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**PCEP Procedures and Protocol Extensions for Using PCE as a Central
Controller (PCECC) of LSPs
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Abstract

The Path Computation Element (PCE) is a core component of Software-Defined Networking (SDN) systems. It can compute optimal paths for traffic across a network and can also update the paths to reflect changes in the network or traffic demands.

PCE was developed to derive paths for MPLS Label Switched Paths (LSPs), which are supplied to the head end of the LSP using the Path Computation Element Communication Protocol (PCEP). But SDN has a broader applicability than signaled MPLS and GMPLS traffic-engineered (TE) networks, and the PCE may be used to determine paths in a range of use cases. PCEP has been proposed as a control protocol for use in these environments to allow the PCE to be fully enabled as a central controller.

A PCE-based Central Controller (PCECC) can simplify the processing of a distributed control plane by blending it with elements of SDN and without necessarily completely replacing it. Thus, the LSP can be calculated/set up/initiated and the label forwarding entries can also be downloaded through a centralized PCE server to each network device along the path, while leveraging the existing PCE technologies as much as possible.

This document specifies the procedures and PCEP extensions for using the PCE as the central controller.

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Table of Contents

1.	Introduction	3
1.1.	Requirements Language	5
2.	Terminology	5
3.	Basic PCECC Mode	5
4.	PCEP Requirements	6
5.	Procedures for Using the PCE as a Central Controller (PCECC)	6
5.1.	Stateful PCE Model	6
5.2.	New LSP Functions	6
5.3.	New PCEP Object	7
5.4.	PCECC Capability Advertisement	7
5.5.	LSP Operations	8
5.5.1.	PCE-Initiated PCECC LSP	8
5.5.2.	PCC-Initiated PCECC LSP	12
5.5.3.	Central Controller Instructions	14
5.5.3.1.	Label Download CCI	15
5.5.3.2.	Label Clean up CCI	15
5.5.4.	PCECC LSP Update	16
5.5.5.	Re-Delegation and Clean up	19
5.5.6.	Synchronization of Central Controllers Instructions .	19

5.5.7.	PCECC LSP State Report	20
5.5.8.	PCC-Based Allocations	20
5.5.9.	Binding Label	20
6.	PCEP Messages	21
6.1.	The PCInitiate Message	22
6.2.	The PCRpt Message	23
7.	PCEP Objects	24
7.1.	OPEN Object	24
7.1.1.	PCECC Capability sub-TLV	25
7.2.	PATH-SETUP-TYPE TLV	25
7.3.	CCI Object	26
7.3.1.	Address TLVs	27
8.	Implementation Status	28
8.1.	Huawei's Proof of Concept based on ONOS	29
9.	Security Considerations	29
9.1.	Malicious PCE	29
10.	Manageability Considerations	30
10.1.	Control of Function and Policy	30
10.2.	Information and Data Models	30
10.3.	Liveness Detection and Monitoring	30
10.4.	Verify Correct Operations	30
10.5.	Requirements On Other Protocols	30
10.6.	Impact On Network Operations	30
11.	IANA Considerations	30
11.1.	PCEP TLV Type Indicators	31
11.2.	PATH-SETUP-TYPE-CAPABILITY Sub-TLV Type Indicators	31
11.3.	PCECC-CAPABILITY sub-TLV's Flag field	31
11.4.	Path Setup Type Registry	31
11.5.	PCEP Object	32
11.6.	CCI Object Flag Field	32
11.7.	PCEP-Error Object	32
12.	Acknowledgments	33
13.	References	33
13.1.	Normative References	33
13.2.	Informative References	34
Appendix A.	Contributor Addresses	37
	Authors' Addresses	38

[1.](#) Introduction

The Path Computation Element (PCE) [[RFC4655](#)] was developed to offload the path computation function from routers in an MPLS traffic-engineered network. Since then, the role and function of the PCE has grown to cover a number of other uses (such as GMPLS [[RFC7025](#)]) and to allow delegated control [[RFC8231](#)] and PCE-initiated use of network resources [[RFC8281](#)].

According to [\[RFC7399\]](#), Software-Defined Networking (SDN) refers to a separation between the control elements and the forwarding components so that software running in a centralized system, called a controller, can act to program the devices in the network to behave in specific ways. A required element in an SDN architecture is a component that plans how the network resources will be used and how the devices will be programmed. It is possible to view this component as performing specific computations to place traffic flows within the network given knowledge of the availability of network resources, how other forwarding devices are programmed, and the way that other flows are routed. This is the function and purpose of a PCE, and the way that a PCE integrates into a wider network control system (including an SDN system) is presented in [\[RFC7491\]](#).

In early PCE implementations, where the PCE was used to derive paths for MPLS Label Switched Paths (LSPs), paths were requested by network elements (known as Path Computation Clients (PCCs)), and the results of the path computations were supplied to network elements using the Path Computation Element Communication Protocol (PCEP) [\[RFC5440\]](#). This protocol was later extended to allow a PCE to send unsolicited requests to the network for LSP establishment [\[RFC8281\]](#).

[\[RFC8283\]](#) introduces the architecture for PCE as a central controller as an extension of the architecture described in [\[RFC4655\]](#) and assumes the continued use of PCEP as the protocol used between PCE and PCC. [\[RFC8283\]](#) further examines the motivations and applicability for PCEP as a Southbound Interface (SBI), and introduces the implications for the protocol. [\[I-D.ietf-teas-pcecc-use-cases\]](#) describes the use cases for the PCECC architecture.

A PCE-based Central Controller (PCECC) can simplify the processing of a distributed control plane by blending it with elements of SDN and without necessarily completely replacing it. Thus, the LSP can be calculated/setup/initiated and the label forwarding entries can also be downloaded through a centralized PCE server to each network devices along the path while leveraging the existing PCE technologies as much as possible.

This document specifies the procedures and PCEP extensions for using the PCE as the central controller for static LSPs, where LSPs can be provisioned as explicit label instructions at each hop on the end-to-end path. Each router along the path must be told what label-forwarding instructions to program and what resources to reserve. The PCE-based controller keeps a view of the network and determines the paths of the end-to-end LSPs, and the controller uses PCEP to communicate with each router along the path of the end-to-end LSP.

