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Transport Layer Security Transport Model for SNMP
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Abstract

This document describes a Transport Model for the Simple Network Management Protocol (SNMP), that uses either the Transport Layer Security protocol or the Datagram Transport Layer Security (DTLS) protocol. The TLS and DTLS protocols provide authentication and privacy services for SNMP applications. This document describes how the TLS Transport Model (TLSTM) implements the needed features of a SNMP Transport Subsystem to make this protection possible in an interoperable way.

This transport model is designed to meet the security and operational needs of network administrators. The TLS mode can make use of TCP's improved support for larger packet sizes and the DTLS mode provides potentially superior operation in environments where a connectionless (e.g. UDP or SCTP) transport is preferred. Both TLS and DTLS integrate well into existing public keying infrastructures.

This document also defines a portion of the Management Information Base (MIB) for use with network management protocols. In particular it defines objects for managing the TLS Transport Model for SNMP.

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1. Introduction

It is important to understand the modular SNMPv3 architecture as defined by [\[RFC3411\]](#) and enhanced by the Transport Subsystem [\[RFC5590\]](#). It is also important to understand the terminology of the SNMPv3 architecture in order to understand where the Transport Model described in this document fits into the architecture and how it interacts with the other architecture subsystems. For a detailed overview of the documents that describe the current Internet-Standard Management Framework, please refer to [Section 7 of \[RFC3410\]](#).

This document describes a Transport Model that makes use of the Transport Layer Security (TLS) [\[RFC5246\]](#) and the Datagram Transport Layer Security (DTLS) Protocol [\[RFC4347\]](#), within a transport subsystem [\[RFC5590\]](#). DTLS is the datagram variant of the Transport Layer Security (TLS) protocol [\[RFC5246\]](#). The Transport Model in this document is referred to as the Transport Layer Security Transport Model (TLSTM). TLS and DTLS take advantage of the X.509 public key infrastructure [\[RFC5280\]](#). This transport model is designed to meet the security and operational needs of network administrators, operate in both environments where a connectionless (e.g. UDP or SCTP) transport is preferred and in environments where large quantities of data need to be sent (e.g. over a TCP based stream). Both TLS and DTLS integrate well into existing public key infrastructures.

This document also defines a portion of the Management Information Base (MIB) for use with network management protocols. In particular it defines objects for managing the TLS Transport Model for SNMP.

For a detailed overview of the documents that describe the current Internet-Standard Management Framework, please refer to [section 7 of RFC \[RFC3410\]](#).

Managed objects are accessed via a virtual information store, termed the Management Information Base or MIB. MIB objects are generally accessed through the Simple Network Management Protocol (SNMP). Objects in the MIB are defined using the mechanisms defined in the Structure of Management Information (SMI). This memo specifies a MIB module that is compliant to the SMIV2, which is described in STD 58: [\[RFC2578\]](#), [\[RFC2579\]](#) and [\[RFC2580\]](#).

The diagram shown below gives a conceptual overview of two SNMP entities communicating using the TLS Transport Model. One entity contains a Command Responder and Notification Originator application, and the other a Command Generator and Notification Responder application. It should be understood that this particular mix of application types is an example only and other combinations are

Network									
\wedge Notifications Commands					\wedge Commands Notifications				
\vee +-----+ +-----+ (D)TLS (D)TLS Service Service (Client) (Server) +-----+ +-----+					\vee +-----+ +-----+ (D)TLS (D)TLS Service Service (Client) (Server) +-----+ +-----+				
\wedge +-----+					\wedge +-----+				
\vee LCD +-----+ +-----+ +-----+ DTLS <-----> Cache TM +-----+ +-----+ Transport Subsystem ^ +-----+ ^ +-----+ v					\vee LCD +-----+ +-----+ +-----+ DTLS <-----> Cache TM +-----+ +-----+ Transport Sub. ^ +-----+ ^ +-----+ v				
+-----+ +-----+ +-----+ Message Sec. Disp. Processing Sub- Subsystem sys. +---+ +-----+ <---> v3MP <---> TSM <--+ +-----+ +-----+ +---+ ^ +-----+ +-----+ +---+-----+ ^ ^ v v					+-----+ +-----+ +-----+ MP Sec. Disp. Sub- Sub- system sys. +---+ +-----+ <---> v3MP <---> TSM <--+ +-----+ +-----+ +---+ ^ +-----+ +-----+ +---+-----+ ^ ^ v v				
+-----+ +-----+ COMMAND NOTIFICATION RESPONDER ORIGINATOR					+-----+ +-----+ COMMAND NOTIFICATION GENERATOR RESPONDER				

application	applications	application	application
+-----+	+-----+	+-----+	+-----+
	SNMP entity		SNMP entity
+-----+	+-----+	+-----+	+-----+

1.1. Conventions

For consistency with SNMP-related specifications, this document favors terminology as defined in STD62 rather than favoring terminology that is consistent with non-SNMP specifications. This is consistent with the IESG decision to not require the SNMPv3 terminology be modified to match the usage of other non-SNMP specifications when SNMPv3 was advanced to Full Standard.

Authentication in this document typically refers to the English meaning of "serving to prove the authenticity of" the message, not data source authentication or peer identity authentication.

Large portions of this document simultaneously refer to both TLS and DTLS when discussing TLSTM components that function equally with either protocol. "(D)TLS" is used in these places to indicate that the statement applies to either or both protocols as appropriate. When a distinction between the protocols is needed they are referred to independently through the use of "TLS" or "DTLS". The Transport Model, however, is named "TLS Transport Model" and refers not to the TLS or DTLS protocol but to the standard defined in this document, which includes support for both TLS and DTLS.

The terms "manager" and "agent" are not used in this document, because in the [RFC 3411](#) architecture [[RFC3411](#)], all SNMP entities have the capability of acting in either manager or agent or in both roles depending on the SNMP application types supported in the implementation. Where distinction is required, the application names of Command Generator, Command Responder, Notification Originator, Notification Receiver, and Proxy Forwarder are used. See "SNMP Applications" [[RFC3413](#)] for further information.

Throughout this document, the terms "client" and "server" are used to refer to the two ends of the (D)TLS transport connection. The client actively opens the (D)TLS connection, and the server passively listens for the incoming (D)TLS connection. Either SNMP entity may act as client or as server.

The User-Based Security Model (USM) [[RFC3414](#)] is a mandatory-to-implement Security Model in STD 62. While (D)TLS and USM frequently refer to a user, the terminology preferred in [RFC3411](#) and in this memo is "principal". A principal is the "who" on whose behalf services are provided or processing takes place. A principal can be,

among other things, an individual acting in a particular role; a set of individuals, with each acting in a particular role; an application or a set of applications, or a combination of these within an administrative domain.

Throughout this document, the term "session" is used to refer to a secure association between two TLS Transport Models that permits the transmission of one or more SNMP messages within the lifetime of the session.

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [[RFC2119](#)].

2. The Transport Layer Security Protocol

(D)TLS provides authentication, data message integrity, and privacy at the transport layer. (See [[RFC4347](#)])

The primary goals of the TLS Transport Model are to provide privacy, source authentication and data integrity between two communicating SNMP entities. The TLS and DTLS protocols provide a secure transport upon which the TLSTM is based. An overview of (D)TLS can be found in section [Appendix A](#). Please refer to [[RFC5246](#)] and [[RFC4347](#)] for complete descriptions of the protocols.

2.1. SNMP requirements of (D)TLS

To properly support the SNMP over TLS Transport Model, the (D)TLS implementation requires the following:

- o The TLS Transport Model SHOULD always use authentication of both the server and the client.
- o At a minimum the TLS Transport Model MUST support authentication of the Command Generator, Notification Originator and Proxy Forwarder principals to guarantee the authenticity of the securityName.
- o The TLS Transport Model SHOULD support the message encryption to protect sensitive data from eavesdropping attacks.

3. How the TLSTM fits into the Transport Subsystem

A transport model is a component of the Transport Subsystem. The TLS Transport Model thus fits between the underlying (D)TLS transport

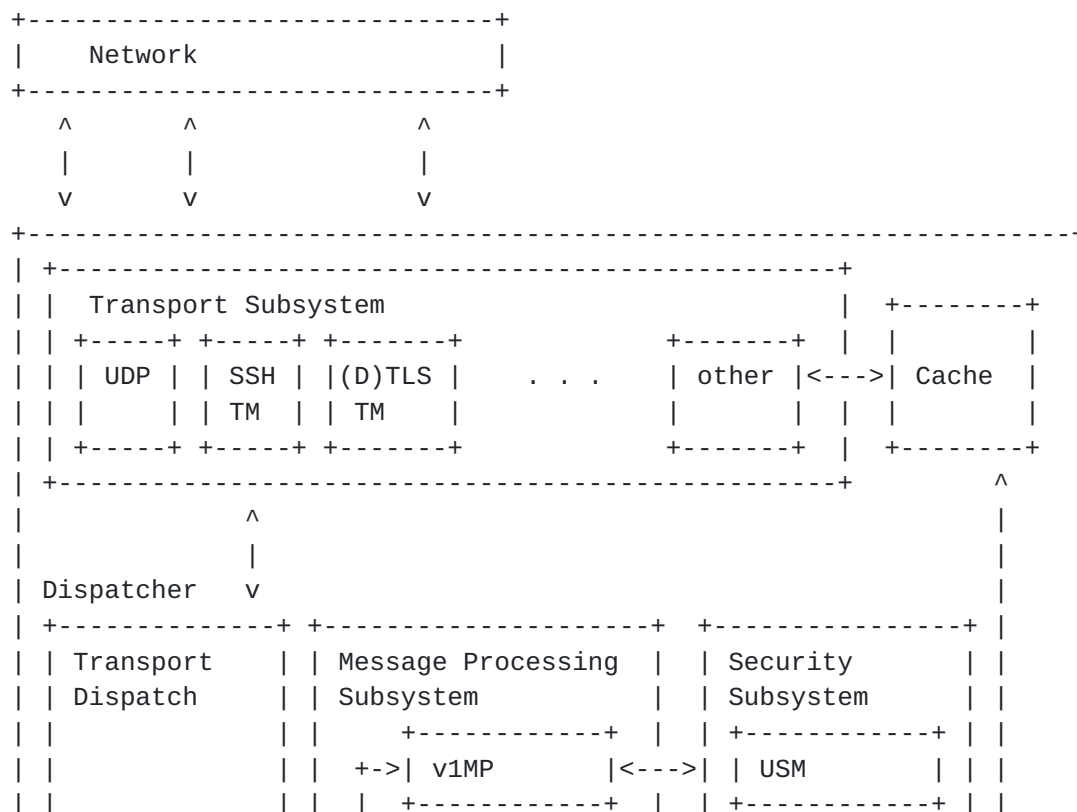
layer and the message dispatcher [[RFC3411](#)] component of the SNMP engine and the Transport Subsystem.

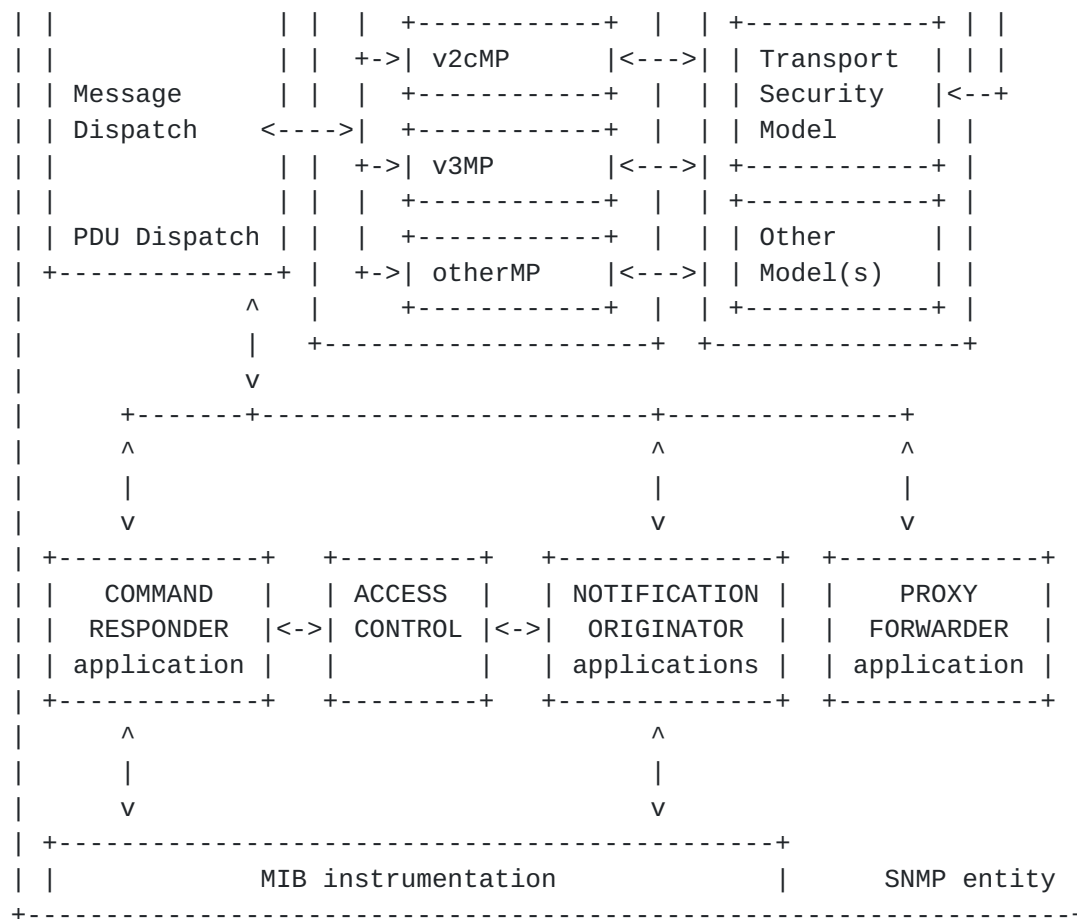
The TLS Transport Model will establish a session between itself and the TLS Transport Model of another SNMP engine. The sending transport model passes unprotected messages from the dispatcher to (D)TLS to be protected, and the receiving transport model accepts decrypted and authenticated/integrity-checked incoming messages from (D)TLS and passes them to the dispatcher.

