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## **Security Policy Protocol**

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Abstract

This document describes a protocol for discovering, accessing and processing security policy information of hosts, subnets or networks of a security domain. The Security Policy Protocol defines how the policy information is exchanged, processed, and protected by clients and servers. The protocol is extensible and flexible. It allows the exchange of complex policy objects between clients and servers.

Sanchez, Condell

[Page 1]

1.2 Policies       5         2. Overview       5         3. SPP Message       6         3.1 SPP Message Format       8         3.2 SPP Payloads       11         3.2.1 Query Payload       12         3.2.2 Record Payload       12         3.2.3 Signature Payload       13         3.3 SPP Messages       14         3.3.1 Query Messages       14         3.3.2 Reply Messages       14         3.3.2 Reply Messages       14	1999
1.1 Definitions.       2         1.2 Policies       5         2. Overview       5         3. SPP Message       7         3.1 SPP Message Format       8         3.2 SPP Payloads       11         3.2.1 Query Payload       11         3.2.2 Record Payload       12         3.2.3 Signature Payload       13         3.3 SPP Messages       14         3.3.1 Query Messages       14         3.3.2 Reply Messages       14	
3. SPP Message	<u>4</u>
3.1 SPP Message Format       8         3.2 SPP Payloads       11         3.2.1 Query Payload       12         3.2.2 Record Payload       12         3.2.3 Signature Payload       13         3.3 SPP Messages       14         3.3.1 Query Messages       14         3.3.2 Reply Messages       14         3.3.3 SPP Messages       14	<u>5</u>
3.3.3 Policy Messages	8 1 1 2 3 4 4 4 5 5 5 5
4. Policy Queries	<u>6</u> 6
5. Policy Records       15         5.1 Security Gateway Record       15         5.2 COMSEC Record       26         5.3 Security Association Record       22         5.4 Policy Server Record       23         5.5 Certificate Record       25	9 0 2 3
6. Transfer Records	<u>5</u>
7. Policy Attribute Encoding	<u>6</u>
8. SPP Message Processing	8 8 2 5

Int	ernet Draft	Secu	ri	ty	Ро	li	СУ	Pr	ot	00	co]	L							Jι	ıly	/ 1999	)
<u>9</u> .	IANA Considerations.																				<u>41</u>	
	9.1 Message Type																				<u>41</u>	
	9.2 Message Code																				<u>41</u>	
	9.3 Identity Type																				<u>41</u>	
	<u>9.4</u> Payload Class																				<u>42</u>	
	<u>9.5</u> Query Type																				<u>42</u>	
	<u>9.6</u> Record Type																				<u>42</u>	
	<u>9.7</u> Signature Type .																					
	9.8 Certificate Type																					
	9.9 Certificate Ident																					
	9.10 Attribute Data 1																					
	9.11 User Name Type.																					
	9.12 System Name Type																					
	9.13 IPsec Action Att																					
	9.14 IKE Action Attri	Lbute		•		•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	<u>43</u>	
<u>10</u> .	Security Consideration	ons.																			<u>44</u>	
Ackı	nowledgments																				<u>44</u>	
Ref	erences																				<u>45</u>	
<u>App</u>	endix A. DATA_TYPE Def	init	io	ns																	<u>46</u>	
<u>App</u>	endix B. An SPP Exampl	Le .																			<u>69</u>	
Dis	claimer																				<u>75</u>	
Autl	nor Information																				<u>75</u>	
San	chez, Condell																		[	Ρá	age 3]	

#### 1. Introduction

The IPSec protocols [Kent98] provide a mechanism for securing communications at the IP layer and IKE [Harkins98] can be used to provide keys for IPSec. Currently practice with these protocols maintains an assumption that communicating hosts have some a-priori knowledge of which communications with particular newtwork entities must be secured. While this assumption is valid in some environments (e.g. some VPN environments), it does not support more general IPSec senarios in a scalable manner.

In order to allow IPSec to scale in general cases, it is necessary to be able to identify which entities involved in a communication will require IPSec to protect the communication and what their policies are regarding it.

The Security Policy Protocol (SPP) defines how the policy information is exchanged, processed, and protected by clients and servers. The protocol also defines what policy information is exchanged and the format used to encode the information. The protocol specifies six different message types used to exchange policy information. An SPP message contains a message header section followed by zero or more SPP payloads, depending on the message type.

SPP is part of the Security Policy System architecture [SPS]. This document uses terms and references functionality described in [SPS].

The remainder of this section defines terms and concepts that will be used throughout this document. <u>Section 2</u> provides and overview of the protocol. The remainder of the document describes the encoding of the protocol and how SPP messages are processed.

## **1.1** Definitions

The following terms are used throughout this document, in addition to the terms defined in section [\*\*\*] of [SPS].

#### Authoritative

Host A is authoritative over host B if host A has the right to represent policy for host B. Host A may assert its relationship to host B using policy server records ( $\underline{\text{section 5.4}}$ ), but MUST be able to cryptographically prove the assertion.

## Transitively Authoritative

A host is transitively authoritative over another host, A, if it is either authoritative over host A or authoritative over a host, B, which is trasitively athoritative over host A. For example, if host X is authoritative for host Y and host Y is authoritative

Sanchez, Condell

[Page 4]

Chain of Trust

A chain of trust is a set of cryptographically-proven authoritative assertions that prove that a policy server is transitively authoritative over the source or destination of a communication. The chain of trust is used to prove that a policy server has a right to be involved in an SPP exchange. See <a href="section 10">section 10</a> for more about the chain of trust.

Keywords "MUST", "MUST NOT", "REQUIRED", "SHOULD", "SHOULD NOT" and "MAY" that appear in this document are to be interpreted as described in [Bra97].

#### 1.2 Policies

Defining and storing policies are beyond the scope if this document. However, this section describes SPP's policy requirements and a brief high-level look at its representation.

## Policy Representation

SPP provides both the comsec record (section 5.2) and the Security Association (SA rec) record (section 5.3) to describe policies. The comsec record defines the selectors that describe a communication along with a permit or deny action. The SA rec defines the actions, specifically the IPSec and IKE security associations, necessary for the communication to proceed. A policy transferred by SPP, therefore, MUST consist of one comsec record to describe the selectors of the communication and zero or more SA recs which describe the security associations that are required to complete the communication.

## Decorrelation

Policies exchanged using SPP MUST be decorrelated as described in [SPS]. Two policies are decorrelated if there exists at least one selector in both policies for which their values do not intersect. Decorrelation is necessary to permit policy servers to properly cache policies.

### 2. Overview

This section provides an overview of the SPP operation. A more detailed and complex example of SPP operation is available in appendix **B**. This overview assumes the policy servers have been loaded with policies for their security domains and the policy has been appropriately decorrelated (as specified in [SPS]).

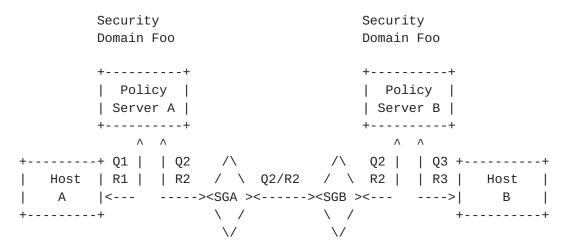


Figure 1: Overview of SPP operation

Host A, wanting to communicate with Host B, invokes its policy client. Host A's client sends a Query (Q1) to its configured local policy server, Policy Server A. Policy Server A looks in its cache for a policy record that matches the query. If it doesn't find one, it sends a Query (Q2) containing the same policy request information to Host B. Q2 is sent to Host B since Policy Server A may not know about the existence of SGB or Policy Server B. This message includes a signature that validates the authenticity and integrity of the query's content.

(Q2) is intercepted by SGB. SGB forwards the message (Q2) to Policy Server B. Policy server B verifies that it can accept queries from Policy Server A and validates the signature in Q2. It searches its database for the appropriate policy information after verifying that it is authoritative over Policy Client B.

Policy Server B merges its local policy with the policy information in (Q2) and it sends a Reply (R2) to Policy Server A. The reply includes the original query information and all policy information needed to allow Policy Client A to establish a secure communication with Host

B. Policy Server B also attaches additional information to the reply asserting its authority over Host B.

When Policy Server A receives the reply (R2) from Policy Server A, it validates the signature in R2 and cryptographically verifies that Policy Server B is authoritative over Host B. It then merges is local policy with the policy information in (R2) and sends a Reply (R1) to Host A. Policy Server A caches the merged policy to use when answering future queries. Host A may then use this information to establish necessary security associations with Host B.

If, however, Policy Server B is not authoritative over Host B, it would query Host B for its policy with respect to this particular communication. Policy Server B would generate a third query (Q3). Host

B would respond with its policy in (R3). Policy Server B merges its policy for this communication and the policy in (R3) before replying to Policy Server A. Policy Server A processes the reply as it did above.

Sanchez, Condell

[Page 6]

SPP accommodates topology changes, hence policy changes, rather easily without the scalability constraints imposed by static reconfiguration of each client. The protocol is extensible and flexible It allows the exchange of complex policy objects between clients and servers.

#### 3. SPP Message

The SPP header is present in every message. It contains fields identifying the message, the type of message, the status of the message, the number of queries and/or record payloads, and the host requesting policy information. The header also includes a timestamp field that provides anti-replay protection. Following the header there might be zero or more SPP payloads. Currently, there are three payload types defined in SPP: Query, Record, and Signature payloads. See section 3.2 for encoding details.

SPP has six distinct message types. Query messages contain a specific request for policy information. Reply messages include policy records that answer specific policy queries. Policy messages include policy information and are utilized for up/downloading security policies to and from a policy server. Policy Acknowledgment messages are utilized to acknowledge corresponding Policy messages but do not themselves contain policy information. Transfer messages, which include policy information, are utilized by policy servers to exchange bulk policy information between servers. Finally, policy servers use keep alive messages to inform security gateways and/or other monitoring devices of the status of the server.

SPP messages MUST be authenticated either using IPSec [Kent98] or another security mechanism. SPP provides a basic security mechanism that can be used to provide authentication and integrity to its messages when other security mechanisms are not in use. The SPP authentication is especially useful when traversing heterogenous domains and the identity of the policy server authoritative for the destination is unknown. These services are provided using digital signatures.

SPP caries signatures in the signature payload. The signature is calculated over the entire SPP message. When this service is used, the entity (host, policy server, or security gateway) verifying the signature must have access to the public key that corresponds to the private key used to sign the SPP message.

Certificate fetching is out of the scope of SPP. However, SPP provides a simple certificate fetching mechanism for entities that elect to use it as an alternative to other mechanisms. SPP suports several Public Key certificates formats.

SPP is modular and extensible (See <a href="section9">section9</a>. for IANA

considerations). New policy queries and records can be defined and incorporated easily. This document defines a minimum set of queries and policy records required in a policy-based security management system.

Sanchez, Condell

[Page 7]

### 3.1 SPP Message Format

An SPP message follows the format depicted in figure 2. It is comprised of a header and zero or more SPP payloads. This section defines the encoding for the SPP header. Sections 3.2 and 3.3 cover the encoding for the SPP payload and message types, respectively.

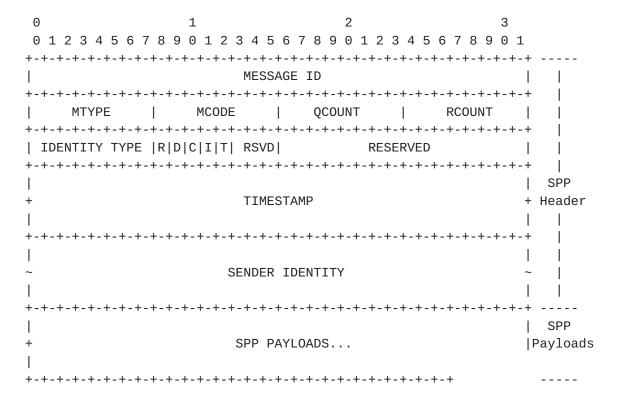


Figure 2: Format of an SPP Message

The SPP header includes the following fields:

# MESSAGE ID

A 4 octet field used to match messages and their responses (e.g. queries to replies and policy to policy acknowledgement messages). This value starts at "zero" and MUST be incremented by (1) with every new message.

## MTYPE

A 1-octet field indicating the SPP message type. The currently defined values are:

Message Type	Va⊥ue
Value Not Assigned	0
SPP-QUERY	1
SPP-REPLY	2
SPP-P0L	3

SPP-POL_ACK	4
SPP-XFR	5
SPP-KEEP_ALIVE	6

Sanchez, Condell [Page 8]

values 7-250 are reserved to IANA. Values 251-255 are for private use among mutually consenting parties.

#### MCODE

A 1-octet field providing information about this message. All MTYPEs share a common MCODE space, although each message type may not use all the defined message codes. See <u>section</u> 3.3 for the codes applicable to each message type.

Action	Code
Туре	Field
Value Not Assigned	0
message accepted	1
denied, administratively prohibited	2
denied, timestamp failed	3
denied, failed signature	4
denied, insufficient resources	5
denied, malformed message	6
denied, unspecified	7
partially available	8
unavailable	9
communication prohibited	10
partially available, server unreachable	11

values 12-250 are reserved to IANA. Values 251-255 are for private use among mutually consenting parties.

## QCOUNT

A 1 octet field indicating the number of Query payloads included in the message.

## RCOUNT

A 1 octet field indicating the number of Record payloads included in the message.

