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## 1. Changes in this version

The following changes were made in this version:

- \*Use of KEMRecipientInfo to communicate algorithm info;
- \*Editorial changes.

## 2. Introduction

In recent years, there has been a substantial amount of research on quantum computers -- machines that exploit quantum mechanical

phenomena to solve mathematical problems that are difficult or intractable for conventional computers. If large-scale quantum computers are ever built, they will be able to break many of the public-key cryptosystems currently in use. This would seriously compromise the confidentiality and integrity of digital communications on the Internet and elsewhere. Under such a threat model, the current key encapsulation mechanisms would be vulnerable.

Post-quantum key encapsulation mechanisms (PQ-KEM) are being developed in order to provide secure key establishment against an adversary with access to a quantum computer.

As the National Institute of Standards and Technology (NIST) is still in the process of selecting the new post-quantum cryptographic algorithms that are secure against both quantum and classical computers, the purpose of this document is to propose a generic "algorithm-agnostic" solution to protect in confidentiality the CMS envelopped-data content against the quantum threat : the KEM-TRANS mechanism.

Although this mechanism could thus be used with any key encapsulation mechanism, including post-quantum KEMs or hybrid KEMs.

This RFC nonetheless specifically specifies the case where the algorithm PQ-KEM algorithm is Kyber.

### 3. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [[RFC2119](#)] [[RFC8174](#)] when, and only when, they appear in all capitals, as shown here.

The following terms are used in this document:

BER: Basic Encoding Rules (BER) as defined in [[X.690](#)].

DER: Distinguished Encoding Rules as defined in [[X.690](#)].

### 4. Design Rationales

The Cryptographic Message Syntax (CMS) [[RFC5652](#)] defines two levels of encryptions in the Envelopped-Data Content section:

- \*the Content-encryption process which protects the data using a symmetric algorithm used with a content encryption key (CEK);
- \*the Key-encryption process which protects this CEK using a key transport mechanism.

One of the typical use case of the CMS Envelopped-Data Content is to randomly generate a CEK, encrypt the data with a symmetric algorithm using this CEK and individually send the CEK to one or more recipients protected by asymmetric cryptography in a RecipientInfo object.

To achieve this scenario with KEM primitives, it is necessary to define a new key transport mechanism that will fulfil the following requirements:

- \*the Key Transport Mechanism SHALL be secure against quantum computers.
- \*the Key Transport Mechanism SHALL take the Content-Encryption Key (CEK) as input.

According to NIST, a KEM generates a random secret and a ciphertext from which the recipient can extract the shared secret, meaning that a KEM can not be used straightforwardly as a key transport mechanism in the CMS "multi-recipients" context. The KEM-TRANS mechanism defined in this document aims to turn a KEM into a key transport scheme allowing the sender to distribute a randomly generated key to several recipients. The KEM-TRANS Key transport mechanism described in the following section fulfils the requirements listed above and is an adaptation of the RSA-KEM algorithm previously specified in [[RFC5990](#)]. The solution is also aligned with the hybrid public key encryption scheme described in [[RFC9180](#)].

## **5. KEM Key Transport Mechanism (KEM-TRANS)**

The KEM Key Transport Mechanism (KEM-TRANS) is a one-pass (store-and-forward) mechanism for transporting keying data to a recipient.

With this type of mechanism, a sender cryptographically encapsulates the keying data using the recipient's public key to obtain encrypted keying data. The recipient can then decapsulate the encrypted keying data using his private key to recover the plaintext keying data.

### **5.1. Underlying Components**

The KEM-TRANS requires use of the following underlying components, which are provided to KEM-TRANS as algorithm parameters.

- \*KEM, a Key Encapsulation Mechanism;
- \*KDF, a Key Derivation Function, which derives keying data of a specified length from a shared secret value;
- \*WRAP, a symmetric key-wrapping scheme, which encrypts keying Data using a key-encrypting key (KEK).

#### 5.1.1. KEM

A KEM is a cryptographic algorithm consisting of three functions :

- \*a key generation function **KEM.keygen** taking as input a security level and returning a key pair (private key and the associated public key) for this security level.
- \*an encapsulation function **KEM.encaps** taking a public key as input and returning a random session key and a ciphertext that is an encapsulation of the session key.
- \*a decapsulation function **KEM.decaps** taking as input a private key and a ciphertext and returning a session key.