After a TLS Transport Model session is established, SNMP messages can conceptually be sent through the session from one SNMP message dispatcher to another SNMP message dispatcher. If multiple SNMP messages are needed to be passed between two SNMP applications they SHOULD be passed through the same session. A TLSTM implementation engine MAY choose to close a (D)TLS session to conserve resources.

The TLS Transport Model of an SNMP engine will perform the translation between (D)TLS-specific security parameters and SNMP-specific, model-independent parameters.

The diagram below depicts where the TLS Transport Model fits into the architecture described in [RFC3411](#) and the Transport Subsystem:





3.1. Security Capabilities of this Model

3.1.1. Threats

The TLS Transport Model provides protection against the threats identified by the [RFC 3411](#) architecture [[RFC3411](#)]:

1. Modification of Information - The modification threat is the danger that some unauthorized entity may alter in-transit SNMP messages generated on behalf of an authorized principal in such a way as to effect unauthorized management operations, including falsifying the value of an object.

(D)TLS provides verification that the content of each received message has not been modified during its transmission through the network, data has not been altered or destroyed in an unauthorized manner, and data sequences have not been altered to an extent greater than can occur non-maliciously.

2. Masquerade - The masquerade threat is the danger that management operations unauthorized for a given principal may be attempted by assuming the identity of another principal that has the appropriate authorizations.

The TLSTM provides for authentication of the Command Generator, Command Responder, Notification Generator, Notification Responder and Proxy Forwarder through the use of X.509 certificates.

The masquerade threat can be mitigated against by using an appropriate Access Control Model (ACM) such as the View-based Access Control Module (VACM) [[RFC3415](#)].

3. Message stream modification - The re-ordering, delay or replay of messages can and does occur through the natural operation of many connectionless transport services. The message stream modification threat is the danger that messages may be maliciously re-ordered, delayed or replayed to an extent which is greater than can occur through the natural operation of connectionless transport services, in order to effect unauthorized management operations.

(D)TLS provides replay protection with a MAC that includes a sequence number. Since UDP provides no sequencing ability DTLS uses a sliding window protocol with the sequence number for replay protection (see [[RFC4347](#)]).

4. Disclosure - The disclosure threat is the danger of eavesdropping on the exchanges between SNMP engines.

Symmetric cryptography (e.g., [[AES](#)], [[DES](#)] etc.) can be used by (D)TLS for data privacy. The keys for this symmetric encryption are generated uniquely for each session and are based on a secret negotiated by another protocol (such as the (D)TLS Handshake Protocol).

5. Denial of Service - the [RFC 3411](#) architecture [[RFC3411](#)] states that denial of service (DoS) attacks need not be addressed by an SNMP security protocol. However, datagram-based security protocols like DTLS are susceptible to a variety of denial of service attacks because it is more vulnerable to spoofed messages.

In order to counter these attacks, DTLS borrows the stateless cookie technique used by Photuris [[RFC2522](#)] and IKEv2 [[RFC4306](#)] and is described fully in [section 4.2.1 of \[RFC4347\]](#). This mechanism, though, does not provide any defense against denial of service attacks mounted from valid IP addresses. DTLS Transport

Model server implementations MUST support DTLS cookies.

Implementations are not required to perform the stateless cookie exchange for every DTLS handshakes but in environments where amplification could be an issue or has been detected it is RECOMMENDED that the cookie exchange is utilized.

See [Section 9](#) for more detail on the security considerations associated with the DTLSTM and these security threats.

3.1.2. Message Protection

The [RFC 3411](#) architecture recognizes three levels of security:

- o without authentication and without privacy (noAuthNoPriv)
- o with authentication but without privacy (authNoPriv)
- o with authentication and with privacy (authPriv)

The TLS Transport Model determines from (D)TLS the identity of the authenticated principal, and the type and address associated with an incoming message. The TLS Transport Model provides this information to (D)TLS for an outgoing message.

When an application requests a session for a message, through the cache, the application requests a security level for that session. The TLS Transport Model MUST ensure that the (D)TLS session provides security at least as high as the requested level of security. How the security level is translated into the algorithms used to provide data integrity and privacy is implementation-dependent. However, the NULL integrity and encryption algorithms MUST NOT be used to fulfill security level requests for authentication or privacy. Implementations MAY choose to force (D)TLS to only allow cipher_suites that provide both authentication and privacy to guarantee this assertion.

If a suitable interface between the TLS Transport Model and the (D)TLS Handshake Protocol is implemented to allow the selection of security level dependent algorithms (for example a security level to cipher_suites mapping table) then different security levels may be utilized by the application.

The authentication, integrity and privacy algorithms used by the (D)TLS Protocols may vary over time as the science of cryptography continues to evolve and the development of (D)TLS continues over time. Implementers are encouraged to plan for changes in operator trust of particular algorithms and implementations should offer

configuration settings for mapping algorithms to SNMPv3 security levels.

3.1.3. (D)TLS Sessions

(D)TLS sessions are opened by the TLS Transport Model during the elements of procedure for an outgoing SNMP message. Since the sender of a message initiates the creation of a (D)TLS session if needed, the (D)TLS session will already exist for an incoming message.

Implementations MAY choose to instantiate (D)TLS sessions in anticipation of outgoing messages. This approach might be useful to ensure that a (D)TLS session to a given target can be established before it becomes important to send a message over the (D)TLS session. Of course, there is no guarantee that a pre-established session will still be valid when needed.

DTLS sessions, when used over UDP, are uniquely identified within the TLS Transport Model by the combination of transportDomain, transportAddress, securityName, and requestedSecurityLevel associated with each session. Each unique combination of these parameters MUST have a locally-chosen unique tlsSessionID associated for active sessions. For further information see [Section 4.4](#) and [Section 5](#). TLS and DTLS over SCTP sessions, on the other hand, do not require a unique pairing of attributes since their lower layer protocols (TCP and SCTP) already provide adequate session framing.

3.2. Security Parameter Passing

For the (D)TLS server-side, (D)TLS-specific security parameters (i.e., cipher_suites, X.509 certificate fields, IP address and port) are translated by the TLS Transport Model into security parameters for the TLS Transport Model and security model (i.e., securityLevel, securityName, transportDomain, transportAddress). The transport-related and (D)TLS-security-related information, including the authenticated identity, are stored in a cache referenced by tmStateReference.

For the (D)TLS client-side, the TLS Transport Model takes input provided by the dispatcher in the sendMessage() Abstract Service Interface (ASI) and input from the tmStateReference cache. The (D)TLS Transport Model converts that information into suitable security parameters for (D)TLS and establishes sessions as needed.

The elements of procedure in [Section 5](#) discuss these concepts in much greater detail.

3.3. Notifications and Proxy

(D)TLS sessions may be initiated by (D)TLS clients on behalf of command generators, notification originators or proxy forwarders. Command generators are frequently operated by a human, but notification originators and proxy forwarders are usually unmanned automated processes. The targets to whom notifications should be sent is typically determined and configured by a network administrator.

The SNMP-TARGET-MIB module [[RFC3413](#)] contains objects for defining management targets, including transportDomain, transportAddress, securityName, securityModel, and securityLevel parameters, for Notification Generator, Proxy Forwarder, and SNMP-controllable Command Generator applications. Transport domains and transport addresses are configured in the snmpTargetAddrTable, and the securityModel, securityName, and securityLevel parameters are configured in the snmpTargetParamsTable. This document defines a MIB module that extends the SNMP-TARGET-MIB's snmpTargetParamsTable to specify a (D)TLS client-side certificate to use for the connection.

When configuring a (D)TLS target, the snmpTargetAddrTDomain and snmpTargetAddrTAddress parameters in snmpTargetAddrTable should be set to the snmpTLSDomain, snmpDTLSUDPDDomain, or snmpDTLSSCTPDDomain object and an appropriate snmpTLSAddress, snmpDTLSUDPAddress or snmpDTLSSCTPAddress value respectively. The snmpTargetParamsMPModel column of the snmpTargetParamsTable should be set to a value of 3 to indicate the SNMPv3 message processing model. The snmpTargetParamsSecurityName should be set to an appropriate securityName value and the tlstmParamsClientFingerprint parameter of the tlstmParamsTable should be set a value that refers to a locally held certificate to be used. The tlstmAddrServerFingerprint must be set to a hash value that refers to a locally held copy of the server's presented identity certificate. Other parameters, for example cryptographic configuration such as cipher suites to use, must come from configuration mechanisms not defined in this document. The securityName defined in the snmpTargetParamsSecurityName column will be used by the access control model to authorize any notifications that need to be sent.

4. Elements of the Model

This section contains definitions required to realize the (D)TLS Transport Model defined by this document.

4.1. X.509 Certificates

(D)TLS makes use of X.509 certificates for authentication of both sides of the transport. This section discusses the use of certificates in the TLSTM. A brief overview of X.509 certificate infrastructure can be found in [Appendix B](#).

4.1.1. Provisioning for the Certificate

Authentication using (D)TLS will require that SNMP entities are provisioned with certificates, which are signed by trusted certificate authorities. Furthermore, SNMP entities will most commonly need to be provisioned with root certificates which represent the list of trusted certificate authorities that an SNMP entity can use for certificate verification. SNMP entities SHOULD also be provisioned with a X.509 certificate revocation mechanism which can be used to verify that a certificate has not been revoked. The certificate trust anchors, being either CA certificates or public keys for use by self-signed certificates, must be installed through an out of band trusted mechanism into the server and its authenticity MUST be verified before access is granted.

Having received a certificate, the authenticated tmSecurityName of the principal is looked up using the tlstmCertToSNTTable. This table either:

- o Maps a certificate's fingerprint type and value to a directly specified tmSecurityName, or
- o Identifies a certificate issuer's fingerprint and allows a child certificate's subjectAltName or CommonName to be mapped to the tmSecurityName.

Implementations MAY choose to discard any connections for which no potential tlstmCertToSNTTable mapping exists before performing certificate verification to avoid expending computational resources associated with certificate verification.

The typical enterprise configuration will map a "subjectAltName" component of the tbsCertificate to the TLSTM specific tmSecurityName. The authenticated identity can be obtained by the TLS Transport Model by extracting the subjectAltName(s) from the peer's certificate. The receiving application will then have an appropriate tmSecurityName for use by other SNMPv3 components like an access control model.

An example of this type of mapping setup can be found in [Appendix C](#)

This tmSecurityName may be later translated from a TLSTM specific

tmSecurityName to a SNMP engine securityName by the security model. A security model, like the TSM security model [[RFC5591](#)], may perform an identity mapping or a more complex mapping to derive the securityName from the tmSecurityName offered by the TLS Transport Model.

[4.2.](#) Messages

As stated in [Section 4.1.1 of \[RFC4347\]](#), each DTLS record must fit within a single DTLS datagram. The TLSTM SHOULD prohibit SNMP messages from being sent that exceeds the maximum DTLS message size. The TLSTM implementation SHOULD return an error when the DTLS message size would be exceeded and the message won't be sent.

[4.3.](#) SNMP Services

This section describes the services provided by the (D)TLS Transport Model with their inputs and outputs. The services are between the Transport Model and the dispatcher.

The services are described as primitives of an abstract service interface (ASI) and the inputs and outputs are described as abstract data elements as they are passed in these abstract service primitives.

