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Abstract

This document updates PKINIT, as defined in <u>RFC 4556</u>, to remove protocol structures tied to specific cryptographic algorithms. The PKINIT key derivation function is made negotiable, and the digest algorithms for signing the pre-authentication data and the client's X.509 certificates are made discoverable.

These changes provide preemptive protection against vulnerabilities discovered in the future against any specific cryptographic algorithm, and allow incremental deployment of newer algorithms.

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1. Introduction

This document updates PKINIT [RFC4556] to remove protocol structures tied to specific cryptographic algorithms. The PKINIT key derivation function is made negotiable, the digest algorithms for signing the pre-authentication data and the client's X.509 certificates are made discoverable.

These changes provide preemptive protection against vulnerabilities discovered in the future against any specific cryptographic algorithm, and allow incremental deployment of newer algorithms.

In August 2004, Xiaoyun Wang's research group reported MD4 [RFC6150] collisions generated using hand calculation [WANG04], alongside attacks on later hash function designs in the MD4, MD5 [RFC1321] and SHA [RFC6234] family. These attacks and their consequences are discussed in [RFC6194]. These discoveries challenged the security of protocols relying on the collision resistance properties of these hashes.

The Internet Engineering Task Force (IETF) called for actions to update existing protocols to provide crypto algorithm agility so that protocols support multiple cryptographic algorithms (including hash functions) and provide clean, tested transition strategies between algorithms, as recommended by BCP 201 [RFC7696].

This document updates PKINIT to provide crypto algorithm agility. Several protocol structures used in the [RFC4556] protocol are either tied to SHA-1, or do not support negotiation or discovery, but are instead based on local policy. The following concerns have been addressed in this update:

- o The checksum algorithm in the authentication request is hardwired to use SHA-1 [RFC6234].
- o The acceptable digest algorithms for signing the authentication data are not discoverable.
- o The key derivation function in <u>Section 3.2.3.1 of [RFC4556]</u> is hardwired to use SHA-1.
- o The acceptable digest algorithms for signing the client X.509 certificates are not discoverable.

To address these concerns, new key derivation functions (KDFs), identified by object identifiers, are defined. The PKINIT client provides a list of KDFs in the request and the Key Distribution Center (KDC) picks one in the response, thus a mutually-supported KDF is negotiated.

Furthermore, structures are defined to allow the client to discover the Cryptographic Message Syntax (CMS) [RFC5652] digest algorithms supported by the KDC for signing the pre-authentication data and signing the client X.509 certificate.

2. Requirements Notation

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

3. paChecksum Agility

The paChecksum defined in Section 3.2.1 of [RFC4556] provides a cryptographic binding between the client's pre-authentication data and the corresponding Kerberos request body. This also prevents the KDC-REQ body from being tampered with. SHA-1 is the only allowed checksum algorithm defined in [RFC4556]. This facility relies on the collision resistance properties of the SHA-1 checksum [RFC6234].

When the reply key delivery mechanism is based on public key encryption as described in <u>Section 3.2.3.2 of [RFC4556]</u>, the asChecksum in the KDC reply provides the binding between the preauthentication and the ticket request and response messages, and integrity protection for the unauthenticated clear text in these messages. However, if the reply key delivery mechanism is based on the Diffie-Hellman key agreement as described in Section 3.2.3.1 of [RFC4556], the security provided by using SHA-1 in the paChecksum is weak, and nothing else cryptographically binds the AS request to the ticket response. In this case, the new KDF selected by the KDC as described in <u>Section 6</u> provides the cryptographic binding and integrity protection.

4. CMS Digest Algorithm Agility

When the KDC_ERR_DIGEST_IN_SIGNED_DATA_NOT_ACCEPTED error is returned as described in <u>Section 3.2.2 of [RFC4556]</u>, implementations conforming to this specification can OPTIONALLY send back a list of supported CMS types signifying the digest algorithms supported by the KDC, in the decreasing preference order. This is accomplished by

including a TD_CMS_DATA_DIGEST_ALGORITHMS typed data element in the error data.

td-cms-digest-algorithms INTEGER ::= 111

The corresponding data for the TD_CMS_DATA_DIGEST_ALGORITHMS contains the ASN.1 Distinguished Encoding Rules (DER) [X680] [X690] encoded TD-CMS-DIGEST-ALGORITHMS-DATA structure defined as follows:

TD-CMS-DIGEST-ALGORITHMS-DATA ::= SEQUENCE OF AlgorithmIdentifier

- -- Contains the list of CMS algorithm [RFC5652]
- -- identifiers indicating the digest algorithms
- -- acceptable to the KDC for signing CMS data in
- -- the order of decreasing preference.