#### 5.1.2. KDF

A key derivation function (KDF) is a cryptographic function that deterministically derives one or more secret keys from a secret value using a pseudorandom function. KDFs can be used to stretch keys into longer keys or to obtain keys of a required format.

If the session key obtained from the KEM algorithm is long enough to fit into the WRAP algorithm, then the KDF could be equal to the identity function.

#### 5.1.3. WRAP

A wrapping algorithm is a symmetric algorithm protecting data in confidentiality and integrity. It is especially designed to transport key material. the WRAP algorithm consists of two functions :

- \*a wrapping function **Wrap** taking a wrapping key and a plaintext key as input and returning a wrapped key.
- \*a decapsulation function **Unwrap** taking as input a wrapping key and a wrapped key and returning the plaintext key.

In the following, *kekLen* denotes the length in bytes of the wrapping key for the underlying symmetric key-wrapping scheme.

In this scheme, the length of the keying data to be transported MUST be among the lengths supported by the underlying symmetric key-wrapping scheme.

## 5.2. Recipient's Key Generation and Distribution

The KEM-TRANS described in the next sections assumes that the recipient has previously generated a key pair (*recipPrivKey* and *recipPubKey*) and has distributed this public key to the sender.

The protocols and mechanisms by which the key pair is securely generated and the public key is securely distributed are out of the scope of this document.

## 5.3. Sender's Operations

This process assumes that the following algorithm parameters have been selected:

\**KEM*: a key encapsulation mechanism, as defined above.

\**KDF*: a key derivation function, as defined above.

\**Wrap*: a symmetric key-wrapping algorithm, as defined above.

\**kekLen*: the length in bits of the key required for the Wrap algorithm.

This process assumes that the following input data has been provided:

\**recipPubKey*: the recipient's public key.

\**K*: the keying data to be transported, assumed to be a length that is compatible with the chosen Wrap algorithm.

This process outputs:

\**EK*: the encrypted keying data, from which the recipient will be able to retrieve *K*.

The sender performs the following operations:

1. Generate a shared secret *SS* and the associated ciphertext *CT* using the KEM encapsulation function and the recipient's public key *recipPubKey*:

$$(SS, CT) = \text{KEM.encaps}(\text{recipPubKey})$$

2. Derive a key-encrypting key *KEK* of length *kekLen* bytes from the shared secret *SS* using the underlying key derivation function:

$$\text{KEK} = \text{KDF}(SS, \text{kekLen})$$

3. Wrap the keying data  $K$  with the key-encrypting key  $KEK$  using the underlying key-wrapping scheme to obtain wrapped keying data  $WK$  of length  $wrappedKekLen$ :

$$WK = \text{Wrap}(KEK, K)$$

4. Concatenate the wrapped keying data  $WK$  of length  $wrappedKekLen$  and the ciphertext  $CT$  to obtain the encrypted keying data  $EK$ :

$$EK = (WK \parallel CT)$$

5. Output the encrypted keying data  $EK$ .

#### 5.4. Recipient's Operations

This process assumes that the following algorithm parameters have been communicated from the sender:

\* $KEM$ : a key encapsulation mechanism, as defined above.

\* $KDF$ : a key derivation function, as defined above.

\* $Wrap$ : a symmetric key-wrapping algorithm, as defined above.

\* $kekLen$ : the length in bits of the key required for the  $Wrap$  algorithm.

This process assumes that the following input data has been provided:

\* $recipPrivKey$ : the recipient's private key.

\* $EK$ : the encrypted keying data.

This process outputs:

\* $K$ : the keying data to be transported.

The recipient performs the following operations:

1. Separate the encrypted keying data  $EK$  into wrapped keying data  $WK$  of length  $wrappedKekLen$  and a ciphertext  $CT$  :

$$(WK \parallel CT) = EK$$

2. Decapsulate the ciphertext  $CT$  using the  $KEM$  decapsulation function and the recipient's private key to retrieve the shared secret  $SS$ :

$$SS = KEM.\text{decaps}(recipPrivKey, CT)$$

If the decapsulation operation outputs an error, output "decryption error", and stop.

3. Derive a key-encrypting key *KEK* of length *kekLen* bytes from the shared secret *SS* using the underlying key derivation function:

$$\text{KEK} = \text{KDF}(\text{SS}, \text{kekLen})$$

4. Unwrap the wrapped keying data *WK* with the key-encrypting key *KEK* using the underlying key-wrapping scheme to recover the keying data *K*:

$$\text{K} = \text{Unwrap}(\text{KEK}, \text{WK})$$

If the unwrapping operation outputs an error, output "decryption error", and stop.

