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Network Management Datastore Architecture
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Abstract

Datastores are a fundamental concept binding the data models written in the YANG data modeling language to network management protocols such as NETCONF and RESTCONF. This document defines an architectural framework for datastores based on the experience gained with the initial simpler model, addressing requirements that were not well supported in the initial model.

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[1.](#) Introduction

This document provides an architectural framework for datastores as they are used by network management protocols such as NETCONF [[RFC6241](#)], RESTCONF [[RFC8040](#)] and the YANG [[RFC7950](#)] data modeling language. Datastores are a fundamental concept binding network management data models to network management protocols. Agreement on a common architectural model of datastores ensures that data models can be written in a network management protocol agnostic way. This architectural framework identifies a set of conceptual datastores but it does not mandate that all network management protocols expose all these conceptual datastores. This architecture is agnostic with regard to the encoding used by network management protocols.

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [BCP 14](#) [[RFC2119](#)] [[RFC8174](#)] when, and only when, they appear in all capitals, as shown here.

[2.](#) Objectives

Network management data objects can often take two different values, the value configured by the user or an application (configuration) and the value that the device is actually using (operational state). These two values may be different for a number of reasons, e.g., system internal interactions with hardware, interaction with protocols or other devices, or simply the time it takes to propagate a configuration change to the software and hardware components of a system. Furthermore, configuration and operational state data objects may have different lifetimes.

The original model of datastores required these data objects to be modeled twice in the YANG schema, as "config true" objects and as "config false" objects. The convention adopted by the interfaces data model ([[RFC7223](#)]) and the IP data model ([[RFC7277](#)]) was using two separate branches rooted at the root of the data tree, one branch for configuration data objects and one branch for operational state data objects.

The duplication of definitions and the ad-hoc separation of operational state data from configuration data leads to a number of problems. Having configuration and operational state data in separate branches in the data model is operationally complicated and impacts the readability of module definitions. Furthermore, the relationship between the branches is not machine readable and filter expressions operating on configuration and on related operational state are different.

With the revised architectural model of datastores defined in this document, the data objects are defined only once in the YANG schema but independent instantiations can appear in two different datastores, one for configured values and one for operational state values. This provides a more elegant and simpler solution to the problem.

The revised architectural model of datastores supports additional datastores for systems that support more advanced processing chains converting configuration to operational state. For example, some systems support configuration that is not currently used (so called inactive configuration) or they support configuration templates that are used to expand configuration data via a common template.

3. Terminology

This document defines the following terminology. Some of the terms are revised definitions of terms originally defined in [\[RFC6241\]](#) and [\[RFC7950\]](#) (see also section [Section 4](#)). The revised definitions are semantically equivalent with the definitions found in [\[RFC6241\]](#) and [\[RFC7950\]](#). It is expected that the revised definitions provided in this section will replace the definitions in [\[RFC6241\]](#) and [\[RFC7950\]](#) when these documents are revised.

- o datastore: A conceptual place to store and access information. A datastore might be implemented, for example, using files, a database, flash memory locations, or combinations thereof. A datastore maps to an instantiated YANG data tree.
- o schema node: A node in the schema tree. The formal definition is in [RFC 7950](#).
- o datastore schema: The combined set of schema nodes for all modules supported by a particular datastore, taking into consideration any deviations and enabled features for that datastore.
- o configuration: Data that is required to get a device from its initial default state into a desired operational state. This data

is modeled in YANG using "config true" nodes. Configuration can originate from different sources.

- o configuration datastore: A datastore holding configuration.
- o running configuration datastore: A configuration datastore holding the current configuration of the device. It may include configuration that requires further transformations before it can be applied. This datastore is referred to as "<running>".
- o candidate configuration datastore: A configuration datastore that can be manipulated without impacting the device's running configuration datastore and that can be committed to the running configuration datastore. This datastore is referred to as "<candidate>".
- o startup configuration datastore: A configuration datastore holding the configuration loaded by the device into the running configuration datastore when it boots. This datastore is referred to as "<startup>".
- o intended configuration: Configuration that is intended to be used by the device. It represents the configuration after all configuration transformations to <running> have been performed and is the configuration that the system attempts to apply.
- o intended configuration datastore: A configuration datastore holding the complete intended configuration of the device. This datastore is referred to as "<intended>".
- o configuration transformation: The addition, modification or removal of configuration between the <running> and <intended> datastores. Examples of configuration transformations include the removal of inactive configuration and the configuration produced through the expansion of templates.
- o conventional configuration datastore: One of the following set of configuration datastores: <running>, <startup>, <candidate>, and <intended>. These datastores share a common datastore schema, and protocol operations allow copying data between these datastores. The term "conventional" is chosen as a generic umbrella term for these datastores.
- o conventional configuration: Configuration that is stored in any of the conventional configuration datastores.
- o dynamic configuration datastore: A configuration datastore holding configuration obtained dynamically during the operation of a

device through interaction with other systems, rather than through one of the conventional configuration datastores.

- o dynamic configuration: Configuration obtained via a dynamic configuration datastore.
- o learned configuration: Configuration that has been learned via protocol interactions with other systems and that is neither conventional nor dynamic configuration.
- o system configuration: Configuration that is supplied by the device itself.
- o default configuration: Configuration that is not explicitly provided but for which a value defined in the data model is used.
- o applied configuration: Configuration that is actively in use by a device. Applied configuration originates from conventional, dynamic, learned, system and default configuration.
- o system state: The additional data on a system that is not configuration, such as read-only status information and collected statistics. System state is transient and modified by interactions with internal components or other systems. System state is modeled in YANG using "config false" nodes.
- o operational state: The combination of applied configuration and system state.
- o operational state datastore: A datastore holding the complete operational state of the device. This datastore is referred to as "<operational>".
- o origin: A metadata annotation indicating the origin of a data item.
- o remnant configuration: Configuration that remains part of the applied configuration for a period of time after it has been removed from the intended configuration or dynamic configuration. The time period may be minimal, or may last until all resources used by the newly-deleted configuration (e.g., network connections, memory allocations, file handles) have been deallocated.

