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**The HMAC-SHA-256-128 Algorithm and Its Use With IPsec**  
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Abstract

This document describes the use of the HMAC algorithm in conjunction with the SHA-256 algorithm as an experimental authentication mechanism within the context of the IPsec AH and ESP protocols. This algorithm is intended to provide data origin authentication and integrity protection. Given the current lack of practical experience with SHA-256, implementations based on this document will be experimental in nature, and implementation is not required in order to claim compliance with the IPsec proposed standards. The version of the HMAC-SHA-256 authenticator described in this document specifies truncation to 128 bits, and is therefore named HMAC-SHA-256-128.



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## **1. Specification of Requirements**

The keywords "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" that appear in this document are to be interpreted as described in [\[RFC-2119\]](#).

## **2. Introduction**

This document specifies the use of SHA-256 [\[SHA2-1\]](#) combined with HMAC [\[HMAC\]](#) as an experimental keyed authentication mechanism within the context of the IPsec AH and ESP protocols. This algorithm is intended to provide data origin authentication and integrity protection. Given the current lack of practical experience with SHA-256, implementations based on this document will be experimental in nature, and implementation is not required in order to claim compliance with the IPsec proposed standards. Furthermore, HMAC-SHA-1-96 [\[HMAC-SHA\]](#) provides sufficient security at a lower computational cost. The version of the HMAC-SHA-256 authenticator described in this document specifies truncation to 128 bits, and is therefore named HMAC-SHA-256-128. For further information on ESP, refer to [\[ESP\]](#) and [\[ROADMAP\]](#). For further information on AH, refer to [\[AH\]](#) and [\[ROADMAP\]](#).

The goal of HMAC-SHA-256-128 is to ensure that the packet is authentic and cannot be modified in transit. Data integrity and data origin authentication as provided by HMAC-SHA-256-128 are dependent upon the scope of the distribution of the secret key. If the key is known only by the source and destination, this algorithm will provide both data origin authentication and data integrity for packets sent between the two parties. In addition, only a party with the identical key can verify the MAC.

## **3. The HMAC-SHA-256-128 Algorithm**

[\[SHA2-1\]](#) and [\[SHA2-2\]](#) describe the underlying SHA-256 algorithm, while [\[HMAC\]](#) describes the HMAC algorithm. The HMAC algorithm provides a framework for inserting various hashing algorithms such as SHA-256.

The following sections contain descriptions of the various characteristics and requirements of the HMAC-SHA-256-128 algorithm.

### **3.1 Keying Material**

HMAC-SHA-256-128 is a secret key algorithm. While no fixed key length is specified in [\[HMAC\]](#), for use with either ESP or AH a fixed key length of 256-bits MUST be supported. Key lengths other than 256-bits MUST NOT be supported (i.e. only 256-bit keys are to be used by

HMAC-SHA-256-128). A key length of 256-bits was chosen based on the recommendations in [\[HMAC\]](#) (i.e. key lengths less than the authenticator length decrease security strength and keys longer than the authenticator length do not significantly increase security strength).

[HMAC] discusses requirements for key material, which includes a discussion on requirements for strong randomness. A strong pseudo-random function **MUST** be used to generate the required 256-bit key.

At the time of this writing there are no specified weak keys for use with HMAC. This does not mean to imply that weak keys do not exist.

[ARCH] describes the general mechanism for obtaining keying material when multiple keys are required for a single SA (e.g. when an ESP SA requires a key for confidentiality and a key for authentication).

In order to provide data origin authentication, the key distribution mechanism must ensure that unique keys are allocated and that they are distributed only to the parties participating in the communication.

[HMAC] makes the following recommendation with regard to rekeying: "Current attacks do not indicate a specific recommended frequency for key changes ... However, periodic key refreshment is a fundamental security practice that helps against potential weaknesses of the function and keys, and limits the damage of an exposed key." Rekeying also reduces the information available to a cryptanalyst.

### **[3.2](#) Padding**

HMAC-SHA-256-128 operates on 512-bit blocks of data. Padding requirements are specified in [[SHA2-1](#)] and are part of the SHA-256 algorithm. If you build SHA-256 according to [[SHA2-1](#)] you do not need to add any additional padding as far as HMAC-SHA-256-128 is concerned. With regard to "implicit packet padding" as defined in [[AH](#)], no implicit packet padding is required.

### **[3.3](#) Truncation**

HMAC-SHA-256-128 produces a 256-bit authenticator value. This 256-bit value can be truncated as described in [[HMAC](#)]. For use with either ESP or AH, a truncated value using the first 128 bits **MUST** be supported. Upon sending, the truncated value is stored within the authenticator field. Upon receipt, the entire 256-bit value is computed and the first 128 bits are compared to the value stored in the authenticator field. No other authenticator value lengths are supported by HMAC-SHA-256-128.

