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Ephemeral Diffie-Hellman Over COSE (EDHOC)

Abstract

This document specifies Ephemeral Diffie-Hellman Over COSE (EDHOC), a very compact and lightweight authenticated Diffie-Hellman key exchange with ephemeral keys. EDHOC provides mutual authentication, forward secrecy, and identity protection. EDHOC is intended for usage in constrained scenarios and a main use case is to establish an OSCORE security context. By reusing COSE for cryptography, CBOR for encoding, and CoAP for transport, the additional code size can be kept very low.

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1. Introduction

1.1. Motivation

Many Internet of Things (IoT) deployments require technologies which
are highly performant in constrained environments [RFC7228]. IoT
devices may be constrained in various ways, including memory,
storage, processing capacity, and power. The connectivity for these
settings may also exhibit constraints such as unreliable and lossy
channels, highly restricted bandwidth, and dynamic topology. The
IETF has acknowledged this problem by standardizing a range of
lightweight protocols and enablers designed for the IoT, including
the Constrained Application Protocol (CoAP, [RFC7252]), Concise

Binary Object Representation (CBOR, [RFC8949]), and Static Context Header Compression (SCHC, [RFC8724]).

The need for special protocols targeting constrained IoT deployments extends also to the security domain [I-D.ietf-lake-reqs]. Important characteristics in constrained environments are the number of round trips and protocol message sizes, which if kept low can contribute to good performance by enabling transport over a small number of radio frames, reducing latency due to fragmentation or duty cycles, etc. Another important criterion is code size, which may be prohibitively large for certain deployments due to device capabilities or network load during firmware update. Some IoT deployments also need to support a variety of underlying transport technologies, potentially even with a single connection.

Some security solutions for such settings exist already. CBOR Object Signing and Encryption (COSE, [RFC9052]) specifies basic application-layer security services efficiently encoded in CBOR. Another example is Object Security for Constrained RESTful Environments (OSCORE, [RFC8613]) which is a lightweight communication security extension to CoAP using CBOR and COSE. In order to establish good quality cryptographic keys for security protocols such as COSE and OSCORE, the two endpoints may run an authenticated Diffie-Hellman key exchange protocol, from which shared secret keying material can be derived. Such a key exchange protocol should also be lightweight; to prevent bad performance in case of repeated use, e.g., due to device rebooting or frequent rekeying for security reasons; or to avoid latencies in a network formation setting with many devices authenticating at the same time.

This document specifies Ephemeral Diffie-Hellman Over COSE (EDHOC), a lightweight authenticated key exchange protocol providing good security properties including forward secrecy, identity protection, and cipher suite negotiation. Authentication can be based on raw public keys (RPK) or public key certificates and requires the application to provide input on how to verify that endpoints are trusted. This specification supports the referencing of credentials in order to reduce message overhead, but credentials may alternatively be embedded in the messages. EDHOC does not currently support pre-shared key (PSK) authentication as authentication with static Diffie-Hellman public keys by reference produces equally small message sizes but with much simpler key distribution and identity protection.

EDHOC makes use of known protocol constructions, such as SIGMA [SIGMA], the Noise XX pattern [Noise], and Extract-and-Expand [RFC5869]. EDHOC uses COSE for cryptography and identification of credentials (including COSE_Key, CBOR Web Token (CWT), CWT Claims Set (CCS), X.509, and CBOR encoded X.509 (C509) certificates, see

Section 3.5.2). COSE provides crypto agility and enables the use of
future algorithms and credential types targeting IoT.

EDHOC is designed for highly constrained settings making it
especially suitable for low-power networks [RFC8376] such as
Cellular IoT, 6TiSCH, and LoRaWAN. A main objective for EDHOC is to
be a lightweight authenticated key exchange for OSCORE, i.e., to
provide authentication and session key establishment for IoT use
cases such as those built on CoAP [RFC7252] involving 'things' with
embedded microcontrollers, sensors, and actuators. By reusing the
same lightweight primitives as OSCORE (CBOR, COSE, CoAP) the
additional code size can be kept very low. Note that while CBOR and
COSE primitives are built into the protocol messages, EDHOC is not
bound to a particular transport.

A typical setting is when one of the endpoints is constrained or in
a constrained network, and the other endpoint is a node on the
Internet (such as a mobile phone). Thing-to-thing interactions over
constrained networks are also relevant since both endpoints would
then benefit from the lightweight properties of the protocol. EDHOC
could, e.g., be run when a device connects for the first time, or to
establish fresh keys which are not revealed by a later compromise of
the long-term keys.

1.2. Message Size Examples

Examples of EDHOC message sizes are shown in Figure 1, using
different kinds of authentication keys and COSE header parameters
for identification: static Diffie-Hellman keys or signature keys,
either in CBOR Web Token (CWT) / CWT Claims Set (CCS) [RFC8392]
identified by a key identifier using 'kid' [RFC9052], or in X.509
certificates identified by a hash value using 'x5t'
[I-D.ietf-cose-x509]. As a comparison, in the case of RPK
authentication, the EDHOC message size when transferred in CoAP can
be less than 1/7 of the DTLS 1.3 handshake [RFC9147] with ECDHE and
connection ID, see Section 2 of
[I-D.ietf-lwig-security-protocol-comparison].

	Static DH Keys		Signature Keys	
	kid	x5t	kid	x5t
message 1	37	37	37	37
message 2	45	58	102	115
message 3	19	33	77	90
Total	101	128	216	242

Figure 1: Examples of EDHOC message sizes in bytes.

1.3. Document Structure

The remainder of the document is organized as follows: Section 2
outlines EDHOC authenticated with signature keys, Section 3
describes the protocol elements of EDHOC, including formatting of
the ephemeral public keys, Section 4 specifies the key derivation,
Section 5 specifies message processing for EDHOC authenticated with
signature keys or static Diffie-Hellman keys, Section 6 describes
the error messages, and Section 7 lists compliance requirements.
Note that normative text is also used in appendices, in particular
Appendix A.

1.4. Terminology and Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT",
"SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and
"OPTIONAL" in this document are to be interpreted as described in
BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all
capitals, as shown here.

Readers are expected to be familiar with the terms and concepts
described in CBOR [RFC8949], CBOR Sequences [RFC8742], COSE
structures and processing [RFC9052], COSE algorithms [RFC9053], CWT
and CWT Claims Set [RFC8392], and the Concise Data Definition
Language (CDDL, [RFC8610]), which is used to express CBOR data
structures. Examples of CBOR and CDDL are provided in Appendix C.1.
When referring to CBOR, this specification always refers to
Deterministically Encoded CBOR as specified in Sections 4.2.1 and
4.2.2 of [RFC8949]. The single output from authenticated encryption
(including the authentication tag) is called "ciphertext", following
[RFC5116].

2. EDHOC Outline

EDHOC specifies different authentication methods of the ephemeral-ephemeral Diffie-Hellman key exchange: signature keys and static Diffie-Hellman keys. This section outlines the signature key based method. Further details of protocol elements and other authentication methods are provided in the remainder of this document.

SIGMA (SIGn-and-MAc) is a family of theoretical protocols with a large number of variants [SIGMA]. Like in IKEv2 [RFC7296] and (D)TLS 1.3 [RFC8446][RFC9147], EDHOC authenticated with signature keys is built on a variant of the SIGMA protocol, SIGMA-I, which provides identity protection against active attacks on the party initiating the protocol. Also like IKEv2, EDHOC implements the MAC-then-Sign variant of the SIGMA-I protocol. The message flow (excluding an optional fourth message) is shown in Figure 2.

```
Initiator                                                  Responder
                              G X
 |─────────────────────────────────────────────────────────────►|
 |                                                               |
 |      G Y. Enc( ID CRED R. Sig( R: MAC( CRED R. G X. G Y ) ) ) |
 |◄─────────────────────────────────────────────────────────────|
 |                                                               |
 |        AEAD( ID CRED I. Sig( I: MAC( CRED I. G Y. G X ) ) )   |
 |─────────────────────────────────────────────────────────────►|
 |                                                               |
```
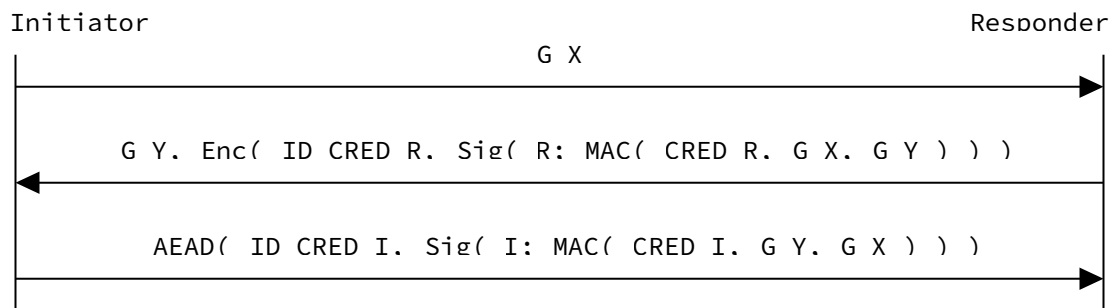
Figure 2: MAC-then-Sign variant of the SIGMA-I protocol used by EDHOC method 0.

The parties exchanging messages in an EDHOC session are called Initiator (I) and Responder (R), where the Initiator sends message_1 (see Section 3). They exchange ephemeral public keys, compute a shared secret session key PRK_out, and derive symmetric application keys used to protect application data.

*G_X and G_Y are the ECDH ephemeral public keys of I and R, respectively.

*CRED_I and CRED_R are the authentication credentials containing the public authentication keys of I and R, respectively.

*ID_CRED_I and ID_CRED_R are used to identify and optionally transport the credentials of the Initiator and the Responder, respectively.

*Sig(I; .) and Sig(R; .) denote signatures made with the private
   authentication key of I and R, respectively.

  *Enc(), AEAD(), and MAC() denotes encryption, authenticated
   encryption with additional data, and message authentication code
   - crypto algorithms applied with keys derived from one or more
   shared secrets calculated during the protocol.

In order to create a "full-fledged" protocol some additional
protocol elements are needed. EDHOC adds:

  *Transcript hashes (hashes of message data) TH_2, TH_3, TH_4 used
   for key derivation and as additional authenticated data.

  *Computationally independent keys derived from the ECDH shared
   secret and used for authenticated encryption of different
   messages.

  *An optional fourth message giving key confirmation to I in
   deployments where no protected application data is sent from R to
   I.

  *A keying material exporter and a key update function with forward
   secrecy.

  *Secure negotiation of cipher suite.

  *Method types, error handling, and padding.

  *Selection of connection identifiers C_I and C_R which may be used
   in EDHOC to identify protocol state.

  *Transport of external authorization data.

EDHOC is designed to encrypt and integrity protect as much
information as possible. Symmetric keys and random material derived
using EDHOC_KDF are derived with as much previous information as
possible, see Figure 8. EDHOC is furthermore designed to be as
compact and lightweight as possible, in terms of message sizes,
processing, and the ability to reuse already existing CBOR, COSE,
and CoAP libraries. Like in (D)TLS, authentication is the
responsibility of the application. EDHOC identifies (and optionally
transports) authentication credentials, and provides proof-of-
possession of the private authentication key.

To simplify for implementors, the use of CBOR and COSE in EDHOC is
summarized in Appendix C. Test vectors including CBOR diagnostic
notation are provided in [I-D.ietf-lake-traces].

3. Protocol Elements

3.1. General

The EDHOC protocol consists of three mandatory messages (message_1, message_2, message_3), an optional fourth message (message_4), and an error message, between an Initiator (I) and a Responder (R). The odd messages are sent by I, the even by R. Both I and R can send error messages. The roles have slightly different security properties which should be considered when the roles are assigned, see Section 8.1. All EDHOC messages are CBOR Sequences [RFC8742], and are deterministically encoded. Figure 3 illustrates an EDHOC message flow with the optional fourth message as well as the content of each message. The protocol elements in the figure are introduced in Section 3 and Section 5. Message formatting and processing are specified in Section 5 and Section 6.

Application data may be protected using the agreed application algorithms (AEAD, hash) in the selected cipher suite (see Section 3.6) and the application can make use of the established connection identifiers C_I and C_R (see Section 3.3). Media types that may be used for EDHOC are defined in Section 9.8.

The Initiator can derive symmetric application keys after creating EDHOC message_3, see Section 4.2.1. Protected application data can therefore be sent in parallel or together with EDHOC message_3. EDHOC message_4 is typically not sent.

```
Initiator                                                      Responder
                    METHOD, SUITES I, G X, C I, EAD 1
|------------------------------------------------------------------->|
                              message 1

           G Y, Enc( ID CRED R, Signature or MAC 2, EAD 2 ), C R
|<-------------------------------------------------------------------|
                              message 2

              AEAD( ID CRED I, Signature or MAC 3, EAD 3 )
|------------------------------------------------------------------->|
                              message 3

                            AEAD( EAD 4 )
|<- - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - - -|
                              message 4
```
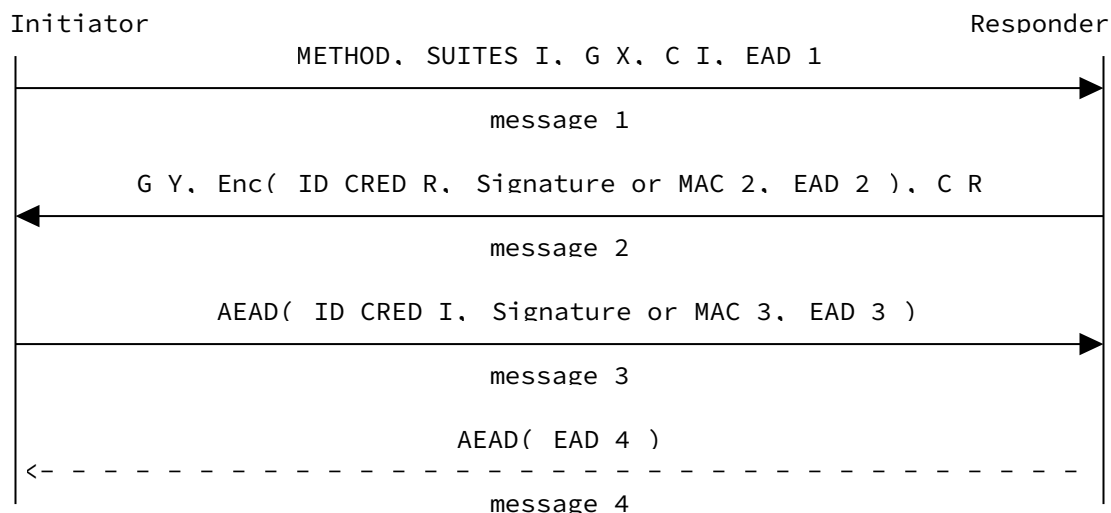
Figure 3: EDHOC message flow including the optional fourth message.

3.2. Method

The data item METHOD in message_1 (see Section 5.2.1), is an integer
specifying the authentication method. EDHOC supports authentication
with signature or static Diffie-Hellman keys, as defined in the four
authentication methods: 0, 1, 2, and 3, see Figure 4. When using a
static Diffie-Hellman key the authentication is provided by a
Message Authentication Code (MAC) computed from an ephemeral-static
ECDH shared secret which enables significant reductions in message
sizes. Note that also in the static Diffie-Hellman based
authentication methods there is an ephemeral-ephemeral Diffie-
Hellman key exchange.

The Initiator and the Responder need to have agreed on a single
method to be used for EDHOC, see Section 3.9.

Method Type Value	Initiator Authentication Key	Responder Authentication Key
0	Signature Key	Signature Key
1	Signature Key	Static DH Key
2	Static DH Key	Signature Key
3	Static DH Key	Static DH Key

Figure 4: Authentication Keys for Method Types

EDHOC does not have a dedicated message field to indicate the
protocol version. Breaking changes to EDHOC can be introduced by
specifying and registering new methods.

3.3. Connection Identifiers

EDHOC includes the selection of connection identifiers (C_I, C_R)
identifying a connection for which keys are agreed.

Connection identifiers may be used to correlate EDHOC messages and
facilitate the retrieval of protocol state during an EDHOC session
(see Section 3.4), or may be used in applications of EDHOC, e.g., in
OSCORE (see Section 3.3.3). The connection identifiers do not have
any cryptographic purpose in EDHOC except facilitating the retrieval
of security data associated with the protocol state.

Connection identifiers in EDHOC are intrinsically byte strings. Most
constrained devices only have a few connections for which short
identifiers may be sufficient. In some cases minimum length
identifiers are necessary to comply with overhead requirements.

However, CBOR byte strings - with the exception of the empty byte
string h'' which encodes as one byte (0x40) - are encoded as two or
more bytes. To enable one-byte encoding of certain byte strings
while maintaining CBOR encoding, EDHOC represents certain
identifiers as CBOR integers on the wire, see Section 3.3.2.

3.3.1. Selection of Connection Identifiers

C_I and C_R are chosen by I and R, respectively. The Initiator
selects C_I and sends it in message_1 for the Responder to use as a
reference to the connection in communication with the Initiator. The
Responder selects C_R and sends it in message_2 for the Initiator to
use as a reference to the connection in communications with the
Responder.

If connection identifiers are used by an application protocol for
which EDHOC establishes keys then the selected connection
identifiers SHALL adhere to the requirements for that protocol, see
Section 3.3.3 for an example.

3.3.2. Representation of Byte String Identifiers

To allow identifiers with minimal overhead on the wire, certain byte
strings are defined to have integer representations.

The integers with one-byte CBOR encoding are -24, ..., 23, see
Figure 5.

```
 Integer:                 -24  -23   ...   -2   -1    0    1   ...   23
 CBOR encoding (1 byte):   37   36   ...   21   20   00   01   ...   17
```

Figure 5: One-Byte CBOR Encoded Integers

The byte strings which coincide with a one-byte CBOR encoding of an
integer MUST be represented by the CBOR encoding of that integer.
Other byte strings are simply encoded as CBOR byte strings.

For example:

  *0x21 is represented by 0x21 (CBOR encoding of the integer -2),
   not by 0x4121 (CBOR encoding of the byte string 0x21).

  *0x0D is represented by 0x0D (CBOR encoding of the integer 13),
   not by 0x410D (CBOR encoding of the byte string 0x0D).

  *0x18 is represented by 0x4118 (CBOR encoding of the byte string
   0x18).

*0x38 is represented by 0x4138 (CBOR encoding of the byte string 0x38).

*0xABCD is represented by 0x42ABCD (CBOR encoding of the byte string 0xABCD).

One way to view this representation of byte strings as a transport encoding: a byte string which parses as the one-byte CBOR encoding of an integer (i.e., integer in the interval -24, ..., 23) is just copied directly into the message, a byte string which does not is encoded as a CBOR byte string during transport.

Implementation Note: When implementing the byte string identifier representation, it can in some programming languages help to define a new type, or other data structure, which (in its user facing API) behaves like a byte string, but when serializing to CBOR produces a CBOR byte string or a CBOR integer depending on its value.

3.3.3. Use of Connection Identifiers with OSCORE

For OSCORE, the choice of connection identifier results in the endpoint selecting its Recipient ID, see Section 3.1 of [RFC8613], for which certain uniqueness requirements apply, see Section 3.3 of [RFC8613]. Therefore, the Initiator and the Responder MUST NOT select connection identifiers such that it results in the same OSCORE Recipient ID. Since the connection identifier is a byte string, it is converted to an OSCORE Recipient ID equal to the byte string.

Examples:

*A connection identifier 0xFF (represented in the EDHOC message as 0x41FF, see Section 3.3.2) is converted to the OSCORE Recipient ID 0xFF.

*A connection identifier 0x21 (represented in the EDHOC message as 0x21, see Section 3.3.2) is converted to the OSCORE Recipient ID 0x21.

3.4. Transport

Cryptographically, EDHOC does not put requirements on the lower layers. EDHOC is not bound to a particular transport layer and can even be used in environments without IP. In addition to the transport of messages including errors, the transport is responsible, where necessary, to handle:

*message loss,

*message reordering,

*message duplication,

*flow control,

*fragmentation and reassembly,

*demultiplex EDHOC messages from other types of messages,

*denial-of-service protection,

*message correlation, see Section 3.4.1.

EDHOC does not require error free transport since a change in message content is detected through the transcript hashes in a subsequent integrity verification, see Section 5.

EDHOC is designed to enable an authenticated key exchange with small messages, where the minimum message sizes are of the order illustrated in the first column of Figure 1. There is no maximum message size specified by the protocol; this is for example dependent on the size of authentication credentials (if they are transported, see Section 3.5).

The use of transport is specified in the application profile, which in particular may specify limitations in message sizes, see Section 3.9.

3.4.1. EDHOC Message Correlation

Correlation between EDHOC messages is needed to facilitate the retrieval of protocol state and security context during an EDHOC session. It is also helpful for the Responder to get an indication that a received EDHOC message is the beginning of a new session, such that no existing protocol state or security context needs to be retrieved.

Correlation may be based on existing mechanisms in the transport protocol, for example, the CoAP Token may be used to correlate EDHOC messages in a CoAP response and in an associated CoAP request. The connection identifiers may also be used to correlate EDHOC messages.

If correlation between consecutive messages is not provided by other means then the transport binding SHOULD mandate prepending of an appropriate connection identifier (when available from the EDHOC protocol) to the EDHOC message. If message_1 indication is not provided by other means, then the transport binding SHOULD mandate prepending of message_1 with the CBOR simple value true (0xf5), as a unique dummy identifier.

