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**Restart Signaling for IS-IS**  
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Abstract

This document describes a mechanism for a restarting router to signal to its neighbors that it is restarting, allowing them to reestablish their adjacencies without cycling through the down state, while still correctly initiating database synchronization.

This document additionally describes a mechanism for a router to signal its neighbors that it is preparing to initiate a restart while maintaining forwarding plane state. This allows the neighbors to maintain their adjacencies until the router has restarted, but also allows the neighbors to bring the adjacencies down in the event of other topology changes.

This document additionally describes a mechanism for a restarting router to determine when it has achieved Link State Protocol Data Unit (LSP) database synchronization with its neighbors and a mechanism to optimize LSP database synchronization, while minimizing transient routing disruption when a router starts.

This document obsoletes [RFC 5306](#).

Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [BCP 14](#) [[RFC2119](#)] [[RFC8174](#)] when, and only when, they appear in all capitals, as shown here.

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## [1. Overview](#)

The Intermediate System to Intermediate System (IS-IS) routing protocol [[RFC1195](#)] [[ISO10589](#)] is a link state intra-domain routing protocol. Normally, when an IS-IS router is restarted, temporary disruption of routing occurs due to events in both the restarting router and the neighbors of the restarting router.

The router that has been restarted computes its own routes before achieving database synchronization with its neighbors. The results of this computation are likely to be non-convergent with the routes computed by other routers in the area/domain.

Neighbors of the restarting router detect the restart event and cycle their adjacencies with the restarting router through the down state. The cycling of the adjacency state causes the neighbors to regenerate their LSPs describing the adjacency concerned. This in turn causes a temporary disruption of routes passing through the restarting router.

In certain scenarios, the temporary disruption of the routes is highly undesirable. This document describes mechanisms to avoid or minimize the disruption due to both of these causes.

When an adjacency is reinitialized as a result of a neighbor restarting, a router does three things:

1. It causes its own LSP(s) to be regenerated, thus triggering SPF runs throughout the area (or in the case of Level 2, throughout the domain).
2. It sets SRMflags on its own LSP database on the adjacency concerned.
3. In the case of a Point-to-Point link, it transmits a complete set of Complete Sequence Number PDUs (CSNPs), over the adjacency.

In the case of a restarting router process, the first of these is highly undesirable, but the second is essential in order to ensure synchronization of the LSP database.

The third action above minimizes the number of LSPs that must be exchanged and, if made reliable, provides a means of determining when the LSP databases of the neighboring routers have been synchronized. This is desirable whether or not the router is being restarted (so



that the overload bit can be cleared in the router's own LSP, for example).

This document describes a mechanism for a restarting router to signal that it is restarting to its neighbors, and allow them to reestablish their adjacencies without cycling through the down state, while still correctly initiating database synchronization.

This document additionally describes a mechanism for a restarting router to determine when it has achieved LSP database synchronization with its neighbors and a mechanism to optimize LSP database synchronization and minimize transient routing disruption when a router starts.

It is assumed that the three-way handshake [[RFC5303](#)] is being used on Point-to-Point circuits.

## **2. Approach**

### **2.1. Timers**

Three additional timers, T1, T2, and T3, are required to support the functionality defined in this document.

An instance of the timer T1 is maintained per interface, and indicates the time after which an unacknowledged (re)start attempt will be repeated. A typical value might be 3 seconds.

An instance of the timer T2 is maintained for each LSP database (LSPDB) present in the system, i.e., for a Level 1/2 system, there will be an instance of the timer T2 for Level 1 and an instance for Level 2. This is the maximum time that the system will wait for LSPDB synchronization. A typical value might be 60 seconds.

A single instance of the timer T3 is maintained for the entire system. It indicates the time after which the router will declare that it has failed to achieve database synchronization (by setting the overload bit in its own LSP). This is initialized to 65535 seconds, but is set to the minimum of the remaining times of received IS-IS Hellos (IIHs) containing a restart TLV with the Restart Acknowledgement (RA) set and an indication that the neighbor has an adjacency in the "UP" state to the restarting router.

NOTE: The timer T3 is only used by a restarting router.



## 2.2. Restart TLV

A new TLV is defined to be included in IIH PDUs. The presence of this TLV indicates that the sender supports the functionality defined in this document and it carries flags that are used to convey information during a (re)start. All IIHs transmitted by a router that supports this capability MUST include this TLV.