# IDENTITY TYPE

This 1 octet field indicates the type of indentity found in the Sender Identity field. Valid values are:

Identity Type	Value
Value Not Assigned	Θ
IPV4_ADDR	1
IPV6_ADDR	2
Host DNS Name	3

values 4-250 are reserved to IANA. Values 251-255 are for private use among mutually consenting parties.

Raw policy flag. When this flag is set, policy servers MUST not resolve the policies that they return.

Sanchez, Condell

[Page 9]

D

Domain flag. Only resolve the policies as far as the last policy server that is transitively authoritative over the host requesting the policy resolution.

С

Dont cache flag. Don't cache the policies generated by the query.

Ι

Ignore cache flag. Ignore any cached policies when processing the query.

Τ

No chain-of-trust. A client indicates to its server that it does not need chain-of-trust information. Policy Servers MUST NOT set this flag. Only Policy Clients have the option to set it.

### **RSVD**

A 4 bit field reserved for future use. Set value to all zeros (0).

#### **RESERVED**

A 2 octet field used for field 32-bit boundary alignment and reserved for future use. Set value to all zeros (0).

#### **TIMESTAMP**

This 8-octet field contains a timestamp used to provide limited protection against replay attacks. The timestamp is formatted as specified by the Network Time Protocol [RFC1305].

## SENDER IDENTITY

A variable length field containing the identity of the sender (host, security gateway, or policy server) of the SPP message. The IDENTITY\_TYPE field indicates the format of the content in this field:

Identity Type Sender Identity

IPV4\_ADDR An IPv4 Address IPV6\_ADDR An IPv6 Address

Host DNS Name A DNS name encoded as described

in [<u>rfc1035</u>]

This field does not allow IP address ranges or wildcards. If this field is not aligned at the 4 octet boundary, the field MUST be padded on the right with (00)hex to align on the next 32-bit boundary.

### 3.2 SPP Payloads

# 3.2.1 Query Payload

The Query payload contains fields to express a particular request for policy information. Hosts, security gateways, or policy servers can generate and transmit Query payloads in SPP messages to policy servers. Figure 3 shows the format of the Query payload.

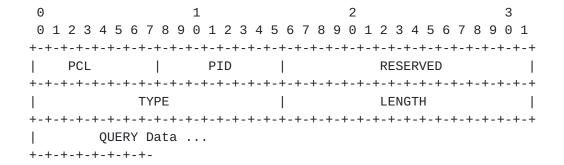


Figure 3: Format of Query Payload

The Query Payload fields are defined as follows:

PCL

A 1 octet field indicating the payload class. Query payloads MUST contain (1) in the PCL field.

PID

A 1 octet field containing the ID number that identifies a particular Query payload within an SPP message. Since one SPP message can contain multiple Query payloads, each one MUST be uniquely identified. This number MUST be unique among the Query payloads within an SPP message.

# RESERVED

A 2 octet field reserved for future use. Set value to all zeros (0).

TYPE

A 2 octet field that specifies the type of query contained in the QUERY Data fields. The currently defined queries are:

Query Payload Type	Value
Value Not Assigned	0
Security Gateway Query	1
Communication Security Query	2
Certificate Query	3

values 4-65000 are reserved to IANA. Values 65001-65535 are for private use among mutually consenting parties.

Sanchez, Condell

[Page 11]

LENGTH

A 2 octet field indicating the length in octets of the query data field.

**QUERY** Data

A variable length field containing a single policy query. See section 7 for encoding format.

## 3.2.2 Record Payload

The Record payload contains fields that assert policy information. Hosts, security gateways, or policy servers can generate and transmit Record payloads in SPP messages. Figure 4 shows the format of the Record payload.

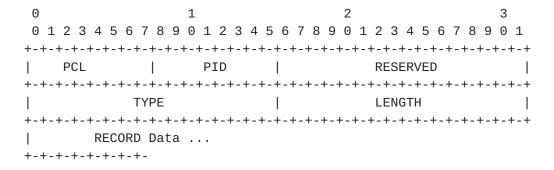


Figure 4: Format of Record Payload

The Record Payload fields are defined as follows:

PCL

A 1 octet field indicating the payload class. Record payloads MUST contain (2) in the PCL field.

PID

This field is used to match queries to Record payloads. If the record is a reply to a query, then the value for this field MUST match the correspondent Query payload PID. If it is not a reply to a query, the value SHOULD be set to zero.

## RESERVED

A 2 octet field reserved for future use. Set value to all zeros (0).

Sanchez, Condell

TYPE

A 2 octet field that specifies the type of Record. The currently defined records are:

Record Type	Value
Value Not Assigned	0
Security Gateway Record	1
Communication Security Record	2
Security Association Record	3
Certificate Record	4
Policy Server Record	5
Transfer Record	6

values 7-65000 are reserved to IANA. Values 65001-65535 are for private use among mutually consenting parties.

#### LENGTH

A 2 octet field indicating the length in octets of the RECORD data field.

#### RECORD Data

A variable length field containing a single policy record. See section 8 for encoding format.

# 3.2.3 Signature Payload

The Signature Payload contains data generated by the digital signature function (selected by the originator), over the entire SPP message, except for part of the Signature payload. This payload is used to verify the integrity of the data in the SPP message. Figure 5 shows the format of the Signature payload.

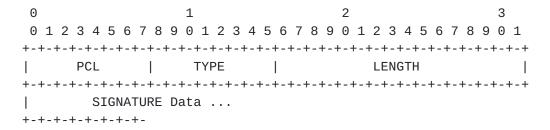


Figure 5: Signature Payload Format

The Signature payload fields are defined as follows:

PCL

A 1 octet field indicating the payload class. Signature payloads MUST contain (3) in the PCL field.

TYPE

A 1 octet field that specifies the signature algorithm employed. The currently defined signature types are:

Sanchez, Condell

[Page 13]

Algorithm		Туре	Value
Value	Not	Assigned	Θ
RSA			1
DSA			2

values 3-250 are reserved to IANA. Values 251-255 are for private use among mutually consenting parties.

#### LENGTH

A 2 octet field indicating the length in octets of the SIGNATURE Data field.

#### SIGNATURE Data

A variable length field that contains the results from applying the digital signature function to the entire SPP message (including the PCL, TYPE, and LENGTH fields of the Signature payload), except for the Signature Data field of the Signature payload.

## 3.3 SPP Messages

## 3.3.1 Query Message

An SPP-QUERY message is comprised of an SPP header, one or more Query payloads, zero or more Record payloads, and a Signature payload, if one is required. Query messages MUST always contain a Query payload. Record payloads may optionally be included to pass policy information along with the query. If the Signature payload is employed it MUST be the last payload in the message. The Query message MTYPE value is (1). The MCODE field must be set to zero (0).

## 3.3.2 Reply Message

An SPP-REPLY message is comprised of an SPP header, one or more Query payloads, zero or more Record payloads which answer the corresponding Query payload, and a Signature payload, if one is required. Reply messages MUST contain a Query payload. Reply messages MUST include a Record payload unless the reply contains an MCODE of values 2-8. If the Signature payload is employed it MUST be the last payload in the message. The MTYPE value for a Reply message is (2). The following MCODE values may be used for Reply messages:

Action	Code
Туре	Field
Value Not Assigned	0
message accepted	1
denied, administratively prohibited	2

denied,	timestamp failed	3
denied,	failed signature	4
denied,	insufficient resources	5
denied,	malformed message	6

Sanchez, Condell [Page 14]

denied, unspecified	7
partially available	8
unavailable	9
communication prohibited	10
partially available, server unreachable	11

### 3.3.3 Policy Message

An SPP-POL message is comprised of an SPP header, one or more Record payloads, and a Signature payload, if one is required. Policy messages MUST NOT include Query payloads. If the Signature payload is employed it MUST be the last payload in the message. The MTYPE value for a Policy message is (3). The MCODE field must be set to zero (0).

## 3.3.4 Policy Acknowledgement Message

An SPP-POL\_ACK message is comprised of an SPP header and a Signature payload, if one is required. These messages MUST NOT contain Query or Record payloads. The status of the associated Policy message is expressed within the MCODE field. If the Signature payload is employed it MUST be the only payload in the message. The MTYPE value for a Policy Acknowledgement message is (4). The following MCODE values may be used for Policy Acknowledgement messages:

Action	Code
Туре	Field
Value Not Assigned	0
message accepted	1
denied, administratively prohibited	2
denied, timestamp failed	3
denied, failed signature	4
denied, insufficient resources	5
denied, malformed message	6
denied, unspecified	7

# 3.3.5 Transfer Message

An SPP-XFR message is comprised of an SPP header, one or more Record payloads, and a Signature payload, if one is required. Transfer messages MUST NOT include Query payloads. If the Signature payload is employed it MUST be the last payload in the message. The MTYPE value for a Transfer message is (5). The MCODE field must be set to zero (0).

# 3.3.6 KeepAlive Message

An SPP-KEEP\_ALIVE message is comprised of an SPP header and a Signature payload, if one is required. These messages MUST NOT contain Query or Record payloads. If the Signature payload is employed it MUST be the only payload in the message. The MTYPE value for a KeepAlive message is (6). The MCODE field must be set to zero (0).

Sanchez, Condell

[Page 15]

### 4. Policy Queries

## 4.1 Security Gateway Query

This basic query provides a dynamic mechanism to determine which relevant security gateways, both primary and backup, are in the path to a particular destination address. Since the answer to a request for information could depend on the identity of the requestor, the host address of the source of the intended communication is included in the query. Figure 6 shows the format of the Security Gateway Query.

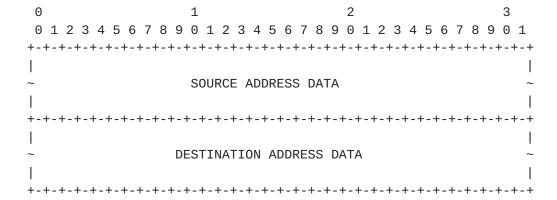


Figure 6: Security Gateway Query Format

The Security Gateway Query fields are defined as follows:

## SOURCE ADDRESS DATA

This variable length field contains a single IP address (unicast) either in IPv4 or IPv6 format. The encoding format is specified in <u>section 7</u>. The acceptable DATA\_TYPE values are 3 and 9.

### DESTINATION ADDRESS DATA

This variable length field contains a single IP address (unicast) either in IPv4 or IPv6 format. The encoding format is specified in <u>section 7</u>. The acceptable DATA\_TYPE values are 6 and 12.

## 4.2 COMSEC Query

The Communication Security Query (or COMSEC query) provides a dynamic mechanism for a host or security gateway to inquire if a communication having a particular set of characteristics is allowed. The communication is described in terms of source and destination addresses, protocols, source port, destination port, and other parameters as defined in <a href="mailto:section7">section 7</a>. These parameters are known as selectors in the IPSec context and are primarily the contents of the

IP, TCP, and UDP headers. Figure 7 shows the format of the COMSEC Query.

Sanchez, Condell

[Page 16]

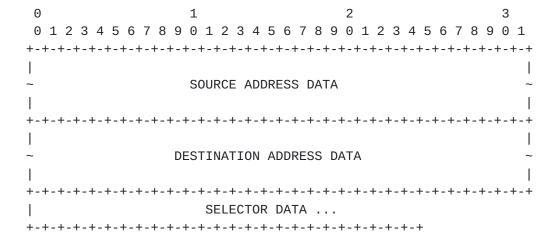


Figure 7: COMSEC Query Format

The COMSEC Query fields are defined as follows:

### SOURCE ADDRESS DATA

This variable length field contains a single IP address (unicast) either in IPv4 or IPv6 format. The encoding format is specified in <u>section 7</u>. The acceptable DATA\_TYPE values are 3 and 9.

## DESTINATION ADDRESS DATA

This variable length field contains a single IP address (unicast) either in IPv4 or IPv6 format. The encoding format is specified in <a href="mailto:section-7">section 7</a>. The acceptable DATA\_TYPE values are 6 and 12.

## SELECTOR DATA

This includes one or more fields following the encoding format specified in <u>section 7</u>. The acceptable DATA\_TYPE values are 15-29, inclusive.

# 4.3 CERT Query

Mechanisms to dispatch and fetch public-key certificates are not part of SPP. However, in the absence of external request/dispatch mechanisms, SPP provides for a certificate request query that allows a host, security gateway, or server to solicit a certificate. Figure 8 shows the format of the CERT Query.

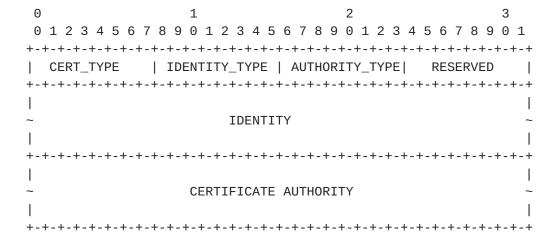


Figure 8: Certificate Query Format

The Certificate query fields are defined as follows:

# CERT\_TYPE

A 1 octet field that contains an encoding of the type of certificate requested. Acceptable values are listed below:

Certificate Type	Value
Value Not Assigned	0
PKCS #7 wrapped X.509 certificate	1
PGP Certificate	2
DNS Signed Key	3
X.509 Certificate - Signature	4
X.509 Certificate - Key Exchange	5
Kerberos Tokens	6
SPKI Certificate	7

values 8-250 are reserved to IANA. Values 251-255 are for private use among mutually consenting parties.

