Internet Engineering Task Force
Internet-Draft
Updates: 4120 (if approved)
Intended status: Standards Track
Expires: December 27, 2015

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June 25, 2015

Kerberos Authorization Data Container Authenticated by Multiple MACs draft-ietf-kitten-cammac-03

Abstract

This document specifies a Kerberos Authorization Data container that supersedes AD-KDC-ISSUED. It allows for multiple Message Authentication Codes (MACs) or signatures to authenticate the contained Authorization Data elements. The multiple MACs are needed to mitigate shortcomings in the existing AD-KDC-ISSUED container. This document updates RFC 4120.

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

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This Internet-Draft will expire on December 27, 2015.

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1. Introduction

This document specifies a new Authorization Data container for Kerberos, called the CAMMAC (Container Authenticated by Multiple MACs). The ASN.1 type implementing the CAMMAC concept is the AD-CAMMAC, which supersedes the AD-KDC-ISSUED Authorization Data type specified in [RFC4120]. This new container allows both the receiving application service and the Key Distribution Center (KDC) itself to verify the authenticity of the contained authorization data. The AD-CAMMAC container can also include additional verifiers that "trusted services" can use to verify the contained authorization data.

2. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

3. Motivations

The Kerberos protocol allows clients to submit arbitrary authorization data for a KDC to insert into a Kerberos ticket. These client-requested authorization data allow the client to express authorization restrictions that the application service will interpret. With few exceptions, the KDC can safely copy these client-requested authorization data to the issued ticket without necessarily inspecting, interpreting, or filtering their contents.

The AD-KDC-ISSUED authorization data container specified in RFC 4120 [RFC4120] is a means for KDCs to include positive or permissive (rather than restrictive) authorization data in service tickets in a way that the service named in a ticket can verify that the KDC has issued the contained authorization data. This capability takes advantage of a shared symmetric key between the KDC and the service to assure the service that the KDC did not merely copy client-requested authorization data to the ticket without inspecting them.

The AD-KDC-ISSUED container works well for situations where the flow of authorization data is from the KDC to the service. However, protocol extensions such as Constrained Delegation (S4U2Proxy [MS-SFU]) require that a service present to the KDC a service ticket that the KDC previously issued, as evidence that the service is authorized to impersonate the client principal named in that ticket. In the S4U2Proxy extension, the KDC uses the evidence ticket as the basis for issuing a derivative ticket that the service can then use to impersonate the client. The authorization data contained within the evidence ticket constitute a flow of authorization data from the application service to the KDC. The properties of the AD-KDC-ISSUED container are insufficient for this use case because the service knows the symmetric key for the checksum in the AD-KDC-ISSUED container. Therefore, the KDC has no way to detect whether the service has tampered with the contents of the AD-KDC-ISSUED container within the evidence ticket.

The new AD-CAMMAC authorization data container specified in this document improves upon AD-KDC-ISSUED by including additional verifier elements. The svc-verifier (service verifier) element of the AD-CAMMAC has the same functional and security properties as the ad-checksum element of AD-KDC-ISSUED; the svc-verifier allows the service to verify the integrity of the AD-CAMMAC contents as it already could with the AD-KDC-ISSUED container. The kdc-verifier and other-verifiers elements are new to AD-CAMMAC and provide its enhanced capabilities.

The kdc-verifier element of the AD-CAMMAC container allows a KDC to verify the integrity of authorization data that it previously inserted into a ticket, by using a key that only the KDC knows. The KDC thus avoids recomputing all of the authorization data for the issued ticket; this recomputation might not always be possible when that data includes ephemeral information such as the strength or type of authentication method used to obtain the original ticket.

The verifiers in the other-verifiers element of the AD-CAMMAC container are not required, but can be useful when a lesser-privileged service receives a ticket from a client and needs to extract the AD-CAMMAC to demonstrate to a higher-privileged "trusted

service" on the same host that it is legitimately acting on behalf of that client. The trusted service can use a verifier in the other-verifiers element to validate the contents of the AD-CAMMAC without further communication with the KDC.

