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**Hash Of Root Key Certificate Extension**  
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**Abstract**

This document specifies the Hash Of Root Key certificate extension. This certificate extension is carried in the self-signed certificate for a trust anchor, which is often called a Root Certification Authority (CA) certificate. This certificate extension unambiguously identifies the next public key that will be used at some point in the future as the next Root CA certificate, eventually replacing the current one.

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## [1.](#) Introduction

This document specifies the Hash Of Root Key X.509 version 3 certificate extension. The extension is an optional addition to the Internet X.509 Public Key Infrastructure Certificate and Certificate Revocation List (CRL) Profile [[RFC5280](#)]. The certificate extension facilitates the orderly transition from one Root Certification Authority (CA) public key to the next. It does so by publishing the hash value of the next generation public key in the current self-signed certificate. This hash value is a commitment to a particular public key in the next generation self-signed certificate. This commitment allows a relying party to unambiguously recognize the next generation self-signed certificate when it becomes available, install the new self-signed certificate in the trust anchor store, and eventually remove the previous one from the trust anchor store.

A Root CA Certificate MAY include the Hashed Root Key certificate extension to provide the hash value of the next public key that will be used by the Root CA.

### [1.1.](#) Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [BCP 14](#) [[RFC2119](#)][[RFC8174](#)] when, and only when, they appear in all capitals, as shown here.



## [1.2.](#) ASN.1

Certificates [[RFC5280](#)] are generated using ASN.1 [[X680](#)]; certificates are always encoded with the Distinguished Encoding Rules (DER) [[X690](#)].

## [2.](#) Overview

Before the initial deployment of the Root CA, the following are generated:

- R1 = The initial Root key pair
- R2 = The second generation Root key pair
- H2 = Thumbprint (hash) of the public key of R2
- C1 = Self-signed certificate for R1, which also contains H2

C1 is a self-signed certificate, and it contains H2 within the HashOfRootKey extension. C1 is distributed as part of the initial the system deployment. The HashOfRootKey certificate extension is described in [Section 3](#).

When the time comes to replace the initial Root CA certificate, R1, the following are generated:

- R3 = The third generation Root key pair
- H3 = Thumbprint (hash) the public key of R3
- C2 = Self-signed certificate for R2, which contains H3

This is an iterative process. That is, R4 and H4 are generated when it is time for C3 to replace C2. And so on.

The successor to the Root CA self-signed certificate can be delivered by any means. Whenever a new Root CA self-signed certificate is received, the recipient is able to verify that the potential Root CA certificate links back to a previously authenticated Root CA certificate with the hashOfRootKey certificate extension. That is, the recipient verifies the signature on the self-signed certificate and verifies that the hash of the DER-encoded SubjectPublicKeyInfo from the potential Root CA certificate matches the value from the HashOfRootKey certificate extension of the current Root CA certificate. Checking the self-signed certificate signature ensures that the certificate contains the subject name, public key algorithm identifier, and public key algorithm parameters intended by the key owner; these are important inputs to certification path validation as defined in [Section 6 of \[RFC5280\]](#). Checking the hash of the SubjectPublicKeyInfo ensures that the certificate contains the intended public key. If either check fails, then the potential Root CA certificate is not a valid replacement, and the recipient



continues to use the current Root CA certificate. If both checks succeed, then the recipient adds the potential Root CA certificate to the trust anchor store. As discussed in [Section 5](#), the recipient can remove the current Root CA certificate immediately in some situations. In other situations, the recipient waits an appropriate amount of time to ensure that existing certification paths continue to validate.

### 3. Hash Of Root Key Certificate Extension

The HashOfRootKey certificate extension MUST NOT be critical.

The following ASN.1 [[X680](#)][X690] syntax defines the HashOfRootKey certificate extension:

```
ext-HashOfRootKey EXTENSION ::= {      -- Only in Root CA certificates
    SYNTAX          HashedRootKey
    IDENTIFIED BY    id-ce-hashOfRootKey
    CRITICALITY      {FALSE} }

HashedRootKey ::= SEQUENCE {
    hashAlg          AlgorithmIdentifier,  -- Hash algorithm used
    hashValue        OCTET STRING }       -- Hash of DER-encoded
                                           -- SubjectPublicKeyInfo

id-ce-hashOfRootKey ::= OBJECT IDENTIFIER { 1 3 6 1 4 1 51483 2 1 }
```

The definitions of EXTENSION and HashAlgorithm can be found in [[RFC5912](#)].

