Mobile Ad hoc Networking (MANET)

Internet-Draft

Updates: <u>7182</u> (if approved)

Intended status: Standards Tra

Intended status: Standards Track

Expires: March 8, 2015

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Identity-Based Signatures for MANET Routing Protocols draft-ietf-manet-ibs-03

Abstract

This document updates RFC7182, which specifies a framework for, and specific examples of, integrity check values (ICVs) for packets and messages using the generalized packet/message format specified in RFC5444. It does so by defining an additional cryptographic function that allows the creation of an ICV that is an identity-based signature, defined according to the ECCSI (Elliptic Curve-Based Certificateless Signatures for Identity-Based Encryption) algorithm specified in RFC6507.

Status of this Memo

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1. Introduction

[RFC7182] defines ICV (integrity check value) TLVs for use in packets and messages that use the generalized MANET packet/message format defined in [RFC5444]. This specification extends the TLV definitions therein by defining two new cryptographic function code points from within the registries set up by [RFC7182]. This allows the use of an identity-based signature (IBS) as an ICV. An IBS has an additional property that is not shared by all of the previously specified ICVs, it not only indicates that the protected packet or message is valid, but also verifies the originator of the packet/message.

This specification assumes that each router (i.e., each originator of [RFC5444] format packets/messages) has an identity that may be tied to the packet or message. The router may have more than one identity, but will only use one for each ICV TLV. The cryptographic strength of the IBS is not dependent on the choice of identity.

Two options for the choice of identity are supported (as reflected by the two code points allocated). In the first the identity can be any octet sequence (up to 255 octets) included in the ICV TLV. In the second, the octet sequence is preceded by an address, either the IP source address for a packet TLV, or the message originator address for a message or address block TLV. In particular, the second option allows just the address to be used as an identity.

Identity-based signatures allow identifying the originator of information in a packet or message. They thus allow additional security functions, such as revocation of an identity, and removing all information with a specific originator, if this is recorded - as it is for OLSRv2 [RFC7181], an expected user of this specification. When applied to messages (rather than packets) this can significantly reduce the damage that a compromised router can inflict on the network.

Identity-based signatures are based on forms of asymmetric (public key) cryptography - identity-based encryption (IBE). Compared to symmetric cryptographic methods (such as HMAC and AES), IBE and IBS methods avoid requiring a shared secret key that results in a single point of failure vulnerability. Compared to more widely used asymmetric (public key) cryptographic methods (such as RSA and ECDSA), IBE and IBS methods have a major advantage, and a major disadvantage.

The advantage referred to is that each router can be configured once (for its key lifetime) by a trusted authority, independently of all other routers. Thus a router can connect to the authority (typically in a secure environment) to receive a private key, or can have a

private key delivered securely (out of band) from the authority. During normal operation of the MANET, there is no need for the trusted authority to be connected to the MANET, or even to still exist. Additional routers can be authorized, with no reference to previously authorized routers (the trusted authority must still exist in this case). A router's public key is its identity, which when tied to a packet or message (as is the case when using an address as, or as part of, the identity) means that there is no need for public key certificates or a certificate authority, and a router need not retain key material for any other routers.

The disadvantage referred to is that the trusted authority has complete authority, even more so than a conventional certificate authority. Routers cannot generate their own private keys, only the trusted authority can do that. Through the master secret held by the trusted authority, it could impersonate any router (existing or not). When used for identity-based encryption (not part of this specification) the trusted authority can decrypt anything. However, note that the shared secret key options described in [RFC7182] also have this limitation.

There are alternative mathematical realizations of identity-based signatures. This specification uses one that has been previously published as [RFC6507], known as ECCSI (Elliptic Curve-Based Certificateless Signatures for Identity-Based Encryption). In common with other identity-based encryption/signature approaches, it is based on the use of elliptic curves. Unlike some, it does not use "pairings" (bilinear maps from a product of two elliptic curve groups to another group). It thus may be easier to implement, and more efficient, than some alternatives, although with a greater signature size than some. This specification allows the use of any elliptic curve that may be used by [RFC6507].

