IANA. The value field of the AIGP Attribute is defined here to be a set of TLVs (elements encoded as "Type/Length/Value"). However, this document defines only a single such TLV, the AIGP TLV, that contains the Accumulated IGP Metric. The AIGP TLV is encoded as shown in Figure 1. An AIGP Attribute MUST NOT contain more than one AIGP TLV.



AIGP Attribute
Figure 1

- Type: A single octet encoding the AIGP Attribute Type. Only type 1 is defined in this document.
- Length: Two octets encoding the length in octets of the attribute, including the type and length fields. The length is encoded as an unsigned binary integer.

The length of the AIGP TLV is always 11.

- Accumulated IGP Metric: For a type 1 AIGP attribute, the value field is always 8 bytes long. IGP metrics are frequently expressed as 4-octet values, and this ensures that the AIGP attribute can be used to hold the sum of an arbitrary number of 4-octet values.

[3.1.](#) Applicability Restrictions and Cautions

This document only considers the use of the AIGP attribute in networks where each router uses tunneling of some sort to deliver a packet to its BGP next hop. Use of the AIGP attribute in networks that do not use tunneling is outside the scope of this document.

If a Route Reflector supports the AIGP attribute, but some of its clients do not, then the routing choices that result may not all reflect the intended routing policy.

[3.2.](#) Restrictions on Sending/Receiving

An implementation that supports the AIGP attribute MUST support a per-session configuration item, AIGP_SESSION, that indicates whether the attribute is enabled or disabled for use on that session.

- The default value of AIGP_SESSION, for EBGp sessions, MUST be "disabled".
- The default value of AIGP_SESSION, for IBGP and confederation-EBGP sessions, MUST be "enabled."

The AIGP attribute MUST NOT be sent on any BGP session for which AIGP_SESSION is disabled.

If an AIGP attribute is received on a BGP session for which AIGP_SESSION is disabled, the attribute MUST be treated exactly as if it were an unrecognized non-transitive attribute. That is, "it MUST be quietly ignored and not passed along to other BGP peers" (see [BGP], [section 5](#)).

3.3. Creating and Modifying the AIGP Attribute

3.3.1. Originating the AIGP Attribute

A BGP speaker MUST NOT add the AIGP attribute to any route whose path leads outside the AS to which the BGP speaker belongs. It may be added only to routes that satisfy one of the following conditions:

- The route is a static route that is being redistributed into BGP
- The route is an IGP route that is being redistributed into BGP
- The route is an IBGP-learned route whose AS_PATH attribute is empty.

An implementation that supports the AIGP attribute MUST support a configuration item, AIGP_ORIGINATE, that enables or disables its creation and attachment to routes. The default value of AIGP_ORIGINATE MUST be "disabled".

It SHOULD be possible to set AIGP_ORIGINATE to "enabled for the routes of a particular IGP that are redistributed into BGP" (where "a particular IGP" might be "OSPF" or "ISIS").

When a BGP speaker R learns a route to address prefix P from an IGP, the IGP will have computed a "distance" from R to P. The value assigned to the AIGP attribute is either the IGP-computed distance, or some other value determined by policy.

In the case of a static route whose next hop matches a BGP route that has an AIGP attribute, the static route MAY inherit the AIGP attribute value of that BGP route.

3.3.2. Modifications by the Originator

If BGP speaker R is the originator of the AIGP attribute of prefix P, and at some point the distance from R to P changes, R SHOULD issue a new BGP update containing the new value of the AIGP attribute. However, if the difference between the new distance and the distance advertised in the AIGP attribute is less than a configurable threshold, the update MAY be suppressed.

3.3.3. Modifications by a Non-Originator

Suppose a BGP speaker R1 receives a route with an AIGP attribute whose value is A, and a Next Hop whose value is R2. Suppose also that R1 is about to redistribute that route on a BGP session that is enabled for sending/receiving the attribute.

If R1 does not change the Next Hop of the route, then R1 MUST NOT change the AIGP attribute value of the route.

If R1 changes the Next Hop of the route from R2 to R1, and if R1's route to R2 is an IGP-learned route, or a static route that does not require recursive next hop resolution, then R1 must increase the value of the AIGP attribute by adding to A the distance from R1 to R2. This distance is either the IGP-computed distance from R1 to R2, or some value determined by policy. However, A MUST be increased by a non-zero amount.

Note that if R1 and R2 above are EBGp neighbors, and there is a direct link between them on which no IGP is running, then when R1 changes the next hop of a route from R2 to R1, the AIGP metric value MUST be increased by a non-zero amount. The amount of the increase SHOULD be such that it is properly comparable to the IGP metrics. E.g., if the IGP metric is a function of latency, then the amount of the increase should be a function of the latency from R1 to R2.