Transport of EDHOC in CoAP payloads is described in [Appendix A.2](#), including how to use connection identifiers and message_1 indication with CoAP. A similar construction is possible for other client-server protocols. Protocols that do not provide any correlation at all can prescribe prepending of the peer's connection identifier to all messages.

Note that correlation between EDHOC messages may be obtained without transport support or connection identifiers, for example if the endpoints only accept a single instance of the protocol at a time, and execute conditionally on a correct sequence of messages.

3.5. Authentication Parameters

EDHOC supports various settings for how the other endpoint's authentication (public) key may be transported, identified, and trusted.

EDHOC performs the following authentication related operations:

* *EDHOC transports information about credentials in ID_CRED_I and ID_CRED_R (described in [Section 3.5.3](#)). Based on this information, the authentication credentials CRED_I and CRED_R (described in [Section 3.5.2](#)) can be obtained. EDHOC may also transport certain authentication related information as External Authorization Data (see [Section 3.8](#)).

* *EDHOC uses the authentication credentials in two ways (see [Section 5.3.2](#) and [Section 5.4.2](#)):

  - The authentication credential is input to the integrity verification using the MAC fields.

  - The authentication key of the authentication credential is used with the Signature_or_MAC field to verify proof-of-possession of the private key.

Other authentication related verifications are out of scope for EDHOC, and is the responsibility of the application. In particular, the authentication credential needs to be validated in the context of the connection for which EDHOC is used, see [Appendix D](#). EDHOC MUST allow the application to read received information about credential (ID_CRED_R, ID_CRED_I). EDHOC MUST have access to the authentication key and the authentication credential.

Note that the type of authentication key, authentication credential, and the identification of the credential have a large impact on the message size. For example, the Signature_or_MAC field is much smaller with a static DH key than with a signature key. A CCS is much smaller than a self-signed certificate/CWT, but if it is

possible to reference the credential with a COSE header like 'kid', then that is in turn much smaller than a CCS.

3.5.1. Authentication Keys

The authentication key (i.e., the public key used for authentication) MUST be a signature key or static Diffie-Hellman key. The Initiator and the Responder MAY use different types of authentication keys, e.g., one uses a signature key and the other uses a static Diffie-Hellman key.

The authentication key algorithm needs to be compatible with the method and the selected cipher suite (see Section 3.6). The authentication key algorithm needs to be compatible with the EDHOC key exchange algorithm when static Diffie-Hellman authentication is used, and compatible with the EDHOC signature algorithm when signature authentication is used.

Note that for most signature algorithms, the signature is determined by the signature algorithm and the authentication key algorithm together. When using static Diffie-Hellman keys the Initiator's and Responder's private authentication keys are denoted as I and R, respectively, and the public authentication keys are denoted G_I and G_R, respectively.

For X.509 certificates the authentication key is represented by a SubjectPublicKeyInfo field. For CWT and CCS (see Section 3.5.2)) the authentication key is represented by a 'cnf' claim [RFC8747] containing a COSE_Key [RFC9052]. In EDHOC, a raw public key (RPK) is an authentication key encoded as a COSE_Key wrapped in a CCS.

3.5.2. Authentication Credentials

The authentication credentials, CRED_I and CRED_R, contain the public authentication key of the Initiator and the Responder, respectively. The authentication credential typically also contains other parameters that needs to be verified by the application, see Appendix D, and in particular information about the identity ("subject") of the endpoint to prevent misbinding attacks, see Appendix D.2.

EDHOC relies on COSE for identification of credentials (see Section 3.5.3), for example X.509 certificates [RFC5280], C509 certificates [I-D.ietf-cose-cbor-encoded-cert], CWTs [RFC8392] and CWT Claims Sets (CCS) [RFC8392]. When the identified credential is a chain or a bag, the authentication credential CRED_x is just the end entity X.509 or C509 certificate / CWT. The Initiator and the Responder MAY use different types of authentication credentials, e.g., one uses an RPK and the other uses a public key certificate.

Since CRED_R is used in the integrity verification, see
Section 5.3.2, it needs to be specified such that it is identical
when used by Initiator or Responder. Similarly for CRED_I, see
Section 5.4.2. The Initiator and Responder are expected to agree on
the specific encoding of the credentials, see Section 3.9.

It is RECOMMENDED that the COSE 'kid' parameter, when used to
identify the authentication credential, refers to a specific
encoding. The Initiator and Responder SHOULD use an available
authentication credential (transported in EDHOC or otherwise
provisioned) without re-encoding. If for some reason re-encoding of
the authentication credential may occur, then a potential common
encoding for CBOR based credentials is bytewise lexicographic order
of their deterministic encodings as specified in Section 4.2.1 of
[RFC8949].

  *When the authentication credential is an X.509 certificate,
   CRED_x SHALL be the DER encoded certificate, encoded as a bstr
   [I-D.ietf-cose-x509].

  *When the authentication credential is a C509 certificate, CRED_x
   SHALL be the C509Certificate [I-D.ietf-cose-cbor-encoded-cert].

  *When the authentication credential is a CWT including a COSE_Key,
   CRED_x SHALL be the untagged CWT.

  *When the authentication credential includes a COSE_Key but is not
   in a CWT, CRED_x SHALL be an untagged CCS. This is how RPKs are
   encoded, see Figure 6 for an example.

     -Naked COSE_Keys are thus dressed as CCS when used in EDHOC, in
      its simplest form by prefixing the COSE_Key with 0xA108A101 (a
      map with a 'cnf' claim). In that case the resulting
      authentication credential contains no other identity than the
      public key itself, see Appendix D.2.

An example of a CRED_x is shown below:

```
{                                           /CCS/
  2 : "42-50-31-FF-EF-37-32-39",            /sub/
  8 : {                                     /cnf/
    1 : {                                   /COSE_Key/
      1 : 1,                                /kty/
      2 : h'00',                            /kid/
     -1 : 4,                                /crv/
     -2 : h'b1a3e89460e88d3a8d54211dc95f0b90   /x/
          3ff205eb71912d6db8f4af980d2db83a'
    }
  }
}
```

Figure 6: CWT Claims Set (CCS) containing an X25519 static Diffie-
Hellman key and an EUI-64 identity.

3.5.3. Identification of Credentials

ID_CRED_R and ID_CRED_I are transported in message_2 and message_3,
respectively, see Section 5.3.2 and Section 5.4.2. They are used to
identify and optionally transport credentials:

* ID_CRED_R is intended to facilitate for the Initiator retrieving
  the authentication credential CRED_R and the authentication key
  of R.

* ID_CRED_I is intended to facilitate for the Responder retrieving
  the authentication credential CRED_I and the authentication key
  of I.

ID_CRED_x may contain the authentication credential CRED_x, but for
many settings it is not necessary to transport the authentication
credential within EDHOC. For example, it may be pre-provisioned or
acquired out-of-band over less constrained links. ID_CRED_I and
ID_CRED_R do not have any cryptographic purpose in EDHOC since the
authentication credentials are integrity protected.

EDHOC relies on COSE for identification of credentials and supports
all credential types for which COSE header parameters are defined
including X.509 certificates ([I-D.ietf-cose-x509]), C509
certificates ([I-D.ietf-cose-cbor-encoded-cert]), CWT (Section 9.6)
and CWT Claims Set (CCS) (Section 9.6).

ID_CRED_I and ID_CRED_R are of type COSE header_map, as defined in
Section 3 of [RFC9052], and contains one or more COSE header
parameters. ID_CRED_I and ID_CRED_R MAY contain different header
parameters. The header parameters typically provide some information
about the format of the credential.

Example: X.509 certificates can be identified by a hash value using
the 'x5t' parameter, see Section 2 of [I-D.ietf-cose-x509]:

  *ID_CRED_x = { 34 : COSE_CertHash }, for x = I or R,

Example: CWT or CCS can be identified by a key identifier using the
'kid' parameter, see Section 3.1 of [RFC9052]:

  *ID_CRED_x = { 4 : kid_x }, where kid_x : kid, for x = I or R.

Note that COSE header parameters in ID_CRED_x are used to identify
the message sender's credential. There is therefore no reason to use
the "-sender" header parameters, such as x5t-sender, defined in
Section 3 of [I-D.ietf-cose-x509]. Instead, the corresponding
parameter without "-sender", such as x5t, SHOULD be used.

As stated in Section 3.1 of [RFC9052], applications MUST NOT assume
that 'kid' values are unique and several keys associated with a
'kid' may need to be checked before the correct one is found.
Applications might use additional information such as 'kid context'
or lower layers to determine which key to try first. Applications
should strive to make ID_CRED_x as unique as possible, since the
recipient may otherwise have to try several keys.

See Appendix C.3 for more examples.

3.5.3.1. Compact Encoding of ID_CRED for 'kid'

To comply with the LAKE message size requirements, see
[I-D.ietf-lake-reqs], two optimizations are made for the case when
ID_CRED_x contains a single 'kid' parameter.

  1. The CBOR map { 4 : kid_x } is replaced by the byte string
     kid_x.

  2. The representation of identifiers specified in Section 3.3.2 is
     applied to kid_x.

These optimizations MUST be applied if and only if ID_CRED_x = { 4 :
kid_x } and ID_CRED_x in PLAINTEXT_y of message_y, y = 2 or 3, see
Section 5.3.2 and Section 5.4.2. Note that these optimizations are
not applied to instances of ID_CRED_x which have no impact on
message size, e.g., context_y, or the COSE protected header.
Examples:

  *For ID_CRED_x = { 4 : h'FF' }, the encoding in PLAINTEXT_y is not
   the CBOR map 0xA10441FF but the CBOR byte string h'FF', i.e.,
   0x41FF.

*For ID_CRED_x = { 4 : h'21' }, the encoding in PLAINTEXT_y is
 neither the CBOR map 0xA1044121, nor the CBOR byte string h'21',
 i.e., 0x4121, but the CBOR integer 0x21.

3.6. Cipher Suites

An EDHOC cipher suite consists of an ordered set of algorithms from
the "COSE Algorithms" and "COSE Elliptic Curves" registries as well
as the EDHOC MAC length. All algorithm names and definitions follow
from COSE algorithms [RFC9053]. Note that COSE sometimes uses
peculiar names such as ES256 for ECDSA with SHA-256, A128 for
AES-128, and Ed25519 for the curve edwards25519. Algorithms need to
be specified with enough parameters to make them completely
determined. The EDHOC MAC length MUST be at least 8 bytes. Any
cryptographic algorithm used in the COSE header parameters in
ID_CRED is selected independently of the selected cipher suite.
EDHOC is currently only specified for use with key exchange
algorithms of type ECDH curves, but any Key Encapsulation Method
(KEM), including Post-Quantum Cryptography (PQC) KEMs, can be used
in method 0, see Section 8.4. Use of other types of key exchange
algorithms to replace static DH authentication (method 1,2,3) would
likely require a specification updating EDHOC with new methods.

EDHOC supports all signature algorithms defined by COSE. Just like
in (D)TLS 1.3 [RFC8446][RFC9147] and IKEv2 [RFC7296], a signature in
COSE is determined by the signature algorithm and the authentication
key algorithm together, see Section 3.5.1. The exact details of the
authentication key algorithm depend on the type of authentication
credential. COSE supports different formats for storing the public
authentication keys including COSE_Key and X.509, which use
different names and ways to represent the authentication key and the
authentication key algorithm.

An EDHOC cipher suite consists of the following parameters:

  *EDHOC AEAD algorithm

  *EDHOC hash algorithm

  *EDHOC MAC length in bytes (Static DH)

  *EDHOC key exchange algorithm (ECDH curve)

  *EDHOC signature algorithm

  *Application AEAD algorithm

  *Application hash algorithm

Each cipher suite is identified with a pre-defined integer label.

EDHOC can be used with all algorithms and curves defined for COSE. Implementations can either use any combination of COSE algorithms and parameters to define their own private cipher suite, or use one of the pre-defined cipher suites. Private cipher suites can be identified with any of the four values -24, -23, -22, -21. The pre-defined cipher suites are listed in the IANA registry ([Section 9.2](#)) with initial content outlined here:

* *Cipher suites 0-3, based on AES-CCM, are intended for constrained IoT where message overhead is a very important factor. Note that AES-CCM-16-64-128 and AES-CCM-16-128-128 are compatible with the IEEE CCM* mode.

    - Cipher suites 1 and 3 use a larger tag length (128-bit) in EDHOC than in the Application AEAD algorithm (64-bit).

* *Cipher suites 4 and 5, based on ChaCha20, are intended for less constrained applications and only use 128-bit tag lengths.

* *Cipher suite 6, based on AES-GCM, is for general non-constrained applications. It consists of high performance algorithms that are widely used in non-constrained applications.

* *Cipher suites 24 and 25 are intended for high security applications such as government use and financial applications. These cipher suites do not share any algorithms. Cipher suite 24 consists of algorithms from the CNSA suite [[CNSA](#)].

The different methods ([Section 3.2](#)) use the same cipher suites, but some algorithms are not used in some methods. The EDHOC signature algorithm is not used in methods without signature authentication.

The Initiator needs to have a list of cipher suites it supports in order of preference. The Responder needs to have a list of cipher suites it supports. SUITES_I contains cipher suites supported by the Initiator, formatted and processed as detailed in [Section 5.2.1](#) to secure the cipher suite negotiation. Examples of cipher suite negotiation are given in [Section 6.3.2](#).

3.7. Ephemeral Public Keys

The ephemeral public keys in EDHOC (G_X and G_Y) use compact representation of elliptic curve points, see [Appendix B](#). In COSE, compact representation is achieved by formatting the ECDH ephemeral public keys as COSE_Keys of type EC2 or OKP according to Sections 7.1 and 7.2 of [[RFC9053](#)], but only including the 'x' parameter in G_X and G_Y. For Elliptic Curve Keys of type EC2, compact representation MAY be used also in the COSE_Key. COSE always uses compact output for Elliptic Curve Keys of type EC2. If the COSE

implementation requires a 'y' parameter, the value y = false or a
calculated y-coordinate can be used, see Appendix B.

3.8. External Authorization Data (EAD)

In order to reduce round trips and the number of messages, or to
simplify processing, external security applications may be
integrated into EDHOC by transporting authorization related data in
the messages.

EDHOC allows processing of external authorization data (EAD) to be
defined in a separate specification, and sent in dedicated fields of
the four EDHOC messages (EAD_1, EAD_2, EAD_3, EAD_4). EAD is opaque
data to EDHOC.

Each EAD field, EAD_x for x = 1, 2, 3 or 4, is a CBOR sequence (see
Appendix C.1) consisting of one or more EAD items. An EAD item ead
is a CBOR sequence of an ead_label and an optional ead_value, see
Figure 7 and Appendix C.2 for the CDDL definitions.

```
ead = (
  ead_label : int,
  ? ead_value : bstr,
)
```

Figure 7: EAD item.

A security application may register one or more EAD labels, see
Section 9.5, and specify the associated processing and security
considerations. The IANA registry contains the absolute value of the
ead_label, |ead_label|; the same ead_value applies independently of
sign of ead_label.

An EAD item can be either critical or non-critical, determined by
the sign of the ead_label in the EAD item transported in the EAD
field. A negative value indicates that the EAD item is critical and
a non-negative value indicates that the EAD item is non-critical.

If an endpoint receives a critical EAD item it does not recognize,
or a critical EAD item that contains information that it cannot
process, then the endpoint MUST send an EDHOC error message back as
defined in Section 6, and the protocol MUST be discontinued. The EAD
item specification defines the error processing. A non-critical EAD
item can be ignored.

The security application registering a new EAD item needs to
describe under what conditions the EAD item is critical or non-
critical, and thus whether the ead_label is used with negative or
positive sign. ead_label = 0 is used for padding, see Section 3.8.1.

The security application may define multiple uses of certain EAD items, e.g., the same EAD item may be used in different EDHOC messages. Multiple occurrences of an EAD item in one EAD field may also be specified, but the criticality of the repeated EAD item is expected to be the same.

The EAD fields of EDHOC MUST NOT be used for generic application data. Examples of the use of EAD are provided in Appendix E.

3.8.1. Padding

EDHOC message_1 and the plaintext of message_2, message_3 and message_4 can be padded with the use of the corresponding EAD_x field, for x = 1, 2, 3, 4. Padding in EAD_1 mitigates amplification attacks (see Section 8.7), and padding in EAD_2, EAD_3, and EAD_4 hides the true length of the plaintext (see Section 8.6). Padding is intended to be discarded by the receiving application.

Padding is obtained by using an EAD item with ead_label = 0 and a (pseudo-)randomly generated byte string of appropriate length as ead_value, noting that the ead_label and the CBOR encoding of ead_value also add bytes. Examples:

  *One byte padding (optional ead_value omitted):

     -EAD_x = 0x00

  *Two bytes padding, using the empty byte string (0x40) as
   ead_value:

     -EAD_x = 0x0040

  *Three bytes padding, constructed from the pseudorandomly
   generated ead_value 0xe9 encoded as byte string:

     -EAD_x = 0x0041e9

Multiple occurrences of EAD items with ead_label = 0 are allowed. Certain padding lengths require the use of at least two such EAD items.

Note that padding is non-critical because the intended behaviour when receiving is to ignore it.

3.9. Application Profile

EDHOC requires certain parameters to be agreed upon between Initiator and Responder. Some parameters can be negotiated through the protocol execution (specifically, cipher suite, see Section 3.6) but other parameters are only communicated and may not be negotiated

(e.g., which authentication method is used, see Section 3.2). Yet other parameters need to be known out-of-band. The application decides which endpoint is Initiator and which is Responder.

The purpose of an application profile is to describe the intended use of EDHOC to allow for the relevant processing and verifications to be made, including things like:

1. How the endpoint detects that an EDHOC message is received. This includes how EDHOC messages are transported, for example in the payload of a CoAP message with a certain Uri-Path or Content-Format; see Appendix A.2.

   *The method of transporting EDHOC messages may also describe data carried along with the messages that are needed for the transport to satisfy the requirements of Section 3.4, e.g., connection identifiers used with certain messages, see Appendix A.2.

2. Authentication method (METHOD; see Section 3.2).

3. Profile for authentication credentials (CRED_I, CRED_R; see Section 3.5.2), e.g., profile for certificate or CCS, including supported authentication key algorithms (subject public key algorithm in X.509 or C509 certificate).

4. Type used to identify credentials (ID_CRED_I, ID_CRED_R; see Section 3.5.3).

5. Use and type of external authorization data (EAD_1, EAD_2, EAD_3, EAD_4; see Section 3.8).

6. Identifier used as the identity of the endpoint; see Appendix D.2.

7. If message_4 shall be sent/expected, and if not, how to ensure a protected application message is sent from the Responder to the Initiator; see Section 5.5.

The application profile may also contain information about supported cipher suites. The procedure for selecting and verifying a cipher suite is still performed as described in Section 5.2.1 and Section 6.3, but it may become simplified by this knowledge.

An example of an application profile is shown in Appendix F.

For some parameters, like METHOD, ID_CRED_x, type of EAD, the receiver of an EDHOC message is able to verify compliance with the application profile, and if it needs to fail because of lack of compliance, to infer the reason why the protocol failed.

For other parameters, like the profile of CRED_x in the case that it is not transported, it may not be possible to verify that lack of compliance with the application profile was the reason for failure: Integrity verification in message_2 or message_3 may fail not only because of wrong credential. For example, in case the Initiator uses a public key certificate by reference (i.e., not transported within the protocol) then both endpoints need to use an identical data structure as CRED_I or else the integrity verification will fail.

Note that it is not necessary for the endpoints to specify a single transport for the EDHOC messages. For example, a mix of CoAP and HTTP may be used along the path, and this may still allow correlation between messages.

The application profile may be dependent on the identity of the other endpoint, or other information carried in an EDHOC message, but it then applies only to the later phases of the protocol when such information is known. (The Initiator does not know the identity of the Responder before having verified message_2, and the Responder does not know the identity of the Initiator before having verified message_3.)

Other conditions may be part of the application profile, such as what is the target application or use (if there is more than one application/use) to the extent that EDHOC can distinguish between them. In case multiple application profiles are used, the receiver needs to be able to determine which is applicable for a given session, for example based on URI to which the EDHOC message is sent, or external authorization data type.

4. Key Derivation

4.1. Keys for EDHOC Message Processing

EDHOC uses Extract-and-Expand [RFC5869] with the EDHOC hash algorithm in the selected cipher suite to derive keys used in message processing. This section defines EDHOC_Extract (Section 4.1.1) and EDHOC_Expand (Section 4.1.2), and how to use them to derive PRK_out (Section 4.1.3) which is the shared secret session key resulting from a successful EDHOC exchange.

EDHOC_Extract is used to derive fixed-length uniformly pseudorandom keys (PRK) from ECDH shared secrets. EDHOC_Expand is used to define EDHOC_KDF for generating MACs and for deriving output keying material (OKM) from PRKs.

In EDHOC a specific message is protected with a certain pseudorandom key, but how the key is derived depends on the authentication method (Section 3.2) as detailed in Section 5.

4.1.1. EDHOC_Extract

The pseudorandom keys (PRKs) used for EDHOC message processing are derived using EDHOC_Extract:

PRK = EDHOC_Extract(salt, IKM)

where the input keying material (IKM) and salt are defined for each PRK below.