The algorithm identifiers in the TD-CMS-DIGEST-ALGORITHMS identify digest algorithms supported by the KDC.

This information sent by the KDC via TD_CMS_DATA_DIGEST_ALGORITHMS can facilitate trouble-shooting when none of the digest algorithms supported by the client is supported by the KDC.

5. X.509 Certificate Signer Algorithm Agility

When the client's X.509 certificate is rejected and the KDC_ERR_DIGEST_IN_SIGNED_DATA_NOT_ACCEPTED error is returned as described in Section 3.2.2 of RFC4556], implementations conforming to this specification can OPTIONALLY send a list of digest algorithms acceptable to the KDC for use by the Certificate Authority (CA) in signing the client's X.509 certificate, in the decreasing preference order. This is accomplished by including a TD_CERT_DIGEST_ALGORITHMS typed data element in the error data. The corresponding data contains the ASN.1 DER encoding of the structure TD-CERT-DIGEST-ALGORITHMS-DATA defined as follows:

```
td-cert-digest-algorithms INTEGER ::= 112
TD-CERT-DIGEST-ALGORITHMS-DATA ::= SEQUENCE {
        allowedAlgorithms [0] SEQUENCE OF AlgorithmIdentifier,
                   -- Contains the list of CMS algorithm [RFC5652]
                   -- identifiers indicating the digest algorithms
                   -- that are used by the CA to sign the client's
                   -- X.509 certificate and are acceptable to the KDC
                   -- in the process of validating the client's X.509
                   -- certificate, in the order of decreasing
                   -- preference.
        rejectedAlgorithm [1] AlgorithmIdentifier OPTIONAL,
                   -- This identifies the digest algorithm that was
                   -- used to sign the client's X.509 certificate and
                   -- has been rejected by the KDC in the process of
                   -- validating the client's X.509 certificate
                   -- [<u>RFC5280</u>].
        . . .
}
```

The KDC fills in the allowedAlgorithm field with the list of algorithm [RFC5652] identifiers indicating digest algorithms that are used by the CA to sign the client's X.509 certificate and are acceptable to the KDC in the process of validating the client's X.509 certificate, in the order of decreasing preference. The rejectedAlgorithm field identifies the signing algorithm for use in signing the client's X.509 certificate that has been rejected by the KDC in the process of validating the client's certificate [RFC5280].

6. KDF agility

Based on [RFC3766] and [X9.42], Section 3.2.3.1 of [RFC4556] defines a Key Derivation Function (KDF) that derives a Kerberos protocol key based on the secret value generated by the Diffie-Hellman key exchange. This KDF requires the use of SHA-1 [RFC6234].

The KDF algorithm described in this document (based on [SP80056A]) can be implemented using any cryptographic hash function.

A new KDF for PKINIT usage is identified by an object identifier. The following KDF object identifiers are defined:

```
id-pkinit OBJECT IDENTIFIER ::=
         { iso(1) identified-organization(3) dod(6) internet(1)
           security(5) kerberosv5(2) pkinit (3) }
    -- Defined in RFC 4556 and quoted here for the reader.
id-pkinit-kdf OBJECT IDENTIFIER ::= { id-pkinit kdf(6) }
    -- PKINIT KDFs
id-pkinit-kdf-ah-sha1 OBJECT IDENTIFIER
    ::= { id-pkinit-kdf sha1(1) }
    -- SP800-56A ASN.1 structured hash-based KDF using SHA-1
id-pkinit-kdf-ah-sha256 OBJECT IDENTIFIER
    ::= { id-pkinit-kdf sha256(2) }
    -- SP800-56A ASN.1 structured hash-based KDF using SHA-256
id-pkinit-kdf-ah-sha512 OBJECT IDENTIFIER
    ::= { id-pkinit-kdf sha512(3) }
    -- SP800-56A ASN.1 structured hash-based KDF using SHA-512
id-pkinit-kdf-ah-sha384 OBJECT IDENTIFIER
    ::= { id-pkinit-kdf sha384(4) }
    -- SP800-56A ASN.1 structured hash-based KDF using SHA-384
```

