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Inter-Area P2MP Segmented LSPs

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Abstract

This document describes procedures for building inter-area point-to-multipoint (P2MP) segmented service LSPs by partitioning such LSPs into intra-area segments and using BGP as the inter-area routing and label distribution protocol. Within each IGP area the intra-area segments are either carried over intra-area P2MP LSPs, using P2MP LSP hierarchy, or instantiated using ingress replication. The intra-area P2MP LSPs may be signaled using P2MP RSVP-TE or P2MP mLDP. If ingress replication is used within an IGP area, then MP2P LDP LSPs or P2P RSVP-TE LSPs may be used in the IGP area. The applications/services that use such inter-area service LSPs may be BGP MVPN, VPLS multicast, or global table multicast over MPLS.

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1. Specification of requirements

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

2. Introduction

This document describes procedures for building inter-area point-to-multipoint (P2MP) segmented service LSPs by partitioning such LSPs into intra-area segments and using BGP as the inter-area routing and label distribution protocol. Within each IGP area the intra-area segments are either carried over intra-area P2MP LSPs, potentially using P2MP LSP hierarchy, or instantiated using ingress replication. The intra-area P2MP LSPs may be signaled using P2MP RSVP-TE or P2MP mLDP. If ingress replication is used in an IGP area then MP2P LDP or P2P RSVP-TE LSPs may be used within the IGP area. The applications/services that use such inter-area service LSPs may be

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BGP MVPN, VPLS multicast, or global table multicast over MPLS.

The primary use case of such segmented P2MP service LSPs is when the PEs are in different areas but in the same AS and thousands or more of PEs require P2MP connectivity. For instance this may be the case when MPLS is pushed further to the metro edge and the metros are in different IGP areas. This may also be the case when a Service Provider's network comprises multiple IGP areas in a single Autonomous System, with a large number of PEs. Seamless MPLS is the industry term to address this case [SEAMLESS-MPLS]. Thus one of the applicabilities of this document is that it describes the multicast procedures for seamless MPLS.

It is to be noted that [RFC6514], [VPLS-P2MP] already specify procedures for building segmented inter-AS P2MP service LSPs. This document complements those procedures, as it extends the segmented P2MP LSP model such that it is applicable to inter-area P2MP service LSPs as well. Infact an inter-AS deployment could use inter-AS segmented P2MP LSPs as specified in [RFC6514, VPLS-P2MP] where each intra-AS segment is constructed using inter-area segmented P2MP LSPs as specified in this document.

3. General Assumptions and Terminology

This document allows ABRs, acting as Route Reflectors, to follow the procedures specified in [SEAMLESS-MPLS] when handling BGP Next-Hop of the routes to the PEs. Specifically, when reflecting such routes from the non-backbone areas into the backbone area, the ABRS MUST set BGP Next-Hop to their own loopback addresses (next-hop-self), while when reflecting such routes from the backbone area into the non-backbone areas, the ABRS SHOULD NOT change the BGP Next-Hop addresses (next-hop-unchanged).

While this document allows ABRs to follow the procedures specified in [SEAMLESS-MPLS], procedures specified in this document are applicable even when ABRs do not follow the procedures specified in [SEAMLESS-MPLS].

One possible way to support the global table multicast service is by relying on the MVPN procedures, as specified in [RFC6514], in which case the MVPN procedures specified in this document would be used to support the global table multicast service. This document specifies an alternative approach to support the global table multicast service. This document refers to this approach as "global table multicast" (although this by no means imply that this alternative approach is the only way to support the global table multicast service).

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This document assumes that in the context of global table multicast ABRs do not carry routes to the destinations external to their own AS. Furthermore, in the context of global table multicast this document assumes that an ASBR, when re-advertising into IBGP routes received from an external speaker (received via EBGP) may not change BGP Next-Hop to self.

Within an AS a P2MP service LSP is partitioned into 3 segments: ingress area segment, backbone area segment, and egress area segment. Within each area a segment is carried over an intra-area P2MP LSP or instantiated using ingress replication.

When intra-area P2MP LSPs are used to instantiate the intra-area segments there could be either 1:1 or n:1 mapping between intra-area segments of the inter-area P2MP service LSP and a given intra-area P2MP LSP. The latter is realized using P2MP LSP hierarchy with upstream-assigned labels [RFC5331]. For simplicity of presentation we assume that P2MP LSP hierarchy is used even with 1:1 mapping, in which case the upstream-assigned label SHOULD be an Implicit NULL.

When intra-area segments of the inter-area P2MP service LSP are instantiated using ingress replication, then multiple such segments may be carried in the same P2P RSVP-TE or MP2P LDP LSP. This can be achieved using downstream-assigned labels alone.

The ingress area segment of a P2MP service LSP is rooted at a PE (or at an ASBR in the case where the P2MP service LSP spans multiple ASes). The leaves of this segment are other PEs/ASBRs and ABRs in the same area as the root PE.

The backbone area segment is rooted at an ABR that is connected to the ingress area (ingress ABR), or at an ASBR if the ASBR is present in the backbone area, or at a PE if the PE is present in the backbone area. The backbone area segment has as its leaves ABRs that are connected to the egress area(s) or PEs in the backbone area, or ASBRs in the backbone area.

The egress area segment is rooted at an ABR in the egress area (egress ABR), and has as its leaves PEs and ASBR in that egress area (the latter covers the case where the P2MP service LSP spans multiple ASes). Note that for a given P2MP service LSP there may be more than one backbone segment, each rooted at its own ingress ABR, and more than one egress area segment, each rooted at its own egress ABR.

This document uses the term "A-D routes" for "auto-discovery routes".

An implementation that supports this document MUST implement the procedures described in the following sections to support inter-area Rekhter [Page 6]

point-to-multipoint (P2MP) segmented service LSPs.

4. Inter-area P2MP Segmented Next-Hop Extended Community

This document defines a new BGP Extended Community "Inter-area P2MP Next-Hop" Extended Community. This is an IP-address specific Extended Community of an extended type and is transitive across AS boundaries [RFC4360].

A PE or an ABR or an ASBR constructs the Inter-area P2MP Segmented Next-Hop Extended Community as follows:

- The Global Administrator field MUST be set to an IP address of the PE or ASBR or ABR that originates or advertises the route, which carries the P2MP Next-Hop Extended Community. For example this address may be the loopback address or the PE, ASBR or ABR that advertises the route.
- The Local Administrator field MUST be set to 0.

The detailed usage of this Extended Community is described in the following sections.

5. Discovering P2MP FEC of Inter-Area P2MP Service LSP

Each inter-area P2MP service LSP has associated with it P2MP FEC. The egress PEs need to learn this P2MP FEC in order to initiate the creation of the egress area segment of the P2MP inter-area service LSP.

The P2MP FEC of the inter-area P2MP LSP is learned by the egress PEs either by configuration, or based on the application-specific procedures (e.g., MVPN-specific procedures, VPLS-specific procedures).

5.1. BGP MVPN

Egress PEs and/or ASBRs discover the P2MP FEC of the service LSPs used by BGP MVPN using the I-PMSI or S-PMSI A-D routes that are originated by the ingress PEs or ASBRs following the procedures of [RFC6514], along with modifications as described in this document. The NLRI of such routes encodes the P2MP FEC. The procedures in this document require that at least one ABR in a given IGP area act as Route Reflector for MVPN A-D routes.

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The "Leaf Information Required" flag MUST be set in the P-Tunnel attribute carried in such routes, when originated by the ingress PEs or ASBRs, except for the case where (a) as a matter of policy (provisioned on the ingress PEs or ASBRs) there is no aggregation of ingress area segments of the service LSPs, and (b) mLDP is used as the protocol to establish intra-area transport LSPs in the ingress area. Before any Leaf A-D route is advertised by a PE or ABR in the same area, as described in the following sections, an I-PMSI/S-PMSI A-D route is advertised either with an explicit Tunnel Type and Tunnel Identifier in the PMSI Tunnel Attribute, if the Tunnel Identifier has already been assigned, or with a special Tunnel Type of "No tunnel information present" otherwise.

When the I-PMSI/S-PMSI A-D routes are re-advertised by an ingress ABR, the "Leaf Information Required" flag MUST be set in the P-Tunnel attribute present in the routes, except for the case where (a) as a matter of policy (provisioned on the ingress ABR) there is no aggregation of backbone area segments of the service LSPs, and (b) mLDP is used as the protocol to establish intra-area transport LSPs in the backbone area. Likewise, when the I-PMSI/S-PMSI A-D routes are re-advertised by an egress ABR, the "Leaf Information Required" flag MUST be set in the P-Tunnel attribute present in the routes, except for the case where (a) as a matter of policy (provisioned on the egress ABR) there is no aggregation of egress area segments of the service LSPs, and (b) mLDP is used as the protocol to establish intra-area transport LSPs in the egress area.

