

Open Pluggable Edge Services  
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A. Beck  
M. Hofmann  
Lucent Technologies  
H. Orman  
Purple Streak Development  
R. Penno  
Nortel Networks  
A. Terzis  
Individual Consultant  
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Requirements for OPES Callout Protocols  
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Abstract

This document specifies the requirements that the OPES (Open Pluggable Edge Services) callout protocol must satisfy in order to support the remote execution of OPES services [1]. The requirements are intended to help evaluating possible protocol candidates and to guide the development of such protocols.

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## 1. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC 2119](#) [2].

## [2](#). Introduction

The Open Pluggable Edge Services (OPES) architecture [[1](#)] enables cooperative application services (OPES services) between a data provider, a data consumer, and zero or more OPES processors. The application services under consideration analyze and possibly transform application-level messages exchanged between the data provider and the data consumer.

The execution of such services is governed by a set of rules installed on the OPES processor. The rules enforcement can trigger the execution of service applications local to the OPES processor. Alternatively, the OPES processor can distribute the responsibility of service execution by communicating and collaborating with one or more remote callout servers. As described in [[1](#)], an OPES processor communicates with and invokes services on a callout server by using a callout protocol. This document presents the requirements for such a protocol.

The requirements in this document are divided into three categories - functional requirements, performance requirements, and security requirements. Each requirement is presented as one or more statements, followed by brief explanatory material as appropriate.

### [3. Functional Requirements](#)

#### [3.1 Callout Transactions](#)

The OPES callout protocol MUST enable an OPES processor and a callout server to perform callout transactions with the purpose of exchanging partial or complete application-level protocol messages (or modifications thereof). More specifically, the callout protocol MUST enable an OPES processor to forward a partial or complete application message to a callout server so that one or more OPES services can process the forwarded application message (or parts thereof). The result of the service operation may be a modified application message. The callout protocol MUST therefore enable the callout server to return a modified application message or the modified parts of an application message to the OPES processor.

A callout transaction is defined as a message exchange between an OPES processor and a callout server consisting of a callout request and a callout response. Both, the callout request as well as the callout response, MAY each consist of one or more callout protocol messages, i.e. a series of protocol messages.

Callout transactions are always initiated by a callout request from an OPES processor and typically terminated by a callout response from a callout server. The OPES callout protocol MUST, however, also allow either endpoint of a callout transaction to terminate a callout transaction prematurely, i.e. before a callout request or response has been completely received by the corresponding endpoint. The callout protocol MAY provide an explicit (e.g. through a termination message) or implicit (e.g. through a connection tear-down) mechanism to terminate a callout transaction prematurely. Such a mechanism MUST ensure, however, that a premature termination of a callout transaction does not result in the loss of application message data.

A premature termination of a callout transaction is required to support OPES services which may terminate even before they have processed the entire application message. Content analysis services, for example, may be able to classify a Web object after having processed just the first few bytes of a Web object.

The callout protocol MUST further enable a callout server to report back to the OPES processor the result of a callout transaction, e.g. in the form of a status code.

### [3.2](#) Callout Channels

The OPES callout protocol MUST enable an OPES processor and a callout server to perform multiple callout transactions over a callout

channel. A callout channel is defined as a logical connection at the application-layer between an OPES processor and a callout server.

Callout channels MUST always be established by an OPES processor. A callout channel MAY be closed by either endpoint of the callout channel provided that all callout transactions associated with the channel have terminated.

A callout channel MAY have certain parameters associated with it, for example parameters that control the fail-over behavior of channel endpoints. Callout channel parameters MAY be negotiated between OPES processors and callout servers (see [Section 3.10](#)).

### [3.3](#) Reliability

The OPES callout protocol MUST be able to provide ordered reliability for the communication between OPES processor and callout server. Additionally, the callout protocol SHOULD be able to provide unordered reliability.

In order to satisfy the reliability requirements, the callout protocol MAY specify that it must be used with a lower-level transport protocol which provides ordered reliability at the transport-layer.

### [3.4](#) Congestion and Flow Control

The OPES callout protocol MUST ensure that congestion and flow control mechanisms are applied on all callout transactions. For this purpose, the callout protocol MAY specify callout protocol-specific mechanisms or refer to a lower-level transport protocol and discuss how its mechanisms provide for congestion and flow control.

### [3.5](#) Support for Keep-Alive Mechanism

The OPES callout protocol MUST provide an optional keep-alive mechanism which, if used, would allow both endpoints of a callout channel to detect a failure of the other endpoint even in the absence of callout transactions. The callout protocol MAY specify that keep-alive messages be exchanged over existing callout channel connections or a separate connection between OPES processor and callout server.

