Workgroup:
Oblivious HTTP Application Intermediation
Internet-Draft: draft-ietf-ohai-ohttp-02
Published: 11 July 2022
Intended Status: Standards Track
Expires: 12 January 2023
Authors: M. Thomson C. A. Wood
Mozilla Cloudflare

Oblivious HTTP

Abstract

This document describes a system for the forwarding of encrypted HTTP messages. This allows a client to make multiple requests of a server without the server being able to link those requests to the client or to identify the requests as having come from the same client.

About This Document

This note is to be removed before publishing as an RFC.

The latest revision of this draft can be found at https://ietf-wg-ohai.github.io/oblivious-http/draft-ietf-ohai-ohttp.html. Status information for this document may be found at https://datatracker.ietf.org/doc/draft-ietf-ohai-ohttp/.

Discussion of this document takes place on the Oblivious HTTP Application Intermediation Working Group mailing list (mailto:ohai@ietf.org), which is archived at https://mailarchive.ietf.org/arch/browse/ohai/.

Source for this draft and an issue tracker can be found at https://github.com/ietf-wg-ohai/oblivious-http.

Status of This Memo

This Internet-Draft will expire on 12 January 2023.

Copyright Notice

Table of Contents

1. Introduction
2. Overview
  2.1. Applicability
  2.2. Conventions and Definitions
3. Key Configuration
  3.1. Key Configuration Encoding
  3.2. Key Configuration Media Type
4. HPKE Encapsulation
  4.1. Encapsulation of Requests
  4.2. Encapsulation of Responses
5. HTTP Usage
  5.1. Informational Responses
  5.2. Errors
6. Media Types
  6.1. message/ohttp-req Media Type
  6.2. message/ohttp-res Media Type
7. Security Considerations
  7.1. Client Responsibilities
  7.2. Relay Responsibilities
    7.2.1. Differential Treatment
    7.2.2. Denial of Service
    7.2.3. Linkability Through Traffic Analysis
  7.3. Server Responsibilities
  7.4. Replay Attacks
    7.4.1. Use of Date for Anti-Replay
  7.5. Forward Secrecy
  7.6. Post-Compromise Security
8. Privacy Considerations
9. Operational and Deployment Considerations
  9.1. Performance Overhead

9.2. Resource Mappings
9.3. Network Management
10. Repurposing the Encapsulation Format
11. IANA Considerations
12. References
12.1. Normative References
12.2. Informative References
Appendix A. Complete Example of a Request and Response
Acknowledgments
Authors' Addresses

1. Introduction

The act of making a request using HTTP reveals information about the client identity to a server. Though the content of requests might reveal information, that is information under the control of the client. In comparison, the source address on the connection reveals information that a client has only limited control over.

Even where an IP address is not directly attributed to an individual, the use of an address over time can be used to correlate requests. Servers are able to use this information to assemble profiles of client behavior, from which they can make inferences about the people involved. The use of persistent connections to make multiple requests improves performance, but provides servers with additional certainty about the identity of clients in a similar fashion.

Use of an HTTP proxy can provide a degree of protection against servers correlating requests. Systems like virtual private networks or the Tor network [Dingledine2004], provide other options for clients.

Though the overhead imposed by these methods varies, the cost for each request is significant. Preventing request linkability requires that each request use a completely new TLS connection to the server. At a minimum, this requires an additional round trip to the server in addition to that required by the request. In addition to having high latency, there are significant secondary costs, both in terms of the number of additional bytes exchanged and the CPU cost of cryptographic computations.

This document describes a method of encapsulation for binary HTTP messages [BINARY] using Hybrid Public Key Encryption (HPKE; [HPKE]). This protects the content of both requests and responses and enables a deployment architecture that can separate the identity of a requester from the request.

Though this scheme requires that servers and proxies (called relays in this document) explicitly support it, this design represents a performance improvement over options that perform just one request in each connection. With limited trust placed in the relay (see Section 7), clients are assured that requests are not uniquely attributed to them or linked to other requests.

2. Overview

A client must initially know the following:

- The identity of an Oblivious Gateway Resource. This might include some information about what Target Resources the Oblivious Gateway Resource supports.

- The details of an HPKE public key that the Oblivious Gateway Resource accepts, including an identifier for that key and the HPKE algorithms that are used with that key.

- The identity of an Oblivious Relay Resource that will accept relay requests carrying an encapsulated request as its content and forward the content in these requests to a single Oblivious Gateway Resource. See Section 9.2 for more information about the mapping between Oblivious Relay and Gateway Resources.

This information allows the client to make a request of a Target Resource with that resource having only a limited ability to correlate that request with the client IP or other requests that the client might make to that server.

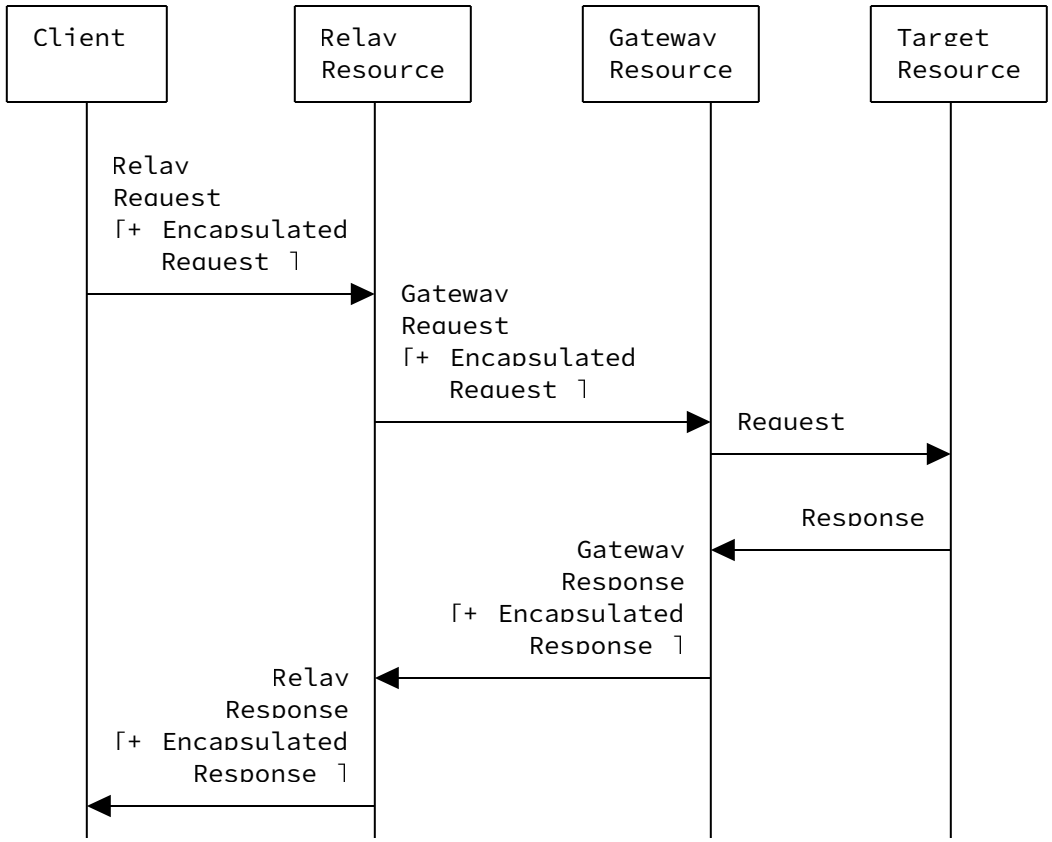

Figure 1: Overview of Oblivious HTTP

In order to make a request to a Target Resource, the following steps occur, as shown in Figure 1:

1. The client constructs an HTTP request for a Target Resource.

2. The client encodes the HTTP request in a binary HTTP message and then encapsulates that message using HPKE and the process from Section 4.1.