Note that the procedures in the above paragraph apply when intra-area segments are realized by either intra-area P2MP LSPs or by ingress replication.

When BGP MVPN I-PMSI or S-PMSI A-D routes are advertised or propagated to signal Inter-area P2MP service LSPs, to indicate that these LSPs should be segmented using the procedures specified in this document, these routes MUST carry the Inter-area P2MP Segmented Next-Hop Extended Community. This Extended Community MUST be included in the I-PMSI/S-PMSI A-D route by the PE or ASBR that originates such a route, and the Global Administrator field MUST be set to the advertising PE or ASBR's IP address. This Extended Community MUST also be included by ABRs as they re-advertise such routes. An ABR MUST set the Global Administrator field of the P2MP Segmented Next-Hop Extended Community to its own IP address. Presence of this community in the I-PMSI/S-PMSI A-D routes indicates to ABRs and PES/ASBRs that they have to follow the procedures in this document when these procedures differ from those in [RFC6514].

To avoid requiring ABRs to participate in the propagation of Cmulticast routes, this document requires ABRs NOT to modify BGP Next Rekhter [Page 8]

Hop when re-advertising Inter-AS I-PMSI A-D routes. For consistency this document requires ABRs to NOT modify BGP Next-Hop when readvertising both Intra-AS and Inter-AS I-PMSI/S-PMSI A-D routes. The egress PEs may advertise the C-multicast routes to RRs that are different than the ABRs. However ABRs still can be configured to be the Route Reflectors for C-multicast routes, in which case they will participate in the propagation of C-multicast routes.

5.2. BGP VPLS or LDP VPLS with BGP auto-discovery

Egress PEs discover the P2MP FEC of the service LSPs used by VPLS, using the VPLS A-D routes that are originated by the ingress PEs [RFC4761, RFC6074], or VPLS S-PMSI A-D routes that are originated by the ingress PEs [VPLS-P2MP]. The NLRI of such routes encodes the P2MP FEC.

The "Leaf Information Required" flag MUST be set in the P-Tunnel attribute carried in such routes, when originated by the ingress PEs or ASBRs, except for the case where (a) as a matter of policy (provisioned on the ingress PEs or ASBRs) there is no aggregation of ingress area segments of the service LSPs, and (b) mLDP is used as the protocol to establish intra-area transport LSPs in the ingress area. Before any Leaf A-D route is advertised by a PE or ABR in the same area, as described in the following sections, an VPLS/S-PMSI A-D route is advertised either with an explicit Tunnel Type and Tunnel Identifier in the PMSI Tunnel Attribute, if the Tunnel Identifier has already been assigned, or with a special Tunnel Type of "No tunnel information present" otherwise.

When the VPLS/S-PMSI A-D routes are re-advertised by an ingress ABR, the "Leaf Information Required" flag MUST be set in the P-Tunnel attribute present in the routes, except for the case where (a) as a matter of policy (provisioned on the ingress ABR) there is no aggregation of backbone area segments of the service LSPs, and (b) mLDP is used as the protocol to establish intra-area transport LSPs in the backbone area. Likewise, when the VPLS/S-PMSI A-D routes are re-advertised by an egress ABR, the "Leaf Information Required" flag MUST be set in the P-Tunnel attribute present in the routes, except for the case where (a) as a matter of policy (provisioned on the egress ABR) there is no aggregation of egress area segments of the service LSPs, and (b) mLDP is used as the protocol to establish intra-area transport LSPs in the egress area.

When VPLS A-D routes or S-PMSI A-D routes are advertised or propagated to signal Inter-area P2MP service LSPs, to indicate that these LSPs should be segmented using the procedures specified in this document, these routes MUST carry the Inter-area P2MP Segmented Next-

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Hop Extended Community. This Extended Community MUST be included in the A-D route by the PE or ASBR that originates such a route and the Global Administrator field MUST be set to the advertising PE or ASBR's IP address. This Extended Community MUST also be included by ABRs as they re-advertise such routes. An ABR MUST set the Global Administrator field of the P2MP Segmented Next-Hop Extended Community to its own IP address. Presence of this community in the I-PMSI/S-PMSI A-D routes indicates to ABRs and PEs/ASBRs that they have to follow the procedures in this document when these procedures differ from those in [VPLS-P2MP].

Note that the procedures in the above paragraph apply when intra-area segments are realized by either intra-area P2MP LSPs or by ingress replication.

The procedures in this document require that at least one ABR in a given area act as Route Reflector for VPLS A-D routes. These ABRs/RRs MUST NOT modify BGP Next Hop when re-advertising these A-D routes.

5.3. Global Table Multicast over MPLS

This section describes how the egress PEs discover the P2MP FEC when the application is global table multicast over an MPLS-capable infrastructure. In the rest of the document we will refer to this application as "global table multicast".

When PIM-SM is used for non-bidirectional ASM ("Any Source Multicast") group addresses, this document refers to this as "PIM-SM in ASM mode".

In the case where global table multicast uses PIM-SM in ASM mode the following assumes that an inter-area P2MP service LSP could be used to either carry traffic on a shared (*,G), or a source (S,G) tree.

An egress PE learns the (S/*, G) of a multicast stream as a result of receiving IGMP or PIM messages on one of its IP multicast interfaces. This (S/*, G) forms the P2MP FEC of the inter-area P2MP service LSP. For each such P2MP FEC there MAY exist a distinct inter-area P2MP service LSP, or multiple such FECs MAY be carried over a single P2MP service LSP using a wildcard (*, *) S-PMSI [RFC6625].

Note that this document does not require the use of (*, G) Inter-area P2MP service LSPs when global table multicast uses PIM-SM in ASM mode. In fact, PIM-SM in ASM mode may be supported entirely by using only (S, G) inter-area P2MP service LSPs.

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6. Egress PE Procedures

This section describes egress PE procedures for constructing segmented inter-area P2MP LSP. The procedures in this section apply irrespective of whether the egress PE is in a leaf IGP area, or the backbone area, or even in the same IGP area as the ingress PE/ASBR.

An egress PE applies procedures specified in this section to MVPN I-PMSI or S-PMSI A-D routes only if these routes carry the Inter-area P2MP Segmented Next-Hop Extended Community. An egress PE applies procedures specified in this section to VPLS A-D routes or VPLS S-PMSI A-D routes only if these routes carry the Inter-area P2MP Segmented Next-Hop Extended Community.

In order to support global table multicast an egress PE MUST autoconfigure an import AS-specific Route Target Extended Community ([RFC4360]) with the Global Administrator field set to the AS of the PE and the Local Administrator field set to 0.

Once an egress PE discovers the P2MP FEC of an inter-area segmented P2MP service LSP, it MUST propagate this P2MP FEC in BGP in order to construct the segmented inter-area P2MP service LSP. This propagation uses BGP Leaf A-D routes.

6.1. Determining the Upstream ABR/PE/ASBR (Upstream Node)

An egress PE discovers the P2MP FEC of an inter-area P2MP Segmented Service LSP as described in section 5 ("Discovering P2MP FEC of the Inter-Area P2MP Service LSP"). Once the egress PE discovers this P2MP FEC, it MUST determine the upstream node to reach such a FEC. If the egress PE is in the egress area, and the ingress PE is not in the that egress area, then this upstream node would be an egress ABR. If the egress PE is in the backbone area and the ingress PE is not in the backbone area, then this upstream node would be an ingress ABR. If the egress PE is in the same area as the ingress PE, then this upstream node would be the ingress PE.

6.1.1. Upstream Node for MVPN or VPLS

If the application is MVPN or VPLS, then the upstream node's IP address is the IP address determined from the Global Administrator field of the Inter-area P2MP Segmented Next-hop Extended Community. As described in section 5 ("Discovering P2MP FEC of the Inter-Area P2MP Service LSP"), this Extended Community MUST be carried in the MVPN or VPLS A-D route from which the P2MP FEC of the inter-area P2MP Segmented Service LSP is determined.