The hashAlg indicates the one-way hash algorithm that was used to compute the hash value.

The hashValue contains the hash value computed from the next generation public key. The public key is DER-encoded SubjectPublicKeyInfo as defined in [[RFC5280](#)].

### 4. IANA Considerations

This document makes no requests of the IANA.

### 5. Operational Considerations

Guidance on the transition from one trust anchor to another is available in [Section 4.4 of \[RFC4210\]](#). In particular, the oldWithNew and newWithOld advice ensures that relying parties are able to validate certificates issued under the current Root CA certificate and the next generation Root CA certificate throughout the



transition. The notAfter field in the oldWithNew certificate MUST cover the validity period of all unexpired certificates issued under the old Root CA private key. Further, this advice SHOULD be followed by Root CAs to avoid the need for all relying parties to make the transition at the same time.

After issuing the oldWithNew and newWithOld certificates, the Root CA MUST stop using the old private key to sign certificates.

In enterprise and application-specific environments where a directory service or certificate repository is available, the oldWithNew and newWithOld certificates SHOULD be published before the successor to the current Root CA self-signed certificate is released. In environments without such a directory service or repository, recipients SHOULD keep both the old and replacement Root CA self-signed certificate in the trust anchor store for some amount of time to ensure that all end-entity certificates can be validated until they expire. The recipient MAY keep the old Root CA self-signed certificate until all of the certificates in the local cache that are subordinate to it have expired.

Certification path construction is more complex when multiple self-signed certificates in the trust anchor store have the same distinguished name. For this reason, the replacement Root CA self-signed certificate SHOULD contain a different distinguished name than the one it is replacing. One approach is to include a number as part of the name that is incremented with each generation, such as "Example CA", "Example CA G2", "Example CA G3", and so on.

Changing names from one generation to another can lead to confusion when reviewing the history of a trust anchor store. To assist with such review, a recipient MAY create an audit entry to capture the old and replacement self-signed certificates.

The Root CA must securely back up the yet-to-be-deployed key pair. If the Root CA stores the key pair in a hardware security module, and that module fails, the Root CA remains committed to the key pair that is no longer available. This leaves the Root CA with no alternative but to deploy a new self-signed certificate that contains a newly-generated key pair in the same manner as the initial self-signed certificate, thus losing the benefits of the Hash Of Root Key certificate extension altogether.

## 6. Security Considerations

The security considerations from [[RFC5280](#)] apply, especially the discussion of self-issued certificates.





The Hash Of Root Key certificate extension facilitates the orderly transition from one Root CA public key to the next by publishing the hash value of the next generation public key in the current certificate. This allows a relying party to unambiguously recognize the next generation public key when it becomes available; however, the full public key is not disclosed until the Root CA releases the next generation certificate. In this way, attackers cannot begin to analyze the public key before the next generation Root CA self-signed certificate is released.

The Root CA needs to ensure that the public key in the next generation certificate is as strong or stronger than the key that it is replacing. Of course, a significant advance in cryptanalytic capability can break the yet-to-be-deployed key pair. Such advances are rare and difficult to predict. If such an advance occurs, the Root CA remains committed to the now broken key. This leaves the Root CA with no alternative but to deploy a new self-signed certificate that contains a newly-generated key pair, most likely using a different signature algorithm, in the same manner as the initial self-signed certificate, thus losing the benefits of the Hash Of Root Key certificate extension altogether.

The Root CA needs to employ a hash function that is resistant to preimage attacks [[RFC4270](#)]. A first-preimage attack against the hash function would allow an attacker to find another input that results published hash value. For the attack to be successful, the input would have to be a valid SubjectPublicKeyInfo that contains a public key that corresponds to a private key known to the attacker. A second-preimage attack becomes possible once the Root CA releases the next generation public key, which makes the input to the hash function available to the attacker and everyone else. Again, the attacker needs to find a valid SubjectPublicKeyInfo that contains the public key that corresponds to a private key known to the attacker.