[4.3.1.](#) SNMP Services for an Outgoing Message

The dispatcher passes the information to the TLS Transport Model using the ASI defined in the transport subsystem:

```
statusInformation =
  sendMessage(
    IN  destTransportDomain      -- transport domain to be used
    IN  destTransportAddress    -- transport address to be used
    IN  outgoingMessage         -- the message to send
    IN  outgoingMessageLength   -- its length
    IN  tmStateReference        -- reference to transport state
  )
```

The abstract data elements returned from or passed as parameters into the abstract service primitives are as follows:

statusInformation: An indication of whether the passing of the message was successful. If not, it is an indication of the problem.

destTransportDomain: The transport domain for the associated **destTransportAddress**. The Transport Model uses this parameter to determine the transport type of the associated **destTransportAddress**. This parameter may also be used by the transport subsystem to route the message to the appropriate Transport Model. This document specifies three TLS and DTLS based Transport Domains for use: the **snmpTLSDomain**, the **snmpDTLSUDPDDomain** and the **snmpDTLSSCTPDDomain**.

destTransportAddress: The transport address of the destination TLS Transport Model in a format specified by the **SnmpTLSAddress**, the **SnmpDTLSUDPAddress** or the **SnmpDTLSSCTPAddress** TEXTUAL-CONVENTIONS.

outgoingMessage: The outgoing message to send to (D)TLS for encapsulation.

outgoingMessageLength: The length of the **outgoingMessage** field.

tmStateReference: A handle/reference to **tmSecurityData** to be used when securing outgoing messages.

4.3.2. SNMP Services for an Incoming Message

The TLS Transport Model processes the received message from the network using the (D)TLS service and then passes it to the dispatcher using the following ASI:

```
statusInformation =
receiveMessage(
IN   transportDomain           -- origin transport domain
IN   transportAddress          -- origin transport address
IN   incomingMessage           -- the message received
IN   incomingMessageLength     -- its length
IN   tmStateReference          -- reference to transport state
)
```

The abstract data elements returned from or passed as parameters into the abstract service primitives are as follows:

statusInformation: An indication of whether the passing of the message was successful. If not, it is an indication of the problem.

transportDomain: The transport domain for the associated **transportAddress**. This document specifies three TLS and DTLS based Transport Domains for use: the **snmpTLSDomain**, the **snmpDTLSUDPDDomain** and the **snmpDTLSSCTPDDomain**.

transportAddress: The transport address of the source of the received message in a format specified by the `SnmpTLSAddress`, the `SnmpDTLSUDPAddress` or the `SnmpDTLSSCTPAddress` TEXTUAL-CONVENTION.

incomingMessage: The whole SNMP message after being processed by (D)TLS and removed of the (D)TLS transport layer data.

incomingMessageLength: The length of the `incomingMessage` field.

tmStateReference: A handle/reference to `tmSecurityData` to be used by the security model.

4.4. (D)TLS Services

This section describes the services provided by the (D)TLS Transport Model with their inputs and outputs. These services are between the TLS Transport Model and the (D)TLS transport layer. The following sections describe services for establishing and closing a session and for passing messages between the (D)TLS transport layer and the TLS Transport Model.

4.4.1. Services for Establishing a Session

The TLS Transport Model provides the following ASI to describe the data passed between the Transport Model and the (D)TLS transport layer for session establishment.

```
statusInformation =          -- errorIndication or success
openSession(
IN  destTransportDomain      -- transport domain to be used
IN  destTransportAddress     -- transport address to be used
IN  securityName             -- on behalf of this principal
IN  securityLevel            -- Level of Security requested
OUT tlsSessionID             -- Session identifier for (D)TLS
)
```

The abstract data elements returned from or passed as parameters into the abstract service primitives are as follows:

statusInformation: An indication of whether the process was successful or not. If not, then the status information will include the error indication provided by (D)TLS.

destTransportDomain: The transport domain for the associated `destTransportAddress`. The TLS Transport Model uses this parameter to determine the transport type of the associated `destTransportAddress`. This document specifies three TLS and DTLS

based Transport Domains for use: the snmpTLSDomain, the snmpDTLSUDPDDomain, and the snmpDTLSSCTPDDomain.

destTransportAddress: The transport address of the destination TLS Transport Model in a format specified by the SnmpTLSAddress, the SnmpDTLSUDPAddress or the SnmpDTLSSCTPAddress TEXTUAL-CONVENTION.

securityName: The security name representing the principal on whose behalf the message will be sent.

securityLevel: The level of security requested by the application.

tlsSessionID: An implementation-dependent session identifier to reference the specific (D)TLS session.

Neither DTLS or UDP provides a session de-multiplexing mechanism and it is possible that implementations will only be able to identify a unique session based on a unique combination of source address, destination address, source UDP port number and destination UDP port number. Because of this, when establishing a new sessions implementations MUST use a different UDP source port number for each connection to a given remote destination IP-address/port-number combination to ensure the remote entity can properly disambiguate between multiple sessions from a host to the same port on a server. TLS and DTLS over SCTP provide session de-multiplexing so this restriction is not needed for TLS or DTLS over SCTP implementations.

The procedural details for establishing a session are further described in [Section 5.3](#).

Upon completion of the process the TLS Transport Model returns status information and, if the process was successful the tlsSessionID for the session. Other implementation-dependent data from (D)TLS may also be returned. The tlsSessionID is formatted and stored in an implementation-dependent manner. It is tied to the tmSecurityData for future use of this session and must remain constant and unique while the session is open.

[4.4.2](#). (D)TLS Services for an Incoming Message

When the TLS Transport Model invokes the (D)TLS record layer to verify proper security for the incoming message, it must use the following ASI:


```
statusInformation =          -- errorIndication or success
tlsRead(
IN   tlsSessionID           -- Session identifier for (D)TLS
IN   wholeTlsMsg             -- as received on the wire
IN   wholeTlsMsgLength       -- length as received on the wire
OUT  incomingMessage        -- the whole SNMP message from (D)TLS
OUT  incomingMessageLength   -- the length of the SNMP message
)
```

The abstract data elements returned from or passed as parameters into the abstract service primitives are as follows:

statusInformation: An indication of whether the process was successful or not. If not, then the status information will include the error indication provided by (D)TLS.

tlsSessionID: An implementation-dependent session identifier to reference the specific (D)TLS session. How the (D)TLS session ID is obtained for each message is implementation-dependent. As an implementation hint for DTLS over UDP, the TLS Transport Model might examine incoming messages to determine the source IP address, source port number, destination IP address, and destination port number and use these values to look up the local `tlsSessionID` in the list of active sessions.

wholeDtlsMsg: The whole message as received on the wire.

wholeDtlsMsgLength: The length of the `wholeDtlsMsg` field.

incomingMessage: The whole SNMP message after being processed by (D)TLS and removed of the (D)TLS transport layer data.

incomingMessageLength: The length of the `incomingMessage` field.

4.4.3. (D)TLS Services for an Outgoing Message

When the TLS Transport Model invokes the (D)TLS record layer to encapsulate and transmit a SNMP message, it must use the following ASI.

```
statusInformation =          -- errorIndication or success
tlsWrite(
IN   tlsSessionID           -- Session identifier for (D)TLS
IN   outgoingMessage        -- the message to send
IN   outgoingMessageLength   -- its length
)
```


The abstract data elements returned from or passed as parameters into the abstract service primitives are as follows:

statusInformation: An indication of whether the process was successful or not. If not, then the status information will include the error indication provided by (D)TLS.

tlsSessionID: An implementation-dependent session identifier to reference the specific (D)TLS session that the message should be sent using.

outgoingMessage: The outgoing message to send to (D)TLS for encapsulation.

outgoingMessageLength: The length of the outgoingMessage field.

4.5. Cached Information and References

When performing SNMP processing, there are two levels of state information that may need to be retained: the immediate state linking a request-response pair, and potentially longer-term state relating to transport and security. "Transport Subsystem for the Simple Network Management Protocol" [[RFC5590](#)] defines general requirements for caches and references.

4.5.1. TLS Transport Model Cached Information

The TLSTM has no specific responsibilities regarding the cached information beyond those discussed in "Transport Subsystem for the Simple Network Management Protocol" [[RFC5590](#)]

5. Elements of Procedure

Abstract service interfaces have been defined by [[RFC3411](#)] and further augmented by [[RFC5590](#)] to describe the conceptual data flows between the various subsystems within an SNMP entity. The TLSTM uses some of these conceptual data flows when communicating between subsystems.

To simplify the elements of procedure, the release of state information is not always explicitly specified. As a general rule, if state information is available when a message gets discarded, the message-state information should also be released. If state information is available when a session is closed, the session state information should also be released. Sensitive information, like cryptographic keys, should be overwritten appropriately first prior to being released.

An error indication in statusInformation will typically include the Object Identifier (OID) and value for an incremented error counter. This may be accompanied by the requested securityLevel and the tmStateReference. Per-message context information is not accessible to Transport Models, so for the returned counter OID and value, contextEngine would be set to the local value of snmpEngineID and contextName to the default context for error counters.

5.1. Procedures for an Incoming Message

This section describes the procedures followed by the (D)TLS Transport Model when it receives a (D)TLS protected packet. The steps are broken into two different sections. [Section 5.1.1](#) describes the needed steps for de-multiplexing multiple DTLS sessions, which is specifically needed for DTLS over UDP sessions. [Section 5.1.2](#) describes the steps specific to transport processing once the (D)TLS processing has been completed. It is assumed that TLS and DTLS/SCP protocol implementations already provide appropriate message demultiplexing and only the processing steps in [Section 5.1.2](#) are needed.

5.1.1. DTLS Processing for Incoming Messages

DTLS is significantly different in terms of session handling than when TLS or DTLS is run over session based streaming protocols like TCP or SCTP. Specifically, the DTLS protocol, when run over UDP, does not have a session identifier that allows implementations to determine through what session a packet arrived. DTLS over SCTP and TLS over TCP streams have built in session demultiplexing and thus the steps in this section are not necessary for those protocol combinations. It is always critical, however, that implementations be able to derive a tlsSessionID from any session demultiplexing process.

A process for demultiplexing multiple DTLS sessions arriving over UDP must be incorporated into the procedures for processing an incoming message. The steps in this section describe one possible method to accomplish this, although any implementation dependent method should be suitable as long as the results are consistently deterministic. The important output results from the steps in this process are the transportDomain, the transportAddress, the wholeMessage, the wholeMessageLength, and a unique implementation-dependent session identifier (tlsSessionID).

This demultiplexing procedure assumes that upon session establishment an entry in a local transport mapping table is created in the Transport Model's Local Configuration Datastore (LCD). The transport mapping table's entry should map a unique combination of the remote

address, remote port number, local address and local port number to an implementation-dependent `tlsSessionID`.

- 1) The TLS Transport Model examines the raw UDP message, in an implementation-dependent manner. If the message is not a DTLS message then it should be discarded. If the message is not a (D)TLS Application Data message (such as a session initialization or session modification message) then the message should be processed by the underlying DTLS framework and no further steps below should be taken by the DTLS Transport.
- 2) The TLS Transport Model queries the LCD using the transport parameters (source and destination addresses and ports) to determine if a session already exists and its `tlsSessionID`.
- 3) If a matching entry in the LCD does not exist then the message is discarded. Increment the `tlstmSessionNoAvailableSessions` counter and stop processing the message.

Note that an entry would already exist if the client and server's session establishment procedures had been successfully completed (as described both above and in [Section 5.3](#)) even if no message had yet been sent through the newly established session. An entry may not exist, however, if a "rogue" message was routed to the SNMP entity by mistake. An entry might also be missing because of a "broken" session (see operational considerations).

- 4) Retrieve the `tlsSessionID` from the LCD.
- 5) The `tlsWholeMsg`, and the `tlsSessionID` are passed to DTLS for integrity checking and decryption using the `tlsRead()` ASI.
- 6) If the message fails integrity checks or other (D)TLS security processing then increment the `tlstmDTLSProtectionErrors` counter, discard and stop processing the message.
- 7) The output of the `tlsRead` ASI results in an `incomingMessage` and an `incomingMessageLength`. These results and the `tlsSessionID` are used below in the [Section 5.1.2](#) to complete the processing of the incoming message.

[5.1.2](#). Transport Processing for Incoming Messages

The procedures in this section describe how the TLS Transport Model should process messages that have already been properly extracted from the (D)TLS stream.

Create a `tmStateReference` cache for the subsequent reference and

assign the following values within it:

tmTransportDomain = snmpTLSDomain, snmpDTLSUDPDDomain or
snmpDTLSSCTPDDomain as appropriate.

tmTransportAddress = The address the message originated from,
determined in an implementation dependent way.

tmSecurityLevel = The derived tmSecurityLevel for the session, as
discussed in [Section 3.1.2](#) and [Section 5.3](#).

tmSecurityName = The derived tmSecurityName for the session as
discussed in [Section 5.3](#). This value MUST be constant during the
lifetime of the (D)TLS session.

tmSessionID = The tlsSessionID, which MUST be a unique session
identifier for this (D)TLS session. The contents and format of
this identifier are implementation dependent as long as it is
unique to the session. A session identifier MUST NOT be reused
until all references to it are no longer in use. The tmSessionID
is equal to the tlsSessionID discussed in [Section 5.1.1](#).
tmSessionID refers to the session identifier when stored in the
tmStateReference and tlsSessionID refers to the session identifier
when stored in the LCD. They MUST always be equal when processing
a given session's traffic.

The wholeMessage and the wholeMessageLength are assigned values from
the incomingMessage and incomingMessageLength values from the (D)TLS
processing.