Type 211

Length: Number of octets in the Value field (1 to (3 + ID Length))

Value

	No. of octets
+-----+    Flags    +-----+	1
+-----+   Remaining Time   +-----+	2
+-----+   Restarting Neighbor ID   +-----+	ID Length

Flags (1 octet)

```

  0  1  2  3  4  5  6  7
+---+---+---+---+---+---+---+
|Reserved|PA|PR|SA|RA|RR|
+---+---+---+---+---+---+---+

```

RR - Restart Request

RA - Restart Acknowledgement

SA - Suppress adjacency advertisement

PR - Restart is planned

PA - Planned restart acknowledgement

(Note: Remaining fields are required when the RA bit is set.)

Remaining Time (2 octets)

Remaining holding time (in seconds)

Restarting Neighbor System ID (ID Length octets)

The System ID of the neighbor to which an RA refers. Note: Implementations based on earlier versions of this document may not include this field in the TLV when the RA is set. In this case, a router that is expecting an RA on a LAN circuit SHOULD assume that the acknowledgement is directed at the local system.





### **2.2.1. Use of RR and RA Bits**

The RR bit is used by a (re)starting router to signal to its neighbors that a (re)start is in progress, that an existing adjacency SHOULD be maintained even under circumstances when the normal operation of the adjacency state machine would require the adjacency to be reinitialized, to request a set of CSNPs, and to request setting of the SRMflags.

The RA bit is sent by the neighbor of a (re)starting router to acknowledge the receipt of a restart TLV with the RR bit set.

When the neighbor of a (re)starting router receives an IIH with the restart TLV having the RR bit set, if there exists on this interface an adjacency in state "UP" with the same System ID, and in the case of a LAN circuit, with the same source LAN address, then, irrespective of the other contents of the "Intermediate System Neighbors" option (LAN circuits) or the "Point-to-Point Three-Way Adjacency" option (Point-to-Point circuits):

- a. the state of the adjacency is not changed. If this is the first IIH with the RR bit set that this system has received associated with this adjacency, then the adjacency is marked as being in "Restart mode" and the adjacency holding time is refreshed -- otherwise, the holding time is not refreshed. The "remaining time" transmitted according to (b) below MUST reflect the actual time after which the adjacency will now expire. Receipt of a normal IIH with the RR bit reset will clear the "Restart mode" state. This procedure allows the restarting router to cause the neighbor to maintain the adjacency long enough for restart to successfully complete, while also preventing repetitive restarts from maintaining an adjacency indefinitely. Whether or not an adjacency is marked as being in "Restart mode" has no effect on adjacency state transitions.
- b. immediately (i.e., without waiting for any currently running timer interval to expire, but with a small random delay of a few tens of milliseconds on LANs to avoid "storms") transmit over the corresponding interface an IIH including the restart TLV with the RR bit clear and the RA bit set, in the case of Point-to-Point adjacencies having updated the "Point-to-Point Three-Way Adjacency" option to reflect any new values received from the (re)starting router. (This allows a restarting router to quickly acquire the correct information to place in its hellos.) The "Remaining Time" MUST be set to the current time (in seconds) before the holding timer on this adjacency is due to expire. If the corresponding interface is a LAN interface, then the Restarting Neighbor System ID SHOULD be set to the System ID of



the router from which the IIH with the RR bit set was received. This is required to correctly associate the acknowledgement and holding time in the case where multiple systems on a LAN restart at approximately the same time. This IIH SHOULD be transmitted before any LSPs or SNPs are transmitted as a result of the receipt of the original IIH.

- c. if the corresponding interface is a Point-to-Point interface, or if the receiving router has the highest LnRouterPriority (with the highest source MAC (Media Access Control) address breaking ties) among those routers to which the receiving router has an adjacency in state "UP" on this interface whose IIHs contain the restart TLV, excluding adjacencies to all routers which are considered in "Restart mode" (note the actual DIS is NOT changed by this process), initiate the transmission over the corresponding interface of a complete set of CSNPs, and set SRMflags on the corresponding interface for all LSPs in the local LSP database.

Otherwise (i.e., if there was no adjacency in the "UP" state to the System ID in question), process the IIH as normal by reinitializing the adjacency and setting the RA bit in the returned IIH.

#### **2.2.2. Use of the SA Bit**

The SA bit is used by a starting router to request that its neighbor suppress advertisement of the adjacency to the starting router in the neighbor's LSPs.