5. Output the keying data *K*.

## 6. Use in CMS

The KEM Key Transport Mechanism MAY be employed for one or more recipients in the CMS envelopped-data content type (Section 6 of [RFC5652]), where the keying data *K* processed by the mechanism is the CMS content-encryption key (*CEK*).

### 6.1. RecipientInfo Conventions

When the KEM Key Transport Mechanism is employed for a recipient, the *RecipientInfo* alternative for that recipient MUST be *KEMRecipientInfo* as defined in draft-ietf-lamps-cms-kemri.

### 6.2. Certificate Conventions

The conventions specified in this section augment [RFC5280].

#### 6.2.1. Key Usage Extension

The intended application for the key MAY be indicated in the key usage certificate extension (see [RFC5280], Section 4.2.1.3). If the *keyUsage* extension is present in a certificate that conveys a public key with the *id-kem* object identifier as discussed above, then the key usage extension MUST contain only the value *keyEncipherment*

*digitalSignature*, *nonRepudiation*, *dataEncipherment*, *keyAgreement*, *keyCertSign*, *cRLSign*, *encipherOnly* and *decipherOnly* SHOULD NOT be present.

A key intended to be employed only with the KEM-TRANS SHOULD NOT also be employed for data encryption. Good cryptographic practice



employs a given key pair in only one scheme. This practice avoids the risk that vulnerability in one scheme may compromise the security of the other, and may be essential to maintain provable security.

### **6.2.2. Subject Public Key Info**

EDITOR'S NOTE' - TODO

Update this section according to the future evolution of draft-ietf-lamp

If the recipient wishes only to employ the KEM-TRANS with a given public key, the recipient **MUST** identify the public key in the certificate using the *id-kem* object identifier.

EDITOR'S NOTE' - TODO

Clarify that id-kem refers to the KEM algo (kyber512, kyber768 or kyber1 while KEM-TRANS refers to the KEM Transport mechanism

### **6.3. SMIME Capabilities Attribute Conventions**

[[RFC8551](#)], Section 2.5.2 defines the SMIMECapabilities signed attribute (defined as a SEQUENCE of SMIMECapability SEQUENCES) to be used to specify a partial list of algorithms that the software announcing the SMIMECapabilities can support. When constructing a signedData object, compliant software **MAY** include the SMIMECapabilities signed attribute announcing that it supports the KEM Key Transport Mechanism.

The SMIMECapability SEQUENCE representing the KEM Key Transport Mechanism **MUST** include the id-kem-trans object identifier in the capabilityID field and **MUST** include a GenericKemTransParameters value in the parameters field identifying the components with which the mechanism is to be employed.

The DER encoding of a SMIMECapability SEQUENCE is the same as the DER encoding of an AlgorithmIdentifier. Example DER encodings for typical sets of components are given in Appendix A.

## **7. Security Considerations**

EDITOR'S NOTE' - TODO

section to be completed

## **8. IANA Considerations**

Within the CMS, algorithms are identified by object identifiers (OIDs). With one exception, all of the OIDs used in this document were assigned in other IETF documents, in ISO/IEC standards documents, by the National Institute of Standards and Technology

(NIST). The two exceptions are the ASN.1 module's identifier and id-kem-transport that are both assigned in this document.

## 9. Acknowledgements

This document incorporates contributions and comments from a large group of experts. The Editors would especially like to acknowledge the expertise and tireless dedication of the following people, who attended many long meetings and generated millions of bytes of electronic mail and VOIP traffic over the past year in pursuit of this document:

EDITOR'S NOTE' - TODO  
section to be completed

We are grateful to all, including any contributors who may have been inadvertently omitted from this list.

This document borrows text from similar documents, including those referenced below. Thanks go to the authors of those documents.

"Copying always makes things easier and less error prone" -  
[[RFC8411](#)].

## 10. Annex A : ASN.1 Syntax

The syntax for the scheme is given in Appendix A.1.

The syntax for selected underlying components including those mentioned above is given in Appendix A.2.