While this document is focused on the procedures for the static LSPs (referred to as basic PCECC mode in [Section 3](#)), the mechanisms and protocol encodings are specified in such a way that extensions for other use cases are easy to achieve. For example, the extensions for PCECC for Segment Routing (SR) are specified in [\[I-D.zhao-pce-pcep-extension-pce-controller-sr\]](#) and [\[I-D.dhody-pce-pcep-extension-pce-controller-srv6\]](#).

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [BCP 14](#) [[RFC2119](#)] [[RFC8174](#)] when, and only when, they appear in all capitals, as shown here.

2. Terminology

The terminology used in this document is the same as that described in the draft [[RFC8283](#)].

3. Basic PCECC Mode

In this mode, LSPs are provisioned as explicit label instructions at each hop on the end-to-end path. Each router along the path must be told what label forwarding instructions to program and what resources to reserve. The controller uses PCEP to communicate with each router along the path of the end-to-end LSP.

Note that the PCE-based controller will take responsibility for managing some part of the MPLS label space for each of the routers that it controls, and may take wider responsibility for partitioning the label space for each router and allocating different parts for different uses. This is also described in [section 3.1.2. of \[RFC8283\]](#). For the purpose of this document, it is assumed that the label range to be used by a PCE is known and set on both PCEP peers. A future extension could add the capability to advertise the range via possible PCEP extensions as well (see [\[I-D.li-pce-controlled-id-space\]](#)). The rest of the processing is similar to the existing stateful PCE mechanism.

This document also allows a case where the label space is maintained by the PCC itself, and the labels are allocated by the PCC, in this case, the PCE should request the allocation from PCC as described in [Section 5.5.8](#).

4. PCEP Requirements

The following key requirements should be considered when designing the PCECC-based solution:

1. A PCEP speaker supporting this draft needs to have the capability to advertise its PCECC capability to its peers.
2. A PCEP speaker need means to identify PCECC-based LSP in the PCEP messages.
3. PCEP procedures need to allow for PCC-based label allocations.
4. PCEP procedures need to provide a mean to update (or clean up) the label-download entry to the PCC.
5. PCEP procedures need to provide a mean to synchronize the labels between the PCE and the PCC via PCEP messages.

5. Procedures for Using the PCE as a Central Controller (PCECC)

5.1. Stateful PCE Model

Active stateful PCE is described in [[RFC8231](#)]. PCE as a central controller (PCECC) reuses the existing active stateful PCE mechanism as much as possible to control LSPs.

5.2. New LSP Functions

Several new functions are required in PCEP to support PCECC. This document extends the existing messages to support the new functions required by PCECC:

(PCInitiate): a PCEP message described in [[RFC8281](#)]. PCInitiate message is used to set up PCE-Initiated LSP based on PCECC mechanism. It is also extended for Central Controller Instructions (CCI) (download or clean up the Label forwarding instructions in the context of this document) on all nodes along the path.

(PCRpt): a PCEP message described in [[RFC8231](#)]. PCRpt message is used to send PCECC LSP Reports. It is also extended to report the set of Central Controller Instructions (CCI) (label forwarding instructions in the context of this document) received from the PCE. See [Section 5.5.6](#) for more details.

(PCUpd): a PCEP message described in [[RFC8231](#)]. PCUpd message is used to send PCECC LSP Update.

The new functions defined in this document are mapped onto the PCEP messages as shown in Table 1.

Function	Message
PCECC Capability advertisement	Open
Label entry Add	PCInitiate
Label entry Clean up	PCInitiate
PCECC Initiated LSP	PCInitiate
PCECC LSP Update	PCUpd
PCECC LSP State Report	PCRpt
PCECC LSP Delegation	PCRpt
PCECC Label Report	PCRpt

Table 1: Functions mapped to the PCEP messages

5.3. New PCEP Object

This document specifies a new PCEP object called CCI (see [Section 7.3](#)) for the encoding of the central controller instructions. In the scope of this document, this is limited to Label forwarding instructions. Future documents can create new CCI object-types for other types of central controller instructions. The CC-ID is the unique identifier for the central controller instructions in PCEP. The PCEP messages are extended in this document to handle the PCECC operations.

5.4. PCECC Capability Advertisement

During the PCEP Initialization Phase, PCEP Speakers (PCE or PCC) advertise their support of PCECC extensions.

This document defines a new Path Setup Type (PST) [[RFC8408](#)] for PCECC, as follows:

- o PST = TBD1: Path is set up via PCECC mode.

A PCEP speaker MUST indicate its support of the function described in this document by sending a PATH-SETUP-TYPE-CAPABILITY TLV in the OPEN object with this new PST included in the PST list.

This document also defines the PCECC Capability sub-TLV [Section 7.1.1](#). PCEP speakers use this sub-TLV to exchange information about their PCECC capability. If a PCEP speaker includes PST=TBD1 in the PST List of the PATH-SETUP-TYPE-CAPABILITY TLV then the receiving peer MUST also include the PCECC Capability sub-TLV (with the L bit set to 1) inside the PATH-SETUP-TYPE-CAPABILITY TLV. If the sub-TLV is absent or the L bit is not set to 1, then the receiving PCEP speaker MUST send a PCErr message with Error-Type 10

(Reception of an invalid object) and Error-Value TBD2 (Missing PCECC Capability sub-TLV) and MUST then close the PCEP session. If a PCEP speaker receives a PATH-SETUP-TYPE-CAPABILITY TLV with a PCECC-CAPABILITY sub-TLV, but the PST list does not contain PST=TBD1, then the PCEP speaker MUST ignore the PCECC-CAPABILITY sub-TLV.

The presence of the PST=TBD1 and PCECC Capability sub-TLV (with the L bit set to 1, see [Section 7.1.1](#)) in a PCC's OPEN Object indicates that the PCC is willing to function as a PCECC client for label download instructions. The presence of the PST=TBD1 and PCECC Capability sub-TLV (with the L bit set to 1) in a PCE's OPEN message indicates that the PCE is interested in function as a PCECC server for label download instructions.