A router that is starting has no maintained forwarding function state. This may or may not be the first time the router has started. If this is not the first time the router has started, copies of LSPs generated by this router in its previous incarnation may exist in the LSP databases of other routers in the network. These copies are likely to appear "newer" than LSPs initially generated by the starting router due to the reinitialization of LSP fragment sequence numbers by the starting router. This may cause temporary blackholes to occur until the normal operation of the update process causes the starting router to regenerate and flood copies of its own LSPs with higher sequence numbers. The temporary blackholes can be avoided if the starting router's neighbors suppress advertising an adjacency to the starting router until the starting router has been able to propagate newer versions of LSPs generated by previous incarnations.

When a router receives an IIH with the restart TLV having the SA bit set, if there exists on this interface an adjacency in state "UP" with the same System ID, and in the case of a LAN circuit, with the same source LAN address, then the router MUST suppress advertisement



of the adjacency to the neighbor in its own LSPs. Until an IIH with the SA bit clear has been received, the neighbor advertisement MUST continue to be suppressed. If the adjacency transitions to the "UP" state, the new adjacency MUST NOT be advertised until an IIH with the SA bit clear has been received.

Note that a router that suppresses advertisement of an adjacency MUST NOT use this adjacency when performing its SPF calculation. In particular, if an implementation follows the example guidelines presented in [[ISO10589](#)], Annex C.2.5, Step 0:b) "pre-load TENT with the local adjacency database", the suppressed adjacency MUST NOT be loaded into TENT.

### **2.2.3. Use of PR and PA Bits**

The PR bit is used by a router which is planning to initiate a restart to signal to its neighbors that it will be restarting.

The PA bit is sent by the neighbor of a router planning to restart to acknowledge receipt of a restart TLV with the PR bit set.

When the neighbor of a router planning a restart receives an IIH with the restart TLV having the PR bit set, if there exists on this interface an adjacency in state "UP" with the same System ID, and in the case of a LAN circuit, with the same source LAN address, then:

- a. if this is the first IIH with the PR bit set that this system has received associated with this adjacency, then the adjacency is marked as being in "Planned Restart state" and the adjacency holding time is refreshed -- otherwise, the holding time is not refreshed. The "remaining time" transmitted according to (b) below MUST reflect the actual time after which the adjacency will now expire. Receipt of a normal IIH with the PR bit reset will clear the "Planned Restart mode" state. This procedure allows the router planning a restart to cause the neighbor to maintain the adjacency long enough for restart to successfully complete. Whether or not an adjacency is marked as being in "Planned Restart mode" has no effect on adjacency state transitions.
- b. immediately (i.e., without waiting for any currently running timer interval to expire, but with a small random delay of a few tens of milliseconds on LANs to avoid "storms") transmit over the corresponding interface an IIH including the restart TLV with the PR bit clear and the PA bit set. The "Remaining Time" MUST be set to the current time (in seconds) before the holding timer on this adjacency is due to expire. If the corresponding interface is a LAN interface, then the Restarting Neighbor System ID SHOULD be set to the System ID of the router from which the IIH with the



PR bit set was received. This is required to correctly associate the acknowledgement and holding time in the case where multiple systems on a LAN are planning a restart at approximately the same time.

While a control plane restart is in progress it is expected that the restarting router will be unable to respond to topology changes. It is therefore useful to signal a planned restart (if the forwarding plane on the restarting router is maintained) so that the neighbors of the restarting router can determine whether it is safe to maintain the adjacency if other topology changes occur prior to the completion of the restart. Signalling a planned restart in the absence of maintained forwarding plane state is likely to lead to significant traffic loss and MUST NOT be done.

Neighbors of the router which has signaled planned restart SHOULD maintain the adjacency in a planned restart state until it receives an IIH with the RR bit set, receives an IIH with both PR and RR bits clear, or the adjacency holding time expires - whichever occurs first.