Where id-pkinit is defined in [RFC4556]. All key derivation functions specified above use the one-step key derivation method described in Section 5.8.2.1 of [SP80056A], using the ASN.1 format for FixedInfo, and Section 4.1 of [SP80056C], using option 1 for the auxiliary function H. id-pkinit-kdf-ah-sha1 uses SHA-1 [RFC6234] as the hash function. id-pkinit-kdf-ah-sha256, id-pkinit-kdf-ah-sha356, and id-pkinit-kdf-ah-sha512 use SHA-256 [RFC6234], SHA-384 ([RFC6234] and SHA-512 [RFC6234] respectively.

To name the input parameters, an abbreviated version of the key derivation method is described below.

- 1. reps = ceiling(L/H_outputBits)
- 2. Initialize a 32-bit, big-endian bit string counter as 1.
- 3. For i = 1 to reps by 1, do the following:
 - Compute Hashi = H(counter || Z || OtherInfo).
 - 2. Increment counter (not to exceed 2^32-1)

- 4. Set key_material = Hash1 || Hash2 || ... so that the length of key_material is L bits, truncating the last block as necessary.
- 5. The above KDF produces a bit string of length L in bits as the keying material. The AS reply key is the output of random-tokey() [RFC3961] using that keying material as the input.

The input parameters for these KDFs are provided as follows:

- o H_outputBits is 160 bits for id-pkinit-kdf-ah-sha1, 256 bits for id-pkinit-kdf-ah-sha256, 384 bits for id-pkinit-kdf-ah-sha384, and 512 bits for id-pkinit-kdf-ah-sha512.
- o max_H_inputBits is 2^64.
- o The secret value (Z) is the shared secret value generated by the Diffie-Hellman exchange. The Diffie-Hellman shared value is first padded with leading zeros such that the size of the secret value in octets is the same as that of the modulus, then represented as a string of octets in big-endian order.
- o The key data length (L) is the key-generation seed length in bits [RFC3961] for the Authentication Service (AS) reply key. The enctype of the AS reply key is selected according to [RFC4120].
- o The algorithm identifier (algorithmID) input parameter is the identifier of the respective KDF. For example, this is id-pkinitkdf-ah-sha1 if the KDF uses SHA-1 as the hash.
- o The initiator identifier (partyUInfo) contains the ASN.1 DER encoding of the KRB5PrincipalName [RFC4556] that identifies the client as specified in the AS-REQ [RFC4120] in the request.
- o The recipient identifier (partyVInfo) contains the ASN.1 DER encoding of the KRB5PrincipalName [RFC4556] that identifies the TGS as specified in the AS-REQ [RFC4120] in the request.
- o The supplemental public information (suppPubInfo) is the ASN.1 DER encoding of the structure PkinitSuppPubInfo as defined later in this section.
- o The supplemental private information (suppPrivInfo) is absent.

OtherInfo is the ASN.1 DER encoding of the following sequence:

```
OtherInfo ::= SEQUENCE {
       algorithmID
                    AlgorithmIdentifier,
        partyUInfo
                       [0] OCTET STRING,
        partyVInfo
                     [1] OCTET STRING,
        suppPubInfo [2] OCTET STRING OPTIONAL,
        suppPrivInfo [3] OCTET STRING OPTIONAL
}
The structure PkinitSuppPubInfo is defined as follows:
PkinitSuppPubInfo ::= SEQUENCE {
      enctype
                         [0] Int32,
           -- The enctype of the AS reply key.
                        [1] OCTET STRING,
           -- The DER encoding of the AS-REQ [RFC4120] from the
      pk-as-rep
                         [2] OCTET STRING,
           -- The DER encoding of the PA-PK-AS-REP [RFC4556] in the
           -- KDC reply.
       . . .
}
```

The PkinitSuppPubInfo structure contains mutually-known public information specific to the authentication exchange. The enctype field is the enctype of the AS reply key as selected according to [RFC4120]. The as-REQ field contains the DER encoding of the type AS-REQ [RFC4120] in the request sent from the client to the KDC. Note that the as-REQ field does not include the wrapping 4 octet length field when TCP is used. The pk-as-rep field contains the DER encoding of the type PA-PK-AS-REP [RFC4556] in the KDC reply. The PkinitSuppPubInfo provides a cryptographic bindings between the preauthentication data and the corresponding ticket request and response, thus addressing the concerns described in Section 3.