The definition of EDHOC_Extract depends on the EDHOC hash algorithm of the selected cipher suite:

  *if the EDHOC hash algorithm is SHA-2, then EDHOC_Extract(salt, IKM) = HKDF-Extract(salt, IKM) [RFC5869]

  *if the EDHOC hash algorithm is SHAKE128, then EDHOC_Extract(salt, IKM) = KMAC128(salt, IKM, 256, "")

  *if the EDHOC hash algorithm is SHAKE256, then EDHOC_Extract(salt, IKM) = KMAC256(salt, IKM, 512, "")

The rest of the section defines the pseudorandom keys PRK_2e, PRK_3e2m and PRK_4e3m; their use is shown in Figure 8. The index of a PRK indicates its use or in what message protection operation it is involved. For example, PRK_3e2m is involved in the encryption of message 3 and in calculating the MAC of message 2.

4.1.1.1. PRK_2e

The pseudorandom key PRK_2e is derived with the following input:

  *The salt SHALL be TH_2.

  *The IKM SHALL be the ephemeral-ephemeral ECDH shared secret G_XY (calculated from G_X and Y or G_Y and X) as defined in Section 6.3.1 of [RFC9053]. The use of G_XY gives forward secrecy, in the sense that compromise of the private authentication keys does not compromise past session keys.

Example: Assuming the use of curve25519, the ECDH shared secret G_XY is the output of the X25519 function [RFC7748]:

G_XY = X25519(Y, G_X) = X25519(X, G_Y)

Example: Assuming the use of SHA-256 the extract phase of HKDF produces PRK_2e as follows:

PRK_2e = HMAC-SHA-256(TH_2, G_XY)

4.1.1.2. PRK_3e2m

The pseudorandom key PRK_3e2m is derived as follows:

If the Responder authenticates with a static Diffie-Hellman key,
then PRK_3e2m = EDHOC_Extract(SALT_3e2m, G_RX), where

  *SALT_3e2m is derived from PRK_2e, see [Section 4.1.2](#), and

  *G_RX is the ECDH shared secret calculated from G_R and X, or G_X
   and R (the Responder's private authentication key, see
   [Section 3.5.1](#)),

else PRK_3e2m = PRK_2e.

4.1.1.3. PRK_4e3m

The pseudorandom key PRK_4e3m is derived as follows:

If the Initiator authenticates with a static Diffie-Hellman key,
then PRK_4e3m = EDHOC_Extract(SALT_4e3m, G_IY), where

  *SALT_4e3m is derived from PRK_3e2m, see [Section 4.1.2](#), and

  *G_IY is the ECDH shared secret calculated from G_I and Y, or G_Y
   and I (the Initiator's private authentication key, see
   [Section 3.5.1](#)),

else PRK_4e3m = PRK_3e2m.

4.1.2. EDHOC_Expand and EDHOC_KDF

The output keying material (OKM) - including keys, IVs, and salts -
are derived from the PRKs using the EDHOC_KDF, which is defined
through EDHOC_Expand:

OKM = EDHOC_KDF(PRK, info_label, context, length)
    = EDHOC_Expand(PRK, info, length)

where info is encoded as the CBOR sequence

```
info = (
  info_label : int,
  context : bstr,
  length : uint,
)
```

  where

    *info_label is an int

*context is a bstr

*length is the length of OKM in bytes

When EDHOC_KDF is used to derive OKM for EDHOC message processing,
then context includes one of the transcript hashes TH_2, TH_3, or
TH_4 defined in Sections 5.3.2 and 5.4.2.

The definition of EDHOC_Expand depends on the EDHOC hash algorithm
of the selected cipher suite:

*if the EDHOC hash algorithm is SHA-2, then EDHOC_Expand(PRK,
 info, length) = HKDF-Expand(PRK, info, length) [RFC5869]

*if the EDHOC hash algorithm is SHAKE128, then EDHOC_Expand(PRK,
 info, length) = KMAC128(PRK, info, L, "")

*if the EDHOC hash algorithm is SHAKE256, then EDHOC_Expand(PRK,
 info, length) = KMAC256(PRK, info, L, "")

where L = 8 · length, the output length in bits.

Figure 8 lists derivations made with EDHOC_KDF during message
processing, where

*hash_length - length of output size of the EDHOC hash algorithm
 of the selected cipher suite

*key_length - length of the encryption key of the EDHOC AEAD
 algorithm of the selected cipher suite

*iv_length - length of the initialization vector of the EDHOC AEAD
 algorithm of the selected cipher suite

Further details of the key derivation and how the output keying
material is used are specified in Section 5.

```
KEYSTREAM_2   = EDHOC_KDF( PRK_2e,   0, TH_2,      plaintext_length )
SALT_3e2m     = EDHOC_KDF( PRK_2e,   1, TH_2,      hash_length )
MAC_2         = EDHOC_KDF( PRK_3e2m, 2, context_2, mac_length_2 )
K_3           = EDHOC_KDF( PRK_3e2m, 3, TH_3,      key_length )
IV_3          = EDHOC_KDF( PRK_3e2m, 4, TH_3,      iv_length )
SALT_4e3m     = EDHOC_KDF( PRK_3e2m, 5, TH_3,      hash_length )
MAC_3         = EDHOC_KDF( PRK_4e3m, 6, context_3, mac_length_3 )
PRK_out       = EDHOC_KDF( PRK_4e3m, 7, TH_4,      hash_length )
K_4           = EDHOC_KDF( PRK_4e3m, 8, TH_4,      key_length )
IV_4          = EDHOC_KDF( PRK_4e3m, 9, TH_4,      iv_length )
```

Figure 8: Key derivations using EDHOC_KDF.

4.1.3. PRK_out

The pseudorandom key PRK_out, derived as shown in Figure 8 is the
output session key of a successful EDHOC exchange.

Keys for applications are derived using EDHOC_Exporter from
PRK_exporter (see Section 4.2.1) which in turn is derived from
PRK_out. For the purpose of generating application keys, it is
sufficient to store PRK_out or PRK_exporter. (Note that the word
"store" used here does not imply that the application has access to
the plaintext PRK_out since that may be reserved for code within a
Trusted Execution Environment, see Section 8.8).

4.2. Keys for EDHOC Applications

This section defines EDHOC_Exporter in terms of EDHOC_KDF and
PRK_out. A key update function is defined in Appendix I.

4.2.1. EDHOC_Exporter

Keying material for the application can be derived using the
EDHOC_Exporter interface defined as:

```
EDHOC_Exporter(exporter_label, context, length)
  = EDHOC_KDF(PRK_exporter, exporter_label, context, length)
```

where

  *exporter_label is a registered uint from the EDHOC_Exporter Label
   registry (Section 9.1)

  *context is a bstr defined by the application

  *length is a uint defined by the application

  *PRK_exporter is derived from PRK_out:

```
PRK_exporter  = EDHOC_KDF( PRK_out, 10, h'', hash_length )
```

where hash_length denotes the output size in bytes of the EDHOC hash
algorithm of the selected cipher suite.

The (exporter_label, context) pair used in EDHOC_Exporter must be
unique, i.e., an (exporter_label, context) MUST NOT be used for two
different purposes. However an application can re-derive the same
key several times as long as it is done in a secure way. For
example, in most encryption algorithms the same key can be reused
with different nonces. The context can for example be the empty CBOR
byte string.

Examples of use of the EDHOC_Exporter are given in Appendix A.

5. Message Formatting and Processing

This section specifies formatting of the messages and processing
steps. Error messages are specified in Section 6. Annotated traces
of EDHOC protocol runs are provided in [I-D.ietf-lake-traces].

An EDHOC message is encoded as a sequence of CBOR data items (CBOR
Sequence, [RFC8742]). Additional optimizations are made to reduce
message overhead.

While EDHOC uses the COSE_Key, COSE_Sign1, and COSE_Encrypt0
structures, only a subset of the parameters is included in the EDHOC
messages, see Appendix C.3. In order to recreate the COSE object,
the recipient endpoint may need to add parameters to the COSE
headers not included in the EDHOC message, for example the parameter
'alg' to COSE_Sign1 or COSE_Encrypt0.

5.1. Message Processing Outline

This section outlines the message processing of EDHOC.

For each new/ongoing EDHOC session, the endpoints are assumed to
keep an associated protocol state containing identifiers, keying
material, etc. used for subsequent processing of protocol related
data. The protocol state is assumed to be associated with an
application profile (Section 3.9) which provides the context for how
messages are transported, identified, and processed.

EDHOC messages SHALL be processed according to the current protocol
state. The following steps are expected to be performed at reception
of an EDHOC message:

1. Detect that an EDHOC message has been received, for example by
   means of port number, URI, or media type (Section 3.9).

2. Retrieve the protocol state according to the message
   correlation, see Section 3.4.1. If there is no protocol state,
   in the case of message_1, a new protocol state is created. The
   Responder endpoint needs to make use of available Denial-of-
   Service mitigation (Section 8.7).

3. If the message received is an error message, then process it
   according to Section 6, else process it as the expected next
   message according to the protocol state.

The message processing steps SHALL be processed in order, unless
otherwise stated. If the processing fails for some reason then,
typically, an error message is sent, the protocol is discontinued,

and the protocol state erased. When the composition and sending of one message is completed and before the next message is received, error messages SHALL NOT be sent.

After having successfully processed the last message (message_3 or message_4 depending on application profile) the protocol is completed, after which no error messages are sent and EDHOC session output MAY be maintained even if error messages are received. Further details are provided in the following subsections and in Section 6.

Different instances of the same message MUST NOT be processed in one session. Note that processing will fail if the same message appears a second time for EDHOC processing in the same session because the state of the protocol has moved on and now expects something else. Message deduplication MUST be done by the transport protocol (see Section 3.4) or, if not supported by the transport, as described in Appendix G.

5.2. EDHOC Message 1

5.2.1. Formatting of Message 1

message_1 SHALL be a CBOR Sequence (see Appendix C.1) as defined below

```
message_1 = (
  METHOD : int,
  SUITES_I : suites,
  G_X : bstr,
  C_I : bstr / -24..23,
  ? EAD_1,
)

suites = [ 2* int ] / int
EAD_1 = 1* ead
```

where:

*METHOD - authentication method, see Section 3.2.

*SUITES_I - array of cipher suites which the Initiator supports constructed as specified in Section 5.2.2.

*G_X - the ephemeral public key of the Initiator

*C_I - variable length connection identifier. Note that connection identifiers are byte strings but certain values are represented as integers in the message, see Section 3.3.2.

*EAD_1 - external authorization data, see Section 3.8.

5.2.2. Initiator Processing of Message 1

The processing steps are detailed below and in Section 6.3.

The Initiator SHALL compose message_1 as follows:

*Construct SUITES_I as an array of cipher suites supported by I in
 order of preference with the first cipher suite in the array
 being the most preferred by I, and the last being the one
 selected by I for this session. If the cipher suite most
 preferred by I is selected then SUITES_I contains only that
 cipher suite and is encoded as an int. All cipher suites, if any,
 preferred by I over the selected one MUST be included. (See also
 Section 6.3.)

  -The selected suite is based on what the Initiator can assume
   to be supported by the Responder; if the Initiator previously
   received from the Responder an error message with error code 2
   containing SUITES_R (see Section 6.3) indicating cipher suites
   supported by the Responder, then the Initiator SHOULD select
   its most preferred supported cipher suite among those (bearing
   in mind that error messages are not authenticated and may be
   forged).

  -The Initiator MUST NOT change its order of preference for
   cipher suites, and MUST NOT omit a cipher suite preferred to
   the selected one because of previous error messages received
   from the Responder.

*Generate an ephemeral ECDH key pair using the curve in the
 selected cipher suite and format it as a COSE_Key. Let G_X be the
 'x' parameter of the COSE_Key.

*Choose a connection identifier C_I and store it during the EDHOC
 session.

*Encode message_1 as a sequence of CBOR encoded data items as
 specified in Section 5.2.1

5.2.3. Responder Processing of Message 1

The Responder SHALL process message_1 in the following order:

*Decode message_1 (see Appendix C.1).

*Process message_1, in particular verify that the selected cipher
 suite is supported and that no prior cipher suite in SUITES_I is
 supported.

*If all processing completed successfully, and if EAD_1 is
 present, then make it available to the application for EAD
 processing.

If any processing step fails, then the Responder MUST send an EDHOC
error message back as defined in [Section 6](#), and the protocol MUST be
discontinued.

5.3. EDHOC Message 2

5.3.1. Formatting of Message 2

message_2 SHALL be a CBOR Sequence (see [Appendix C.1](#)) as defined
below

```
message_2 = (
  G_Y_CIPHERTEXT_2 : bstr,
  C_R : bstr / -24..23,
)
```

where:

*G_Y_CIPHERTEXT_2 - the concatenation of G_Y (i.e., the ephemeral
 public key of the Responder) and CIPHERTEXT_2.

*C_R - variable length connection identifier. Note that connection
 identifiers are byte strings but certain values are represented
 as integers in the message, see [Section 3.3.2](#).

5.3.2. Responder Processing of Message 2

The Responder SHALL compose message_2 as follows:

*Generate an ephemeral ECDH key pair using the curve in the
 selected cipher suite and format it as a COSE_Key. Let G_Y be the
 'x' parameter of the COSE_Key.

*Choose a connection identifier C_R and store it for the length of
 the protocol.

*Compute the transcript hash TH_2 = H(G_Y, C_R, H(message_1))
 where H() is the EDHOC hash algorithm of the selected cipher
 suite. The input to the hash function is a CBOR Sequence. Note
 that H(message_1) can be computed and cached already in the
 processing of message_1.

*Compute MAC_2 as in [Section 4.1.2](#) with context_2 = << ID_CRED_R,
 TH_2, CRED_R, ? EAD_2 >> (see [Appendix C.1](#) for notation)

  -If the Responder authenticates with a static Diffie-Hellman
   key (method equals 1 or 3), then mac_length_2 is the EDHOC MAC
   length of the selected cipher suite. If the Responder
   authenticates with a signature key (method equals 0 or 2),
   then mac_length_2 is equal to the output size of the EDHOC
   hash algorithm of the selected cipher suite.

  -ID_CRED_R - identifier to facilitate the retrieval of CRED_R,
   see [Section 3.5.3](#)

  -CRED_R - CBOR item containing the authentication credential of
   the Responder, see [Section 3.5.2](#)

  -EAD_2 - external authorization data, see [Section 3.8](#)

*If the Responder authenticates with a static Diffie-Hellman key
 (method equals 1 or 3), then Signature_or_MAC_2 is MAC_2. If the
 Responder authenticates with a signature key (method equals 0 or
 2), then Signature_or_MAC_2 is the 'signature' field of a
 COSE_Sign1 object, computed as specified in Section 4.4 of
 [[RFC9053](#)] using the signature algorithm of the selected cipher
 suite, the private authentication key of the Responder, and the
 following parameters as input (see [Appendix C.3](#) for an overview
 of COSE and [Appendix C.1](#) for notation):

  -protected = << ID_CRED_R >>

  -external_aad = << TH_2, CRED_R, ? EAD_2 >>

  -payload = MAC_2

*CIPHERTEXT_2 is calculated by using the EDHOC_Expand function as
 a binary additive stream cipher over the following plaintext:

  -PLAINTEXT_2 = (ID_CRED_R / bstr / -24..23,
   Signature_or_MAC_2, ? EAD_2)

     oIf ID_CRED_R contains a single 'kid' parameter, i.e.,
       ID_CRED_R = { 4 : kid_R }, then the compact encoding is
       applied, see [Section 3.5.3.1](#).

  -Compute KEYSTREAM_2 as in [Section 4.1.2](#), where
   plaintext_length is the length of PLAINTEXT_2. For the case of
   plaintext_length exceeding the EDHOC_KDF output size, see
   [Appendix H](#).

  -CIPHERTEXT_2 = PLAINTEXT_2 XOR KEYSTREAM_2

*Encode message_2 as a sequence of CBOR encoded data items as
 specified in Section 5.3.1.

5.3.3. Initiator Processing of Message 2

The Initiator SHALL process message_2 in the following order:

*Decode message_2 (see Appendix C.1).

*Retrieve the protocol state using available message correlation
 (e.g., the CoAP Token, the 5-tuple, or the prepended C_I, see
 Section 3.4.1).

*Decrypt CIPHERTEXT_2, see Section 5.3.2.

*If all processing completed successfully, then make ID_CRED_R and
 (if present) EAD_2 available to the application for
 authentication- and EAD processing.

*Obtain the authentication credential (CRED_R) and the
 authentication key of R from the application (or by other means).

*Verify Signature_or_MAC_2 using the algorithm in the selected
 cipher suite. The verification process depends on the method, see
 Section 5.3.2. Make the result of the verification available to
 the application.

If any processing step fails, then the Initiator MUST send an EDHOC
error message back as defined in Section 6, and the protocol MUST be
discontinued.

5.4. EDHOC Message 3

5.4.1. Formatting of Message 3

message_3 SHALL be a CBOR Sequence (see Appendix C.1) as defined
below

```
message_3 = (
  CIPHERTEXT_3 : bstr,
)
```

5.4.2. Initiator Processing of Message 3

The Initiator SHALL compose message_3 as follows:

*Compute the transcript hash TH_3 = H(TH_2, PLAINTEXT_2, CRED_R)
 where H() is the EDHOC hash algorithm of the selected cipher
 suite. The input to the hash function is a CBOR Sequence. Note

that TH_3 can be computed and cached already in the processing of
message_2.

*Compute MAC_3 as in Section 4.1.2, with context_3 = << ID_CRED_I,
 TH_3, CRED_I, ? EAD_3 >>

  -If the Initiator authenticates with a static Diffie-Hellman
   key (method equals 2 or 3), then mac_length_3 is the EDHOC MAC
   length of the selected cipher suite. If the Initiator
   authenticates with a signature key (method equals 0 or 1),
   then mac_length_3 is equal to the output size of the EDHOC
   hash algorithm of the selected cipher suite.

  -ID_CRED_I - identifier to facilitate the retrieval of CRED_I,
   see Section 3.5.3

  -CRED_I - CBOR item containing the authentication credential of
   the Initiator, see Section 3.5.2

  -EAD_3 - external authorization data, see Section 3.8

*If the Initiator authenticates with a static Diffie-Hellman key
 (method equals 2 or 3), then Signature_or_MAC_3 is MAC_3. If the
 Initiator authenticates with a signature key (method equals 0 or
 1), then Signature_or_MAC_3 is the 'signature' field of a
 COSE_Sign1 object, computed as specified in Section 4.4 of
 [RFC9052] using the signature algorithm of the selected cipher
 suite, the private authentication key of the Initiator, and the
 following parameters as input (see Appendix C.3):

  -protected = << ID_CRED_I >>

  -external_aad = << TH_3, CRED_I, ? EAD_3 >>

  -payload = MAC_3

*Compute a COSE_Encrypt0 object as defined in Sections 5.2 and 5.3
 of [RFC9052], with the EDHOC AEAD algorithm of the selected
 cipher suite, using the encryption key K_3, the initialization
 vector IV_3 (if used by the AEAD algorithm), the plaintext
 PLAINTEXT_3, and the following parameters as input (see
 Appendix C.3):

  -protected = h''

  -external_aad = TH_3

  -K_3 and IV_3 are defined in Section 4.1.2

```
      -PLAINTEXT_3 = ( ID_CRED_I / bstr / -24..23,
        Signature_or_MAC_3, ? EAD_3 )

          oIf ID_CRED_I contains a single 'kid' parameter, i.e.,
           ID_CRED_I = { 4 : kid_I }, then the compact encoding is
           applied, see Section 3.5.3.1.
```

CIPHERTEXT_3 is the 'ciphertext' of COSE_Encrypt0.

*Compute the transcript hash TH_4 = H(TH_3, PLAINTEXT_3, CRED_I) where H() is the EDHOC hash algorithm of the selected cipher suite. The input to the hash function is a CBOR Sequence. Note that TH_4 can be computed and cached already in the processing of message_3.

*Calculate PRK_out as defined in Figure 8. The Initiator can now derive application keys using the EDHOC_Exporter interface, see Section 4.2.1.

*Encode message_3 as a CBOR data item as specified in Section 5.4.1.

*Make the connection identifiers (C_I, C_R) and the application algorithms in the selected cipher suite available to the application.

The Initiator SHOULD NOT persistently store PRK_out or application keys until the Initiator has verified message_4 or a message protected with a derived application key, such as an OSCORE message, from the Responder. This is similar to waiting for an acknowledgement (ACK) in a transport protocol.

5.4.3. Responder Processing of Message 3

The Responder SHALL process message_3 in the following order:

*Decode message_3 (see Appendix C.1).

*Retrieve the protocol state using available message correlation (e.g., the CoAP Token, the 5-tuple, or the prepended C_R, see Section 3.4.1).

*Decrypt and verify the COSE_Encrypt0 as defined in Sections 5.2 and 5.3 of [RFC9052], with the EDHOC AEAD algorithm in the selected cipher suite, and the parameters defined in Section 5.4.2.

*If all processing completed successfully, then make ID_CRED_I and (if present) EAD_3 available to the application for authentication- and EAD processing.

*Obtain the authentication credential (CRED_I) and the
 authentication key of I from the application (or by other means).

*Verify Signature_or_MAC_3 using the algorithm in the selected
 cipher suite. The verification process depends on the method, see
 [Section 5.4.2](). Make the result of the verification available to
 the application.

*Make the connection identifiers (C_I, C_R) and the application
 algorithms in the selected cipher suite available to the
 application.

After verifying message_3, the Responder can compute PRK_out, see
[Section 4.1.3](), derive application keys using the EDHOC_Exporter
interface, see [Section 4.2.1](), persistently store the keying
material, and send protected application data.

