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Extending the Opening of Files in NFSv4.2
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Abstract

The Network File System v4 (NFSv4) allows a client to both open a file and be granted a delegation of that file. This provides the client the right to cache metadata on the file locally. This document presents several refinements to both the opening and delegating of the file to the client.

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[1.](#) Introduction

In the Network File System version4 (NFSv4) a client may be granted a delegation for a file. This allows the client to act as the authority of the file's metadata and data. In this document, we introduce some new semantics to both the open and the delegation process which allows the client to:

- o detect an offline file, which may be located off premise.
- o determine the extension of OPEN (see [Section 18.16 of \[RFC5661\]](#)) flags.
- o during the OPEN procedure, get either the open or delegation stateids, but not both.
- o cache both the access and modify times, reducing the number of times the client needs to go to the server to get that information.
- o for clients using Parallel NFS (pNFS) (see [Section 12 of \[RFC5661\]](#)), periodically report the attributes of the data files to the metadata server.

Using the process detailed in [[RFC8178](#)], the revisions in this document become an extension of NFSv4.2 [[RFC7862](#)]. They are built on top of the external data representation (XDR) [[RFC4506](#)] generated from [[RFC7863](#)].

1.1. Definitions

delegation: A file delegation, which is a recall-able lock that assures the holder that inconsistent opens and file changes cannot occur so long as the delegation is held.

stateid: A stateid is a 128-bit quantity returned by a server that uniquely defines state held by the server for the client. (See [Section 8 of \[RFC5661\]](#))

weak cache consistency (WCC): In NFSv3, operations are not sent in a compound, hence the client would have to perform two round trips to the server in order to determine the result of modification to the state of a file or directory. With WCC, the server can return post-operation attributes on such operations. As these do not provide a strict consistency between the server and client, the client is free to ignore the data. (See [Section 2.6 of \[RFC1813\]](#))

1.2. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [[RFC2119](#)].

2. Offline Files

If a file is offline, then the server locally has the file's attributes, but not the file's content. It has to be able to present to the client enough information to describe the file, but the content is not readily available. The cost of retrieving the data content is expensive, to the extent that the content should only be retrieved if it is going to be used. A graphical file manager (such as OSX's Finder) may want to access the beginning of the file to preview it for an user who is hovering his pointer over the file name. If the file is retrieved, it will most likely either be immediately thrown away or returned.

A compound with a GETATTR or READDIR can report the file's attributes without bringing the file online. However, either an OPEN or a LAYOUTGET might cause the file server to retrieve the archived data contents, bringing the file online. For non-pNFS systems, the OPEN operation requires a filehandle to the data content. For pNFS systems, the filehandle retrieved from an OPEN need not cause the

data content to be retrieved. But when the LAYOUTGET operation is processed, a layout type specific mapping will cause the data content to be retrieved from offline storage.

If an operating system is not aware that the file is offline, it might inadvertently open the file to determine what type of file it is accessing. By adding the new attribute FATTR4_OFFLINE, a client can predetermine the availability of the file, avoiding the need to open it at all. Being offline might also mean that the file is archived in the cloud, i.e., there can be an expense in both retrieving the file to bring online and in sending the file back to offline status.

<CODE BEGINS>

```
///  
/// typedef bool          fattr4_offline;  
///  
  
///  
/// const FATTR4_OFFLINE      = 83;  
///
```

<CODE ENDS>

3. Determining the Arguments to OPEN

The OPEN (See [Section 18.16 of \[RFC5661\]](#)) procedure returns an open stateid to the client to reference the state of the file. The client could also request a delegation stateid in the OPEN arguments. The file is said to be "open" to the client as long as the count of open and delegated stateids is greater than 0. Either type of stateid is sufficient to keep the file open, which allows READ (See [Section 18.22 of \[RFC5661\]](#)), WRITE (See [Section 18.2 of \[RFC5661\]](#)), LOCK (See [Section 18.10 of \[RFC5661\]](#)), and LAYOUTGET (see [Section 18.43 of \[RFC5661\]](#)) operations to proceed. If the client gets both a open and a delegation stateid as part of the OPEN, then it has to return them both. And during each operation, the client can send a costly GETATTR (See [Section 18.7 of \[RFC5661\]](#)).