# IDENTITY\_TYPE

This 1 octet field indicates the type of indentity found in the Identity field. Valid values are listed below:

Identity Type
Value Not Assigned
IPV4_ADDR
IPV6_ADDR
DNS Name
X.500 Distinguished Name

values 5-250 are reserved to IANA. Values 251-255 are for private use among mutually consenting parties.

Sanchez, Condell

[Page 18]

#### AUTHORITY\_TYPE

This 1 octet field indicates the type of authority found in the Certificate Authority field. Valid values are the same as IDENTITY\_TYPE.

#### **IDENTITY**

This variable length field contains the identity of the principal by which the certificate should be located. The value MUST be of the type stated in IDENTITY\_TYPE.

## CERTIFICATE AUTHORITY

A variable length field containing an encoding of an acceptable certificate authority for the type of certificate requested. The value MUST be of the type stated in AUTHORITY\_TYPE.

## 5. Policy Records

## 5.1 Security Gateway Record

This record contains information that indicates the IP addresses of the interfaces for the the primary and secondary security gateways protecting a host or group of hosts. The record contains the primary and secondary gateways at one point in the communication path between the source and destination addresses listed in the Security Gateway query. If the IP datagram must traverse multiple gateways, a Security Gateway Record must be included for each gateway. The list of secondary security gateways is optional. Figure 9 shows the format of the Security Gateway Record.

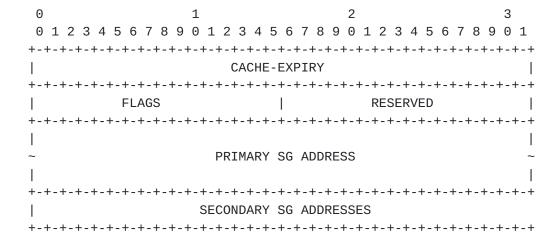


Figure 9: Security Gateway Record Format

The Security Gateway Record fields are defined as follows:

#### CACHE-EXPIRY

A 4 octet field indicating the maximum amount of time, in seconds, this policy record MAY be cached.

#### **FLAGS**

A 2 octet field indicating different options to aid in interpreting the security gateway data. If not in use, set value to all zeros (00)hex. Currently, no flag values are defined so this field MUST be set to (00)hex.

## RESERVED

A 2 octet field reserved for future use. Set value to all zeros (0).

### PRIMARY SG ADDRESS

A variable length field containing the IP address of the primary security gateway for protecting a particular host. This variable length field contains a single unicast IP address. The encoding format is specified in <a href="mailto:section7">section 7</a>. The acceptable DATA\_TYPE values are 1 and 2.

## SECONDARY SG ADDRESSES

This variable length field contains the IP addresses of one or more secondary security gateways protecting a particular host. This field may contain a list of single unicast IP addresses. The encoding format is specified in <a href="mailto:section7">section 7</a>. The acceptable DATA\_TYPE values are 1 and 2.

## 5.2 COMSEC Record

The COMSEC record indicates if a communication having a particular set of characteristics is allowed or not. The communication is described in terms of source and destination addresses, protocols, source ports, destination ports, and other attributes defined in <a href="mailto:section7">section 7</a>. Figure 10 shows the format of the COMSEC Record.

Sanchez, Condell

[Page 20]

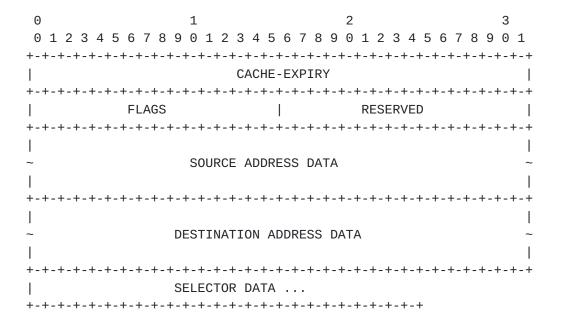


Figure 10: COMSEC Record Format

The COMSEC Record fields are defined as follows:

#### CACHE-EXPIRY

A 4 octet field indicating the maximum amount of time, in seconds, this policy record MAY be cached.

## **FLAGS**

A 2 octet field indicating different options to aid in interpreting the selector data. If not in use, set value to all zeros (0). Currently, no flag values are defined so this field MUST be set to zero (0).

#### **RESERVED**

A 2 octet field reserved for future use. Set value to all zeros (0).

## SOURCE ADDRESS DATA

This variable length field contains a single IP address (unicast, anycast, broadcast (IPv4 only), or multicast group), range of addresses (low and high values, inclusive), address + mask, or a wildcard address. The encoding format is specified in <a href="mailto:section 7">section 7</a>. The acceptable DATA\_TYPE values are 3-5 and 9-11, inclusive.

#### DESTINATION ADDRESS DATA

This variable length field contains a single IP address (unicast, anycast, broadcast (IPv4 only), or multicast group), range of addresses (low and high values, inclusive), address + mask, or a wildcard address. The encoding format is specified in <a href="mailto:section">section</a> 7. The acceptable DATA\_TYPE values are 6-8 and 12-14, inclusive.

#### SELECTOR DATA

This includes one or more fields following the encoding format specified in <u>section 7</u>. The acceptable DATA\_TYPE values are 15-29, inclusive.

# **5.3** Security Association Record

Security Association Records contain selectors and security association attributes (APPLIERS) that characterize a particular Security Association between the source and destination addresses listed in the record. This record contains data types as defined in the section 7. Figure 11 shows the format of the SA Record.

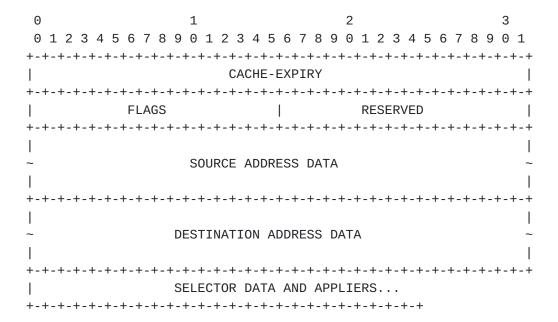


Figure 11: SA Record Format

The SA record fields are defined as follows:

# CACHE-EXPIRY

A 4 octet field indicating the maximum amount of time, in seconds, this policy record MAY be cached.

**FLAGS** 

A 2 octet field indicating different options to aid in interpreting the selector data. If not in use, set value to all zeros (0). Currently, no flag values are defined so this field MUST be set to zero(0).

### RESERVED

A 2 octet field reserved for future use. Set value to all zeros (0).

### SOURCE ADDRESS DATA

This variable length field contains a single IP address (unicast, anycast, broadcast (IPv4 only), or multicast group), range of addresses (low and high values, inclusive), address + mask, or a wildcard address. The encoding format is specified in <a href="mailto:section7">section 7</a>. The acceptable DATA\_TYPE values are 3-5 and 9-11, inclusive.

### DESTINATION ADDRESS DATA

This variable length field contains a single IP address (unicast, anycast, broadcast (IPv4 only), or multicast group), range of addresses (low and high values, inclusive), address + mask, or a wildcard address. The encoding format is specified in <a href="mailto:section 7">section 7</a>. The acceptable DATA\_TYPE values are 6-8 and 12-14, inclusive.

### SELECTOR DATA AND APPLIERS

This includes one or more fields following the encoding format specified in <u>section 7</u>. The acceptable DATA\_TYPE values are 15-29 and 50-51, inclusive.

# **5.4** Policy Server Record

The Policy Server record indicates the host, security gateway, or policy server for which a particular policy server is authoritative. It represents an assertion, typically made by a policy server, with repect to a member of a security domain that the server represents. The record includes the Identity of the policy server and the identity of a node (host, security gateway, another server, etc.). Figure 12 shows the format of the Policy Server Record.

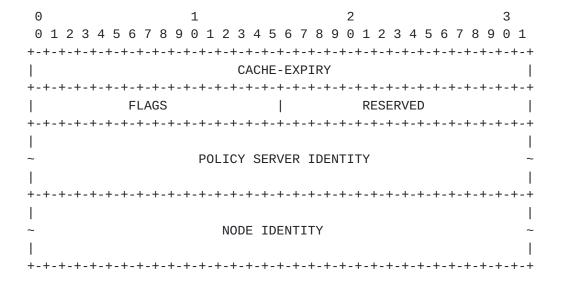


Figure 12: Policy Server record format

The Policy Server Record fields are defined as follows:

### CACHE-EXPIRY

A 4 octet field indicating the maximum amount of time, in seconds, this policy record MAY be cached.

# **FLAGS**

A 2 octet field indicating different options to aid in interpreting the server and node data. If not in use, set value to all zeros (0). Currently, no flag values are defined so this field MUST be set to zero (0).

# RESERVED

A 2 octet field reserved for future use. Set value to all zeros (0).

### POLICY SERVER IDENTITY

This variable length field contains the identity of the policy server. It may contain an IP address (unicast) either in IPv4 or IPv6 format. The encoding format is specified in <u>section 7</u>. The acceptable DATA\_TYPE values are 1 and 2.

### NODE IDENTITY

This variable length field contains the identity of a node for which the policy server is authoritative. It may contain an IP address (unicast) either in IPv4 or IPv6 format. The encoding format is specified in  $\underline{\text{section 7}}.$  The acceptable DATA\_TYPE values are 1 and 2.

Sanchez, Condell

[Page 24]

### **5.5** CERT Record

The CERT record contains one public key certificate. This record is provided in SPP as an alternate mechanism for certificate dispatching. Figure 13 shows the format of the CERT Record.

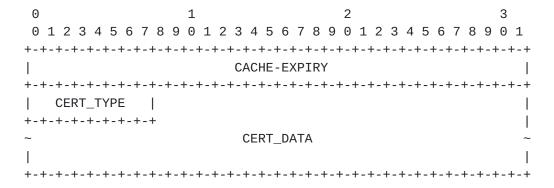


Figure 13: Certificate Record Format

CACHE-EXPIRY

A 4 octet field indicating the maximum amount of time, in seconds, this policy record MAY be cached.

CERT\_TYPE

This 1 octet field indicates the type of certificate or certificate-related information contained in the Certificate Data field. The values for this field are described in Section 4.3.

CERT\_DATA

This variable length field contains the actual encoding of certificate data. The type of certificate is indicated by the Certificate Type field.

# 6. Transfer Records

This record contains the text of the master file that is used to configure the primary policy server.

Θ	:	L	2		3
0 1 2	2 3 4 5 6 7 8 9	1 2 3 4 5	6 7 8 9 0 1 2 3	3 4 5 6 7 8 9	0 1
+-+-+-	+-+-+-+-+-+-+	+-+-+-+-+	+-+-+-+-+-	-+-+-+-+-	+-+-+
					- 1
~		MASTER FIL	E TEXT		~
					- 1
+-+-+-	+-+-+-+-+-+	+-+-+-+-+	+-+-+-+-+-	-+-+-+-+-	+-+-+

Figure 14: Security Gateway Record Format

The Transfer Record field is defined as follows:

Sanchez, Condell [Page 25]

MASTER FILE TEXT

This variable length field contains the text of the master file that is used to configure the policy server.

## 7. Policy Attribute Encoding

Query and Record payloads include several different selector types and SA attributes with their associated values. These data are encoded following a Type/Length/Value (TLV) format to provide flexibility for representing different kinds of data within a payload. Certain Data\_Types with values of length equal to 2 octets follow the Type/Value (T/V) format. The first bit of the DATA\_TYPE field is used to distinguished between the two formats. A value of (0) indicates a TLV format while a value of (1) indicates TV format. This generic encoding format is depicted in figure 15.

Figure 15: Generic Data Attribute Formats

The generic data attribute fields are defined as follows:

Χ

X = 0:

This bit indicates if the DATA\_TYPE follows the TLV(0) or the TV(1) format.

DATA\_TYPE | DATA\_VALUE |

DATA\_TYPE

A 2 octet field indicating the selector type. The currently defined values are:

Sanchez, Condell

DATA	DATA_TYPE	Χ
IPV4_ADDR	1	0
IPV6_ADDR	2	0
SRC_IPV4_ADDR	3	0
SRC_IPV4_ADDR_SUBNET	4	0
SRC_IPV4_ADDR_RANGE	5	0
DST_IPV4_ADDR	6	0
DST_IPV4_ADDR_SUBNET	7	0
DST_IPV4_ADDR_RANGE	8	0
SRC_IPV6_ADDR	9	0
SRC_IPV6_ADDR_SUBNET	10	0
SRC_IPV6_ADDR_RANGE	11	0
DST_IPV6_ADDR	12	0
DST_IPV6_ADDR_SUBNET	13	0
DST_IPV6_ADDR_RANGE	14	0
DIRECTION	15	1
USER_NAME	16	0
SYSTEM_NAME	17	0
XPORT_PROTOCOL	18	0
SRC_PORT	19	0
SRC_PORT_DYNAMIC	20	0
DST_PORT	21	0
DST_PORT_DYNAMIC	22	0
SEC_LABELS	23	0
V6CLASS	24	1
V6FL0W	25	0
V4T0S	26	1
ACTION	27	1
SRC_PORT_RANGE	28	0
DST_PORT_RANGE	29	0
IPSEC_ACTION	50	0
ISAKMP_ACTION	51	0

values 30-49 and 52-3200 are reserved to IANA. Values 3200-32767 are for private use among mutually consenting parties.

### LENGTH

A 2 octet field indicating the length of the selector value in octets, not including any trailing padding added to the DATA\_VALUE field. The padding length is implicit.