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6.1.2. Upstream Node for Global Table Multicast

If the application is global table multicast, then the unicast routes to multicast sources/RPs SHOULD carry the VRF Route Import Extended Community [RFC6514] where the IP address in the Global Administrator field is set to the IP address of the PE or ASBR advertising the unicast route. The Local Administrator field of this community MUST be set to 0 (note, that this is in contrast to the case of MVPN, where the Global Administrator field carries a non-zero value that identifies a particular VRF on a PE that originates VPN-IP routes). If it is not desirable to advertise the VRF Route Import Extended Community in unicast routes, then unicast routes to multicast sources/RPs MUST be advertised using the multicast SAFI i.e. SAFI 2, and such routes MUST carry the VRF Route Import Extended Community.

Further, if the application is global table multicast, then the BGP unicast routes that advertise the routes to the IP addresses of PEs/ASBRs/ABRs SHOULD carry the Inter-area P2MP Segmented Next-Hop Extended Community. The IP address in the Global Administrator field of this community MUST be set to the IP address of the PE or ASBR or ABR advertising the unicast route. The Local Administrator field of this community MUST be set to 0. If it is not desirable to advertise the P2MP Segmented Next-Hop Extended Community in BGP unicast routes, then the BGP unicast routes to the IP addresses of PEs/ASBRs/ABRs MUST be advertised using the multicast SAFI i.e. SAFI 2, and such routes MUST carry the Inter-area P2MP Segmented Next-hop Extended Community. The procedures for handling the BGP Next-Hop attribute of SAFI 2 routes are the same as those of handling regular Unicast routes and MAY follow [SEAMLESS-MPLS].

If the application is global table multicast, then in order to determine the upstream node address the egress PE first determines the ingress PE. In order to determine the ingress PE, the egress PE determines the best route to reach S/RP. The ingress PE address is the IP address determined from the Global Administrator field of the VRF Route Import Extended Community that is carried in this route. The egress PE then finds the best unicast route to reach the ingress PE. The upstream node address is the IP address determined from the Global Administrator field of the Inter-area P2MP Segmented Next-Hop Extended Community that is carried in this route.

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6.2. Originating a Leaf A-D Route

If the P2MP FEC was derived from a MVPN or VPLS A-D route then the egress PE MUST originate a Leaf A-D route if the MVPN or VPLS A-D route carries a P-Tunnel Attribute with the "Leaf Information" Required" flag set.

If the P2MP FEC was derived from a global table multicast (S/*, G), and the upstream node's address is not the same as the egress PE, then the egress PE MUST originate a Leaf A-D route.

6.2.1. Leaf A-D Route for MVPN and VPLS

If the P2MP FEC was derived from MVPN or VPLS A-D routes then the Route Key field of the Leaf A-D route contains the NLRI of the A-D route from which the P2MP FEC was derived. This follows procedures for constructing Leaf A-D routes described in [RFC6514, VPLS-P2MP].

6.2.2. Leaf A-D Route for Global Table Multicast

If the application is global table multicast, then the MCAST-VPN NLRI of the Leaf A-D route is constructed as follows:

The Route Key field of MCAST-VPN NLRI has the following format:

++
RD (8 octets)
Multicast Source Length (1 octet)
Multicast Source (Variable)
Multicast Group Length (1 octet)
Multicast Group (Variable)
Ingress PE's IP address
,

RD is set to 0 for (S,G) state and all 1s for (*,G) state, Multicast Source is set to S for (S,G) state or RP for (*,G) state, Multicast Group is set to G, Multicast Source Length and Multicast Group Length is set to either 4 or 16 (depending on whether S/RP and G are IPv4 or IPv6 addresses).

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The Ingress PE's IP address is determined as described in the section 6.1 ("Determining the Upstream ABR/PE/ASBR (Upstream Node)").

The Originating Router's IP address field of MCAST-VPN NLRI is set to the address of the local PE (PE that originates the route).

Thus the entire MCAST-VPN NLRI of the route has the following format:

++
Route Type = 4 (1 octet)
Length (1 octet)
RD (8 octets)
++
Multicast Source Length (1 octet)
Multicast Source (Variable)
++
Multicast Group Length (1 octet)
++
Multicast Group (Variable)
Ingress PE's IP address
1 Originating Douterla ID address
Originating Router's IP address +

Note that the encoding of MCAST-VPN NLRI for the Leaf A-D routes used for global table multicast is different from the encoding used by the Leaf A-D routes originated in response to S-PMSI or I-PMSI A-D routes. A router that receives a Leaf A-D route can distinguish between these two cases by examining the third octet of the MCAST-VPN NLRI of the route. If the value of this octet is 0x01 or 0x02, or 0x03 then this Leaf A-D route was originated in response to an S-PMSI or I-PMSI A-D route. If the value of this octet is either 0x00 or 0xff, and octets 3 through 10 contain either all 0x00, or all 0xff then this is a Leaf A-D route used for global table multicast.

When the PE deletes (S,G)/(*,G) state that was created as a result of receiving PIM or IGMP messages on one of its IP multicast interfaces, if the PE previously originated a Leaf A-D route for that state, then the PE SHOULD withdraw that route.

An Autonomous System with an IPv4 network may provide global table multicast service for customers that use IPv6, and an Autonomous System with an IPv6 network may provide global table multicast service for customers that use IPv4. Therefore the address family of Rekhter [Page 14]

the Ingress PE's IP address and Originating Router's IP address in the Leaf A-D routes used for global table multicast MUST NOT be inferred from the AFI field of the associated MP_REACH_NLRI/MP_UNREACH_NLRI attribute of these routes. The address family is determined from the length of the address (a length of 4 octets for IPv4 addresses, a length of 16 octets for IPv6 addresses).

For example if an Autonomous System with an IPv4 network is providing IPv6 multicast service to a customer, the Ingress PE's IP address and Originating Router's IP address in the Leaf A-D routes used for IPv6 global table multicast will be a four-octet IPv4 address, even though the AFI of those routes will have the value 2.

Note that the Ingress PE's IP address and the Originating Router's IP address must be either both IPv4 or both IPv6 addresses, and thus must be of the same length. Since the two variable length fields (Multicast Source and Multicast Group) in the Leaf A-D routes used for global table multicast have their own length field, from these two length fields, and the Length field of the MCAST-VPN NLRI, one can compute length of the Ingress PE's IP address and the Originating Router's IP address fields. If the computed length of these fields is neither 4 nor 16, the MP_REACH_NLRI attribute MUST be considered to be "incorrect", and MUST be handled as specified in section 7 of [RFC4760].

6.2.3. Constructing the Rest of the Leaf A-D Route

The Next Hop field of the MP_REACH_NLRI attribute of the route SHOULD be set to the same IP address as the one carried in the Originating Router's IP Address field of the route.

When Ingress Replication is used to instantiate the egress area segment then the Leaf A-D route MUST carry a downstream assigned label in the P-Tunnel Attribute where the P-Tunnel type is set to Ingress Replication. A PE MUST assign a distinct MPLS label for each Leaf A-D route originated by the PE.

To constrain distribution of this route, the originating PE constructs an IP-based Route Target community by placing the IP address of the upstream node in the Global Administrator field of the community, with the Local Administrator field of this community set to 0. The originating PE then adds this Route Target Extended Community to this Leaf A-D route. The upstream node's address is as determined in section 6.1 ("Determining the Upstream ABR/PE/ASBR (Upstream Node)").

The PE then advertises this route to the upstream node.

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6.3. PIM-SM in ASM mode for Global Table Multicast

This specification allows two options for supporting global table multicast with PIM-SM in ASM mode. The first option does not transit IP multicast shared trees over the MPLS network. The second option does transit shared trees over the MPLS network and relies on shared tree to source tree switchover.

6.3.1. Option 1

This option does not transit IP multicast shared trees over the MPLS network. Therefore, when an (egress) PE creates (*, G) state (as a result of receiving PIM or IGMP messages on one of its IP multicast interfaces), the PE does not propagate this state using Leaf A-D routes.

6.3.1.1. Originating Source Active A-D Routes

Whenever as a result of receiving PIM Register or MSDP messages an RP that is co-located with a PE discovers a new multicast source, the RP/PE SHOULD originate a BGP Source Active A-D route. Similarly whenever as a result of receiving MSDP messages a PE, that is not configured as a RP, discovers a new multicast source the PE SHOULD originate a BGP Source Active A-D route. The BGP Source Active A-D route carries a single MCAST-VPN NLRI constructed as follows:

- + The RD in this NLRI is set to 0.
- + The Multicast Source field MUST be set to S. The Multicast Source Length field is set appropriately to reflect this.
- + The Multicast Group field MUST be set to G. The Multicast Group Length field is set appropriately to reflect this.