3. The client sends a POST request to the Oblivious Relay Resource with the Encapsulated Request as the content of that message.

4. The Oblivious Relay Resource forwards this request to the Oblivious Gateway resource.

5. The Oblivious Gateway Resource receives this request and removes the HPKE protection to obtain an HTTP request.

6. The Oblivious Gateway Resource makes an HTTP request that includes the target URI, method, fields, and content of the request it acquires.

7. The Target Resource answers this HTTP request with an HTTP response.

8. The Oblivious Gateway Resource encapsulates the HTTP response following the process in Section 4.2 and sends this in response to the request from the Oblivious Relay Resource.

9. The Oblivious Relay Resource forwards this response to the client.

10. The client removes the encapsulation to obtain the response to the original request.

2.1. Applicability

Oblivious HTTP has limited applicability. Many uses of HTTP benefit from being able to carry state between requests, such as with cookies ([RFC6265]), authentication (Section 11 of [HTTP]), or even alternative services ([RFC7838]). Oblivious HTTP removes linkage at the transport layer, which must be used in conjunction with applications that do not carry state between requests.

Oblivious HTTP is primarily useful where privacy risks associated with possible stateful treatment of requests are sufficiently large that the cost of deploying this protocol can be justified. Oblivious HTTP is simpler and less costly than more robust systems, like Prio ([PRIO]) or Tor ([Dingledine2004]), which can provide stronger guarantees at higher operational costs.

Oblivious HTTP is more costly than a direct connection to a server. Some costs, like those involved with connection setup, can be amortized, but there are several ways in which Oblivious HTTP is more expensive than a direct request:

* Each request requires at least two regular HTTP requests, which adds latency.

* Each request is expanded in size with additional HTTP fields, encryption-related metadata, and AEAD expansion.

* Deriving cryptographic keys and applying them for request and response protection takes non-negligible computational resources.

Examples of where preventing the linking of requests might justify these costs include:

* DNS queries. DNS queries made to a recursive resolver reveal information about the requester, particularly if linked to other queries.

* Telemetry submission. Applications that submit reports about their usage to their developers might use Oblivious HTTP for some types of moderately sensitive data.

These are examples of requests where there is information in a request that - if it were connected to the identity of the user - might allow a server to learn something about that user even if the identity of the user is pseudonymous. Other examples include the submission of anonymous surveys, making search queries, or requesting location-specific content (such as retrieving tiles of a map display).

2.2. Conventions and Definitions

The key words "**MUST**", "**MUST NOT**", "**REQUIRED**", "**SHALL**", "**SHALL NOT**", "**SHOULD**", "**SHOULD NOT**", "**RECOMMENDED**", "**NOT RECOMMENDED**", "**MAY**", and "**OPTIONAL**" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

Client: This document uses its own definition of client. When referring to the HTTP definition of client (Section 3.3 of [HTTP]), the term "HTTP client" is used; see Section 5.

Encapsulated Request: An HTTP request that is encapsulated in an HPKE-encrypted message; see Section 4.1.

Encapsulated Response: An HTTP response that is encapsulated in an HPKE-encrypted message; see Section 4.2.

Oblivious Relay Resource: An intermediary that forwards encapsulated requests and responses between clients and a single Oblivious Gateway Resource.

Oblivious Gateway Resource: A resource that can receive an encapsulated request, extract the contents of that request, forward it to a Target Resource, receive a response, encapsulate that response, then return that response.

Target Resource: The resource that is the target of an encapsulated request. This resource logically handles only regular HTTP requests and responses and so might be ignorant of the use of Oblivious HTTP to reach it.

Relay Request: An HTTP request from Client to Relay that contains an encapsulated request as the content.

Relay Response: An HTTP response from Relay to Client that contains an encapsulated response as the content.

Gateway Request:
An HTTP request from Relay to Gateway that contains an encapsulated request as the content.

Gateway Response: An HTTP response from Gateway to Relay that contains an encapsulated response as the content.

This draft includes pseudocode that uses the functions and conventions defined in [HPKE].

Encoding an integer to a sequence of bytes in network byte order is described using the function encode(n, v), where n is the number of bytes and v is the integer value. ASCII [ASCII] encoding of a string s is indicated using the function encode_str(s). The function len() returns the length of a sequence of bytes.

Formats are described using notation from Section 1.3 of [QUIC].

3. Key Configuration

A client needs to acquire information about the key configuration of the Oblivious Gateway Resource in order to send encapsulated requests. In order to ensure that clients do not encapsulate messages that other entities can intercept, the key configuration **MUST** be authenticated and have integrity protection.

This document does not define how that acquisition occurs. However, in order to help facilitate interoperability, it does specify a format for the keys. This ensures that different client implementations can be configured in the same way and also enables advertising key configurations in a consistent format. This format might be used, for example with HTTPS, as part of a system for configuring or discovering key configurations. Note however that such a system needs to consider the potential for key configuration to be used to compromise client privacy; see Section 8.

A client might have multiple key configurations to select from when encapsulating a request. Clients are responsible for selecting a preferred key configuration from those it supports. Clients need to consider both the key encapsulation method (KEM) and the combinations of key derivation function (KDF) and authenticated encryption with associated data (AEAD) in this decision.

3.1. Key Configuration Encoding

A single key configuration consists of a key identifier, a public key, an identifier for the KEM that the public key uses, and a set HPKE symmetric algorithms. Each symmetric algorithm consists of an identifier for a KDF and an identifier for an AEAD.

Figure 2 shows a single key configuration.

```
HPKE Symmetric Algorithms {
  HPKE KDF ID (16),
  HPKE AEAD ID (16),
}

OHTTP Key Config {
  Key Identifier (8),
  HPKE KEM ID (16),
  HPKE Public Key (Npk * 8),
  HPKE Symmetric Algorithms Length (16),
  HPKE Symmetric Algorithms (32..262140),
}
```

Figure 2: A Single Key Configuration

The definitions for the identifiers used in HPKE and the semantics of the algorithms they identify can be found in [HPKE]. The Npk parameter is determined by the choice of HPKE KEM, which can also be found in [HPKE].