The computational load imposed by ECCSI (and, perhaps more so, other IBS methods) is not trivial, though depending significantly on the quality of implementation of the required elliptic curve and other mathematical functions. For a security level of 128 bits, the ICV data length is 129 octets, which is longer than for alternative ICVs specified in [RFC7182] (e.g., 32 octets for the similar strength HMAC-SHA-256). The signature format used could have been slightly shortened (to 97 octets) by using a compressed representation of an elliptic curve point, however at the expense of some additional work when verifying a signature, and loss of direct compatibility with [RFC6507], and implementations thereof.

The trusted authority is referred to in [RFC6507] as the KMS (Key Management Service). That term will be used in the rest of this specification.

2. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

Additionally, this document uses the terminology of $[\underline{RFC5444}]$, $[\underline{RFC6507}]$, and $[\underline{RFC7182}]$.

3. Applicability Statement

This specification adds an additional option to the framework specified in [RFC7182] for use by [RFC5444] formatted packets and messages. It is applicable as described in [RFC7182], and subject to the additional comments in Section 6, particularly regarding the role of the trusted authority (KMS).

Specific examples of protocols for which this specification is suitable are NHDP [RFC6130] and OLSRv2 [RFC7181].

4. Specification

4.1. Cryptographic Function

This specification defines a cryptographic function named ECCSI that is implemented as specified as the "sign" function in <u>Section 5.2.1</u> of [RFC6507]. To use that specification:

- o The ICV is not calculated as cryptographic-function(hash-function(content)) as defined in [RFC7182], but (like the HMAC ICVs defined in [RFC7182]) uses the hash function within the cryptographic function. The option "none" is not permitted for hash-function, and the hash function must have a known fixed length of N octets, as specified in Section 4.2.
- o M in [RFC6507] is "content" as specified in in [RFC7182].
- o ID, used in $[\underline{RFC6507}]$, is as specified in $\underline{Section 4.3}$.
- o KPAK, SSK and PVT, used in [RFC6507], are as specified in Sections 4.2 and 5.1.1 of [RFC6507], provided by the KMS.

The length of the signature is 4N+1 octets, as specified in $[\mbox{RFC6507}]$, whose affine coordinate format (including an octet valued 0x04 to identify this) is used unchanged.

Verification of the ICV is not implemented by the receiver recalculating the ICV and comparing with the received ICV, as it is necessarily incapable of doing so. Instead the receiver evaluates the "verify" function described in Section 5.2.2 of [RFC6507], which may pass or fail.

To use that function M, KPAK, SSK and PVT are as specified above, while ID is deduced from the received packet or message, as specified in <u>Section 4.3</u>, using the <key-id> element in the <ICV-value>. This element need not match that used by the receiver, and thus when using this cryptographic function, multiple ICV TLVs differing only in their <key-id>, or in the choice of cryptographic function from the two defined in this specification, SHOULD NOT be used unless routers are administratively configured to recognize which to verify.

Routers MAY be administratively configured to reject a packet or message ICV TLV using ECCSI based on part or all of <key-id>; for example if this encodes a time after which this identity is no longer valid.

4.2. ECCSI parameters

Section 4.1 of [RFC6507] specifies parameters n, N, p, E, B, G, and q. The first of these, n, is specified as "A security parameter; the size in bits of the prime p over which elliptic curve cryptography is to be performed." For typical security levels (e.g., 128, 192 and 256 bits), n must be at least twice the required bits of security, see Section 5.6.1 of [NIST-SP-800-57].

Selection of an elliptic curve, and all related parameters, MUST be by administrative means, and known to all routers. This specification follows [RFC6507] with a RECOMMENDED selection to follow Appendix D.1.2 of [NIST-FIPS-186-4]. (Note that n in that document is q in [RFC6507].)

The parameter that is required by this specification is N, which is defined as Ceiling(n/8). The hash function used must create an output of size N octets. In particular for 128 bit security, and hence n=256, N=32, and the RECOMMENDED hash function is SHA-256. The signature (i.e. <ICV-data>) length is 4N+1 octets, i.e., 129 octets for N=32.

Note: [RFC6507] actually refers to the predecessor to [NIST-FIPS-186-4], but the latest version is specified here; there are no significant differences in this regard.

4.3. Identity

There are two options for the identity ID used by [RFC6507], which are indicated by there being two code points allocated for this cryptographic function, see Section 5.