The following additional terms are not datastore specific but commonly used and thus defined here as well:

- o client: An entity that can access YANG-defined data on a server, over some network management protocol.
- o server: An entity that provides access to YANG-defined data to a client, over some network management protocol.
- o notification: A server-initiated message indicating that a certain event has been recognized by the server.
- o remote procedure call: An operation that can be invoked by a client on a server.

4. Background

NETCONF [[RFC6241](#)] provides the following definitions:

- o datastore: A conceptual place to store and access information. A datastore might be implemented, for example, using files, a database, flash memory locations, or combinations thereof.
- o configuration datastore: The datastore holding the complete set of configuration that is required to get a device from its initial default state into a desired operational state.

YANG 1.1 [[RFC7950](#)] provides the following refinements when NETCONF is used with YANG (which is the usual case but note that NETCONF was defined before YANG existed):

- o datastore: When modeled with YANG, a datastore is realized as an instantiated data tree.
- o configuration datastore: When modeled with YANG, a configuration datastore is realized as an instantiated data tree with configuration.

[RFC6244] defined operational state data as follows:

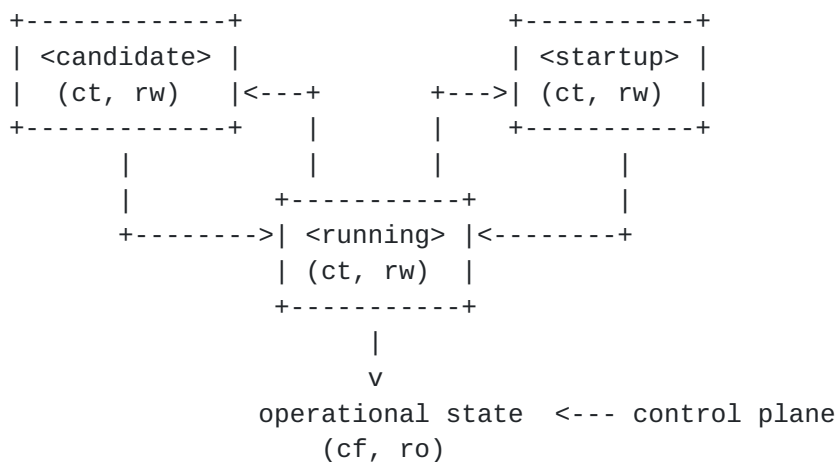
- o Operational state data is a set of data that has been obtained by the system at runtime and influences the system's behavior similar to configuration data. In contrast to configuration data, operational state is transient and modified by interactions with internal components or other systems via specialized protocols.

[Section 4.3.3 of \[RFC6244\]](#) discusses operational state and among other things mentions the option to consider operational state as being stored in another datastore. [Section 4.4](#) of this document then concludes that at the time of the writing, modeling state as distinct leafs and distinct branches is the recommended approach.

Implementation experience and requests from operators [[I-D.ietf-netmod-opstate-reqs](#)], [[I-D.openconfig-netmod-opstate](#)] indicate that the datastore model initially designed for NETCONF and refined by YANG needs to be extended. In particular, the notion of intended configuration and applied configuration has developed.

4.1. Original Model of Datastores

The following drawing shows the original model of datastores as it is currently used by NETCONF [[RFC6241](#)]:



ct = config true; cf = config false
 rw = read-write; ro = read-only
 boxes denote datastores

Note that this diagram simplifies the model: read-only (ro) and read-write (rw) is to be understood at a conceptual level. In NETCONF, for example, support for <candidate> and <startup> is optional and <running> does not have to be writable. Furthermore, <startup> can only be modified by copying <running> to <startup> in the standardized NETCONF datastore editing model. The RESTCONF protocol does not expose these differences and instead provides only a writable unified datastore, which hides whether edits are done through <candidate> or by directly modifying <running> or via some other implementation specific mechanism. RESTCONF also hides how configuration is made persistent. Note that implementations may also have additional datastores that can propagate changes to <running>. NETCONF explicitly mentions so called named datastores.