3.2. Key Configuration Media Type

The "application/ohttp-keys" format is a media type that identifies a serialized collection of key configurations. The content of this media type comprises one or more key configuration encodings (see Section 3.1) that are concatenated.

Evolution of the key configuration format is supported through the definition of new formats that are identified by new media types.

Type name: application

Subtype name: ohttp-keys

Required parameters: N/A

Optional parameters: None

Encoding considerations: only "8bit" or "binary" is permitted

Security considerations: see Section 7

Interoperability considerations: N/A

Published specification: this specification

Applications that use this media type: N/A

Fragment identifier considerations:
N/A

Additional information:

Magic number(s): N/A

Deprecated alias names for this type: N/A

File extension(s): N/A

Macintosh file type code(s): N/A

Person and email address to contact for further information: see Authors' Addresses section

Intended usage: COMMON

Restrictions on usage: N/A

Author: see Authors' Addresses section

Change controller: IESG

4. HPKE Encapsulation

HTTP message encapsulation uses HPKE for request and response encryption.

An encapsulated HTTP request contains a binary-encoded HTTP message [BINARY] and no other fields; see Figure 3.

```
Request {
  Binary HTTP Message (..),
}
```

Figure 3: Plaintext Request Content

An Encapsulated Request is comprised of a key identifier and a HPKE-protected request message. HPKE protection includes an encapsulated KEM shared secret (or enc), plus the AEAD-protected request message. An Encapsulated Request is shown in Figure 4. Section 4.1 describes the process for constructing and processing an Encapsulated Request.

```
Encapsulated Request {
  Key Identifier (8),
  KEM Identifier (16),
  KDF Identifier (16),
  AEAD Identifier (16),
  Encapsulated KEM Shared Secret (8*Nenc),
  AEAD-Protected Request (..),
}
```

Figure 4: Encapsulated Request

The Nenc parameter corresponding to the HpkeKdfId can be found in Section 7.1 of [HPKE].

An encrypted HTTP response includes a binary-encoded HTTP message [BINARY] and no other content; see Figure 5.

```
Response {
  Binary HTTP Message (..),
}
```

Figure 5: Plaintext Response Content

Responses are bound to responses and so consist only of AEAD-protected content. Section 4.2 describes the process for constructing and processing an Encapsulated Response.

```
Encapsulated Response {
  Nonce (Nk),
  AEAD-Protected Response (..),
}
```

Figure 6: Encapsulated Response

The Nenc and Nk parameters corresponding to the HpkeKdfId can be found in [HPKE]. Nenc refers to the size of the encapsulated KEM shared secret, in bytes; Nk refers to the size of the AEAD key for the HPKE ciphersuite, in bits.

4.1. Encapsulation of Requests

Clients encapsulate a request request using values from a key configuration:

- the key identifier from the configuration, keyID, with the corresponding KEM identified by kemID,

- the public key from the configuration, pkR, and

*a selected combination of KDF, identified by kdfID, and AEAD, identified by aeadID.

The client then constructs an Encapsulated Request, enc_request, from a binary encoded HTTP request, request, as follows:

1. Construct a message header, hdr, by concatenating the values of keyID, kemID, kdfID, and aeadID, as one 8-bit integer and three 16-bit integers, respectively, each in network byte order.

2. Build info by concatenating the ASCII-encoded string "message/bhttp request", a zero byte, and the header.

3. Create a sending HPKE context by invoking SetupBaseS() (Section 5.1.1 of [HPKE]) with the public key of the receiver pkR and info. This yields the context sctxt and an encapsulation key enc.

4. Encrypt request by invoking the Seal() method on sctxt (Section 5.2 of [HPKE]) with empty associated data aad, yielding ciphertext ct.

5. Concatenate the values of hdr, enc, and ct, yielding an Encrypted Request enc_request.

Note that enc is of fixed-length, so there is no ambiguity in parsing this structure.

In pseudocode, this procedure is as follows:

```
hdr = concat(encode(1, keyID),
             encode(2, kemID),
             encode(2, kdfID),
             encode(2, aeadID))
info = concat(encode_str("message/bhttp request"),
              encode(1, 0),
              hdr)
enc, sctxt = SetupBaseS(pkR, hdr)
ct = sctxt.Seal([], request)
enc_request = concat(hdr, enc, ct)
```

Servers decrypt an Encapsulated Request by reversing this process. Given an Encapsulated Request enc_request, a server:

1. Parses enc_request into keyID, kemID, kdfID, aeadID, enc, and ct (indicated using the function parse() in pseudocode). The server is then able to find the HPKE private key, skR, corresponding to keyID.

a. If keyID does not identify a key matching the type of kemID, the server returns an error.

b. If kdfID and aeadID identify a combination of KDF and AEAD that the server is unwilling to use with skR, the server returns an error.

2. Build info by concatenating the ASCII-encoded string "message/bhttp request", a zero byte, keyID as an 8-bit integer, plus kemID, kdfID, and aeadID as three 16-bit integers.

3. Create a receiving HPKE context by invoking SetupBaseR() (Section 5.1.1 of [HPKE]) with skR, enc, and info. This produces a context rctxt.

4. Decrypt ct by invoking the Open() method on rctxt (Section 5.2 of [HPKE]), with an empty associated data aad, yielding request or an error on failure. If decryption fails, the server returns an error.

In pseudocode, this procedure is as follows:

```
keyID, kemID, kdfID, aeadID, enc, ct = parse(enc_request)
info = concat(encode_str("message/bhttp request"),
              encode(1, 0),
              encode(1, keyID),
              encode(2, kemID),
              encode(2, kdfID),
              encode(2, aeadID))
rctxt = SetupBaseR(enc, skR, info)
request, error = rctxt.Open([], ct)
```

4.2. Encapsulation of Responses

Given an HPKE context context, a request message request, and a response response, servers generate an Encapsulated Response enc_response as follows:

1. Export a secret secret from context, using the string "message/bhttp response" as context. The length of this secret is max(Nn, Nk), where Nn and Nk are the length of AEAD key and nonce associated with context.

2. Generate a random value of length max(Nn, Nk) bytes, called response_nonce.

3. Extract a pseudorandom key prk using the Extract function provided by the KDF algorithm associated with context. The ikm input to this function is secret; the salt input is the concatenation of enc (from enc_request) and response_nonce

4. Use the Expand function provided by the same KDF to extract an AEAD key key, of length Nk - the length of the keys used by the AEAD associated with context. Generating key uses a label of "key".

5. Use the same Expand function to extract a nonce nonce of length Nn - the length of the nonce used by the AEAD. Generating nonce uses a label of "nonce".

6. Encrypt response, passing the AEAD function Seal the values of key, nonce, empty aad, and a pt input of request, which yields ct.

7. Concatenate response_nonce and ct, yielding an Encapsulated Response enc_response. Note that response_nonce is of fixed-length, so there is no ambiguity in parsing either response_nonce or ct.