The TLS Transport Model passes the transportDomain, transportAddress,
wholeMessage, and wholeMessageLength to the dispatcher using the
receiveMessage ASI:

```
statusInformation =  
receiveMessage(  
IN    transportDomain    -- snmpTLSDomain, snmpDTLSUDPDDomain,  
                                -- or snmpDTLSSCTPDomain  
IN    transportAddress    -- address for the received message  
IN    wholeMessage        -- the whole SNMP message from (D)TLS  
IN    wholeMessageLength  -- the length of the SNMP message  
IN    tmStateReference    -- transport info  
)
```


5.2. Procedures for an Outgoing Message

The dispatcher sends a message to the TLS Transport Model using the following ASI:

```
statusInformation =
sendMessage(
  IN  destTransportDomain      -- transport domain to be used
  IN  destTransportAddress     -- transport address to be used
  IN  outgoingMessage          -- the message to send
  IN  outgoingMessageLength    -- its length
  IN  tmStateReference         -- transport info
)
```

This section describes the procedure followed by the TLS Transport Model whenever it is requested through this ASI to send a message.

- 1) Extract the tmSessionID, tmTransportAddress, tmSecurityName, tmRequestedSecurityLevel, and tmSameSecurity values from the tmStateReference. Note: The tmSessionID value may be undefined if no session exists yet over which the message can be sent.
- 2) If tmSameSecurity is true and either tmSessionID is undefined or refers to a session that is no longer open then increment the tlstmSessionNoAvailableSessions counter, discard the message and return the error indication in the statusInformation. Processing of this message stops.
- 3) If tmSameSecurity is false and tmSessionID refers to a session that is no longer available then an implementation SHOULD open a new session using the openSession() ASI (described in greater detail in step 4b). An implementation MAY choose to return an error to the calling module and stop processing of the message.
- 4) If tmSessionID is undefined, then use tmTransportAddress, tmSecurityName and tmRequestedSecurityLevel to see if there is a corresponding entry in the LCD suitable to send the message over.
 - 4a) If there is a corresponding LCD entry, then this session will be used to send the message.
 - 4b) If there is not a corresponding LCD entry, then open a session using the openSession() ASI (discussed further in [Section 4.4.1](#)). Implementations MAY wish to offer message buffering to prevent redundant openSession() calls for the same cache entry. If an error is returned from OpenSession(), then discard the message, increment the tlstmSessionOpenErrors, return an error indication to the

calling module and stop processing of the message.

- 5) Using either the session indicated by the tmSessionID if there was one or the session resulting from a previous step (3 or 4), pass the outgoingMessage to (D)TLS for encapsulation and transmission.

5.3. Establishing a Session

The TLS Transport Model provides the following primitive to establish a new (D)TLS session (previously discussed in [Section 4.4.1](#)):

```
statusInformation =          -- errorIndication or success
openSession(
IN   destTransportDomain    -- transport domain to be used
IN   destTransportAddress   -- transport address to be used
IN   securityName           -- on behalf of this principal
IN   securityLevel          -- Level of Security requested
OUT  tlsSessionID          -- Session identifier for (D)TLS
)
```

The following describes the procedure to follow when establishing a SNMP over (D)TLS session between SNMP engines for exchanging SNMP messages. This process is followed by any SNMP engine establishing a session for subsequent use.

This MAY be done automatically for SNMP messages which are not Response or Report messages.

- 1) The client selects the appropriate certificate and cipher_suites for the key agreement based on the tmSecurityName and the tmRequestedSecurityLevel for the session. For sessions being established as a result of a SNMP-TARGET-MIB based operation, the certificate will potentially have been identified via the tlstmParamsTable mapping and the cipher_suites will have to be taken from system-wide or implementation-specific configuration. Otherwise, the certificate and appropriate cipher_suites will need to be passed to the openSession() ASI as supplemental information or configured through an implementation-dependent mechanism. It is also implementation-dependent and possibly policy-dependent how tmRequestedSecurityLevel will be used to influence the security capabilities provided by the (D)TLS session. However this is done, the security capabilities provided by (D)TLS MUST be at least as high as the level of security indicated by the tmRequestedSecurityLevel parameter. The actual security level of the session is reported in the tmStateReference cache as tmSecurityLevel. For (D)TLS to provide

strong authentication, each principal acting as a Command Generator SHOULD have its own certificate.

- 2) Using the destTransportDomain and destTransportAddress values, the client will initiate the (D)TLS handshake protocol to establish session keys for message integrity and encryption.

If the attempt to establish a session is unsuccessful, then tlstmSessionOpenErrors is incremented, an error indication is returned, and processing stops.

- 3) Once a (D)TLS secured session is established and both sides have performed any appropriate certificate authentication verification (e.g. [\[RFC5280\]](#)) then each side will determine and/or check the identity of the remote entity using the procedures described below.

- a) The (D)TLS server side of the connection identifies the authenticated identity from the (D)TLS client's principal certificate using configuration information from the tlstmCertToSNTTable mapping table. The resulting derived securityName is recorded in the tmStateReference cache as tmSecurityName. The details of the lookup process are fully described in the DESCRIPTION clause of the tlstmCertToSNTTable MIB object. If any verification fails in any way (for example because of failures in cryptographic verification or because of the lack of an appropriate row in the tlstmCertToSNTTable) then the session establishment MUST fail, the tlstmSessionInvalidClientCertificates object is incremented and processing stops.
- b) The (D)TLS client side of the connection MUST verify that authenticated identity of the (D)TLS server's certificate is the certificate expected. This can be done using the configuration fingerprints found in the tlstmAddrTable if the client is establishing the connection based on SNMP-TARGET-MIB configuration or based on external certificate path validation processes (e.g. [\[RFC5280\]](#)).

Methods for verifying that the proper destination was reached based on the presented certificate are described in [\[I-D.saintandre-tls-server-id-check\]](#). Matching the server's naming against SubjectAltName extension values SHOULD be the preferred mechanism for comparison, but matching the CommonName MAY be used.

(D)TLS provides assurance that the authenticated identity has been signed by a trusted configured certificate authority.

If verification of the server's certificate fails in any way (for example because of failures in cryptographic verification or the presented identity was not the expected identity) then the session establishment MUST fail, the `tlstmSessionInvalidServerCertificates` object is incremented and processing stops.

- 4) The (D)TLS-specific session identifier is passed to the TLS Transport Model and associated with the `tmStateReference` cache entry to indicate that the session has been established successfully and to point to a specific (D)TLS session for future use.

(D)TLS provides no explicit manner for transmitting an identity the client wishes to connect to during or prior to key exchange to facilitate certificate selection at the server (e.g. at a Notification Receiver). I.E., there is no available mechanism for sending notifications to a specific principal at a given TCP, UDP or SCTP port. Therefore, an implementation that wishes to support multiple identities MAY use separate TCP, UDP or SCTP port numbers to indicate the desired principal or some other implementation-dependent solution.

5.4. Closing a Session

The TLS Transport Model provides the following primitive to close a session:

```
statusInformation =  
closeSession(  
IN  tmStateReference      -- transport info  
)
```

The following describes the procedure to follow to close a session between a client and server. This process is followed by any SNMP engine closing the corresponding SNMP session.

- 1) Look up the session in the cache and the LCD using the `tmStateReference` and its contents.
- 2) If there is no session open associated with the `tmStateReference`, then `closeSession` processing is completed.
- 3) Delete the entry from the cache and any other implementation-dependent information in the LCD.

- 4) Have (D)TLS close the specified session. This SHOULD include sending a close_notify TLS Alert to inform the other side that session cleanup may be performed.

6. MIB Module Overview

This MIB module provides management of the TLS Transport Model. It defines needed textual conventions, statistical counters, notifications and configuration infrastructure necessary for session establishment. Example usage of the configuration tables can be found in [Appendix C](#).

6.1. Structure of the MIB Module

Objects in this MIB module are arranged into subtrees. Each subtree is organized as a set of related objects. The overall structure and assignment of objects to their subtrees, and the intended purpose of each subtree, is shown below.

6.2. Textual Conventions

Generic and Common Textual Conventions used in this module can be found summarized at <http://www.ops.ietf.org/mib-common-tcs.html>

This module defines the following new Textual Conventions:

- o New TransportDomain and TransportAddress formats for describing (D)TLS connection addressing requirements.
- o Public certificate fingerprint allowing MIB module objects to generically refer to a stored X.509 certificate using a cryptographic hash as a reference pointer.

6.3. Statistical Counters

The TLSTM-MIB defines some statical counters that can provide network managers with feedback about (D)TLS session usage and potential errors that a MIB-instrumented device may be experiencing.

6.4. Configuration Tables

The TLSTM-MIB defines configuration tables that a manager can use for configuring a MIB-instrumented device for sending and receiving SNMP messages over (D)TLS. In particular, there is are MIB tables that extend the SNMP-TARGET-MIB for configuring (D)TLS certificate usage and a MIB table for mapping incoming (D)TLS client certificates to SNMPv3 securityNames.

6.4.1. Notifications

The TLSTM-MIB defines notifications to alert management stations when a (D)TLS connection fails because the server's presented certificate did not meet an expected value, according to the `tlstmAddrTable`.

6.5. Relationship to Other MIB Modules

Some management objects defined in other MIB modules are applicable to an entity implementing the TLS Transport Model. In particular, it is likely that an entity implementing the TLSTM-MIB will implement the SNMPv2-MIB [[RFC3418](#)], the SNMP-FRAMEWORK-MIB [[RFC3411](#)], the SNMP-TARGET-MIB [[RFC3413](#)], the SNMP-NOTIFICATION-MIB [[RFC3413](#)] and the SNMP-VIEW-BASED-ACM-MIB [[RFC3415](#)].

The TLSTM-MIB module contained in this document is for managing TLS Transport Model information.

6.5.1. MIB Modules Required for IMPORTS

The TLSTM-MIB module imports items from SNMPv2-SMI [[RFC2578](#)], SNMPv2-TC [[RFC2579](#)], SNMP-FRAMEWORK-MIB [[RFC3411](#)], SNMP-TARGET-MIB [[RFC3413](#)] and SNMPv2-CONF [[RFC2580](#)].

7. MIB Module Definition

```
TLSTM-MIB DEFINITIONS ::= BEGIN
```

```
IMPORTS
```

```
    MODULE-IDENTITY, OBJECT-TYPE,
    OBJECT-IDENTITY, snmpModules, snmpDomains,
    Counter32, Unsigned32, NOTIFICATION-TYPE
    FROM SNMPv2-SMI
    TEXTUAL-CONVENTION, TimeStamp, RowStatus, StorageType
    FROM SNMPv2-TC
    MODULE-COMPLIANCE, OBJECT-GROUP, NOTIFICATION-GROUP
    FROM SNMPv2-CONF
    SnmpAdminString
    FROM SNMP-FRAMEWORK-MIB
    snmpTargetParamsName, snmpTargetAddrName
    FROM SNMP-TARGET-MIB
    ;
```

```
tlstmMIB MODULE-IDENTITY
```

```
    LAST-UPDATED "200807070000Z"
    ORGANIZATION "ISMS Working Group"
```


CONTACT-INFO "WG-EMail: isms@lists.ietf.org
Subscribe: isms-request@lists.ietf.org

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"

DESCRIPTION "

The TLS Transport Model MIB

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This version of this MIB module is part of RFC XXXX;
see the RFC itself for full legal notices."

-- NOTE to RFC editor: replace XXXX with actual RFC number
-- for this document and remove this note

REVISION "200807070000Z"

DESCRIPTION "The initial version, published in RFC XXXX."

-- NOTE to RFC editor: replace XXXX with actual RFC number
-- for this document and remove this note

::= { snmpModules xxxx }

-- RFC Ed.: replace xxxx with IANA-assigned number and
-- remove this note

-- *****

-- subtrees of the TLSTM-MIB

-- *****

tlstmNotifications OBJECT IDENTIFIER ::= { tlstmMIB 0 }

tlstmObjects OBJECT IDENTIFIER ::= { tlstmMIB 1 }

tlstmConformance OBJECT IDENTIFIER ::= { tlstmMIB 2 }

-- *****

-- tlstmObjects - Objects

-- *****

snmpTLSDomain OBJECT-IDENTITY

STATUS current

DESCRIPTION

"The SNMP over TLS transport domain. The corresponding

transport address is of type SnmpTLSAddress.

The securityName prefix to be associated with the snmpTLSDomain is 'tls'. This prefix may be used by security models or other components to identify which secure transport infrastructure authenticated a securityName."

::= { snmpDomains xx }

-- RFC Ed.: replace xx with IANA-assigned number and
-- remove this note

-- RFC Ed.: replace 'tls' with the actual IANA assigned prefix string
-- if 'tls' is not assigned to this document.

snmpDTLSUDPDDomain OBJECT-IDENTITY

STATUS current

DESCRIPTION

"The SNMP over DTLS/UDP transport domain. The corresponding transport address is of type SnmpDTLSUDPAddress.

When an SNMP entity uses the snmpDTLSUDPDDomain transport model, it must be capable of accepting messages up to the maximum MTU size for an interface it supports, minus the needed IP, UDP, DTLS and other protocol overheads.

