

NSIS  
Internet-Draft  
Intended status: Informational  
Expires: January 27, 2010

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July 27, 2009

**GIST State Machine**  
**draft-ietf-nsis-ntlp-statemachine-08.txt**

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Abstract

This document describes the state machines for the General Internet Signaling Transport (GIST). The states of GIST nodes for a given flow and their transitions are presented in order to illustrate how GIST



may be implemented.

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## **1. Introduction**

The state machines described in this document are illustrative of how the GIST protocol defined in [1] may be implemented for the GIST nodes in different locations of a flow path. Where there are differences - [1] is authoritative. The state machines are informative only. Implementations may achieve the same results using different methods.

There are two types of possible entities for GIST signaling:

- GIST querying node - GIST node that initiates the discovery of the next peer;
- GIST responding node - GIST node that is the discovered next peer;

We describe a set of state machines for these entities to illustrate how GIST may be implemented.

## **2. Terminology**

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [2].

## **3. Notational conventions used in state diagrams**

The following text is reused from [3] and the state diagrams are based on the conventions specified in [4], Section 8.2.1. Additional state machine details are taken from [5].

The complete text is reproduced here:

State diagrams are used to represent the operation of the protocol by a number of cooperating state machines each comprising a group of connected, mutually exclusive states. Only one state of each machine can be active at any given time.

All permissible transitions between states are represented by arrows, the arrowhead denoting the direction of the possible transition. Labels attached to arrows denote the condition(s) that must be met in order for the transition to take place. All conditions are expressions that evaluate to TRUE or FALSE; if a condition evaluates to TRUE, then the condition is met. The label UCT denotes an unconditional transition (i.e., UCT always evaluates to TRUE). A transition that is global in nature (i.e., a transition that occurs from any of the possible states if the condition attached to the arrow is met) is denoted by an open arrow; i.e., no specific state is



identified as the origin of the transition. When the condition associated with a global transition is met, it supersedes all other exit conditions including UCT. The special global condition BEGIN supersedes all other global conditions, and once asserted remains asserted until all state blocks have executed to the point that variable assignments and other consequences of their execution remain unchanged.

On entry to a state, the procedures defined for the state (if any) are executed exactly once, in the order that they appear on the page. Each action is deemed to be atomic; i.e., execution of a procedure completes before the next sequential procedure starts to execute. No procedures execute outside of a state block. The procedures in only one state block execute at a time, even if the conditions for execution of state blocks in different state machines are satisfied, and all procedures in an executing state block complete execution before the transition to and execution of any other state block occurs, i.e., the execution of any state block appears to be atomic with respect to the execution of any other state block and the transition condition to that state from the previous state is TRUE when execution commences. The order of execution of state blocks in different state machines is undefined except as constrained by their transition conditions. A variable that is set to a particular value in a state block retains this value until a subsequent state block executes a procedure that modifies the value.

On completion of all of the procedures within a state, all exit conditions for the state (including all conditions associated with global transitions) are evaluated continuously until one of the conditions is met. The label ELSE denotes a transition that occurs if none of the other conditions for transitions from the state are met (i.e., ELSE evaluates to TRUE if all other possible exit conditions from the state evaluate to FALSE). Where two or more exit conditions with the same level of precedence become TRUE simultaneously, the choice as to which exit condition causes the state transition to take place is arbitrary.

In addition to the above notation, there are a couple of clarifications specific to this document. First, all boolean variables are initialized to FALSE before the state machine execution begins. Second, the following notational shorthand is specific to this document:

<variable> = <expression1> | <expression2> | ...

Execution of a statement of this form will result in <variable> having a value of exactly one of the expressions. The logic for which of those expressions gets executed is outside of the state





machine and could be environmental, configurable, or based on another state machine such as that of the method.

#### 4. State Machine Symbols

( )

Used to force the precedence of operators in Boolean expressions and to delimit the argument(s) of actions within state boxes.

;

Used as a terminating delimiter for actions within state boxes. Where a state box contains multiple actions, the order of execution follows the normal English language conventions for reading text.