In pseudocode, this procedure is as follows:

```
secret = context.Export("message/bhttp response", Nk)
response_nonce = random(max(Nn, Nk))
salt = concat(enc, response_nonce)
prk = Extract(salt, secret)
aead_key = Expand(prk, "key", Nk)
aead_nonce = Expand(prk, "nonce", Nn)
ct = Seal(aead_key, aead_nonce, "", response)
enc_response = concat(response_nonce, ct)
```

Clients decrypt an Encapsulated Response by reversing this process. That is, they first parse enc_response into response_nonce and ct. They then follow the same process to derive values for aead_key and aead_nonce.

The client uses these values to decrypt ct using the Open function provided by the AEAD. Decrypting might produce an error, as follows:

```
reponse, error = Open(aead_key, aead_nonce, "", ct)
```

5. HTTP Usage

A client interacts with the Oblivious Relay Resource by constructing an Encapsulated Request. This Encapsulated Request is included as the content of a POST request to the Oblivious Relay Resource. This request **MUST** only contain those fields necessary to carry the Encapsulated Request: a method of POST, a target URI of the Oblivious Relay Resource, a header field containing the content type (see (Section 6), and the Encapsulated Request as the request content. In the request to the Oblivious Relay Resource, clients **MAY** include additional fields. However, those fields **MUST** be independent

of the Encapsulated Request and **MUST** be fields that the Oblivious Relay Resource will remove before forwarding the Encapsulated Request towards the target, such as the Connection or Proxy-Authorization header fields [SEMANTICS].

The client role in this protocol acts as an HTTP client both with respect to the Oblivious Relay Resource and the Target Resource. For the request the clients makes to the Target Resource, this diverges from typical HTTP assumptions about the use of a connection (see Section 3.3 of [HTTP]) in that the request and response are encrypted rather than sent over a connection. The Oblivious Relay Resource and the Oblivious Gateway Resource also act as HTTP clients toward the Oblivious Gateway Resource and Target Resource respectively.

The Oblivious Relay Resource interacts with the Oblivious Gateway Resource as an HTTP client by constructing a request using the same restrictions as the client request, except that the target URI is the Oblivious Gateway Resource. The content of this request is copied from the client. The Oblivious Relay Resource **MUST NOT** add information to the request without the client being aware of the type of information that might be added; see Section 7.2 for more information on relay responsibilities.

When a response is received from the Oblivious Gateway Resource, the Oblivious Relay Resource forwards the response according to the rules of an HTTP proxy; see Section 7.6 of [HTTP].

An Oblivious Gateway Resource, if it receives any response from the Target Resource, sends a single 200 response containing the encapsulated response. Like the request from the client, this response **MUST** only contain those fields necessary to carry the encapsulated response: a 200 status code, a header field indicating the content type, and the encapsulated response as the response content. As with requests, additional fields **MAY** be used to convey information that does not reveal information about the encapsulated response.

An Oblivious Gateway Resource acts as a gateway for requests to the Target Resource (see Section 7.6 of [HTTP]). The one exception is that any information it might forward in a response **MUST** be encapsulated, unless it is responding to errors it detects before removing encapsulation of the request; see Section 5.2.

5.1. Informational Responses

This encapsulation does not permit progressive processing of responses. Though the binary HTTP response format does support the inclusion of informational (1xx) status codes, the AEAD

encapsulation cannot be removed until the entire message is received.

In particular, the Expect header field with 100-continue (see Section 10.1.1 of [HTTP]) cannot be used. Clients **MUST NOT** construct a request that includes a 100-continue expectation; the Oblivious Gateway Resource **MUST** generate an error if a 100-continue expectation is received.

5.2. Errors

A server that receives an invalid message for any reason **MUST** generate an HTTP response with a 4xx status code.

Errors detected by the Oblivious Relay Resource and errors detected by the Oblivious Gateway Resource before removing protection (including being unable to remove encapsulation for any reason) result in the status code being sent without protection in response to the POST request made to that resource.

Errors detected by the Oblivious Gateway Resource after successfully removing encapsulation and errors detected by the Target Resource **MUST** be sent in an Encapsulated Response.

6. Media Types

Media types are used to identify Encapsulated Requests and Responses.

Evolution of the format of Encapsulated Requests and Responses is supported through the definition of new formats that are identified by new media types.

6.1. message/ohttp-req Media Type

The "message/ohttp-req" identifies an encrypted binary HTTP request. This is a binary format that is defined in Section 4.1.

Type name: message

Subtype name: ohttp-req

Required parameters: N/A

Optional parameters: None

Encoding considerations: only "8bit" or "binary" is permitted

Security considerations: see Section 7

Interoperability considerations:
N/A

Published specification: this specification

Applications that use this media type: N/A

Fragment identifier considerations: N/A

Additional information:

Magic number(s): N/A

Deprecated alias names for this type: N/A

File extension(s): N/A

Macintosh file type code(s): N/A

Person and email address to contact for further information: see Authors' Addresses section

Intended usage: COMMON

Restrictions on usage: N/A

Author: see Authors' Addresses section

Change controller: IESG

6.2. message/ohttp-res Media Type

The "message/ohttp-res" identifies an encrypted binary HTTP response. This is a binary format that is defined in Section 4.2.

Type name: message

Subtype name: ohttp-res

Required parameters: N/A

Optional parameters: None

Encoding considerations: only "8bit" or "binary" is permitted

Security considerations: see Section 7

Interoperability considerations: N/A

Published specification: this specification

Applications that use this media type:
N/A

Fragment identifier considerations: N/A

Additional information:

Magic number(s): N/A

Deprecated alias names for this type: N/A

File extension(s): N/A

Macintosh file type code(s): N/A

Person and email address to contact for further information: see Authors' Addresses section

Intended usage: COMMON

Restrictions on usage: N/A

Author: see Authors' Addresses section

Change controller: IESG

7. Security Considerations

In this design, a client wishes to make a request of a server that is authoritative for the Target Resource. The client wishes to make this request without linking that request with either:

1. The identity at the network and transport layer of the client (that is, the client IP address and TCP or UDP port number the client uses to create a connection).

2. Any other request the client might have made in the past or might make in the future.

In order to ensure this, the client selects a relay (that serves the Oblivious Relay Resource) that it trusts will protect this information by forwarding the Encapsulated Request and Response without passing it to the server (that serves the Oblivious Gateway Resource).

In this section, a deployment where there are three entities is considered:

*A client makes requests and receives responses

*A relay operates the Oblivious Relay Resource

*A server operates both the Oblivious Gateway Resource and the Target Resource

To achieve the stated privacy goals, the Oblivious Relay Resource cannot be operated by the same entity as the Oblivious Gateway Resource. However, colocation of the Oblivious Gateway Resource and Target Resource simplifies the interactions between those resources without affecting client privacy.