The KDF is negotiated between the client and the KDC. The client sends an unordered set of supported KDFs in the request, and the KDC picks one from the set in the reply.

To accomplish this, the AuthPack structure in [RFC4556] is extended as follows:

```
AuthPack ::= SEQUENCE {
       pkAuthenticator
                         [0] PKAuthenticator,
       clientPublicValue [1] SubjectPublicKeyInfo OPTIONAL,
       supportedCMSTypes [2] SEQUENCE OF AlgorithmIdentifier
                OPTIONAL,
       clientDHNonce
                         [3] DHNonce OPTIONAL,
       supportedKDFs [4] SEQUENCE OF KDFAlgorithmId OPTIONAL,
           -- Contains an unordered set of KDFs supported by the
           -- client.
       . . .
}
KDFAlgorithmId ::= SEQUENCE {
       kdf-id
                         [0] OBJECT IDENTIFIER,
           -- The object identifier of the KDF
       . . .
}
```

The new field supportedKDFs contains an unordered set of KDFs supported by the client.

The KDFAlgorithmId structure contains an object identifier that identifies a KDF. The algorithm of the KDF and its parameters are defined by the corresponding specification of that KDF.

The DHRepInfo structure in [RFC4556] is extended as follows:

The new field kdf in the extended DHRepInfo structure identifies the KDF picked by the KDC. If the supportedKDFs field is present in the request, a KDC conforming to this specification MUST choose one of the KDFs supported by the client and indicate its selection in the kdf field in the reply. If the supportedKDFs field is absent in the request, the KDC MUST omit the kdf field in the reply and use the key derivation function from Section 3.2.3.1 of [RFC4556]. If none of the KDFs supported by the client is acceptable to the KDC, the KDC MUST reply with the new error code KDC_ERR_NO_ACCEPTABLE_KDF:

o KDC_ERR_NO_ACCEPTABLE_KDF 100

If the client fills the supportedKDFs field in the request, but the kdf field in the reply is not present, the client can deduce that the KDC is not updated to conform with this specification, or that the exchange was subjected to a downgrade attack. It is a matter of local policy on the client whether to reject the reply when the kdf field is absent in the reply; if compatibility with non-updated KDCs is not a concern, the reply should be rejected.

Implementations conforming to this specification MUST support idpkinit-kdf-ah-sha256.

7. Interoperability

An old client interoperating with a new KDC will not include the supportedKDFs field in the request. The KDC MUST omit the kdf field in the reply and use the [RFC4556] KDF as expected by the client, or reject the request if local policy forbids use of the old KDF.

A new client interoperating with an old KDC will include the supportedKDFs field in the request; this field will be ignored as an unknown extension by the KDC. The KDC will omit the kdf field in the reply and will use the [RFC4556] KDF. The client can deduce from the omitted kdf field that the KDC is not updated to conform to this specification, or that the exchange was subjected to a downgrade attack. The client MUST use the [RFC4556] KDF, or reject the reply if local policy forbids the use of the old KDF.

8. Test vectors

This section contains test vectors for the KDF defined above.