=

Assignment action. The value of the expression to the right of the operator is assigned to the variable to the left of the operator. Where this operator is used to define multiple assignments, e.g., `a = b = X` the action causes the value of the expression following the right-most assignment operator to be assigned to all of the variables that appear to the left of the right-most assignment operator.

!

Logical NOT operator.

&&

Logical AND operator.

||

Logical OR operator.

if...then...

Conditional action. If the Boolean expression following the if evaluates to TRUE, then the action following the then is executed.

{ statement 1, ... statement N }

Compound statement. Braces are used to group statements that are executed together as if they were a single statement.

!=

Inequality. Evaluates to TRUE if the expression to the left of the operator is not equal in value to the expression to the right.

==

Equality. Evaluates to TRUE if the expression to the left of the operator is equal in value to the expression to the right.



>  
Greater than. Evaluates to TRUE if the value of the expression to the left of the operator is greater than the value of the expression to the right.

<=  
Less than or equal to. Evaluates to TRUE if the value of the expression to the left of the operator is either less than or equal to the value of the expression to the right.

++  
Increment the preceding integer operator by 1.

+  
Arithmetic addition operator.

&  
Bitwise AND operator.

## 5. Common Rules

Throughout the document we use terms defined in the [1], such as Query, Response, Confirm.

State machine represents handling of GIST messages that match a Message Routing State's MRI, NSLPID and SID and with no protocol errors. Separate parallel instances of the state machines should handle messages for different Message Routing States.

The state machine states represent the upstream/downstream peers states of the Message Routing State.

For simplification not all objects included in a message are shown. Only those that are significant for the case are shown. State machines do not present handling of messages that are not significant for management of the states.

Presented in this document state machines do not cover all functions of a GIST node. Functionality of message forwarding, transmission of NSLP data without MRS establishment and providing of the received messages to the appropriate MRS, we refer as "Lower level pre-processing" step. Pre-processing provides to the appropriate MRS FSM only the messages which are matched against waiting Query/Response cookies, or established MRS MRI+NSLPID+SID primary key. This is presented by "rx\_\*" events in the state machines.

Management of a MA is considered in the document (e.g., tg\_Establish\_MA, tg\_MA\_established events), but its FSM is not



explicitly presented.

## **5.1 Common Procedures**

**Tg\_SendMsg:**

NSLP/GIST API message that request transmission of a NSLP message.

**Tg\_SetStateLifetime(time\_period):**

NSLP/GIST API message providing info for the Lifetime of an RS, required by the application. "Time\_period = 0" represents the cancellation of established RSs/MAS (invoked by NSLP application).

**Tg\_MessageStatus:**

NSLP/GIST API message informing NSLP application of unsuccessful delivery of a message

**Tg\_RecvMsg:**

NSLP/GIST API message that provides received message to the NSLP

**Tg\_NetworkNotification:**

NSLP/GIST API message that informs NSLP for change in MRS

**Tx\_Query:**

Transmit of Query message

**Tx\_Response:**

Transmit of Response message

**Tx\_Confirm:**

Transmit of Confirm message

**Rx\_Query:**

Receive of Query message

**Rx\_Response:**

Receive of Response message

**Rx\_Confirm:**

Receive of Confirm message

**Tx\_Error:**

Transmit of Error message

**Rx\_Error:**

Receive of Error message

**Queue NSLP info:**

Save NSLP messages in a queue until a required MA association is



established

Tx\_Data:

Transmit of Data message

Rx\_Data:

Receive of Data message

T\_Inactive\_QNode:

Message Routing State lifetime timer in Querying Node

T\_Expired\_RNode:

Message Routing State lifetime timer in Responding Node

T\_Refresh\_QNode:

Message Routing State refresh timer in Querying Node

T\_No\_Response:

Timer for the waiting period for Response message in Querying Node

T\_No\_Confirm:

Timer for the waiting period for Confirm message in Responding Node

Install downstream/upstream MRS:

Install new Message Routing State and save the corresponding peer state info (IP address and UDP port or pointer to the used MA) for the current Message Routing State or update the corresponding peer state info.