The following object identifier prefixes are used in the definitions below:

```
nistAlgorithm OID ::= {  
    joint-iso-itu-t(2) country(16) us(840) organization(1)  
    gov(101) csor(3) nistAlgorithm(4)  
}  
  
smimeAlgorithm OID ::= { iso(1) member-body(2)  
    us(840) rsadsi(113549) pkcs(1) pkcs-9(9) smime(16) alg(3)  
}
```

### 10.1. Annex A1 : KEM-TRANS Key Transport Mechanism

The object identifier for the KEM Key Transport Mechanism is id-kem-trans, which is defined in this document as:

```
id-kem-trans OID ::= { TBD }
```

When `id-kem-trans` is used in an `AlgorithmIdentifier`, the parameters MUST employ the `GenericKemTransParameters` syntax. The syntax for `GenericKemTransParameters` is as follows:

```
GenericKemTransParameters ::= {  
    kem  KeyEncapsulationMechanism,  
    kdf  KeyDerivationFunction,  
    wrap KeyWrappingMechanism  
}
```

The fields of type `GenericKemTransParameters` have the following meanings:

\*`kem` identifies the underlying key encapsulation mechanism (KEM). This can be Kyber.

\*`kdf` identifies the underlying key derivation function (KDF). This can be any KDF from [[SP-800-56C-r2](#)]. `kdf` can be equal to *null* if the key encapsulation mechanism outputs a shared secret *SS* of size *kekLen*.

\*`wrap` identifies the underlying key wrapping mechanism (WRAP). This can be any wrapping mechanism from [[RFC5649](#)].

## 10.2. Annex A2 : Underlying Components

### 10.2.1. Key Encapsulation Mechanisms

KEM-TRANS can support any NIST KEM, including the post-quantum KEM Kyber. This RFC only specifies the use of Kyber.

The object identifier for KEM depends on the security level (128 bits, 192 bits or 256 bits)

```
id-kyber512 OID ::= { nistAlgorithm TBD }  
id-kyber768 OID ::= { nistAlgorithm TBD }  
id-kyber1024 OID ::= { nistAlgorithm TBD }
```

These object identifiers have no associated parameters.

```
kyber512 ALGORITHM ::= { OID id-kyber512 }  
kyber768 ALGORITHM ::= { OID id-kyber768 }  
kyber1024 ALGORITHM ::= { OID id-kyber1024 }
```

### 10.2.2. Key Derivation Functions

KEM-TRANS can support any KDF from [[RFC5869](#)]. This RFC only specifies the use of HKDF. The KDF can be bypassed if the key encapsulation mechanism outputs a shared secret *SS* of size *kekLen*. `kdf` is then equal to *null*.

The object identifier for KDF depends on the security level (128 bits, 192 bits or 256 bits)

```
id-alg-hkdf-with-sha256 OID ::= {  
    smimeAlgorithm id-alg-hkdf-with-sha256(28)  
}
```

```
id-alg-hkdf-with-sha384 OID ::= {  
    smimeAlgorithm id-alg-hkdf-with-sha384(29)  
}
```

```
id-alg-hkdf-with-sha512 OID ::= {  
    smimeAlgorithm id-alg-hkdf-with-sha512(30)  
}
```

EDITOR'S NOTE' - TO BE DISCUSSED :

Is there a RFC defining KDF with SHA3?

Should we extend the compatible KDFs to [SP-800-56C-r2]?

### 10.2.3. Key Wrapping Schemes

KEM-TRANS can support any wrapping mechanism from [[RFC5649](#)]. This RFC only specifies the use of aes256-Wrap.

The object identifiers for the AES Key Wrap depend on the size of the key-encrypting key. There are three object identifiers (see [[RFC5649](#)]):

```
id-aes128-Wrap OID ::= { nistAlgorithm aes(1) aes128-Wrap(5) }  
id-aes192-Wrap OID ::= { nistAlgorithm aes(1) aes192-Wrap(25) }  
id-aes256-Wrap OID ::= { nistAlgorithm aes(1) aes256-Wrap(45) }
```

These object identifiers have no associated parameters.

```
aes128-Wrap ALGORITHM ::= { OID id-aes128-Wrap }  
aes192-Wrap ALGORITHM ::= { OID id-aes192-Wrap }  
aes256-Wrap ALGORITHM ::= { OID id-aes256-Wrap }
```

### 10.2.4. Appendix A : Examples

EDITOR'S NOTE' - TODO

section to be completed

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