If an early release of the next generation public key occurs and the Root CA is concerned that attackers were given too much lead time to analyze that public key, then the Root CA can transition to a freshly generated key pair by rapidly performing two transitions. The first transition takes the Root CA to the key pair that suffered the early release, and it causes the Root CA to generate the subsequent Root key pair. The second transition occurs when the Root CA is confident that the population of relying parties have completed the first transition, and it takes the Root CA to the freshly generated key pair. Of course, the second transition also causes the Root CA to generate another key pair that is reserved for future use.

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## **7. Acknowledgements**

The Secure Electronic Transaction (SET) [[SET](#)] specification published by MasterCard and VISA in 1997 includes a very similar certificate extension. The SET certificate extension has essentially the same semantics, but the syntax fairly different.

CTIA - The Wireless Association is developing a public key infrastructure that will make use of the certificate extension described in this document.

Many thanks to Stefan Santesson, Jim Schaad, Daniel Kahn Gillmor, Joel Halpern, Paul Hoffman, and Rich Salz. Their review and comments have greatly improved the document, especially the Operational Considerations and Security Considerations sections.

## **8. References**

### **8.1. Normative References**

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- [RFC8174] Leiba, B., "Ambiguity of Uppercase vs Lowercase in [RFC 2119](#) Key Words", [BCP 14](#), [RFC 8174](#), DOI 10.17487/RFC8174, May 2017, <<https://www.rfc-editor.org/info/rfc8174>>.
- [X680] ITU-T, "Information technology -- Abstract Syntax Notation One (ASN.1): Specification of basic notation", ITU-T Recommendation X.680, 2015.
- [X690] ITU-T, "Information Technology -- ASN.1 encoding rules: Specification of Basic Encoding Rules (BER), Canonical Encoding Rules (CER) and Distinguished Encoding Rules (DER)", ITU-T Recommendation X.690, 2015.

## **[8.2.](#) Informative References**

- [SET] MasterCard and VISA, "SET Secure Electronic Transaction Specification -- Book 2: Programmer's Guide, Version 1.0", May 1997.

## **[Appendix A.](#) ASN.1 Module**

The following ASN.1 module provides the complete definition of the HashOfRootKey certificate extension.



```
HashedRootKeyCertExtn { 1 3 6 1 4 1 51483 0 1 }
```

```
DEFINITIONS IMPLICIT TAGS ::=
```

```
BEGIN
```

```
-- EXPORTS All
```

```
IMPORTS
```

```
AlgorithmIdentifier{}, DIGEST-ALGORITHM
```

```
FROM AlgorithmInformation-2009 -- [RFC5912]
```

```
{ iso(1) identified-organization(3) dod(6) internet(1)
  security(5) mechanisms(5) pkix(7) id-mod(0)
  id-mod-algorithmInformation-02(58) }
```

```
EXTENSION
```

```
FROM PKIX-CommonTypes-2009
```

```
{ iso(1) identified-organization(3) dod(6) internet(1)
  security(5) mechanisms(5) pkix(7) id-mod(0)
  id-mod-pkixCommon-02(57) } ;
```

```
--
```

```
-- Expand the certificate extensions list in [RFC5912]
```

```
--
```

```
CertExtensions EXTENSION ::= {
  ext-HashOfRootKey, ... }
```

```
--
```

```
-- HashOfRootKey Certificate Extension
```

```
--
```

```
ext-HashOfRootKey EXTENSION ::= { -- Only in Root CA certificates
  SYNTAX      HashedRootKey
  IDENTIFIED BY id-ce-hashOfRootKey
  CRITICALITY {FALSE} }
```

```
HashedRootKey ::= SEQUENCE {
  hashAlg      HashAlgorithmId, -- Hash algorithm used
  hashValue    OCTET STRING } -- Hash of DER-encoded
                                -- SubjectPublicKeyInfo
```

```
HashAlgorithmId ::= AlgorithmIdentifier {DIGEST-ALGORITHM, { ... }}
```

```
id-ce-hashOfRootKey OBJECT IDENTIFIER ::= { 1 3 6 1 4 1 51483 2 1 }
```

```
END
```





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