The PCEP extensions for PCECC for label download MUST NOT be used if one or both PCEP Speakers have not included the PST=TBD1 or the PCECC Capability sub-TLV (with the L bit set to 1) in their respective OPEN message. If a PCEP speaker which supports the extensions of this draft but did not advertise this capability attempts a PCECC operation, then a PCErr message with Error-Type=19 (Invalid Operation) and Error-Value=TBD3 (Attempted PCECC operations when PCECC capability was not advertised) MUST be generated by its peer and the PCEP session will be terminated. If a PCEP speaker does not recognize the PCECC Capability sub-TLV, it will ignore the sub-TLV in accordance with [\[RFC8408\]](#) and [\[RFC5440\]](#).

A PCC or a PCE MUST include both the PCECC-CAPABILITY sub-TLV (with the L bit set to 1) and the STATEFUL-PCE-CAPABILITY TLV ([\[RFC8231\]](#)) (with the I flag set [\[RFC8281\]](#)) in the OPEN Object to support the extensions defined in this document. If the PCECC-CAPABILITY sub-TLV is advertised and the STATEFUL-PCE-CAPABILITY TLV is not advertised in the OPEN Object, it MUST send a PCErr message with Error-Type=19 (Invalid Operation) and Error-value=TBD4 (stateful PCE capability was not advertised) and terminate the session. This error is also triggered if the PCECC-CAPABILITY sub-TLV is advertised and the I flag in the STATEFUL-PCE-CAPABILITY TLV is not set.

5.5. LSP Operations

The PCEP messages pertaining to a PCECC MUST include PATH-SETUP-TYPE TLV [\[RFC8408\]](#) in the SRP object to clearly identify that PCECC LSP is intended.

5.5.1. PCE-Initiated PCECC LSP

The LSP Instantiation operation is the same as defined in [\[RFC8281\]](#).

In order to set up a PCE-Initiated LSP based on the PCECC mechanism, a PCE sends PCInitiate message with Path Setup Type set for PCECC (see [Section 7.2](#)) to the ingress PCC.

An LSP-IDENTIFIER TLV MUST be included for PCECC LSPs, the tuple uniquely identifies the LSP in the network. The LSP object is included in the central controller instructions (label download [Section 7.3](#)) to identify the PCECC LSP for this instruction. The PLSP-ID is the original identifier used by the ingress PCC, so a transit/egress LSR could have multiple central controller instructions that have the same PLSP-ID. The PLSP-ID in combination with the source (in LSP-IDENTIFIER TLV) MUST be unique. The PLSP-ID is included for maintainability reasons to ease debugging. As per [\[RFC8281\]](#), the LSP object could also include the SPEAKER-ENTITY-ID TLV to identify the PCE that initiated these instructions. Also, the CC-ID is unique in the PCEP session as described in [Section 7.3](#).

The ingress PCC MUST also set D (Delegate) flag (see [\[RFC8231\]](#)) and C (Create) flag (see [\[RFC8281\]](#)) in the LSP object of the PCRpt message to the PCE. The PCC responds with a PCRpt message with the status set to "GOING-UP" and carrying the assigned PLSP-ID. When the PCE receives this PCRpt message with the PLSP-ID, it assigns labels along the path; and sets up the path by sending a PCInitiate message to each node along the path of the LSP as per the PCECC technique. The CC-ID uniquely identifies the central controller instruction within a PCEP session. Each PCC further responds with the PCRpt messages including the central controller instruction (CCI) and the LSP objects.

Note that the label forwarding instructions (see [Section 5.5.3](#)) from PCECC are sent after the initial PCInitiate and PCRpt message exchange with the ingress PCC. This is done so that the PLSP-ID and other LSP identifiers can be obtained from the ingress and can be included in the label forwarding instruction in the next set of PCInitiate messages.

The ingress node would receive one CCI object with 0 bit (out-label) set. The transit node(s) would receive two CCI objects with the in-label CCI without an 0 bit set and the out-label CCI with 0 bit set. The egress node would receive one CCI object without 0 bit set. A node can determine its role based on the setting of the 0 bit in the CCI object(s) and the LSP-IDENTIFIER TLV in the LSP object.

The LSP deletion operation for PCE-Initiated PCECC LSP is the same as defined in [\[RFC8281\]](#). The PCE should further perform Label entry clean up operation as described in [Section 5.5.3.2](#) for the corresponding LSP.

The PCE-Initiated PCECC LSP setup sequence is shown in Figure 1.