To constrain distribution of the Source Active A-D route to the AS of the advertising RP this route SHOULD carry the NO_EXPORT Community ([RFC1997]).

Using the normal BGP procedures the Source Active A-D route is propagated to all other PEs within the AS.

Whenever the RP/PE discovers that the source is no longer active, the RP MUST withdraw the Source Active A-D route, if such a route was previously advertised by the RP.

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6.3.1.2. Receiving BGP Source Active A-D Route by PE

When as a result of receiving PIM or IGMP messages on one of its IP multicast interfaces an (egress) PE creates in its Tree Information Base (TIB) a new (*, G) entry with a non-empty outgoing interface list that contains one or more IP multicast interfaces, the PE MUST check if it has any Source Active A-D routes for that G. If there is such a route, S of that route is reachable via an MPLS interface, and the PE does not have (S, G) state in its TIB for (S, G) carried in the route, then the PE originates a Leaf A-D route carrying that (S, G), as specified in section 6.2.2 ("Leaf A-D Route for Global Table Multicast").

When an (egress) PE receives a new Source Active A-D route, the PE MUST check if its TIB contains an (*, G) entry with the same G as carried in the Source Active A-D route. If such an entry is found, S is reachable via an MPLS interface, and the PE does not have (S, G) state in its TIB for (S, G) carried in the route, then the PE originates a Leaf A-D route carrying that (S, G), as specified in section 6.2.2 ("Leaf A-D Route for Global Table Multicast").

6.3.1.3. Handling (S, G, RPTbit) state

Creation and deletion of (S, G, RPTbit) state on a PE that resulted from receiving PIM messages on one of its IP multicast interfaces does not result in any BGP actions by the PE.

6.3.2. Option 2

This option does transit IP multicast shared trees over the MPLS network. Therefore, when an (egress) PE creates (*, G) state (as a result of receiving PIM or IGMP messages on one of its IP multicast interfaces), the PE does propagate this state using Leaf A-D routes.

6.3.2.1. Originating Source Active A-D Routes

Whenever a PE creates an (S, G) state as a result of receiving Leaf A-D routes associated with the global table multicast service, if S is reachable via one of the IP multicast capable interfaces, and the PE determines that G is in the PIM-SM in ASM mode range, the PE MUST originate a BGP Source Active A-D route. The route carries a single MCAST-VPN NLRI constructed as follows:

+ The RD in this NLRI is set to 0.

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- + The Multicast Source field MUST be set to S. The Multicast Source Length field is set appropriately to reflect this.
- + The Multicast Group field MUST be set to G. The Multicast Group Length field is set appropriately to reflect this.

To constrain distribution of the Source Active A-D route to the AS of the advertising PE this route SHOULD carry the NO_EXPORT Community ([RFC1997]).

Using the normal BGP procedures the Source Active A-D route is propagated to all other PEs within the AS.

Whenever the PE deletes the (S, G) state that was previously created as a result of receiving a Leaf A-D route for (S, G), the PE that deletes the state MUST also withdraw the Source Active A-D route, if such a route was advertised when the state was created.

6.3.2.2. Receiving BGP Source Active A-D Route

Procedures for receiving BGP Source Active A-D routes are the same as with Option 1.

6.3.2.3. Pruning Sources off the Shared Tree

If after receiving a new Source Active A-D route for (S,G) a PE determines that (a) it has the (*,G) entry in its TIB, (b) the incoming interface list (iif) for that entry contains one of the IP interfaces, (c) a MPLS LSP is in the outgoing interface list (oif) for that entry, and (d) the PE does not originate a Leaf A-D route for (S,G), then the PE MUST transition the (S,G,rpt) downstream state to the Prune state. [Conceptually the PIM state machine on the PE will act "as if" it had received Prune(S,G,Rpt) from some other PE, without actually having received one.] Depending on the (S,G,rpt) state on the iifs, this may result in the PE using PIM procedures to prune S off the Shared (*,G) tree.

Transitioning the state machine to the Prune state SHOULD be done after a delay that is controlled by a timer. The value of the timer MUST be configurable. The purpose of this timer is to ensure that S is not pruned off the shared tree until all PEs have had time to receive the Source Active A-D route for (S,G).

The PE MUST keep the (S,G,rpt) downstream state machine in the Prune state for as long as (a) the outgoing interface list (oif) for (*, G) contains a MPLS LSP, and (b) the PE has at least one Source Active A-

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D route for (S,G), and (c) the PE does not originate the Leaf A-D route for (S,G). Once either of these conditions become no longer valid, the PE MUST transition the (S,G,rpt) downstream state machine to the NoInfo state.

Note that except for the scenario described in the first paragraph of this section, in all other scenarios relying solely on PIM procedures on the PE is sufficient to ensure the correct behavior when pruning sources off the shared tree.

6.3.2.4. More on handling (S, G, RPTbit) state

Creation and deletion of (S, G, RPTbit) state on a PE that resulted from receiving PIM messages on one of its IP multicast interfaces does not result in any BGP actions by the PE.

7. Egress ABR Procedures

This section describes Egress ABR Procedures for constructing segmented inter-area P2MP LSP.

7.1. Handling Leaf A-D route on Egress ABR

When an egress ABR receives a Leaf A-D route and the Route Target Extended Community carried by the route contains the IP address of this ABR, then the following procedures will be executed.

If the value of the third octet of the MCAST-VPN NLRI of the received Leaf A-D route is either 0x01, or 0x02, or 0x03, this indicates that the Leaf A-D route was originated in response to an S-PMSI or I-PMSI A-D route (see section "Leaf A-D Route for Global Table Multicast"). In this case the egress ABR MUST find a S-PMSI or I-PMSI route whose NLRI has the same value as the Route Key field of the received Leaf A-D route. If such a matching route is found then the Leaf A-D route MUST be accepted, else it MUST be discarded. If the Leaf A-D route is accepted and if its the first Leaf A-D route update for the Route Key field in the route, or the withdrawl of the last Leaf A-D route for the Route Key field then the following procedures will be executed.

If the RD of the received Leaf A-D route is set to all 0s or all 1s then the received Leaf A-D route is for the global table multicast service.

If the received Leaf A-D route is the first Leaf A-D route update for

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the Route Key field carried in the route, then the egress ABR originates a Leaf A-D route, whose MCAST-VPN NLRI is constructed as follows.

The Route Key field of MCAST-VPN NLRI is the same as the Route Key field of MCAST-VPN NLRI of the received Leaf A-D route. The Originating Router's IP address field of MCAST-VPN NLRI is set to the address of the local ABR (the ABR that originates the route).

The Next Hop field of the MP REACH NLRI attribute of the route SHOULD be set to the same IP address as the one carried in the Originating Router's IP Address field of the route.

To constrain distribution of this route the originating egress ABR constructs an IP-based Route Target Extended Community by placing the IP address of the upstream node in the Global Administrator field of the community, with the Local Administrator field of this community set to 0, and sets the Extended Communities attribute of this Leaf A-D route to that community.

The upstream node's IP address is the IP address determined from the Global Administrator field of the Inter-area P2MP Segmented Next-hop Extended Community, where this Extended Community is obtained as follows. When the Leaf A-D route is for MVPN or VPLS, then this Extended Community is the one carried in the I-PMSI/S-PMSI A-D route that matches the Leaf A-D route. When the Leaf A-D route is for global table multicast, then this Extended Community is the one carried in the best unicast route to the Ingress PE. The Ingress PE address is determined from the received Leaf A-D route. The best unicast route MUST first be determined from multicast SAFI i.e., SAFI 2 routes, if present.

The ABR then advertises this Leaf A-D route to the upstream node in the backbone area.

Mechanisms specific in [RFC4684] for constrained BGP route distribution can be used along with this specification to ensure that only the needed PE/ABR will have to process a said Leaf A-D route.

When Ingress Replication is used to instantiate the backbone area segment then the Leaf A-D route originated by the egress ABR MUST carry a downstream assigned label in the P-Tunnel Attribute where the P-Tunnel type is set to Ingress Replication. The ABR MUST assign a distinct MPLS label for each Leaf A-D route that it originates.

In order to support global table multicast an egress ABR MUST autoconfigure an import AS-based Route Target Extended Community with the Global Administrator field set to the AS of the ABR and the Local

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Administrator field set to 0.

If the received Leaf A-D route is the withdrawl of the last Leaf A-D route for the Route Key carried in the route, then the egress ABR must withdraw the Leaf A-D route associated with that Route Key that has been previously advertised by the egress ABR in the backbone area.