As a consequence of this configuration, Oblivious HTTP prevents linkability described above. Informally, this means:

1. Requests and responses are known only to clients and targets in possession of the corresponding response encapsulation key and HPKE keying material. In particular, the Oblivious Relay knows the origin and destination of an Encapsulated Request and Response, yet does not know the decrypted contents. Likewise, targets know only the Oblivious Gateway origin, i.e., the relay, and the decrypted request. Only the client knows both the plaintext request and response.

2. Targets cannot link requests from the same client in the absence of unique per-client keys.

Traffic analysis that might affect these properties are outside the scope of this document; see Section 7.2.3.

A formal analysis of Oblivious HTTP is in [OHTTP-ANALYSIS].

7.1. Client Responsibilities

Clients **MUST** ensure that the key configuration they select for generating Encapsulated Requests is integrity protected and authenticated so that it can be attributed to the Oblivious Gateway Resource; see Section 3.

Since clients connect directly to the relay instead of the target, application configurations wherein clients make policy decisions about target connections, e.g., to apply certificate pinning, are incompatible with Oblivious HTTP. In such cases, alternative technologies such as HTTP CONNECT (Section 9.3.6 of [HTTP]) can be used. Applications could implement related policies on key configurations and relay connections, though these might not provide the same properties as policies enforced directly on target connections. When this difference is relevant, applications can instead connect directly to the target at the cost of either privacy or performance.

Clients **MUST NOT** include identifying information in the request that is encrypted. Identifying information includes cookies [COOKIES],

authentication credentials or tokens, and any information that might reveal client-specific information such as account credentials.

Clients cannot carry connection-level state between requests as they only establish direct connections to the relay responsible for the Oblivious Relay resource. However, clients need to ensure that they construct requests without any information gained from previous requests. Otherwise, the server might be able to use that information to link requests. Cookies [COOKIES] are the most obvious feature that **MUST NOT** be used by clients. However, clients need to include all information learned from requests, which could include the identity of resources.

Clients **MUST** generate a new HPKE context for every request, using a good source of entropy ([RANDOM]) for generating keys. Key reuse not only risks requests being linked, reuse could expose request and response contents to the relay.

The request the client sends to the Oblivious Relay Resource only requires minimal information; see Section 5. The request that carries the Encapsulated Request and is sent to the Oblivious Relay Resource **MUST NOT** include identifying information unless the client ensures that this information is removed by the relay. A client **MAY** include information only for the Oblivious Relay Resource in header fields identified by the Connection header field if it trusts the relay to remove these as required by Section 7.6.1 of [HTTP]. The client needs to trust that the relay does not replicate the source addressing information in the request it forwards.

Clients rely on the Oblivious Relay Resource to forward Encapsulated Requests and responses. However, the relay can only refuse to forward messages, it cannot inspect or modify the contents of Encapsulated Requests or responses.

7.2. Relay Responsibilities

The relay that serves the Oblivious Relay Resource has a very simple function to perform. For each request it receives, it makes a request of the Oblivious Gateway Resource that includes the same content. When it receives a response, it sends a response to the client that includes the content of the response from the Oblivious Gateway Resource.

When forwarding a request, the relay **MUST** follow the forwarding rules in Section 7.6 of [HTTP]. A generic HTTP intermediary implementation is suitable for the purposes of serving an Oblivious Relay Resource, but additional care is needed to ensure that client privacy is maintained.

Firstly, a generic implementation will forward unknown fields. For Oblivious HTTP, a Oblivious Relay Resource **SHOULD NOT** forward unknown fields. Though clients are not expected to include fields that might contain identifying information, removing unknown fields removes this privacy risk.

Secondly, generic implementations are often configured to augment requests with information about the client, such as the Via field or the Forwarded field [FORWARDED]. A relay **MUST NOT** add information when forwarding requests that might be used to identify clients, with the exception of information that a client is aware of.

Finally, a relay can also generate responses, though it assumed to not be able to examine the content of a request (other than to observe the choice of key identifier, KDF, and AEAD), so it is also assumed that it cannot generate an Encapsulated Response.

7.2.1. Differential Treatment

A relay **MAY** add information to requests if the client is aware of the nature of the information that could be added. The client does not need to be aware of the exact value added for each request, but needs to know the range of possible values the relay might use. It is important to note that information added by the relay can reduce the size of the anonymity set of clients at a gateway.

Moreover, relays **MAY** apply differential treatment to clients that engage in abusive behavior, e.g., by sending too many requests in comparison to other clients, or as a response to rate limits signalled from the gateway. Any such differential treatment can reveal information to the gateway that would not be revealed otherwise and therefore reduce the size of the anonymity set of clients using a gateway. For example, if a relay chooses to rate limit or block an abusive client, this means that any client requests which are not treated this way are known to be non-abusive by the gateway. Clients should consider the likelihood of such differential treatment and the privacy risks when using a relay.

Some patterns of abuse cannot be detected without access to the request that is made to the target. This means that only the gateway or target are in a position to identify abuse. A gateway **MAY** send signals toward the relay to provide feedback about specific requests. A relay that acts on this feedback could - either inadvertently or by design - lead to clients being deanonymized.

7.2.2. Denial of Service

As there are privacy benefits from having a large rate of requests forwarded by the same relay (see Section 7.2.3), servers that operate the Oblivious Gateway Resource might need an arrangement

with Oblivious Relay Resources. This arrangement might be necessary to prevent having the large volume of requests being classified as an attack by the server.

If a server accepts a larger volume of requests from a relay, it needs to trust that the relay does not allow abusive levels of request volumes from clients. That is, if a server allows requests from the relay to be exempt from rate limits, the server might want to ensure that the relay applies a rate limiting policy that is acceptable to the server.

Servers that enter into an agreement with a relay that enables a higher request rate might choose to authenticate the relay to enable the higher rate.

7.2.3. Linkability Through Traffic Analysis

This document assumes that all communication between different entities is protected by HTTPS. This protects information about which resources are the subject of request and prevents a network observer from being able to trivially correlate messages on either side of a relay.

As the time at which Encapsulated Request or response messages are sent can reveal information to a network observer. Though messages exchanged between the Oblivious Relay Resource and the Oblivious Gateway Resource might be sent in a single connection, traffic analysis could be used to match messages that are forwarded by the relay.

A relay could, as part of its function, add delays in order to increase the anonymity set into which each message is attributed. This could latency to the overall time clients take to receive a response, which might not be what some clients want.

A relay can use padding to reduce the effectiveness of traffic analysis. Padding is a capability provided by binary HTTP messages; see Section 3.8 of [BINARY].

A relay that forwards large volumes of exchanges can provide better privacy by providing larger sets of messages that need to be matched.

7.3. Server Responsibilities

A server that operates both Oblivious Gateway and Target Resources is responsible for removing request encryption, generating a response to the Encapsulated Request, and encrypting the response.

Servers should account for traffic analysis based on response size or generation time. Techniques such as padding or timing delays can help protect against such attacks; see Section 7.2.3.

If separate entities provide the Oblivious Gateway Resource and Target Resource, these entities might need an arrangement similar to that between server and relay for managing denial of service; see Section 7.2.2. It is also necessary to provide confidentiality protection for the unprotected requests and responses, plus protections for traffic analysis; see Section 7.2.3.