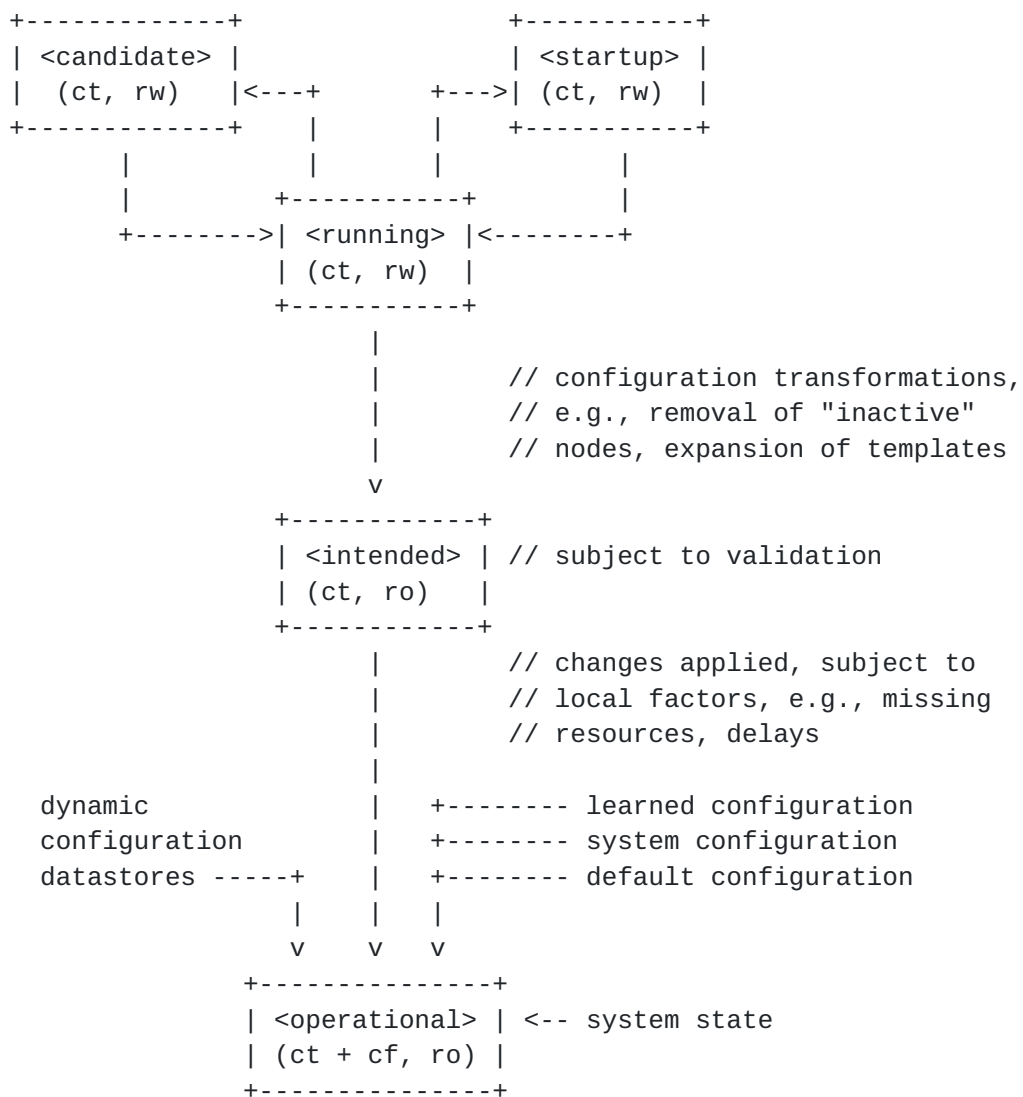
Some observations:

- o Operational state has not been defined as a datastore although there were proposals in the past to introduce an operational state datastore.

- o The NETCONF <get> operation returns the contents of <running> together with the operational state. It is therefore necessary that "config false" data is in a different branch than the "config true" data if the operational state can have a different lifetime compared to configuration or if configuration is not immediately or successfully applied.
- o Several implementations have proprietary mechanisms that allow clients to store inactive data in <running>. Inactive data is conceptually removed before validation.
- o Some implementations have proprietary mechanisms that allow clients to define configuration templates in <running>. These templates are expanded automatically by the system, and the resulting configuration is applied internally.
- o Some operators have reported that it is essential for them to be able to retrieve the configuration that has actually been successfully applied, which may be a subset or a superset of the <running> configuration.

5. Architectural Model of Datastores

Below is a new conceptual model of datastores extending the original model in order to reflect the experience gained with the original model.



```
ct = config true; cf = config false
rw = read-write; ro = read-only
boxes denote named datastores
```

5.1. Conventional Configuration Datastores

The conventional configuration datastores are a set of configuration datastores that share exactly the same datastore schema, allowing data to be copied between them. The term is meant as a generic umbrella description of these datastores. The set of datastores include:

- ```
0 <running>
0 <candidate>
```



- o <startup>
- o <intended>

Other conventional configuration datastores may be defined in future documents.

The flow of data between these datastores is depicted in [Section 5](#).

The specific protocols may define explicit operations to copy between these datastores, e.g., NETCONF defines the <copy-config> operation.

#### **[5.1.1](#). The Startup Configuration Datastore (<startup>)**

The startup configuration datastore (<startup>) is a configuration datastore holding the configuration loaded by the device when it boots. <startup> is only present on devices that separate the startup configuration from the running configuration datastore.

The startup configuration datastore may not be supported by all protocols or implementations.

On devices that support non-volatile storage, the contents of <startup> will typically persist across reboots via that storage. At boot time, the device loads the saved startup configuration into <running>. To save a new startup configuration, data is copied to <startup>, either via implicit or explicit protocol operations.

#### **[5.1.2](#). The Candidate Configuration Datastore (<candidate>)**

The candidate configuration datastore (<candidate>) is a configuration datastore that can be manipulated without impacting the device's current configuration and that can be committed to <running>.

The candidate configuration datastore may not be supported by all protocols or implementations.

<candidate> does not typically persist across reboots, even in the presence of non-volatile storage. If <candidate> is stored using non-volatile storage, it is reset at boot time to the contents of <running>.

#### **[5.1.3](#). The Running Configuration Datastore (<running>)**

The running configuration datastore (<running>) is a configuration datastore that holds the complete current configuration on the device. It MAY include configuration that requires further



transformation before it can be applied, e.g., inactive configuration, or template-mechanism-oriented configuration that needs further expansion. However, <running> MUST always be a valid configuration data tree, as defined in [Section 8.1 of \[RFC7950\]](#).

<running> MUST be supported if the device can be configured via conventional configuration datastores.

If a device does not have a distinct <startup> and non-volatile storage is available, the device will typically use that non-volatile storage to allow <running> to persist across reboots.

#### **[5.1.4. The Intended Configuration Datastore \(<intended>\)](#)**

The intended configuration datastore (<intended>) is a read-only configuration datastore. It represents the configuration after all configuration transformations to <running> are performed (e.g., template expansion, removal of inactive configuration), and is the configuration that the system attempts to apply.

<intended> is tightly coupled to <running>. Whenever data is written to <running>, then <intended> MUST also be immediately updated by performing all necessary configuration transformations to the contents of <running> and then <intended> is validated.