The length of 128 bits was selected because it meets the security requirements described in [[HMAC](#)]. [[HMAC](#)] discusses the potential additional security which is provided by the truncation of the resulting MAC. Specifications which include HMAC are strongly encouraged to perform this MAC truncation.

### **3.4 Interaction with the ESP Cipher Mechanism**

As of this writing, there are no known issues which preclude the use of the HMAC-SHA-256-128 with any specific cipher algorithm.

### 3.5 Performance

[HASH] states that "(HMAC) performance is essentially that of the underlying hash function". As of this writing no detailed performance analysis has been done of SHA-256, HMAC or HMAC combined with SHA-256.

[HMAC] outlines an implementation modification which can improve per-packet performance without affecting interoperability.

### 3.6 Test Vectors

The following test cases for HMAC-SHA-256 and HMAC-SHA-256-128 include the key, the data, and the resulting HMAC. The values of keys and data are either hexadecimal numbers (prefixed by "0x") or ASCII character strings (surrounded by double quotes). If a value is an ASCII character string, then the HMAC computation for the corresponding test case DOES NOT include the trailing null character ('\0') of the string. The computed HMAC values are all hexadecimal numbers.

These test cases were verified using 3 independent implementations: an HMAC wrapper on top of Aaron Gifford's SHA256 implementation ([www.aarongifford.com/computers/sha.html](http://www.aarongifford.com/computers/sha.html)), the BeeCrypt crypto library ([www.virtualunlimited.com/products/beecrypt](http://www.virtualunlimited.com/products/beecrypt)) and the Nettle cryptographic library ([www.lysator.liu.se/~nisse/nettle](http://www.lysator.liu.se/~nisse/nettle)). Partial blocks were padded as specified in [[SHA2-1](#)].

Test cases 1 and 2 were taken from the SHA-2 FIPS [[SHA2-1](#)] and test cases 4-10 were borrowed from [[HMAC-TEST](#)] with some key sizes adjusted for HMAC-SHA-256. These test cases illustrate HMAC-SHA-256 with various combinations of input and keysize. All test cases include the computed HMAC-SHA-256; only those with a keysize of 32 bytes (256 bits) also include the truncated HMAC-SHA-256-128.

Test Case #1: HMAC-SHA-256 with 3-byte input and 32-byte key

```
Key_len      : 32
Key          : 0x0102030405060708090a0b0c0d0e0f10
              1112131415161718191a1b1c1d1e1f20
Data_len     : 3
Data         : "abc"
HMAC-SHA-256 : 0xa21b1f5d4cf4f73a4dd939750f7a066a
              7f98cc131cb16a6692759021cfab8181
HMAC-SHA-256-128: 0xa21b1f5d4cf4f73a4dd939750f7a066a
```

Test Case #2: HMAC-SHA-256 with 56-byte input and 32-byte key

```
Key_len      : 32
Key          : 0x0102030405060708090a0b0c0d0e0f10
              1112131415161718191a1b1c1d1e1f20
Data_len     : 56
```

Data : "abcdbcdecdefdefgefghfghighijhijk  
ijkljklmklmnlmnomnopq"  
HMAC-SHA-256 : 0x104fdc1257328f08184ba73131c53cae  
e698e36119421149ea8c712456697d30  
HMAC-SHA-256-128: 0x104fdc1257328f08184ba73131c53cae

Test Case #3: HMAC-SHA-256 with 112-byte (multi-block) input  
and 32-byte key

Key\_len : 32  
Key : 0x0102030405060708090a0b0c0d0e0f10  
1112131415161718191a1b1c1d1e1f20  
Data\_len : 112  
Data : "abcdbcdecdefdefgefghfghighijhijk  
ijkljklmklmnlmnomnopnopqabcdbcde  
cdefdefgefghfghighijhijkijkljklm  
klmnlmnomnopnopq"  
HMAC-SHA-256 : 0x470305fc7e40fe34d3eeb3e773d95aab  
73acf0fd060447a5eb4595bf33a9d1a3  
HMAC-SHA-256-128: 0x470305fc7e40fe34d3eeb3e773d95aab