The securityName prefix to be associated with the snmpDTLSUDPDDomain is 'dudp'. This prefix may be used by security models or other components to identify which secure transport infrastructure authenticated a securityName."

::= { snmpDomains yy }

-- RFC Ed.: replace yy with IANA-assigned number and
-- remove this note

-- RFC Ed.: replace 'dudp' with the actual IANA assigned prefix string
-- if 'dtls' is not assigned to this document.

snmpDTLSSCTPDDomain OBJECT-IDENTITY

STATUS current

DESCRIPTION

"The SNMP over DTLS/SCTP transport domain. The corresponding transport address is of type SnmpDTLSSCTPAddress.

The securityName prefix to be associated with the

snmpDTLSSCTPDomain is 'dsct'. This prefix may be used by security models or other components to identify which secure transport infrastructure authenticated a securityName."

::= { snmpDomains zz }

-- RFC Ed.: replace zz with IANA-assigned number and
-- remove this note

-- RFC Ed.: replace 'dsct' with the actual IANA assigned prefix string
-- if 'dtls' is not assigned to this document.

SnmpTLSAddress ::= TEXTUAL-CONVENTION

DISPLAY-HINT "1a"

STATUS current

DESCRIPTION

"Represents a TCP connection address for an IPv4 address, an IPv6 address or an US-US-ASCII encoded hostname and port number.

An IPv4 address must be in dotted decimal format followed by a colon ':' (US-ASCII character 0x3A) and a decimal port number in US-ASCII.

An IPv6 address must be a colon separated format, surrounded by square brackets ('[', US-ASCII character 0x5B, and ']', US-ASCII character 0x5D), followed by a colon ':' (US-ASCII character 0x3A) and a decimal port number in US-ASCII.

A hostname is always in US-US-ASCII (as per [RFC1033](#)); internationalized hostnames are encoded in US-US-ASCII as specified in [RFC 3490](#). The hostname is followed by a colon ':' (US-US-ASCII character 0x3A) and a decimal port number in US-US-ASCII. The name SHOULD be fully qualified whenever possible.

Values of this textual convention may not be directly usable as transport-layer addressing information, and may require run-time resolution. As such, applications that write them must be prepared for handling errors if such values are not supported, or cannot be resolved (if resolution occurs at the time of the management operation).

The DESCRIPTION clause of TransportAddress objects that may have snmpTLSAddress values must fully describe how (and when) such names are to be resolved to IP addresses and vice versa.

This textual convention SHOULD NOT be used directly in object definitions since it restricts addresses to a specific format. However, if it is used, it MAY be used either on its own or in conjunction with TransportAddressType or TransportDomain as a pair.

When this textual convention is used as a syntax of an index object, there may be issues with the limit of 128 sub-identifiers specified in SMIV2 (STD 58). It is RECOMMENDED that all MIB documents using this textual convention make explicit any limitations on index component lengths that management software must observe. This may be done either by including SIZE constraints on the index components or by specifying applicable constraints in the conceptual row DESCRIPTION clause or in the surrounding documentation."

REFERENCE

"[RFC 1033](#): DOMAIN ADMINISTRATORS OPERATIONS GUIDE
[RFC 3490](#): Internationalizing Domain Names in Applications
[RFC 3986](#): Uniform Resource Identifier (URI): Generic Syntax
[RFC 5246](#): The Transport Layer Security (TLS) Protocol Version 1.2
"

SYNTAX OCTET STRING (SIZE (1..255))

SnmpDTLSUDPAddress ::= TEXTUAL-CONVENTION

DISPLAY-HINT "1a"

STATUS current

DESCRIPTION

"Represents a UDP connection address for an IPv4 address, an IPv6 address or an US-ASCII encoded hostname and port number.

An IPv4 address must be a dotted decimal format followed by a colon ':' (US-ASCII character 0x3A) and a decimal port number in US-ASCII.

An IPv6 address must be a colon separated format, surrounded by square brackets ('[', US-ASCII character 0x5B, and ']', US-ASCII character 0x5D), followed by a colon ':' (US-ASCII character 0x3A) and a decimal port number in US-ASCII.

A hostname is always in US-US-ASCII (as per [RFC1033](#)); internationalized hostnames are encoded in US-US-ASCII as specified in [RFC 3490](#). The hostname is followed by a colon ':' (US-US-ASCII character 0x3A) and a decimal port number in US-US-ASCII. The name SHOULD be fully qualified whenever possible.

Values of this textual convention may not be directly usable as transport-layer addressing information, and may require

run-time resolution. As such, applications that write them must be prepared for handling errors if such values are not supported, or cannot be resolved (if resolution occurs at the time of the management operation).

The DESCRIPTION clause of TransportAddress objects that may have snmpDTLSUDPAddress values must fully describe how (and when) such names are to be resolved to IP addresses and vice versa.

This textual convention SHOULD NOT be used directly in object definitions since it restricts addresses to a specific format. However, if it is used, it MAY be used either on its own or in conjunction with TransportAddressType or TransportDomain as a pair.

When this textual convention is used as a syntax of an index object, there may be issues with the limit of 128 sub-identifiers specified in SMIV2 (STD 58). It is RECOMMENDED that all MIB documents using this textual convention make explicit any limitations on index component lengths that management software must observe. This may be done either by including SIZE constraints on the index components or by specifying applicable constraints in the conceptual row DESCRIPTION clause or in the surrounding documentation."

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[RFC 3490](#): Internationalizing Domain Names in Applications
[RFC 3986](#): Uniform Resource Identifier (URI): Generic Syntax
[RFC 4347](#): Datagram Transport Layer Security
[RFC 5246](#): The Transport Layer Security (TLS) Protocol Version 1.2

"

SYNTAX OCTET STRING (SIZE (1..255))

SnmpDTLSSCTPAddress ::= TEXTUAL-CONVENTION

DISPLAY-HINT "1a"

STATUS current

DESCRIPTION

"Represents a SCTP connection address for an IPv4 address, an IPv6 address or an US-ASCII encoded hostname and port number.

An IPv4 address must be a dotted decimal format followed by a colon ':' (US-ASCII character 0x3A) and a decimal port number in US-ASCII.

An IPv6 address must be a colon separated format, surrounded by square brackets ('[', US-ASCII character 0x5B, and ']', US-ASCII character 0x5D), followed by a colon ':' (US-ASCII

character 0x3A) and a decimal port number in US-ASCII.

A hostname is always in US-US-ASCII (as per [RFC1033](#)); internationalized hostnames are encoded in US-US-ASCII as specified in [RFC 3490](#). The hostname is followed by a colon ':' (US-US-ASCII character 0x3A) and a decimal port number in US-US-ASCII. The name SHOULD be fully qualified whenever possible.

Values of this textual convention may not be directly usable as transport-layer addressing information, and may require run-time resolution. As such, applications that write them must be prepared for handling errors if such values are not supported, or cannot be resolved (if resolution occurs at the time of the management operation).

The DESCRIPTION clause of TransportAddress objects that may have snmpDTLSSCTPAddress values must fully describe how (and when) such names are to be resolved to IP addresses and vice versa.

This textual convention SHOULD NOT be used directly in object definitions since it restricts addresses to a specific format. However, if it is used, it MAY be used either on its own or in conjunction with TransportAddressType or TransportDomain as a pair.

When this textual convention is used as a syntax of an index object, there may be issues with the limit of 128 sub-identifiers specified in SMIV2 (STD 58). It is RECOMMENDED that all MIB documents using this textual convention make explicit any limitations on index component lengths that management software must observe. This may be done either by including SIZE constraints on the index components or by specifying applicable constraints in the conceptual row DESCRIPTION clause or in the surrounding documentation."

SYNTAX OCTET STRING (SIZE (1..255))

Fingerprint ::= TEXTUAL-CONVENTION

DISPLAY-HINT "1x:254x"

STATUS current

DESCRIPTION

"A Fingerprint value that can be used to uniquely reference other data of potentially arbitrary length.

A Fingerprint value is composed of a 1-octet hashing algorithm type. The octet value encoded is taken from the IANA TLS HashAlgorithm Registry ([RFC5246](#)). The remaining octets are

filled using the results of the hashing algorithm.

This TEXTUAL-CONVENTION SHOULD NOT be used as a form of cryptographic verification and a data source with a matching fingerprint should not be considered authenticated because the value matches. This TEXTUAL-CONVENTION is only intended for use as a reference to a stored copy of a longer data source. The contents of full data source referenced by this fingerprint needs to be compared against to assure collisions have not resulted."

REFERENCE

"[RFC 1033](#): DOMAIN ADMINISTRATORS OPERATIONS GUIDE
[RFC 3490](#): Internationalizing Domain Names in Applications
[RFC 3986](#): Uniform Resource Identifier (URI): Generic Syntax
[RFC 4347](#): Datagram Transport Layer Security
[RFC 5246](#): The Transport Layer Security (TLS) Protocol Version 1.2

"

SYNTAX OCTET STRING (SIZE (1..255))

-- The tlstmSession Group

tlstmSession OBJECT IDENTIFIER ::= { tlstmObjects 1 }

tlstmSessionOpens OBJECT-TYPE

SYNTAX Counter32

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"The number of times an openSession() request has been executed as an (D)TLS client, whether it succeeded or failed."

::= { tlstmSession 1 }

tlstmSessionCloses OBJECT-TYPE

SYNTAX Counter32

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"The number of times a closeSession() request has been executed as an (D)TLS client, whether it succeeded or failed."

::= { tlstmSession 2 }

tlstmSessionOpenErrors OBJECT-TYPE

SYNTAX Counter32

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"The number of times an openSession() request failed to open a session as a (D)TLS client, for any reason."


```
::= { tlstmSession 3 }
```

tlstmSessionNoAvailableSessions OBJECT-TYPE

SYNTAX Counter32

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"The number of times an outgoing message was dropped because the session associated with the passed tmStateReference was no longer (or was never) available."

```
::= { tlstmSession 4 }
```

tlstmSessionInvalidClientCertificates OBJECT-TYPE

SYNTAX Counter32

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"The number of times an incoming session was not established on an (D)TLS server because the presented client certificate was invalid. Reasons for invalidation includes, but is not limited to, cryptographic validation failures and lack of a suitable mapping row in the tlstmCertToSNTTable."

```
::= { tlstmSession 5 }
```

tlstmSessionInvalidServerCertificates OBJECT-TYPE

SYNTAX Counter32

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"The number of times an outgoing session was not established on an (D)TLS client because the presented server certificate was invalid. Reasons for invalidation includes, but is not limited to, cryptographic validation failures and an unexpected presented certificate identity."

```
::= { tlstmSession 6 }
```

tlstmTLSProtectionErrors OBJECT-TYPE

SYNTAX Counter32

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"The number of times (D)TLS processing resulted in a message being discarded because it failed its integrity test, decryption processing or other (D)TLS processing."

```
::= { tlstmSession 7 }
```

-- Configuration Objects


```
tlstmConfig          OBJECT IDENTIFIER ::= { tlstmObjects 2 }

-- Certificate mapping

tlstmCertificateMapping OBJECT IDENTIFIER ::= { tlstmConfig 1 }

tlstmCertToSNCount OBJECT-TYPE
    SYNTAX      Unsigned32
    MAX-ACCESS   read-only
    STATUS      current
    DESCRIPTION
        "A count of the number of entries in the tlstmCertToSNTable"
    ::= { tlstmCertificateMapping 1 }

tlstmCertToSNTableLastChanged OBJECT-TYPE
    SYNTAX      TimeStamp
    MAX-ACCESS   read-only
    STATUS      current
    DESCRIPTION
        "The value of sysUpTime.0 when the tlstmCertToSNTable
        was last modified through any means, or 0 if it has not been
        modified since the command responder was started."
    ::= { tlstmCertificateMapping 2 }

tlstmCertToSNTable OBJECT-TYPE
    SYNTAX      SEQUENCE OF TlstmCertToSNEntry
    MAX-ACCESS   not-accessible
    STATUS      current
    DESCRIPTION
        "A table listing the X.509 certificates known to the entity
        and the associated method for determining the SNMPv3 security
        name from a certificate.
```

On an incoming (D)TLS/SNMP connection the client's presented certificate must be examined and validated based on an established trusted path from a CA certificate or self-signed public certificate (e.g. [RFC5280](#)). This table provides a mapping from a validated certificate to a SNMPv3 securityName. This table does not provide any mechanisms for uploading trusted certificates; the transfer of any needed trusted certificates for path validation is expected to occur through an out-of-band transfer.

Once the authenticity of a certificate has been verified, this table is consulted to determine the appropriate securityName to identify with the remote connection. This is done by considering each active row from this table in prioritized order according to its tlstmCertToSNID value. Each row's

tlstmCertToSNFingerprint value determines whether the row is a match for the incoming connection:

- 1) If the row's tlstmCertToSNFingerprint value identifies the presented certificate and the contents of the presented certificate match a locally cached copy of the certificate then consider the row as a successful match.
- 2) If the row's tlstmCertToSNFingerprint value identifies a locally held copy of a trusted CA certificate and that CA certificated was used to validate the path to the presented certificate then consider the row as a successful match.