<intended> MAY also be updated independently of <running> if the effect of a configuration transformation changes, but <intended> MUST always be a valid configuration data tree, as defined in [Section 8.1 of \[RFC7950\]](#).

For simple implementations, <running> and <intended> are identical.

The contents of <intended> are also related to the "config true" subset of <operational>, and hence a client can determine to what extent the intended configuration is currently in use by checking whether the contents of <intended> also appear in <operational>.

<intended> does not persist across reboots; its relationship with <running> makes that unnecessary.

Currently there are no standard mechanisms defined that affect <intended> so that it would have different content than <running>, but this architecture allows for such mechanisms to be defined.

One example of such a mechanism is support for marking nodes as inactive in <running>. Inactive nodes are not copied to <intended>. A second example is support for templates, which can perform



transformations on the configuration from <running> to the configuration written to <intended>.

### **5.2. Dynamic Configuration Datastores**

The model recognizes the need for dynamic configuration datastores that are, by definition, not part of the persistent configuration of a device. In some contexts, these have been termed ephemeral datastores since the information is ephemeral, i.e., lost upon reboot. The dynamic configuration datastores interact with the rest of the system through <operational>.

### **5.3. The Operational State Datastore (<operational>)**

The operational state datastore (<operational>) is a read-only datastore that consists of all "config true" and "config false" nodes defined in the datastore's schema. In the original NETCONF model the operational state only had "config false" nodes. The reason for incorporating "config true" nodes here is to be able to expose all operational settings without having to replicate definitions in the data models.

<operational> contains system state and all configuration actually used by the system. This includes all applied configuration from <intended>, learned configuration, system-provided configuration, and default values defined by any supported data models. In addition, <operational> also contains applied configuration from dynamic configuration datastores.

The datastore schema for <operational> MUST be a superset of the combined datastore schema used in all configuration datastores except that YANG nodes supported in a configuration datastore MAY be omitted from <operational> if a server is not able to accurately report them.

Requests to retrieve nodes from <operational> always return the value in use if the node exists, regardless of any default value specified in the YANG module. If no value is returned for a given node, then this implies that the node is not used by the device.

The interpretation of what constitutes as being "in use" by the system is dependent on both the schema definition and the device implementation. Generally, functionality that is enabled and operational on the system would be considered as being "in use". Conversely, functionality that is neither enabled nor operational on the system is considered as not being "in use", and hence SHOULD be omitted from <operational>.





<operational> SHOULD conform to any constraints specified in the data model, but given the principal aim of returning "in use" values, it is possible that constraints MAY be violated under some circumstances, e.g., an abnormal value is "in use", the structure of a list is being modified, or due to remnant configuration (see [Section 5.3.1](#)). Note, that deviations SHOULD be used when it is known in advance that a device does not fully conform to the <operational> schema.

Only semantic constraints MAY be violated, these are the YANG "when", "must", "mandatory", "unique", "min-elements", and "max-elements" statements; and the uniqueness of key values.

Syntactic constraints MUST NOT be violated, including hierarchical organization, identifiers, and type-based constraints. If a node in <operational> does not meet the syntactic constraints then it MUST NOT be returned, and some other mechanism should be used to flag the error.

<operational> does not persist across reboots.

#### **[5.3.1.](#) Remnant Configuration**

Changes to configuration may take time to percolate through to <operational>. During this period, <operational> may contain nodes for both the previous and current configuration, as closely as possible tracking the current operation of the device. Such remnant configuration from the previous configuration persists until the system has released resources used by the newly-deleted configuration (e.g., network connections, memory allocations, file handles).

Remnant configuration is a common example of where the semantic constraints defined in the data model cannot be relied upon for <operational>, since the system may have remnant configuration whose constraints were valid with the previous configuration and that are not valid with the current configuration. Since constraints on "config false" nodes may refer to "config true" nodes, remnant configuration may force the violation of those constraints.

#### **[5.3.2.](#) Missing Resources**

Configuration in <intended> can refer to resources that are not available or otherwise not physically present. In these situations, these parts of <intended> are not applied. The data appears in <intended> but does not appear in <operational>.

A typical example is an interface configuration that refers to an interface that is not currently present. In such a situation, the



interface configuration remains in <intended> but the interface configuration will not appear in <operational>.

Note that configuration validity cannot depend on the current state of such resources, since that would imply that removing a resource might render the configuration invalid. This is unacceptable, especially given that rebooting such a device would cause it to restart with an invalid configuration. Instead we allow configuration for missing resources to exist in <running> and <intended>, but it will not appear in <operational>.

### **5.3.3. System-controlled Resources**

Sometimes resources are controlled by the device and the corresponding system controlled data appears in (and disappears from) <operational> dynamically. If a system controlled resource has matching configuration in <intended> when it appears, the system will try to apply the configuration, which causes the configuration to appear in <operational> eventually (if application of the configuration was successful).

### **5.3.4. Origin Metadata Annotation**

As configuration flows into <operational>, it is conceptually marked with a metadata annotation ([[RFC7952](#)]) that indicates its origin. The origin applies to all configuration nodes except non-presence containers. The "origin" metadata annotation is defined in [Section 7](#). The values are YANG identities. The following identities are defined:

- o origin: abstract base identity from which the other origin identities are derived.
- o intended: represents configuration provided by <intended>.
- o dynamic: represents configuration provided by a dynamic configuration datastore.
- o system: represents configuration provided by the system itself. Examples of system configuration include applied configuration for an always existing loopback interface, or interface configuration that is auto-created due to the hardware currently present in the device.
- o learned: represents configuration that has been learned via protocol interactions with other systems, including protocols such as link-layer negotiations, routing protocols, DHCP, etc.



- o `default`: represents configuration using a default value specified in the data model, using either values in the `"default"` statement or any values described in the `"description"` statement. The default origin is only used when the configuration has not been provided by any other source.
- o `unknown`: represents configuration for which the system cannot identify the origin.