An Oblivious Gateway Resource needs to have a plan for replacing keys. This might include regular replacement of keys, which can be assigned new key identifiers. If an Oblivious Gateway Resource receives a request that contains a key identifier that it does not understand or that corresponds to a key that has been replaced, the server can respond with an HTTP 422 (Unprocessable Content) status code.

A server can also use a 422 status code if the server has a key that corresponds to the key identifier, but the Encapsulated Request cannot be successfully decrypted using the key.

A server **MUST** ensure that the HPKE keys it uses are not valid for any other protocol that uses HPKE with the "message/bhttp request" label. Designers of protocols that reuse this encryption format, especially new versions of this protocol, can ensure key diversity by choosing a different label in their use of HPKE. The "message/bhttp response" label was chosen for symmetry only as it provides key diversity only within the HPKE context created using the "message/bhttp request" label; see Section 10.

A server is responsible for either rejecting replayed requests or ensuring that the effect of replays does not adversely affect clients or resources; see Section 7.4.

7.4. Replay Attacks

Encrypted requests can be copied and replayed by the Oblivious Relay resource. The threat model for Oblivious HTTP allows the possibility that an Oblivious Relay Resource might replay requests. Furthermore, if a client sends an Encapsulated Request in TLS early data (see Section 8 of [TLS] and [RFC8470]), a network-based adversary might be able to cause the request to be replayed. In both cases, the effect of a replay attack and the mitigations that might be employed are similar to TLS early data.

It is the responsibility of the application that uses Oblivious HTTP to either reject replayed requests or to ensure that replayed requests have no adverse affects on their operation. This section

describes some approaches that are universally applicable and suggestions for more targeted techniques.

A client or Oblivious Relay Resource **MUST NOT** automatically attempt to retry a failed request unless it receives a positive signal indicating that the request was not processed or forwarded. The HTTP/2 REFUSED_STREAM error code (Section 8.1.4 of [RFC7540]), the HTTP/3 H3_REQUEST_REJECTED error code (Section 8.1 of [QUIC-HTTP]), or a GOAWAY frame with a low enough identifier (in either protocol version) are all sufficient signals that no processing occurred. Connection failures or interruptions are not sufficient signals that no processing occurred.

The anti-replay mechanisms described in Section 8 of [TLS] are generally applicable to Oblivious HTTP requests. The encapsulated keying material (or enc) can be used in place of a nonce to uniquely identify a request. This value is a high-entropy value that is freshly generated for every request, so two valid requests will have different values with overwhelming probability.

The mechanism used in TLS for managing differences in client and server clocks cannot be used as it depends on being able to observe previous interactions. Oblivious HTTP explicitly prevents such linkability.

The considerations in [RFC8470] as they relate to managing the risk of replay also apply, though there is no option to delay the processing of a request.

Limiting requests to those with safe methods might not be satisfactory for some applications, particularly those that involve the submission of data to a server. The use of idempotent methods might be of some use in managing replay risk, though it is important to recognize that different idempotent requests can be combined to be not idempotent.

Even without replay prevention, the server-chosen response_nonce field ensures that responses have unique AEAD keys and nonces even when requests are replayed.

7.4.1. Use of Date for Anti-Replay

Clients **SHOULD** include a Date header field in Encapsulated Requests. Though HTTP requests often do not include a Date header field, the value of this field might be used by a server to limit the amount of requests it needs to track if it needs to prevent replay attacks.

A server can maintain state for requests for a small window of time over which it wishes to accept requests. The server then rejects requests if the request is the same as one that was previously

answered within that time window. Servers can reject requests outside of this window and signal that clients might retry with a different Date header field; see Section 4 of [REQUEST-DATE]. Servers can identify duplicate requests using the encapsulation (enc) value.

Servers **SHOULD** allow for the time it takes requests to arrive from the client, with a time window that is large enough to allow for differences in the clock of clients and servers. How large a time window is needed could depend on the population of clients that the server needs to serve.

Servers **MUST NOT** treat the time window as secret information. An attacker can actively probe the server with specially crafted request timestamps to determine the time window over which the server will accept responses.

[REQUEST-DATE] contains further considerations for the use of the Date request header field. This includes the way in which clients might correct for clock skew and the privacy considerations arising from that usage. Servers that reject requests on the basis of the Date request header field **SHOULD** implement the feedback mechanism in Section 4 of [REQUEST-DATE] to support clock correction by clients.

7.5. Forward Secrecy

This document does not provide forward secrecy for either requests or responses during the lifetime of the key configuration. A measure of forward secrecy can be provided by generating a new key configuration then deleting the old keys after a suitable period.

7.6. Post-Compromise Security

This design does not provide post-compromise security for responses.

A client only needs to retain keying material that might be used compromise the confidentiality and integrity of a response until that response is consumed, so there is negligible risk associated with a client compromise.

A server retains a secret key that might be used to remove protection from messages over much longer periods. A server compromise that provided access to the Oblivious Gateway Resource secret key could allow an attacker to recover the plaintext of all requests sent toward affected keys and all of the responses that were generated.

Even if server keys are compromised, an adversary cannot access messages exchanged by the client with the Oblivious Relay Resource as messages are protected by TLS. Use of a compromised key also

requires that the Oblivious Relay Resource cooperate with the attacker or that the attacker is able to compromise these TLS connections.

The total number of affected messages affected by server key compromise can be limited by regular rotation of server keys.

8. Privacy Considerations

One goal of this design is that independent client requests are only linkable by the chosen key configuration. The Oblivious Relay and Gateway resources can link requests using the same key configuration by matching KeyConfig.key_id, or, if the Target Resource is willing to use trial decryption, a limited set of key configurations that share an identifier. An Oblivious Relay Resource can link requests using the public key corresponding to KeyConfig.key_id.

Request resources are capable of linking requests depending on how KeyConfigs are produced by servers and discovered by clients. Specifically, servers can maliciously construct key configurations to track individual clients. A specific method for a client to acquire key configurations is not included in this specification. Clients need to consider these tracking vectors when choosing a discovery method. Applications using this design should provide accommodations to mitigate tracking using key configurations. [CONSISTENCY] provides an analysis of the options for ensuring the key configurations are consistent between different clients.

9. Operational and Deployment Considerations

This section discusses various operational and deployment considerations.

9.1. Performance Overhead

Using Oblivious HTTP adds both cryptographic and latency to requests relative to a simple HTTP request-response exchange. Deploying relay services that are on path between clients and servers avoids adding significant additional delay due to network topology. A study of a similar system [ODoH] found that deploying proxies close to servers was most effective in minimizing additional latency.

9.2. Resource Mappings

This protocol assumes a fixed, one-to-one mapping between the Oblivious Relay Resource and the Oblivious Gateway Resource. This means that any encrypted request sent to the Oblivious Relay Resource will always be forwarded to the Oblivious Gateway Resource. This constraint was imposed to simplify relay configuration and mitigate against the Oblivious Relay Resource being used as a

generic relay for unknown Oblivious Gateway Resources. The relay will only forward for Oblivious Gateway Resources that it has explicitly configured and allowed.