Once a matching row has been found, the tlstmCertToSNMapType value can be used to determine how the securityName to associate with the session should be determined. See the tlstmCertToSNMapType column's DESCRIPTION for details on determining the securityName value. If it is impossible to determine a securityName from the row's data combined with the data presented in the certificate then additional rows MUST be searched looking for another potential match. If a resulting securityName mapped from a given row is not compatible with the needed requirements of a securityName (e.g., VACM imposes a 32-octet-maximum length and the certificate derived securityName could be longer) then it must be considered an invalid match and additional rows MUST be searched looking for another potential match.

Missing values of tlstmCertToSNID are acceptable and implementations should continue to the next highest numbered row. E.G., the table may legally contain only two rows with tlstmCertToSNID values of 10 and 20.

Users are encouraged to make use of certificates with subjectAltName fields that can be used as securityNames so that a single root CA certificate can allow all child certificate's subjectAltName to map directly to a securityName via a 1:1 transformation. However, this table is flexible to allow for situations where existing deployed certificate infrastructures do not provide adequate subjectAltName values for use as SNMPv3 securityNames. Certificates may also be mapped to securityNames using the CommonName portion of the Subject field but usage of the CommonName field is deprecated. Direct mapping from each individual certificate fingerprint to a securityName is also possible but requires one entry in the table per securityName and requires more management operations

to completely configure a device."
::= { tlstmCertificateMapping 3 }

tlstmCertToSNEntry OBJECT-TYPE

SYNTAX TlstmCertToSNEntry

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"A row in the tlstmCertToSNTable that specifies a mapping for an incoming (D)TLS certificate to a securityName to use for a connection."

INDEX { tlstmCertToSNID }

::= { tlstmCertToSNTable 1 }

TlstmCertToSNEntry ::= SEQUENCE {

tlstmCertToSNID Unsigned32,

tlstmCertToSNFingerprint Fingerprint,

tlstmCertToSNMapType INTEGER,

tlstmCertToSNSecurityName SnmpAdminString,

tlstmCertToSNSANType INTEGER,

tlstmCertToSNStorageType StorageType,

tlstmCertToSNRowStatus RowStatus

}

tlstmCertToSNID OBJECT-TYPE

SYNTAX Unsigned32

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"A unique, prioritized index for the given entry."

::= { tlstmCertToSNEntry 1 }

tlstmCertToSNFingerprint OBJECT-TYPE

SYNTAX Fingerprint

MAX-ACCESS read-create

STATUS current

DESCRIPTION

"A cryptographic hash of a X.509 certificate. The results of a successful matching fingerprint to either the trusted CA in the certificate validation path or to the certificate itself is dictated by the tlstmCertToSNMapType column."

::= { tlstmCertToSNEntry 2 }

tlstmCertToSNMapType OBJECT-TYPE

SYNTAX INTEGER { specified(1), bySubjectAltName(2), byCN(3) }

MAX-ACCESS read-create

STATUS current

DESCRIPTION

"The mapping type used to obtain the securityName from the certificate. The possible values of use and their usage methods are defined as follows:

specified(1): The securityName that should be used to associate with the session is directly specified in the tlstmCertToSNSecurityName column from this table. Note: The tlstmCertToSNSecurityName column's value is ignored for all other tlstmCertToSNMapType values.

bySubjectAltName(2):

The securityName that should be used to associate with the session should be taken from the subjectAltName(s) portion of the client's X.509 certificate. The subjectAltName used MUST be the first encountered subjectAltName type indicated by the tlstmCertToSNSANType column.

If the resulting mapped value from the subjectAltName component is not compatible with the needed requirements of a securityName (e.g., VACM imposes a 32-octet-maximum length and the certificate derived securityName could be longer) then the next appropriate subjectAltName of the correct type should be used if available.

If no appropriate subjectAltName of the given type is found within the certificate then additional rows in the tlstmCertToSNTable must be searched for additional tlstmCertToSNFingerprint matches.

byCN(3):

The securityName that should be used to associate with the session should be taken from the CommonName portion of the Subject field from the client's presented X.509 certificate.

If the value of the CommonName component is not compatible with the needed requirements of a securityName (e.g., VACM imposes a 32-octet-maximum length and the certificate derived securityName could be longer) then additional rows in the tlstmCertToSNTable must be searched for additional tlstmCertToSNFingerprint matches."

DEFVAL { specified }


```
::= { tlstmCertToSNEntry 3 }
```

tlstmCertToSNSecurityName OBJECT-TYPE

SYNTAX SnmpAdminString (SIZE(0..32))

MAX-ACCESS read-create

STATUS current

DESCRIPTION

"The securityName that the session should use if the tlstmCertToSNMapType is set to specified(1), otherwise the value in this column should be ignored. If tlstmCertToSNMapType is set to specified(1) and this column contains a zero-length string (which is not a legal securityName value) this row is effectively disabled and the match will not be considered successful and other rows in the table will need to be searched for a proper match."

DEFVAL { "" }

```
::= { tlstmCertToSNEntry 4 }
```

tlstmCertToSNSANType OBJECT-TYPE

SYNTAX INTEGER { any(1), rfc822Name(2), dNSName(3),
ipAddress(4), otherName(5) }

MAX-ACCESS read-create

STATUS current

DESCRIPTION

"Specifies the subjectAltName type that may be used to extract the securityName from.

The any(1) value indicates the (D)TLS server should use the first value found for any of the following subjectAltName value types for the securityName: rfc822Name, dNSName, and ipAddress.

When multiple types for a given subjectAltName type are encountered within a certificate the first legally usable value is the one selected.

Values for type ipAddress(4) are converted to a valid securityName by:

- 1) for IPv4 the value is converted into a decimal dotted quad address (e.g. '192.0.2.1')
- 2) for IPv6 addresses the value is converted into a 32-character hexadecimal string without any colon separators.

Note that the resulting length is the maximum length supported by the View-Based Access Control Model

(VACM). Note that using both the Transport Security Model's support for transport prefixes (see the SNMP-TSM-MIB::snmpTsmConfigurationUsePrefix object for details) will result in securityName lengths that exceed what VACM can handle.

Values for type otherName(5) are converted to a valid securityName by using only the decoded value portion of the OtherName sequence. I.E. the OBJECT IDENTIFIER portion of the OtherName sequence is not included as part of the resulting securityName."

DEFVAL { any }
::= { tlstmCertToSNEntry 5 }

tlstmCertToSNStorageType OBJECT-TYPE

SYNTAX StorageType
MAX-ACCESS read-create
STATUS current

DESCRIPTION

"The storage type for this conceptual row. Conceptual rows having the value 'permanent' need not allow write-access to any columnar objects in the row."

DEFVAL { nonVolatile }
::= { tlstmCertToSNEntry 6 }

tlstmCertToSNRowStatus OBJECT-TYPE

SYNTAX RowStatus
MAX-ACCESS read-create
STATUS current

DESCRIPTION

"The status of this conceptual row. This object may be used to create or remove rows from this table.

To create a row in this table, a manager must set this object to either createAndGo(4) or createAndWait(5).

Until instances of all corresponding columns are appropriately configured, the value of the corresponding instance of the tlstmParamsRowStatus column is 'notReady'.

In particular, a newly created row cannot be made active until the corresponding tlstmCertToSNFingerprint, tlstmCertToSNMapType, tlstmCertToSNSecurityName, and tlstmCertToSNSANType columns have been set.

The following objects may not be modified while the value of this object is active(1):

- tlstmCertToSNFingerprint
- tlstmCertToSNMapType
- tlstmCertToSNSecurityName
- tlstmCertToSNSANType

An attempt to set these objects while the value of tlstmParamsRowStatus is active(1) will result in an inconsistentValue error."

::= { tlstmCertToSNEnter 7 }

-- Maps securityNames to certificates for use by the SNMP-TARGET-MIB

tlstmParamsCount OBJECT-TYPE

SYNTAX Unsigned32

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"A count of the number of entries in the tlstmParamsTable"

::= { tlstmCertificateMapping 4 }

tlstmParamsTableLastChanged OBJECT-TYPE

SYNTAX TimeStamp

MAX-ACCESS read-only

STATUS current

DESCRIPTION

"The value of sysUpTime.0 when the tlstmParamsTable was last modified through any means, or 0 if it has not been modified since the command responder was started."

::= { tlstmCertificateMapping 5 }

tlstmParamsTable OBJECT-TYPE

SYNTAX SEQUENCE OF TlstmParamsEntry

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"This table extends the SNMP-TARGET-MIB's snmpTargetParamsTable with an additional (D)TLS client-side certificate fingerprint identifier to use when establishing new (D)TLS connections."

::= { tlstmCertificateMapping 6 }

tlstmParamsEntry OBJECT-TYPE

SYNTAX TlstmParamsEntry

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"A conceptual row containing a fingerprint hash of a locally held certificate for a given snmpTargetParamsEntry. The values in this row should be ignored if the connection that

needs to be established, as indicated by the SNMP-TARGET-MIB infrastructure, is not a certificate and (D)TLS based connection. The connection SHOULD NOT be established if the certificate fingerprint stored in this entry does not point to a valid locally held certificate or if it points to an usable certificate (such as might happen when the certificate's expiration date has been reached)."

```
INDEX      { IMPLIED snmpTargetParamsName }  
::= { tlstmParamsTable 1 }
```

```
TlstmParamsEntry ::= SEQUENCE {  
    tlstmParamsClientFingerprint Fingerprint,  
    tlstmParamsStorageType       StorageType,  
    tlstmParamsRowStatus         RowStatus  
}
```

tlstmParamsClientFingerprint OBJECT-TYPE

```
SYNTAX      Fingerprint  
MAX-ACCESS  read-create  
STATUS      current
```

DESCRIPTION

"A cryptographic hash of a X.509 certificate. This object should store the hash of a locally held X.509 certificate that should be used when initiating a (D)TLS connection as a (D)TLS client."

```
::= { tlstmParamsEntry 1 }
```

tlstmParamsStorageType OBJECT-TYPE

```
SYNTAX      StorageType  
MAX-ACCESS  read-create  
STATUS      current
```

DESCRIPTION

"The storage type for this conceptual row. Conceptual rows having the value 'permanent' need not allow write-access to any columnar objects in the row."

```
DEFVAL      { nonVolatile }
```

```
::= { tlstmParamsEntry 2 }
```

tlstmParamsRowStatus OBJECT-TYPE

```
SYNTAX      RowStatus  
MAX-ACCESS  read-create  
STATUS      current
```

DESCRIPTION

"The status of this conceptual row. This object may be used to create or remove rows from this table."

To create a row in this table, a manager must set this object

to either `createAndGo(4)` or `createAndWait(5)`.

Until instances of all corresponding columns are appropriately configured, the value of the corresponding instance of the `tlstmParamsRowStatus` column is 'notReady'.

In particular, a newly created row cannot be made active until the corresponding `tlstmParamsClientFingerprint` column has been set.

The `tlstmParamsClientFingerprint` object may not be modified while the value of this object is `active(1)`.

An attempt to set these objects while the value of `tlstmParamsRowStatus` is `active(1)` will result in an `inconsistentValue` error.

If this row is deleted it has no effect on the corresponding row in the `targetParamsTable`.

If the corresponding row in the `targetParamsTable` is deleted then this row must be automatically removed."

```
::= { tlstmParamsEntry 3 }
```

```
-- Lists expected certificate fingerprints to be presented by a DTLS
-- server
```

```
tlstmAddrCount OBJECT-TYPE
```

```
    SYNTAX      Unsigned32
```

```
    MAX-ACCESS  read-only
```

```
    STATUS      current
```

```
    DESCRIPTION
```

```
        "A count of the number of entries in the tlstmAddrTable"
```

```
 ::= { tlstmCertificateMapping 7 }
```

```
tlstmAddrTableLastChanged OBJECT-TYPE
```

```
    SYNTAX      TimeStamp
```

```
    MAX-ACCESS  read-only
```

```
    STATUS      current
```

```
    DESCRIPTION
```

```
        "The value of sysUpTime.0 when the tlstmAddrTable
        was last modified through any means, or 0 if it has not been
        modified since the command responder was started."
```

```
 ::= { tlstmCertificateMapping 8 }
```

```
tlstmAddrTable OBJECT-TYPE
```

```
    SYNTAX      SEQUENCE OF TlstmAddrEntry
```

```
    MAX-ACCESS  not-accessible
```


STATUS current

DESCRIPTION

"This table extends the SNMP-TARGET-MIB's snmpTargetAddrTable with an expected (D)TLS server-side certificate identifier to expect when establishing a new (D)TLS connections. If a matching row in this table exists and the row is active then a local copy of the certificate matching the fingerprint identifier should be compared against the certificate being presented by the server. If the certificate presented by the server does not match the locally held copy then the connection MUST NOT be established. If no matching row exists in this table then the connection SHOULD still proceed if another certificate validation path algorithm (e.g. [RFC5280](#)) can be followed to a configured trust anchor. "

::= { tlstmCertificateMapping 9 }

tlstmAddrEntry OBJECT-TYPE

SYNTAX TlstmAddrEntry

MAX-ACCESS not-accessible

STATUS current

DESCRIPTION

"A conceptual row containing a copy of a locally held certificate's fingerprint for a given snmpTargetAddrEntry. The values in this row should be ignored if the connection that needs to be established, as indicated by the SNMP-TARGET-MIB infrastructure, is not a (D)TLS based connection. If an tlstmAddrEntry exists for a given snmpTargetAddrEntry then the presented server certificate MUST match or the connection MUST NOT be established. If a row in this table does not exist to match a snmpTargetAddrEntry row then the connection SHOULD still proceed if some other certificate validation path algorithm (e.g. [RFC5280](#)) can be followed to a configured trust anchor."