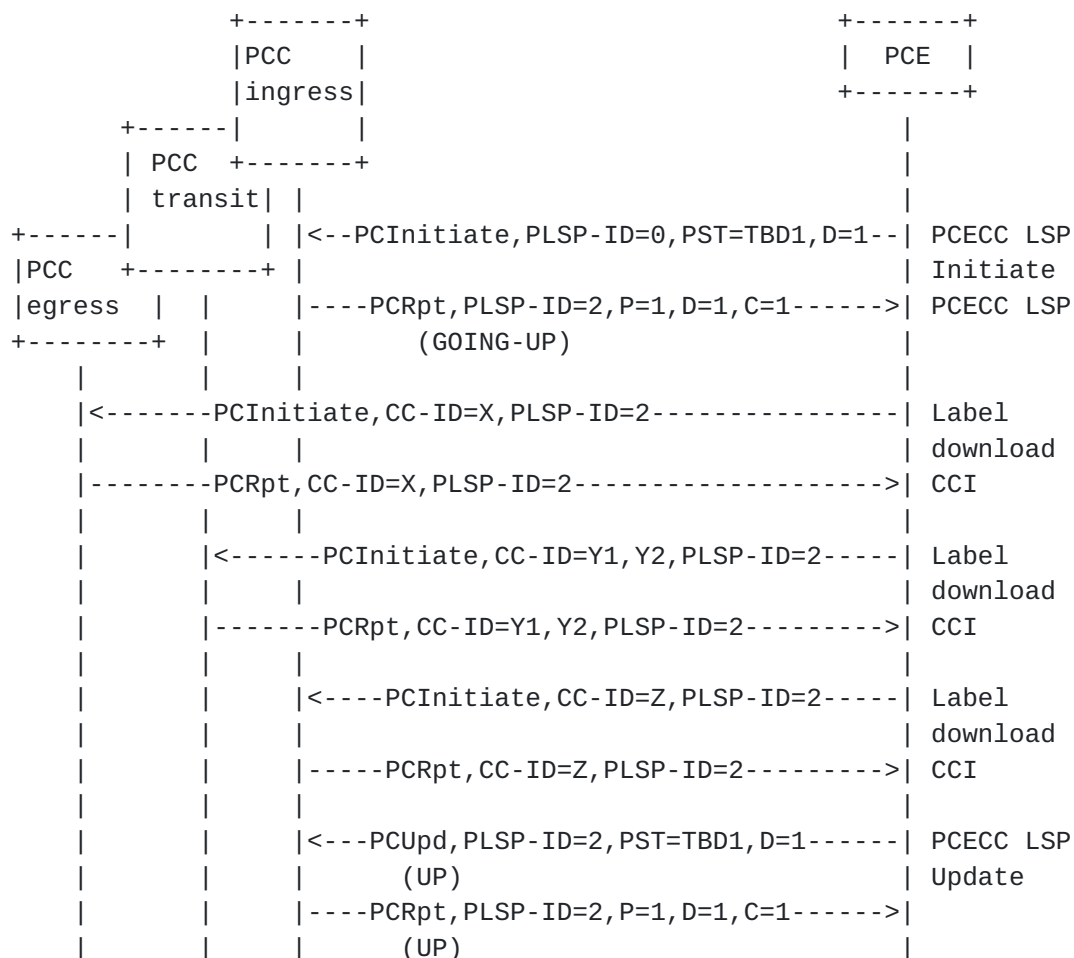


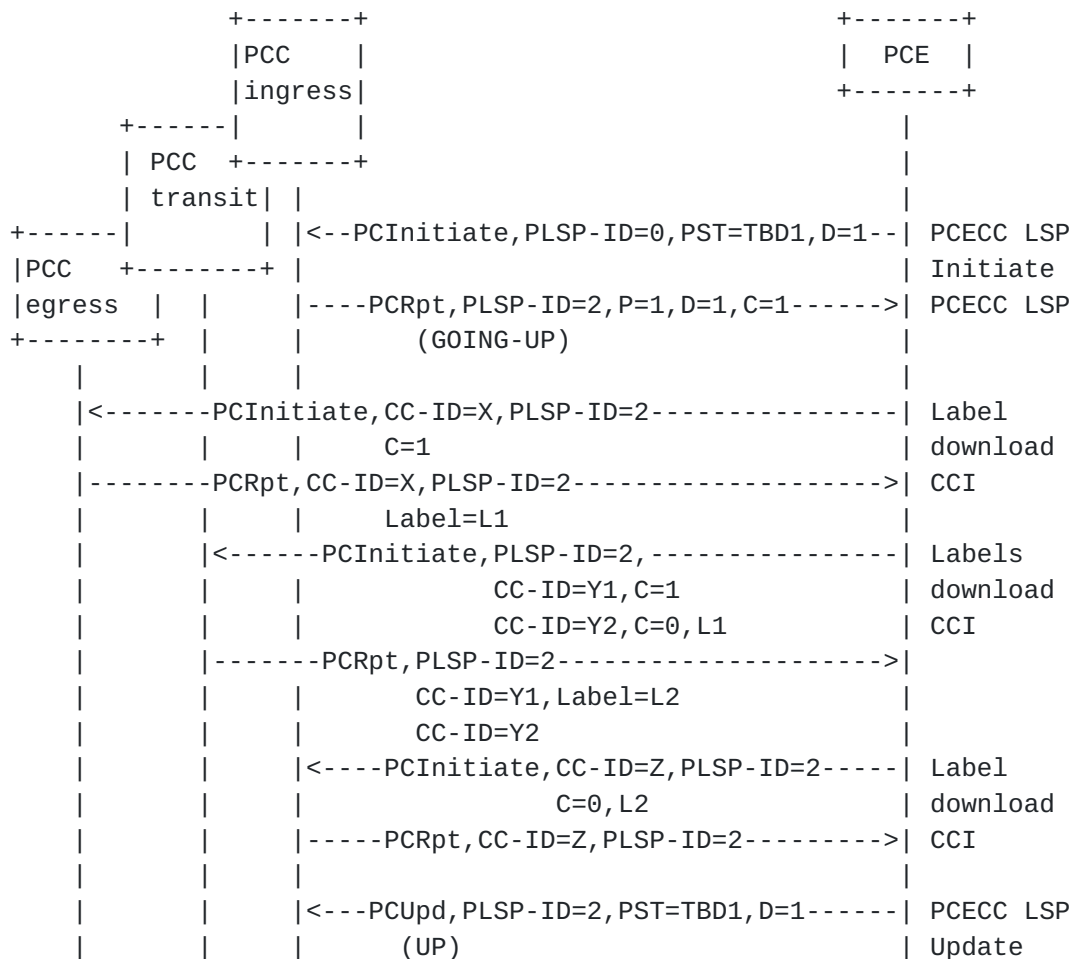
Figure 1: PCE-Initiated PCECC LSP

Once the label operations are completed, the PCE SHOULD send a PCUpd message to the ingress PCC. The PCUpd message is as per [\[RFC8231\]](#).

The PCECC LSPs are considered to be 'up' by default (on receipt of PCUpd message from PCE). The ingress MAY further choose to deploy a data plane check mechanism and report the status back to the PCE via a PCRpt message to make sure that the correct label instructions are made along the path of the PCECC LSP (and it is ready to carry traffic).