While the adjacency is in planned restart state the following actions MAY be taken:

- a. If additional topology changes occur, the adjacency which is in planned restart state MAY be brought down even though the hold time has not yet expired. Given that the neighbor which has signaled a planned restart is not expected to update its forwarding plane in response to signaling of the topology changes (since it is restarting) traffic which transits that node is at risk of being improperly forwarded. On a LAN circuit, if the router in planned restart state is the DIS at any supported level, the adjacency(ies) SHOULD be brought down whenever any LSP update is either generated or received so as to trigger a new DIS election. Failure to do so will compromise the reliability of the Update Process on that circuit. What other criteria are used to determine what topology changes will trigger bringing the adjacency down is a local implementation decision.
- b. If a BFD session to the neighbor which signals a planned restart is in the UP state and subsequently goes DOWN, the event MAY be ignored since it is possible this is an expected side effect of the restart. Use of the Control Plane Independent state as signalled in BFD control packets [[RFC5880](#)] SHOULD be considered in the decision to ignore a BFD Session DOWN event





- c. On a Point-to-Point circuit, transmission of LSPs, CSNPs, and PSNPs MAY be suppressed. It is expected that the PDUs will not be received.

### **2.3. Adjacency (Re)Acquisition**

Adjacency (re)acquisition is the first step in (re)initialization. Restarting and starting routers will make use of the RR bit in the restart TLV, though each will use it at different stages of the (re)start procedure.

#### **2.3.1. Adjacency Reacquisition during Restart**

The restarting router explicitly notifies its neighbor that the adjacency is being reacquired, and hence that it SHOULD NOT reinitialize the adjacency. This is achieved by setting the RR bit in the restart TLV. When the neighbor of a restarting router receives an IIH with the restart TLV having the RR bit set, if there exists on this interface an adjacency in state "UP" with the same System ID, and in the case of a LAN circuit, with the same source LAN address, then the procedures described in [Section 3.2.1](#) are followed.

A router that does not support the restart capability will ignore the restart TLV and reinitialize the adjacency as normal, returning an IIH without the restart TLV.

On restarting, a router initializes the timer T3, starts the timer T2 for each LSPDB, and for each interface (and in the case of a LAN circuit, for each level) starts the timer T1 and transmits an IIH containing the restart TLV with the RR bit set.

On a Point-to-Point circuit, the restarting router SHOULD set the "Adjacency Three-Way State" to "Init", because the receipt of the acknowledging IIH (with RA set) MUST cause the adjacency to enter the "UP" state immediately.

On a LAN circuit, the LAN-ID assigned to the circuit SHOULD be the same as that used prior to the restart. In particular, for any circuits for which the restarting router was previously DIS, the use of a different LAN-ID would necessitate the generation of a new set of pseudonode LSPs, and corresponding changes in all the LSPs referencing them from other routers on the LAN. By preserving the LAN-ID across the restart, this churn can be prevented. To enable a restarting router to learn the LAN-ID used prior to restart, the LAN-ID specified in an IIH with RR set MUST be ignored.

Transmission of "normal" IIHs is inhibited until the conditions described below are met (in order to avoid causing an unnecessary



adjacency initialization). Upon expiry of the timer T1, it is restarted and the IIH is retransmitted as above.

When a restarting router receives an IIH a local adjacency is established as usual, and if the IIH contains a restart TLV with the RA bit set (and on LAN circuits with a Restart Neighbor System ID that matches that of the local system), the receipt of the acknowledgement over that interface is noted. When the RA bit is set and the state of the remote adjacency is "UP", then the timer T3 is set to the minimum of its current value and the value of the "Remaining Time" field in the received IIH.

On a Point-to-Point link, receipt of an IIH not containing the restart TLV is also treated as an acknowledgement, since it indicates that the neighbor is not restart capable. However, since no CSNP is guaranteed to be received over this interface, the timer T1 is cancelled immediately without waiting for a complete set of CSNPs. Synchronization may therefore be deemed complete even though there are some LSPs which are held (only) by this neighbor (see [Section 3.4](#)). In this case, we also want to be certain that the neighbor will reinitialize the adjacency in order to guarantee that the SRMflags have been set on its database, thus ensuring eventual LSPDB synchronization. This is guaranteed to happen except in the case where the Adjacency Three-Way State in the received IIH is "UP" and the Neighbor Extended Local Circuit ID matches the extended local circuit ID assigned by the restarting router. In this case, the restarting router MUST force the adjacency to reinitialize by setting the local Adjacency Three-Way State to "DOWN" and sending a normal IIH.

In the case of a LAN interface, receipt of an IIH not containing the restart TLV is unremarkable since synchronization can still occur so long as at least one of the non-restarting neighboring routers on the LAN supports restart. Therefore, T1 continues to run in this case. If none of the neighbors on the LAN are restart capable, T1 will eventually expire after the locally defined number of retries.