7.2. P2MP LSP as the Intra-Area LSP in the Egress Area

This section describes procedures for using intra-area P2MP LSPs in the egress area. The procedures that are common to both P2MP RSVP-TE and P2MP LDP are described first, followed by procedures that are specific to the signaling protocol.

When P2MP LSPs are used as the intra-area LSPs, note that an existing intra-area P2MP LSP may be used solely for a particular inter-area P2MP service LSP, or for other inter-area P2MP service LSPs as well. The choice between the two options is purely local to the egress ABR. The first option provides one-to-one mapping between inter-area P2MP service LSPs and intra-area P2MP LSPs; the second option provides many-to-one mapping, thus allowing to aggregate forwarding state.

7.2.1. Received Leaf A-D route is for MVPN or VPLS

If the value of the third octet of the MCAST-VPN NLRI of the received Leaf A-D route is either 0x01, or 0x02, or 0x03, this indicates that the Leaf A-D route was originated in response to an MVPN or VPLS S-PMSI or I-PMSI A-D route (see section "Leaf A-D Route for Global Table Multicast"). In this case the ABR MUST re-advertise in the egress area the MVPN/VPLS A-D route that matches the Leaf A-D route to signal the binding of the intra-area P2MP LSP to the inter-area P2MP service LSP. This must be done ONLY if (a) such a binding hasn't already been advertised, or (b) the binding has changed. The readvertised route MUST carry the Inter-area P2MP Segmented Next-Hop Extended Community.

The PMSI Tunnel attribute of the re-advertised route specifies either an intra-area P2MP RSVP-TE LSP or an intra-area P2MP LDP LSP rooted at the ABR and MUST also carry an upstream assigned MPLS label. The upstream-assigned MPLS label MUST be set to implicit NULL if the mapping between the inter-area P2MP service LSP and the intra-area P2MP LSP is one-to-one. If the mapping is many-to-one the intra-area segment of the inter-area P2MP service LSP (referred to as the "inner" P2MP LSP) is constructed by nesting the inter-area P2MP

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service LSP in an intra-area P2MP LSP (referred to as the "outer" intra-area P2MP LSP), by using P2MP LSP hierarchy based on upstreamassigned MPLS labels [RFC 5332].

If segments of multiple MVPN or VPLS S-PMSI service LSPs are carried over a given intra-area P2MP LSP, each of these segments MUST carry a distinct upstream-assigned label, even if all these service LSPs are for (C-S/*, C-G/*)s from the same MVPN/VPLS. Therefore, an ABR maintains an LFIB state for each of the (C-S/*, C-G/*)s carried over S-PMSIs traversing this ABR (that applies to both the ingress and the egress ABRs).

7.2.2. Received Leaf A-D route is for global table multicast

When the RD of the received Leaf A-D route is set to all 0s or all 1s, then this is the case of inter-area P2MP service LSP being associated with the global table multicast service. The procedures for this are described below.

7.2.2.1. Global Table Multicast and S-PMSI A-D Routes

This section applies only if it is desirable to send a particular (S, G) or (*, G) global table multicast flow only to those egress PEs that have receivers for that multicast flow.

If the egress ABR have not previousely received (and re-advertised) an S-PMSI A-D route for (S, G) or (*, G) that has been originated by an ingress PE/ASBR (see section "P2MP LSP as the Intra-Area LSP in the Ingress Area"), then the egress ABR MUST originate a S-PMSI A-D route. The PMSI Tunnel attribute of the route MUST contain the identity of the intra-area P2MP LSP and an upstream assigned MPLS label (although this label may be an Implicit NULL - see section "General Assumptions and Terminology"). The RD, Multicast Source Length, Multicast Source, Multicast Group Length (1 octet), and Multicast Group fields of the NLRI of this route are the same as of the received Leaf A-D route. The Originating Router's IP address field in the S-PMSI A-D route is the same as the Ingress PE's IP address field in the received Leaf A-D route. The Route Target of this route is an AS-specific Route Target Extended Community with the Global Administrator field set to the AS of the advertising ABR, and the Local Administrator field is set to 0. The route MUST carry the Inter-area P2MP Segmented Next-Hop Extended Community. This community is constructed following the procedures in section 4 ("Inter-area P2MP Segmented Next-Hop Extended Community").

The egress ABR MUST advertise this route into the egress area. PEs in

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the egress area that participate in the global table multicast will import this route.

A PE in the egress area that originated the Leaf A-D route SHOULD join the P2MP LSP advertised in the PMSI Tunnel attribute of the S-PMSI A-D route.

7.2.2.2. Global Table Multicast and Wildcard S-PMSI A-D Routes

It may be desirable for an ingress PE to carry multiple multicast flows associated with the global table multicast over the same interarea P2MP service LSP. This can be achieved using wildcard, i.e., (*,*) S-PMSI A-D routes [RFC6625]. An ingress PE MAY advertise a wildcard S-PMSI A-D route as described in section 9 ("Ingress PE/ASBR Procedures").

If the ingress PE originates a wildcard S-PMSI A-D route, and the egress ABR receives this route from the ingress ABR, then the egress ABR either (a) MUST re-advertise this route into the egress area with the PMSI Tunnel Attribute containing the identifier of the intra-area P2MP LSP in the egress area and an upstream assigned label (note that this label may be an Implicit NULL - see section "General Assumptions and Terminology") assigned to the inter-area wildcard S-PMSI, or (b) MUST be able to disaggregate traffic carried over the wildcard S-PMSI onto the egress area (S, G) or (*, G) S-PMSIs. The procedures for such disaggregation require IP processing on the egress ABRs.

If the egress ABR advertises a wildcard S-PMSI A-D route into the egress area, this route MUST carry AS-specific Route Target Extended Community with the Global Administrator field set to the AS of the advertising ABR, and the Local Administrator field set to 0. PEs in the egress area that participate in the global table multicast will import this route.

A PE in the egress area SHOULD join the P2MP LSP advertised in the PMSI Tunnel attribute of the wildcard S-PMSI A-D route if (a) the Originating Router's IP Address field in the S-PMSI A-D route has the same value as the Ingress PE's IP address in at least one of the Leaf A-D routes for global table multicast originated by the PE, and (b) the upstream ABR for the Ingress PE's IP address in that Leaf A-D route is the (egress) ABR that advertises the wildcard S-PMSI A-D route.

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7.2.3. Global Table Multicast and the Expected Upstream Node

If the mapping between the inter-area P2MP service LSP for global table multicast service and the intra-area P2MP LSP is many-to-one then an egress PE must be able to determine whether a given multicast packet for a particular (S, G) is received from the "expected" upstream node. The expected node is the node towards which the Leaf A-D route is sent by the egress PE. Packets received from another upstream node for that (S, G) MUST be dropped. To allow the egress PE to determine the sender upstream node, the intra-area P2MP LSP must be signaled with no PHP, when the mapping between the inter-area P2MP service LSP for global table multicast service and the intra-area P2MP LSP is many-to-one.

Further the egress ABR MUST first push onto the label stack the upstream assigned label advertised in the S-PMSI A-D route, if the label is not the Implicit NULL.

7.2.4. P2MP LDP LSP as the Intra-Area P2MP LSP

The above procedures are sufficient if P2MP LDP LSPs are used as the Intra-area P2MP LSP in the Egress area.

7.2.5. P2MP RSVP-TE LSP as the Intra-Area P2MP LSP

If P2MP RSVP-TE LSP is used as the the intra-area LSP in the egress area, then the egress ABR can either (a) graft the leaf (whose IP address is specified in the received Leaf A-D route) into an existing P2MP LSP rooted at the egress ABR, and use that LSP for carrying traffic for the inter-area segmented P2MP service LSP, or (b) originate a new P2MP LSP to be used for carrying (S,G).

When the RD of the received Leaf A-D route is all 0s or all 1s, then the procedures are as described in section 7.1.2 ("RD of the received Leaf A-D route is all 0s or all 1s").

Note also that the SESSION object that the egress ABR would use for the intra-area P2MP LSP need not encode the P2MP FEC from the received Leaf A-D route.

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7.3. Ingress Replication in the Egress Area

When Ingress Replication is used to instantiate the egress area segment then the Leaf A-D route advertised by the egress PE MUST carry a downstream assigned label in the P-Tunnel Attribute where the P-Tunnel type is set to Ingress Replication. We will call this label the egress PE downstream assigned label.