### DATA\_VALUE

A variable length field containing the value of the selector specified by DATA\_TYPE. If the Selector value is not aligned at

the 4 octet boundary the field MUST be padded on the right with (00)hex to align on the next 32-bit boundary.

Sanchez, Condell

[Page 27]

### 8. SPP Message Processing

SPP messages use UDP or TCP as their transport protocol. Messages carried by UDP are restricted to 512 bytes (not counting the IP or UDP headers). Fragmentation is allowed for messages containing more than 512 bytes. The SPP-XFR message SHOULD use TCP to transfer the contents of the SPS Database from a primary to secondary policy server. SPP uses UDP and TCP ports 501.

# 8.1 General Message processing

For an SPP-Query or SPP-Pol message, the transmitting entity MUST do the following:

- 1. Set a timer and initialize a retry counter.
- If an SPP-Reply or SPP-Pol\_Ack message corresponding to the appropriate SPP-Query or SPP-Pol message is received within the time interval, or before the retry counter reaches 0, the transmitting entity continues normal operation.
- 3. If an SPP-Reply or SPP-Pol\_Ack message corresponding to the appropriate SPP-Query or SPP-Pol message is not received within the time interval and a secondary policy server, which has not been attempted on this value of the retry counter, is available, the message is sent to the secondary server. If all the secondary servers fail to respond within the time interval, decrement the retry counter and resend the message to the primary server.
- 4. If the retry counter reaches zero (0) the event MAY be logged in the appropriate system audit file.
- 5. Following step 4, processing is more specific:
  - For hosts and security gateways:
    - o the transmitting entity will clear state for this peer and revert to using conventional security mechanisms.
  - For policy servers:
    - o For SPP-Pol messages, clear state for this peer.
    - o For SPP-Query messages, clear state for this peer, lookup policy in the server's SPS database and send an SPP-Reply message per <u>section 8.3</u> with the policy and MCODE 11.

# 8.2 Query Message (SPP-Query) Processing

When creating a SPP-Query message, the entity (host, security gateway, or policy server) MUST do the following:

Sanchez, Condell

[Page 28]

- 1. Generate the Message ID value. This value starts at zero (0) and MUST be incremented by (1) with every new message.
- 2. Set the value of the MTYPE field to 1
- 3. Set the MCODE value to zero (0).
- 4. Set the QCOUNT and RCOUNT fields. All fields MUST be set to zero (0) when their corresponding payloads do no exist.
- 5. Set the flag bits accordingly and set the RESERVED field to zero (0).
- 6. Set the IDENTITY\_TYPE and IDENTITY of the Sender of the SPP message.
- 7. Construct the SPP data payloads. A Query payload MUST be present in this message.
- 8. Local policy information related to the query MAY be included as Record payloads following the Query payloads.
- 9. A Policy Server record and a Cert Record SHOULD also be included in the message. A Cert record MUST be included if the query is a Cert Query to avoid a possible certificate query loop.
- 10. Calculate and place the timestamp value used for anti-replay attack protection.
- 11. If the Signature payload is required for message integrity and authentication, calculate a signature over the SPP header, SPP payloads, the PCL, TYPE, and LENGTH fields of the Signature payload. If required, the Signature payload MUST be the last payload in the message.

When a policy server receives an SPP-Query message it MUST do the following:

- 1. Check for SPP access control. If enabled, read the IP address in the Sender's field of the SPP header.
  - Verify whether or not the message is allowed. If the message is not allowed then:
    - o send an SPP-Reply message with the MCODE 2 or 7 o discard message and the event MAY be logged.
  - If the message is allowed, continue.
- 2. Check timestamp field for anti-replay protection. If a replayed

# message is detected:

- send an SPP-Reply message with the MCODE 3 or 7  $\,$
- discard message and the event MAY be logged.

Sanchez, Condell

[Page 29]

- 3. If the message requires signature validation.
  - If a certificate record is present, the server MUST process it, however, if the server already has a valid key for the host or server identified in the certificate, the certificate MAY be ignored.
  - Fetch certificate or key corresponding to the IP address found in the sender field of the SPP header.
  - If a certificate or key is not available the entity MAY, depending on configuration:
    - o choose to abort validation process, send SPP-Reply message with MCODE 5 or 7, discard the message, and MAY log the event.
    - o send an SPP-Query message to the source of the IP address found in the sender field of the SPP header with a CERT Query payload. Keep the SPP-Query message until the SPP-Reply returns. Extract certificate or key, validate it and proceed.
  - Once a validated certificate or key is available then validate signature.
    - o If validation fails, send SPP-Reply message with MCODE 5 or 7, discard the message, and the event MAY be logged.
- 4. Parse the Query records.
  - If the SPP-Query message only contains cert queries:
    - o If the Identity field identifies the server or a member of the server's security domain, send an SPP-Reply message per section 8.3 with one or more cert records with the certificates in the certification chain between the requested Identity and Certificate\_Authority and MCODE 1.
    - o Otherwise, send an SPP-Reply message per <u>section 8.3</u> with with MCODE 9 or 7.
  - If the SPP-Query message does not only contain cert queries:
    - o Check the Destination Address Data field to determine if the message received was intended for a node that is a member of the server's security domain.
    - o Continue processing.

5. If the message is for a node that is a member of the server's security domain or the D bit in the SPP header is set and the server is the outermost server that is authoritative over the client or server that sent the message, then:

Sanchez, Condell

[Page 30]

- Using src, dst, and any other applicable fields found in the Query Payload, search the SPS database for an appropriate policy.
  - o If a policy is found then construct an SPP-Reply message per <u>section 8.3</u> with appropriate payloads and MCODE 1.
  - o If a policy is not found then construct an SPP-Reply message per <u>section 8.3</u> with appropriate payloads and MCODE 9 or 7.
- 6. If the message is for a node that is not part of the server's security domain then:
  - If the I and R bits are not set in the SPP header, check the Cache database for any relevant policies that apply to this communication.
    - o If a policy is found then construct a SPP-Reply message per section 8.3 with appropriate payloads and MCODE 1.
    - o If a policy is not found then continue.
  - Check the Local database for any relevant policies that apply to this communication.
  - If the server is authoritative for the source address of the query or a matching policy is found (A matching policy is defined as a policy that either produces a comsec record with an action attribute with the value "deny", or a policy that would produce an SA record with one or more IPSec action and IKE action attributes. A policy that only produces a comsec record with an action attribute with the value "permit" is not considered matching for this step.):
    - o Generate a new SPP-Query message. The new message MUST use the same query payload as the old message. If the new query will include policy from the server, then the policy used SHOULD be the server's local policy merged with any policies received with the query message. The Sender Address will be the address of the server. The destination for this message MUST be the one in the original SPP-Query received.
    - o Keep state. Upon reception of the corresponding SPP-Reply the policy server MUST send an SPP-Reply addressed to sender of the original SPP-Query.
  - If the server is not authoritative for the source address of the query and a matching policy is not found then:

o The policy server MUST send the SPP-Query message unchanged. The SPP-Query message MUST use the same source port that was used to send it to the server so the next server that processes the query will return it to the correct port. This SPP-Query message MUST NOT be added to the retry queue (section 8.1).

# 8.3 Reply Message (SPP-Reply) Processing

When creating a SPP-Reply message, the policy server MUST do the following:

- 1. Copy the Message ID value from the corresponding SPP-Query message into the Message ID field.
- 2. Set the value of the MTYPE field to 2
- 3. Set the MCODE value.
- 4. Set the QCOUNT and RCOUNT fields. All fields MUST be set to zero (0) when their corresponding payloads do no exist.
- 5. Set the flag bits accordingly and set the RESERVED field to (0).
- 6. Set the IDENTITY\_TYPE and IDENTITY of the Sender of the SPP message.
- 7. Copy the Query payloads from the SPP-Query message to the SPP-Reply message.
- 8. Construct the SPP data payload. Unless there is an error, at least one record corresponding to each Query payload MUST be present.
- 9. A policy server record and a CERT record MUST be added to the SPP-Reply message if the the query to which this is a reply did NOT have the T bit set. If the T bit is set, the records SHOULD NOT be added.
- 10. Calculate and place the timestamp value used for anti-replay attack protection.
- 11. If the Signature payload is required for message integrity and authentication, calculate a signature over the SPP header, SPP payloads, the PCL, TYPE, and LENGTH fields of the Signature payload. If present, the Signature payload MUST be the last payload in the message.
- 12. The SPP-Reply message is sent to the host that is listed in the Sender ID field of the SPP-Query to which this is a reply.

When a host or security gateway receives an SPP-Reply message it  ${\tt MUST}$  do the following:

Sanchez, Condell

[Page 32]

- 1. Read the Message ID field. If the value does not match the value of an outstanding SPP-Query message from a policy server then:
  - silently discard the message and the event MAY be logged.
- 2. If Message ID matches, Check the timestamp field for anti-replay protection. If a replayed message is detected the message is silently discarded and the event MAY be logged.
- 3. Establish if the message requires signature validation by searching for a Signature payload:
  - if signature validation is required proceed with step 4.
  - if signature validation is not required proceed with step 6.
- 4. Validate the signature on the message.
  - If a certificate record is present, the server MUST process it, however, if the server already has a valid key for the host or server identified in the certificate, the certificate MAY be ignored.
  - Fetch certificate or key corresponding to the IP address found in the sender field of the SPP header.
  - If a certificate or key is not available the entity MAY:
    - o choose, depending on configuration, to abort validation process, discard the message and MAY log the event.
    - o send an SPP-Query message to the source of the IP address found in the sender field of the SPP header with a CERT query payload. Keep SPP-Reply message until the corresponding SPP-Reply returns. Extract certificate or key, validate it and proceed.
  - 5. Once a validated certificate or key is available then validate signature.
    - If validation fails:
    - the message is silently discarded and the event MAY be logged
    - If validation succeeds, continue processing.
- 6. For Policy Servers, validate the chain of trust:
  - For each Policy Server record, verify that the Policy Server is authoritative over the Node. This may be done using information contained in certificates.
  - Use the Policy Server records to Create a chain of hosts from the destination host to this policy server where two records

Sanchez, Condell [Page 33]

- If the chain has no breaks, then this policy server MUST be authoritative over the sender of the reply, otherwise drop the message and stop processing it.
- If the chain has one break, then the last policy server on the chain MUST be the sender of the reply, otherwise drop the message and stop processing it.
- If the chain has two or more breaks, then there is an error in the chain so drop the message and stop processing it.
- Verify that the Policy Server that is authoritative over the destination host is authoritative over ALL destination hosts in any comsec records.
- 7. If MCODE value is 2-7, 9 or 10:

For hosts or security gateways:

- clear all the state and stop processing For policy servers:
- create an SPP-Reply message using the same MCODE value.
- 8. If the reply received correponds to a Cert query and the MCODE is either (1), (8) (accept or partially unavailable) process message that was held up waiting for the cert.
- 9. For Policy Servers:
  - Search the Local Policy Database for a policy entry that matches the policy expressed in Query payload.
  - If the R bit is not set, merge the local and non-local policy information using the policy resolution process outlined in section [\*\*\*] in [SPS].
  - If the R bit is set, include both the policies found in the Local Policy Database and the policies in the reply to send in the new reply.
  - If the merge fails send an SPP-Reply message with MCODE 10 or 7 and cache the failure.
  - If the merge succeeds or the R bit is set:
    - o If the R and C bits are not set, cache policy information. This includes all Record payloads.
    - o send an SPP-Reply message with the resulting policy records and MCODE 1.
    - o If the R and D bits are not set and the merge indicated that policies should be sent to one or more security

gateways, construct a signal for each gateway by creating an SPP-Pol message with the appropriate policy from the merge.

Sanchez, Condell

[Page 34]

- 10. For hosts or security gateways:
  - verify that the information in the Record payload corresponds to the information in the Query payload.
  - extract the records and create a policy entry in the SPD according to local format. The policy is entered in the SPD ONLY if local configuration allows.

### 8.4 Policy Message (SPP-Pol) Processing

When creating a SPP-Pol message, the entity (host, security gateway, or policy server) MUST do the following:

- 1. Generate the Message ID value. This value starts at zero (0) and MUST be incremented by (1) with every new message.
- 2. Set the value of the MTYPE field to 3.
- 3. Set the MCODE value to zero (0).
- 4. Set the RCOUNT field. The QCOUNT field MUST be set to zero (0) since no query payloads exist.
- 5. Set the flag bits accordingly and set the RESERVED field to (0).
- 6. Set the IDENTITY\_TYPE and IDENTITY of the Sender of the SPP message.
- 7. Construct the SPP data payloads. A Record payload MUST be present in this message. Query payloads MUST NOT be present.
- 8. Calculate and place the timestamp value used for anti-replay attack protection.
- 9. If the Signature payload is required for message integrity and authentication, calculate a signature over the SPP header, SPP payloads, the PCL, TYPE, and LENGTH fields of the Signature payload. If required, the Signature payload MUST be the last payload in the message.

When a policy server receives an SPP-Pol message it MUST do the following:

- 1. Check for SPP access control. If enabled, read the IP address in the Sender's field of the SPP header.
  - Verify whether or not the message is allowed. If the message is not allowed then:
    - o send an SPP-Pol\_Ack message with the MCODE 2 or 7

- o discard message and the event MAY be logged.
- If the message is allowed, continue.