If the client knows that the server supports the OPEN4_SHARE_ACCESS_WANT_OPEN_XOR_DELEGATION flag (as determined by an earlier GETATTR operation which queried for the FATTR4_OPEN_ARGUMENTS attribute), then the client can supply that flag during the OPEN and only get either an open or delegation stateid.

The client is already prepared to not get a delegation stateid even if requested. In order to not send an open stateid, the server can

indicate that fact with the result flag of OPEN4_RESULT_NO_OPEN_STATEID. The open stateid field, OPEN4resok.stateid (see [Section 18.16.2 of \[RFC5661\]](#)), should also be set to the special all zero stateid.

3.1. XDR Modifications to OPEN

[RFC8178] (see [Section 4.4.2](#)) allows for extending the microversion of the NFSv4.x protocol without increasing the microversion. The client can probe the capabilities of the server and based on that result, determine if both it and the server support features not specified in the main microversion document.

The XDR extensions presented in this section allow for the OPEN procedure to be extended in such a fashion. It models all of the parameters via bitmap4 data structures, which allows for the addition of a new flag to any of the OPEN arguments (see [Section 18.16.1 of \[RFC5661\]](#)). Two new flags are provided:

- o OPEN4_SHARE_ACCESS_WANT_OPEN_XOR_DELEGATION (see [Section 4](#))
- o OPEN4_SHARE_ACCESS_WANT_DELEG_TIMESTAMPS

Subsequent documents can use this framework to introduce new functionality to OPEN.

<CODE BEGINS>

```
///  
/// struct open_arguments4 {  
///     bitmap4  oa_share_access;  
///     bitmap4  oa_share_deny;  
///     bitmap4  oa_share_access_want;  
///     bitmap4  oa_open_claim;  
///     bitmap4  oa_create_mode;  
/// };  
///  
  
///  
/// enum open_args_share_access4 = {  
///     OPEN_ARGS_SHARE_ACCESS_READ  = 0;  
///     OPEN_ARGS_SHARE_ACCESS_WRITE = 1;  
///     OPEN_ARGS_SHARE_ACCESS_BOTH  = 2;  
/// };  
///
```



```
///  
/// enum open_args_share_deny4 = {  
///     OPEN_ARGS_SHARE_DENY_NONE    = 0;  
///     OPEN_ARGS_SHARE_DENY_READ    = 1;  
///     OPEN_ARGS_SHARE_DENY_WRITE    = 2;  
/// };  
///  
  
///  
/// enum open_args_share_access4 = {  
///     OPEN_ARGS_SHARE_ACCESS_WANT_ANY_DELEG        = 0;  
///     OPEN_ARGS_SHARE_ACCESS_WANT_NO_DELEG         = 1;  
///     OPEN_ARGS_SHARE_ACCESS_WANT_CANCEL           = 2;  
///     OPEN_ARGS_SHARE_ACCESS_WANT_SIGNAL_DELEG_WHEN_RESRC_AVAIL  
///                                                  = 3;  
///     OPEN_ARGS_SHARE_ACCESS_WANT_PUSH_DELEG_WHEN_UNCONTENDED  
///                                                  = 4;  
///     OPEN_ARGS_SHARE_ACCESS_WANT_DELEG_TIMESTAMPS = 5;  
///     OPEN_ARGS_SHARE_ACCESS_WANT_OPEN_XOR_DELEGATION = 6;  
/// };  
///  
  
///  
/// enum open_args_share_access4 = {  
///     OPEN_ARGS_CLAIM_NULL            = 0;  
///     OPEN_ARGS_CLAIM_PREVIOUS        = 1;  
///     OPEN_ARGS_CLAIM_DELEGATE_CUR    = 2;  
///     OPEN_ARGS_CLAIM_DELEGATE_PREV   = 3;  
///     OPEN_ARGS_CLAIM_FH              = 4;  
///     OPEN_ARGS_CLAIM_DELEG_CUR_FH    = 5;  
///     OPEN_ARGS_CLAIM_DELEG_PREV_FH   = 6;  
/// };  
///  
  
///  
/// enum open_args_share_access4 = {  
///     OPEN_ARGS_CREATE_MODE_GUARDED    = 0;  
///     OPEN_ARGS_CREATE_MODE_EXCLUSIVE  = 1;  
/// };  
///  
  
///  
/// typedef open_arguments4 fattr4_open_arguments;  
///
```



```
///  
/// %/*  
/// % * Determine what OPEN4 supports.  
/// % */  
/// const FATTR4_OPEN_ARGUMENTS      = 86;  
///  
  
///  
/// const OPEN4_SHARE_ACCESS_WANT_OPEN_XOR_DELEGATION = 0x200000;  
///  
  
///  
/// const OPEN4_RESULT_NO_OPEN_STATEID = 0x00000010;  
///  
  
<CODE ENDS>
```