8.1. Common Inputs

```
Z: Length = 256 bytes, Hex Representation = (All Zeros)
client: Length = 9 bytes, ASCII Representation = lha@SU.SE
server: Length = 18 bytes, ASCII Representation = krbtgt/SU.SE@SU.SE
as-req: Length = 10 bytes, Hex Representation =
AAAAAAA AAAAAAA AAAA
pk-as-rep: Length = 9 bytes, Hex Representation =
BBBBBBB BBBBBBB BB
ticket: Length = 55 bytes, Hex Representation =
61353033 A0030201 05A1071B 0553552E 5345A210 300EA003 020101A1 0730051B
036C6861 A311300F A0030201 12A20804 0668656A 68656A
```

8.2. Test Vector for SHA-1, enctype 18

8.2.1. Specific Inputs

```
algorithm-id: (id-pkinit-kdf-ah-sha1) Length = 8 bytes, Hex
Representation = 28060105 02030601
enctype: (aes256-cts-hmac-sha1-96) Length = 1 byte, Decimal
Representation = 18
```

8.2.2. Outputs

```
key-material: Length = 32 bytes, Hex Representation =
E6AB38C9 413E035B B079201E D0B6B73D 8D49A814 A737C04E E6649614 206F73AD
```

key: Length = 32 bytes, Hex Representation = E6AB38C9 413E035B B079201E D0B6B73D 8D49A814 A737C04E E6649614 206F73AD

8.3. Test Vector for SHA-256, enctype

8.3.1. Specific Inputs

algorithm-id: (id-pkinit-kdf-ah-sha256) Length = 8 bytes, Hex
Representation = 2B060105 02030602
enctype: (aes256-cts-hmac-sha1-96) Length = 1 byte, Decimal

Representation = 18

8.3.2. Outputs

key-material: Length = 32 bytes, Hex Representation = 77EF4E48 C420AE3F EC75109D 7981697E ED5D295C 90C62564 F7BFD101 FA9bC1D5

key: Length = 32 bytes, Hex Representation =
77EF4E48 C420AE3F EC75109D 7981697E ED5D295C 90C62564 F7BFD101 FA9bC1D5

8.4. Test Vector for SHA-512, enctype

8.4.1. Specific Inputs

algorithm-id: (id-pkinit-kdf-ah-sha512) Length = 8 bytes, Hex Representation = 2B060105 02030603

enctype: (des3-cbc-sha1-kd) Length = 1 byte, Decimal Representation = 16

8.4.2. Outputs

key-material: Length = 24 bytes, Hex Representation = D3C78A79 D65213EF E9A826F7 5DFB01F7 2362FB16 FB01DAD6

key: Length = 32 bytes, Hex Representation = D3C78A79 D65213EF E9A826F7 5DFB01F7 2362FB16 FB01DAD6

9. Security Considerations

This document describes negotiation of checksum types, key derivation functions and other cryptographic functions. If a given negotiation is unauthenticated, care must be taken to accept only secure values; to do otherwise allows an active attacker to perform a downgrade attack.

10. Acknowledgements

Jeffery Hutzelman, Shawn Emery, Tim Polk, Kelley Burgin, Ben Kaduk, and Scott Bradner reviewed the document and provided suggestions for improvements.

11. IANA Considerations

IANA is requested to update the following registrations in the Kerberos Pre-authentication and Typed Data Registry created by section 7.1 of RFC 6113 to refer to this specification. These values were reserved for this specification in the initial registrations.

TD-CMS-DIGEST-ALGORITHMS 111 [ALG-AGILITY]
TD-CERT-DIGEST-ALGORITHMS 112 [ALG-AGILITY]