In the case of a Point-to-Point circuit, the "LocalCircuitID" and "Extended Local Circuit ID" information contained in the IIH can be used immediately to generate an IIH containing the correct three-way handshake information. The presence of "Neighbor Extended Local Circuit ID" information that does not match the value currently in use by the local system is ignored (since the IIH may have been transmitted before the neighbor had received the new value from the restarting router), but the adjacency remains in the initializing state until the correct information is received.



In the case of a LAN circuit, the source neighbor information (e.g., SNPAAddress) is recorded and used for adjacency establishment and maintenance as normal.

When BOTH a complete set of CSNPs (for each active level, in the case of a Point-to-Point circuit) and an acknowledgement have been received over the interface, the timer T1 is cancelled.

Once the timer T1 has been cancelled, subsequent IIHs are transmitted according to the normal algorithms, but including the restart TLV with both RR and RA clear.

If a LAN contains a mixture of systems, only some of which support the new algorithm, database synchronization is still guaranteed, but the "old" systems will have reinitialized their adjacencies.

If an interface is active, but does not have any neighboring router reachable over that interface, the timer T1 would never be cancelled, and according to [Section 3.4.1.1](#), the SPF would never be run. Therefore, timer T1 is cancelled after some predetermined number of expirations (which MAY be 1).

### **2.3.2. Adjacency Acquisition during Start**

The starting router wants to ensure that in the event that a neighboring router has an adjacency to the starting router in the "UP" state (from a previous incarnation of the starting router), this adjacency is reinitialized. The starting router also wants neighboring routers to suppress advertisement of an adjacency to the starting router until LSP database synchronization is achieved. This is achieved by sending IIHs with the RR bit clear and the SA bit set in the restart TLV. The RR bit remains clear and the SA bit remains set in subsequent transmissions of IIHs until the adjacency has reached the "UP" state and the initial T1 timer interval (see below) has expired.

Receipt of an IIH with the RR bit clear will result in the neighboring router utilizing normal operation of the adjacency state machine. This will ensure that any old adjacency on the neighboring router will be reinitialized.

Upon receipt of an IIH with the SA bit set, the behavior described in [Section 3.2.2](#) is followed.

Upon starting, a router starts timer T2 for each LSPDB.

For each interface (and in the case of a LAN circuit, for each level), when an adjacency reaches the "UP" state, the starting router



starts a timer T1 and transmits an IIH containing the restart TLV with the RR bit clear and SA bit set. Upon expiry of the timer T1, it is restarted and the IIH is retransmitted with both RR and SA bits set (only the RR bit has changed state from earlier IIHs).

Upon receipt of an IIH with the RR bit set (regardless of whether or not the SA bit is set), the behavior described in [Section 2.2.1](#) is followed.

When an IIH is received by the starting router and the IIH contains a restart TLV with the RA bit set (and on LAN circuits with a Restart Neighbor System ID that matches that of the local system), the receipt of the acknowledgement over that interface is noted.

On a Point-to-Point link, receipt of an IIH not containing the restart TLV is also treated as an acknowledgement, since it indicates that the neighbor is not restart capable. Since the neighbor will have reinitialized the adjacency, this guarantees that SRMflags have been set on its database, thus ensuring eventual LSPDB synchronization. However, since no CSNP is guaranteed to be received over this interface, the timer T1 is cancelled immediately without waiting for a complete set of CSNPs. Synchronization may therefore be deemed complete even though there are some LSPs that are held (only) by this neighbor (see [Section 2.4](#)).

In the case of a LAN interface, receipt of an IIH not containing the restart TLV is unremarkable since synchronization can still occur so long as at least one of the non-restarting neighboring routers on the LAN supports restart. Therefore, T1 continues to run in this case. If none of the neighbors on the LAN are restart capable, T1 will eventually expire after the locally defined number of retries. The usual operation of the update process will ensure that synchronization is eventually achieved.

When BOTH a complete set of CSNPs (for each active level, in the case of a Point-to-Point circuit) and an acknowledgement have been received over the interface, the timer T1 is cancelled. Subsequent IIHs sent by the starting router have the RR and RA bits clear and the SA bit set in the restart TLV.

Timer T1 is cancelled after some predetermined number of expirations (which MAY be 1).

When the T2 timer(s) are cancelled or expire, transmission of "normal" IIHs (with RR, RA, and SA bits clear) will begin.