The egress ABR MUST forward packets received from the backbone area intra-area segment, for a particular inter-area P2MP LSP, to all the egress PEs from which the egress ABR has imported a Leaf A-D route for the inter-area P2MP LSP. A packet to a particular egress PE is encapsulated, by the egress ABR, using a MPLS label stack the bottom label of which is the egress PE downstream assigned label. The top label is the P2P RSVP-TE or the MP2P LDP label to reach the egress PE.

Note that these procedures ensures that an egress PE always receives packets only from the upstream node expected by the egress PE.

8. Ingress ABR Procedures

When an ingress ABR receives a Leaf A-D route and the Route Target Extended Community carried by the route contains the IP address of this ABR, then the ingress ABR follows the same procedures as in section 7 ("Egress ABR Procedures"), with egress ABR replaced by ingress ABR, backbone area replaced by ingress area, and backbone area segment replaced by ingress area segment.

In order to support global table multicast the ingress ABR MUST be auto-configured with an import AS-based Route Target Extended Community whose Global Administrator field is set to the AS of the ABR and the Local Administrator field is set to 0.

8.1. P2MP LSP as the Intra-Area LSP in the Backbone Area

The procedures for binding the backbone area segment of an inter-area P2MP LSP to the intra-area P2MP LSP in the backbone area are the same as in section 7 ("Egress ABR Procedures") and sub-section 7.1 ("P2MP LSP as the Intra-Area LSP in the Egress Area"), with egress PE being replace by egress ABR, egress ABR being replaced by ingress ABR, and egress area being replaced by backbone area. This applies to the inter-area P2MP LSPs associated with either MVPN, or VPLS, or global table multicast.

It is to be noted that in the case of global table multicast, if the

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backbone area uses wildcard S-PMSI, then the egress area also SHOULD use wildcard S-PMSI for global table multicast, or the egress ABRs MUST be able to disaggregate traffic carried over the wildcard S-PMSI onto the egress area (S, G) or (*, G) S-PMSIs. The procedures for such disaggregation require IP processing on the egress ABRs.

8.2. Ingress Replication in the Backbone Area

When Ingress Replication is used to instantiate the backbone area segment then the Leaf A-D route advertised by the egress ABR MUST carry a downstream assigned label in the P-Tunnel Attribute where the P-Tunnel type is set to Ingress Replication. We will call this the egress ABR downstream assigned label. The egress ABR MUST assign a distinct MPLS label for each Leaf A-D route originated by the ABR.

The ingress ABR MUST forward packets received from the ingress area intra-area segment, for a particular inter-area P2MP LSP, to all the egress ABRs from which the ingress ABR has imported a Leaf A-D route for the inter-area P2MP LSP. A packet to a particular egress ABR is encapsulated, by the ingress ABR, using a MPLS label stack the bottom label of which is the egress ABR downstream assigned label. The top label is the P2P RSVP-TE or the MP2P LDP label to reach the egress ABR.

9. Ingress PE/ASBR Procedures

This section describes Ingress PE/ASBR procedures for constructing segmented inter-area P2MP LSP.

When an ingress PE/ASBR receives a Leaf A-D route and the Route Target Extended Community carried by the route contains the IP address of this PE/ASBR, then the following procedures will be executed.

If the value of the third octet of the MCAST-VPN NLRI of the received Leaf A-D route is either 0x01, or 0x02, or 0x03, this indicates that the Leaf A-D route was originated in response to an S-PMSI or I-PMSI A-D route (see section "Leaf A-D Route for Global Table Multicast"). In this case the ingress PE/ASBR MUST find a S-PMSI or I-PMSI route whose NLRI has the same value as the Route Key field of the received Leaf A-D route. If such a matching route is found then the Leaf A-D route MUST be accepted else it MUST be discarded. If the Leaf A-D route is accepted then it MUST be processed as per MVPN or VPLS procedures.

If the RD of the received A-D route is set to all 0s or all 1s, then

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the received Leaf A-D route is for the global table multicast service. If this is the first Leaf A-D route for the Route Key carried in the route, the PIM implementation MUST set its upstream (S/RP,G) state machine to Joined state for the (S/RP, G) received via a Leaf A-D route update. Likewise, if this is the withdrawal of the last Leaf A-D route whose Route Key matches a particular (S/RP, G) state, the PIM implementation MUST set its upstream (S/RP, G) state machine to Pruned state for the (S/RP, G).

9.1. P2MP LSP as the Intra-Area LSP in the Ingress Area

If the value of the third octet of the MCAST-VPN NLRI of the received Leaf A-D route is either 0x01, or 0x02, or 0x03 (which indicates that the Leaf A-D route was originated in response to an S-PMSI or I-PMSI A-D route), and P2MP LSP is used as the the intra-area LSP in the ingress area, then the procedures for binding the ingress area segment of the inter-area P2MP LSP to the intra-area P2MP LSP in the ingress area, are the same as in section 7 ("Egress ABR Procedures") and sub-section 7.1.1 ("P2MP LSP as the Intra-Area LSP in the Egress Area").

When the RD of the received Leaf A-D route is all 0s or all 1s, as is the case where the inter-area service P2MP LSP is associated with the global table multicast service, then the ingress PE MAY originate a S-PMSI A-D route with the RD, multicast source, multicast group fields being the same as those in the received Leaf A-D route.

Further in the case of global table multicast an ingress PE MAY originate a wildcard S-PMSI A-D route as per the procedures in [RFC6625] with the RD set to 0. This route may be originated by the ingress PE based on configuration or based on the import of a Leaf A-D route with RD set to 0. If an ingress PE originates such a route, then the ingress PE MAY decide not to originate (S, G) or (*, G) S-PMSI A-D routes.

The wildcard S-PMSI A-D route MUST carry the Inter-area P2MP Segmented Next-Hop Extended Community. This community is constructed following the procedures in section 4 ("Inter-area P2MP Segmented Next-Hop Extended Community").

It is to be noted that in the case of global table multicast, if the ingress area uses wildcard S-PMSI, then the backbone area also SHOULD use wildcard S-PMSI for global table multicast, or the ingress ABRS MUST be able to disaggregate traffic carried over the wildcard S-PMSI onto the backbone area (S, G) or (*, G) S-PMSIs. The procedures for such disaggregation require IP processing on the ingress ABRs.

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9.2. Ingress Replication in the Ingress Area

When Ingress Replication is used to instantiate the ingress area segment then the Leaf A-D route advertised by the ingress ABR MUST carry a downstream assigned label in the P-Tunnel Attribute where the P-Tunnel type is set to Ingress Replication. We will call this the ingress ABR downstream assigned label. The ingress ABR MUST assign a distinct MPLS label for each Leaf A-D route originated by the ABR.

The ingress PE/ASBR MUST forward packets received from the CE, for a particular inter-area P2MP LSP, to all the ingress ABRs from which the ingress PE/ASBR has imported a Leaf A-D route for the inter-area P2MP LSP. A packet to a particular ingress ABR is encapsulated, by the ingress PE/ASBR, using a MPLS label stack the bottom label of which is the ingress ABR downstream assigned label. The top label is the P2P RSVP-TE or the MP2P LDP label to reach the ingress ABR.

10. Common Tunnel Type in the Ingress and Egress Areas

For a given inter-area service P2MP LSP, the PE/ASBR that is the root of that LSP controls the tunnel type of the intra-area P-tunnel that carries the ingress area segment of that LSP. However, the tunnel type of the intra-area P-tunnel that carries the backbone area segment of that LSP may be different from the tunnel type of the intra-area P-tunnels that carry the ingress area segment and the egress area segment of that LSP. In that situation if for a given inter-area P2MP LSP it is desirable/necessary to use the same tunnel type for the intra-area P-tunnels that carry the ingress area segment and the egress area segment of that LSP, then the following procedures on the ingress ABR and egress ABR provide this functionality.

When an ingress ABR re-advertises into the backbone area a BGP MVPN I-PMSI, or S-PMSI A-D route, or VPLS A-D route, the ingress ABR places the PMSI Tunnel attribute of this route into the ATTR_SET BGP Attribute [RFC6368], adds this attribute to the re-advertised route, and then replaces the original PMSI Tunnel attribute with a new one (note, that the Tunnel type of the new attribute may be different from the Tunnel type of the original attribute).

When an egress ABR re-advertises into the egress area a BGP MVPN I-PMSI or S-PMSI A-D route, or VPLS A-D route, if the route carries the ATTR_SET BGP attribute [RFC6368], then the ABR sets the Tunnel type of the PMSI Tunnel attribute in the re-advertised route to the Tunnel type of the PMSI Tunnel attribute carried in the ATTR SET BGP attribute, and removes the ATTR_SET from the route.