If any processing step fails, then the Responder MUST send an EDHOC
error message back as defined in [Section 6](), and the protocol MUST be
discontinued.

5.5. EDHOC Message 4

This section specifies message_4 which is OPTIONAL to support. Key
confirmation is normally provided by sending an application message
from the Responder to the Initiator protected with a key derived
with the EDHOC_Exporter, e.g., using OSCORE (see [Appendix A]()). In
deployments where no protected application message is sent from the
Responder to the Initiator, message_4 MUST be supported and MUST be
used. Two examples of such deployments are:

1. When EDHOC is only used for authentication and no application
   data is sent.

2. When application data is only sent from the Initiator to the
   Responder.

Further considerations about when to use message_4 are provided in
[Section 3.9]() and [Section 8.1]().

5.5.1. Formatting of Message 4

message_4 SHALL be a CBOR Sequence (see [Appendix C.1]()) as defined
below

```
message_4 = (
  CIPHERTEXT_4 : bstr,
)
```

5.5.2. Responder Processing of Message 4

The Responder SHALL compose message_4 as follows:

* Compute a COSE_Encrypt0 as defined in Sections 5.2 and 5.3 of
  [RFC9052], with the EDHOC AEAD algorithm of the selected cipher
  suite, using the encryption key K_4, the initialization vector
  IV_4 (if used by the AEAD algorithm), the plaintext PLAINTEXT_4,
  and the following parameters as input (see Appendix C.3):

    - protected = h''

    - external_aad = TH_4

    - K_4 and IV_4 are defined in Section 4.1.2

    - PLAINTEXT_4 = (? EAD_4)

        o EAD_4 - external authorization data, see Section 3.8.

  CIPHERTEXT_4 is the 'ciphertext' of COSE_Encrypt0.

* Encode message_4 as a CBOR data item as specified in
  Section 5.5.1.

5.5.3. Initiator Processing of Message 4

The Initiator SHALL process message_4 as follows:

* Decode message_4 (see Appendix C.1).

* Retrieve the protocol state using available message correlation
  (e.g., the CoAP Token, the 5-tuple, or the prepended C_I, see
  Section 3.4.1).

* Decrypt and verify the COSE_Encrypt0 as defined in Sections 5.2
  and 5.3 of [RFC9052], with the EDHOC AEAD algorithm in the
  selected cipher suite, and the parameters defined in
  Section 5.5.2.

* Make (if present) EAD_4 available to the application for EAD
  processing.

If any processing step fails, then the Initiator MUST send an EDHOC
error message back as defined in Section 6, and the protocol MUST be
discontinued.

After verifying message_4, the Initiator is assured that the
Responder has calculated the key PRK_out (key confirmation) and that
no other party can derive the key.

6. Error Handling

This section defines the format for error messages, and the processing associated with the currently defined error codes. Additional error codes may be registered, see Section 9.4.

There are many kinds of errors that can occur during EDHOC processing. As in CoAP, an error can be triggered by errors in the received message or internal errors in the receiving endpoint. Except for processing and formatting errors, it is up to the implementation when to send an error message. Sending error messages is essential for debugging but MAY be skipped if, for example, a session cannot be found or due to denial-of-service reasons, see Section 8.7. Error messages in EDHOC are always fatal. After sending an error message, the sender MUST discontinue the protocol. The receiver SHOULD treat an error message as an indication that the other party likely has discontinued the protocol. But as the error message is not authenticated, a received error message might also have been sent by an attacker and the receiver MAY therefore try to continue the protocol.

An EDHOC error message can be sent by either endpoint as a reply to any non-error EDHOC message. How errors at the EDHOC layer are transported depends on lower layers, which need to enable error messages to be sent and processed as intended.

error SHALL be a CBOR Sequence (see Appendix C.1) as defined below

```
error = (
  ERR_CODE : int,
  ERR_INFO : any,
)
```

Figure 9: EDHOC error message.

where:

   *ERR_CODE - error code encoded as an integer. The value 0 is used for success, all other values (negative or positive) indicate errors.

   *ERR_INFO - error information. Content and encoding depend on error code.

The remainder of this section specifies the currently defined error codes, see Figure 10. Additional error codes and corresponding error information may be specified.

ERR CODE	ERR INFO Type	Description
0	any	Success
1	tstr	Unspecified error
2	suites	Wrong selected cipher suite

Figure 10: Error codes and error information included in the EDHOC
error message.

6.1. Success

Error code 0 MAY be used internally in an application to indicate
success, i.e., as a standard value in case of no error, e.g., in
status reporting or log files. ERR_INFO can contain any type of CBOR
item, the content is out of scope for this specification. Error code
0 MUST NOT be used as part of the EDHOC message exchange flow. If an
endpoint receives an error message with error code 0, then it MUST
discontinue the protocol and MUST NOT send an error message.

6.2. Unspecified Error

Error code 1 is used for errors that do not have a specific error
code defined. ERR_INFO MUST be a text string containing a human-
readable diagnostic message written in English, for example "Method
not supported". The diagnostic text message is mainly intended for
software engineers that during debugging need to interpret it in the
context of the EDHOC specification. The diagnostic message SHOULD be
provided to the calling application where it SHOULD be logged.

6.3. Wrong Selected Cipher Suite

Error code 2 MUST only be used when replying to message_1 in case
the cipher suite selected by the Initiator is not supported by the
Responder, or if the Responder supports a cipher suite more
preferred by the Initiator than the selected cipher suite, see
Section 5.2.3. In this case, ERR_INFO = SUITES_R and is of type
suites, see Section 5.2.1. If the Responder does not support the
selected cipher suite, then SUITES_R MUST include one or more
supported cipher suites. If the Responder supports a cipher suite in
SUITES_I other than the selected cipher suite (independently of if
the selected cipher suite is supported or not) then SUITES_R MUST
include the supported cipher suite in SUITES_I which is most
preferred by the Initiator. SUITES_R MAY include a single cipher
suite, in which case it is encoded as an int. If the Responder does

not support any cipher suite in SUITES_I, then it SHOULD include all
its supported cipher suites in SUITES_R.

In contrast to SUITES_I, the order of the cipher suites in SUITES_R
has no significance.

6.3.1. Cipher Suite Negotiation

After receiving SUITES_R, the Initiator can determine which cipher
suite to select (if any) for the next EDHOC run with the Responder.

If the Initiator intends to contact the Responder in the future, the
Initiator SHOULD remember which selected cipher suite to use until
the next message_1 has been sent, otherwise the Initiator and
Responder will likely run into an infinite loop where the Initiator
selects its most preferred cipher suite and the Responder sends an
error with supported cipher suites. After a successful run of EDHOC,
the Initiator MAY remember the selected cipher suite to use in
future EDHOC sessions. Note that if the Initiator or Responder is
updated with new cipher suite policies, any cached information may
be outdated.

Note that the Initiator's list of supported cipher suites and order
of preference is fixed (see Section 5.2.1 and Section 5.2.2).
Furthermore, the Responder SHALL only accept message_1 if the
selected cipher suite is the first cipher suite in SUITES_I that the
Responder supports (see Section 5.2.3). Following this procedure
ensures that the selected cipher suite is the most preferred (by the
Initiator) cipher suite supported by both parties.

If the selected cipher suite is not the first cipher suite which the
Responder supports in SUITES_I received in message_1, then the
Responder MUST discontinue the protocol, see Section 5.2.3. If
SUITES_I in message_1 is manipulated, then the integrity
verification of message_2 containing the transcript hash TH_2 will
fail and the Initiator will discontinue the protocol.

6.3.2. Examples

Assume that the Initiator supports the five cipher suites 5, 6, 7,
8, and 9 in decreasing order of preference. Figures 11 and 12 show
two examples of how the Initiator can format SUITES_I and how
SUITES_R is used by Responders to give the Initiator information
about the cipher suites that the Responder supports.

In Example 1 (Figure 11), the Responder supports cipher suite 6 but
not the initially selected cipher suite 5. The Responder rejects the
first message_1 with an error indicating support for suite 6 in
SUITES_R. The Initiator also supports suite 6, and therefore selects
suite 6 in the second message_1. The Initiator prepends in SUITES_I

the selected suite 6 with the more preferred suites, in this case
suite 5, to mitigate a potential attack on the cipher suite
negotiation.

```
Initiator                                                     Responder
            METHOD, SUITES I = 5. G X. C I. EAD 1
 |------------------------------------------------------------------>|
                            message 1

                ERR CODE = 2. SUITES R = 6
 |<------------------------------------------------------------------|
                              error

          METHOD, SUITES I = ⌈5. 6⌉. G X. C I. EAD 1
 |------------------------------------------------------------------>|
                            message 1
```
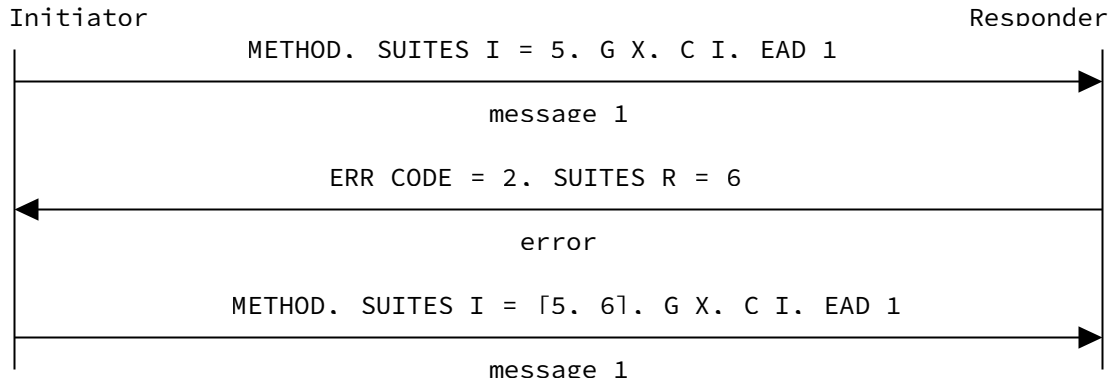
Figure 11: Cipher Suite Negotiation Example 1.

In Example 2 (Figure 12), the Responder supports cipher suites 8 and
9 but not the more preferred (by the Initiator) cipher suites 5, 6
or 7. To illustrate the negotiation mechanics we let the Initiator
first make a guess that the Responder supports suite 6 but not suite
5. Since the Responder supports neither 5 nor 6, it rejects the
first message_1 with an error indicating support for suites 8 and 9
in SUITES_R (in any order). The Initiator also supports suites 8 and
9, and prefers suite 8, so therefore selects suite 8 in the second
message_1. The Initiator prepends in SUITES_I the selected suite 8
with the more preferred suites in order of preference, in this case
suites 5, 6 and 7, to mitigate a potential attack on the cipher
suite negotiation.

Note 1. If the Responder had supported suite 5, it would have
included it in SUITES_R of the response, and it would in that case
have become the selected and only suite in the second message_1.

Note 2. The content of the fields of message_1 may be different when
sent the second time, in particular the ephemeral key MUST be
different.

```
Initiator                                             Responder
            METHOD. SUITES I = ⌈5. 6⌉. G X. C I. EAD 1
  |————————————————————————————————————————————————————————————>|
  |                        message 1                             |
  |                                                              |
  |              ERR CODE = 2. SUITES R = ⌈9. 8⌉                 |
  |<————————————————————————————————————————————————————————————|
  |                         error                                |
  |                                                              |
  |         METHOD. SUITES I = ⌈5. 6. 7. 8⌉. G X. C I. EAD 1     |
  |————————————————————————————————————————————————————————————>|
  |                        message 1                             |
```
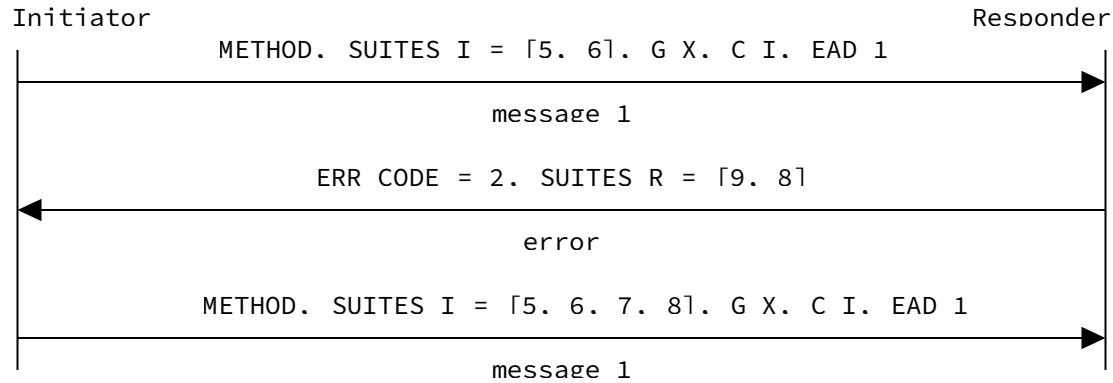
Figure 12: Cipher Suite Negotiation Example 2.

7. Compliance Requirements

In the absence of an application profile specifying otherwise:

An implementation MAY support only Initiator or only Responder.

An implementation MAY support only a single method. None of the methods are mandatory-to-implement.

Implementations MUST support 'kid' parameters. None of the other COSE header parameters are mandatory-to-implement.

An implementation MAY support only a single credential type (CCS, CWT, X.509, C509). None of the credential types are mandatory-to-implement.

Implementations MUST support the EDHOC_Exporter.

Implementations MAY support message_4. Error codes (ERR_CODE) 1 and 2 MUST be supported.

Implementations MUST support EAD.

Implementations MUST support cipher suite 2 and 3. Cipher suites 2 (AES-CCM-16-64-128, SHA-256, 8, P-256, ES256, AES-CCM-16-64-128, SHA-256) and 3 (AES-CCM-16-128-128, SHA-256, 16, P-256, ES256, AES-CCM-16-64-128, SHA-256) only differ in the size of the MAC length, so supporting one or both of these is not significantly different. Implementations only need to implement the algorithms needed for their supported methods.

8. Security Considerations

8.1. Security Properties

EDHOC has similar security properties as can be expected from the
theoretical SIGMA-I protocol [SIGMA] and the Noise XX pattern
[Noise], which are similar to methods 0 and 3, respectively. Proven
security properties are detailed in the security analysis
publications referenced at the end of this section.

Using the terminology from [SIGMA], EDHOC provides forward secrecy,
mutual authentication with aliveness, consistency, and peer
awareness. As described in [SIGMA], message_3 provides peer
awareness to the Responder while message_4 provides peer awareness
to the Initiator. By including the authentication credentials in the
transcript hash, EDHOC protects against Duplicate Signature Key
Selection (DSKS)-like identity mis-binding attack that the MAC-then-
Sign variant of SIGMA-I is otherwise vulnerable to.

As described in [SIGMA], different levels of identity protection are
provided to the Initiator and the Responder. EDHOC provides identity
protection of the Initiator against active attacks and identity
protection of the Responder against passive attacks. An active
attacker can get the credential identifier of the Responder by
eavesdropping on the destination address used for transporting
message_1 and then sending its own message_1 to the same address.
The roles should be assigned to protect the most sensitive identity/
identifier, typically that which is not possible to infer from
routing information in the lower layers.

EDHOC messages might change in transit due to a noisy channel or
through modification by an attacker. Changes in message_1 and
message_2 (except Signature_or_MAC_2 when the signature scheme is
not strongly unforgeable) are detected when verifying
Signature_or_MAC_2. Changes to not strongly unforgeable
Signature_or_MAC_2, and message_3 are detected when verifying
CIPHERTEXT_3. Changes to message_4 are detected when verifying
CIPHERTEXT_4.

Compared to [SIGMA], EDHOC adds an explicit method type and expands
the message authentication coverage to additional elements such as
algorithms, external authorization data, and previous plaintext
messages. This protects against an attacker replaying messages or
injecting messages from another session.

EDHOC also adds selection of connection identifiers and downgrade
protected negotiation of cryptographic parameters, i.e., an attacker
cannot affect the negotiated parameters. A single session of EDHOC
does not include negotiation of cipher suites, but it enables the

Responder to verify that the selected cipher suite is the most preferred cipher suite by the Initiator which is supported by both the Initiator and the Responder, and to discontinue the session if not.

As required by [RFC7258], IETF protocols need to mitigate pervasive monitoring when possible. EDHOC therefore only supports methods with ephemeral Diffie-Hellman and provides a key update function (see Appendix I) for lightweight application protocol rekeying. Either of these provides forward secrecy, in the sense that compromise of the private authentication keys does not compromise past session keys (PRK_out), and compromise of a session key does not compromise past session keys. Frequently re-running EDHOC with ephemeral Diffie-Hellman forces attackers to perform dynamic key exfiltration where the attacker must have continuous interactions with the collaborator, which is a significant complication.

To limit the effect of breaches, it is important to limit the use of symmetric group keys for bootstrapping. EDHOC therefore strives to make the additional cost of using raw public keys and self-signed certificates as small as possible. Raw public keys and self-signed certificates are not a replacement for a public key infrastructure but SHOULD be used instead of symmetric group keys for bootstrapping.

Compromise of the long-term keys (private signature or static DH keys) does not compromise the security of completed EDHOC exchanges. Compromising the private authentication keys of one party lets an active attacker impersonate that compromised party in EDHOC exchanges with other parties but does not let the attacker impersonate other parties in EDHOC exchanges with the compromised party. Compromise of the long-term keys does not enable a passive attacker to compromise future session keys (PRK_out). Compromise of the HDKF input parameters (ECDH shared secret) leads to compromise of all session keys derived from that compromised shared secret. Compromise of one session key does not compromise other session keys. Compromise of PRK_out leads to compromise of all keying material derived with the EDHOC_Exporter.

Based on the cryptographic algorithms requirements Section 8.3, EDHOC provides a minimum of 64-bit security against online brute force attacks and a minimum of 128-bit security against offline brute force attacks. To break 64-bit security against online brute force an attacker would on average have to send 4.3 billion messages per second for 68 years, which is infeasible in constrained IoT radio technologies. A forgery against a 64-bit MAC in EDHOC breaks the security of all future application data, while a forgery against a 64-bit MAC in the subsequent application protocol (e.g., OSCORE

[RFC8613]) typically only breaks the security of the data in the forged packet.

As the EDHOC protocol is discontinued when verification fails, the security against online attacks is given by the sum of the strength of the verified signatures and MACs (including MAC in AEAD). As an example, if EDHOC is used with method 3, cipher suite 2, and message_4, the Responder is authenticated with 128-bit security against online attacks (the sum of the 64-bit MACs in message_2 and message_4). The same principle applies for MACs in an application protocol keyed by EDHOC as long as EDHOC is rerun when verification of the first MACs in the application protocol fails. As an example, if EDHOC with method 3 and cipher suite 2 is used as in Figure 2 of [I-D.ietf-core-oscore-edhoc], 128-bit mutual authentication against online attacks can be achieved after completion of the first OSCORE request and response.

After sending message_3, the Initiator is assured that no other party than the Responder can compute the key PRK_out. While the Initiator can securely send protected application data, the Initiator SHOULD NOT persistently store the keying material PRK_out until the Initiator has verified message_4 or a message protected with a derived application key, such as an OSCORE message, from the Responder. After verifying message_3, the Responder is assured that an honest Initiator has computed the key PRK_out. The Responder can securely derive and store the keying material PRK_out, and send protected application data.

External authorization data sent in message_1 (EAD_1) or message_2 (EAD_2) should be considered unprotected by EDHOC, see Section 8.5. EAD_2 is encrypted but the Responder has not yet authenticated the Initiator and the encryption does not provide confidentiality against active attacks.

External authorization data sent in message_3 (EAD_3) or message_4 (EAD_4) is protected between Initiator and Responder by the protocol, but note that EAD fields may be used by the application before the message verification is completed, see Section 3.8. Designing a secure mechanism that uses EAD is not necessarily straightforward. This document only provides the EAD transport mechanism, but the problem of agreeing on the surrounding context and the meaning of the information passed to and from the application remains. Any new uses of EAD should be subject to careful review.

Key compromise impersonation (KCI): In EDHOC authenticated with signature keys, EDHOC provides KCI protection against an attacker having access to the long-term key or the ephemeral secret key. With static Diffie-Hellman key authentication, KCI protection would be

provided against an attacker having access to the long-term Diffie-Hellman key, but not to an attacker having access to the ephemeral secret key. Note that the term KCI has typically been used for compromise of long-term keys, and that an attacker with access to the ephemeral secret key can only attack that specific session.

Repudiation: If an endpoint authenticates with a signature, the other endpoint can prove that the endpoint performed a run of the protocol by presenting the data being signed as well as the signature itself. With static Diffie-Hellman key authentication, the authenticating endpoint can deny having participated in the protocol.

Earlier versions of EDHOC have been formally analyzed [Bruni18] [Norrman20] [CottierPointcheval22] [Jacomme23] [GuentherIlunga22] and the specification has been updated based on the analysis.

8.2. Cryptographic Considerations

The SIGMA protocol requires that the encryption of message_3 provides confidentiality against active attackers and EDHOC message_4 relies on the use of authenticated encryption. Hence the message authenticating functionality of the authenticated encryption in EDHOC is critical: authenticated encryption MUST NOT be replaced by plain encryption only, even if authentication is provided at another level or through a different mechanism.