### **2.3.3. Multiple Levels**

A router that is operating as both a Level 1 and a Level 2 router on a particular interface **MUST** perform the above operations for each level.

On a LAN interface, it **MUST** send and receive both Level 1 and Level 2 IIHs and perform the CSNP synchronizations independently for each level.

On a Point-to-Point interface, only a single IIH (indicating support for both levels) is required, but it **MUST** perform the CSNP synchronizations independently for each level.

### **2.4. Database Synchronization**

When a router is started or restarted, it can expect to receive a complete set of CSNPs over each interface. The arrival of the CSNP(s) is now guaranteed, since an IIH with the RR bit set will be retransmitted until the CSNP(s) are correctly received.

The CSNPs describe the set of LSPs that are currently held by each neighbor. Synchronization will be complete when all these LSPs have been received.

When (re)starting, a router starts an instance of timer T2 for each LSPDB as described in [Section 3.3.1](#) or [Section 3.3.2](#). In addition to normal processing of the CSNPs, the set of LSPIDs contained in the first complete set of CSNPs received over each interface is recorded, together with their remaining lifetime. In the case of a LAN interface, a complete set of CSNPs **MUST** consist of CSNPs received from neighbors that are not restarting. If there are multiple interfaces on the (re)starting router, the recorded set of LSPIDs is the union of those received over each interface. LSPs with a remaining lifetime of zero are **NOT** so recorded.

As LSPs are received (by the normal operation of the update process) over any interface, the corresponding LSPID entry is removed (it is also removed if an LSP arrives before the CSNP containing the reference). When an LSPID has been held in the list for its indicated remaining lifetime, it is removed from the list. When the list of LSPIDs is empty and the timer T1 has been cancelled for all the interfaces that have an adjacency at this level, the timer T2 is cancelled.

At this point, the local database is guaranteed to contain all the LSP(s) (either the same sequence number or a more recent sequence number) that were present in the neighbors' databases at the time of



(re)starting. LSPs that arrived in a neighbor's database after the time of (re)starting may or may not be present, but the normal operation of the update process will guarantee that they will eventually be received. At this point, the local database is deemed to be "synchronized".

Since LSPs mentioned in the CSNP(s) with a zero remaining lifetime are not recorded, and those with a short remaining lifetime are deleted from the list when the lifetime expires, cancellation of the timer T2 will not be prevented by waiting for an LSP that will never arrive.

#### **2.4.1. LSP Generation and Flooding and SPF Computation**

The operation of a router starting, as opposed to restarting, is somewhat different. These two cases are dealt with separately below.

##### **2.4.1.1. Restarting**

In order to avoid causing unnecessary routing churn in other routers, it is highly desirable that the router's own LSPs generated by the restarting system are the same as those previously present in the network (assuming no other changes have taken place). It is important therefore not to regenerate and flood the LSPs until all the adjacencies have been re-established and any information required for propagation into the local LSPs is fully available. Ideally, the information is loaded into the LSPs in a deterministic way, such that the same information occurs in the same place in the same LSP (and hence the LSPs are identical to their previous versions). If this can be achieved, the new versions may not even cause SPF to be run in other systems. However, provided the same information is included in the set of LSPs (albeit in a different order, and possibly different LSPs), the result of running the SPF will be the same and will not cause churn to the forwarding tables.

In the case of a restarting router, none of the router's own LSPs are transmitted, nor are the router's own forwarding tables updated while the timer T3 is running.

Redistribution of inter-level information MUST be regenerated before this router's LSP is flooded to other nodes. Therefore, the Level-n non-pseudonode LSP(s) MUST NOT be flooded until the other level's T2 timer has expired and its SPF has been run. This ensures that any inter-level information that is to be propagated can be included in the Level-n LSP(s).

During this period, if one of the router's own (including pseudonodes) LSPs is received, which the local router does not



currently have in its own database, it is NOT purged. Under normal operation, such an LSP would be purged, since the LSP clearly should not be present in the global LSP database. However, in the present circumstances, this would be highly undesirable, because it could cause premature removal of a router's own LSP -- and hence churn in remote routers. Even if the local system has one or more of the router's own LSPs (which it has generated, but not yet transmitted), it is still not valid to compare the received LSP against this set, since it may be that as a result of propagation between Level 1 and Level 2 (or vice versa), a further router's own LSP will need to be generated when the LSP databases have synchronized.