It is possible for a server to be configured with multiple Oblivious Relay Resources, each for a different Oblivious Gateway Resource as needed. If the goal is to support a large number of Oblivious Gateway Resources, clients might be provided with a URI template [TEMPLATE], from which multiple Oblivious Relay Resources could be constructed.

9.3. Network Management

Oblivious HTTP might be incompatible with network interception regimes, such as those that rely on configuring clients with trust anchors and intercepting TLS connections. While TLS might be intercepted successfully, interception middleboxes devices might not receive updates that would allow Oblivious HTTP to be correctly identified using the media types defined in Section 6.

Oblivious HTTP has a simple key management design that is not trivially altered to enable interception by intermediaries. Clients that are configured to enable interception might choose to disable Oblivious HTTP in order to ensure that content is accessible to middleboxes.

10. Repurposing the Encapsulation Format

The encrypted payload of an OHTTP request and response is a binary HTTP message [BINARY]. Client and target agree on this encrypted payload type by specifying the media type "message/bhttp" in the HPKE info string and HPKE export context string for request and response encryption, respectively.

Future specifications may repurpose the encapsulation mechanism described in Section 4, provided that the content type of the encrypted payload is appropriately reflected in the HPKE info and context strings. For example, if a future specification were to use the encryption mechanism in this specification for DNS messages, identified by the "application/dns-message" media type, then the HPKE info string **SHOULD** be "application/dns-message request" for request encryption, and the HPKE export context string should be "application/dns-message response" for response encryption.

11. IANA Considerations

Please update the "Media Types" registry at https://www.iana.org/assignments/media-types with the registration information in Section 6 for the media types "message/ohttp-req", "message/ohttp-res", and "application/ohttp-keys".

12. References

12.1. Normative References

[ASCII] Cerf, V., "ASCII format for network interchange", STD 80, RFC 20, DOI 10.17487/RFC0020, October 1969, <https://www.rfc-editor.org/rfc/rfc20>.

[BINARY] Thomson, M. and C. A. Wood, "Binary Representation of HTTP Messages", Work in Progress, Internet-Draft, draft-ietf-httpbis-binary-message-06, 6 July 2022, <https://datatracker.ietf.org/doc/html/draft-ietf-httpbis-binary-message-06>.

[HPKE] Barnes, R. L., Bhargavan, K., Lipp, B., and C. A. Wood, "Hybrid Public Key Encryption", Work in Progress, Internet-Draft, draft-irtf-cfrg-hpke-12, 2 September 2021, <https://datatracker.ietf.org/doc/html/draft-irtf-cfrg-hpke-12>.

[HTTP] Fielding, R., Ed., Nottingham, M., Ed., and J. Reschke, Ed., "HTTP Semantics", STD 97, RFC 9110, DOI 10.17487/RFC9110, June 2022, <https://www.rfc-editor.org/rfc/rfc9110>.

[QUIC] Iyengar, J., Ed. and M. Thomson, Ed., "QUIC: A UDP-Based Multiplexed and Secure Transport", RFC 9000, DOI 10.17487/RFC9000, May 2021, <https://www.rfc-editor.org/rfc/rfc9000>.

[QUIC-HTTP] Bishop, M., "HTTP/3", Work in Progress, Internet-Draft, draft-ietf-quic-http-34, 2 February 2021, <https://datatracker.ietf.org/doc/html/draft-ietf-quic-http-34>.

[REQUEST-DATE] Thomson, M., "Using The Date Header Field In HTTP Requests", Work in Progress, Internet-Draft, draft-thomson-httpapi-date-requests-00, 8 February 2022, <https://datatracker.ietf.org/doc/html/draft-thomson-httpapi-date-requests-00>.

[RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", BCP 14, RFC 2119, DOI 10.17487/RFC2119, March 1997, <https://www.rfc-editor.org/rfc/rfc2119>.

[RFC7540] Belshe, M., Peon, R., and M. Thomson, Ed., "Hypertext Transfer Protocol Version 2 (HTTP/2)", RFC 7540, DOI 10.17487/RFC7540, May 2015, <https://www.rfc-editor.org/rfc/rfc7540>.

[RFC8174]
Leiba, B., "Ambiguity of Uppercase vs Lowercase in RFC 2119 Key Words", BCP 14, RFC 8174, DOI 10.17487/RFC8174, May 2017, <https://www.rfc-editor.org/rfc/rfc8174>.

[RFC8470] Thomson, M., Nottingham, M., and W. Tarreau, "Using Early Data in HTTP", RFC 8470, DOI 10.17487/RFC8470, September 2018, <https://www.rfc-editor.org/rfc/rfc8470>.

[TLS] Rescorla, E., "The Transport Layer Security (TLS) Protocol Version 1.3", RFC 8446, DOI 10.17487/RFC8446, August 2018, <https://www.rfc-editor.org/rfc/rfc8446>.

12.2. Informative References

[CONSISTENCY] Davidson, A., Finkel, M., Thomson, M., and C. A. Wood, "Key Consistency and Discovery", Work in Progress, Internet-Draft, draft-wood-key-consistency-02, 4 March 2022, <https://datatracker.ietf.org/doc/html/draft-wood-key-consistency-02>.

[COOKIES] Barth, A., "HTTP State Management Mechanism", RFC 6265, DOI 10.17487/RFC6265, April 2011, <https://www.rfc-editor.org/rfc/rfc6265>.

[Dingledine2004] Dingledine, R., Mathewson, N., and P. Syverson, "Tor: The Second-Generation Onion Router", August 2004, <https://svn.torproject.org/svn/projects/design-paper/tor-design.html>.

[FORWARDED] Petersson, A. and M. Nilsson, "Forwarded HTTP Extension", RFC 7239, DOI 10.17487/RFC7239, June 2014, <https://www.rfc-editor.org/rfc/rfc7239>.

[ODoH]
Singanamalla, S., Chunhapanya, S., Vavrusa, M., Verma, T., Wu, P., Fayed, M., Heimerl, K., Sullivan, N., and C. A. Wood, "Oblivious DNS over HTTPS (ODoH): A Practical Privacy Enhancement to DNS", 7 January 2021, <https://www.petsymposium.org/2021/files/papers/issue4/popets-2021-0085.pdf>.

[ODOH] Kinnear, E., McManus, P., Pauly, T., Verma, T., and C. A. Wood, "Oblivious DNS over HTTPS", Work in Progress, Internet-Draft, draft-pauly-dprive-oblivious-doh-11, 17

February 2022, <https://datatracker.ietf.org/doc/html/draft-pauly-dprive-oblivious-doh-11>.

[OHTTP-ANALYSIS] Hoyland, J., "Tamarin Model of Oblivious HTTP", 23 August 2021, <https://github.com/cloudflare/ohttp-analysis>.