Sanchez, Condell

[Page 35]

- 2. Check timestamp field for anti-replay protection. If a replayed message is detected:
  - send an SPP-Pol\_Ack message with the MCODE 3 or 7
  - discard message and the event MAY be logged.
- 3. If the message requires signature validation.
  - If a certificate record is present, the server MUST process it, however, if the server already has a valid key for the host or server identified in the certificate, the certificate MAY be ignored.
  - Fetch certificate or key corresponding to the IP address found in the sender field of the SPP header.
  - If a certificate or key is not available the entity MAY, depending on configuration:
    - o choose to abort validation process, send SPP-Pol\_Ack message with MCODE 5 or 7, discard the message and MAY log the event.
    - o send an SPP-Query message to the source of the IP address found in the sender field of the SPP header with a CERT Query payload. Keep SPP-Pol message until the SPP-Reply returns. Extract certificate or key, validate it and proceed.
  - Once a validated certificate or key is available then validate signature.
    - o If validation fails the message is silently discarded and the event MAY be logged
- 4. If signature was not required or upon successful validation of a signature parse the payloads.
- 5. For hosts and security gateways:
  - extract the records and create a policy entry in the cache database. The policy MAY be entered in the SPD, also, ONLY if configuration allows.
- 6. For Policy Servers:
  - extract the records, find corresponding policies in the server's SPS database, merge the two sets of policies, and create a policy entry in the cache database.

7. Send an SPP-Pol\_Ack message with MCODE 1.

Sanchez, Condell [Page 36]

### 8.5 Policy Acknowledgement Message (SPP-Pol\_Ack) Processing

When creating a SPP-Pol\_Ack message, the policy server MUST do the following:

- 1. Copy the Message ID value from the corresponding SPP-Pol message into the Message ID field.
- 2. Set the value of the MTYPE field to 4
- 3. Set the MCODE value.
- 4. Set the QCOUNT and RCOUNT fields. All fields MUST be set to zero (0).
- 5. Set the flag bits accordingly and set the RESERVED field to (0).
- 6. Set the IDENTITY\_TYPE and IDENTITY of the Sender of the SPP message.
- 7. Query or Record payloads MUST NOT be present.
- 8. Calculate and place the timestamp value used for anti-replay attack protection.
- 9. If the Signature payload is required for message integrity and authentication, calculate a signature over the SPP header, the PCL, TYPE, and LENGTH fields of the Signature payload.

When a host, security gateway, or policy server receives an SPP-Pol\_Ack message the entity MUST do the following:

- 1. Read the Message ID field. If the value does not match the value of an outstanding SPP-Pol message from a policy server then:
  - silently discard the message and the event MAY be logged.
- If Message ID matches then check the timestamp field for anti-replay protection. If a replayed message is detected the message is silently discarded and the event MAY be logged.
- 3. If the message is original (not replayed) and the message requires signature validation then:
  - If a certificate record is present, the server MUST process it, however, if the server already has a valid key for the host or server identified in the certificate, the certificate MAY be ignored.
  - Fetch certificate or key corresponding to the IP address found in the sender field of the SPP header.

- If a certificate or key is not available the entity MAY:

Sanchez, Condell

[Page 37]

- o choose, depending on configuration, to abort validation process, discard the message and MAY log the event.
- o send an SPP-Query message to the source of the IP address found in the sender field of the SPP header with a CERT Query payload. Keep SPP-Pol\_ack message until the SPP-Reply returns. Extract certificate or key, validate it and proceed.
- 4. Once a validated certificate or key is available then validate the signature.
  - If validation fails:
    - the message is silently discarded and the event MAY be logged
  - If validation succeeds:
  - read the MCODE field and proceed accordingly. If no errors, clear all the state for this message and proceed.

### 8.6 Transfer Message (SPP-XFER) Processing

When creating an SPP-Xfer message, the policy server MUST do the following:

- Generate the Message ID value. This value starts at zero (0) and MUST be incremented by (1) with every new message.
- 2. Set the value of the MTYPE field to 5.
- 3. Set the MCODE value to (0).
- 4. Set the RCOUNT field. The QCOUNT field MUST be set to zero (0) since no query payloads exist.
- 5. Set the flag bits accordingly and set the RESERVED field to zero (0).
- 6. Set the IDENTITY\_TYPE and IDENTITY of the Sender of the SPP message.
- 7. Construct the SPP data payloads. A single Transfer Record MUST be present in this payload and MUST contain the master file used to configure this policy server.
- 8. Calculate and place the timestamp value used for anti-replay attack protection.
- 9. If the Signature payload is required for message integrity and authentication, calculate a signature over the SPP header, SPP payloads, the PCL, TYPE, and LENGTH fields of the Signature payload. If required, the Signature payload MUST be the last payload in the message.

Sanchez, Condell

[Page 38]

- 1. Check for SPP access control. If enabled, read the IP address in the Sender's field of the SPP header.
  - Verify whether or not the message is allowed. If the message is not allowed then:
    - o discard message and the event MAY be logged.
  - If the message is allowed, continue.
- 2. Check timestamp field for anti-replay protection. If a replayed message is detected:
  - discard message and the event MAY be logged.
- 3. If the message requires signature validation.
  - If a certificate record is present, the server MUST process it, however, if the server already has a valid key for the host or server identified in the certificate, the certificate MAY be ignored.
  - Fetch certificate or key corresponding to the IP address found in the sender field of the SPP header.
  - If a certificate or key is not available the entity MAY, depending on configuration:
    - o choose to discard the message, and MAY log the event.
    - o send an SPP-Query message to the source of the IP address found in the sender field of the SPP header with a CERT Query payload. Keep the SPP-Query message until the SPP-Reply returns. Extract certificate or key, validate it and proceed.
  - Once a validated certificate or key is available then validate signature.
    - o discard the message, and the event MAY be logged.
- 4. If signature was not required or upon successful validation of a signature parse the payload.
  - extract the Transfer Record and save the master file that it contains.
  - Flush the contents of the SPS database, domain database, and cache.

- Load the new information from the transferred master file into the databases.

Sanchez, Condell

[Page 39]

### **8.7** KeepAlive Message (SPP-KEEP\_ALIVE) Processing

When creating an SPP-Keep\_Alive message, the policy server MUST do the following:

- 1. Generate the Message ID value. This value starts at zero (0) and MUST be incremented by (1) with every new message.
- 2. Set the value of the MTYPE field to 6.
- 3. Set the MCODE value to zero (0).
- 4. Set the QCOUNT and RCOUNT fields. All fields MUST be set to zero (0).
- 5. Set the flag bits accordingly and set the RESERVED field to (0).
- 6. Set the IDENTITY\_TYPE and IDENTITY of the Sender of the SPP message.
- 7. Query or Record payloads MUST NOT be present.
- 8. Calculate and place the timestamp value used for anti-replay attack protection.
- 9. If the Signature payload is required for message integrity and authentication, calculate a signature over the SPP header, the PCL, TYPE, and LENGTH fields of the Signature payload.

When a host or security gateway receives an SPP-Keep\_Alive message it MUST do the following:

- 1. Check for SPP access control. If enabled, read the IP address in the Sender's field of the SPP header.
  - Verify whether or not the message is allowed. If the message is not allowed then discard message and the event MAY be logged.
- 2. Check timestamp field for anti-replay protection. If a replayed message is detected discard message and the event MAY be logged.
- 3. If the message requires signature validation then:
  - If a certificate record is present, the server MUST process it, however, if the server already has a valid key for the host or server identified in the certificate, the certificate MAY be ignored.
  - Fetch certificate or key corresponding to the IP address

found in the sender field of the SPP header.

- If a certificate or key is not available the entity MAY depending on configuration:

Sanchez, Condell

[Page 40]

- o choose to abort validation process, discard the message and the event MAY be logged.
- o send an SPP-Query message to the source of the IP address found in the sender field of the SPP header with a CERT Query payload. Keep SPP-Keep\_Alive message until the SPP-Reply returns. Extract certificate or key, validate it and proceed.
- Once a validated certificate or key is available then validate signature.
  - o If validation fails the message is silently discarded and the event MAY be logged
- 4. If signature validation was not required or upon successful validation of a signature, the event MAY be logged.

### 9. IANA Considerations

This document contains many "magic numbers" to be maintained by the IANA. This section explains the criteria to be used by the IANA to assign numbers beyond the ones defined in this document.

### 9.1 Message Type

The MTYPE field of the SPP Header (<u>section 3.1</u>) defines message exchange types for SPP. Requests for assignment of new message type values 7-250 must be accompanied by a reference to a standards-track or Informational RFC which describes the new message type and how it should be processed. Values 251-255 are for private use.

# 9.2 Message Code

The MCODE field of the SPP Header (<u>section 3.1</u>) defines the acceptable return codes for an SPP message. Requests for assignment of new message code values 12-250 must be accompanied by a description of the conditions under which the code is returned. Values 251-255 are for private use.

## 9.3 Identity Type

The Identity Type field of the SPP Header (<u>section 3.1</u>) defines the acceptable formats for identifying the sender of an SPP message. Requests for assignment of new identity types 4-250 must be accompanied by a description of the format of the corresponding SENDER IDENTITY field in the header. Values 251-255 are for private use.

### 9.4 Payload Class

The first octect of each payload header (<u>section 3.2</u>) defines the type of payload that follows it. Requests for assignment of new message type values 4-250 must be accompanied by a reference to a standards-track or Informational RFC which describes the format of the payload's header and data. Values 251-255 are for private use.

### 9.5 Query Type

The query type (section 3.2.1) defines how the payload data will be interpreted and answered. Requests for assignment of new query type values 4-65000 must be accompanied by a reference to a standards-track or Informational RFC which describes the format of the data and how it should be used. Values 65001-65535 are for private use.

# 9.6 Record Type

The record type (<u>section 3.2.2</u>) defines how the payload data will be interpreted. Requests for assignment of new record type values 4-65000 must be accompanied by a reference to a standards-track or Informational RFC which describes the format of the data and how it should be used. Values 65001-65535 are for private use.

### 9.7 Signature Type

The signature type (section 3.2.3) defines the signature algorithm used to sign the SPP message. Requests for assignment of new signature type values 3-250 must be accompanied by a reference to a standards-track or Informational RFC or a reference to published cryptographic literature which describes this algorithm. Values 251-255 are for private use.

### 9.8 Certificate Type

The Cert Type field of the Certificate query and record (<u>section 3.1</u>) defines the type of certificate requested or included in the payload. Requests for assignment of new certificate types 8-250 must be accompanied by a description of certificate and its encoding. Values 251-255 are for private use.

# 9.9 Certificate Identity Type

The Identity Type and Authority Type fields of the certificate query (section 4.3) define the acceptable formats for identifying the host and its certificate authority for which a certificate is requested. Requests for assignment of new certificate identity types 5-250 must be accompanied by a description of the format of the corresponding IDENTITY and CERTIFICATE AUTHORITY fields in the payload. Values 251-255 are for private use.

#### 9.10 Attribute Data Type

The Data\_Type field of the attribute encoding (<u>section 7</u>) defines the type of attribute included in the data\_value field. Requests for assignment of new attribute data types 30-49 and 52-3200 must be accompanied by a description of the X bit indicating if it is in TLV or TV format, a detailed description of the Data\_Value field corresponding to the attribute type, and in which record and query data fields the type may be used. Values 3200-32767 are for private use.

# 9.11 User Name Type

The Name\_Type field of the user name attribute (section A.16) defines the data in the Name field of the attribute. Requests for assignment of new user name types 2-250 must be accompanied by a description of the corresponding Name field. Values 251-255 are for private use.

#### 9.12 System Name Type

The Name\_Type field of the system name attribute (section A.17) defines the data in the Name field of the attribute. Requests for assignment of new system name types 9-249 must be accompanied by a description of the corresponding Name field. Values 251-255 are for private use.

# 9.13 IPsec Action Attribute

The assigned values of Lifetime\_Type, Cipher\_Alg, Int\_Alg\_Esp, Int\_Alg\_Ah, and Ipcomp\_Alg use the values of their associated fields in [Piper98] and are updated when the IANA updates their values in [Piper98].

The Loc\_Type field of the IPSec action attribute (section A.30) defines the type of location address in the Loc\_Src and Loc\_Dst fields. Requests for assignment of new location types 5-250 must be accompanied by a description of the corresponding Loc\_Src and Loc\_Dst field. Values 251-255 are for private use.

The Loc\_Src and Loc\_Dst fields of the IPSec action attribute (section A.30) may define a general location type. Requests for assignment of new general location values 5-250 must be accompanied by a description of the general location type. Values 251-255 are for private use.

## 9.14 IKE Action Attribute

The assigned values of Group Description, Group\_Type, Auth\_Method, PRF, Lifetime\_Type, Cipher\_Alg, and Hash\_Alg use the values of their associated fields in [Harkins98] and are updated when the IANA updates their values in [Harkins98].

The Mode field of the IKE action attribute (section A.31) defines the IKE Mode. Requests for assignment of new Modes 3-250 must only be done when new modes are added to the IKE protocol. Values 251-255 are for private use.

### 10. Security Considerations

All SPP messages MUST be authenticated to prove which policy server sent the message and that it hasn't been modified en-route. The authentication MAY be provided using the signature payload provided by SPP or some other mechanism such as IPSec.

However, since the policy data may change during SPP exchanges, the messages cannot maintain a signature from every policy server that is involved in the policy exchange. SPP depends on a chain-of-trust for end-to-end authentication. Messages between policy servers are authenticated and contain policy server records and certificates which include proof that the server is authoritative over a set of nodes. The receiving server can use this information to create a chain of servers involved in the policy exchange from itself to the server authoritative over the destination of the query. This chain is used to prove that only servers that have authorization to be involved in the communication were involved.