During this period, a restarting router SHOULD send CSNPs as it normally would. Information about the router's own LSPs MAY be included, but if it is included it MUST be based on LSPs that have been received, not on versions that have been generated (but not yet transmitted). This restriction is necessary to prevent premature removal of an LSP from the global LSP database.

When the timer T2 expires or is cancelled indicating that synchronization for that level is complete, the SPF for that level is run in order to derive any information that is required to be propagated to another level, but the forwarding tables are not yet updated.

Once the other level's SPF has run and any inter-level propagation has been resolved, the router's own LSPs can be generated and flooded. Any own LSPs that were previously ignored, but that are not part of the current set of own LSPs (including pseudonodes), MUST then be purged. Note that it is possible that a Designated Router change may have taken place, and consequently the router SHOULD purge those pseudonode LSPs that it previously owned, but that are now no longer part of its set of pseudonode LSPs.

When all the T2 timers have expired or been cancelled, the timer T3 is cancelled and the local forwarding tables are updated.

If the timer T3 expires before all the T2 timers have expired or been cancelled, this indicates that the synchronization process is taking longer than the minimum holding time of the neighbors. The router's own LSP(s) for levels that have not yet completed their first SPF computation are then flooded with the overload bit set to indicate that the router's LSPDB is not yet synchronized (and therefore other routers MUST NOT compute routes through this router). Normal operation of the update process resumes, and the local forwarding tables are updated. In order to prevent the neighbor's adjacencies from expiring, IIHs with the normal interface value for the holding time are transmitted over all interfaces with neither RR nor RA set



in the restart TLV. This will cause the neighbors to refresh their adjacencies. The router's own LSP(s) will continue to have the overload bit set until timer T2 has expired or been cancelled.

#### **2.4.1.2. Starting**

In the case of a starting router, as soon as each adjacency is established, and before any CSNP exchanges, the router's own zeroth LSP is transmitted with the overload bit set. This prevents other routers from computing routes through the router until it has reliably acquired the complete set of LSPs. The overload bit remains set in subsequent transmissions of the zeroth LSP (such as will occur if a previous copy of the router's own zeroth LSP is still present in the network) while any timer T2 is running.

When all the T2 timers have been cancelled, the router's own LSP(s) MAY be regenerated with the overload bit clear (assuming the router is not in fact overloaded, and there is no other reason, such as incomplete BGP convergence, to keep the overload bit set) and flooded as normal.

Other LSPs owned by this router (including pseudonodes) are generated and flooded as normal, irrespective of the timer T2. The SPF is also run as normal and the Routing Information Base (RIB) and Forwarding Information Base (FIB) updated as routes become available.

To avoid the possible formation of temporary blackholes, the starting router sets the SA bit in the restart TLV (as described in [Section 3.3.2](#)) in all IIHs that it sends.

When all T2 timers have been cancelled, the starting router MUST transmit IIHs with the SA bit clear.

### **3. State Tables**

This section presents state tables that summarize the behaviors described in this document. Other behaviors, in particular adjacency state transitions and LSP database update operation, are NOT included in the state tables except where this document modifies the behaviors described in [[ISO10589](#)] and [[RFC5303](#)].

The states named in the columns of the tables below are a mixture of states that are specific to a single adjacency (ADJ suppressed, ADJ Seen RA, ADJ Seen CSNP) and states that are indicative of the state of the protocol instance (Running, Restarting, Starting, SPF Wait).

Three state tables are presented from the point of view of a running router, a restarting router, and a starting router.





**3.1. Running Router**

Event	Running	ADJ suppressed
RX PR	Set Planned Restart state. Send PA	
RX PR clr and RR clr	Clear Planned Restart State	
RX RR	Maintain ADJ State Send RA Set SRM, send CSNP (Note 1) Update Hold Time, set Restart Mode (Note 2)	
RX RR clr	Clr Restart mode	
RX SA	Suppress IS neighbor TLV in LSP(s) Goto ADJ Suppressed	
RX SA clr		Unsuppress IS neighbor TLV in LSP(s) Goto Running

Note 1: CSNPs are sent by routers in accordance with [Section 2.2.1c](#)