Test Case #4: HMAC-SHA-256 with 8-byte input and 32-byte key

Key\_len : 32  
Key : 0x0b repeated 32 times  
Data\_len : 8  
Data : 0x4869205468657265  
Data : "Hi There"  
HMAC-SHA-256 : 0x198a607eb44bfbcb69903a0f1cf2bbdc5  
ba0aa3f3d9ae3c1c7a3b1696a0b68cf7  
HMAC-SHA-256-128: 0x198a607eb44bfbcb69903a0f1cf2bbdc5

Test Case #5: HMAC-SHA-256 with 28-byte input and 4-byte key

Key\_len : 4  
Key : "Jefe"  
Data\_len : 28  
Data : "what do ya want for nothing?"  
HMAC-SHA-256 : 0x5bdcc146bf60754e6a042426089575c7  
5a003f089d2739839dec58b964ec3843

Test Case #6: HMAC-SHA-256 with 50-byte input and 32-byte key

Key\_len : 32  
Key : 0xaa repeated 32 times  
Data\_len : 50  
Data : 0xdd repeated 50 times  
HMAC-SHA-256 : 0xcdcb1220d1ecccea91e53aba3092f962  
e549fe6ce9ed7fdc43191fbde45c30b0  
HMAC-SHA-256-128: 0xcdcb1220d1ecccea91e53aba3092f962

Test Case #7: HMAC-SHA-256 with 50-byte input and 37-byte key

Key\_len : 37  
Key : 0x0102030405060708090a0b0c0d0e0f10  
1112131415161718191a1b1c1d1e1f20  
2122232425  
Data\_len : 50  
Data : 0xcd repeated 50 times

HMAC-SHA-256 : 0xd4633c17f6fb8d744c66dee0f8f07455  
6ec4af55ef07998541468eb49bd2e917

Test Case #8: HMAC-SHA-256 with 20-byte input and 32-byte key

Key\_len : 32

Key : 0x0c repeated 32 times

```
Data_len      : 20
Data          : "Test With Truncation"
HMAC-SHA-256  : 0x7546af01841fc09b1ab9c3749a5f1c17
               d4f589668a587b2700a9c97c1193cf42
HMAC-SHA-256-128: 0x7546af01841fc09b1ab9c3749a5f1c17
```

Test Case #9: HMAC-SHA-256 with 54-byte input and 80-byte key

```
Key_len       : 80
Key           : 0xaa repeated 80 times
Data_len      : 54
Data          : "Test Using Larger Than Block-Size Key -
               Hash Key First"
HMAC-SHA-256  : 0x6953025ed96f0c09f80a96f78e6538db
               e2e7b820e3dd970e7ddd39091b32352f
```

Test Case #10: HMAC-SHA-256 with 73-byte (multi-block) input  
and 80-byte key

```
Key_len       : 80
Key           : 0xaa repeated 80 times
Data_len      : 73
Data          : "Test Using Larger Than Block-Size Key and
               Larger Than One Block-Size Data"
HMAC-SHA-256  : 0x6355ac22e890d0a3c8481a5ca4825bc8
               84d3e7a1ff98a2fc2ac7d8e064c3b2e6
```

## **4. IKE Interactions**

### **4.1 Phase 1 Identifier**

For Phase 1 negotiations, IANA has assigned a Hash Algorithm ID of 4 for SHA2-256.

For further information on the use of Hash Algorithm IDs within IKE, see [[IKE](#)].

### **4.2 Phase 2 Identifier**

For Phase 2 negotiations, IANA has assigned an AH Transform Identifier of 5 for AH\_SHA2-256.

For Phase 2 negotiations, IANA has assigned an AH/ESP Authentication Algorithm Attribute Value of 5 for HMAC-SHA2-256.

For further information on the use of Transform Identifiers and Attributes Value within IKE, see [[IKE](#)] and [[DOI](#)].

## **5. Security Considerations**

The security provided by HMAC-SHA-256-128 is based upon the strength

of SHA-256. At the time of this writing there are no practical cryptographic attacks against SHA-256.

As is true with any cryptographic algorithm, part of its strength lies in the correctness of the algorithm implementation, the security

of the key management mechanism and its implementation, the strength of the associated secret key, and upon the correctness of the implementation in all of the participating systems. This draft contains test vectors to assist in verifying the correctness of HMAC-SHA-256-128 code.

## **6. IANA Considerations**

IANA has assigned Hash Algorithm ID 4 to SHA2-256.  
IANA has assigned AH Transform Identifier 5 to AH\_SHA2-256.  
IANA has assigned AH/ESP Authentication Algorithm Attribute Value 5 to HMAC-SHA2-256.

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## **8. Acknowledgments**

Portions of this text were unabashedly borrowed from [[HMAC-SHA](#)].

Thanks to Hugo Krawczyk for his comments and recommendations.

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