In the case where the label allocations are made by the PCC itself (see [Section 5.5.8](#)), the PCE could request an allocation to be made by the PCC, and then the PCC would send a PCRpt with the allocated label encoded in the CC-ID object as shown in Figure 2 in the

configuration sequence from the egress towards the ingress along the path.



- The 0 bit is set as before (and thus not included)

Figure 2: PCE-Initiated PCECC LSP (PCC allocation)

It should be noted that in this example, the request is made to the egress node with the C bit set in the CCI object to indicate that the label allocation needs to be done by the egress and the egress responds with the allocated label to the PCE. The PCE further inform the transit PCC without setting the C bit to 1 in the CCI object for out-label but the C bit is set to 1 for in-label so the transit node make the label allocation (for the in-label) and report to the PCE. Similarly, the C bit is unset towards the ingress to complete all the label allocation for the PCECC LSP.

5.5.2. PCC-Initiated PCECC LSP

In order to set up an LSP based on the PCECC mechanism where the LSP is configured at the PCC, a PCC MUST delegate the LSP by sending a PCRpt message with PST set for PCECC (see [Section 7.2](#)) and D (Delegate) flag (see [\[RFC8231\]](#)) set in the LSP object.

When a PCE receives the initial PCRpt message with D flag and PST Type set to TBD1, it SHOULD calculate the path and assigns labels along the path; and sets up the path by sending a PCInitiate message to each node along the path of the LSP as per the PCECC technique. The CC-ID uniquely identifies the central controller instruction within a PCEP session. Each PCC further responds with the PCRpt messages including the central controller instruction (CCI) and the LSP objects.

Once the central controller instructions (label operations) are completed, the PCE MUST send the PCUpd message to the ingress PCC. Per [\[RFC8231\]](#), this PCUpd message should include the path information calculated by the PCE.

Note that the PCECC LSPs MUST be delegated to a PCE at all times.

The LSP deletion operation for PCECC LSPs is the same as defined in [\[RFC8231\]](#). If the PCE receives a PCRpt message for LSP deletion then it does label clean up operation as described in [Section 5.5.3.2](#) for the corresponding LSP.

The Basic PCECC LSP setup sequence is as shown in Figure 3.

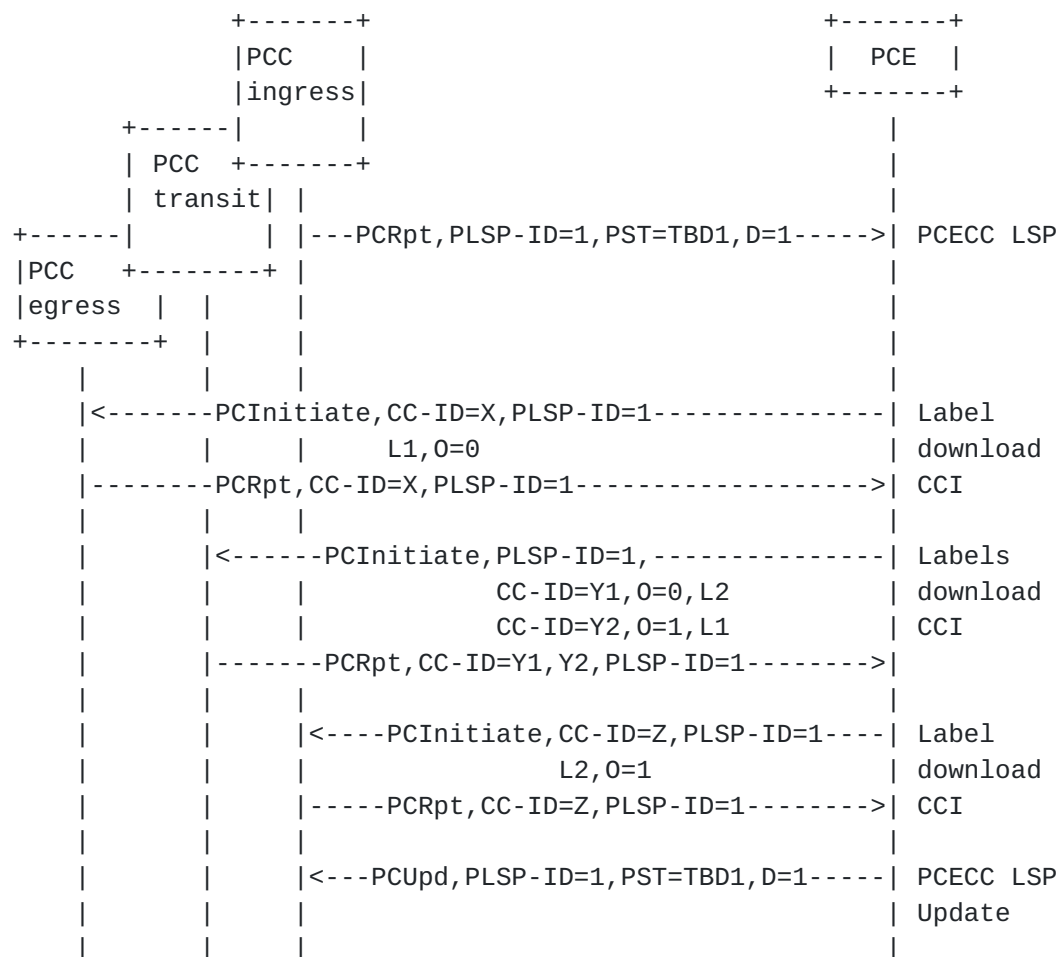
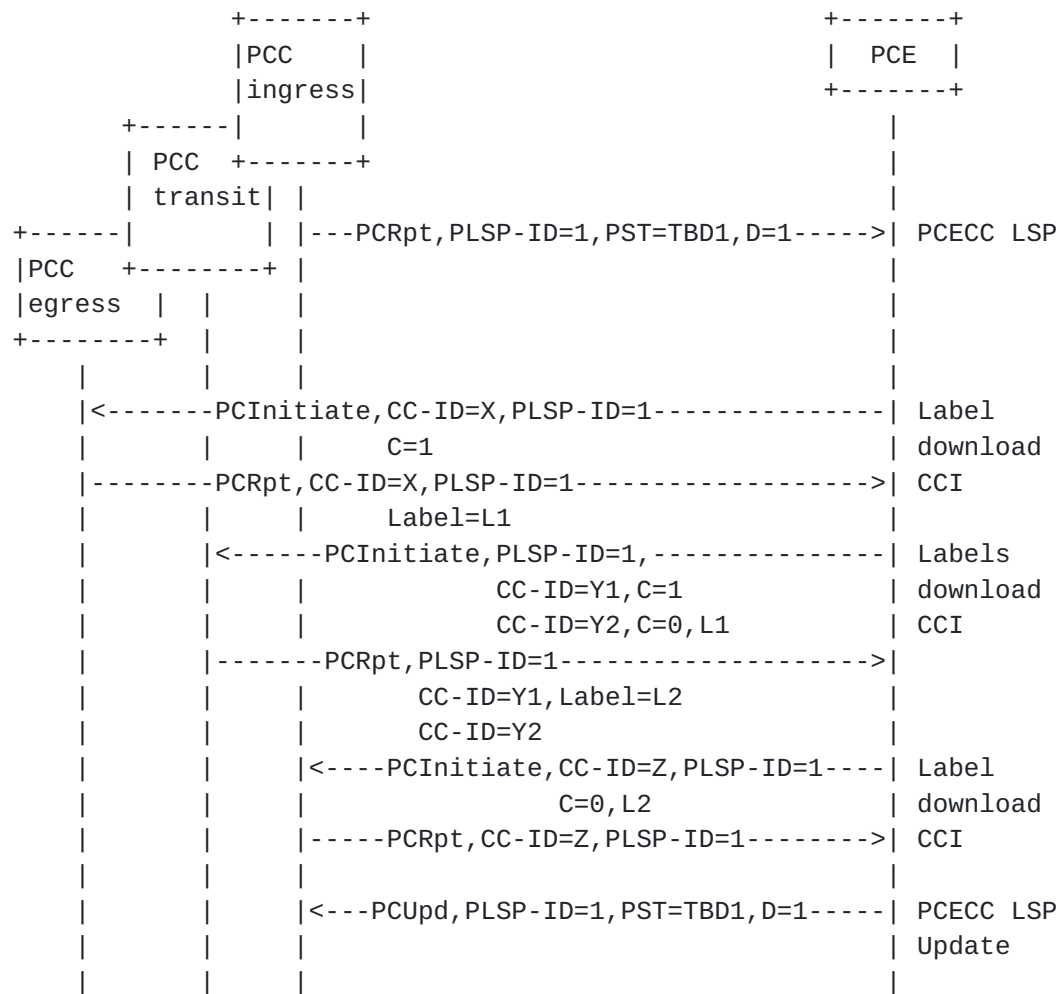