[PRIO] Corrigan-Gibbs, H. and D. Boneh, "Prio: Private, Robust, and Scalable Computation of Aggregate Statistics", 14 March 2017, <https://crypto.stanford.edu/prio/paper.pdf>.

[RANDOM] Eastlake 3rd, D., Schiller, J., and S. Crocker, "Randomness Requirements for Security", BCP 106, RFC 4086, DOI 10.17487/RFC4086, June 2005, <https://www.rfc-editor.org/rfc/rfc4086>.

[RFC6265] Barth, A., "HTTP State Management Mechanism", RFC 6265, DOI 10.17487/RFC6265, April 2011, <https://www.rfc-editor.org/rfc/rfc6265>.

[RFC7838] Nottingham, M., McManus, P., and J. Reschke, "HTTP Alternative Services", RFC 7838, DOI 10.17487/RFC7838, April 2016, <https://www.rfc-editor.org/rfc/rfc7838>.

[SEMANTICS] Fielding, R., Ed., Nottingham, M., Ed., and J. Reschke, Ed., "HTTP Semantics", STD 97, RFC 9110, DOI 10.17487/RFC9110, June 2022, <https://www.rfc-editor.org/rfc/rfc9110>.

[TEMPLATE] Gregorio, J., Fielding, R., Hadley, M., Nottingham, M., and D. Orchard, "URI Template", RFC 6570, DOI 10.17487/RFC6570, March 2012, <https://www.rfc-editor.org/rfc/rfc6570>.

[X25519] Langley, A., Hamburg, M., and S. Turner, "Elliptic Curves for Security", RFC 7748, DOI 10.17487/RFC7748, January 2016, <https://www.rfc-editor.org/rfc/rfc7748>.

Appendix A. Complete Example of a Request and Response

A single request and response exchange is shown here. Binary values (key configuration, secret keys, the content of messages, and intermediate values) are shown in hexadecimal. The request and response here are minimal; the purpose of this example is to show the cryptographic operations. In this example, the client is configured with the Oblivious Relay Resource URI of https://proxy.example.org/request.example.net/proxy, and the proxy is configured to map requests to this URI to the Oblivious Gateway URI https://example.com/oblivious/request. The Target Resource URI,

i.e., the resource the client ultimately wishes to query, is https://example.com.

To begin the process, the Oblivious Gateway Resource generates a key pair. In this example the server chooses DHKEM(X25519, HKDF-SHA256) and generates an X25519 key pair [X25519]. The X25519 secret key is:

```
3c168975674b2fa8e465970b79c8dcf09f1c741626480bd4c6162fc5b6a98e1a
```

The Oblivious Gateway Resource constructs a key configuration that includes the corresponding public key as follows:

```
01002031e1f05a740102115220e9af918f738674aec95f54db6e04eb705aae8e
79815500080001000100010003
```

This key configuration is somehow obtained by the client. Then when a client wishes to send an HTTP request of a GET request to the target https://example.com, it constructs the following binary HTTP message:

```
00034745540568747470730b6578616d706c652e636f6d012f
```

The client then reads the Oblivious Gateway Resource key configuration and selects a mutually supported KDF and AEAD. In this example, the client selects HKDF-SHA256 and AES-128-GCM. The client then generates an HPKE sending context that uses the server public key. This context is constructed from the following ephemeral secret key:

```
bc51d5e930bda26589890ac7032f70ad12e4ecb37abb1b65b1256c9c48999c73
```

The corresponding public key is:

```
4b28f881333e7c164ffc499ad9796f877f4e1051ee6d31bad19dec96c208b472
```

And an info parameter of:

```
6d6573736167652f626874747020726571756573740001002000010001
```

Applying the Seal operation from the HPKE context produces an encrypted message, allowing the client to construct the following Encapsulated Request:

```
010020000100014b28f881333e7c164ffc499ad9796f877f4e1051ee6d31bad1
9dec96c208b4726374e469135906992e1268c594d2a10c695d858c40a026e796
5e7d86b83dd440b2c0185204b4d63525
```

The client then sends this to the Oblivious Relay Resource in a POST request, which might look like the following HTTP/1.1 request:

```
POST /request.example.net/proxy HTTP/1.1
Host: proxy.example.org
Content-Type: message/ohttp-req
Content-Length: 78

<content is the Encapsulated Request above>
```

The Oblivious Relay Resource receives this request and forwards it to the Oblivious Gateway Resource, which might look like:

```
POST /oblivious/request HTTP/1.1
Host: example.com
Content-Type: message/ohttp-req
Content-Length: 78

<content is the Encapsulated Request above>
```

The Oblivous Gateway Resource receives this request, selects the key it generated previously using the key identifier from the message, and decrypts the message. As this request is directed to the same server, the Oblivious Gateway Resource does not need to initiate an HTTP request to the Target Resource. The request can be served directly by the Target Resource, which generates a minimal response (consisting of just a 200 status code) as follows:

```
0140c8
```

The response is constructed by extracting a secret from the HPKE context:

```
62d87a6ba569ee81014c2641f52bea36
```

The key derivation for the Encapsulated Response uses both the encapsulated KEM key from the request and a randomly selected nonce. This produces a salt of:

```
4b28f881333e7c164ffc499ad9796f877f4e1051ee6d31bad19dec96c208b472
c789e7151fcba46158ca84b04464910d
```

The salt and secret are both passed to the Extract function of the selected KDF (HKDF-SHA256) to produce a pseudorandom key of:

```
979aaeae066cf211ab407b31ae49767f344e1501e475c84e8aff547cc5a683db
```

The pseudorandom key is used with the Expand function of the KDF and an info field of "key" to produce a 16-byte key for the selected AEAD (AES-128-GCM):

```
5d0172a080e428b16d298c4ea0db620d
```

With the same KDF and pseudorandom key, an info field of "nonce" is used to generate a 12-byte nonce:

```
f6bf1aeb88d6df87007fa263
```

The AEAD Seal() function is then used to encrypt the response, which is added to the randomized nonce value to produce the Encapsulated Response:

```
c789e7151fcba46158ca84b04464910d86f9013e404feea014e7be4a441f234f
857fbd
```

The Oblivious Gateway Resource constructs a response with the same content:

```
HTTP/1.1 200 OK
Date: Wed, 27 Jan 2021 04:45:07 GMT
Cache-Control: private, no-store
Content-Type: message/ohttp-res
Content-Length: 38

<content is the Encapsulated Response>
```

The same response might then be generated by the Oblivious Relay Resource which might change as little as the Date header. The client is then able to use the HPKE context it created and the nonce from the Encapsulated Response to construct the AEAD key and nonce and decrypt the response.

Acknowledgments

This design is based on a design for Oblivious DoH, described in [ODOH]. David Benjamin and Eric Rescorla made technical contributions.

Authors' Addresses

Martin Thomson
Mozilla

Email: mt@lowentropy.net

Christopher A. Wood
Cloudflare

Email: caw@heapingbits.net