To reduce message overhead EDHOC does not use explicit nonces and instead relies on the ephemeral public keys to provide randomness to each session. A good amount of randomness is important for the key generation, to provide liveness, and to protect against interleaving attacks. For this reason, the ephemeral keys MUST NOT be used in more than one EDHOC message, and both parties SHALL generate fresh random ephemeral key pairs. Note that an ephemeral key may be used to calculate several ECDH shared secrets. When static Diffie-Hellman authentication is used the same ephemeral key is used in both ephemeral-ephemeral and ephemeral-static ECDH.

As discussed in [SIGMA], the encryption of message_2 does only need to protect against passive attacker as active attackers can always get the Responder's identity by sending their own message_1. EDHOC uses the EDHOC_Expand function (typically HKDF-Expand) as a binary additive stream cipher which is proven secure as long as the expand function is a PRF. HKDF-Expand is not often used as a stream cipher as it is slow on long messages, and most applications require both IND-CCA confidentiality as well as integrity protection. For the encryption of message_2, any speed difference is negligible, IND-CCA does not increase security, and integrity is provided by the inner MAC (and signature depending on method).

Requirements for how to securely generate, validate, and process the ephemeral public keys depend on the elliptic curve. For X25519 and X448, the requirements are defined in [RFC7748]. For secp256r1, secp384r1, and secp521r1, the requirements are defined in Section 5 of [SP-800-56A]. For secp256r1, secp384r1, and secp521r1, at least partial public-key validation MUST be done.

As noted in Section 12 of [RFC9052] the use of a single key for multiple algorithms is strongly discouraged unless proven secure by a dedicated cryptographic analysis. In particular this recommendation applies to using the same private key for static Diffie-Hellman authentication and digital signature authentication. A preliminary conjecture is that a minor change to EDHOC may be sufficient to fit the analysis of secure shared signature and ECDH key usage in [Degabriele11] and [Thormarker21].

The property that a completed EDHOC exchange implies that another identity has been active is upheld as long as the Initiator does not have its own identity in the set of Responder identities it is allowed to communicate with. In Trust on first use (TOFU) use cases, see Appendix D.5, the Initiator should verify that the Responder's identity is not equal to its own. Any future EDHOC methods using e.g., pre-shared keys might need to mitigate this in other ways. However, an active attacker can gain information about the set of identities an Initiator is willing to communicate with. If the Initiator is willing to communicate with all identities except its own an attacker can determine that a guessed Initiator identity is correct. To not leak any long-term identifiers, using a freshly generated authentication key as identity in each initial TOFU exchange is RECOMMENDED.

8.3. Cipher Suites and Cryptographic Algorithms

When using private cipher suite or registering new cipher suites, the choice of key length used in the different algorithms needs to be harmonized, so that a sufficient security level is maintained for authentication credentials, the EDHOC protocol, and the protection of application data. The Initiator and the Responder should enforce a minimum security level.

The output size of the EDHOC hash algorithm MUST be at least 256-bits, i.e., the hash algorithms SHA-1 and SHA-256/64 (SHA-256 truncated to 64-bits) SHALL NOT be supported for use in EDHOC except for certificate identification with x5t and c5t. For security considerations of SHA-1, see [RFC6194]. As EDHOC integrity protects the whole authentication credentials, the choice of hash algorithm in x5t and c5t does not affect security, and using the same hash algorithm as in the cipher suite, but with as much truncation as possible, is RECOMMENDED. That is, when the EDHOC hash algorithm is

SHA-256, using SHA-256/64 in x5t and c5t is RECOMMENDED. The EDHOC
MAC length MUST be at least 8 bytes and the tag length of the EDHOC
AEAD algorithm MUST be at least 64-bits. Note that secp256k1 is only
defined for use with ECDSA and not for ECDH. Note that some COSE
algorithms are marked as not recommended in the COSE IANA registry.

8.4. Post-Quantum Considerations

As of the publication of this specification, it is unclear when or
even if a quantum computer of sufficient size and power to exploit
public key cryptography will exist. Deployments that need to
consider risks decades into the future should transition to Post-
Quantum Cryptography (PQC) in the not-too-distant future. Many other
systems should take a slower wait-and-see approach where PQC is
phased in when the quantum threat is more imminent. Current PQC
algorithms have limitations compared to Elliptic Curve Cryptography
(ECC) and the data sizes would be problematic in many constrained
IoT systems.

Symmetric algorithms used in EDHOC such as SHA-256 and AES-
CCM-16-64-128 are practically secure against even large quantum
computers. EDHOC supports all signature algorithms defined by COSE,
including PQC signature algorithms such as HSS-LMS. EDHOC is
currently only specified for use with key exchange algorithms of
type ECDH curves, but any Key Encapsulation Method (KEM), including
PQC KEMs, can be used in method 0. While the key exchange in method
0 is specified with terms of the Diffie-Hellman protocol, the key
exchange adheres to a KEM interface: G_X is then the public key of
the Initiator, G_Y is the encapsulation, and G_XY is the shared
secret. Use of PQC KEMs to replace static DH authentication would
likely require a specification updating EDHOC with new methods.

8.5. Unprotected Data and Privacy

The Initiator and the Responder must make sure that unprotected data
and metadata do not reveal any sensitive information. This also
applies for encrypted data sent to an unauthenticated party. In
particular, it applies to EAD_1, ID_CRED_R, EAD_2, and error
messages. Using the same EAD_1 in several EDHOC sessions allows
passive eavesdroppers to correlate the different sessions. Note that
even if ead_value is encrypted outside of EDHOC, the ead_label in
EAD_1 is revealed to passive attackers and the ead_label in EAD_2 is
revealed to active attackers. Another consideration is that the list
of supported cipher suites may potentially be used to identify the
application. The Initiator and the Responder must also make sure
that unauthenticated data does not trigger any harmful actions. In
particular, this applies to EAD_1 and error messages.

An attacker observing network traffic may use connection identifiers sent in clear in EDHOC or the subsequent application protocol to correlate packets sent on different paths or at different times. The attacker may use this information for traffic flow analysis or to track an endpoint. Application protocols using connection identifiers from EDHOC SHOULD provide mechanisms to update the connection identifiers and MAY provide mechanisms to issue several simultaneously active connection identifiers. See [RFC9000] for a non-constrained example of such mechanisms. Connection identifiers can e.g., be chosen randomly among the set of unused 1-byte connection identifiers. Connection identity privacy mechanisms are only useful when there are not fixed identifiers such as IP address or MAC address in the lower layers.

8.6. Updated Internet Threat Model Considerations

Since the publication of [RFC3552] there has been an increased awareness of the need to protect against endpoints that are compromised, malicious, or whose interests simply do not align with the interests of users [I-D.arkko-arch-internet-threat-model-guidance]. [RFC7624] describes an updated threat model for Internet confidentiality, see Section 8.1. [I-D.arkko-arch-internet-threat-model-guidance] further expands the threat model. Implementations and users SHOULD consider these threat models. In particular, even data sent protected to the other endpoint such as ID_CRED and EAD can be used for tracking, see Section 2.7 of [I-D.arkko-arch-internet-threat-model-guidance].

The fields ID_CRED_I, ID_CRED_R, EAD_2, EAD_3, and EAD_4 have variable length, and information regarding the length may leak to an attacker. A passive attacker may, e.g., be able to differentiate endpoints using identifiers of different length. To mitigate this information leakage an implementation may ensure that the fields have fixed length or use padding. An implementation may, e.g., only use fixed length identifiers like 'kid' of length 1. Alternatively, padding may be used (see Section 3.8.1) to hide the true length of, e.g., certificates by value in 'x5chain' or 'c5c'.

8.7. Denial-of-Service

EDHOC itself does not provide countermeasures against Denial-of-Service attacks. In particular, by sending a number of new or replayed message_1 an attacker may cause the Responder to allocate state, perform cryptographic operations, and amplify messages. To mitigate such attacks, an implementation SHOULD rely on lower layer mechanisms. For instance, when EDHOC is transferred as an exchange of CoAP messages, the CoAP server can use the Echo option defined in [RFC9175] which forces the CoAP client to demonstrate reachability at its apparent network address. To avoid an additional roundtrip

the Initiator can reduce the amplification factor by padding
message_1, i.e., using EAD_1, see Section 3.8.1.

An attacker can also send faked message_2, message_3, message_4, or
error in an attempt to trick the receiving party to send an error
message and discontinue the protocol. EDHOC implementations MAY
evaluate if a received message is likely to have been forged by an
attacker and ignore it without sending an error message or
discontinuing the protocol.

8.8. Implementation Considerations

The availability of a secure random number generator is essential
for the security of EDHOC. If no true random number generator is
available, a random seed MUST be provided from an external source
and used with a cryptographically secure pseudorandom number
generator. As each pseudorandom number must only be used once, an
implementation needs to get a unique input to the pseudorandom
number generator after reboot, or continuously store state in
nonvolatile memory. Appendix B.1.1 in [RFC8613] describes issues and
solution approaches for writing to nonvolatile memory. Intentionally
or unintentionally weak or predictable pseudorandom number
generators can be abused or exploited for malicious purposes.
[RFC8937] describes a way for security protocol implementations to
augment their (pseudo)random number generators using a long-term
private key and a deterministic signature function. This improves
randomness from broken or otherwise subverted random number
generators. The same idea can be used with other secrets and
functions such as a Diffie-Hellman function or a symmetric secret
and a PRF like HMAC or KMAC. It is RECOMMENDED to not trust a single
source of randomness and to not put unaugmented random numbers on
the wire.

Implementations might consider deriving secret and non-secret
randomness from different PRNG/PRF/KDF instances to limit the damage
if the PRNG/PRF/KDF turns out to be fundamentally broken. NIST
generally forbids deriving secret and non-secret randomness from the
same KDF instance, but this decision has been criticized by Krawczyk
[HKDFpaper] and doing so is common practice. In addition to IVs,
other examples are the challenge in EAP-TTLS, the RAND in 3GPP AKAs,
and the Session-Id in EAP-TLS 1.3. Note that part of KEYSTREAM_2 is
also non-secret randomness as it is known or predictable to an
attacker. As explained by Krawczyk, if any attack is mitigated by
the NIST requirement it would mean that the KDF is fully broken and
would have to be replaced anyway.

For many constrained IoT devices it is problematic to support
several crypto primitives. Existing devices can be expected to
support either ECDSA or EdDSA. If ECDSA is supported, "deterministic

ECDSA" as specified in [RFC6979] MAY be used. Pure deterministic elliptic-curve signatures such as deterministic ECDSA and EdDSA have gained popularity over randomized ECDSA as their security do not depend on a source of high-quality randomness. Recent research has however found that implementations of these signature algorithms may be vulnerable to certain side-channel and fault injection attacks due to their determinism. See e.g., Section 1 of [I-D.irtf-cfrg-det-sigs-with-noise] for a list of attack papers. As suggested in Section 2.1.1 of [RFC9053] this can be addressed by combining randomness and determinism.

Appendix D of [I-D.ietf-lwig-curve-representations] describes how Montgomery curves such as X25519 and X448 and (twisted) Edwards curves as curves such as Ed25519 and Ed448 can mapped to and from short-Weierstrass form for implementation on platforms that accelerate elliptic curve group operations in short-Weierstrass form.

All private keys, symmetric keys, and IVs MUST be secret. Implementations should provide countermeasures to side-channel attacks such as timing attacks. Intermediate computed values such as ephemeral ECDH keys and ECDH shared secrets MUST be deleted after key derivation is completed.

The Initiator and the Responder are responsible for verifying the integrity and validity of certificates. The selection of trusted CAs should be done very carefully and certificate revocation should be supported. The choice of revocation mechanism is left to the application. For example, in case of X.509 certificates, Certificate Revocation Lists [RFC5280] or OCSP [RFC6960] may be used. Verification of validity may require the use of a Real-Time Clock (RTC). The private authentication keys MUST be kept secret, only the Responder SHALL have access to the Responder's private authentication key and only the Initiator SHALL have access to the Initiator's private authentication key.

The Initiator and the Responder are allowed to select the connection identifier C_I and C_R, respectively, for the other party to use in the ongoing EDHOC protocol as well as in a subsequent application protocol (e.g., OSCORE [RFC8613]). The choice of connection identifier is not security critical in EDHOC but intended to simplify the retrieval of the right security context in combination with using short identifiers. If the wrong connection identifier of the other party is used in a protocol message it will result in the receiving party not being able to retrieve a security context (which will discontinue the protocol) or retrieve the wrong security context (which also discontinues the protocol as the message cannot be verified).

If two nodes unintentionally initiate two simultaneous EDHOC message exchanges with each other even if they only want to complete a single EDHOC message exchange, they MAY discontinue the exchange with the lexicographically smallest G_X. Note that in cases where several EDHOC exchanges with different parameter sets (method, COSE headers, etc.) are used, an attacker can affect which parameter set will be used by blocking some of the parameter sets.

If supported by the device, it is RECOMMENDED that at least the long-term private keys are stored in a Trusted Execution Environment (TEE, see for example [I-D.ietf-teep-architecture]) and that sensitive operations using these keys are performed inside the TEE. To achieve even higher security it is RECOMMENDED that additional operations such as ephemeral key generation, all computations of shared secrets, and storage of the PRK keys can be done inside the TEE. The use of a TEE aims at preventing code within that environment to be tampered with, and preventing data used by such code to be read or tampered with by code outside that environment.

Note that HKDF-Expand has a relatively small maximum output length of 255 * hash_length, where hash_length is the output size in bytes of the EDHOC hash algorithm of the selected cipher suite. This means that when SHA-256 is used as hash algorithm, PLAINTEXT_2 cannot be longer than 8160 bytes. This is probably not a limitation for most intended applications, but to be able to support for example long certificate chains or large external authorization data, there is a backwards compatible method specified in Appendix H.

The sequence of transcript hashes in EDHOC (TH_2, TH_3, TH_4) does not make use of a so called running hash. This is a design choice as running hashes are often not supported on constrained platforms.

When parsing a received EDHOC message, implementations MUST discontinue the protocol if the message does not comply with the CDDL for that message. Discontinuing the protocol if the received EDHOC message is not deterministic CBOR is RECOMMENDED.

9. IANA Considerations

This Section gives IANA Considerations and, unless otherwise noted, conforms with [RFC8126].

9.1. EDHOC Exporter Label Registry

IANA is requested to create a new registry under the new registry group "Ephemeral Diffie-Hellman Over COSE (EDHOC)" as follows:

Registry Name: EDHOC Exporter Label

Reference: [[this document]]

Label	Description	Reference
0	Derived OSCORE Master Secret	[[this document]]
1	Derived OSCORE Master Salt	[[this document]]
2-22	Unassigned	
23	Reserved	[[this document]]
24-32767	Unassigned	
32768-65535	Private Use	

Figure 13: EDHOC Exporter Label

```
+-------------+------------------------------------+
| Range       | Registration Procedures            |
+=============+====================================+
| 0-23        | Standards Action                   |
+-------------+------------------------------------+
| 24-32767    | Expert Review                      |
+-------------+------------------------------------+
```

9.2. EDHOC Cipher Suites Registry

IANA is requested to create a new registry under the new registry
group "Ephemeral Diffie-Hellman Over COSE (EDHOC)" as follows:

Registry Name: EDHOC Cipher Suites

Reference: [[this document]]

The columns of the registry are Value, Array and Description, where
Value is an integer and the other columns are text strings. The
initial contents of the registry are:

Value: -24
Array: N/A
Description: Private Use
Reference: [[this document]]

Value: -23
Array: N/A
Description: Private Use
Reference: [[this document]]

```
Value: -22
Array: N/A
Description: Private Use
Reference: [[this document]]

Value: -21
Array: N/A
Description: Private Use
Reference: [[this document]]

Value: 0
Array: 10, -16, 8, 4, -8, 10, -16
Description: AES-CCM-16-64-128, SHA-256, 8, X25519, EdDSA,
      AES-CCM-16-64-128, SHA-256
Reference: [[this document]]

Value: 1
Array: 30, -16, 16, 4, -8, 10, -16
Description: AES-CCM-16-128-128, SHA-256, 16, X25519, EdDSA,
      AES-CCM-16-64-128, SHA-256
Reference: [[this document]]

Value: 2
Array: 10, -16, 8, 1, -7, 10, -16
Description: AES-CCM-16-64-128, SHA-256, 8, P-256, ES256,
      AES-CCM-16-64-128, SHA-256
Reference: [[this document]]

Value: 3
Array: 30, -16, 16, 1, -7, 10, -16
Description: AES-CCM-16-128-128, SHA-256, 16, P-256, ES256,
      AES-CCM-16-64-128, SHA-256
Reference: [[this document]]

Value: 4
Array: 24, -16, 16, 4, -8, 24, -16
Description: ChaCha20/Poly1305, SHA-256, 16, X25519, EdDSA,
      ChaCha20/Poly1305, SHA-256
Reference: [[this document]]

Value: 5
Array: 24, -16, 16, 1, -7, 24, -16
Description: ChaCha20/Poly1305, SHA-256, 16, P-256, ES256,
      ChaCha20/Poly1305, SHA-256
Reference: [[this document]]
```

```
Value: 6
Array: 1, -16, 16, 4, -7, 1, -16
Description: A128GCM, SHA-256, 16, X25519, ES256,
     A128GCM, SHA-256
Reference: [[this document]]

Value: 23
Reserved
Reference: [[this document]]

Value: 24
Array: 3, -43, 16, 2, -35, 3, -43
Description: A256GCM, SHA-384, 16, P-384, ES384,
     A256GCM, SHA-384
Reference: [[this document]]

Value: 25
Array: 24, -45, 16, 5, -8, 24, -45
Description: ChaCha20/Poly1305, SHAKE256, 16, X448, EdDSA,
     ChaCha20/Poly1305, SHAKE256
Reference: [[this document]]
```

```
+----------------+------------------------------------+
| Range          | Registration Procedures            |
+================+====================================+
| -65536 to -25  | Specification Required             |
+----------------+------------------------------------+
| -20 to 23      | Standards Action with Expert Review |
+----------------+------------------------------------+
| 24 to 65535    | Specification Required             |
+----------------+------------------------------------+
```

9.3. EDHOC Method Type Registry

IANA is requested to create a new registry under the new registry
group "Ephemeral Diffie-Hellman Over COSE (EDHOC)" as follows:

Registry Name: EDHOC Method Type

Reference: [[this document]]

The columns of the registry are Value, Initiator Authentication Key,
and Responder Authentication Key, and Reference, where Value is an
integer and the key columns are text strings describing the
authentication keys.

The initial contents of the registry are shown in Figure 4. Method
23 is Reserved.

```
+----------------+-------------------------------------+
| Range          | Registration Procedures             |
+================+=====================================+
| -65536 to -25  | Specification Required              |
+----------------+-------------------------------------+
| -24 to 23      | Standards Action with Expert Review |
+----------------+-------------------------------------+
| 24 to 65535    | Specification Required              |
+----------------+-------------------------------------+
```

9.4. EDHOC Error Codes Registry

IANA is requested to create a new registry under the new registry
group "Ephemeral Diffie-Hellman Over COSE (EDHOC)" as follows:

Registry Name: EDHOC Error Codes

Reference: [[this document]]

The columns of the registry are ERR_CODE, ERR_INFO Type,
Description, and Reference, where ERR_CODE is an integer, ERR_INFO
is a CDDL defined type, and Description is a text string. The
initial contents of the registry are shown in Figure 10. Error code
23 is Reserved.

```
+----------------+-------------------------------------+
| Range          | Registration Procedures             |
+================+=====================================+
| -65536 to -25  | Expert Review                       |
+----------------+-------------------------------------+
| -24 to 23      | Standards Action                    |
+----------------+-------------------------------------+
| 24 to 65535    | Expert Review                       |
+----------------+-------------------------------------+
```

9.5. EDHOC External Authorization Data Registry

IANA is requested to create a new registry under the new registry
group "Ephemeral Diffie-Hellman Over COSE (EDHOC)" as follows:

Registry Name: EDHOC External Authorization Data

Reference: [[this document]]

The columns of the registry are Name, Label, Description, and
Reference, where Label is a non-negative integer and the other
columns are text strings. The initial contents of the registry is
shown in Figure 14. EAD label 23 is Reserved.

Name	Label	Description	Reference
Padding	0	Randomly generated CBOR byte string	[[this document]] Section 3.8.1

Figure 14: EAD Labels

```
+----------------+-------------------------------------+
| Range          | Registration Procedures             |
+================+=====================================+
| -65536 to -25  | Specification Required              |
+----------------+-------------------------------------+
| -24 to 23      | Standards Action with Expert Review |
+----------------+-------------------------------------+
| 24 to 65535    | Specification Required              |
+----------------+-------------------------------------+
```

9.6. COSE Header Parameters Registry

IANA is requested to register the following entries in the "COSE Header Parameters" registry under the registry group "CBOR Object Signing and Encryption (COSE)" (see Figure 15): The value of the 'kcwt' header parameter is a COSE Web Token (CWT) [RFC8392], and the value of the 'kccs' header parameter is a CWT Claims Set (CCS), see Section 1.4. The CWT/CCS must contain a COSE_Key in a 'cnf' claim [RFC8747]. The Value Registry for this item is empty and omitted from the table below.

Name	Label	Value Type	Description
kcwt	TBD1	COSE Messages	A CBOR Web Token (CWT) containing a COSE Key in a 'cnf' claim
kccs	TBD2	map / #6(map)	A CWT Claims Set (CCS) containing a COSE Key in a 'cnf' claim

Figure 15: COSE Header Parameter Labels

9.7. The Well-Known URI Registry

IANA is requested to add the well-known URI "edhoc" to the "Well-Known URIs" registry.