Figure 3: PCC-Initiated PCECC LSP

In the case where the label allocations are made by the PCC itself (see [Section 5.5.8](#)), the PCE could request an allocation to be made by the PCC, and then the PCC would send a PCRpt with the allocated label encoded in the CC-ID object as shown in Figure 4.



- The 0 bit is set as before (and thus not included)

Figure 4: PCC-Initiated PCECC LSP (PCC allocation)

In the case where the label allocations are made by the PCC itself (see [Section 5.5.8](#)), the procedure remains the same, with just an additional constraint on the configuration sequence.

The rest of the PCC-Initiated PCECC LSP setup operations are the same as those described in [Section 5.5.1](#).

5.5.3. Central Controller Instructions

The new central controller instructions (CCI) for the label operations in PCEP is done via the PCInitiate message, by defining a new PCEP Object for CCI operations. The local label range of each PCC is assumed to be known by both the PCC and the PCE.

5.5.3.1. Label Download CCI

In order to set up an LSP based on PCECC, the PCE sends a PCInitiate message to each node along the path to download the Label instruction as described in [Section 5.5.1](#) and [Section 5.5.2](#).

The CCI object MUST be included, along with the LSP object in the PCInitiate message. The LSP-IDENTIFIER TLV MUST be included in the LSP object. The SPEAKER-ENTITY-ID TLV SHOULD be included in the LSP object.

If a node (PCC) receives a PCInitiate message which includes a Label to download, as part of CCI, that is out of the range set aside for the PCE, it MUST send a PCErr message with Error-type=TBD5 (PCECC failure) and Error-value=TBD6 (Label out of range) and MUST include the SRP object to specify the error is for the corresponding label update via PCInitiate message. If a PCC receives a PCInitiate message but fails to download the Label entry, it MUST send a PCErr message with Error-type=TBD5 (PCECC failure) and Error-value=TBD7 (instruction failed) and MUST include the SRP object to specify the error is for the corresponding label update via PCInitiate message.

A new PCEP object for central controller instructions (CCI) is defined in [Section 7.3](#).

5.5.3.2. Label Clean up CCI

In order to delete an LSP based on PCECC, the PCE sends a central controller instructions via a PCInitiate message to each node along the path of the LSP to clean up the Label forwarding instruction.

If the PCC receives a PCInitiate message but does not recognize the label in the CCI, the PCC MUST generate a PCErr message with Error-Type 19(Invalid operation) and Error-Value=TBD8, "Unknown Label" and MUST include the SRP object to specify the error is for the corresponding label clean up (via PCInitiate message).

The R flag in the SRP object defined in [[RFC8281](#)] specifies the deletion of Label Entry in the PCInitiate message.