These identities can be further refined, e.g., there could be separate identities for particular types or instances of dynamic configuration datastores derived from `"dynamic"`.

For all configuration data nodes in `<operational>`, the device **SHOULD** report the origin that most accurately reflects the source of the configuration that is in use by the system.

In cases where it could be ambiguous as to which origin should be used, i.e. where the same data node value has originated from multiple sources, then the description statement in the YANG module **SHOULD** be used as guidance for choosing the appropriate origin. For example:

If for a particular configuration node, the associated YANG description statement indicates that a protocol negotiated value overrides any configured value, then the origin would be reported as `"learned"`, even when a learned value is the same as the configured value.

Conversely, if for a particular configuration node, the associated YANG description statement indicates that a protocol negotiated value does not override an explicitly configured value, then the origin would be reported as `"intended"` even when a learned value is the same as the configured value.

In the case that a device cannot provide an accurate origin for a particular configuration data node then it **SHOULD** use the origin `"unknown"`.

## **6. Implications on YANG**

### **6.1. XPath Context**

This section updates [section 6.4.1 of RFC 7950](#).

If a server implements the architecture defined in this document, the accessible trees for some XPath contexts are refined as follows:



- o If the XPath expression is defined in a substatement to a data node that represents system state, the accessible tree is all operational state in the server. The root node has all top-level data nodes in all modules as children.
- o If the XPath expression is defined in a substatement to a "notification" statement, the accessible tree is the notification instance and all operational state in the server. If the notification is defined on the top level in a module, then the root node has the node representing the notification being defined and all top-level data nodes in all modules as children. Otherwise, the root node has all top-level data nodes in all modules as children.
- o If the XPath expression is defined in a substatement to an "input" statement in an "rpc" or "action" statement, the accessible tree is the RPC or action operation instance and all operational state in the server. The root node has top-level data nodes in all modules as children. Additionally, for an RPC, the root node also has the node representing the RPC operation being defined as a child. The node representing the operation being defined has the operation's input parameters as children.
- o If the XPath expression is defined in a substatement to an "output" statement in an "rpc" or "action" statement, the accessible tree is the RPC or action operation instance and all operational state in the server. The root node has top-level data nodes in all modules as children. Additionally, for an RPC, the root node also has the node representing the RPC operation being defined as a child. The node representing the operation being defined has the operation's output parameters as children.

## **6.2. Invocation of Actions and RPCs**

This section updates [section 7.15 of RFC 7950](#).

Actions are always invoked in the context of the operational state datastore. The node for which the action is invoked MUST exist in the operational state datastore.

Note that this document does not constrain the result of invoking an RPC or action in any way. For example, an RPC might be defined to modify the contents of some datastore.





## 7. YANG Modules

```
<CODE BEGINS> file "ietf-datastores@2017-08-17.yang"

module ietf-datastores {
 yang-version 1.1;
 namespace "urn:ietf:params:xml:ns:yang:ietf-datastores";
 prefix ds;

 organization
 "IETF Network Modeling (NETMOD) Working Group";

 contact
 "WG Web: <https://datatracker.ietf.org/wg/netmod/>

 WG List: <mailto:netmod@ietf.org>

 Author: Martin Bjorklund
 <mailto:mbj@tail-f.com>

 Author: Juergen Schoenwaelder
 <mailto:j.schoenwaelder@jacobs-university.de>

 Author: Phil Shafer
 <mailto:phil@juniper.net>

 Author: Kent Watsen
 <mailto:kwatsen@juniper.net>

 Author: Rob Wilton
 <rwilton@cisco.com>";

 description
 "This YANG module defines two sets of identities for datastores.
 The first identifies the datastores themselves, the second
 identifies datastore properties.

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 authors of the code. All rights reserved.

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 forth in Section 4.c of the IETF Trust's Legal Provisions
 Relating to IETF Documents
 (http://trustee.ietf.org/license-info).

 This version of this YANG module is part of RFC XXXX
```



(<http://www.rfc-editor.org/info/rfcxxxx>); see the RFC itself for full legal notices.";

```
revision 2017-08-17 {
 description
 "Initial revision."
 reference
 "RFC XXXX: Network Management Datastore Architecture";
}
```

```
/*
 * Identities
 */
```

```
identity datastore {
 description
 "Abstract base identity for datastore identities."
}
```

```
identity conventional {
 base datastore;
 description
 "Abstract base identity for conventional configuration
 datastores."
}
```

```
identity running {
 base conventional;
 description
 "The running configuration datastore."
}
```

```
identity candidate {
 base conventional;
 description
 "The candidate configuration datastore."
}
```

```
identity startup {
 base conventional;
 description
 "The startup configuration datastore."
}
```

```
identity intended {
 base conventional;
 description
 "The intended configuration datastore."
}
```



```
}

identity dynamic {
 base datastore;
 description
 "Abstract base identity for dynamic configuration datastores.";
}

identity operational {
 base datastore;
 description
 "The operational state datastore.";
}

/*
 * Type definitions
 */

typedef datastore-ref {
 type identityref {
 base datastore;
 }
 description
 "A datastore identity reference.";
}

}

<CODE ENDS>

<CODE BEGINS> file "ietf-origin@2017-08-17.yang"

module ietf-origin {
 yang-version 1.1;
 namespace "urn:ietf:params:xml:ns:yang:ietf-origin";
 prefix or;

 import ietf-yang-metadata {
 prefix md;
 }

 organization
 "IETF Network Modeling (NETMOD) Working Group";

 contact
 "WG Web: <https://datatracker.ietf.org/wg/netmod/>

 WG List: <mailto:netmod@ietf.org>
```



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description

"This YANG module defines an 'origin' metadata annotation, and a set of identities for the origin value.