   *URI suffix: edhoc

   *Change controller: IETF

   *Specification document(s): [[this document]]

   *Related information: None

9.8. Media Types Registry

IANA is requested to add the media types "application/edhoc+cbor-seq" and "application/cid-edhoc+cbor-seq" to the "Media Types" registry.

9.8.1. application/edhoc+cbor-seq Media Type Registration

   *Type name: application

   *Subtype name: edhoc+cbor-seq

   *Required parameters: N/A

   *Optional parameters: N/A

   *Encoding considerations: binary

   *Security considerations: See Section 7 of this document.

   *Interoperability considerations: N/A

   *Published specification: [[this document]] (this document)

   *Applications that use this media type: To be identified

   *Fragment identifier considerations: N/A

   *Additional information:

      -Magic number(s): N/A

      -File extension(s): N/A

      -Macintosh file type code(s): N/A

   *Person & email address to contact for further information: See
    "Authors' Addresses" section.

*Intended usage: COMMON

    *Restrictions on usage: N/A

    *Author: See "Authors' Addresses" section.

    *Change Controller: IESG

9.8.2. application/cid-edhoc+cbor-seq Media Type Registration

    *Type name: application

    *Subtype name: cid-edhoc+cbor-seq

    *Required parameters: N/A

    *Optional parameters: N/A

    *Encoding considerations: binary

    *Security considerations: See Section 7 of this document.

    *Interoperability considerations: N/A

    *Published specification: [[this document]] (this document)

    *Applications that use this media type: To be identified

    *Fragment identifier considerations: N/A

    *Additional information:

       -Magic number(s): N/A

       -File extension(s): N/A

       -Macintosh file type code(s): N/A

    *Person & email address to contact for further information: See
     "Authors' Addresses" section.

    *Intended usage: COMMON

    *Restrictions on usage: N/A

    *Author: See "Authors' Addresses" section.

    *Change Controller: IESG

9.9. CoAP Content-Formats Registry

IANA has added the media types "application/edhoc+cbor-seq" and "application/cid-edhoc+cbor-seq" to the "CoAP Content-Formats" registry under the registry group "Constrained RESTful Environments (CoRE) Parameters".

Media Type	Encoding	ID	Reference
application/edhoc+cbor-seq	–	TBD5	[[this document]]
application/cid-edhoc+cbor-seq	–	TBD6	[[this document]]

Figure 16: CoAP Content-Format IDs

9.10. Resource Type (rt=) Link Target Attribute Values Registry

IANA is requested to add the resource type "core.edhoc" to the "Resource Type (rt=) Link Target Attribute Values" registry under the registry group "Constrained RESTful Environments (CoRE) Parameters".

   *Value: "core.edhoc"

   *Description: EDHOC resource.

   *Reference: [[this document]]

9.11. Expert Review Instructions

The IANA Registries established in this document are defined as "Expert Review", "Specification Required" or "Standards Action with Expert Review". This section gives some general guidelines for what the experts should be looking for, but they are being designated as experts for a reason so they should be given substantial latitude.

Expert reviewers should take into consideration the following points:

   *Clarity and correctness of registrations. Experts are expected to
    check the clarity of purpose and use of the requested entries.
    Expert needs to make sure the values of algorithms are taken from
    the right registry, when that is required. Experts should
    consider requesting an opinion on the correctness of registered
    parameters from relevant IETF working groups. Encodings that do
    not meet these objective of clarity and completeness should not
    be registered.

*Experts should take into account the expected usage of fields
   when approving code point assignment. The length of the encoded
   value should be weighed against how many code points of that
   length are left, the size of device it will be used on, and the
   number of code points left that encode to that size.

  *Even for "Expert Review" specifications are recommended. When
   specifications are not provided for a request where Expert Review
   is the assignment policy, the description provided needs to have
   sufficient information to verify the code points above.

10. References

10.1. Normative References

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Appendix A. Use with OSCORE and Transfer over CoAP

This appendix describes how to derive an OSCORE security context
when EDHOC is used to key OSCORE, and how to transfer EDHOC messages
over CoAP. The use of CoAP or OSCORE with EDHOC is optional, but if
you are using CoAP or OSCORE, then certain normative requirements
apply as detailed in the subsections.

A.1. Deriving the OSCORE Security Context

This section specifies how to use EDHOC output to derive the OSCORE
security context.

After successful processing of EDHOC message_3, Client and Server
derive Security Context parameters for OSCORE as follows (see
Section 3.2 of [RFC8613]):

  *The Master Secret and Master Salt SHALL be derived by using the
   EDHOC_Exporter interface, see Section 4.2.1:

     -The EDHOC Exporter Labels for deriving the OSCORE Master
      Secret and the OSCORE Master Salt, are the uints 0 and 1,
      respectively.

     -The context parameter is h'' (0x40), the empty CBOR byte
      string.

     -By default, oscore_key_length is the key length (in bytes) of
      the application AEAD Algorithm of the selected cipher suite

for the EDHOC session. Also by default, oscore_salt_length has value 8. The Initiator and Responder MAY agree out-of-band on a longer oscore_key_length than the default and on a different oscore_salt_length.

```
Master Secret = EDHOC_Exporter( 0, h'', oscore_key_length )
Master Salt   = EDHOC_Exporter( 1, h'', oscore_salt_length )
```

   *The AEAD Algorithm SHALL be the application AEAD algorithm of the selected cipher suite for the EDHOC session.

   *The HKDF Algorithm SHALL be the one based on the application hash algorithm of the selected cipher suite for the EDHOC session. For example, if SHA-256 is the application hash algorithm of the selected cipher suite, HKDF SHA-256 is used as HKDF Algorithm in the OSCORE Security Context.

   *The relationship between identifiers in OSCORE and EDHOC is specified in Section 3.3.3. The OSCORE Sender ID and Recipient ID SHALL be determined by the EDHOC connection identifiers C_R and C_I for the EDHOC session as shown in Figure 17.

EDHOC \ OSCORE	Sender ID	Recipient ID
Initiator	C R	C I
Responder	C I	C R

Figure 17: Usage of connection identifiers in OSCORE

Client and Server SHALL use the parameters above to establish an OSCORE Security Context, as per Section 3.2.1 of [RFC8613].

From then on, Client and Server retrieve the OSCORE protocol state using the Recipient ID, and optionally other transport information such as the 5-tuple.

A.2. Transferring EDHOC over CoAP

This section specifies how EDHOC can be transferred as an exchange of CoAP [RFC7252] messages. CoAP provides a reliable transport that can preserve packet ordering, flow control, and handle message duplication. CoAP can also perform fragmentation and protect against denial-of-service attacks. The underlying CoAP transport should be used in reliable mode, in particular when fragmentation is used, to

avoid, e.g., situations with hanging endpoints waiting for each other.

EDHOC may run with the Initiator either being CoAP client or CoAP server. We denote the former by the "forward message flow" (see Appendix A.2.1) and the latter by the "reverse message flow" (see Appendix A.2.2). By default we assume the forward message flow, but the roles SHOULD be chosen to protect the most sensitive identity, see Section 8.

According to this specification, EDHOC is transferred in POST requests to the Uri-Path: "/.well-known/edhoc" (see Section 9.7), and 2.04 (Changed) responses. An application may define its own path that can be discovered, e.g., using a resource directory [RFC9176]. Client applications can use the resource type "core.edhoc" to discover a server's EDHOC resource, i.e., where to send a request for executing the EDHOC protocol, see Section 9.10. An alternative transfer of the forward message flow is specified in [I-D.ietf-core-oscore-edhoc].

In order for the server to correlate a message received from a client to a message previously sent in the same EDHOC session over CoAP, messages sent by the client SHALL be prepended with the CBOR serialization of the connection identifier which the server has selected, see Section 3.4.1. This applies both to the forward and the reverse message flows. To indicate a new EDHOC session in the forward message flow, message_1 SHALL be prepended with the CBOR simple value true (0xf5), as a dummy identifier. Even if CoAP is carried over a reliable transport protocol such as TCP, the prepending of identifiers specified here SHALL be practiced to enable interoperability independent of how CoAP is transported.

The prepended identifiers are encoded in CBOR and thus self-delimiting. The representation of identifiers described in Section 3.3.2 SHALL be used. They are sent in front of the actual EDHOC message to keep track of messages in an EDHOC session, and only the part of the body following the identifier is used for EDHOC processing. In particular, the connection identifiers within the EDHOC messages are not impacted by the prepended identifiers.

An EDHOC message has media type application/edhoc+cbor-seq, whereas an EDHOC message prepended by a connection identifier has media type application/cid-edhoc+cbor-seq, see Section 9.9.

To protect against denial-of-service attacks, the CoAP server MAY respond to the first POST request with a 4.01 (Unauthorized) containing an Echo option [RFC9175]. This forces the Initiator to demonstrate reachability at its apparent network address. If message

fragmentation is needed, the EDHOC messages may be fragmented using the CoAP Block-Wise Transfer mechanism [RFC7959].

EDHOC error messages need to be transported in response to a message that failed (see Section 6). EDHOC error messages transported with CoAP are carried in the payload.

Note that the transport over CoAP can serve as a blueprint for other client-server protocols:

  *The client prepends the connection identifier selected by the server (or, for message_1, the CBOR simple value true) to any request message it sends.

  *The server does not send any such indicator, as responses are matched to request by the client-server protocol design.

A.2.1. The Forward Message Flow

In the forward message flow the CoAP client is the Initiator and the CoAP server is the Responder. This flow protects the client identity against active attackers and the server identity against passive attackers.

In the forward message flow, the CoAP Token enables correlation on the Initiator (client) side, and the prepended C_R enables correlation on the Responder (server) side.

  *EDHOC message_1 is sent in the payload of a POST request from the client to the server's resource for EDHOC, prepended with the dummy identifier true (0xf5) indicating a new session.

  *EDHOC message_2 or the EDHOC error message is sent from the server to the client in the payload of the response, in the former case with response code 2.04 (Changed), in the latter with response code as specified in Appendix A.2.3.

  *EDHOC message_3 or the EDHOC error message is sent from the client to the server's resource in the payload of a POST request, prepended with the connection identifier C_R.

  *If EDHOC message_4 is used, or in case of an error message, it is sent from the server to the client in the payload of the response, with response codes analogously to message_2. In case of an error message sent in response to message_4, it is sent analogously to error message sent in response to message_2.

An example of a successful EDHOC exchange over CoAP in the forward message flow is shown in Figure 18.

```
Client     Server
   |————————————▶|  Header: POST (Code=0.02)
   |    POST     |  Uri-Path: "/.well-known/edhoc"
   |             |  Content-Format: application/cid-edhoc+cbor-seq
   |             |  Payload: true, EDHOC message 1
   |             |
   |◀————————————|  Header: 2.04 Changed
   |    2.04     |  Content-Format: application/edhoc+cbor-seq
   |             |  Payload: EDHOC message 2
   |             |
   |————————————▶|  Header: POST (Code=0.02)
   |    POST     |  Uri-Path: "/.well-known/edhoc"
   |             |  Content-Format: application/cid-edhoc+cbor-seq
   |             |  Payload: C_R, EDHOC message 3
   |             |
   |◀————————————|  Header: 2.04 Changed
   |    2.04     |  Content-Format: application/edhoc+cbor-seq
   |             |  Payload: EDHOC message 4
   |             |
```
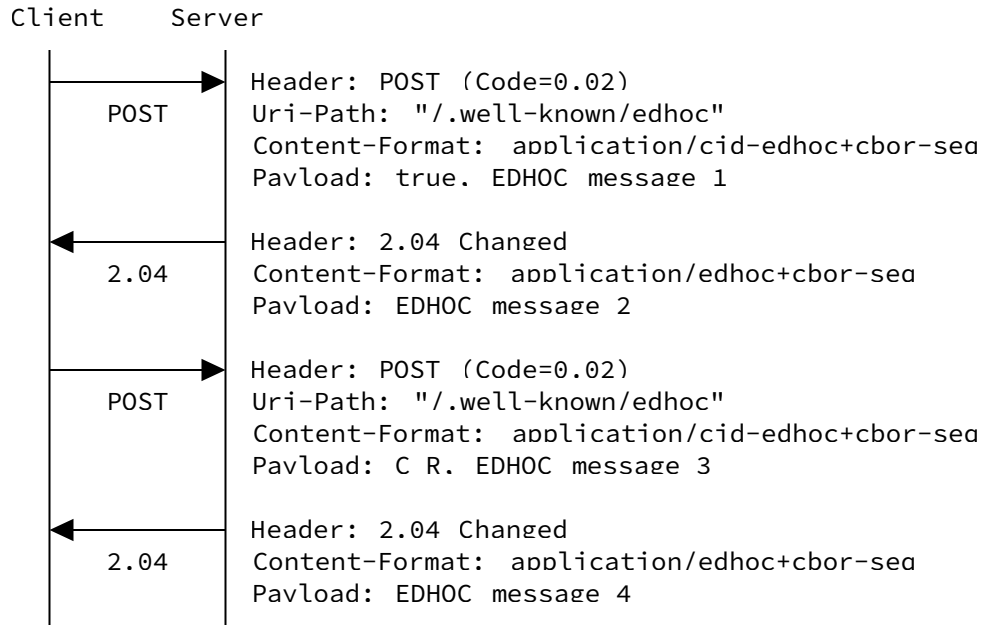
Figure 18: Example of the forward message flow.

The forward message flow of EDHOC can be combined with an OSCORE
exchange in a total of two round-trips, see
[I-D.ietf-core-oscore-edhoc].

A.2.2. The Reverse Message Flow

In the reverse message flow the CoAP client is the Responder and the
CoAP server is the Initiator. This flow protects the server identity
against active attackers and the client identity against passive
attackers.

In the reverse message flow, the CoAP Token enables correlation on
the Responder (client) side, and the prepended C_I enables
correlation on the Initiator (server) side.

  *To trigger a new session, the client makes an empty POST request
   to the server's resource for EDHOC.

  *EDHOC message_1 is sent from the server to the client in the
   payload of the response with response code 2.04 (Changed).

  *EDHOC message_2 or the EDHOC error message is sent from the
   client to the server's resource in the payload of a POST request,
   prepended with the connection identifier C_I.

  *EDHOC message_3 or the EDHOC error message is sent from the
   server to the client in the payload of the response, in the
```

former case with response code 2.04 (Changed), in the latter with
response code as specified in Appendix A.2.3.

*If EDHOC message_4 is used, or in case of an error message, it is
 sent from the client to the server's resource in the payload of a
 POST request, prepended with the connection identifier C_I. In
 case of an error message sent in response to message_4, it is
 sent analogously to an error message sent in response to
 message_2.

An example of a successful EDHOC exchange over CoAP in the reverse
message flow is shown in Figure 19.

```
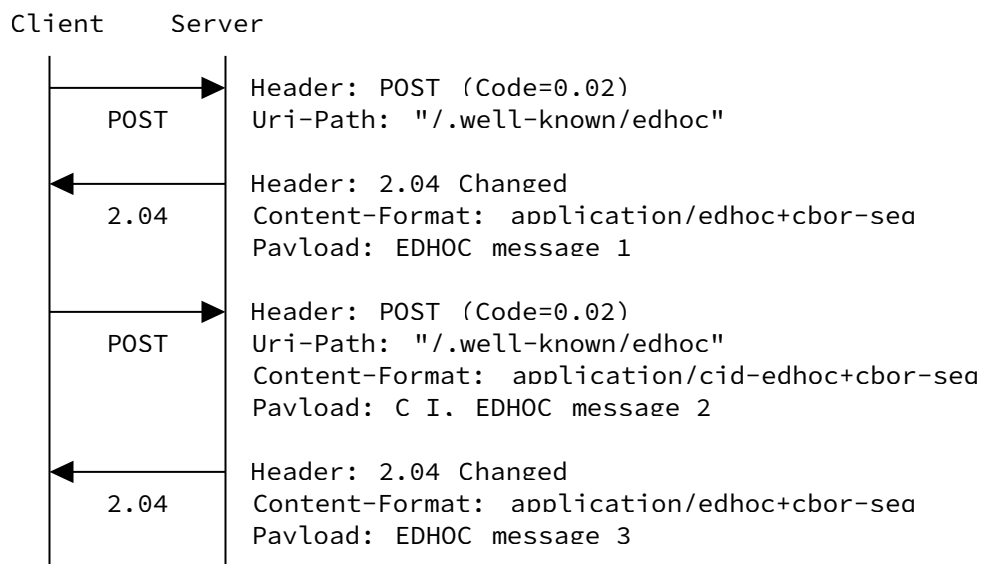Client     Server
   |          |
   |--------->|    Header: POST (Code=0.02)
   |  POST    |    Uri-Path: "/.well-known/edhoc"
   |          |
   |<---------|    Header: 2.04 Changed
   |  2.04    |    Content-Format: application/edhoc+cbor-seq
   |          |    Payload: EDHOC message 1
   |          |
   |--------->|    Header: POST (Code=0.02)
   |  POST    |    Uri-Path: "/.well-known/edhoc"
   |          |    Content-Format: application/cid-edhoc+cbor-seq
   |          |    Payload: C_I, EDHOC message 2
   |          |
   |<---------|    Header: 2.04 Changed
   |  2.04    |    Content-Format: application/edhoc+cbor-seq
   |          |    Payload: EDHOC message 3
   |          |
```

Figure 19: Example of the reverse message flow.

A.2.3. Errors in EDHOC over CoAP

When using EDHOC over CoAP, EDHOC error messages sent as CoAP
responses MUST be sent in the payload of error responses, i.e., they
MUST specify a CoAP error response code. In particular, it is
RECOMMENDED that such error responses have response code either 4.00
(Bad Request) in case of client error (e.g., due to a malformed
EDHOC message), or 5.00 (Internal Server Error) in case of server
error (e.g., due to failure in deriving EDHOC keying material). The
Content-Format of the error response MUST be set to application/
edhoc+cbor-seq, see Section 9.9.

Appendix B. Compact Representation

This section defines a format for compact representation based on the Elliptic-Curve-Point-to-Octet-String Conversion defined in Section 2.3.3 of [SECG].

As described in Section 4.2 of [RFC6090] the x-coordinate of an elliptic curve public key is a suitable representative for the entire point whenever scalar multiplication is used as a one-way function. One example is ECDH with compact output, where only the x-coordinate of the computed value is used as the shared secret.

In EDHOC, compact representation is used for the ephemeral public keys (G_X and G_Y), see Section 3.7. Using the notation from [SECG], the output is an octet string of length ceil((log2 q) / 8), where ceil(x) is the smallest integer not less than x. See [SECG] for a definition of q, M, X, xp, and ~yp. The steps in Section 2.3.3 of [SECG] are replaced by:

1. Convert the field element xp to an octet string X of length ceil((log2 q) / 8) octets using the conversion routine specified in Section 2.3.5 of [SECG].

2. Output M = X

The encoding of the point at infinity is not supported.

Compact representation does not change any requirements on validation, see Section 8.2. Using compact representation has some security benefits. An implementation does not need to check that the point is not the point at infinity (the identity element). Similarly, as not even the sign of the y-coordinate is encoded, compact representation trivially avoids so called "benign malleability" attacks where an attacker changes the sign, see [SECG].

The following may be needed for validation or compatibility with APIs that do not support compact representation or do not support the full [SECG] format:

  *If a compressed y-coordinate is required, then the value ~yp set to zero can be used. The compact representation described above can in such a case be transformed into the SECG point compressed format by prepending it with the single byte 0x02 (i.e., M = 0x02 || X).

  *If a uncompressed y-coordinate is required, then a y-coordinate has to be calculated following Section 2.3.4 of [SECG] or Appendix C of [RFC6090]. Any of the square roots (see [SECG] or

[RFC6090]) can be used. The uncompressed SECG format is M = 0x04
|| X || Y.

For example: The curve P-256 has the parameters (using the notation
in [RFC6090])

*p = $2^{256} - 2^{224} + 2^{192} + 2^{96} - 1$

*a = -3

*b = 410583637251521421293261297800472684091144410159937255
54835256314039467401291

Given an example x:

*x = 115792089183396302095546807154740558443406795108653336
39897069777278879 9766525

we can calculate y as the square root w = $(x^3 + a \cdot x + b)^{((p + 1)/4)}$
(mod p)

*y = 834387180070192806820075864918626005281451259964015754
16632522940595860276856

Note that this does not guarantee that (x, y) is on the correct
elliptic curve. A full validation according to Section 5.6.2.3.3 of
[SP-800-56A] can be achieved by also checking that 0 ≤ x < p and
that $y^2 \equiv x^3 + a \cdot x + b$ (mod p).

Appendix C. Use of CBOR, CDDL, and COSE in EDHOC

This Appendix is intended to help implementors not familiar with
CBOR [RFC8949], CDDL [RFC8610], COSE [RFC9052], and HKDF [RFC5869].

C.1. CBOR and CDDL

The Concise Binary Object Representation (CBOR) [RFC8949] is a data
format designed for small code size and small message size. CBOR
builds on the JSON data model but extends it by e.g., encoding
binary data directly without base64 conversion. In addition to the
binary CBOR encoding, CBOR also has a diagnostic notation that is
readable and editable by humans. The Concise Data Definition
Language (CDDL) [RFC8610] provides a way to express structures for
protocol messages and APIs that use CBOR. [RFC8610] also extends the
diagnostic notation.

CBOR data items are encoded to or decoded from byte strings using a
type-length-value encoding scheme, where the three highest order
bits of the initial byte contain information about the major type.
CBOR supports several different types of data items, in addition to

integers (int, uint), simple values, byte strings (bstr), and text strings (tstr), CBOR also supports arrays [] of data items, maps {} of pairs of data items, and sequences [RFC8742] of data items. Some examples are given below.

The EDHOC specification sometimes use CDDL names in CBOR diagnostic notation as in e.g., << ID_CRED_R, ? EAD_2 >>. This means that EAD_2 is optional and that ID_CRED_R and EAD_2 should be substituted with their values before evaluation. I.e., if ID_CRED_R = { 4 : h'' } and EAD_2 is omitted then << ID_CRED_R, ? EAD_2 >> = << { 4 : h'' } >>, which encodes to 0x43a10440. We also make use of the occurrence symbol "*", like in e.g., 2* int, meaning two or more CBOR integers.

For a complete specification and more examples, see [RFC8949] and [RFC8610]. We recommend implementors get used to CBOR by using the CBOR playground [CborMe].

Diagnostic	Encoded	Type
1	0x01	unsigned integer
24	0x1818	unsigned integer
-24	0x37	negative integer
-25	0x3818	negative integer
true	0xf5	simple value
h''	0x40	byte string
h'12cd'	0x4212cd	byte string
'12cd'	0x4431326364	byte string
"12cd"	0x6431326364	text string
{ 4 : h'cd' }	0xa10441cd	map
<< 1, 2, true >>	0x430102f5	byte string
[1, 2, true]	0x830102f5	array
(1, 2, true)	0x0102f5	sequence
1, 2, true	0x0102f5	sequence

Figure 20: Examples of use of CBOR and CDDL

C.2. CDDL Definitions

This sections compiles the CDDL definitions for ease of reference.