4. Proxying of Times

When a client is granted a write delegation on a file, it is the authority for the file. If the server queries the client as to the state of the file via a CB_GETATTR (see [Section 20.1 of \[RFC5661\]](#)), then it can only determine the size of the file. Likewise, if the client holding the delegation wants to know either of the access, modify, or change times, it has to send a GETATTR to the server. While it is the authority for these values, it has no way to guarantee these values after the delegation has been returned. And as such, it can not pass these times up to an application expecting posix compliance.

With the addition of the new flag:

OPEN4_SHARE_ACCESS_WANT_DELEG_TIMESTAMPS, the client and server can negotiate that the client will be the authority for these values and upon return of the delegation stateid via a DELEGRETURN (see [section 18.6 of \[RFC5661\]](#)), the times will be passed back to the server. If the server is queried by another client for either the size or the times, it will need to use a CB_GETATTR to query the client which holds the delegation (see [Section 20.1 of \[RFC5661\]](#)).

If a server informs the client via the FATTR4_OPEN_ARGUMENTS attribute that it supports

OPEN_ARGS_SHARE_ACCESS_WANT_DELEG_TIMESTAMPS and it returns a valid delegation stateid for an OPEN operation which sets the OPEN4_SHARE_ACCESS_WANT_DELEG_TIMESTAMPS flag, then it MUST be able to query the client via a CB_GETATTR for the FATTR4_TIME_DELEG_ACCESS attribute and FATTR4_TIME_DELEG_MODIFY attribute. (The change time can be derived from the modify time.) Further, when it gets a SETATTR (see [Section 18.30 of \[RFC5661\]](#)) in the same compound as the

DELEGRETURN, then it MUST accept those FATTR4_TIME_DELEG_ACCESS attribute and FATTR4_TIME_DELEG_MODIFY attribute changes and derive the change time or reject the changes with NFS4ERR_DELAY.

A key prerequisite of this approach is that the server and client are in time synchronization with each other. Note that while the base NFSv4.2 does not require such synchronization, the use of RPCSEC_GSS typically makes such a requirement. When the client presents either FATTR4_TIME_DELEG_ACCESS or FATTR4_TIME_DELEG_MODIFY attributes to the server, the server MUST decide whether the times presented are before the old times or past the current time. If the time presented is before the original time, then the update is ignored. If the time presented is in the future, the server can either clamp the new time to the current time, or it may return NFS4ERR_DELAY to the client, allowing it to retry. Note that if the clock skew is large, this policy will result in access to the file being denied until such time that the clock skew is exceeded.