Policy information may be considered sensitive, since examining policies may expose expoitable weaknesses in the policies. The distribution of policies may be limited to reduce this risk. Policy distribution MAY be limited to those nodes that need to know the information. Limiting distribution any further negates the purpose of the protocol so is not allowed for proper use of SPP.

Additional protections, such as privacy protection, may be desired by some domains. This can be achieved by encrypting SPP data. Encrypting SPP messages is out of scope of this document and may be discussed elsewhere.

SPP uses timestamps to protect against replay attacks. This requires that nodes have adequately synchronized time-of-day clocks. It is necessary to choose an appropriately sized window of time in which timestamps may be accepted. If the window is too small, valid messages may be discarded. On the other hand, if the window is too large it may leave the server open to replay attacks.

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Sanchez, Condell

[Page 44]

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Sanchez, Condell

[Page 45]

#### APPENDIX A

### DATA\_TYPE Definitions

The encoding of each selector and SA attribute is decribed here. Each attribute is described using the following set of data:

X The value of the X bit in the attribute encoding. DATA\_TYPE The value of the DATA\_TYPE field in the attribute

encoding.

LENGTH The length of the data to use if X = 0.

list 'yes' indicates the attribute may be used as a list

as described below.

DATA\_VALUE Encoding of the DATA\_VALUE field in the attribute

encoding.

Attributes generally encode "any" in one of two ways. If using the TLV format (X = 0) then the length is set to 0 to indicate any. If the TV format (X = 1) is used, then the value is set to 0;

Attributes that may be expressed as lists provide the DATA\_VALUE encoding for one element of the list. Multiple list elements may be expressed by concatenating multiple list elements. The LENGTH of attribute is the number of elements present times the length of one list element. Therefore, when reading an attribute that can be expressed as a list, the number of list elements may be determined by dividing the length by the size of a single list element.

The remainder of this appendix describes the values and encoding for each selector and SA attribute specified in  $\frac{1}{2}$ .

### A.1 IPV4\_ADDR

X 0 DATA\_TYPE 1

LENGTH 4 if an IP address is present,

0 if no IP address is present.

list No.

DATA\_VALUE

## **ADDRESS**

#### A.2 IPV6\_ADDR

X 0 DATA\_TYPE 2

LENGTH 16 if an IP address is present,

0 if no IP address is present.

list No

DATA\_VALUE

**ADDRESS** 

An IPV6 address

### A.3 SRC\_IPV4\_ADDR

X 0 DATA\_TYPE 3

LENGTH 4 times the number of addresses in the list.

A length of 0 indicates any address.

list Yes

DATA\_VALUE

SRC ADDRESS

An IPV4 address representing the source host of a  $\operatorname{communication}$ 

### A.4 SRC\_IPV4\_ADDR\_SUBNET

X 0 DATA\_TYPE 4

LENGTH 8 times the number of subnets in the list.

A length of 0 indicates any subnet.

list Yes

DATA\_VALUE

0	1	2	3											
0 1 2 3 4 5	6 7 8 9 0 1 2 3 4 5	5 6 7 8 9 0 1 2 3 4	5 6 7 8 9 0 1											
+-+-+-+-	+-+-+-+-+-	.+-+-+-+-+-+-+-+-	+-+-+-+-+-+-+											
SUBNET ADDRESS														
+-														
1	SUBNET MASK													
+-+-+-+-+-	+-+-+-+-+-	-+-+-+-+-+-+-+-+	+-+-+-+-+-+-+											

### SUBNET ADDRESS

An IPV4 address representing the source subnet of a communication

#### SUBNET MASK

An IPV4 address representing the source subnet mask of a communication

### A.5 SRC\_IPV4\_ADDR\_RANGE

Χ

DATA\_TYPE 5
LENGTH 8 times the number of address ranges in the list.

A length of 0 indicates any address.

list Yes

DATA\_VALUE

0	1	2	3											
0 1 2 3 4	5 6 7 8 9 0 1 2 3 4 5	6 7 8 9 0 1 2 3 4	5 6 7 8 9 0 1											
+-+-+-+-+	-+-+-+-+-+-+-+-+-	+-+-+-+-+-+-+-+-+	-+-+-+-+-+											
LOWER BOUND SRC ADDRESS														
+-+-+-+-+	+-													
1	UPPER BOUND SRC	ADDRESS												
+-+-+-+-+	-+-+-+-+-	+-+-+-+-+-+-+-+	+-											

# LOWER BOUND SRC ADDRESS

An IPV4 address representing the includsive lower-bound of a range of source addresses of a communication.

### UPPER BOUND SRC ADDRESS

An IPV4 address representing the includsive upper-bound of a range of source addresses of a communication.

# A.6 DST\_IPV4\_ADDR

Χ	0
DATA_TYPE	6

4 times the number of addresses in the list. LENGTH

A length of 0 indicates any address.

list Yes

DATA\_VALUE

Sanchez, Condell [Page 48]

1  $\begin{smallmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 0 & 1 \\ \end{smallmatrix}$ DST ADDRESS 

DST ADDRESS

An IPV4 address representing the destination host of a communication

### A.7 DST\_IPV4\_ADDR\_SUBNET

Χ 0 DATA\_TYPE

8 times the number of subnets in the list. LENGTH

A length of 0 indicates any subnet.

list

DATA\_VALUE

1 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 SUBNET ADDRESS SUBNET MASK 

#### SUBNET ADDRESS

An IPV4 address representing the destination subnet of a communication

SUBNET MASK

An IPV4 address representing the destination subnet mask of a communication

# A.8 DST\_IPV4\_ADDR\_RANGE

0 DATA\_TYPE

LENGTH 8 times the number of address ranges in the list.

A length of 0 indicates any address.

list Yes

DATA\_VALUE

0 1 2 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 

	LOWER BOUND DST ADDRESS		
+-+-+-+-+-+-	+-	+-+	
	UPPER BOUND DST ADDRESS		
+-+-+-+-+-	+-	+-+	
Sanchez, Condell		[Page	49]

### LOWER BOUND DST ADDRESS

An IPV4 address representing the includsive lower-bound of a range of destination addresses of a communication.

#### UPPER BOUND DST ADDRESS

An IPV4 address representing the includsive upper-bound of a range of destination addresses of a communication.

### A.9 SRC\_IPV6\_ADDR

X 0 DATA\_TYPE 9

LENGTH 16 times the number of addresses in the list.

A length of 0 indicates any address.

list Yes

DATA\_VALUE

SRC ADDRESS

An IPV6 address representing the source host of a communication

# A.10 SRC\_IPV6\_ADDR\_SUBNET

X 0 DATA\_TYPE 10

LENGTH 32 times the number of subnets in the list.

A length of 0 indicates any subnet.

list Yes

DATA\_VALUE

Sanchez, Condell

[Page 50]

0				1							2	2								3	
0 1 2	3 4 5 6	5 7	8 9	9 0	1	2 3	3 4	5	6 7	8	9 (	9 1	2	3 -	4 5	6	7	7 8	9	0	1
+-													-+								
					5	SUBI	NET														
					Αľ	DDRE	ESS														
+-+-+-	+-																				
					5	SUBI	NET														
						MAS	SK														
+-+-+-	+-+-	-+-+	+	-+	<b>⊢</b> – +	+ <b>-</b> + -	+-	+-+	-+-	+	+-+	-+-	+	- <b>+</b>	-+-	+	+-	-+-	+ - +	+	-+
SUBNET	ADDRESS	6																			
	An IP\				5 ľ	сері	res	ent	ing	tŀ	ne s	sou	rce	e s	ubn	et	(	of :	a		

SUBNET MASK

Χ

An IPV6 address representing the source subnet mask of a communication

# A.11 SRC\_IPV6\_ADDR\_RANGE

0

11

DATA\_TYPE 32 times the number of address ranges in the list. LENGTH A length of 0 indicates any address. list Yes DATA\_VALUE 1  $0\ 1\ 2\ 3\ 4\ 5\ 6\ 7\ 8\ 9\ 0\ 1\ 2\ 3\ 4\ 5\ 6\ 7\ 8\ 9\ 0\ 1\ 2\ 3\ 4\ 5\ 6\ 7\ 8\ 9\ 0\ 1$ LOWER BOUND SRC ADDRESS UPPER BOUND SRC ADDRESS

LOWER BOUND SRC ADDRESS

An IPV6 address representing the includsive lower-bound

Sanchez, Condell [Page 51]

UPPER BOUND SRC ADDRESS

An IPV6 address representing the includsive upper-bound of a range of source addresses of a communication.

### A.12 DST\_IPV6\_ADDR

X 0 DATA\_TYPE 12

LENGTH 16 times the number of addresses in the list.

A length of 0 indicates any address.

list Yes

DATA\_VALUE

DST ADDRESS

An IPV6 address representing the destination host of a  $\operatorname{communication}$ 

# A.13 DST\_IPV6\_ADDR\_SUBNET

X 0 DATA\_TYPE 13

LENGTH 32 times the number of subnets in the list.

A length of 0 indicates any subnet.

list Yes

DATA\_VALUE

### SUBNET ADDRESS

An IPV6 address representing the destination subnet of a communication

#### SUBNET MASK

An IPV6 address representing the destination subnet mask of a communication

### A.14 DST\_IPV6\_ADDR\_RANGE

X 0 DATA\_TYPE 14

LENGTH 32 times the number of address ranges in the list.

A length of 0 indicates any address.

list Yes

DATA\_VALUE

0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 4 5 6 7 8 9 0 1 5 6 7 8 9 0 1 1 2 3 4 5 6 7 8 9 0 1 1 2 3 4 5 6 7 8 9 0

## LOWER BOUND DST ADDRESS

An IPV6 address representing the includsive lower-bound of a range of destination addresses of a communication.

### UPPER BOUND DST ADDRESS

An IPV6 address representing the includsive upper-bound of a range of destination addresses of a communication.

## A.15 DIRECTION

X 1 DATA\_TYPE 15

LENGTH TV attribute, no length

list No

DATA\_VALUE

#### DIRECTION

In/Outbound 0
Inbound 1
Outbound 2

Direction is with respect to the senders interface.

### A.16 USER\_NAME

X 0 DATA\_TYPE 16

LENGTH 1 plus the length of NAME

A length of 0 indicates any name.

list No

DATA\_VALUE

NAME\_TYPE

822\_EMAIL 0 DIST NAME 1

values 2-250 are reserved to IANA. Values 251-255 are for private use among mutually consenting parties.

NAME

Name of type NAME.

[\*\*\*\* probably should describe in more detail \*\*\*\*]

#### A.17 SYSTEM\_NAME

X 0 DATA\_TYPE 17

LENGTH 1 plus the length of NAME

A length of 0 indicates any name.

list No

DATA\_VALUE

Sanchez, Condell

```
1
\begin{smallmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 0 & 1 \\ \end{smallmatrix}
NAME_TYPE |
                         NAME
NAME_TYPE
      DNS_NAME
      DIST_NAME
      822_NAME
      X400_ADDR
      DIR_NAME
      EDI_PARTY_NAME 5
      URI
      IPADDR
                   7
      REGID
      OTHER
                   250
      values 9-249 are reserved to IANA. Values 251-255 are for
      private use among mutually consenting parties.
NAME
      Name of type NAME.
      [***** probably should describe in more detail ****]
A.18 XPORT_PROTOCOL
DATA_TYPE
            18
LENGTH
            1 plus length of pdata
            A length of 0 indicates any address.
            No (see below)
list
DATA_VALUE
                 1
                                 2
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
PDATA
 PTYPE
Describes the rest of the data:
PTYPE
      ANY
      OPAQUE 1
      LIST
             2
      RANGE
             3
PDATA
      Not used if PTYPE is ANY or OPAQUE.
      LIST
```

Sanchez, Condell [Page 55]

The length of pdata to be used as part of the LENGTH field is 1 times the number of elements in the list.

#### RANGE

indicates a range of protocol values whose inclusive lower-bound is LOWER, and inclusive upper-bound is UPPER.

The length of pdata to be used as part of the LENGTH field is 2.

# A.19 SRC\_PORT

X 0 DATA\_TYPE 19

LENGTH 2 times the number of ports in the list.

A length of 0 indicates any port.

list Yes

DATA\_VALUE

PORT

Port that the communication must be initiated with. This may be a list of ports.

# A.20 SRC\_PORT\_DYNAMIC

X 0 DATA\_TYPE 20

LENGTH 4 plus 2 times the number of ports in the list.

A length of 4 indicates any port.

list See Below

DATA\_VALUE

```
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4
```

The use of this attribute indicates that dynamic port allocation is permitted. Communications that are intitiated with any of the ports indicated, may then dynamically allocate any of the ports listed within the LOWER and UPPER BOUNDS, inclusive.

#### DYNAMIC LOWER BOUND

Lower bound of the range of ports that may be dynamically allocated. If this and DYNAMIC UPPER BOUND are both 0, then any port may be dynamically allocated.

### DYNAMIC UPPER BOUND

Upper bound of the range of ports that may be dynamically allocated. If this and DYNAMIC LOWER BOUND are both 0, then any port may be dynamically allocated.

**PORT** 

Port that the communication must be initiated with. This may be a list of ports.

# A.21 DST\_PORT

X 0 DATA\_TYPE 21

LENGTH 2 times the number of ports in the list.

A length of 0 indicates any port.

list Yes

DATA\_VALUE

PORT

Port that the communication must be initiated with. This may be a list of ports.