The detection of a callout server failure may enable an OPES processor to establish a channel connection with a stand-by callout server so that future callout transactions do not result in the loss of application message data. The detection of the failure of an OPES processor may enable a callout server to release resources which would otherwise not be available for callout transactions with other

OPES processors.

### [3.6](#) Operation in NAT Environments

The OPES protocol SHOULD be NAT-friendly, i.e. its operation should not be compromised by the presence of one or more NAT devices in the path between an OPES processor and a callout server.

### [3.7](#) Multiple Callout Servers

The OPES callout protocol MUST allow an OPES processor to simultaneously communicate with more than one callout server.

In larger networks, OPES services are likely to be hosted by different callout servers. Therefore, an OPES processor will likely have to communicate with multiple callout servers. The protocol design MUST enable an OPES processor to do so.

### [3.8](#) Multiple OPES Processors

The OPES callout protocol MUST allow a callout server to simultaneously communicate with more than one OPES processor.

The protocol design MUST support a scenario in which multiple OPES processors use the services of a single callout server.

### [3.9](#) Support for Different Application Protocols

The OPES callout protocol MUST be application protocol-agnostic, i.e. it MUST not make any assumptions about the characteristics of the application-layer protocol used on the data path between data provider and data consumer.

The OPES entities on the data path may use different application-layer protocols, including, but not limited to, HTTP [\[3\]](#) and RTP [\[4\]](#). It would be desirable to be able to use the same OPES callout protocol for any such application-layer protocol.

### [3.10](#) Capability and Parameter Negotiations

The OPES callout protocol MUST support the negotiation of capabilities and callout channel parameters between an OPES processor and a callout server. This implies that the OPES processor and the callout server MUST be able to exchange their capabilities and preferences and engage into a deterministic negotiation process at the end of which the two endpoints have either agreed on the capabilities and parameters to be used for future callout channel transactions or determined that their capabilities are incompatible.



Capabilities and parameters that could be negotiated between an OPES processor and a callout server include (but are not limited to): callout protocol version, transport-layer protocol, fail-over behavior, heartbeat rate for keep-alive messages, security-related parameters etc.

Channel parameters may also pertain to the characteristics of OPES callout services if, for example, callout channels are associated with one or more specific OPES services. An OPES service-specific parameter may, for example, specify which parts of an application message an OPES service requires for its operation.

Callout channel parameters **MUST** be negotiated on a per-callout channel basis and before any callout transactions are performed over the corresponding channel. Other parameters and capabilities, such as the fail-over behavior, **MAY** be negotiated between the two endpoints independently of callout channels.

The parties to a callout protocol **MAY** use callout channels to negotiate all or some of their capabilities and parameters. Alternatively, a separate control connection **MAY** be used for this purpose.

### [3.11](#) Meta Data and Instructions

The OPES callout protocol **MUST** provide a mechanism for the endpoints of a particular callout transaction to include in callout requests and responses meta data and instructions for the OPES processor or callout server.

Specifically, the callout protocol **MUST** enable an OPES processor to include information about the forwarded application message in a callout request, e.g. in order to specify the type of the forwarded application message or to specify what part(s) of the application message are forwarded to the callout server. Likewise, the callout server **MUST** be able to include information about the returned application message.

The OPES processor **MUST** further be able to include an ordered list of one or more uniquely specified OPES services which are to be performed on the forwarded application message in the specified order. However, as the callout protocol **MAY** also choose to associate callout channels with specific OPES services, there may not be a need to identify OPES service on a per-callout transaction basis.

Additionally, the OPES callout protocol **MUST** allow the callout server to indicate to the OPES processor the cacheability of callout responses. This implies that callout responses may have to carry

cache-control instructions for the OPES processor.

The OPES callout protocol MUST further enable the OPES processor to indicate to the callout server if it has kept a local copy of the forwarded application message (or parts thereof). This information enables the callout server to determine whether the forwarded application message must be returned to the OPES processor even it has not been modified by an OPES service.

The OPES callout protocol MUST also allow OPES processors to comply with the tracing requirements of the OPES architecture as laid out in [1] and [5]. This implies that the callout protocol MUST enable a callout server to convey to the OPES processor information about the OPES service operations performed on the forwarded application message.

### [3.12](#) Asynchronous Message Exchange

The OPES callout protocol MUST support an asynchronous message exchange between an OPES processor and a callout server.

In order to allow asynchronous processing on the OPES processor and callout server, it MUST be possible to separate request issuance from response processing. The protocol MUST therefore allow multiple outstanding requests and provide a method to correlate responses to requests.

Additionally, the callout protocol MUST enable a callout server to respond to a callout request before it has received the entire request.

### [3.13](#) Message Segmentation

The OPES callout protocol MUST allow an OPES processor to forward an application message to a callout server in a series of smaller message fragments. The callout protocol MUST further enable the receiving callout server to assemble the fragmented application message.