DELETE MRS:

Delete installed downstream/upstream peer's info for the current Message Routing State and delete the Message Routing State if required.

Established MA:

A Message Association (MA) is established between the current node and its upstream peer. The initiator for the establishment is the upstream peer.

Re-use existing MA:

An existing MA between the current node and its peer is re-used.

DELETE MA:

Delete/disconnect used MA.

Stop using shared MA:

Stop using shared MA. If the shared MA is no more used by any





other MRSs, it depends on the local policy whether it is deleted or kept.

**REFRESH MRS:**

Refreshes installed MRS.

**Tg\_MA\_Error:**

Error event with used MA.

**Tg\_InvalidRoutingState:**

Notification from NSLP application for path change

**Tg\_Establish\_MA:**

Triggers establishment of MA.

**Tg\_MA\_Established:**

MA has been successfully established.

**Tg\_ERROR:**

General Error event / system level error.

**No\_MRS\_Installed:**

Error response, send by the Responding node indicating lost Confirm message.

## **5.2 Common Variables**

It is assumed that the type of mode and destination info (which need to be taken from the application parameters and local GIST policy) is provided. This is represented by the common variables Dmode, Cmode, MAinfo, MApresent and Refresh.

**Cmode:**

The message MUST be transmitted in Cmode. This is specified by "Message transfer attributes" set to any of the following values:

"Reliability" is set to TRUE.

"Security" is set to values that request secure handling of a message.

"Local processing" is set to values that require services offered by Cmode (e.g., congestion control). [[1](#)]

**Dmode:**

The message MUST be transmitted in Dmode. This is specified by local policy rules and in case that the "Message transfer attributes" are not set to any of the following values:



"Reliability" is set to TRUE.

"Security" is set to values that request special security handling of a message.

"Local processing" is set to values that require services offered by Cmode [\[1\]](#)

**MAinfo:**

GIST message parameters describing the required MA or proposed MA e.g. "Stack-proposal" and "Stack-Configuration-Data".

**NSLPdata:**

NSLP application data.

**RespCookie:**

Responder Cookie that is being sent by the Responding node with the Response message in case that its local policy requires a confirmation from the querying node.

**ConfirmRequired:**

Confirm message is required by the local policy rule for installation of the new MRS.

**NewPeer:**

Response message is received from new responding peer.

**MAexist:**

Existing MA will be reused.

**CheckPeerInfo:**

The sender of the received data message is matched against the installed peer info in the MRS.

**UpstreamPeerInstalled:**

Upstream peer info is installed in the MRS.

## **[6.](#) State machines**

The following section presents the state machine diagrams of GIST peers.

### **[6.1](#) Diagram notations**

(see the .pdf version for missing diagram or refer to [Appendix A](#) if reading the .txt version)



Figure 1: Diagram notations

## 6.2 State machine for GIST querying node

The following is a diagram of the GIST querying node state machine. Also included is clarification of notation.

(see the .pdf version for missing diagram or refer to [Appendix A.1](#) if reading the .txt version)

Figure 1: GIST Querying Node State Machine

- \*) Response and Comfirm messages might be send either in Dmode or Cmode, before or after MA establishment depending on node s local 3-way handshake policy and the availability of MAs to be reused. See draft for details.
- \*\*) Depending on the local policy NSLPdata might be send as payload of Query and Confirm messages. (piggybacking)
- 1) Initial request from NSLP is received, which triggers Query messages requesting either D\_mode or C\_mode. Depending on node s local policy NSLP data might be piggybacked in the Query requesting D\_mode. Query may carry Mainfo if C\_mode transport is needed.
- 2) Response message is received. If C\_mode connection must be established and there is no available MA to be reused, MA establishment is initiated and waited to be completed.
- 3) Response message is received. If D\_mode connection is requested or available MA can be reused for requested C\_mode, the MRS is established.
- 4) No\_Response timer expires. Query is resent.
- 5) No\_Response timer expires and maximum number of retries has been reached. NSLP application is notified for the GIST peer discovery failure.
- 6) NSLP data is queued, because downstream peer is not discovered or required MA is still not established.
- 7) Data message is received. It is checked if its sender matches the installed downstream peer info in the MRS and then processed. In WaitResponse state, this event might happen in the process of MA upgrade, when the downstream peer is still not aware of establishment of the new MA.
- 8) Provided NSLP data is sent via Data message towards downstream GIST peer.
- 9) Refresh\_QNode timer expires. Query message is sent.
- 10) Response message from the downstream GIST peer is received. The peer is not changed. MRS is refreshed (Refresh\_QNode timer is restarted).
- 11) Path change detected. Response message from a new downstream GIST peer is received. D\_mode is requested or existing MA can be reused



for requested C\_mode.