Note 2: If Restart Mode clear

**3.2. Restarting Router**

Event	Restarting	ADJ Seen RA	ADJ Seen CSNP	SPF Wait
Restart planned	Send PR			
Planned restart canceled	Send PR clr			
Router	Send IIH/RR			



restarts	ADJ Init			
	Start T1,T2,T3			
-----+				
RX RR	Send RA			
-----+				
RX RA	Adjust T3		Cancel T1	
	Goto ADJ Seen RA		Adjust T3	
-----+				
RX CSNP set	Goto ADJ Seen CSNP	Cancel T1		
-----+				
RX IIH w/o Restart TLV	Cancel T1 (Point- to-point only)			
-----+				
T1 expires	Send IIH/RR	Send IIH/RR	Send IIH/RR	
	Restart T1	Restart T1	Restart T1	
-----+				
T1 expires nth time	Send IIH/ normal	Send IIH/ normal	Send IIH/ normal	
-----+				
T2 expires	Trigger SPF			
	Goto SPF Wait			
-----+				
T3 expires	Set overload bit			
	Flood local LSPs			
	Update fwd plane			
-----+				
LSP DB Sync	Cancel T2, and T3			
	Trigger SPF			
	Goto SPF wait			
-----+				
All SPF done				Clear overload bit Update fwd plane Flood local LSPs Goto Running
=====				

### 3.3. Starting Router



Event	Starting	ADJ Seen RA	ADJ Seen CSNP
Router starts	Send IIH/SA Start T1,T2		
RX RR	Send RA		
RX RA	Goto ADJ Seen RA		Cancel T1
RX CSNP Set	Goto ADJ Seen CSNP	Cancel T1	
RX IIH w no Restart TLV	Cancel T1 (Point-to-Point only)		
ADJ UP	Start T1 Send local LSPs with overload bit set		
T1 expires	Send IIH/RR and SA Restart T1	Send IIH/RR and SA Restart T1	Send IIH/RR and SA Restart T1
T1 expires nth time	Send IIH/SA	Send IIH/SA	Send IIH/SA
T2 expires	Clear overload bit Send IIH normal Goto Running		
LSP DB Sync	Cancel T2 Clear overload bit Send IIH normal		

#### 4. IANA Considerations

This document defines the following IS-IS TLV that is listed in the IS-IS TLV codepoint registry:

Type	Description	IIH	LSP	SNP
211	Restart TLV	y	n	n



## 5. Security Considerations

Any new security issues raised by the procedures in this document depend upon the ability of an attacker to inject a false but apparently valid IIH, the ease/difficulty of which has not been altered.

If the RR bit is set in a false IIH, neighbors who receive such an IIH will continue to maintain an existing adjacency in the "UP" state and may (re)send a complete set of CSNPs. While the latter action is wasteful, neither action causes any disruption in correct protocol operation.

If the RA bit is set in a false IIH, a (re)starting router that receives such an IIH may falsely believe that there is a neighbor on the corresponding interface that supports the procedures described in this document. In the absence of receipt of a complete set of CSNPs on that interface, this could delay the completion of (re)start procedures by requiring the timer T1 to time out the locally defined maximum number of retries. This behavior is the same as would occur on a LAN where none of the (re)starting router's neighbors support the procedures in this document and is covered in Sections [2.3.1](#) and [2.3.2](#).

If an SA bit is set in a false IIH, this could cause suppression of the advertisement of an IS neighbor, which could either continue for an indefinite period or occur intermittently with the result being a possible loss of reachability to some destinations in the network and/or increased frequency of LSP flooding and SPF calculation.

The possibility of IS-IS PDU spoofing can be reduced by the use of authentication as described in [[RFC1195](#)] and [[ISO10589](#)], and especially the use of cryptographic authentication as described in [[RFC5304](#)] and [[RFC5310](#)].

## 6. Manageability Considerations

These extensions that have been designed, developed, and deployed for many years do not have any new impact on management and operation of the IS-IS protocol via this standardization process.

## 7. Acknowledgements

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## **Appendix A.** Summary of Changes from [RFC 5306](#)

This document extends [RFC 5306](#) by introducing support for signalling the neighbors of a restarting router that a planned restart is about to occur. This allows the neighbors to be aware of the state of the restarting router so that appropriate action may be taken if other topology changes occur while the planned restart is in progress. Since the forwarding plane of the restarting router is maintained based upon the pre-restart state of the network, additional topology changes introduce the possibility that traffic may be lost if paths via the restarting router continue to be used while the restart is in progress.

In support of this new functionality two new flags have been introduced:

- PR - Restart is planned
- PA - Planned restart acknowledgement

No changes to the post restart exchange between the restarting router and its neighbors have been introduced.

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