```
suites = [ 2* int ] / int

ead = (
  ead_label : int,
  ? ead_value : bstr,
)

EAD_1 = 1* ead
EAD_2 = 1* ead
EAD_3 = 1* ead
EAD_4 = 1* ead

message_1 = (
  METHOD : int,
  SUITES_I : suites,
  G_X : bstr,
  C_I : bstr / -24..23,
  ? EAD_1,
)

message_2 = (
  G_Y_CIPHERTEXT_2 : bstr,
  C_R : bstr / -24..23,
)

message_3 = (
  CIPHERTEXT_3 : bstr,
)

message_4 = (
  CIPHERTEXT_4 : bstr,
)

error = (
  ERR_CODE : int,
  ERR_INFO : any,
)

info = (
  info_label : int,
  context : bstr,
  length : uint,
)
```

C.3. COSE

CBOR Object Signing and Encryption (COSE) [RFC9052] describes how to
create and process signatures, message authentication codes, and
encryption using CBOR. COSE builds on JOSE, but is adapted to allow
more efficient processing in constrained devices. EDHOC makes use of

COSE_Key, COSE_Encrypt0, and COSE_Sign1 objects in the message
processing:

  *ECDH ephemeral public keys of type EC2 or OKP in message_1 and
   message_2 consist of the COSE_Key parameter named 'x', see
   Section 7.1 and 7.2 of [RFC9053]

  *The ciphertexts in message_3 and message_4 consist of a subset of
   the single recipient encrypted data object COSE_Encrypt0, which
   is described in Sections 5.2-5.3 of [RFC9052]. The ciphertext is
   computed over the plaintext and associated data, using an
   encryption key and an initialization vector. The associated data
   is an Enc_structure consisting of protected headers and
   externally supplied data (external_aad). COSE constructs the
   input to the AEAD [RFC5116] for message_i (i = 3 or 4, see
   Section 5.4 and Section 5.5, respectively) as follows:

     -Secret key K = K_i

     -Nonce N = IV_i

     -Plaintext P for message_i

     -Associated Data A = ["Encrypt0", h'', TH_i]

  *Signatures in message_2 of method 0 and 2, and in message_3 of
   method 0 and 1, consist of a subset of the single signer data
   object COSE_Sign1, which is described in Sections 4.2-4.4 of
   [RFC9052]. The signature is computed over a Sig_structure
   containing payload, protected headers and externally supplied
   data (external_aad) using a private signature key and verified
   using the corresponding public signature key. For COSE_Sign1, the
   message to be signed is:

    ["Signature1", protected, external_aad, payload]

   where protected, external_aad and payload are specified in
   Section 5.3 and Section 5.4.

Different header parameters to identify X.509 or C509 certificates
by reference are defined in [I-D.ietf-cose-x509] and
[I-D.ietf-cose-cbor-encoded-cert]:

  *by a hash value with the 'x5t' or 'c5t' parameters, respectively:

     -ID_CRED_x = { 34 : COSE_CertHash }, for x = I or R,

     -ID_CRED_x = { TBD3 : COSE_CertHash }, for x = I or R;

    *or by a URI with the 'x5u' or 'c5u' parameters, respectively:

      -ID_CRED_x = { 35 : uri }, for x = I or R,

      -ID_CRED_x = { TBD4 : uri }, for x = I or R.

When ID_CRED_x does not contain the actual credential, it may be very short, e.g., if the endpoints have agreed to use a key identifier parameter 'kid':

  *ID_CRED_x = { 4 : kid_x }, where kid_x : kid, for x = I or R. For further optimization, see [Section 3.5.3](#).

Note that ID_CRED_x can contain several header parameters, for example { x5u, x5t } or { kid, kid_context }.

ID_CRED_x MAY also identify the credential by value. For example, a certificate chain can be transported in ID_CRED_x with COSE header parameter c5c or x5chain, defined in [[I-D.ietf-cose-cbor-encoded-cert](#)] and [[I-D.ietf-cose-x509](#)] and credentials of type CWT and CCS can be transported with the COSE header parameters registered in [Section 9.6](#).

Appendix D. Authentication Related Verifications

EDHOC performs certain authentication related operations, see [Section 3.5](#), but in general it is necessary to make additional verifications beyond EDHOC message processing. What verifications are needed depend on the deployment, in particular the trust model and the security policies, but most commonly it can be expressed in terms of verifications of credential content.

EDHOC assumes the existence of mechanisms (certification authority or other trusted third party, pre-provisioning, etc.) for generating and distributing authentication credentials and other credentials, as well as the existence of trust anchors (CA certificates, trusted public keys, etc.). For example, a public key certificate or CWT may rely on a trusted third party whose public key is pre-provisioned, whereas a CCS or a self-signed certificate/CWT may be used when trust in the public key can be achieved by other means, or in the case of Trust on first use, see [Appendix D.5](#).

In this section we provide some examples of such verifications. These verifications are the responsibility of the application but may be implemented as part of an EDHOC library.

D.1. Validating the Authentication Credential

The authentication credential may contain, in addition to the authentication key, other parameters that needs to be verified. For example:

   *In X.509 and C509 certificates, signature keys typically have key usage "digitalSignature" and Diffie-Hellman public keys typically have key usage "keyAgreement" [RFC3279][RFC8410].

   *In X.509 and C509 certificates validity is expressed using Not After and Not Before. In CWT and CCS, the "exp" and "nbf" claims have similar meanings.

D.2. Identities

The application must decide on allowing a connection or not depending on the intended endpoint, and in particular whether it is a specific identity or in a set of identities. To prevent misbinding attacks, the identity of the endpoint is included in a MAC verified through the protocol. More details and examples are provided in this section.

Policies for what connections to allow are typically set based on the identity of the other endpoint, and endpoints typically only allow connections from a specific identity or a small restricted set of identities. For example, in the case of a device connecting to a network, the network may only allow connections from devices which authenticate with certificates having a particular range of serial numbers and signed by a particular CA. Conversely, a device may only be allowed to connect to a network which authenticates with a particular public key.

   *When a Public Key Infrastructure (PKI) is used with certificates, the identity is the subject whose unique name, e.g., a domain name, a Network Access Identifier (NAI), or an Extended Unique Identifier (EUI), is included in the endpoint's certificate.

   *Similarly, when a PKI is used with CWTs, the identity is the subject identified by the relevant claim(s), such as 'sub' (subject).

   *When PKI is not used (e.g., CCS, self-signed certificate/CWT) the identity is typically directly associated with the authentication key of the other party. For example, if identities can be expressed in the form of unique subject names assigned to public keys, then a binding to identity is achieved by including both public key and associated subject name in the authentication credential: CRED_I or CRED_R may be a self-signed certificate/CWT or CCS containing the authentication key and the subject name,

see [Section 3.5.2](#). Each endpoint thus needs to know the specific authentication key/unique associated subject name, or set of public authentication keys/unique associated subject names, which it is allowed to communicate with.

To prevent misbinding attacks in systems where an attacker can register public keys without proving knowledge of the private key, SIGMA [[SIGMA](#)] enforces a MAC to be calculated over the "identity". EDHOC follows SIGMA by calculating a MAC over the whole authentication credential, which in case of an X.509 or C509 certificate includes the "subject" and "subjectAltName" fields, and in the case of CWT or CCS includes the "sub" claim.

(While the SIGMA paper only focuses on the identity, the same principle is true for other information such as policies associated with the public key.)

D.3. Certification Path and Trust Anchors

When a Public Key Infrastructure (PKI) is used with certificates, the trust anchor is a Certification Authority (CA) certificate. Each party needs at least one CA public key certificate, or just the CA public key. The certification path contains proof that the subject of the certificate owns the public key in the certificate. Only validated public-key certificates are to be accepted.

Similarly, when a PKI is used with CWTs, each party needs to have at least one trusted third party public key as trust anchor to verify the end entity CWTs. The trusted third party public key can, e.g., be stored in a self-signed CWT or in a CCS.

The signature of the authentication credential needs to be verified with the public key of the issuer. X.509 and C509 certificates includes the "Issuer" field. In CWT and CCS, the "iss" claim has a similar meaning. The public key is either a trust anchor or the public key in another valid and trusted credential in a certification path from trust anchor to authentication credential.

Similar verifications as made with the authentication credential (see [Appendix D.1](#)) are also needed for the other credentials in the certification path.

When PKI is not used (CCS, self-signed certificate/CWT), the trust anchor is the authentication key of the other party, in which case there is no certification path.

D.4. Revocation Status

The application may need to verify that the credentials are not revoked, see [Section 8.8](#). Some use cases may be served by short-

lived credentials, for example, where the validity of the credential
is on par with the interval between revocation checks. But, in
general, credential lifetime and revocation checking are
complementary measures to control credential status. Revocation
information may be transported as External Authentication Data
(EAD), see Appendix E.

D.5. Unauthenticated Operation

EDHOC might be used without authentication by allowing the Initiator
or Responder to communicate with any identity except its own. Note
that EDHOC without mutual authentication is vulnerable to man-in-
the-middle attacks and therefore unsafe for general use. However, it
is possible to later establish a trust relationship with an unknown
or not-yet-trusted endpoint. Some examples:

*The EDHOC authentication credential can be verified out-of-band
 at a later stage.

*The EDHOC session key can be bound to an identity out-of-band at
 a later stage.

*Trust on first use (TOFU) can be used to verify that several
 EDHOC connections are made to the same identity. TOFU combined
 with proximity is a common IoT deployment model which provides
 good security if done correctly. Note that secure proximity based
 on short range wireless technology requires very low signal
 strength or very low latency.

Appendix E. Use of External Authorization Data

In order to reduce the number of messages and round trips, or to
simplify processing, external security applications may be
integrated into EDHOC by transporting related external authorization
data (EAD) in the messages.

The EAD format is specified in Section 3.8, this section contains
examples and further details of how EAD may be used with an
appropriate accompanying specification.

*One example is third party assisted authorization, requested with
 EAD_1, and an authorization artifact ("voucher", cf. [RFC8366])
 returned in EAD_2, see [I-D.selander-lake-authz].

*Another example is remote attestation, requested in EAD_2, and an
 Entity Attestation Token (EAT, [I-D.ietf-rats-eat]) returned in
 EAD_3.

*A third example is certificate enrolment, where a Certificate
 Signing Request (CSR, [RFC2986]) is included EAD_3, and the

issued public key certificate (X.509 [RFC5280], C509
[I-D.ietf-cose-cbor-encoded-cert]) or a reference thereof is
returned in EAD_4.

External authorization data should be considered unprotected by
EDHOC, and the protection of EAD is the responsibility of the
security application (third party authorization, remote attestation,
certificate enrolment, etc.). The security properties of the EAD
fields (after EDHOC processing) are discussed in Section 8.1.

The content of the EAD field may be used in the EDHOC processing of
the message in which they are contained. For example, authentication
related information like assertions and revocation information,
transported in EAD fields may provide input about trust anchors or
validity of credentials relevant to the authentication processing.
The EAD fields (like ID_CRED fields) are therefore made available to
the application before the message is verified, see details of
message processing in Section 5. In the first example above, a
voucher in EAD_2 made available to the application can enable the
Initiator to verify the identity or public key of the Responder
before verifying the signature. An application allowing EAD fields
containing authentication information thus may need to handle
authentication related verifications associated with EAD processing.

Conversely, the security application may need to wait for EDHOC
message verification to complete. In the third example above, the
validation of a CSR carried in EAD_3 is not started by the Responder
before EDHOC has successfully verified message_3 and proven the
possession of the private key of the Initiator.

The security application may reuse EDHOC protocol fields which
therefore need to be available to the application. For example, the
security application may use the same crypto algorithms as in the
EDHOC session and therefore needs access to the selected cipher
suite (or the whole SUITES_I). The application may use the ephemeral
public keys G_X and G_Y, as ephemeral keys or as nonces, see
[I-D.selander-lake-authz].

The processing of the EAD item (ead_label, ? ead_value) by the
security application needs to be described in the specification
where the ead_label is registered, see Section 9.5, including the
optional ead_value for each message and actions in case of errors.
An application may support multiple security applications that make
use of EAD, which may result in multiple EAD items in one EAD field,
see Section 3.8. Any dependencies on security applications with
previously registered EAD items needs to be documented, and the
processing needs to consider their simultaneous use.

Since data carried in EAD may not be protected, or be processed by the application before the EDHOC message is verified, special considerations need to be made such that it does not violate security and privacy requirements of the service which uses this data, see Section 8.5. The content in an EAD item may impact the security properties provided by EDHOC. Security applications making use of the EAD items must perform the necessary security analysis.

Appendix F. Application Profile Example

This appendix contains a rudimentary example of an application profile, see Section 3.9.

For use of EDHOC with application X the following assumptions are made:

1. Transfer in CoAP as specified in Appendix A.2 with requests expected by the CoAP server (= Responder) at /app1-edh, no Content-Format needed.

2. METHOD = 1 (I uses signature key, R uses static DH key.)

3. CRED_I is an IEEE 802.1AR IDevID encoded as a C509 certificate of type 0 [I-D.ietf-cose-cbor-encoded-cert].

   *R acquires CRED_I out-of-band, indicated in EAD_1.

   *ID_CRED_I = {4: h''} is a 'kid' with value the empty CBOR byte string.

4. CRED_R is a CCS of type OKP as specified in Section 3.5.2.

   *The CBOR map has parameters 1 (kty), -1 (crv), and -2 (x-coordinate).

   *ID_CRED_R is {TBD2 : CCS}. Editor's note: TBD2 is the COSE header parameter value of 'kccs', see Section 9.6

5. External authorization data is defined and processed as specified in [I-D.selander-lake-authz].

6. EUI-64 is used as the identity of the endpoint (see example in Section 3.5.2).

7. No use of message_4: the application sends protected messages from R to I.

Appendix G. EDHOC Message Deduplication

EDHOC by default assumes that message duplication is handled by the transport, in this section exemplified with CoAP.

Deduplication of CoAP messages is described in Section 4.5 of [RFC7252]. This handles the case when the same Confirmable (CON) message is received multiple times due to missing acknowledgement on the CoAP messaging layer. The recommended processing in [RFC7252] is that the duplicate message is acknowledged (ACK), but the received message is only processed once by the CoAP stack.

Message deduplication is resource demanding and therefore not supported in all CoAP implementations. Since EDHOC is targeting constrained environments, it is desirable that EDHOC can optionally support transport layers which do not handle message duplication. Special care is needed to avoid issues with duplicate messages, see Section 5.1.

The guiding principle here is similar to the deduplication processing on the CoAP messaging layer: a received duplicate EDHOC message SHALL NOT result in another instance of the next EDHOC message. The result MAY be that a duplicate next EDHOC message is sent, provided it is still relevant with respect to the current protocol state. In any case, the received message MUST NOT be processed more than once in the same EDHOC session. This is called "EDHOC message deduplication".

An EDHOC implementation MAY store the previously sent EDHOC message to be able to resend it.

In principle, if the EDHOC implementation would deterministically regenerate the identical EDHOC message previously sent, it would be possible to instead store the protocol state to be able to recreate and resend the previously sent EDHOC message. However, even if the protocol state is fixed, the message generation may introduce differences which compromise security. For example, in the generation of message_3, if I is performing a (non-deterministic) ECDSA signature (say, method 0 or 1, cipher suite 2 or 3) then PLAINTEXT_3 is randomized, but K_3 and IV_3 are the same, leading to a key and nonce reuse.

The EDHOC implementation MUST NOT store previous protocol state and regenerate an EDHOC message if there is a risk that the same key and IV are used for two (or more) distinct messages.

The previous message or protocol state MUST NOT be kept longer than what is required for retransmission, for example, in the case of CoAP transport, no longer than the EXCHANGE_LIFETIME (see Section 4.8.2 of [RFC7252]).

Appendix H. Long PLAINTEXT_2

By the definition of encryption of PLAINTEXT_2 with KEYSTREAM_2, it
is limited to lengths of PLAINTEXT_2 not exceeding the output of
EDHOC_KDF, see Section 4.1.2. If the EDHOC hash algorithm is SHA-2
then HKDF-Expand is used, which limits the length of the EDHOC_KDF
output to 255 · hash_length, where hash_length is the length of the
output of the EDHOC hash algorithm given by the cipher suite. For
example, with SHA-256 as EDHOC hash algorithm, the length of the
hash output is 32 bytes and the maximum length of PLAINTEXT_2 is 255
· 32 = 8160 bytes.

While PLAINTEXT_2 is expected to be much shorter than 8 kB for the
intended use cases, it seems nevertheless prudent to specify a
solution for the event that this should turn out to be a limitation.

A potential work-around is to use a cipher suite with a different
hash function. In particular, the use of KMAC removes all practical
limitations in this respect.

This section specifies a solution which works with any hash
function, by making use of multiple invocations of HKDF-Expand and
negative values of info_label.

Consider the PLAINTEXT_2 partitioned in parts P(i) of length equal
to M = 255 · hash_length, except possibly the last part P(last)
which has 0 < length ≤ M.

PLAINTEXT_2 = P(0) | P(1) | ... | P(last)

where | indicates concatenation.

The object is to define a matching KEYSTREAM_2 of the same length
and perform the encryption in the same way as defined in
Section 5.3.2:

CIPHERTEXT_2 = PLAINTEXT_2 XOR KEYSTREAM_2

Define the keystream as:

KEYSTREAM_2 = OKM(0) | OKM(1) | ... | OKM(last)

where

OKM(i) = EDHOC_KDF(PRK_2e, -i, TH_2, length(P(i)))

Note that if length(PLAINTEXT_2) ≤ M then P(0) = PLAINTEXT_2 and the
definition of KEYSTREAM_2 = OKM(0) coincides with Figure 8.

This describes the processing of the Responder when sending message_2. The Initiator makes the same calculations when receiving message_2, but interchanging PLAINTEXT_2 and CIPHERTEXT_2.

An application profile may specify if it supports or not the method described in this appendix.

Appendix I. EDHOC_KeyUpdate

To provide forward secrecy in an even more efficient way than re-running EDHOC, this section specifies the optional function EDHOC_KeyUpdate in terms of EDHOC_KDF and PRK_out.

When EDHOC_KeyUpdate is called, a new PRK_out is calculated as a "hash" of the old PRK_out using the EDHOC_Expand function as illustrated by the following pseudocode. The change of PRK_out causes a change to PRK_exporter which enables the derivation of new application keys superseding the old ones, using EDHOC_Exporter, see [Section 4.2.1](#).