- 12) Path change detected. Response message from a new downstream GIST peer is received. A new MA must be established for requested C\_mode.
- 13) Requested by NSLP application transport parameters requires upgrade of established MRS from D\_mode/C\_mode to C\_mode. NSLP application notifies GIST for path change. Downstream GIST peer discovery is initiated.
- 14) Sent Confirm message has not been received by downstream GIST peer. Confirm message is resent.
- 15) MRS lifetime expires. Notification by NSLP application that MRS is no longer needed.
- 16) MA is established.
- 17) MA establishment failure.

### **6.3 State machine for GIST responding node**

The following is a diagram of the GIST responding node state machine. Also included is clarification of notation.

(see the .pdf version for missing diagram or refer to [Appendix A.2](#) if reading the .txt version)

Figure 3: GIST Responding Node State Machine

- 1) A Query message is received. Explicit Confirm message is required for MRS installation, based on the local policy. Query message might carry piggybacked NSLP data which is provided to the NSLP application.
- 2) A Query message is received. MRS is installed immediately, based on the local policy. Query message might carry piggybacked NSLP data which is provided to the NSLP application.
- 3) Confirm message is received which causes installation of the complete MRS or just installation of the used MA as a upstream peer info.
- 4) Sent Response message has not been received by upstream GIST peer. Response message is resent.
- 5) In case of lost Confirm message, data messages might be received from the upstream GIST node (it is unaware of the lost Confirm message). Response indicating the loss of the Confirm is sent back to the upstream GIST node.
- 6) No\_Confirm timer expires. Note that all cases of lost handshake GIST messages are handled only by GIST querying node via resend of Query message.
- 7) NSLP data is sent if discovery process is successfully accomplished or is queued if Confirm message is still expected to confirm establishment of MA.
- 8) Data messages are accepted only if complete MRS is installed,





e.g., there is installed upstream peer info. If not, then Confirm message is expected and data message won't be accepted. Response indicating the loss of the Confirm is sent back to the upstream GIST node.

- 9) Change of the upstream GIST node (e.g., path change). Local policy does not need explicit Confirm message for MRS installation. MRS data is updated.
- 10) Change of the upstream GIST node or request for change of the used connection mode (from D\_mode/C\_mode to better C\_mode). Local policy requires explicit Confirm message for MRS installation.
- 11) Request for change of the used connection mode (from D\_mode/C\_mode to better C\_mode). Local policy does not need explicit Confirm message for MRS installation. MRS data is updated.
- 12) MRS lifetime expires. Notification by NSLP application that MRS is no longer needed.

## **7. Security Considerations**

This document does not raise new security considerations. Any security concerns with GIST are likely reflected in security related NSIS work already (such as [\[1\]](#) or [\[6\]](#)).

## **8. IANA Considerations**

This document has no actions for IANA.

## **9. Contributors**

Christian Dickmann contributed to refining of the state machine since 01 version.

## **10. Acknowledgments**

The authors would like to thank Robert Hancock, Ingo Juchem, Andreas Westermaier, Alexander Zrim, Julien Abeille Youssef Abidi and Bernd Schloer for their insightful comments.



## **11. References**

### **11.1. Normative References**

- [1] Schulzrinne, H., "GIST: General Internet Signaling Transport", [draft-ietf-nsis-ntlp-20](#) (work in progress), June 2009.
- [2] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", [BCP 14](#), [RFC 2119](#), March 1997.