4. Encoding

The Kerberos protocol is defined in [RFC4120] using Abstract Syntax Notation One (ASN.1) [X.680] and using the ASN.1 Distinguished Encoding Rules (DER) [X.690]. For consistency, this specification also uses ASN.1 for specifying the layout of AD-CAMMAC. The ad-data of the AD-CAMMAC authorization data element is the ASN.1 DER encoding of the AD-CAMMAC ASN.1 type specified below.

```
KerberosV5CAMMAC {
        iso(1) identified-organization(3) dod(6) internet(1)
        security(5) kerberosV5(2) modules(4) cammac(7)
} DEFINITIONS EXPLICIT TAGS ::= BEGIN

IMPORTS
      AuthorizationData, PrincipalName, Checksum, UInt32, Int32
        FROM KerberosV5Spec2 { iso(1) identified-organization(3)
          dod(6) internet(1) security(5) kerberosV5(2)
          modules(4) krb5spec2(2) };
          -- as defined in RFC 4120.

AD-CAMMAC                   ::= SEQUENCE {
      elements              [0] AuthorizationData,
      kdc-verifier          [1] Verifier-MAC OPTIONAL,
      svc-verifier          [2] Verifier-MAC OPTIONAL,
      other-verifiers       [3] SEQUENCE (SIZE (1..MAX))
                                OF Verifier OPTIONAL
}

Verifier              ::= CHOICE {
      mac             Verifier-MAC,
      ...
}

Verifier-MAC          ::= SEQUENCE {
      identifier      [0] PrincipalName OPTIONAL,
      kvno            [1] UInt32 OPTIONAL,
      enctype         [2] Int32 OPTIONAL,
      mac             [3] Checksum
}

END
```

elements:
A sequence of authorization data elements issued by the KDC. These elements are the authorization data that the verifier fields authenticate.

Verifier:
A CHOICE type that currently contains only one alternative: Verifier-MAC. Future extensions might add support for public-key signatures.

Verifier-MAC:
Contains an RFC 3961 [RFC3961] Checksum (MAC) computed over the ASN.1 DER encoding of the AuthorizationData value in the elements field of the AD-CAMMAC. The identifier, kvno, and enctype fields help the recipient locate the key required for verifying the MAC. For the kdc-verifier and the svc-verifier, the identifier, kvno and enctype fields are often obvious from context and MAY be omitted. For the kdc-verifier, the MAC is computed differently than for the svc-verifier and the other-verifiers, as described later. The key usage number for computing the MAC (Checksum) is 64.

kdc-verifier:
A Verifier-MAC where the key is a long-term key of the local Ticket-Granting Service (TGS). The checksum type is the required checksum type for the enctype of the TGS key. In contrast to the other Verifier-MAC elements, the KDC computes the MAC in the kdc-verifier over the ASN.1 DER encoding of the EncTicketPart of the surrounding ticket, but where the AuthorizationData value in the EncTicketPart contains the AuthorizationData value contained in the AD-CAMMAC instead of the AuthorizationData value that would otherwise be present in the ticket. This altered Verifier-MAC computation binds the kdc-verifier to the other contents of the ticket, assuring the KDC that a malicious service has not substituted a mismatched AD-CAMMAC received from another ticket.

svc-verifier:
A Verifier-MAC where the key is the same long-term service key that the KDC uses to encrypt the surrounding ticket. The checksum type is the required checksum type for the enctype of the service key used to encrypt the ticket. This field MUST be present if the service principal of the ticket is not the local TGS, including when the ticket is a cross-realm Ticket-Granting Ticket (TGT).

other-verifiers:
A sequence of additional verifiers. In each additional Verifier-MAC, the key is a long-term key of the principal name specified in the identifier field. The PrincipalName MUST be present and be a

valid principal in the realm. KDCs MAY add one or more "trusted service" verifiers. Unless otherwise administratively configured, the KDC SHOULD determine the "trusted service" principal name by replacing the service identifier component of the sname of the surrounding ticket with "host". The checksum is computed using a long-term key of the identified principal, and the checksum type is the required checksum type for the enctype of that long-term key. The kvno and enctype SHOULD be specified to disambiguate which of the long-term keys of the trusted service is used.

5. Usage

Application servers and KDCs MAY ignore the AD-CAMMAC container and the authorization data elements it contains. For compatibility with older Kerberos implementations, a KDC issuing an AD-CAMMAC SHOULD enclose it in an AD-IF-RELEVANT container [RFC4120] unless the KDC knows that the application server is likely to recognize it.

6. Assigned numbers

RFC 4120 is updated in the following ways:

- The ad-type number 96 is assigned for AD-CAMMAC, updating the table in Section 7.5.4 of [RFC4120].

- The table in Section 5.2.6 of [RFC4120] is updated to map the ad-type 96 to "DER encoding of AD-CAMMAC".

- The key usage number 64 is assigned for the Verifier-MAC checksum, updating the table in Section 7.5.1 of [RFC4120].

7. IANA Considerations

[RFC Editor: please remove this section prior to publication.]

There are no IANA considerations in this document. Any numbers assigned in this document are not in IANA-controlled number spaces.

8. Security Considerations

The CAMMAC provides data origin authentication for authorization data contained in it, attesting that it originated from the KDC. This section describes the precautions required to maintain the integrity of that data origin authentication through various information flows involving a Kerberos ticket containing a CAMMAC.