```
EDHOC_KeyUpdate( context ):
   new PRK_out = EDHOC_KDF( old PRK_out, 11, context, hash_length )
   new PRK_exporter = EDHOC_KDF( new PRK_out, 10, h'', hash_length )
```

where hash_length denotes the output size in bytes of the EDHOC hash algorithm of the selected cipher suite.

The EDHOC_KeyUpdate takes a context as input to enable binding of the updated PRK_out to some event that triggered the key update. The Initiator and the Responder need to agree on the context, which can, e.g., be a counter or a pseudorandom number such as a hash. To provide forward secrecy the old PRK_out and keys derived from it (old PRK_exporter and old application keys) must be deleted as soon as they are not needed. When to delete the old keys and how to verify that they are not needed is up to the application.

An application using EDHOC_KeyUpdate needs to store PRK_out. Compromise of PRK_out leads to compromise of all keying material derived with the EDHOC_Exporter since the last invocation of the EDHOC_KeyUpdate function.

While this key update method provides forward secrecy it does not give as strong security properties as re-running EDHOC. EDHOC_KeyUpdate can be used to meet cryptographic limits and provide partial protection against key leakage, but it provides significantly weaker security properties than re-running EDHOC with ephemeral Diffie-Hellman. Even with frequent use of EDHOC_KeyUpdate, compromise of one session key compromises all future session keys, and an attacker therefore only needs to perform static key

exfiltration [RFC7624], which is less complicated and has a lower
risk profile than the dynamic case, see Section 8.1.

A similar method to do key update for OSCORE is KUDOS, see
[I-D.ietf-core-oscore-key-update].

Appendix J. Example Protocol State Machine

This appendix describes an example protocol state machine for the
Initiator and for the Responder. States are denoted in all capitals
and parentheses denote actions taken only in some circumstances.

Note that this state machine is just an example, and that details of
processing are omitted, for example:

  *When error messages are being sent (with one exception)

  *How credentials and EAD are processed by EDHOC and the
   application in the RCVD state

  *What verifications are made, which includes not only MACs and
   signatures

J.1. Initiator State Machine

The Initiator sends message_1, triggering the state machine to
transition from START to WAIT_M2, and waits for message_2.

If the incoming message is an error message then the Initiator
transitions from WAIT_M2 to ABORTED. In case of error code 2 (Wrong
Selected Cipher Suite), the Initiator remembers the supported cipher
suites for this particular Responder and transitions from ABORTED to
START. The message_1 that the Initiator subsequently sends takes
into account the cipher suites supported by the Responder.

Upon receiving a non-error message, the Initiator transitions from
WAIT_M2 to RCVD_M2 and processes the message. If a processing error
occurs on message_2, then the Initiator transitions from RCVD_M2 to
ABORTED. In case of successful processing of message_2, the
Initiator transitions from RCVD_M2 to VRFD_M2.

The Initiator prepares and processes message_3 for sending. If any
processing error is encountered, the Initiator transitions from
VRFD_M2 to ABORTED. If message_3 is successfully sent, the Initiator
transitions from VRFD_M2 to COMPLETED.

If the application profile includes message_4, then the Initiator
waits for message_4. If the incoming message is an error message
then the Initiator transitions from COMPLETED to ABORTED. Upon
receiving a non-error message, the Initiator transitions from

COMPLETED (="WAIT_M4") to RCVD_M4 and processes the message. If a
processing error occurs on message_4, then the Initiator transitions
from RCVD_M4 to ABORTED. In case of successful processing of
message_4, the Initiator transitions from RCVD_M4 to PERSISTED
(="VRFD_M4").

If the application profile does not include message_4, then the
Initiator waits for an incoming application message. If the
decryption and verification of the application message is
successful, then the the Initiator transitions from COMPLETED to
PERSISTED.

```
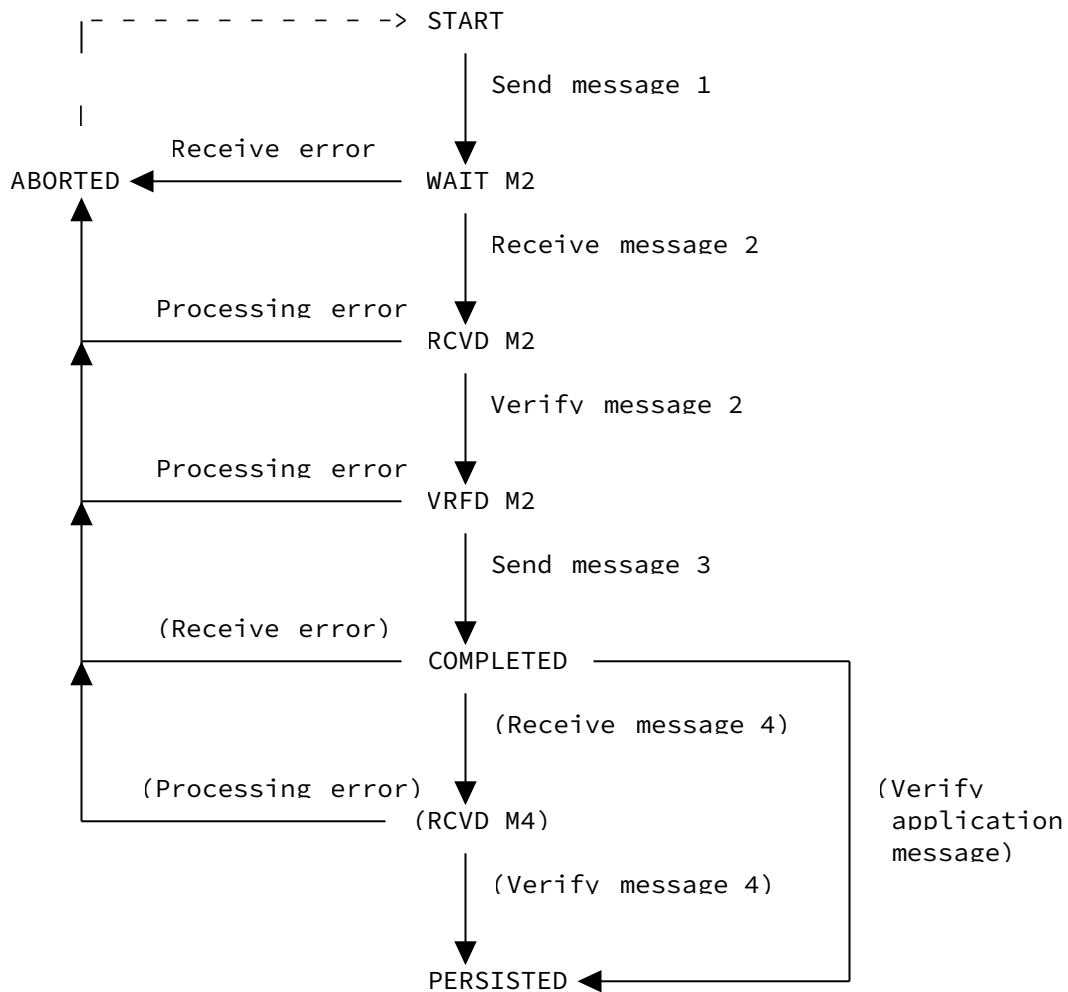      |- - - - - - - - -> START
      |                      |
      |                      | Send message 1
      |       Receive error  |
ABORTED <--------------------- WAIT M2
      ^                      |
      |                      | Receive message 2
      |   Processing error   v
      ^--------------------- RCVD M2
      |                      |
      |                      | Verify message 2
      |   Processing error   v
      ^--------------------- VRFD M2
      |                      |
      |                      | Send message 3
      |   (Receive error)    v
      ^--------------------- COMPLETED -----------------
      |                      |                         |
      |                      | (Receive message 4)     |
      |  (Processing error)  v               (Verify   |
      ^--------------------- (RCVD M4)        application
                             |                message)  |
                             | (Verify message 4)       |
                             v                          |
                          PERSISTED <-------------------
```

J.2. Responder State Machine

Upon receiving message_1, the Responder transitions from START to
RCVD_M1.

If a processing error occurs on message_1, the Responder transitions
from RCVD_M1 to ABORTED. This includes sending error message with
error code 2 (Wrong Selected Cipher Suite) if the selected cipher

suite in message_1 is not supported. In case of successful processing of message_1, the Responder transitions from RCVD_M1 to VRFD_M1.

The Responder prepares and processes message_2 for sending. If any processing error is encountered, the Responder transitions from VRFD_M1 to ABORTED. If message_2 is successfully sent, the Initiator transitions from VRFD_M2 to WAIT_M3, and waits for message_3.

If the incoming message is an error message then the Responder transitions from WAIT_M3 to ABORTED.

Upon receiving message_3, the Responder transitions from WAIT_M3 to RCVD_M3. If a processing error occurs on message_3, the Responder transitions from RCVD_M3 to ABORTED. In case of successful processing of message_3, the Responder transitions from RCVD_M3 to COMPLETED (="VRFD_M3").

If the application profile includes message_4, the Responder prepares and processes message_4 for sending. If any processing error is encountered, the Responder transitions from COMPLETED to ABORTED.

If message_4 is successfully sent, or if the application profile does not include message_4, the Responder transitions from COMPLETED to PERSISTED.

```
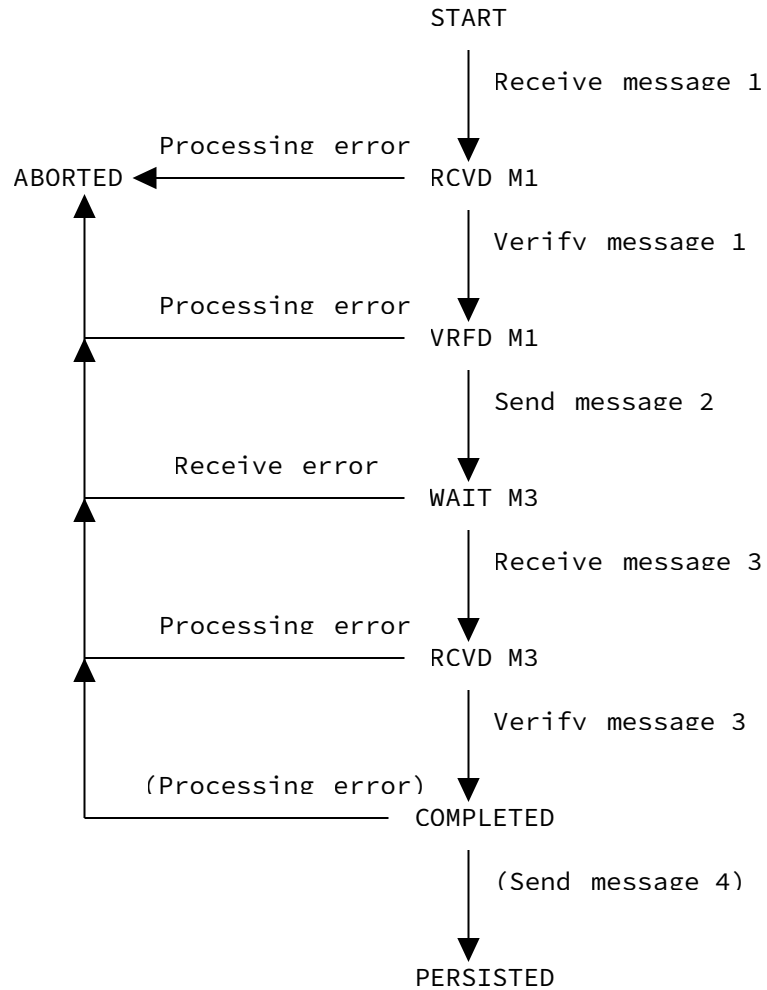                                        START
                                          │
                                          │  Receive message 1
                                          │
                   Processing error       ▼
         ABORTED ◀─────────────────────  RCVD M1
           ▲                              │
           │                              │  Verify message 1
           │                              │
           │       Processing error       ▼
           ├──────────────────────────── VRFD M1
           │                              │
           │                              │  Send message 2
           │                              │
           │       Receive error          ▼
           ├──────────────────────────── WAIT M3
           │                              │
           │                              │  Receive message 3
           │                              │
           │       Processing error       ▼
           ├──────────────────────────── RCVD M3
           │                              │
           │                              │  Verify message 3
           │                              │
           │    (Processing error)        ▼
           └──────────────────────────── COMPLETED
                                          │
                                          │  (Send message 4)
                                          │
                                          ▼
                                       PERSISTED
```

Appendix K. Change Log

RFC Editor: Please remove this appendix.

*From -18 to -19

  -Clarifications:

    oRelation to SIGMA

    oRole of Static DH

    oInitiator and Responder roles

    oTransport properties

    oConstruction of SUITES_I

    oMessage correlation, new subsection 3.4.1, replacing former
     appendix H

oRole of description about long PLAINTEXT_2

oead_label and ead_value

oMessage processing (Section 5)

oPadding

oCipher suite negotiation example

-Other updates:

oImproved and stricter normative text in Appendix A

oNaming and separate sections for the two message flows in
 Appendix A: Forward/Reverse message flow,

oTable index style captions

oAligning with COSE terminology: header map -> header_map

oAligning terminology, use of "_" instead of "-"

oPrefixing "EDHOC_" to functions

oUpdated list of security analysis papers

oNew appendix with example state machine

oAcknowledgements

oLanguage improvements by native English speakers

oUpdated IANA section with registration procedures

oNew and updated references

oRemoved appendix H

*From -17 to -18

-Padding realised as EAD with ead_label = 0, PAD fields removed

-Revised EAD syntax; ead is now EAD item; ead_value is now
 optional

-Clarifications of

oIdentifier representation

oAuthentication credentials

oRPK

oEncoding of ID_CRED with kid

oRepresentation of public keys, y-coordinate of ephemeral
 keys and validation

oProcessing after completed protocol

oMaking verifications available to the application

oRelation between EDHOC and OSCORE identifiers

-Terminology alignment in particular session / protocol;
 discontinue / terminate

-Updated CDDL

-Additional unicode encodings

-Large number of nits from WGLC

*From -16 to -17

-EDHOC-KeyUpdate moved to appendix

-Updated peer awareness properties based on SIGMA

-Clarify use of random connection identifiers

-Editorials related to appendix about messages with long
 PLAINTEXT_2

-Updated acknowledgments (have we forgotten someone else?
 please send email)

*From -15 to -16

-TH_2 used as salt in the derivation of PRK_2e

-CRED_R/CRED_I included in TH_3/TH_4

-Distinguish label used in info, exporter or elsewhere

-New appendix for optional handling arbitrarily large message_2

oinfo_label type changed to int to support this

-Updated security considerations

-Implementation note about identifiers which are bstr/int

-Clarifications, especifically about compact representation

-Type bug fix in CDDL section

*From -14 to -15

   -Connection identifiers and key identifiers are now byte
    strings

      oRepresented as CBOR bstr in the EDHOC message

         oUnless they happen to encode a one-byte CBOR int

      oMore examples

   -EAD updates and details

      oDefinition of EAD item

      oDefinition of critical / non-critical EAD item

   -New section in Appendix D: Unauthenticated Operation

   -Clarifications

      oLengths used in EDHOC-KDF

      oKey derivation from PRK_out

         oEDHOC-KeyUpdate and EDHOC-Exporter

      oPadding

   -Security considerations

      oWhen a change in a message is detected

      oConfidentiality in case of active attacks

      oConnection identifiers should be unpredictable

      oMaximum length of message_2

   -Minor bugs

*From -13 to -14

   -Merge of section 1.1 and 1.2

   -Connection and key identifiers restricted to be byte strings

-Representation of byte strings as one-byte CBOR ints (-24..23)

-Simplified mapping between EDHOC and OSCORE identifiers

-Rewrite of 3.5

   oClarification of authentication related operations
    performed by EDHOC

   oAuthentication related verifications, including old section
    3.5.1, moved to new appendix D

-Rewrite of 3.8

   oMove content about use of EAD to new appendix E

   oead_value changed to bstr

-EDHOC-KDF updated

   otranscript_hash argument removed

   oTH included in context argument

   olabel argument is now type uint, all labels replaced

-Key schedule updated

   oNew salts derived to avoid reuse of same key with expand
    and extract

   oPRK_4x3m renamed PRK_4e3m

   oK_4 and IV_4 derived from PRK_4e3m

   oNew PRK: PRK_out derived from PRK_4e3m and TH_4

   oClarified main output of EDHOC is the shared secret PRK_out

   oExporter defined by EDHOC-KDF and new PRK PRK_exporter
    derived from PRK_out

   oKey update defined by Expand instead of Extract

-All applications of EDHOC-KDF in one place

-Update of processing

   oEAD and ID_CRED passed to application when available

oidentity verification and credential retrieval omitted in
      protocol description

     oTranscript hash defined by plaintext messages instead of
      ciphertext

     oChanged order of input to TH_2

     oRemoved general G_X checking against selfie-attacks

   -Support for padding of plaintext

   -Updated compliance requirements

   -Updated security considerations

     oUpdated and more clear requirements on MAC length

     oClarification of key confirmation

     oForbid use of same key for signature and static DH

   -Updated appendix on message deduplication

   -Clarifications of

     oconnection identifiers

     ocipher suites, including negotiation

     oEAD

     oError messages

   -Updated media types

   -Applicability template renamed application profile

   -Editorials

*From -12 to -13

   -no changes

*From -12:

   -Shortened labels to derive OSCORE key and salt

   -ead_value changed to bstr

   -Removed general G_X checking against selfie-attacks

-Updated and more clear requirements on MAC length

-Clarifications from Kathleen, Stephen, Marco, Sean, Stefan,

-Authentication Related Verifications moved to appendix

-Updated MTI section and cipher suite

-Updated security considerations

*From -11 to -12:

-Clarified applicability to KEMs

-Clarified use of COSE header parameters

-Updates on MTI

-Updated security considerations

-New section on PQC

-Removed duplicate definition of cipher suites

-Explanations of use of COSE moved to Appendix C.3

-Updated internal references

*From -10 to -11:

-Restructured section on authentication parameters

-Changed UCCS to CCS

-Changed names and description of COSE header parameters for
 CWT/CCS

-Changed several of the KDF and Exporter labels

-Removed edhoc_aead_id from info (already in transcript_hash)

-Added MTI section

-EAD: changed CDDL names and added value type to registry

-Updated Figures 1, 2, and 3

-Some correction and clarifications

-Added core.edhoc to CoRE Resource Type registry

*From -09 to -10:

   -SUITES_I simplified to only contain the selected and more
    preferred suites

   -Info is a CBOR sequence and context is a bstr

   -Added kid to UCCS example

   -Separate header parameters for CWT and UCCS

   -CWT Confirmation Method kid extended to bstr / int

*From -08 to -09:

   -G_Y and CIPHERTEXT_2 are now included in one CBOR bstr

   -MAC_2 and MAC_3 are now generated with EDHOC-KDF

   -Info field "context" is now general and explicit in EDHOC-KDF

   -Restructured Section 4, Key Derivation

   -Added EDHOC MAC length to cipher suite for use with static DH

   -More details on the use of CWT and UCCS

   -Restructured and clarified Section 3.5, Authentication
    Parameters

   -Replaced 'kid2' with extension of 'kid'

   -EAD encoding now supports multiple ead types in one message

   -Clarified EAD type

   -Updated message sizes

   -Replaced "perfect forward secrecy" with "forward secrecy"

   -Updated security considerations

   -Replaced prepended 'null' with 'true' in the CoAP transport of
    message_1

   -Updated CDDL definitions

   -Expanded on the use of COSE

*From -07 to -08:

   -Prepended C_x moved from the EDHOC protocol itself to the
    transport mapping

   -METHOD_CORR renamed to METHOD, corr removed

   -Removed bstr_identifier and use bstr / int instead; C_x can
    now be int without any implied bstr semantics

   -Defined COSE header parameter 'kid2' with value type bstr /
    int for use with ID_CRED_x

   -Updated message sizes

   -New cipher suites with AES-GCM and ChaCha20 / Poly1305

   -Changed from one- to two-byte identifier of CNSA compliant
    suite

   -Separate sections on transport and connection id with further
    sub-structure

   -Moved back key derivation for OSCORE from draft-ietf-core-
    oscore-edhoc

   -OSCORE and CoAP specific processing moved to new appendix

   -Message 4 section moved to message processing section

*From -06 to -07:

   -Changed transcript hash definition for TH_2 and TH_3

   -Removed "EDHOC signature algorithm curve" from cipher suite

   -New IANA registry "EDHOC Exporter Label"

   -New application defined parameter "context" in EDHOC-Exporter

   -Changed normative language for failure from MUST to SHOULD
    send error

   -Made error codes non-negative and 0 for success

   -Added detail on success error code

   -Aligned terminology "protocol instance" -> "session"

   -New appendix on compact EC point representation

-Added detail on use of ephemeral public keys

-Moved key derivation for OSCORE to draft-ietf-core-oscore-
 edhoc

-Additional security considerations

-Renamed "Auxililary Data" as "External Authorization Data"

-Added encrypted EAD_4 to message_4

*From -05 to -06:

-New section 5.2 "Message Processing Outline"

-Optional inital byte C_1 = null in message_1

-New format of error messages, table of error codes, IANA
 registry

-Change of recommendation transport of error in CoAP

-Merge of content in 3.7 and appendix C into new section 3.7
 "Applicability Statement"

-Requiring use of deterministic CBOR

-New section on message deduplication

-New appendix containin all CDDL definitions

-New appendix with change log

-Removed section "Other Documents Referencing EDHOC"

-Clarifications based on review comments

*From -04 to -05:

-EDHOC-Rekey-FS -> EDHOC-KeyUpdate

-Clarification of cipher suite negotiation

-Updated security considerations

-Updated test vectors

-Updated applicability statement template

*From -03 to -04:

   -Restructure of section 1

   -Added references to C509 Certificates

   -Change in CIPHERTEXT_2 -> plaintext XOR KEYSTREAM_2 (test
    vector not updated)

   -"K_2e", "IV_2e" -> KEYSTREAM_2

   -Specified optional message 4

   -EDHOC-Exporter-FS -> EDHOC-Rekey-FS

   -Less constrained devices SHOULD implement both suite 0 and 2

   -Clarification of error message

   -Added exporter interface test vector

*From -02 to -03:

   -Rearrangements of section 3 and beginning of section 4

   -Key derivation new section 4

   -Cipher suites 4 and 5 added

   -EDHOC-EXPORTER-FS - generate a new PRK_4x3m from an old one

   -Change in CIPHERTEXT_2 -> COSE_Encrypt0 without tag (no change
    to test vector)

   -Clarification of error message

   -New appendix C applicability statement

*From -01 to -02:

   -New section 1.2 Use of EDHOC

   -Clarification of identities

   -New section 4.3 clarifying bstr_identifier

   -Updated security considerations

   -Updated text on cipher suite negotiation and key confirmation

   -Test vector for static DH

*From -00 to -01:

      -Removed PSK method

      -Removed references to certificate by value

Acknowledgments

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