A change in the access time MUST not advance the change time, also known as the time_metadata attribute (see [Section 5.8.2.42 of \[RFC5661\]](#)), but a change in the modify time might advance the change time (and in turn the change attribute (See [Section 5.8.1.4 of \[RFC5661\]](#)). If the modify time is greater than the change time and before the current time, then the change time is adjusted to the modify time and not the current time (as is most likely done on most SETATTR calls that change the metadata). If the modify time is in the future, it will be clamped to the current time.

Note that each of the possible times, access, modify, and change, are compared to the current time. They should all be compared against the same time value for the current time. I.e., do not retrieve a different value of the current time for each calculation.

If the client sets the OPEN4_SHARE_ACCESS_WANT_DELEG_TIMESTAMPS flag in an OPEN operation, then it MUST support the FATTR4_TIME_DELEG_ACCESS and FATTR4_TIME_DELEG_MODIFY attributes both in the CB_GETATTR and SETATTR operations.

[4.1.](#) Use case

When a server is a proxy for a NFSv4 server, it is a client to the NFSv4 server and during file I/O, it may get a delegation on a file. The client of the proxy would be querying the proxy for attributes and not the NFSv4 server. Each GETATTR from that client would result in at least one additional GETATTR being sent across the wire.

4.2. XDR for Proxying of Times

```

<CODE BEGINS>

///
/// /*
///  * attributes for the delegation times being
///  * cached and served by the "client"
///  */
/// typedef nfstime4      fattr4_time_deleg_access;
/// typedef nfstime4      fattr4_time_deleg_modify;
///

///
/// %/*
/// % * New RECOMMENDED Attribute for
/// % * delegation caching of times
/// % */
/// const FATTR4_TIME_DELEG_ACCESS  = 84;
/// const FATTR4_TIME_DELEG_MODIFY  = 85;
///

///
/// const OPEN4_SHARE_ACCESS_WANT_DELEG_TIMESTAMPS = 0x100000;
///

<CODE ENDS>

```

5. Operation 77: LAYOUT_WCC - Layout Weak Cache Consistency

5.1. ARGUMENT

```

<CODE BEGINS>
/// struct LAYOUT_WCC4args {
///     stateid4      lowa_stateid;
///     layouttype4   lowa_type;
///     opaque         lowa_body<>;
/// };
<CODE ENDS>

```

5.2. RESULT

```

<CODE BEGINS>
/// struct LAYOUT_WCC4res {
///     nfsstat4      lowr_status;
/// };
<CODE ENDS>

```


5.3. DESCRIPTION

When using pNFS (See [Section 12 of \[RFC5661\]](#)), the client is most likely to be performing file operations to the storage device and not the metadata server. With some layout types (most notably the flexible files layout type in [\[RFC8435\]](#)) there is no control protocol between the metadata server and the storage device. In order to update the metadata state of the file, the metadata server will need to track the metadata state of the data file - once the layout is issued, it is not able to see the NFSv3 file operations from the client to the storage device. Thus the metadata server will be required to query the storage device for the data file attributes.

For example, with a flexible files layout type, the metadata server would issue a NFSv3 GETATTR to the storage device. These queries are most likely triggered in response to a NFSv4 GETATTR to the metadata server. Not only are these GETATTRs to the storage device individually expensive, the storage device can become inundated by a storm of such requests. NFSv3 solved a similar issue by having the READ and WRITE operations employ a post-operation attribute to report the weak cache consistency (WCC) data (See [Section 2.6 of \[RFC1813\]](#)).

Each NFSv3 operation corresponds to one round trip between the client and server. So a WRITE followed by a GETATTR would require two round trips. In that scenario, the attribute information retrieved is considered to be strict server-client consistency for a cache consistency protocol. For NFSv4, the WRITE and GETATTR can be issued together inside a compound, which only requires one round trip between the client and server. And this is also considered to be a strict server-client consistency. In essence, the NFSv4 READ and WRITE operations drop the post-operation attributes, allowing the client to decide if it needs that information.