Likewise, the callout protocol MUST enable a callout server to return an application message to an OPES processor in a series of smaller message fragments. The callout protocol MUST enable the receiving OPES processor to assemble the fragmented application message.

Depending on the application-layer protocol used on the data path, application messages may be very large in size (for example in the case of audio/video streams) or of unknown size. In both cases, the OPES processor has to initiate a callout transaction before it has

received the entire application message to avoid long delays for the data consumer. The OPES processor MUST therefore be able to forward fragments or chunks of an application message to a callout server as it receives them from the data provider or consumer. Likewise, the callout server MUST be able to process and return application message fragments as it receives them from the OPES processor.

## [4.](#) Performance Requirements

### [4.1](#) Protocol Efficiency

The OPES callout protocol SHOULD be efficient in that it minimizes the additionally introduced latency, for example by minimizing the protocol overhead.

As OPES callout transactions introduce additional latency to application protocol transactions on the data path, callout protocol efficiency is crucial.

## [5. Security Requirements](#)

In the absence of any security mechanisms, sensitive information might be communicated between the OPES processor and the callout server in violation of either endpoint's security and privacy policy through misconfiguration or a deliberate insider attack. By using strong authentication, message encryption, and integrity checks, this threat can be minimized to a smaller set of insiders and/or operator configuration errors.

The OPES processor and the callout servers SHOULD have enforceable policies that limit the parties they communicate with, that determine the protections to use based on identities of the endpoints and other data (such as enduser policies). In order to enforce the policies, they MUST be able to authenticate the callout protocol endpoints using cryptographic methods.

### [5.1 Authentication, Confidentiality, and Integrity](#)

The parties to the callout protocol MUST have a sound basis for binding authenticated identities to the protocol endpoints, and they MUST verify that these identities are consistent with their security

policies.

The OPES callout protocol MUST provide for message authentication, confidentiality, and integrity between the OPES processor and the callout server. It MUST provide mutual authentication. For this purpose, the callout protocol SHOULD use existing security mechanisms. The callout protocol specification is not required to specify the security mechanisms, but it MAY instead refer to a lower-level security protocol and discuss how its mechanisms are to be used with the callout protocol.

## [5.2](#) Hop-by-Hop Confidentiality

If end-to-end encryption is a requirement for the content path, then this confidentiality MUST be extended to the communication between the callout servers and the OPES processor. In order to minimize data exposure, the callout protocol MUST use a different encryption key for each encrypted content stream.

## [5.3](#) Operation Across Un-trusted Domains

The OPES callout protocol MUST operate securely across un-trusted domains between the OPES processor and the callout server.

If the communication channels between the OPES processor and callout server cross outside of the organization taking responsibility for

the OPES services, then endpoint authentication and message protection (confidentiality and integrity) MUST be used.

## [5.4](#) Privacy

Any communication carrying information relevant to privacy policies MUST protect the data using encryption.

## [6](#). Security Considerations

The security requirements for the OPES callout protocol are discussed in [Section 5](#).

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#### Authors' Addresses

Andre Beck  
Lucent Technologies  
101 Crawfords Corner Road  
Holmdel, NJ 07733  
US

E-Mail: [abeck@bell-labs.com](mailto:abeck@bell-labs.com)

Markus Hofmann  
Lucent Technologies  
Room 4F-513  
101 Crawfords Corner Road  
Holmdel, NJ 07733  
US

Phone: +1 732 332 5983  
E-Mail: [hofmann@bell-labs.com](mailto:hofmann@bell-labs.com)

Hilarie Orman  
Purple Streak Development

E-Mail: [ho@alum.mit.edu](mailto:ho@alum.mit.edu)  
URI: <http://www.purplestreak.com>

Reinaldo Penno  
Nortel Networks  
2305 Mission College Boulevard  
San Jose, CA 95134  
US

E-Mail: [rpenno@nortelnetworks.com](mailto:rpenno@nortelnetworks.com)

Andreas Terzis  
Individual Consultant  
150 Golf Course Dr.  
Rohnert Park, CA 94928  
US

Phone: +1 707 586 8864  
E-Mail: [aterzis@sbcglobal.net](mailto:aterzis@sbcglobal.net)

[Appendix A](#). Acknowledgments

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[Appendix B](#). Change Log

Changes from [draft-ietf-opes-protocol-reqs-01.txt](#)

- o Reworded and clarified several statements of the draft

Changes from [draft-ietf-opes-protocol-reqs-00.txt](#)

- o Aligned terminology with [\[1\]](#)
- o Clarified in [Section 3.11](#) that OCP requests not only have to identify one or more OPES services, but also the order in which the services are to be executed
- o Removed requirement from [Section 4.1](#) that OCP must satisfy performance requirements of the application-layer protocol used between data consumer and provider

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