### A.22 DST\_PORT\_DYNAMIC

X 0 DATA\_TYPE 22 LENGTH 4 plus 2 times the number of ports in the list.

A length of 4 indicates any port.

list See Below

DATA\_VALUE

Sanchez, Condell [Page 57]

0								1										2										3	
0 1	1 2 3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1
+-+	-+-+-	+	+	<del> </del>	<b>-</b> - +	<del> </del>	+		+	+	+	+	+	+	+	+	+	+ - +	<b>-</b> - +	- <del>-</del> +	+	<del> </del>	<del> </del>	+	+	<b>-</b> - +		+	+
	D	ΥN	ΔM:	IC	LC	OWE	R	В	IUC	ND						[	IYC	NAI	110	C	JPF	PEF	R E	301	JNI	)			
+-																													
				PC	DR1	Γ																							~
+-+	-+-+-	+	+	<del> </del>	H – H	<del> </del>	+	<b>-</b> -	+	<b>+</b>	+	+	+	+	<b>+</b>	<del>-</del>	+	<b>+</b> - +	H	<del>-</del>	+	<b>+</b>	<b>+</b>	+	+	H – H	+	+	+

The use of this attribute indicates that dynamic port allocation is permitted. Communications that are intitiated with any of the ports indicated, may then dynamically allocate any of the ports listed within the LOWER and UPPER BOUNDS, inclusive.

#### DYNAMIC LOWER BOUND

Lower bound of the range of ports that may be dynamically allocated. If this and DYNAMIC UPPER BOUND are both 0, then any port may be dynamically allocated.

### DYNAMIC UPPER BOUND

Upper bound of the range of ports that may be dynamically allocated. If this and DYNAMIC LOWER BOUND are both 0, then any port may be dynamically allocated.

**PORT** 

Port that the communication must be initiated with. This may be a list of ports.

# A.23 SEC\_LABELS

X 0 DATA\_TYPE 23

LENGTH Variable.

list No

DATA\_VALUE

#### SECURITY LABEL

[\*\*\*\* ????? \*\*\*\*\*]

## A.24 V6CLASS

X 1 DATA\_TYPE 24

LENGTH TV attribute, no length

list No DATA\_VALUE

Sanchez, Condell [Page 58]

```
Internet Draft Security Policy Protocol
```

July 1999

```
1 2
6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
PADDING | CLASS |
PADDING
   set to 0
CLASS
    class value
A.25 V6FLOW
DATA_TYPE
       25
LENGTH
        4
list
        No
DATA_VALUE
           1
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
PADDING |
                   FLOW
PADDING
   set to 0
FLOW
   set to flow value
A.26 V4T0S
Χ
        1
DATA_TYPE
        26
LENGTH
       TV attribute, no length
list
        No
DATA_VALUE
1 2
6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
PADDING | TOS
```

**PADDING** 

Sanchez, Condell

[Page 59]

TOS

type of service value

### A.27 ACTION

Χ 1 DATA\_TYPE 27

LENGTH TV attribute, no length

list

DATA\_VALUE

1 2 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 ACTION | 

ACTION

Deny Permit 1

# A.28 SRC\_PORT\_RANGE

0 DATA\_TYPE 28

LENGTH 4 times the number of port ranges in the list.

A length of 0 indicates any port.

list Yes

DATA\_VALUE

0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 PORT LOWER BOUND | PORT UPPER BOUND | 

#### PORT LOWER BOUND

Inclusive lower-bound of a range of port numbers that the communication must be initiated with.

### PORT UPPER BOUND

Inclusive upper-bound of a range of port numbers that the communication must be initiated with.

Sanchez, Condell

### A.29 DST\_PORT\_RANGE

X 0 DATA\_TYPE 29

LENGTH 4 times the number of port ranges in the list.

A length of 0 indicates any port.

list Yes

DATA\_VALUE

#### PORT LOWER BOUND

Inclusive lower-bound of a range of port numbers that the communication must be initiated with.

# PORT UPPER BOUND

Inclusive upper-bound of a range of port numbers that the communication must be initiated with.

# A.30 IPSEC\_ACTION

X 0 DATA\_TYPE 50

LENGTH Variable

list Yes

DATA\_VALUE

Sanchez, Condell

[Page 61]

**ESP** 

0	1		2	3	3									
0 1 2 3 4 5 6 7														
+-+-+-+-+-+-+   ESP +-+-+-+-+-+-+-	RESERVED	1	LIFETIME_	TYPE										
1	LIFETIME													
AH	IPCOMP	1	LIFETIME_	TYPE										
1	LI	FETIME				İ								
+-+-+-+-+-+	CIPHER_ALG	1	CIPHER_KEYLE	NGTH										
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-														
+-+-+-+-+-+-+	+-+-+-+-+-   INT_ALG_AH +-+-+-+-+-   IPCOMP_ALG . +-+-+-+-+-   LOC_SRC +-+-+-+-+-   LOC_DST +-+-+-+-+-	+-+-+   +-+-+-+	INT_KEYLEN	+-+-+-+-+- GTH +-+-+-+-+-+- +-+-+-+-+-+-+-+- +-+-+-+-	·+-+ ·+-+ ·+-+									
+-+-+-+-+-	+-+-+-+-+-	+-+-+-	+-+-+-+-+-	+-+-+-+-	.+-+									

This octet indicates if ESP is to be used and in what mode. NOT REQUIRED means that ESP is not necessary but if used it MUST be negotiated based on the parameters defined below. TUNNEL\_MODE or TRANSP\_MODE means that ESP MUST be negotiatiated in this mode. ANY\_MODE means that ESP MUST be negotited and that any mode (Tunnel or transport) will suffice. NOT ALLOWED means that ESP SHOULD not be negotiated

NOT\_REQUIRED 0
TUNNEL\_MODE 1
TRANSP\_MODE 2
TUNNEL\_MODE\_OPT 3
TRANSP\_MODE\_OPT 4
ANY\_MODE 5
NOT\_ALLOWED 6

and it MUST not be part of this SA.

LIFETIME\_TYPE

This 2 octet field indicates type of lifetime.

Sanchez, Condell [Page 62]

RESERVED 0
SECONDS 1
KILOBYTES 2

These values are assigned in section 4.5 of [Piper98] and are updated when those assigned values change.

#### RESERVED

This 1 octet field primarily used for alignment purposes. Its value is always 0.

### LIFETIME

This 4 octet field indicates the SA lifetime. For a given "Lifetime\_Type" the value of the "Lifetime" attribute defines the actual length of the SA life--either a number of seconds, or a number of kilobytes protected. 0 is not used.

ΑН

This octet indicates if AH is to be used and in what mode. NOT REQUIRED means that AH is not necessary but if used it MUST be negotiated based on the parameters defined below. TUNNEL\_MODE or TRANSP\_MODE means that AH MUST be negotiatiated in this mode. ANY\_MODE means that AH MUST be negotited and that any mode (Tunnel or transport) will suffice. NOT ALLOWED means that AH SHOULD not be negotiated and it MUST not be part of this SA.

NOT\_REQUIRED 0
TUNNEL\_MODE 1
TRANSP\_MODE 2
TUNNEL\_MODE\_OPT 3
TRANSP\_MODE\_OPT 4
ANY\_MODE 5
NOT\_ALLOWED 6

#### IPCOMP

This field indicates if IP Compression is to be used. NOT REQUIRED means that IPCOMP is not necessary but if used it MUST be negotiated based on the parameters defined below. REQUIRED means that IPCOMP MUST be negotiated as part of this SA. NOT ALLOWED means that IPCOMP MUST not be part of this SA.

NOT\_REQUIRED 0 REQUIRED 1 NOT ALLOWED 2

# N\_OF\_CIPHERS

This octet indicates the number of CIPHER\_ALG fields in octets that will follow this field and that could be used during an IKE phase 2 negotiation. If the value of the ESP field is (04)hex this field MUST be set to 0.

Sanchez, Condell

[Page 63]

### CIPHER\_ALG

This octet indicates which ciphers should be used for the IKE phase 2 negotiation. If the ANY identifier is used, it MUST be the only identifier in the list, and its meaning does not include the NULL cipher. If the value of the N\_OF\_CIPHERS field is 0 the CIPHER\_ALG, the CIPHER\_KEYLENTH and the ROUNDS fields are ignored.

ANY	0
NULL	1
RFC1829_IV64	2
DES	3
DES3	4
RC5	5
IDEA	6
CAST	7
BLOWFISH	8
3IDEA	9
RFC1829_IV32	10
RC4	11

These values are assigned in section 4.4.4 of [Piper98], with the exception of 0 being defined as ANY, and are updated when those assigned values change.

# CIPHER\_KEYLENGTH

The first octet corresponds to the minimum value and the second octet corresponds to the maximum value. If no range exist the first octet indicates the keylength. The second octet contains a value of (00)hex.

#### ROUNDS

The first octet corresponds to the minimum value and the second octet corresponds to the maximum value. If no range exist the first octet indicates the rounds. The second octet contains a value of (00)hex.

# N\_OF\_INT\_ESP

This octet indicates the number of INTEGRITY\_ALG fields in octets that will follow this field and that could be used during an IKE phase 1 negotiation. If this field is 0 no authentication/integrity is used with ESP.

## INT\_ALG\_ESP

This octet indicates which algorithm should be used for the

IKE phase 2 negotiation. If the ANY identifier is used, it MUST be the only identifier in the list. If the value of the N\_OF\_INT\_ESP field is 0 this INT\_ALG\_ESP and ESP\_INT\_KEYLENGTH are ignored.

Sanchez, Condell

[Page 64]

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Security	POLICY	Protocol

Internet Draft

July 1999

ANY HMAC\_MD5 1 2 HMAC\_SHA1 3 DES\_MAC KPDK 4

These values are assigned in section 4.5 of [Piper98], with the exception of 0 being defined as ANY, and are updated when those assigned values change.

# ESP\_INT\_KEYLENGTH

The first octet corresponds to the minimum value and the second octet corresponds to the maximum value. If no range exist the first octet indicates the keylength. The second octet contains a value of (00)hex.

### N\_OF\_INT\_AH

This octet indicates the number of INTEGRITY\_ALG fields in octets that will follow this field and that could be used during an IKE phase 1 negotiation. If the value of the AH field is (04)hex this field MUST be set to 0.

# INT\_ALG\_AH

This octet indicates which algorithm should be used for the IKE phase 2 negotiation. If the value of the N\_OF\_INT\_AH field is 0 the INT\_ALG\_AH and the INT\_KEYLENGTH fields are ignored.

ANY HMAC MD5 1 HMAC\_SHA1 2 DES\_MAC 3 KPDK

These values are assigned in section 4.5 of [Piper98], with the exception of 0 being defined as ANY, and are updated when those assigned values change.

## INT\_ KEYLENGTH

The first octet corresponds to the minimum value and the second octet corresponds to the maximum value. If no range exist the first octet indicates the keylength. The second octet contains a value of (00)hex.

## N\_OF\_IPCOMP

This octet indicates the number of IPCOMP\_ALG fields in octets that will follow this field and that could be used during an IKE phase 2 negotiation. If the value of the IPCOMP field is (04)hex this field MUST be set to 0.

## IPCOMP\_ALG

This octet indicates which algorithm should be used for the IKE phase 2 negotiation. If the ANY identifier is used, it MUST be the only identifier in the list. If the value of the  $N_0F_IPCOMP$  field is 0 this field is ignored.

ANY 0
0UI 1
DEFLATE 2
LZS 3

These values are assigned in section 4.4.5 of [Piper98], with the exception of 0 being defined as ANY, and are updated when those assigned values change.

## LOC\_TYPE

This 1 octet field indicates the contents of the LOC\_SRC or LOC\_DST field. If this field is 0 then the LOC\_SRC or LOC\_DST will be omitted.

NONE 0
IPv4 address 1
IPv6 address 2
DNS Name 3
General 4

values 5-250 are reserved to IANA. Values 251-255 are for private use among mutually consenting parties.

# LOC\_SRC

Variable length field depending on LOC\_TYPE.

IF LOC\_TYPE is (04) then this field is 1 octet in length an it may only take the following values:

ANY 0
DEST 1
HOST 2
LOCAL-SG 3
REMOTE-SG 4

values 5-250 are reserved to IANA. Values 251-255 are for private use among mutually consenting parties.

Sanchez, Condell

[Page 66]

LOC\_DST

See LOC\_SRC.

# A.31 IKE\_ACTION

X 0 DATA\_TYPE 51

LENGTH Variable

list No

DATA\_VALUE

0	1		2	2		3
0 1 2 3 4 5 6	7 8 9 0 1	2 3 4 5 6	6 7 8 9 0	9 1 2 3 4	5 6 7 8	9 0 1
+-+-+-+-+-+	+-+-+-	+-+-+-+-	-+-+-+	-+-+-+-+	-+-+	+-+-+-+
MODE	PF	-s	GRO	OUP DESCRI	PTION	
+-+-+-+-+-+-+	+-+-+-	+-+-+-+-	-+-+-+	-+-+-+-+	-+-+	+-+-+
1	PRF	1	I	_IFETIME_T	YPE	- 1
+-+-+-+-+-+	+-+-+-	+-+-+-+-	+-+-+-	-+-+-+-+	-+-+	+-+-+-+
1		LIFE	ГІМЕ			- 1
+-+-+-+-+-+-+	+-+-+-	+-+-+-+-	+-+-+-	-+-+-+-+	-+-+	+-+-+
N_OF_AUTH	AUTH_ME	ETHOD				
+-+-+-+-+-+-+	+-+-+-	+-+-+-+-	+-+-+-	-+-+-+-+	-+-+	+-+-+
N_OF_CIPHERS	CIPHER_	_ALG		KEYLENGTH		
+-+-+-+-+-+	+-+-+-	+-+-+-+-	+-+-+-	-+-+-+-+	-+-+	+-+-+-+
N_OF_HASH	HASH_AL	_G				
+-+-+-+-+-+	+-+-+-	+-+-+-+-	+-+-+-	-+-+-+-+	-+-+	+-+-+-+

MODE

This octet indicates the IKE mode of operation.