With the flexible files layout type, the client can leverage the NFSv3 WCC to service the proxying of times (See [Section 4](#)). But the granularity of this data is limited. With client side mirroring (See [Section 8 of \[RFC8435\]](#)), the client has to aggregate the N mirrored files in order to send one piece of information instead of N pieces of information. Also, the client is limited to sending that information only when it returns the delegation.

The current filehandle and the `lowa_stateid` identifies the particular layout for the LAYOUT_WCC operation. The `lowa_type` indicates how to unpack the layout type specific payload inside the `lowa_body` field. The `lowa_type` is defined to be a value from the IANA registry for "pNFS Layout Types Registry".

The `lowa_body` will contain the data file attributes. The client will be responsible for mapping the NFSv3 post-operation attributes to those in a `fattr4`. Just as the post-operation attributes may be ignored by the client, the server may ignore the attributes inside the `LAYOUT_WCC`. But the server can also use those attributes to avoid querying the storage device for the data file attributes. Note that as these attributes are optional and there is nothing the client can do if the server ignores one, there is no need to return a `bitmap4` of which attributes were accepted in the result of the `LAYOUT_WCC`.

5.4. Allowed Errors

The `LAYOUT_WCC` operation can raise the errors in Table 1. When an error is encountered, the metadata server can decide to ignore the entire operation or depending on the layout type specific payload, it could decide to apply a portion of the payload.

Valid Error Returns for `LAYOUT_WCC`

+-----+	
Errors	
+-----+	
NFS4ERR_ADMIN_REVOKED, NFS4ERR_BADXDR, NFS4ERR_BAD_STATEID,	
NFS4ERR_DEADSESSION, NFS4ERR_DELAY, NFS4ERR_DELEG_REVOKED,	
NFS4ERR_EXPIRED, NFS4ERR_FHEXPIRED, NFS4ERR_GRACE, NFS4ERR_INVAL,	
NFS4ERR_ISDIR, NFS4ERR_MOVED, NFS4ERR_NOFILEHANDLE,	
NFS4ERR_NOTSUPP, NFS4ERR_NO_GRACE, NFS4ERR_OLD_STATEID,	
NFS4ERR_OP_NOT_IN_SESSION, NFS4ERR_REP_TOO_BIG,	
NFS4ERR_REP_TOO_BIG_TO_CACHE, NFS4ERR_REQ_TOO_BIG,	
NFS4ERR_RETRY_UNCACHED_REP, NFS4ERR_SERVERFAULT, NFS4ERR_STALE,	
NFS4ERR_TOO_MANY_OPS, NFS4ERR_UNKNOWN_LAYOUTTYPE,	
NFS4ERR_WRONG_CRED, NFS4ERR_WRONG_TYPE	
+-----+	

Table 1

5.5. Flex Files Layout Type


```
<CODE BEGINS>
/// struct ff_data_server_wcc4 {
///     deviceid4          ffdsw_deviceid;
///     stateid4           ffdsw_stateid;
///     nfs_fh4            ffdsw_fh_vers<>;
///     fattr4             ffdsw_attributes;
/// };
///
/// struct ff_mirror_wcc4 {
///     ff_data_server_wcc4  ffdsw_data_servers<>;
/// };
///
/// struct ff_layout_wcc4 {
///     ff_mirror_wcc4       ffdsw_mirrors<>;
/// };
<CODE ENDS>
```

The flex file layout type specific results SHOULD correspond to the ff_layout4 data structure as defined in [Section 5.1 of \[RFC8435\]](#). There SHOULD be a one-to-one correspondence between:

- o ff_data_server4 -> ff_data_server_wcc4
- o ff_mirror4 -> ff_mirror_wcc4
- o ff_layout4 -> ff_layout_wcc4

Each ff_layout4 has an array of ff_mirror4, which have an array of ff_data_server4. Based on the current filehandle and the lowa_stateid, the server can match the reported attributes.