- o For the cryptographic function ECCSI, ID is the element <key-id> defined in <u>Section 12.1 of [RFC7182]</u>. This MUST NOT be empty.
- o For the cryptographic function ECCSI-ADDR, ID is the concatenation of an address (in network byte order) and the element <key-id> defined in Section 12.1 of RFC7182, where the latter MAY be empty.
 - * For a packet TLV this address is the IP source address of the IP datagram in which this packet is included.
 - * For a message TLV or an address block TLV this address is the message originator address (the element <msg-orig-addr> defined in [RFC5444]) if that address is present, if not present and the message is known to have travelled only one hop, then the IP source address of the IP datagram in which this message is included is used, otherwise no address is defined and the message MUST be rejected. (Note that HELLO messages specified in NHDP [RFC6130] and used in OLSRv2 [RFC7181] always only travel one hop, and hence their IP source address SHOULD be used if no originator address is present.)

Note that this identity is formatted by [RFC6507], and thus does not need a length field incorporated into it by this specification.

5. IANA Considerations

IANA has, in accordance with [RFC7182], defined a registry "Cryptographic Functions" under "Mobile Ad Hoc NETwork Parameters". IANA is requested to make two new allocations from this registry, and modify the unassigned range, as indicated.

Value Algorithm Description Reference +	++			++
specification specification 8 ECCSI-ADDR ECCSI [RFC6507] with an This address (source or specification originator) joined to identity	Value	Algorithm	Description	Reference
address (source or specification originator) joined to identity	7	ECCSI	ECCSI [RFC6507]	1=
		ECCSI-ADDR	address (source or originator) joined to identity	1

Table 1: Cryptographic Function Registry

6. Security Considerations

This specification extends the security framework for MANET routing protocols specified in [RFC7182] by the addition of an additional cryptographic function, in two forms according to how identity is specified.

This cryptographic function implements a form of identity-based signature (IBS), a stronger form of integrity check value (ICV) that verifies not just that the received packet or message is valid but that the packet or message originated at a router that was assigned a private key for the specified identity.

For a message the identity is, and for a packet it is recommended that it is, either the originator address of the router (i.e., an address unique to that router), or the originator address with additional information appended. The use of that additional information is outside the scope of this specification, a typical use may be to indicate an expiry time for signatures created using that identity.

In common with other forms of IBS, a feature of the form of IBS (known as ECCSI) used in this specification is that it requires a trusted authority (KMS) that issues all private keys, and has complete cryptographic information about all possible private keys. However to set against that, the solution is scalable, as all routers can be independently keyed, and does not need the KMS in the network. If no future keys will be required, then the KMS's master secret can be destroyed. As routers are individually keyed, key revocation (by blacklist and time expiry of keys) is possible, but is beyond the scope of this specification.

ECCSI is based on elliptic curve mathematics. This specification

follows [RFC6507] in its recommendation of elliptic curves, but any suitable (prime power) elliptic curve may be used; this must be administratively specified. Implementation of this specification will require an available implementation of suitable mathematical functions. Unlike some other forms of IBS, ECCSI requires only basic elliptic curve operations, it does not require "pairings" (bilinear functions of a product of two elliptic curve groups). This increases the available range of suitable mathematical libraries.

Acknowledgments

The author would like to thank his colleagues who have been involved in identity-based security for ad hoc networks, including (in alphabetical order) Alan Cullen, Peter Smith and Bill Williams. He would also like to thank Benjamin Smith (INRIA/Ecole Polytechnique) for independently recreating the signature and other values in Appendix A to ensure their correctness, and Thomas Clausen (Ecole Polytechnique) for additional comments.

8. References

8.1. Normative References

- [RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", <u>BCP 14</u>, <u>RFC 2119</u>, March 1997.
- [RFC5444] Clausen, T., Dearlove, C., Dean, J., and C. Adjih,
 "Generalized Mobile Ad Hoc Network (MANET) Packet/Message
 Format", RFC 5444, February 2009.
- [RFC6507] Groves, M., "Elliptic Curve-Based Certificateless Signatures for Identity-Based Encryption (ECCSI)", RFC 6507, February 2012.
- [RFC7182] Herberg, U., Clausen, T., and C. Dearlove, "Integrity Check Value and Timestamp TLV Definitions for Mobile Ad Hoc Networks (MANETs)", RFC 7182, April 2014.