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11. Placement of Ingress and Egress PEs

As described in the earlier sections, procedures in this document allow the placement of ingress and egress PEs in the backbone area. They also allow the placement of egress PEs in the ingress area or the placement of ingress PEs in the egress area.

For instance, ABRs in the backbone area may act as ingress and egress PEs for global table multicast, as per the ingress and egress PE definition in this document. This may be the case if the service is global table multicast and relies on global table multicast in the ingress and egress areas and its desirable to carry global table multicast over MPLS in the backbone area. This may also be the case if the service is MVPN and the P-tunnel technology in the ingress and egress areas uses PIM based IP/GRE P-tunnels. As far as the ABRs are concerned PIM signaling for such P-Tunnels is handled as per the ingress/egress PE global table multicast procedures specified in this document. To facilitate this the ABRs may advertise their loopback addresses in BGP using multicast-SAFI i.e., SAFI 2, if non-congruence between unicast and multicast is desired.

12. Data Plane

This section describes the data plane procedures on the ABRs, ingress PEs, egress PEs and transit routers.

12.1. Data Plane Procedures on ABRs

When procedures in this document are followed to signal inter-area P2MP Segmented LSPs, then ABRs are required to perform only MPLS switching. When an ABR receives a MPLS packet from an "incoming" intra-area segment of the inter-area P2MP Segmented LSP, it forwards the packet, based on MPLS switching, onto another "outgoing" intraarea segment of the inter-area P2MP Segmented LSP.

If the outgoing intra-area segment is instantiated using a P2MP LSP, and if there is a one-to-one mapping between the outgoing intra-area segment and the P2MP LSP, then the ABR MUST pop the incoming segment's label stack and push the label stack of the outgoing P2MP LSP. If there is a many-to-one mapping between outgoing intra-area segments and the P2MP LSP then the ABR MUST pop the incoming segment's label stack and first push the upstream assigned label corresponding to the outgoing intra-area segment, if such a label has been assigned, and then push the label stack of the outgoing P2MP LSP.

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If the outgoing intra-area segment is instantiated using ingress replication then the ABR must pop the incoming segment's label stack and replicate the packet once to each leaf ABR or PE of the outgoing intra-area segment. The label stack of the packet sent to each such leaf MUST first include a downstream assigned label assigned by the leaf to the segment, followed by the label stack of the P2P or MP2P LSP to the leaf.

12.2. Data Plane Procedures on Egress PEs

An egress PE must first identify the inter-area P2MP segmented LSP based on the incoming label stack. After this identification the egress PE must forward the packet using the application that is bound to the inter-area P2MP segmented LSP.

Note that the application specific forwarding for MVPN service may require the egress PE to determine whether the packets were received from the expected sender PE. When the application is MVPN then the FEC of an inter-area P2MP Segmented LSP is at the granularity of the sender PE. Note that MVPN intra-AS I-PMSI A-D routes and S-PMSI A-D routes both carry the Originating Router IP Address. Thus an egress PE could associate the data arriving on P-tunnels advertised by these routes with the Originating Router IP Address carried by these routes which is the same as the ingress PE. Since a unique label stack is associated with each such FEC, the egress PE can determine the sender PE from the label stack.

Likewise for VPLS service for the purposes of MAC learning the egress PE must be able to determine the "VE-ID" from which the packets have been received. The FEC of the VPLS A-D routes carries the VE-ID. Thus an egress PE could associate the data arriving on P-tunnels advertised by these routes with the VE-ID carried by these routes. Since a unique label stack is associated with each such FEC, the egress PE can perform MAC learning for packets received from a given VE-ID.

When the application is global table multicast it is sufficient for the label stack to include identification of the sender upstream node. When P2MP LSPs are used this requires that PHP MUST be turned off. When Ingress Replication is used the egress PE knows the incoming downstream assigned label to which it has bound a particular $(S/^*, G)$ and must accept packets with only that label for that $(S/^*, G)$.

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12.3. Data Plane Procedures on Ingress PEs

An Ingress PE must perform application specific forwarding procedures to identify the outgoing inta-area segment of an incoming packet.

If the outgoing intra-area segment is instantiated using a P2MP LSP, and if there is a one-to-one mapping between the outgoing intra-area segment and the P2MP LSP, then the ingress PE MUST encapsulate the packet in the label stack of the outgoing P2MP LSP. If there is a many-to-one mapping between outgoing intra-area segments and the P2MP LSP then the PE MUST first push the upstream assigned label corresponding to the outgoing intra-area segment, if such a label has been assigned, and then push the label stack of the outgoing P2MP LSP.

If the outgoing intra-area segment is instantiated using ingress replication then the PE must replicate the packet once to each leaf ABR or PE of the outgoing intra-area segment. The label stack of the packet sent to each such leaf MUST first include a downstream assigned label assigned by the leaf to the segment, followed by the label stack of the P2P or MP2P LSP to the leaf.

12.4. Data Plane Procedures on Transit Routers

When procedures in this document are followed to signal inter-area P2MP Segmented LSPs then transit routers in each area perform only MPLS switching.

13. Support for Inter-Area Transport LSPs

This section describes OPTIONAL procedures that allow to aggregate multiple (inter-area) P2MP service LSPs into a single inter-area P2MP transport LSPs, and then apply the segmentation procedures, as specified in this document, to these inter-area P2MP transport LSPs (rather than applying these procedures directly to the inter-area P2MP service LSPs).

<u>13.1</u>. Transport Tunnel Tunnel Type

For the PMSI Tunnel Attribute we define a new Tunnel type, called "Transport Tunnel", whose Tunnel Identifier is a tuple <Source PE Address, Local Number>. This Tunnel type is assigned a value of 8. The Source PE Address is the address of the PE that originates the (service) A-D route that carries this attribute, and the Local Number is a number that is unique to the Source PE. The length of the Local

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Number part is the same as the length of the Source PE Address. Thus if the Source PE Address is an IPv4 address, then the Local Number part is 4 octets, and if the Source PE Address is an IPv6 address, then the Local Number part is 16 octets.

13.2. Discovering Leaves of the Inter-Area P2MP Service LSP

When aggregating multiple P2MP LSPs using P2MP LSP hierarchy, determining the leaf nodes of the LSPs being aggregated is essential for being able to tradeoff the overhead due to the P2MP LSPs vs suboptimal use of bandwidth due to the partial congruency of the LSPs being aggregated.

Therefore, if a PE that is a root of a given service P2MP LSP wants to aggregate this LSP with other (service) P2MP LSPs rooted at the same PE into an inter-area P2MP transport LSP, the PE should first determine all the leaf nodes of that service LSP, as well as those of the other service LSPs being aggregated).

To accomplish this the PE sets the PMSI Tunnel attribute of the service A-D route (an I-PMSI or S-PMSI A-D route for MVPN or VPLS service) associated with that LSP as follows. The Tunnel Type is set to "No tunnel information present", Leaf Information Required flag is set to 1, the route MUST NOT carry the P2MP Segmented Next-Hop Extended Community. In contrast to the procedures for the MVPN and VPLS A-D routes described so far, the Route Reflectors that participate in the distribution of this route need not be ABRs.

13.3. Discovering P2MP FEC of P2MP Transport LSP

Once the ingress PE determines all the leaves of a given P2MP service LSP, the PE (using some local to the PE criteria) selects a particular inter-area transport P2MP LSP to be used for carrying the (inter-area) service P2MP LSP.

Once the PE selects the transport P2MP LSP, the PE (re)originates the service A-D route. The PMSI Tunnel attribute of this route now carries the Tunnel Identifier of the selected transport LSP, with the Tunnel Type set to "Transport Tunnel". If the transport LSP carries multiple P2MP service LSPs, then the MPLS Label field in the attribute carries an upstream-assigned label assigned by the PE that is bound to the P2MP FEC of the inter-area P2MP service LSP. Otherwise, this field is set to Implicit NULL.

Just as described earlier, this service A-D route MUST NOT carry the P2MP Segmented Next-Hop Extended Community. Just as described

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earlier, the Route Reflectors that participate in the distribution of this route need not be ABRs.