### **11.2. Informative References**

- [3] Vollbrecht, J., Eronen, P., Petroni, N., and Y. Ohba, "State Machines for Extensible Authentication Protocol (EAP) Peer and Authenticator", [RFC4137](#), August 2005.
- [4] Institute of Electrical and Electronics Engineers, "Standard for Local and Metropolitan Area Networks: Port-Based Network Access Control", IEEE 802-1X-2004, December 2004.
- [5] Fajardo, V., Ohba, Y. and R. Lopez, "State Machines for Protocol for Carrying Authentication for Network Access (PANA)", [draft-ietf-pana-statemachine-12](#) (work in progress), April 2009.
- [6] Tschofenig, H. and D. Kroeselberg, "Security Threats for NSIS", [RFC 4081](#), June 2005.



## Appendix A. ASCII versions of state diagrams

This appendix contains the state diagrams in ASCII format. Please use the PDF version whenever possible: it is much easier to understand.

For each state there is a separate table that lists in each row:

- an event that triggers a transition,
- actions taken as a result of the incoming event,
- and the new state at which the transitions ends.

### A.1. State machine for GIST querying node (Figure 2)

-----  
State: IDLE  
-----

Condition	Action	State	Note
tg_SendMsg	tx_Query	Wait	1)
	start T_No_Response	Response	**
	Queue NSLP data		
Tg_ERROR	Delete MRS	IDLE	
	IF (MA is used)		
	((Delete MA)		
	(Stop using shared MA))		
	Tg_NetworkNotification		

-----  
State: WaitResponse  
-----

Condition	Action	State	Note
rx_Response(MAinfo)&& (!MAexist)	tg_Establish_MA	Wait MA	*
	(tx_Confirm)	Establish.	2)
rx_Response)	Install MRS	Established	3)
(rx_Response(MAinfo)&& (MAexist))	IF (RespCookie)	Downstream	
	tx_Confirm(RespCookie)	MRS	
	tx_Data(Queued NSLP data)		



(timeout T_No_Response)	Tx_Query	Wait	4)
&&(!MaxRetry)	restart T_No_Response	Response	
(timeout T_No_Response)	tg_MessageStatus	IDLE	5)
&&(MaxRetry)			
tg_SendMsg	Queue NSLP data	Wait	6)
		Response	
rx_Data	IF(CheckPeerInfo)	Wait	7)
	tg_RecvMsg to Appl.	Response	
Tg_ERROR	Delete MRS)	IDLE	
	IF (MA is used)		
	((Delete MA)		
	(Stop using shared MA))		
	Tg_NetworkNotification		
-----+-----+-----+-----			

-----  
State: Established Downstream MRS  
-----

Condition	Action	State	Note
-----+-----+-----+-----			
tg_SendMsg	tx_Data	Established	8)
	restart T_Inactive_QNode	Downstream	
		MRS	
timeout T_Refresh_QNode	tx_Query	Established	9)
		Downstream	
		MRS	
(rx_Response)&&	Refresh MRS	Established	10)
(!NewPeer)	restart T_Inactive_QNode	Downstream	
		MRS	
(rx_Response)	IF (MA is used)	Established	11)
(rx_Response(Mainfo)&&	(Delete MA)	Downstream	
(MAexist)))&&(NewPeer)	(Stop using shared MA)	MRS	
	Install MRS		
	restart T_Inactive_QNode		
	IF (RespCookie)		





	tx_Confirm(RespCookie)		
(rx_Response(MAinfo)&& (NewPeer)&&(!MA_exist))	((Delete MA)   (Stop using shared MA))	Wait MA  Establish.	12)  *
	tg_Establish_MA  (tx_Confirm)		
((tg_SendMsg)&&(Cmode)&& (!MAexist))   (tg_MA_error)   (tg_InvalidRoutingState)	tx_Query  Queue NSLP data	Wait  Response	13) 
rx_Response(No_MRS_ installed)	tx_Confirm(RespCookie)  tx_Data(Queued NSLP data)	Established  Downstream	14) 
		MRS	
(timeout T_Inactive_ QNode)   (tg_SetStateLifetime(0))	Delete MRS  IF (MA is used)  (Delete MA)    (Stop using shared MA)   Tg_NetworkNotification	IDLE	15) 
rx_Data	IF(CheckPeerInfo)  tg_RecvMsg to Appl.	Established  Downstream	7) 
		MRS	
Tg_ERROR	Delete MRS)  IF (MA is used)  ((Delete MA)    (Stop using shared MA))  Tg_NetworkNotification	IDLE	
-----+-----+-----+-----			