8.2. Informative References

[NIST-FIPS-186-4]

National Institute of Standards and Technology, "Digital Signature Standard (DSS)", FIPS 186-4, July 2013.

[NIST-SP-800-57]

National Institute of Standards and Technology,

- "Recommendation for Key Management Part 1: General (Revision 3)", SP 800-57, Part 1, Revision 3, July 2012.
- [RFC5497] Clausen, T. and C. Dearlove, "Representing Multi-Value Time in Mobile Ad Hoc Networks (MANETs)", RFC 5497, March 2009.
- [RFC6130] Clausen, T., Dearlove, C., and J. Dean, "Mobile Ad Hoc Network (MANET) Neighborhood Discovery Protocol (NHDP)", RFC 6130, April 2011.
- [RFC7181] Clausen, T., Dearlove, C., Jacquet, P., and U. Herberg, "The Optimized Link State Routing Protocol Version 2", RFC 7181, April 2014.

Appendix A. Example

Appendix C of [RFC6130] contains this example of a HELLO message. (Note that normally, a TIMESTAMP ICV would also be added before the ICV TLV, but for simplicity that step has been omitted here.)

```
0
                   2
                             3
          1
\begin{smallmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 0 & 1 \\ \end{smallmatrix}
MF=7 | MAL=3 | Message Length = 45
| Hop Limit = 1 | Hop Count = 0 |
                Message Sequence Number
| Message TLV Block Length = 8 | VALIDITY_TIME | MTLVF = 16
| Value Len = 1 | Value (Time) | INTERVAL_TIME | MTLVF = 16 |
| Value Len = 1 | Value (Time) | Num Addrs = 5 | ABF = 128 |
| Head Len = 3 |
                  Head
Mid 0
         Mid 1
                 Mid 2
                      Mid 3
| Address TLV Block Length = 14 | LOCAL_IF |
| ATLVF = 80 | Index = 0 | Value Len = 1 |
                        THIS_IF
| LINK_STATUS | ATLV = 52 | Strt Indx = 1 | Stop Indx = 4 |
HEARD |
                 HEARD |
| Value Len = 4 |
                        SYMMETRIC |
LOST
+-+-+-+-+-+-+
```

In order to provide an example of an ECCSI ICV Message TLV that may be added to this message, the fields shown need to all have numerical values, both by inserting defined numerical values (e.g., 0 for HELLO) and by selecting example values where needed. The latter consists of:

- o The message sequence number will be zero.
- o The five addresses will be 192.0.2.1 to 192.0.2.5.
- o The message validity time will be 6 seconds, and the message interval time will be 2 seconds, each encoded with a constant value C = 1/1024 seconds, as described in [RFC5497], as referenced from [RFC6130].

In addition, when calculating an ICV, the hop count and hop limit are both set to zero. This results in the message:

```
0
        2
            3
    1
\begin{smallmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 0 & 1 \\ \end{smallmatrix}
|0 0 0 0 0 0 0 0 0 0 1 1 1 0 0 1 1 0 0 0 0 0 0 0 0 0 0 1 0 1 1 0 1 1
|0 0 0 0 0 1 0 1 0 0 0 0 0 0 0 0 0 0 1 1 1 0 0 0 0 0 0 1 0 0
10 0 0 0 0 0 0 0 0
+-+-+-+-+-+-+
```

Or in hexadecimal form

M := 0x 0073002D 00000000 00080110 01640010 01580580 03C00002 01020304 05000E02 50000100 03340104 04020201 00

The ICV TLV that will be added will have cryptographic function ECCSI-ADDR, and hash function SHA-256. This message has no originator address, but it travels a single hop, and its IP source address can be used. This will be assumed to be 192.0.2.0, with an empty <key-id>, thus the sender's identity will be, in hexadecimal form:

ID := 0x C0000200

Parameters for [RFC6507] will thus be n = 256, N = 32. The same parameters and master key will be used as in Appendix A of [RFC6507], i.e., the elliptic curve P-256, with parameters:

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p	:= 0x			00000000 FFFFFFF	
В	:= 0x	0.100000		B3EBBD55 3BCE3C3E	
q	:= 0x			FFFFFFF F3B9CAC2	
G	:= 0x	77037D81 4FE342E2	2DEB33A0 FE1A7F9B	F8BCE6E5 F4A13945 8EE7EB4A CBB64068	D898C296 7C0F9E16
KSAK	:= 0x	12345;			
KPAK	:= 0x	7A72686D DBDD3755	4522D4C8 1AFD263B	28D2838A 273FB644 5DFD617F 66DCE7D4	2AEBFA93 3960C65A

The remaining steps to creating a private key for ID use the same "random" value v as Appendix A of [RFC6507] and are:

:= 0x 23456

PVT := 0x 04

> 758A1427 79BE89E8 29E71984 CB40EF75 8CC4AD77 5FC5B9A3 E1C8ED52 F6FA36D9 A79D2476 92F4EDA3 A6BDAB77 D6AA6474 A464AE49 34663C52 65BA7018 BA091F79

HS := hash(0x 04

> 6B17D1F2 E12C4247 F8BCE6E5 63A440F2 77037D81 2DEB33A0 F4A13945 D898C296 4FE342E2 FE1A7F9B 8EE7EB4A 7C0F9E16 2BCE3357 6B315ECE CBB64068 37BF51F5

50D4670B DE75244F 28D2838A 0D25558A 7A72686D 4522D4C8 273FB644 2AEBFA93 DBDD3755 1AFD263B 5DFD617F 3960C65A 8C298850 FF99F203 66DCE7D4 367217F4

C0000200 04

758A1427 79BE89E8 29E71984 CB40EF75 8CC4AD77 5FC5B9A3 E1C8ED52 F6FA36D9 A79D2476 92F4EDA3 A6BDAB77 D6AA6474 A464AE49 34663C52 65BA7018 BA091F79)

= 0x F64FFD76 D2EC3E87 BA670866 C0832B80 B740C2BA 016034C8 1A6F5E5B 5F9AD8F3

The remaining steps to creating a signature for M use the same "random" value j as Appendix A of [RFC6507] and are:

```
:= 0x 34567
j
J
        := 0x 04
                269D4C8F DEB66A74 E4EF8C0D 5DCC597D
                DFE6029C 2AFFC493 6008CD2C C1045D81
                6DDA6A13 10F4B067 BD5DABDA D741B7CE
                F36457E1 96B1BFA9 7FD5F8FB B3926ADB
         := 0x 269D4C8F DEB66A74 E4EF8C0D 5DCC597D
                DFE6029C 2AFFC493 6008CD2C C1045D81
ΗE
          := hash( 0x
                   F64FFD76 D2EC3E87 BA670866 C0832B80
                   B740C2BA 016034C8 1A6F5E5B 5F9AD8F3
                   269D4C8F DEB66A74 E4EF8C0D 5DCC597D
                   DFE6029C 2AFFC493 6008CD2C C1045D81
                   0073002D 00000000 00080110 01640010
                   01580580 03C00002 01020304 05000E02
                   50000100 03340104 04020201 00
                                                 )
          = 0x FE236B30 CF72E060 28E229ED 5751D796
                91DED33C 24D2F661 28EA0804 30D8A832
s'
         := 0x C8C739D5 FB3EFB75 221CB818 8CAAB86A
                2E2669CF 209EA622 7D7072BA A83C2509
S
         := 0x C8C739D5 FB3EFB75 221CB818 8CAAB86A
                2E2669CF 209EA622 7D7072BA A83C2509
Signature := 0x 269D4C8F DEB66A74 E4EF8C0D 5DCC597D
                DFE6029C 2AFFC493 6008CD2C C1045D81
                C8C739D5 FB3EFB75 221CB818 8CAAB86A
                2E2669CF 209EA622 7D7072BA A83C2509
                04
                758A1427 79BE89E8 29E71984 CB40EF75
                8CC4AD77 5FC5B9A3 E1C8ED52 F6FA36D9
                A79D2476 92F4EDA3 A6BDAB77 D6AA6474
                A464AE49 34663C52 65BA7018 BA091F79
```

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