When handling a TGS-REQ containing a CAMMAC, a KDC makes a policy decision on how to produce the CAMMAC contents of the newly issued

ticket based on properties of the ticket(s) accompanying the TGS-REQ. This policy decision can involve filtering, transforming, or verbatim copying of the original CAMMAC contents. The following paragraphs provide some guidance on formulating such policies.

A KDC verifies a CAMMAC as originating from a local realm KDC when at least one of following the criteria is true:

1. The KDC successfully verifies the kdc-verifier; or

2. The KDC successfully verifies the svc-verifier, and the svc-verifier uses a key known only to the local realm KDCs; or

3. No verifiers are present, the ticket-encrypting key is known only to local realm KDCs, and all local realm KDCs properly filter out client-submitted CAMMACs. (This can require particular caution in a realm that has KDCs with mixed CAMMAC support, as might happen when incrementally upgrading KDCs in a realm to support CAMMAC.)

A CAMMAC that originates from a local realm KDC might contain information that originates from elsewhere. Originating from a local realm KDC means that a local realm KDC attests that the CAMMAC contents conform to the local realm's policy, regardless of the ultimate origin of the information in the CAMMAC (which could be a remote realm in the case of a CAMMAC contained in a cross-realm TGT).

Local policy determines when a KDC can apply a kdc-verifier to a CAMMAC (or otherwise creates a CAMMAC that satisfies the local origin criteria listed above). Semantically, a CAMMAC that a KDC verifies as originating from a local realm KDC attests that the CAMMAC contents conformed to local policy at the time of creation of the CAMMAC. Such a local policy can include allowing verbatim copying of CAMMAC contents from cross-realm TGTs from designated remote realms and applying a kdc-verifier to the new CAMMAC.

Usually, when a KDC verifies a CAMMAC as originating from a local realm KDC, the KDC can assume that the CAMMAC contents continue to conform to the local realm's policies. It is generally safe for a KDC to make verbatim copies of the contents of such a CAMMAC into a new CAMMAC when handling a TGS-REQ. Particularly strict implementations might conduct additional policy checks on the contents of a CAMMAC originating from a local realm KDC if the local realm's policy materially changes during the life of the CAMMAC.

A KDC MAY omit the kdc-verifier from the CAMMAC when it is not needed, according to how realm policies will subsequently treat the containing ticket. An implementation might choose to do this

omission to reduce the size of tickets it issues. Some examples of when such an omission is safe are:

1. For a local realm TGT, if all local realm KDCs correctly filter out client-submitted CAMMACs, the local realm origin criteria listed above allow omission of the kdc-verifier.

2. An application service might not use the S4U2Proxy extension, or the realm policy might disallow the use of S4U2Proxy by that service. In such situations where there is no flow of authorization data from the service to the KDC, the application service could modify the CAMMAC contents, but such modifications would have no effect on other services. Because of the lack of security impact to other application services, the KDC MAY omit the kdc-verifier from a CAMMAC contained in a ticket for that service.

Extracting a CAMMAC from a ticket for use as a credential removes it from the context of the ticket. In the general case, this could turn it into a bearer token, with all of the associated security implications. Also, the CAMMAC does not itself necessarily contain sufficient information to identify the client principal. Therefore, application protocols that rely on extracted CAMMACs might need to duplicate a substantial portion of the ticket contents and include that duplicated information in the authorization data contained within the CAMMAC. The extent of this duplication would depend on the security properties required by the application protocol.

The method for computing the kdc-verifier binds it only to the authorization data contained within the CAMMAC; it does not bind the CAMMAC to any authorization data within the containing ticket but outside the CAMMAC. At least one (non-standard) authorization data type, AD-SIGNEDPATH, attempts to bind to other authorization data in a ticket, and it is very difficult for two such authorization data types to coexist.

The kdc-verifier in CAMMAC does not bind the service principal name to the CAMMAC contents, because the service principal name is not part of the EncTicketPart. An entity that has access to the keys of two different service principals can decrypt a ticket for one service and encrypt it in the key of the other service, altering the svc-verifier to match. Both the kdc-verifier and the svc-verifier would still validate, but the KDC never issued this fabricated ticket. The impact of this manipulation is minor if the CAMMAC contents only communicate attributes related to the client. If an application requires an authenticated binding between the service principal name and the CAMMAC or ticket contents, the KDC MUST include in the CAMMAC some authorization data element that names the service principal.

9. Acknowledgements

Shawn Emery, Sam Hartman, Greg Hudson, Ben Kaduk, Barry Leiba, Meral Shirazipour, Zhanna Tsitkov, Qin Wu, and Kai Zheng provided helpful technical and editorial feedback on earlier versions of this document.

10. References

10.1. Normative References

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10.2. Informative References

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