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Appendix A. PKINIT ASN.1 Module

```
KerberosV5-PK-INIT-Agility-SPEC {
       iso(1) identified-organization(3) dod(6) internet(1)
       security(5) kerberosV5(2) modules(4) pkinit(5) agility (1)
} DEFINITIONS EXPLICIT TAGS ::= BEGIN
IMPORTS
  AlgorithmIdentifier, SubjectPublicKeyInfo
       FROM PKIX1Explicit88 { iso (1)
         identified-organization (3) dod (6) internet (1)
         security (5) mechanisms (5) pkix (7) id-mod (0)
         id-pkix1-explicit (18) }
         -- As defined in <u>RFC 5280</u>.
  Ticket, Int32, Realm, EncryptionKey, Checksum
       FROM KerberosV5Spec2 { iso(1) identified-organization(3)
         dod(6) internet(1) security(5) kerberosV5(2)
         modules(4) krb5spec2(2) }
         -- as defined in RFC 4120.
  PKAuthenticator, DHNonce, id-pkinit
       FROM KerberosV5-PK-INIT-SPEC {
         iso(1) identified-organization(3) dod(6) internet(1)
         security(5) kerberosV5(2) modules(4) pkinit(5) };
```

```
-- as defined in RFC 4556.
id-pkinit-kdf OBJECT IDENTIFIER ::= { id-pkinit kdf(6) }
    -- PKINIT KDFs
id-pkinit-kdf-ah-sha1 OBJECT IDENTIFIER
    ::= { id-pkinit-kdf sha1(1) }
    -- SP800-56A ASN.1 structured hash-based KDF using SHA-1
id-pkinit-kdf-ah-sha256 OBJECT IDENTIFIER
    ::= { id-pkinit-kdf sha256(2) }
    -- SP800-56A ASN.1 structured hash-based KDF using SHA-256
id-pkinit-kdf-ah-sha512 OBJECT IDENTIFIER
    ::= { id-pkinit-kdf sha512(3) }
    -- SP800-56A ASN.1 structured hash-based KDF using SHA-512
id-pkinit-kdf-ah-sha384 OBJECT IDENTIFIER
    ::= { id-pkinit-kdf sha384(4) }
    -- SP800-56A ASN.1 structured hash-based KDF using SHA-384
TD-CMS-DIGEST-ALGORITHMS-DATA ::= SEQUENCE OF
    AlgorithmIdentifier
        -- Contains the list of CMS algorithm [RFC5652]
        -- identifiers indicating the digest algorithms
        -- acceptable to the KDC for signing CMS data in
        -- the order of decreasing preference.
TD-CERT-DIGEST-ALGORITHMS-DATA ::= SEQUENCE {
       allowedAlgorithms [0] SEQUENCE OF AlgorithmIdentifier,
           -- Contains the list of CMS algorithm [RFC5652]
           -- identifiers indicating the digest algorithms
           -- that are used by the CA to sign the client's
           -- X.509 certificate and are acceptable to the KDC
           -- in the process of validating the client's X.509
           -- certificate, in the order of decreasing
           -- preference.
       rejectedAlgorithm [1] AlgorithmIdentifier OPTIONAL,
           -- This identifies the digest algorithm that was
           -- used to sign the client's X.509 certificate and
           -- has been rejected by the KDC in the process of
           -- validating the client's X.509 certificate
           -- [RFC5280].
       . . .
}
OtherInfo ::= SEQUENCE {
        algorithmID AlgorithmIdentifier,
```

```
partyUInfo
                       [0] OCTET STRING,
        partyVInfo
                       [1] OCTET STRING,
        suppPubInfo
                      [2] OCTET STRING OPTIONAL,
        suppPrivInfo [3] OCTET STRING OPTIONAL
}
PkinitSuppPubInfo ::= SEQUENCE {
       enctype
                         [0] Int32,
          -- The enctype of the AS reply key.
                         [1] OCTET STRING,
       as-REQ
           -- The DER encoding of the AS-REQ [RFC4120] from the
           -- client.
                         [2] OCTET STRING,
       pk-as-rep
           -- The DER encoding of the PA-PK-AS-REP [RFC4556] in the
           -- KDC reply.
       . . .
}
AuthPack ::= SEQUENCE {
       pkAuthenticator
                        [0] PKAuthenticator,
       clientPublicValue [1] SubjectPublicKeyInfo OPTIONAL,
       supportedCMSTypes [2] SEQUENCE OF AlgorithmIdentifier
                OPTIONAL,
       clientDHNonce
                         [3] DHNonce OPTIONAL,
       supportedKDFs [4] SEQUENCE OF KDFAlgorithmId OPTIONAL,
           -- Contains an unordered set of KDFs supported by the
           -- client.
}
KDFAlgorithmId ::= SEQUENCE {
       kdf-id
                         [0] OBJECT IDENTIFIER,
           -- The object identifier of the KDF
       . . .
}
DHRepInfo ::= SEQUENCE {
       dhSignedData
                       [0] IMPLICIT OCTET STRING,
       serverDHNonce [1] DHNonce OPTIONAL,
       . . . ,
       kdf
                         [2] KDFAlgorithmId OPTIONAL,
           -- The KDF picked by the KDC.
}
END
```

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