MAIN 0 AGRESSIVE 1 QUICK 2

values 3-250 are reserved to IANA. Values 251-255 are for private use among mutually consenting parties.

PFS Indicates if PFS is to be used for the SA negotiation.

FALSE 0 TRUE 1

# GROUP DESCRIPTION

This 2 octet field indicates which group should be used during the ISAKMP phase 1 or phase 2 negotiation.

Sanchez, Condell

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TH	ıteri	ıeı	ומוט	L	

# Security Policy Protocol

July 1999

Not Used	0
default 768-bit MODP group	1
alternate 1024-bit MODP group	2
EC2N group on GP[2^155]	3
EC2N group on GP[2^185]	4

These values are assigned in <u>Appendix A</u> of [<u>Harkins98</u>] and are updated when those assigned values change.

### PRF

There are currently no pseudo-random functions defined.

These values are assigned in <u>Appendix A</u> of [<u>Harkins98</u>] and are updated when those assigned values change.

### LIFETIME TYPE

This 2 octet field indicates type of lifetime.

seconds	1
kilobytes	2

These values are assigned in <u>Appendix A</u> of [<u>Harkins98</u>] and are updated when those assigned values change.

## LIFETIME

This 4 octet field indicates the SA lifetime. For a given "Lifetime\_Type" the value of the "Lifetime" attribute defines the actual length of the SA life-- either a number of seconds, or a number of kilobytes protected.

## N\_OF\_AUTH

This octet indicates the number of AUTH\_METHOD fields in octets that will follow this field and that could be used during an IKE phase 1 negotiation.

# AUTH\_METHOD

This octet indicates which authentication methods should be used. The number of auth\_methods that could be used is N\_OF\_AUTH

pre-shared key	1
DSS signatures	2
RSA signatures	3
Encryption with RSA	4
Revised encryption with RSA	5

These values are assigned in Appendix A of [Harkins98]

Sanchez, Condell [Page 68]

## N\_OF\_CIPHERS

This octet indicates the number of CIPHER\_ALG fields in octets that will follow this field and that could be used during an IKE phase 1 negotiation.

### **KEYLENGTH**

The first octet corresponds to the minimum value and the second octet corresponds to the maximum value. If no range exist the first octet indicates the keylength. The second octet contains a value of (00)hex.

### CIPHER\_ALG

This octet indicates which ciphers should be used for the IKE phase 1 negotiation. For IKE phase 2 negotiations this field is ignored. The number of ciphers that could be used is N OF CIPHERS

ANY 0
DES 1
IDEA 2
BLOWFISH 3
RC5 4
DES3 5
CAST 6

These values are assigned in  $\underline{\mathsf{Appendix}\ \mathsf{A}}$  of  $[\underline{\mathsf{Harkins98}}]$ , with the exception of 0 being defined as ANY, and are updated when those assigned values change.

# N\_OF\_HASH

This octet indicates the number of HASH\_ALG fields in octets that will follow this field and that could be used during an IKE phase 1 negotiation.

## HASH\_ALG

This octet indicates which algorithm should be used for the IKE phase 1 negotiation. For IKE phase 2 negotiations this field is ignored.

ANY 0 MD5 1 SHA1 2 TIGER 3

These values are assigned in Appendix A of [Harkins98], with

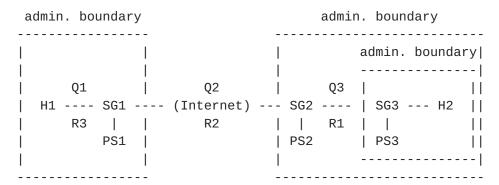
Sanchez, Condell

[Page 69]

#### APPENDIX B

An SPP Example

This appendix provides a detailed example of SPP in use. This example expands on the one provided in section [\*\*\*] of  $[\underline{SPS}]$ .



ESP Tunnel

|========|

ESP Tunnel

ESP Transport

|==| = security association required by policy

---- = connectivity (or if so labeled, administrative boundary)

Hx = host x

SGx = security gateway x

PSx = policy server x

Qx = query x

Rx = reply x

The following entities have these policies for a communication between H1 and H2 for UDP port 79:

H1: requires an ESP Transport SA with H2

PS1: requires an ESP Tunnel SA between SG1 and SG2

PS2: requires an ESP Tunnel SA between SG1 and SG2

PS3: requires an ESP Tunnel SA between H1 and SG3

H2: requires an ESP Transport SA with H1

PS1, PS2, PS3 also have policies allowing ESP to pass through their respective Security Gateways.

- The policy client at H1 is asked for a policy for a communication: H1 to H2 using UDP port 79.
- 2. H1's policy client does not have an answer so it creates an SPP query, Q1:

SPP Header [Query, Sender H1, qcount 1, rcount 2]

Query Payload [comsec]:
 src H1, dst H2, UDP, 79
Record Payload [comsec]:
 src H1, dst H2, UDP, 79, permit

Sanchez, Condell

[Page 70]

Record Payload [SA rec]:
src H1, dst H2, UDP, 79, permit, ESP transport H1->H2
Signature Payload
H1 sends Q1 to PS1, its configured policy server.

3. PS1 receives the query and verifies the signature. Its domain database indicates that it is not authoritative over H2 so it checks its cache to see if it has a cached answer. For this example, it does not, so it creates a new SPP query, Q2, with the query and records formed by merging the local policy with the policy from Q1:

SPP Header [Query, Sender PS1, qcount 1, rcount 3]
Query Payload [comsec]:
 src H1, dst H2, UDP, 79
Record Payload [comsec]:
 src H1, dst H2, UDP, 79, permit
Record Payload [SA rec]:
 src SG1, dst SG2, UDP, 79, permit, ESP tunnel SG1->SG2
Record Payload [SA rec]:
 src H1, dst H2, UDP, 79, permit, ESP transport H1->H2
Signature Payload

PS1 sends Q2 to H2.

- 4. SG2 intercepts Q2 and passes it to PS2.
- 5. PS2 receives the query and verifies the signature. Its domain database indicates that it is not authoritative over H2 so it checks its cache to see if it has a cached answer. For this example, it does not, so it creates a new SPP query, Q3, with the query and records formed by merging the local policy with the policy from Q2:

SPP Header [Query, Sender PS1, qcount 1, rcount 3]
Query Payload [comsec]:
 src H1, dst H2, UDP, 79
Record Payload [comsec]:
 src H1, dst H2, UDP, 79, permit
Record Payload [SA rec]:
 src SG1, dst SG2, UDP, 79, permit, ESP tunnel SG1->SG2
Record Payload [SA rec]:
 src H1, dst H2, UDP, 79, permit, ESP transport H1->H2
Signature Payload

PS2 sends Q3 to H2.

- 6. SG3 intercepts Q3 and passes it to PS3.
- 7. PS3 receives the query and verifies the signature. Its domain database indicates that it is authoritative over H2 so it will send a reply. It checks its cache to see if it has a cached answer. For this example, it does have one cached from previous

information sent to it by H2. PS3 merges the cached policy with the policy it received from Q3. The merge indicates that a signal and a reply will be needed. PS3 caches the merged policy.

Sanchez, Condell

[Page 71]

PS3 creates a reply with the query payload from Q3, the merged policy and policy server and cert records: SPP Header [Reply, Sender PS3, gcount 1, rcount 6] Query Payload [comsec]: src H1, dst H2, UDP, 79 Record Payload [comsec]: src H1, dst H2, UDP, 79, permit Record Payload [SA rec]: src SG1, dst SG2, UDP, 79, permit, ESP tunnel SG1->SG2 Record Payload [SA rec]: src H1, dst SG3, UDP, 79, permit, ESP tunnel H1->SG3 Record Payload [SA rec]: src H1, dst H2, UDP, 79, permit, ESP transport H1->H2 Record Payload [policy server]: policy server PS3, node H1 Record Payload [cert]: cert for PS3

PS3 sends R1 to PS2.

Signature Payload

PS3 creates a signal with a comsec record derived from knowing the traffic that will pass through SG3 and, the part of the merged policy that terminates at SG3:

SPP Header [Pol, Sender PS3, qcount 0, rcount 2]
Record Payload [comsec]:
src H1, dst H2, ESP, OPAQUE, permit
Record Payload [SA rec]:
src H1, dst SG3, UDP, 79, permit, ESP tunnel H1->SG3
Signature Payload
PS3 sends the signal to SG3.

- 8. SG3 receives the signal and verifies the signature. SG3 creates an Ack message to indicate that it has received the policy message: SPP Header [Pol-Ack, Sender SG3, qcount 0, rcount 0] Signature Payload SG3 sends the signal to PS3.
- 9. PS3 receives the Pol-Ack and verifies the signature. PS3 removes the corresponding policy message from its retry queue.
- Meanwhile, PS2 receives the reply R1 and verifies the signature and the chain-of-trust to verify the policy came from a server authoritative for H2. It matches an outstanding query message, so it will send a reply. PS2 merges the policy received in R1 with its local policy and the policy information it received from Q2. The merge indicates that a signal and a reply will be needed. PS2 caches the merged policy.

PS2 creates a reply with the query payload from R1, the merged

Sanchez, Condell [Page 72]

```
SPP Header [Reply, Sender PS2, qcount 1, rcount 8]
Query Payload [comsec]:
  src H1, dst H2, UDP, 79
Record Payload [comsec]:
  src H1, dst H2, UDP, 79, permit
Record Payload [SA rec]:
  src SG1, dst SG2, UDP, 79, permit, ESP tunnel SG1->SG2
Record Payload [SA rec]:
  src H1, dst SG3, UDP, 79, permit, ESP tunnel H1->SG3
Record Payload [SA rec]:
  src H1, dst H2, UDP, 79, permit, ESP transport H1->H2
Record Payload [policy server]:
  policy server PS3, node H1
Record Payload [cert]:
  cert for PS3
Record Payload [policy server]:
  policy server PS2, node PS3
Record Payload [cert]:
  cert for PS2
Signature Payload
```

PS2 sends R2 to PS1.

PS2 creates a signal with a comsec record derived from knowing the traffic that will pass through SG2 and, the part of the merged policy that terminates at SG2:

SPP Header [Pol, Sender PS2, qcount 0, rcount 2]
Record Payload [comsec]:
 src H1, dst SG3, ESP, OPAQUE, permit
Record Payload [SA rec]:
 src SG1, dst SG2, UDP, 79, permit, ESP tunnel SG1->SG2
Signature Payload

PS2 sends the signal to SG2.

- 11. SG2 receives the signal and verifies the signature. SG2 creates an Ack message to indicate that it has received the policy message: SPP Header [Pol-Ack, Sender SG2, qcount 0, rcount 0] Signature Payload SG2 sends the signal to PS2.
- <u>12</u>. **PS2** receives the Pol-Ack and verifies the signature. PS2 removes the corresponding policy message from its retry queue.
- Meanwhile, PS1 receives the reply R2 and verifies the signature and the chain-of-trust to verify the policy came from a server authoritative for H2. R2 matches an outstanding query message, so it will send a reply. PS1 merges the policy received in R2 with its local policy and the policy information it received from Q1. The merge indicates that a signal and a reply will be needed. PS1 caches the merged policy.

PS1 creates a reply with the query payload from R2 and the merged policy. Policy server and cert records are not necessary since PS1 is authoritative for H1:

Sanchez, Condell

[Page 73]

```
SPP Header [Reply, Sender PS1, qcount 1, rcount 3]
Query Payload [comsec]:
    src H1, dst H2, UDP, 79
Record Payload [comsec]:
    src H1, dst H2, UDP, 79, permit
Record Payload [SA rec]:
    src H1, dst SG3, UDP, 79, permit, ESP tunnel H1->SG3
Record Payload [SA rec]:
    src H1, dst H2, UDP, 79, permit, ESP transport H1->H2
    Signature Payload
PS1 sends R3 to H1.
```

PS1 creates a signal with a comsec record derived from knowing the traffic that will pass through SG1 and, the part of the merged policy that terminates at SG1:

SPP Header [Pol, Sender PS1, qcount 0, rcount 2]
Record Payload [comsec]:
src H1, dst SG3, ESP, OPAQUE, permit
Record Payload [SA rec]:
src SG1, dst SG2, UDP, 79, permit, ESP tunnel SG1->SG2
Signature Payload
PS1 sends the signal to SG1.

- 12. SG1 receives the signal and verifies the signature. SG1 creates an Ack message to indicate that it has received the policy message:

  SPP Header [Pol-Ack, Sender SG1, qcount 0, rcount 0]

  Signature Payload

  SG1 sends the signal to PS1.
- 13. PS1 receives the Pol-Ack and verifies the signature. PS1 removes the corresponding policy message from its retry queue.
- 14. Meanwhile, H1 receives the reply R3 and verifies the signature. The client can now use the policy as it is needed.

Sanchez, Condell

[Page 74]

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Sanchez, Condell

[Page 75]