INDEX { IMPLIED snmpTargetAddrName }

::= { tlstmAddrTable 1 }

TlstmAddrEntry ::= SEQUENCE {

tlstmAddrServerFingerprint Fingerprint,

tlstmAddrStorageType StorageType,

tlstmAddrRowStatus RowStatus

}

tlstmAddrServerFingerprint OBJECT-TYPE

SYNTAX Fingerprint

MAX-ACCESS read-create

STATUS current

DESCRIPTION

"A cryptographic hash of a public X.509 certificate. This

object should store the hash of a local copy of the public X.509 certificate that the remote server should present during the (D)TLS connection setup. The presented certificate and the locally held copy, referred to by this hash value, MUST match exactly or the connection MUST NOT be established."

::= { tlstmAddrEntry 1 }

tlstmAddrStorageType OBJECT-TYPE

SYNTAX StorageType

MAX-ACCESS read-create

STATUS current

DESCRIPTION

"The storage type for this conceptual row. Conceptual rows having the value 'permanent' need not allow write-access to any columnar objects in the row."

DEFVAL { nonVolatile }

::= { tlstmAddrEntry 2 }

tlstmAddrRowStatus OBJECT-TYPE

SYNTAX RowStatus

MAX-ACCESS read-create

STATUS current

DESCRIPTION

"The status of this conceptual row. This object may be used to create or remove rows from this table.

To create a row in this table, a manager must set this object to either createAndGo(4) or createAndWait(5).

Until instances of all corresponding columns are appropriately configured, the value of the corresponding instance of the tlstmAddrRowStatus column is 'notReady'.

In particular, a newly created row cannot be made active until the corresponding tlstmAddrServerFingerprint column has been set.

The tlstmAddrServerFingerprint object may not be modified while the value of this object is active(1).

An attempt to set these objects while the value of tlstmAddrRowStatus is active(1) will result in an inconsistentValue error.

If this row is deleted it has no effect on the corresponding

row in the targetAddrTable.

If the corresponding row in the targetAddrTable is deleted
then this row must be automatically removed."

::= { tlstmAddrEntry 3 }

```
-- *****
-- tlstmNotifications - Notifications Information
-- *****
```

tlstmServerCertNotFound NOTIFICATION-TYPE

STATUS current

DESCRIPTION

"Notification that the server certificate presented by a SNMP
over (D)TLS server could not be found in the tlstmAddrTable."

::= { tlstmNotifications 1 }

tlstmServerAuthFailure NOTIFICATION-TYPE

OBJECTS { tlstmAddrServerFingerprint }

STATUS current

DESCRIPTION

"Notification that the server certificate presented by an SNMP
over (D)TLS server was found, but the connection could not be
established because of a cryptographic validation failure."

::= { tlstmNotifications 2 }

```
-- *****
-- tlstmCompliances - Conformance Information
-- *****
```

tlstmCompliances OBJECT IDENTIFIER ::= { tlstmConformance 1 }

tlstmGroups OBJECT IDENTIFIER ::= { tlstmConformance 2 }

```
-- *****
-- Compliance statements
-- *****
```

tlstmCompliance MODULE-COMPLIANCE

STATUS current

DESCRIPTION

"The compliance statement for SNMP engines that support the
TLSTM-MIB"

MODULE

MANDATORY-GROUPS { tlstmStatsGroup,


```
        tlstmIncomingGroup,
        tlstmOutgoingGroup,
        tlstmNotificationGroup }
 ::= { tlstmCompliances 1 }

-- *****
-- Units of conformance
-- *****
tlstmStatsGroup OBJECT-GROUP
  OBJECTS {
    tlstmSessionOpens,
    tlstmSessionCloses,
    tlstmSessionOpenErrors,
    tlstmSessionNoAvailableSessions,
    tlstmSessionInvalidClientCertificates,
    tlstmSessionInvalidServerCertificates,
    tlstmTLSProtectionErrors
  }
  STATUS      current
  DESCRIPTION
    "A collection of objects for maintaining
    statistical information of an SNMP engine which
    implements the SNMP TLS Transport Model."
  ::= { tlstmGroups 1 }

tlstmIncomingGroup OBJECT-GROUP
  OBJECTS {
    tlstmCertToSNCount,
    tlstmCertToSNTableLastChanged,
    tlstmCertToSNFingerprint,
    tlstmCertToSNMapType,
    tlstmCertToSNSecurityName,
    tlstmCertToSNSANType,
    tlstmCertToSNStorageType,
    tlstmCertToSNRowStatus
  }
  STATUS      current
  DESCRIPTION
    "A collection of objects for maintaining
    incoming connection certificate mappings to
    securityNames of an SNMP engine which implements the
    SNMP TLS Transport Model."
  ::= { tlstmGroups 2 }

tlstmOutgoingGroup OBJECT-GROUP
  OBJECTS {
    tlstmParamsCount,
    tlstmParamsTableLastChanged,
```



```
        tlstmParamsClientFingerprint,
        tlstmParamsStorageType,
        tlstmParamsRowStatus,
        tlstmAddrCount,
        tlstmAddrTableLastChanged,
        tlstmAddrServerFingerprint,
        tlstmAddrStorageType,
        tlstmAddrRowStatus
    }
    STATUS current
    DESCRIPTION
        "A collection of objects for maintaining
        outgoing connection certificates to use when opening
        connections as a result of SNMP-TARGET-MIB settings."
    ::= { tlstmGroups 3 }

tlstmNotificationGroup NOTIFICATION-GROUP
    NOTIFICATIONS {
        tlstmServerCertNotFound,
        tlstmServerAuthFailure
    }
    STATUS current
    DESCRIPTION
        "Notifications"
    ::= { tlstmGroups 4 }

END
```

8. Operational Considerations

This section discusses various operational aspects of deploying TLSTM.

8.1. Sessions

A session is discussed throughout this document as meaning a security association between the (D)TLS client and the (D)TLS server. State information for the sessions are maintained in each TLSTM implementation and this information is created and destroyed as sessions are opened and closed. A "broken" session (one side up and one side down) can result if one side of a session is brought down abruptly (i.e., reboot, power outage, etc.). Whenever possible, implementations SHOULD provide graceful session termination through the use of disconnect messages. Implementations SHOULD also have a system in place for dealing with "broken" sessions. Implementations SHOULD support the session resumption feature of TLS.

To simplify session management it is RECOMMENDED that implementations use separate ports for Notification sessions and for Command sessions. If this implementation recommendation is followed, (D)TLS clients will always send REQUEST messages and (D)TLS servers will always send RESPONSE messages. With this assertion, implementations may be able to simplify "broken" session handling, session resumption, and other aspects of session management such as guaranteeing that Request- Response pairs use the same session.

Implementations SHOULD limit the lifetime of established sessions depending on the algorithms used for generation of the master session secret, the privacy and integrity algorithms used to protect messages, the environment of the session, the amount of data transferred, and the sensitivity of the data.

8.2. Notification Receiver Credential Selection

When an SNMP engine needs to establish an outgoing session for notifications, the `snmpTargetParamsTable` includes an entry for the `snmpTargetParamsSecurityName` of the target. However, the receiving SNMP engine (Server) does not know which (D)TLS certificate to offer to the Client so that the `tmSecurityName` identity-authentication will be successful.

One solution is to maintain a one-to-one mapping between certificates and incoming ports for notification receivers. This can be handled at the Notification Originator by configuring the `snmpTargetAddrTable` (`snmpTargetAddrTDomain` and `snmpTargetAddrTAddress`) and requiring the receiving SNMP engine to monitor multiple incoming static ports based on which principals are capable of receiving notifications.

Implementations MAY also choose to designate a single Notification Receiver Principal to receive all incoming TRAPS and INFORMS or select an implementation specific method of selecting a server certificate to present to clients.

8.3. contextEngineID Discovery

Most Command Responders have `contextEngineIDs` that are identical to the USM `securityEngineID`. USM provides a discovery service that allows Command Generators to determine a `securityEngineID` and thus a default `contextEngineID` to use. Because the TLS Transport Model does not make use of a `securityEngineID`, it may be difficult for Command Generators to discover a suitable default `contextEngineID`. Implementations should consider offering another `engineID` discovery mechanism to continue providing Command Generators with a suitable `contextEngineID` mechanism. A recommended discovery solution is documented in [[RFC5343](#)].

9. Security Considerations

This document describes a transport model that permits SNMP to utilize (D)TLS security services. The security threats and how the (D)TLS transport model mitigates these threats are covered in detail throughout this document. Security considerations for DTLS are covered in [\[RFC4347\]](#) and security considerations for TLS are described in [Section 11](#) and Appendices D, E, and F of TLS 1.2 [\[RFC5246\]](#). DTLS adds to the security considerations of TLS only because it is more vulnerable to denial of service attacks. A random cookie exchange was added to the handshake to prevent anonymous denial of service attacks. [RFC 4347](#) recommends that the cookie exchange is utilized for all handshakes and therefore this specification also RECOMMENDEDs that implementers also support this cookie exchange.

9.1. Certificates, Authentication, and Authorization

Implementations are responsible for providing a security certificate installation and configuration mechanism. Implementations SHOULD support certificate revocation lists.

(D)TLS provides for authentication of the identity of both the (D)TLS server and the (D)TLS client. Access to MIB objects for the authenticated principal MUST be enforced by an access control subsystem (e.g. the VACM).

Authentication of the Command Generator principal's identity is important for use with the SNMP access control subsystem to ensure that only authorized principals have access to potentially sensitive data. The authenticated identity of the Command Generator principal's certificate is mapped to an SNMP model-independent securityName for use with SNMP access control.

The (D)TLS handshake only provides assurance that the certificate of the authenticated identity has been signed by an configured accepted Certificate Authority. (D)TLS has no way to further authorize or reject access based on the authenticated identity. An Access Control Model (such as the VACM) provides access control and authorization of a Command Generator's requests to a Command Responder and a Notification Responder's authorization to receive Notifications from a Notification Originator. However to avoid man-in-the-middle attacks both ends of the (D)TLS based connection MUST check the certificate presented by the other side against what was expected. For example, Command Generators must check that the Command Responder presented and authenticated itself with a X.509 certificate that was expected. Not doing so would allow an impostor, at a minimum, to present false data, receive sensitive information and/or provide a

false belief that configuration was actually received and acted upon. Authenticating and verifying the identity of the (D)TLS server and the (D)TLS client for all operations ensures the authenticity of the SNMP engine that provides MIB data.

The instructions found in the DESCRIPTION clause of the `tlstmCertToSNTTable` object must be followed exactly. It is also important that the rows of the table be searched in prioritized order starting with the row containing the lowest numbered `tlstmCertToSNID` value.

9.2. Use with SNMPv1/SNMPv2c Messages

The SNMPv1 and SNMPv2c message processing described in [RFC3484](#) ([BCP 74](#)) [[RFC3584](#)] always selects the SNMPv1(1) Security Model for an SNMPv1 message, or the SNMPv2c(2) Security Model for an SNMPv2c message. When running SNMPv1/SNMPv2c over a secure transport like the TLS Transport Model, the `securityName` and `securityLevel` used for access control decisions are then derived from the community string, not the authenticated identity and `securityLevel` provided by the TLS Transport Model.

9.3. MIB Module Security

There are a number of management objects defined in this MIB module with a MAX-ACCESS clause of read-write and/or read-create. Such objects may be considered sensitive or vulnerable in some network environments. The support for SET operations in a non-secure environment without proper protection can have a negative effect on network operations. These are the tables and objects and their sensitivity/vulnerability:

- o The `tlstmParamsTable` can be used to change the outgoing X.509 certificate used to establish a (D)TLS connection. Modification to objects in this table need to be adequately authenticated since modification to values in this table will have profound impacts to the security of outbound connections from the device. Since knowledge of authorization rules and certificate usage mechanisms may be considered sensitive, protection from disclosure of the SNMP traffic via encryption is also highly recommended.
- o The `tlstmAddrTable` can be used to change the expectations of the certificates presented by a remote (D)TLS server. Modification to objects in this table need to be adequately authenticated since modification to values in this table will have profound impacts to the security of outbound connections from the device. Since knowledge of authorization rules and certificate usage mechanisms may be considered sensitive, protection from disclosure of the

SNMP traffic via encryption is also highly recommended.

- o The `tlstmCertToSNTTable` is used to specify the mapping of incoming X.509 certificates to SNMPv3 securityNames. Modification to objects in this table need to be adequately authenticated since modification to values in this table will have profound impacts to the security of incoming connections to the device. Since knowledge of authorization rules and certificate usage mechanisms may be considered sensitive, protection from disclosure of the SNMP traffic via encryption is also highly recommended.