13.4. Egress PE Procedures for P2MP Transport LSP

When an egress PE receives and accepts an MVPN or VPLS service A-D route, if the Leaf Information Required flag in the PMSI Tunnel attribute of the received A-D route is set to 1, and the route does not carry the P2MP Segmented Next-Hop Extended Community, then the egress PE, following the "regular" MVPN or VPLS procedures associated with the received A-D route (as specified in [RFC6514] and [VPLS-P2MP]), originates a Leaf A-D route.

In addition, if the Tunnel Type in the PMSI Tunnel attribute of the received service A-D route is set to "Transport Tunnel", the egress PE originates yet another Leaf A-D route.

The format of the Route Key field of MCAST-VPN NLRI of this additional Leaf A-D route is the same as defined in section 6.2.2 ("Leaf A-D Route for Global Table Multicast"). The Route Key field of MCAST-VPN NLRI of this route is constructed as follows:

RD (8 octets) - set to 0

Multicast Source Length, Multicast Source - constructed from the Source PE address part of the Tunnel Identifier carried in the PMSI Tunnel attribute of the received service S-PMSI A-D route.

Multicast Group Length, Multicast Group - constructed from Local Number part of the Tunnel Identifier carried in the PMSI Tunnel attribute of the received service S-PMSI A-D route.

Ingress PE IP Address is set to the Source PE address part of the Tunnel Identifier carried in the PMSI Tunnel attribute of the received service S-PMSI A-D route.

The egress PE, when determining the upstream ABR, follows the procedures specified in section 6.1 ("Determining the Upstream ABR/PE/ASBR (Upstream Node)") for global table multicast.

The egress PE constructs the rest of the Leaf A-D route following the procedures specified in $\underline{\text{section } 6.2.3}$ ("Constructing the Rest of the Leaf A-D Route").

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From that point on we follow the procedures used for the Leaf A-D routes for global table multicast, as outlined below.

13.5. Egress/Ingress ABR, Ingress PE procedures for P2MP Transport LSP

When an egress ABR receives the Leaf A-D route, and P2MP LSP is used to instantiate the egress area segment of the inter-area transport LSP, the egress ABR will advertise into the egress area an S-PMSI A-D route. This route is constructed following the procedures in section 7.1.2.1 ("Global Table Multicast and S-PMSI A-D Routes"). The egress PE(s) will import this route.

The egress ABR will also propagate, with appropriate modifications, the received Leaf A-D route into the backbone area. This is irrespective of whether the egress area segment is instantiated using P2MP LSP or ingress replication.

If P2MP LSP is used to instantiate the backbone area segment of the inter-area transport LSP, then an ingress ABR will advertise into the backbone area an S-PMSI A-D route. This route is constructed following the procedures in section 7.1.2.1 ("Global Table Multicast and S-PMSI A-D Routes"). The egress ABR(s) will import this route.

The ingress ABR will also propagate, with appropriate modifications, the received Leaf A-D route into the ingress area towards the ingress/root PE. This is irrespective of whether the backbone area segment is instantiated using P2MP LSP or ingress replication.

Finally, if P2MP LSP is used to instantiate the ingress area segment, the ingress PE will advertise into the ingress area an S-PMSI A-D route with the RD, multicast source, and multicast group fields being the same as those in the received Leaf A-D route. The PMSI Tunnel attribute of this route contains the identity of a the intra-area P2MP LSP used to instantiate the ingress area segment, and an upstream-assigned MPLS label. The ingress ABR(s) and other PE(s) in the ingress area, if any, will import this route. The ingress PE will use the (intra-area) P2MP LSP advertised in this route for carrying traffic associated with the original service A-D route advertised by the PE.

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13.6. Discussion

Use of inter-area transport P2MP LSPs, as described in this section, creates a level of indirection between (inter-area) P2MP service LSPs, and intra-area transport LSPs that carry the service LSPs. Rather than segmenting (inter-area) service P2MP LSPs, and then aggregating (intra-area) segments of these service LSPs into intra-area transport LSPs, this approach first aggregates multiple (inter-area) service P2MP LSPs into a single inter-area transport P2MP LSP, then segments such inter-area transport P2MP LSPs, and then aggregates (intra-area) segments of these inter-area transport LSPs into intra-area transport LSPs.

On one hand this approach could result in reducing the state (and the overhead associated with maintaining the state) on ABRs. This is because instead of requiring ABRs to maintain state for individual P2MP service LSPs, ABRs would need to maintain state only for the inter-area P2MP transport LSPs. Note however, that this reduction is possible only if a single inter-area P2MP transport LSP aggregates more than one (inter-area) service LSP. In the absence of such aggregation, use of inter-area transport LSPs provides no benefits, yet results in extra overhead. And while such aggregation does allow to reduce state on ABRs, it comes at a price, as described below.

As we mentioned before, aggregating multiple P2MP service LSPs into a single inter-area P2MP transport LSP requires the PE rooted at these LSPs to determine all the leaf nodes of the service LSPs to be aggregated. This means that the root PE has to track all the leaf PEs of these LSPs. In contrast, when one applies segmentation procedures directly to the P2MP service LSPs, the root PE has to track only the leaf PEs in its own IGP area, plus the ingress ABR(s). Likewise, an ingress ABR has to track only the egress ABRs. Finally, an egress ABR has to track only the leaf PEs in its own area. Therefore, while the total overhead of leaf tracking due to the P2MP service LSPs is about the same in both approaches, the distribution of this overhead is different. Specifically, when one uses inter-area P2MP transport LSPs, this overhead is concentrated on the ingress PE. When one applies segmentation procedures directly to the P2MP service LSPs, this overhead is distributed among ingress PE and ABRs.

Moreover, when one uses inter-area P2MP transport LSPs, ABRs have to bear the overhead of leaf tracking for inter-area P2MP transport LSPs. In contrast, when one applies segmentation procedures directly to the P2MP service LSPs, there is no such overhead (as there are no inter-area P2MP transport LSPs).

Use of inter-area P2MP transport LSPs may also result in more bandwidth inefficiency, as compared to applying segmentation

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procedures directly to the P2MP service LSPs. This is because with inter-area P2MP transport LSPs the ABRs aggregate segments of interarea P2MP transport LSPs, rather than segments of (inter-area) P2MP service LSPs. To illustrate this consider the following example.

Assume PE1 in Area 1 is an ingress PE, with two multicast streams, (C-S1, C-G1) and (C-S2, C-G2), originated by an MVPN site connected to PE1. Assume that PE2 in Area 2 has an MVPN site with receivers for (C-S1, C-G1), PE3 and PE4 in Area 3 have an MVPN site with receivers both for (C-S1, C-G1) and for (C-S2, C-G2). Finally, assume that PE5 in Area 4 has an MVPN site with receivers for (C-S2, C-G2).

When segmentation applies directly to the P2MP service LSPs then Area 2 would have just one intra-area transport LSP which would carry the egress area segment of the P2MP service LSP for (C-S1, C-G1); Area 3 would have just one intra-area transport LSP which would carry the egress area segments of both the P2MP service LSP for (C-S1, C-G1) and the P2MP service LSP for (C-S2, C-G2); Area 4 would have just one intra-area transport LSP which would carry the egress area segment of the P2MP service LSP for (C-S2, C-G2). Note that there is no bandwidth inefficiency in this case at all.

When using inter-area P2MP transport LSPs, to achieve the same state overhead on the routers within each of the egress areas (except for egress ABRs), PE1 would need to aggregate the P2MP service LSP for (C-S1, C-G1) and the P2MP service LSP for (C-S2, C-G2) into the same inter-area P2MP transport LSP. While such aggregation would reduce state on ABRs, it would also result in bandwidth inefficiency, as (C-S1, C-G1) will be delivered not just to PE2, PE3, and PE4, but also to PE5, which has no receivers for this stream. Likewise, (C-S2, C-G2) will be delivered not just to PE3, PE4, and PE5, but also to PE2, which has no receivers for this stream.

14. IANA Considerations

This document defines a new BGP Extended Community called "Inter-area P2MP Segmented Next-Hop" (see section "Inter-area P2MP Segmented Next-Hop Extended Community"). This community is IP Address Specific, of an extended type, and is transitive. A codepoint for this community should be assigned both from the IPv4 Address Specific Extended Community registry, and from the IPv6 Address Specific Extended Community registry. The same code point should be assigned from both registries.

This document also assigns a new Tunnel Type in the PMSI Tunnel Attribute, called the "Transport Tunnel" (see section "Transport Rekhter [Page 36]

Tunnel Type"). This Tunnel Type is assigned a value of 8.

15. Security Considerations

These will be spelled out in a future revision.

16. Acknowledgements

We would like to thank Eric Rosen for his comments.

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