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This version of this YANG module is part of RFC XXXX (<http://www.rfc-editor.org/info/rfcxxxx>); see the RFC itself for full legal notices.";

```
revision 2017-08-17 {
 description
 "Initial revision.";
 reference
 "RFC XXXX: Network Management Datastore Architecture";
}

/*
 * Identities
 */

identity origin {
 description
 "Abstract base identity for the origin annotation.";
}
```





```
identity intended {
 base origin;
 description
 "Denotes configuration from the intended configuration
 datastore";
}

identity dynamic {
 base origin;
 description
 "Denotes configuration from a dynamic configuration
 datastore.";
}

identity system {
 base origin;
 description
 "Denotes configuration originated by the system itself.

 Examples of system configuration include applied configuration
 for an always existing loopback interface, or interface
 configuration that is auto-created due to the hardware
 currently present in the device.";
}

identity learned {
 base origin;
 description
 "Denotes configuration learned from protocol interactions with
 other devices, instead of via either the intended
 configuration datastore or any dynamic configuration
 datastore.

 Examples of protocols that provide learned configuration
 include link-layer negotiations, routing protocols, and
 DHCP.";
}

identity default {
 base origin;
 description
 "Denotes configuration that does not have an configured or
 learned value, but has a default value in use. Covers both
 values defined in a 'default' statement, and values defined
 via an explanation in a 'description' statement.";
}

identity unknown {
```



```
 base origin;
 description
 "Denotes configuration for which the system cannot identify the
 origin.";
}

/*
 * Type definitions
 */

typedef origin-ref {
 type identityref {
 base origin;
 }
 description
 "An origin identity reference.";
}

/*
 * Metadata annotations
 */

md:annotation origin {
 type origin-ref;
 description
 "The 'origin' annotation can be present on any configuration
 data node in the operational datastore. It specifies from
 where the node originated. If not specified for a given
 configuration data node then the origin is the same as the
 origin of its parent node in the data tree. The origin for
 any top level configuration data nodes must be specified.";
}
}

<CODE ENDS>
```

## **8. IANA Considerations**

### **8.1. Updates to the IETF XML Registry**

This document registers two URIs in the IETF XML registry [[RFC3688](#)]. Following the format in [[RFC3688](#)], the following registrations are requested:



URI: urn:ietf:params:xml:ns:yang:ietf-datastores  
Registrant Contact: The IESG.  
XML: N/A, the requested URI is an XML namespace.

URI: urn:ietf:params:xml:ns:yang:ietf-origin  
Registrant Contact: The IESG.  
XML: N/A, the requested URI is an XML namespace.

## **8.2. Updates to the YANG Module Names Registry**

This document registers two YANG modules in the YANG Module Names registry [[RFC6020](#)]. Following the format in [[RFC6020](#)], the the following registrations are requested:

|            |                                             |
|------------|---------------------------------------------|
| name:      | ietf-datastores                             |
| namespace: | urn:ietf:params:xml:ns:yang:ietf-datastores |
| prefix:    | ds                                          |
| reference: | RFC XXXX                                    |
| name:      | ietf-origin                                 |
| namespace: | urn:ietf:params:xml:ns:yang:ietf-origin     |
| prefix:    | or                                          |
| reference: | RFC XXXX                                    |

## **9. Security Considerations**

This document discusses an architectural model of datastores for network management using NETCONF/RESTCONF and YANG. It has no security impact on the Internet.

Although this document specifies several YANG modules, these modules only define identities and meta-data, hence the "YANG module security guidelines" do not apply.

## **10. Acknowledgments**

This document grew out of many discussions that took place since 2010. Several Internet-Drafts ([[I-D.bjorklund-netmod-operational](#)], [[I-D.wilton-netmod-opstate-yang](#)], [[I-D.ietf-netmod-opstate-reqs](#)], [[I-D.kwatsen-netmod-opstate](#)], [[I-D.openconfig-netmod-opstate](#)]) and [[RFC6244](#)] touched on some of the problems of the original datastore model. The following people were authors to these Internet-Drafts or otherwise actively involved in the discussions that led to this document:

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## **Appendix A. Guidelines for Defining Datastores**

The definition of a new datastore in this architecture should be provided in a document (e.g., an RFC) purposed to the definition of the datastore. When it makes sense, more than one datastore may be defined in the same document (e.g., when the datastores are logically connected). Each datastore's definition should address the points specified in the sections below.

### **A.1. Define which YANG modules can be used in the datastore**

Not all YANG modules may be used in all datastores. Some datastores may constrain which data models can be used in them. If it is desirable that a subset of all modules can be targeted to the datastore, then the documentation defining the datastore must indicate this.

### **A.2. Define which subset of YANG-modeled data applies**

By default, the data in a datastore is modeled by all YANG statements in the available YANG modules. However, it is possible to specify criteria that YANG statements must satisfy in order to be present in a datastore. For instance, maybe only "config true" nodes, or "config false" nodes that also have a specific YANG extension, are present in the datastore.

### **A.3. Define how data is actualized**

The new datastore must specify how it interacts with other datastores.

For example, the diagram in [Section 5](#) depicts dynamic configuration datastores feeding into <operational>. How this interaction occurs has to be defined by the particular dynamic configuration datastores. In some cases, it may occur implicitly, as soon as the data is put into the dynamic configuration datastore while, in other cases, an explicit action (e.g., an RPC) may be required to trigger the application of the datastore's data.



#### **[A.4.](#) Define which protocols can be used**

By default, it is assumed that both the NETCONF and RESTCONF protocols can be used to interact with a datastore. However, it may be that only a specific protocol can be used (e.g., ForCES) or that a subset of all protocol operations or capabilities are available (e.g., no locking or no XPath-based filtering).