Some of the readable objects in this MIB module (i.e., objects with a MAX-ACCESS other than not-accessible) may be considered sensitive or vulnerable in some network environments. It is thus important to control even GET and/or NOTIFY access to these objects and possibly to even encrypt the values of these objects when sending them over the network via SNMP. These are the tables and objects and their sensitivity/vulnerability:

- o This MIB contains a collection of counters that monitor the (D)TLS connections being established with a device. Since knowledge of connection and certificate usage mechanisms may be considered sensitive, protection from disclosure of the SNMP traffic via encryption is also highly recommended.

SNMP versions prior to SNMPv3 did not include adequate security. Even if the network itself is secure (for example by using IPsec), even then, there is no control as to who on the secure network is allowed to access and GET/SET (read/change/create/delete) the objects in this MIB module.

It is RECOMMENDED that implementers consider the security features as provided by the SNMPv3 framework (see [\[RFC3410\]](#), [section 8](#)), including full support for the SNMPv3 cryptographic mechanisms (for authentication and privacy).

Further, deployment of SNMP versions prior to SNMPv3 is NOT RECOMMENDED. Instead, it is RECOMMENDED to deploy SNMPv3 and to enable cryptographic security. It is then a customer/operator responsibility to ensure that the SNMP entity giving access to an instance of this MIB module is properly configured to give access to the objects only to those principals (users) that have legitimate rights to indeed GET or SET (change/create/delete) them.

10. IANA Considerations

IANA is requested to assign:

1. a TCP port number in the range 1..1023 in the <http://www.iana.org/assignments/port-numbers> registry which will be the default port for receipt of SNMP command messages over a TLS Transport Model as defined in this document,
2. a TCP port number in the range 1..1023 in the <http://www.iana.org/assignments/port-numbers> registry which will be the default port for receipt of SNMP notification messages over a TLS Transport Model as defined in this document,
3. a UDP port number in the range 1..1023 in the <http://www.iana.org/assignments/port-numbers> registry which will be the default port for receipt of SNMP command messages over a DTLS/UDP connection as defined in this document,
4. a UDP port number in the range 1..1023 in the <http://www.iana.org/assignments/port-numbers> registry which will be the default port for receipt of SNMP notification messages over a DTLS/UDP connection as defined in this document,
5. a SCTP port number in the range 1..1023 in the <http://www.iana.org/assignments/port-numbers> registry which will be the default port for receipt of SNMP command messages over a DTLS/SCTP connection as defined in this document,
6. a SCTP port number in the range 1..1023 in the <http://www.iana.org/assignments/port-numbers> registry which will be the default port for receipt of SNMP notification messages over a DTLS/SCTP connection as defined in this document,
7. an SMI number under snmpDomains for the snmpTLSDomain object identifier,
8. an SMI number under snmpDomains for the snmpDTLSUDPDDomain object identifier,
9. an SMI number under snmpDomains for the snmpDTLSSCTPDdomain object identifier,
10. a SMI number under snmpModules, for the MIB module in this document,
11. "tls" as the corresponding prefix for the snmpTLSDomain in the SNMP Transport Model registry,
12. "dudp" as the corresponding prefix for the snmpDTLSUDPDdomain in the SNMP Transport Model registry,

13. "dsct" as the corresponding prefix for the snmpDTLSSCTPDomain in the SNMP Transport Model registry;

If possible, IANA is requested to use matching port numbers for all assignments for SNMP Commands being sent over TLS, DTLS/UDP, DTLS/SCTP.

If possible, IANA is requested to use matching port numbers for all assignments for SNMP Notifications being sent over TLS, DTLS/UDP, DTLS/SCTP.

Editor's note: this section should be replaced with appropriate descriptive assignment text after IANA assignments are made and prior to publication.

11. Acknowledgements

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Appendix A. (D)TLS Overview

The (D)TLS protocol is composed of two layers: the (D)TLS Record Protocol and the (D)TLS Handshake Protocol. The following subsections provide an overview of these two layers. Please refer to [\[RFC4347\]](#) for a complete description of the protocol.

A.1. The (D)TLS Record Protocol

At the lowest layer, layered on top of the transport control protocol or a datagram transport protocol (e.g. UDP or SCTP) is the (D)TLS Record Protocol.

The (D)TLS Record Protocol provides security that has three basic properties:

- o The session can be confidential. Symmetric cryptography is used for data encryption (e.g., [\[AES\]](#), [\[DES\]](#) etc.). The keys for this symmetric encryption are generated uniquely for each session and are based on a secret negotiated by another protocol (such as the (D)TLS Handshake Protocol). The Record Protocol can also be used without encryption.
- o Messages can have data integrity. Message transport includes a message integrity check using a keyed MAC. Secure hash functions (e.g., SHA, MD5, etc.) are used for MAC computations. The Record Protocol can operate without a MAC, but is generally only used in this mode while another protocol is using the Record Protocol as a transport for negotiating security parameters.
- o Messages are protected against replay. (D)TLS uses explicit sequence numbers and integrity checks. DTLS uses a sliding window to protect against replay of messages within a session.

(D)TLS also provides protection against replay of entire sessions. In a properly-implemented keying material exchange, both sides will generate new random numbers for each exchange. This results in different encryption and integrity keys for every session.

A.2. The (D)TLS Handshake Protocol

The (D)TLS Record Protocol is used for encapsulation of various higher-level protocols. One such encapsulated protocol, the (D)TLS Handshake Protocol, allows the server and client to authenticate each other and to negotiate an integrity algorithm, an encryption algorithm and cryptographic keys before the application protocol transmits or receives its first octet of data. Only the (D)TLS client can initiate the handshake protocol. The (D)TLS Handshake

Protocol provides security that has four basic properties:

- o The peer's identity can be authenticated using asymmetric (public key) cryptography (e.g., RSA [[RSA](#)], DSS [[DSS](#)], etc.). This authentication can be made optional, but is generally required by at least one of the peers.

(D)TLS supports three authentication modes: authentication of both the server and the client, server authentication with an unauthenticated client, and total anonymity. For authentication of both entities, each entity provides a valid certificate chain leading to an acceptable certificate authority. Each entity is responsible for verifying that the other's certificate is valid and has not expired or been revoked. See [[I-D.saintandre-tls-server-id-check](#)] for further details on standardized processing when checking server certificate identities.

- o The negotiation of a shared secret is secure: the negotiated secret is unavailable to eavesdroppers, and for any authenticated handshake the secret cannot be obtained, even by an attacker who can place himself in the middle of the session.
- o The negotiation is not vulnerable to malicious modification: it is infeasible for an attacker to modify negotiation communication without being detected by the parties to the communication.
- o DTLS uses a stateless cookie exchange to protect against anonymous denial of service attacks and has retransmission timers, sequence numbers, and counters to handle message loss, reordering, and fragmentation.

[Appendix B](#). PKIX Certificate Infrastructure

Users of a public key from a PKIX / X.509 certificate can be confident that the associated private key is owned by the correct remote subject (person or system) with which an encryption or digital signature mechanism will be used. This confidence is obtained through the use of public key certificates, which are data structures that bind public key values to subjects. The binding is asserted by having a trusted CA digitally sign each certificate. The CA may base this assertion upon technical means (i.e., proof of possession through a challenge- response protocol), presentation of the private key, or on an assertion by the subject. A certificate has a limited valid lifetime which is indicated in its signed contents. Because a certificate's signature and timeliness can be independently checked by a certificate-using client, certificates can be distributed via

untrusted communications and server systems, and can be cached in unsecured storage in certificate-using systems.

ITU-T X.509 (formerly CCITT X.509) or ISO/IEC/ITU 9594-8, which was first published in 1988 as part of the X.500 Directory recommendations, defines a standard certificate format [[x509](#)] which is a certificate which binds a subject (principal) to a public key value. This was later further documented in [[RFC5280](#)].

A X.509 certificate is a sequence of three required fields:

tbsCertificate: The tbsCertificate field contains the names of the subject and issuer, a public key associated with the subject, a validity period, and other associated information. This field may also contain extension components.

signatureAlgorithm: The signatureAlgorithm field contains the identifier for the cryptographic algorithm used by the certificate authority (CA) to sign this certificate.

signatureValue: The signatureValue field contains a digital signature computed by the CA upon the ASN.1 DER encoded tbsCertificate field. The ASN.1 DER encoded tbsCertificate is used as the input to the signature function. This signature value is then ASN.1 DER encoded as a BIT STRING and included in the Certificate's signature field. By generating this signature, the CA certifies the validity of the information in the tbsCertificate field. In particular, the CA certifies the binding between the public key material and the subject of the certificate.

The basic X.509 authentication procedure is as follows: A system is initialized with a number of root certificates that contain the public keys of a number of trusted CAs. When a system receives a X.509 certificate, signed by one of those CAs, the certificate has to be verified. It first checks the signatureValue field by using the public key of the corresponding trusted CA. Then it compares the decrypted information with a digest of the tbsCertificate field. If they match, then the subject in the tbsCertificate field is authenticated.

[Appendix C](#). Target and Notificaton Configuration Example

Configuring the SNMP-TARGET-MIB and NOTIFICATION-MIB along with access control settings for the SNMP-VIEW-BASED-ACM-MIB can be a daunting task without an example to follow. The following section describes an example of what pieces must be in place to accomplish this configuration.

The `isAccessAllowed()` ASI requires configuration to exist in the following SNMP-VIEW-BASED-ACM-MIB tables:

```
vacmSecurityToGroupTable
vacmAccessTable
vacmViewTreeFamilyTable
```

The only table that needs to be discussed as particularly different here is the `vacmSecurityToGroupTable`. This table is indexed by both the SNMPv3 security model and the security name. The security model, when TLSTM is in use, should be set to the value of 4, corresponding to the TSM [[RFC5591](#)]. An example `vacmSecurityToGroupTable` row might be filled out as follows (using a single SNMP SET request):

```
vacmSecurityModel          = 4 (TSM)
vacmSecurityName           = "blueberry"
vacmGroupName              = "administrators"
vacmSecurityToGroupStorageType = 3 (nonVolatile)
vacmSecurityToGroupStatus   = 4 (createAndGo)
```

This example will assume that the "administrators" group has been given proper permissions via rows in the `vacmAccessTable` and `vacmViewTreeFamilyTable`.

Depending on whether this VACM configuration is for a Command Responder or a Command Generator the security name "blueberry" will come from a few different locations.

C.1. Configuring the Notification Generator

For Notification Generator's performing authorization checks, the server's certificate must be verified against the expected certificate before proceeding to send the notification. The expected certificate from the server may be listed in the `tlstmAddrTable` or may be determined through other X.509 path validation mechanisms. The securityName to use for VACM authorization checks is set by the SNMP-TARGET-MIB's `snmpTargetParamsSecurityName` column.

The certificate that the notification generator should present to the server is taken from the `tlstmParamsClientFingerprint` column from the appropriate entry in the `tlstmParamsTable` table.

C.2. Configuring the Command Responder

For Command Responder applications, the `vacmSecurityName` "blueberry" value is a value that needs derive from an incoming (D)TLS session. The mapping from a received (D)TLS client certificate to a

securityName is done with the tlstmCertToSNTable. The certificates must be loaded into the device so that a tlstmCertToSNEntry may refer to it. As an example, consider the following entry which will provide a mapping from a client's public X.509's hash fingerprint directly to the "blueberry" securityName:

```
tlstmCertToSNID          = 1          (chosen by ordering preference)
tlstmCertToSNFingerprint = HASH       (appropriate fingerprint)
tlstmCertToSNMapType     = 1          (specified)
tlstmCertToSNSecurityName = "blueberry"
tlstmCertToSNStorageType = 3          (nonVolatile)
tlstmCertToSNRowStatus   = 4          (createAndGo)
```

The above is an example of how to map a particular certificate to a particular securityName. It is recommended, however, that users make use of direct subjectAltName or CommonName mappings where possible as it provides a more scalable approach to certificate management. This entry provides an example of using a subjectAltName mapping:

```
tlstmCertToSNID          = 1          (chosen by ordering preference)
tlstmCertToSNFingerprint = HASH       (appropriate fingerprint)
tlstmCertToSNMapType     = 2          (bySubjectAltName)
tlstmCertToSNSANType     = 1          (any)
tlstmCertToSNStorageType = 3          (nonVolatile)
tlstmCertToSNRowStatus   = 4          (createAndGo)
```

The above entry indicates the subjectAltName field for certificates created by an Issuing certificate with a corresponding fingerprint will be trusted to always produce common names that are directly 1 to 1 mappable into SNMPv3 securityNames. This type of configuration should only be used when the certificate authorities naming conventions are carefully controlled.

In the example, if the incoming (D)TLS client provided certificate contained a subjectAltName where the first listed subjectAltName in the extension is the rfc822Name of "blueberry", the certificate was signed by a certificate matching the tlstmCertToSNFingerprint value and the CA's certificate was properly installed on the device then the string "blueberry" would be used as the securityName for the session.

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