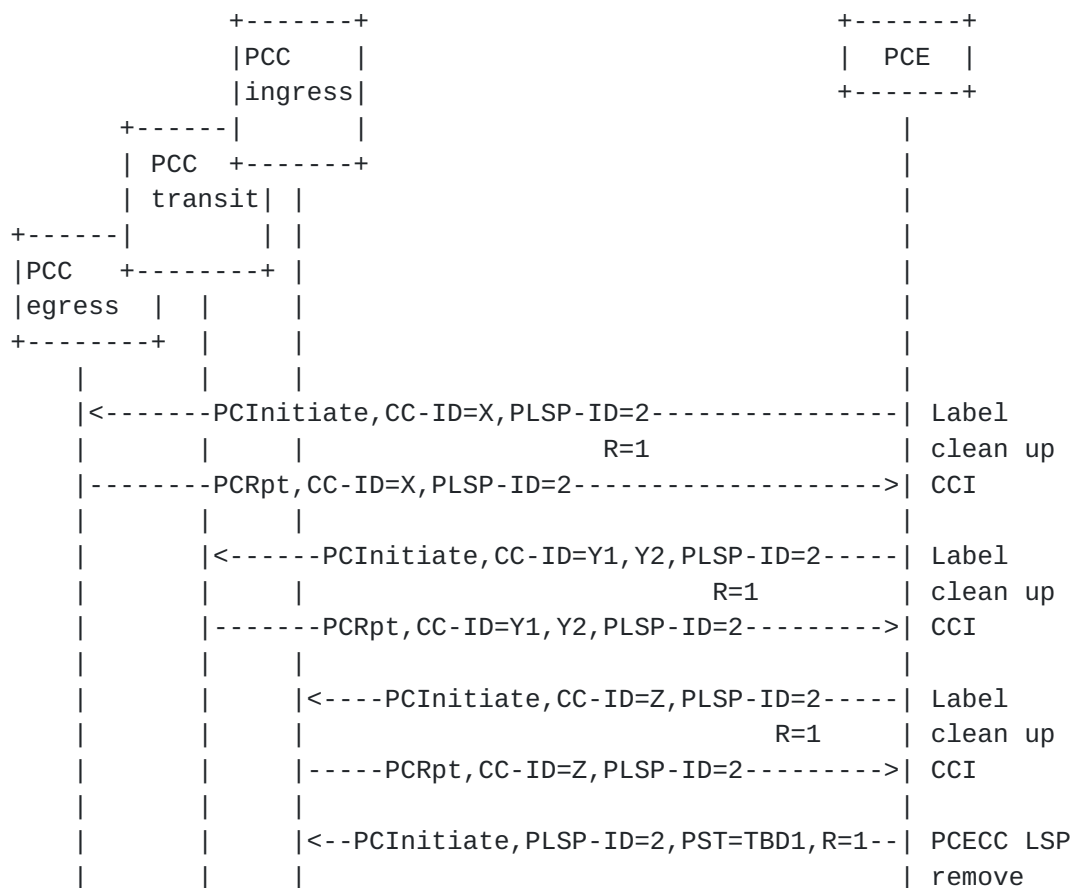


Figure 5: Label Cleanup

As per [RFC8281], following the removal of the Label forwarding instruction, the PCC MUST send a PCRpt message. The SRP object in the PCRpt MUST include the SRP-ID-number from the PCInitiate message that triggered the removal. The R flag in the SRP object MUST be set.

In the case where the label allocation is made by the PCC itself (see [Section 5.5.8](#)), the removal procedure remains the same, adding the sequence constraint.

5.5.4. PCECC LSP Update

In case of a modification of a PCECC LSP with a new path, a PCE sends a PCUpd message to the ingress PCC. But to follow the make-before-break procedures, the PCECC first updates new instructions based on the updated LSP and then update to the ingress to switch traffic, before cleaning up the former instructions. A new CC-ID is used to identify the updated instruction, the existing identifiers in the LSP object identify the existing LSP. Once new instructions are

downloaded, the PCE further updates the new path at the ingress which triggers the traffic switch on the updated path. The ingress PCC acknowledges with a PCRpt message, on receipt of the PCRpt message, the PCE does clean up operation for the former LSP as described in [Section 5.5.3.2](#).

The PCECC LSP Update sequence is shown in Figure 6.