But the positional correspondence between the elements is not sufficient to determine the attributes to update. Consider the case where a layout had three mirrors and two of them had updated attributes, but the third did not. A client could decide to present all three mirrors, with one mirror having an attribute mask with no attributes present. Or it could decide to present only the two mirrors which had been changed.

In either case, the combination of ffdsw_deviceid, ffdsw_stateid, and ffdsw_fh_vers will uniquely identify the attributes to be updated. All three arguments are required. A layout might have multiple data files on the same storage device, in which case the ffdsw_deviceid and ffdsw_stateid would match, but the ffdsw_fh_vers would not.

The ffdsw_attributes are processed similar to the obj_attributes in the SETATTR arguments (See [Section 18.30 of \[RFC5661\]](#)).

6. Extraction of XDR

This document contains the external data representation (XDR) [RFC4506] description of the new open flags for delegating the file to the client. The XDR description is embedded in this document in a way that makes it simple for the reader to extract into a ready-to-compile form. The reader can feed this document into the following shell script to produce the machine readable XDR description of the new flags:

```
<CODE BEGINS>
```

```
#!/bin/sh
grep '^ *///' $* | sed 's?^ */// ??' | sed 's?^ *///$??'
```

```
<CODE ENDS>
```

That is, if the above script is stored in a file called "extract.sh", and this document is in a file called "spec.txt", then the reader can do:

```
sh extract.sh < spec.txt > delstid_prot.x
```

The effect of the script is to remove leading white space from each line, plus a sentinel sequence of "///". XDR descriptions with the sentinel sequence are embedded throughout the document.

Note that the XDR code contained in this document depends on types from the NFSv4.2 nfs4_prot.x file (generated from [RFC7863]). This includes both nfs types that end with a 4, such as offset4, length4, etc., as well as more generic types such as uint32_t and uint64_t.

While the XDR can be appended to that from [RFC7863], the various code snippets belong in their respective areas of the that XDR.

6.1. Code Components Licensing Notice

Both the XDR description and the scripts used for extracting the XDR description are Code Components as described in [Section 4](#) of "Legal Provisions Relating to IETF Documents" [LEGAL]. These Code Components are licensed according to the terms of that document.

7. Security Considerations

There are no new security considerations beyond those in [RFC7862].

8. IANA Considerations

There are no IANA considerations.

9. Normative References

- [LEGAL] IETF Trust, "Legal Provisions Relating to IETF Documents", November 2008, <<http://trustee.ietf.org/docs/IETF-Trust-License-Policy.pdf>>.
- [RFC1813] IETF, "NFS Version 3 Protocol Specification", [RFC 1813](#), June 1995.
- [RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", [BCP 14](#), [RFC 2119](#), March 1997.
- [RFC4506] Eisler, M., "XDR: External Data Representation Standard", STD 67, [RFC 4506](#), May 2006.
- [RFC5661] Shepler, S., Ed., Eisler, M., Ed., and D. Noveck, Ed., "Network File System (NFS) Version 4 Minor Version 1 Protocol", [RFC 5661](#), January 2010.
- [RFC7862] Haynes, T., "NFS Version 4 Minor Version 2", [RFC 7862](#), November 2016.
- [RFC7863] Haynes, T., "Network File System (NFS) Version 4 Minor Version 2 External Data Representation Standard (XDR) Description", [RFC 7863](#), November 2016.
- [RFC8178] Noveck, D., "Rules for NFSv4 Extensions and Minor Versions", [RFC 8178](#), July 2017.
- [RFC8435] Halevy, B. and T. Haynes, "Parallel NFS (pNFS) Flexible File Layout", [RFC 8435](#), August 2018.

Appendix A. Acknowledgments

Trond Myklebust and David Flynn all worked on the prototype at Hammerspace.

Appendix B. RFC Editor Notes

[RFC Editor: please remove this section prior to publishing this document as an RFC]

[RFC Editor: prior to publishing this document as an RFC, please replace all occurrences of RFCTBD10 with RFCxxxx where xxxx is the RFC number of this document]

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