#### **[A.5.](#) Define YANG identities for the datastore**

The datastore must be defined with a YANG identity that uses the "ds:datastore" identity, or one of its derived identities, as its base. This identity is necessary so that the datastore can be referenced in protocol operations (e.g., <get-data>).

The datastore may also be defined with an identity that uses the "or:origin" identity or one of its derived identities as its base. This identity is needed if the datastore interacts with <operational> so that data originating from the datastore can be identified as such via the "origin" metadata attribute defined in [Section 7](#).

An example of these guidelines in use is provided in [Appendix B](#).

#### **[Appendix B.](#) Ephemeral Dynamic Configuration Datastore Example**

The section defines documentation for an example dynamic configuration datastore using the guidelines provided in [Appendix A](#). While this example is very terse, it is expected to be that a standalone RFC would be needed when fully expanded.

This example defines a dynamic configuration datastore called "ephemeral", which is loosely modeled after the work done in the I2RS working group.

| +-----+-----+ |                                                   |
|---------------|---------------------------------------------------|
| Name          | Value                                             |
| +-----+-----+ |                                                   |
| Name          | ephemeral                                         |
| YANG modules  | all (default)                                     |
| YANG nodes    | all "config true" data nodes                      |
| How applied   | changes automatically propagated to <operational> |
| Protocols     | NC/RC (default)                                   |
| YANG Module   | (see below)                                       |
| +-----+-----+ |                                                   |

The example "ephemeral" datastore properties



```
module example-ds-ephemeral {
 yang-version 1.1;
 namespace "urn:example:ds-ephemeral";
 prefix eph;

 import ietf-datastores {
 prefix ds;
 }
 import ietf-origin {
 prefix or;
 }

 // datastore identity
 identity ds-ephemeral {
 base ds:dynamic;
 description
 "The ephemeral dynamic configuration datastore.";
 }

 // origin identity
 identity or-ephemeral {
 base or:dynamic;
 description
 "Denotes data from the ephemeral dynamic configuration
 datastore.";
 }
}
```

## [Appendix C](#). Example Data

The use of datastores is complex, and many of the subtle effects are more easily presented using examples. This section presents a series of example data models with some sample contents of the various datastores.

### [C.1](#). System Example

In this example, the following fictional module is used:

```
module example-system {
 yang-version 1.1;
 namespace urn:example:system;
 prefix sys;

 import ietf-inet-types {
 prefix inet;
 }
}
```





```
 container system {
 leaf hostname {
 type string;
 }

 list interface {
 key name;

 leaf name {
 type string;
 }

 container auto-negotiation {
 leaf enabled {
 type boolean;
 default true;
 }
 leaf speed {
 type uint32;
 units mbps;
 description
 "The advertised speed, in mbps.";
 }
 }
 }

 leaf speed {
 type uint32;
 units mbps;
 config false;
 description
 "The speed of the interface, in mbps.";
 }

 list address {
 key ip;

 leaf ip {
 type inet:ip-address;
 }
 leaf prefix-length {
 type uint8;
 }
 }
 }
 }
}
```



The operator has configured the host name and two interfaces, so the contents of <intended> are:

```
<system xmlns="urn:example:system">

 <hostname>foo</hostname>

 <interface>
 <name>eth0</name>
 <auto-negotiation>
 <speed>1000</speed>
 </auto-negotiation>
 <address>
 <ip>2001:db8::10</ip>
 <prefix-length>64</prefix-length>
 </address>
 </interface>

 <interface>
 <name>eth1</name>
 <address>
 <ip>2001:db8::20</ip>
 <prefix-length>64</prefix-length>
 </address>
 </interface>

</system>
```

The system has detected that the hardware for one of the configured interfaces ("eth1") is not yet present, so the configuration for that interface is not applied. Further, the system has received a host name and an additional IP address for "eth0" over DHCP. In addition to a default value, a loopback interface is automatically added by the system, and the result of the "speed" auto-negotiation. All of this is reflected in <operational>. Note how the origin metadata attribute for several "config true" data nodes is inherited from their parent data nodes.



```
<system
 xmlns="urn:example:system"
 xmlns:or="urn:ietf:params:xml:ns:yang:ietf-origin">

 <hostname or:origin="or:dynamic">bar</hostname>

 <interface or:origin="or:intended">
 <name>eth0</name>
 <auto-negotiation>
 <enabled or:origin="or:default">true</enabled>
 <speed>1000</speed>
 </auto-negotiation>
 <speed>100</speed>
 <address>
 <ip>2001:db8::10</ip>
 <prefix-length>64</prefix-length>
 </address>
 <address or:origin="or:dynamic">
 <ip>2001:db8::1:100</ip>
 <prefix-length>64</prefix-length>
 </address>
 </interface>

 <interface or:origin="or:system">
 <name>lo0</name>
 <address>
 <ip>::1</ip>
 <prefix-length>128</prefix-length>
 </address>
 </interface>

</system>
```

## **C.2. BGP Example**

Consider the following fragment of a fictional BGP module:



```
container bgp {
 leaf local-as {
 type uint32;
 }
 leaf peer-as {
 type uint32;
 }
 list peer {
 key name;
 leaf name {
 type inet:ip-address;
 }
 leaf local-as {
 type uint32;
 description
 ".... Defaults to ../local-as";
 }
 leaf peer-as {
 type uint32;
 description
 "... Defaults to ../peer-as";
 }
 leaf local-port {
 type inet:port;
 }
 leaf remote-port {
 type inet:port;
 default 179;
 }
 leaf state {
 config false;
 type enumeration {
 enum init;
 enum established;
 enum closing;
 }
 }
 }
}
```

In this example model, both `bgp/peer/local-as` and `bgp/peer/peer-as` have complex hierarchical values, allowing the user to specify default values for all peers in a single location.

The model also follows the pattern of fully integrating state ("config false") nodes with configuration ("config true") nodes. There is no separate "bgp-state" hierarchy, with the accompanying





repetition of containment and naming nodes. This makes the model simpler and more readable.