If R1 changes the Next Hop of the route from R2 to R1, and if R1's route to R2 is a BGP-learned route, or a static route that requires recursive next hop resolution, then the AIGP attribute value needs to be increased in several steps:

1. Let Xattr be the new AIGP attribute value.
2. Initialize Xattr to A.
3. Set the XNH to R2.
4. Find the route to XNH.
5. If the route to XNH does not require recursive next hop resolution, get the distance D from R1 to XNH. If D is above a configurable threshold, set the AIGP attribute value to Xattr+D. If D is below a configurable threshold, set the AIGP attribute value to Xattr. In either case, exit this procedure.

6. If the route to XNH is a BGP-learned route, and the route does NOT have an AIGP attribute, then exit this procedure and do not pass on any AIGP attribute.
7. If the route to XNH is a BGP-learned route, and the route has an AIGP attribute value of Y, then set $Xattr = Xattr + Y$, and set XNH to the next hop of this route. (The intention here is that Y is the AIGP value of the route as it was received by R1, without having been modified by R1.)
8. Go to step 4.

The AIGP value of a given route depends on (a) the AIGP values of all the next hops that are recursively resolved during this procedure, and (b) the IGP distance to any next hop that is not recursively resolved. Any change due to (a) in any of these values MUST trigger a new AIGP computation for that route. Whether a change due to (b) triggers a new AIGP computation depends upon whether the change in IGP distance exceeds a configurable threshold.

Note that the overall shortest path may not be selected if the next hop has to be recursively resolved more than once.

If the AIGP attribute is carried across several ASes, each with its own IGP domain, it is clear that these procedures are unlikely to give a sensible result if the IGPs are different (e.g., some OSPF and some IS-IS), or if the meaning of the metrics is different in the different IGPs (e.g., if the metric represents bandwidth in some IGP domains but represents latency in others). These procedures also are unlikely to give a sensible result if the metric assigned to inter-AS BGP links (on which no IGP is running) or to static routes is not comparable to the IGP metrics. All such cases are outside the scope of the current document.

4. Decision Process

4.1. When a Route has an AIGP Attribute

Use of the AIGP attribute involves several modifications to the BGP "phase 2" decision process as described in [BGP], [section 9.1.2.2](#). The procedures defined in this section MUST be executed BEFORE any of the tie breaking procedures described therein are executed.

If any routes have an AIGP attribute, remove from consideration all routes that do not have an AIGP attribute.

If router R is considering route T, where T has an AIGP attribute,

- then R must compute the value A, defined as follows: set A to the sum of (a) T's AIGP attribute value and (b) the IGP distance from R to T's next hop.
- remove from consideration all routes that are not tied for the lowest value of A.

4.2. When the Route to the Next Hop has an AIGP attribute

Suppose that a given router R1 is comparing two routes, neither of which has an AIGP attribute. The BGP decision process as specified in [BGP] makes use, in its tie breaker procedures, of "interior cost", defined as follows:

"interior cost of a route is determined by calculating the metric to the NEXT_HOP for the route using the Routing Table."

Suppose route T has a next hop of N. We modify the notion of the "interior cost" from node R to node N as follows:

- If the route to N has an AIGP attribute, set A to the AIGP value of the route to N, computing the AIGP value of the route according to the procedure of [section 3.3.3](#). (This will have been computed at the time the route to N was installed.)
- If the route to N does not have an AIGP value, set A to 0. (This can only be the case if there is no route to N that does have an AIGP value.)
- Let R2 be the next hop of the route to N, after all recursive resolution of the next hop is done. Let m be the IGP distance (or in the case of a static route, the configured distance) from R1 to R2.
- The "interior cost" of route T is the quantity A+m.

5. Deployment Considerations

Using the AIGP attribute to achieve a desired routing policy will be more effective if each BGP speaker can use it to choose from among multiple routes. Thus is it highly recommended that the procedures of [BESTEXT] and [[ADDPATH](#)] be used in conjunction with the AIGP Attribute.

If a Route Reflector does not pass all paths to its clients, then it will tend to pass the paths for which the IGP distance from the Route Reflector itself to the next hop is smallest. This may result in a non-optimal choice by the clients.

6. IANA Considerations

IANA shall assign a codepoint for the AIGP attribute. This codepoint will come from the "BGP Path Attributes" registry.

IANA shall create a registry for "BGP AIGP Attribute Types". Type 1 should be defined as "AIGP", and should refer to this document.

7. Security Considerations

The spurious introduction, though error or malfeasance, of an AIGP attribute, could result in the selection of paths other than those desired.

Improper configuration on both ends of an EBGP connection could result in an AIGP attribute being passed from one service provider to another. This would likely result in an unsound selection of paths.

8. Acknowledgments

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