-----  
State: Wait MA Establishment  
-----

Condition	Action	State	Note
-----+-----+-----+-----			
tg_MA_Established	Install MRS  (tx_Confirm)  tx_Data(Queued NSLP data)	Established  Downstream  MRS	16)  * 
tg_MA_error	Delete MRS  tg_MessageStatus	IDLE	17) 



tg_SendMsg	Queue NSLP data	Wait MA	6)
		Establish.	
Tg_ERROR	Delete MRS	IDLE	
	IF (MA is used)		
	((Delete MA))		
	(Stop using shared MA))		
	Tg_NetworkNotification		
-----+-----+-----			

Figure 4

**A.2. State Machine for GIST responding node (Figure 3)**

-----  
 State: IDLE  
 -----

Condition	Action	State	Note
-----+-----+-----			
rx_Query&& (ConfirmRequired)	tx_Response	Wait	1)
	start T_No_Confirm	Confirm	
	IF(NSLPdata)		
	tg_RecvMsg(NSLPdata)		
	to Appl.		
rx_Query&& (!ConfirmRequired)	tx_Response	Established	2)
	Install MRS	Upstream	
	IF(NSLPdata)	MRS	
	tg_RecvMsg(NSLPdata)		
	to Appl.		
-----+-----+-----			

-----  
 State: WAIT CONFIRM  
 -----

Condition	Action	State	Note
-----+-----+-----			
rx_Confirm	Install Upstream MRS	Established	3)
		Upstream	
		MRS	



rx_Query&& (ConfirmRequired)	tx_Response	Wait	4)
	start T_No_Confirm	Confirm	
	IF(NSLPdata)		
	tg_RecvMsg(NSLPdata)		
	to Appl.		
rx_Data	tx_Error(No_MRS_	Wait	5)
	installed)	Confirm	
timeout T_No_Confirm		IDLE	6)
-----+-----+-----+-----			

-----  
State: Established Upstream MRS  
-----

Condition	Action	State	Note
-----+-----+-----+-----			
tg_SendMsg	IF(!UpstreamPeerInfo)	Established	7)
	Queue NSLP data	Upstream	
	ELSE tx_Data	MRS	
rx_Data	IF(UpstreamPeerInfo)	Established	8)
	(tg_RecvMsg to Appl.)	Upstream	
	&&(restart_T_Expire_	MRS	
	RNode)		
	ELSE		
rx_Query	tx_Error(No_MRS_		
	installed)		
	IF (NewPeer)	Established	9)
	Update UpstreamPeerInfo	Upstream	
(rx_Query)&& (ConfirmRequired)	tx_Response	MRS	
	restart T_Expire_RNode		
	Delete MRS	Wait	
	tx_Response	Confirm	
	start T_No_Confirm		
	IF(MA is used)		
	(Delete MA)		
	(Stop using shared MA)		
	IF(NSLPdata)		
	tg_RecvMsg(NSLPdata)		
	to Appl.		



rx_Query(MAinfo)&& (!ConfirmRequired)	Delete UpstreamPeerInfo restart T_Expire_RNode tx_Response(MAinfo)	Established Upstream MRS	11)
(timeout T_Expire_RNode)    (tg_SetStateLifetime(0))	Delete MRS tg_NetworkNotification IF(MA is used) (Delete MA)   (Stop using shared MA)	IDLE	12)
rx_Confirm	Install UpstreamPeerInfo tx_Data(queued_NSLP_data)	Established Upstream MRS	3)
Tg_ERROR	(Delete MRS) IF (MA is used) ((Delete MA)   (Stop using shared MA)) Tg_NetworkNotification	IDLE	
-----+-----+-----+			

Figure 5





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