### **C.2.1. Datastores**

Each datastore represents differing views of these nodes. <running> will hold the configuration provided by the operator, for example a single BGP peer. <intended> will conceptually hold the data as validated, after the removal of data not intended for validation and after any local template mechanisms are performed. <operational> will show data from <intended> as well as any "config false" nodes.

### **C.2.2. Adding a Peer**

If the user configures a single BGP peer, then that peer will be visible in both <running> and <intended>. It may also appear in <candidate>, if the server supports the candidate configuration datastore. Retrieving the peer will return only the user-specified values.

No time delay should exist between the appearance of the peer in <running> and <intended>.

In this scenario, we've added the following to <running>:

```
<bgp>
 <local-as>64501</local-as>
 <peer-as>64502</peer-as>
 <peer>
 <name>10.1.2.3</name>
 </peer>
</bgp>
```

#### **C.2.2.1. <operational>**

The operational datastore will contain the fully expanded peer data, including "config false" nodes. In our example, this means the "state" node will appear.

In addition, <operational> will contain the "currently in use" values for all nodes. This means that local-as and peer-as will be populated even if they are not given values in <intended>. The value of bgp/local-as will be used if bgp/peer/local-as is not provided; bgp/peer-as and bgp/peer/peer-as will have the same relationship. In the operational view, this means that every peer will have values for their local-as and peer-as, even if those values are not explicitly configured but are provided by bgp/local-as and bgp/peer-as.



Each BGP peer has a TCP connection associated with it, using the values of local-port and remote-port from <intended>. If those values are not supplied, the system will select values. When the connection is established, <operational> will contain the current values for the local-port and remote-port nodes regardless of the origin. If the system has chosen the values, the "origin" attribute will be set to "system". Before the connection is established, one or both of the nodes may not appear, since the system may not yet have their values.

```
<bgp or:origin="or:intended">
 <local-as>64501</local-as>
 <peer-as>64502</peer-as>
 <peer>
 <name>10.1.2.3</name>
 <local-as or:origin="or:default">64501</local-as>
 <peer-as or:origin="or:default">64502</peer-as>
 <local-port or:origin="or:system">60794</local-port>
 <remote-port or:origin="or:default">179</remote-port>
 <state>established</state>
 </peer>
</bgp>
```

### **C.2.3. Removing a Peer**

Changes to configuration may take time to percolate through the various software components involved. During this period, it is imperative to continue to give an accurate view of the working of the device. <operational> will contain nodes for both the previous and current configuration, as closely as possible tracking the current operation of the device.

Consider the scenario where a client removes a BGP peer. When a peer is removed, the operational state will continue to reflect the existence of that peer until the peer's resources are released, including closing the peer's connection. During this period, the current data values will continue to be visible in <operational>, with the "origin" attribute set to indicate the origin of the original data.



```
<bgp or:origin="or:intended">
 <local-as>64501</local-as>
 <peer-as>64502</peer-as>
 <peer>
 <name>10.1.2.3</name>
 <local-as or:origin="or:default">64501</local-as>
 <peer-as or:origin="or:default">64502</peer-as>
 <local-port or:origin="or:system">60794</local-port>
 <remote-port or:origin="or:default">179</remote-port>
 <state>closing</state>
 </peer>
</bgp>
```

Once resources are released and the connection is closed, the peer's data is removed from <operational>.

### **[C.3.](#) Interface Example**

In this section, we will use this simple interface data model:

```
container interfaces {
 list interface {
 key name;
 leaf name {
 type string;
 }
 leaf description {
 type string;
 }
 leaf mtu {
 type uint16;
 }
 leaf-list ip-address {
 type inet:ip-address;
 }
 }
}
```

#### **[C.3.1.](#) Pre-provisioned Interfaces**

One common issue in networking devices is the support of Field Replaceable Units (FRUs) that can be inserted and removed from the device without requiring a reboot or interfering with normal operation. These FRUs are typically interface cards, and the devices support pre-provisioning of these interfaces.



If a client creates an interface "et-0/0/0" but the interface does not physically exist at this point, then <intended> might contain the following:

```
<interfaces>
 <interface>
 <name>et-0/0/0</name>
 <description>Test interface</description>
 </interface>
</interfaces>
```

Since the interface does not exist, this data does not appear in <operational>.

When a FRU containing this interface is inserted, the system will detect it and process the associated configuration. <operational> will contain the data from <intended>, as well as nodes added by the system, such as the current value of the interface's MTU.

```
<interfaces or:origin="or:intended">
 <interface>
 <name>et-0/0/0</name>
 <description>Test interface</description>
 <mtu or:origin="or:system">1500</mtu>
 </interface>
</interfaces>
```

If the FRU is removed, the interface data is removed from <operational>.

### **C.3.2. System-provided Interface**

Imagine if the system provides a loopback interface (named "lo0") with a default ip-address of "127.0.0.1" and a default ip-address of "::1". The system will only provide configuration for this interface if there is no data for it in <intended>.

When no configuration for "lo0" appears in <intended>, then <operational> will show the system-provided data:

```
<interfaces or:origin="or:intended">
 <interface or:origin="or:system">
 <name>lo0</name>
 <ip-address>127.0.0.1</ip-address>
 <ip-address>::1</ip-address>
 </interface>
</interfaces>
```





When configuration for "lo0" does appear in <intended>, then <operational> will show that data with the origin set to "intended". If the "ip-address" is not provided, then the system-provided value will appear as follows:

```
<interfaces or:origin="or:intended">
 <interface>
 <name>lo0</name>
 <description>loopback</description>
 <ip-address or:origin="or:system">127.0.0.1</ip-address>
 <ip-address>::1</ip-address>
 </interface>
</interfaces>
```

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