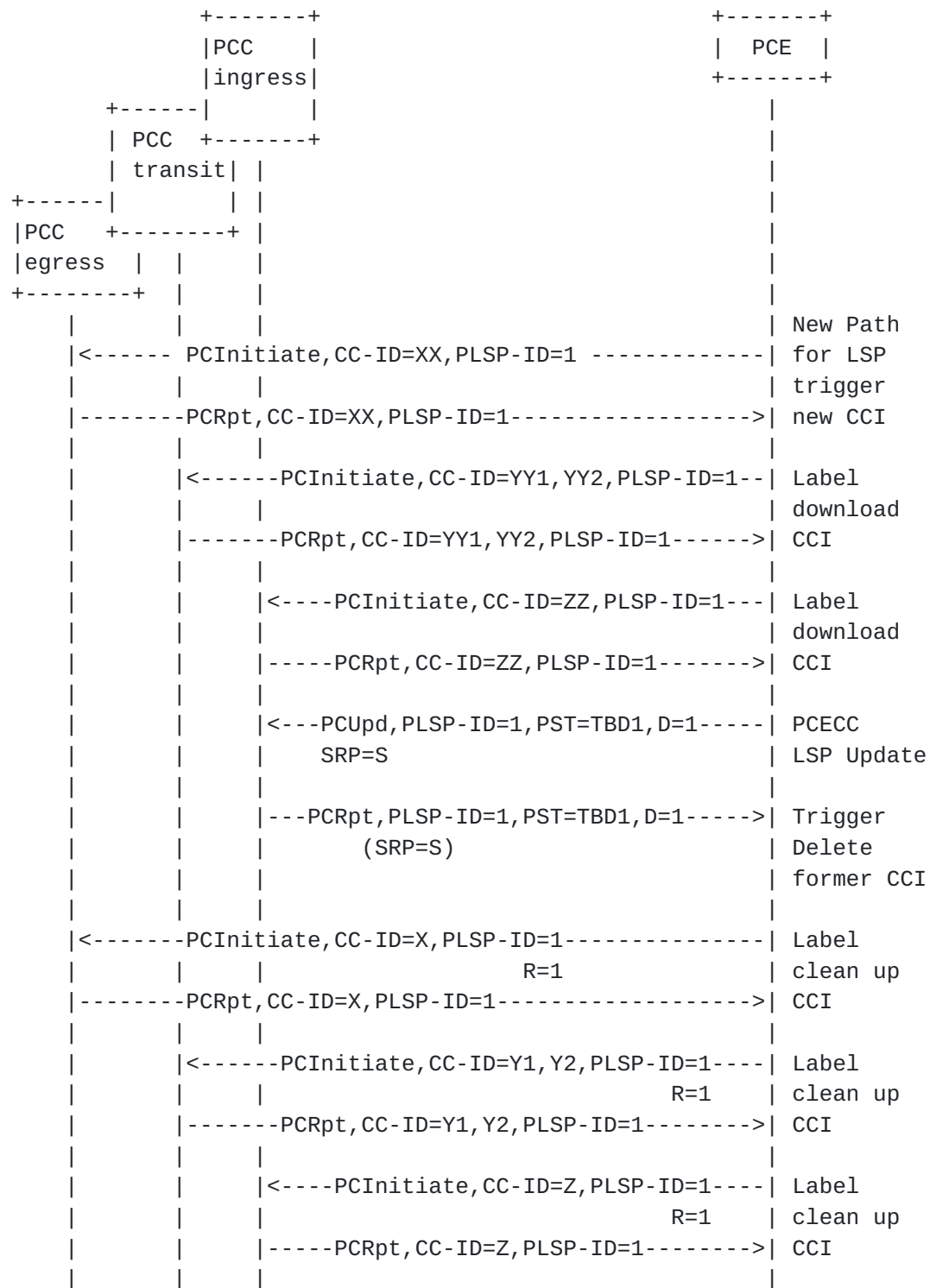


Figure 6: PCECC LSP Update

The modified PCECC LSPs are considered to be 'up' by default. The ingress MAY further choose to deploy a data plane check mechanism and report the status back to the PCE via a PCRpt message.

In the case where the label allocations are made by the PCC itself (see [Section 5.5.8](#)), the procedure remains the same.

5.5.5. Re-Delegation and Clean up

As described in [[RFC8281](#)], a new PCE can gain control over an orphaned LSP. In the case of a PCECC LSP, the new PCE MUST also gain control over the central controller instructions in the same way by sending a PCInitiate message that includes the SRP, LSP, and CCI objects and carries the CC-ID and PLSP-ID identifying the instruction that it wants to take control of.

Further, as described in [[RFC8281](#)], the State Timeout Interval timer ensures that a PCE crash does not result in automatic and immediate disruption for the services using PCE-initiated LSPs. Similarly the central controller instructions are not removed immediately upon PCE failure. Instead, they are cleaned up on the expiration of this timer. This allows for network clean up without manual intervention. The PCC MUST support the removal of CCI as one of the behaviors applied on expiration of the State Timeout Interval timer.

In case of PCC-initiated PCECC LSP, the control over the orphaned LSP at the ingress PCC is taken over by the mechanism specified in [[RFC8741](#)] to request delegation. The control over the central controller instructions is described above using [[RFC8281](#)].

5.5.6. Synchronization of Central Controllers Instructions

The purpose of Central Controllers Instructions synchronization (labels in the context of this document) is to make sure that the PCE's view of CCI (Labels) matches with the PCC's Label allocation. This synchronization is performed as part of the LSP state synchronization as described in [[RFC8231](#)] and [[RFC8232](#)].

As per LSP State Synchronization [[RFC8231](#)], a PCC reports the state of its LSPs to the PCE using PCRpt messages and as per [[RFC8281](#)], PCE would initiate any missing LSPs and/or remove any LSPs that are not wanted. The same PCEP messages and procedures are also used for the Central Controllers Instructions synchronization. The PCRpt message includes the CCI and the LSP object to report the label forwarding instructions. The PCE would further remove any unwanted instructions or initiate any missing instructions.

5.5.7. PCECC LSP State Report

As mentioned before, an ingress PCC MAY choose to apply any OAM mechanism to check the status of LSP in the Data plane and MAY further send its status in a PCRpt message to the PCE.

5.5.8. PCC-Based Allocations

The PCE can request the PCC to allocate the label using the PCInitiate message. The C flag in the CCI object is set to 1 to indicate that the allocation needs to be done by the PCC. The PCC SHOULD allocate the Label and SHOULD report to the PCE using the PCRpt message.

If the value of the Label is 0 and the C flag is set to 1, it indicates that the PCE is requesting the allocation to be done by the PCC. If the Label is 'n' and the C flag is set to 1 in the CCI object, it indicates that the PCE requests a specific value 'n' for the Label. If the allocation is successful, the PCC MUST report via the PCRpt message with the CCI object. Else, it MUST send a PCErr message with Error-Type = TBD5 ("PCECC failure") and Error Value = TBD9 ("Invalid CCI"). If the value of the Label in the CCI object is valid, but the PCC is unable to allocate it, it MUST send a PCErr message with Error-Type = TBD5 ("PCECC failure") and Error Value = TBD10 ("Unable to allocate the specified CCI").

If the PCC wishes to withdraw or modify the previously assigned label, it MUST send a PCRpt message without any Label or with the Label containing the new value respectively in the CCI object. The PCE would further trigger the removal of the central controller instruction as per this document.

5.5.9. Binding Label

As per [[I-D.ietf-pce-binding-label-sid](#)], when a stateful PCE is deployed for setting up TE paths, it may be desirable to report the binding label to the stateful PCE for the purpose of enforcing end-to-end TE. In the case of the PCECC, the binding label may be allocated by the PCE itself as described in this section. This procedure is thus applicable for all path setup types including PCECC.

A P flag in the LSP object is introduced in [[I-D.ietf-pce-sr-path-segment](#)] to indicate the allocation needs to be made by the PCE. A PCC would set this bit to 1 (and carry the TE-PATH-BINDING TLV [[I-D.ietf-pce-binding-label-sid](#)] in the LSP object) to request for allocation of the binding label by the PCE in the PCReq or PCRpt message. A PCE would also set this bit to 1 to

indicate that the binding label is allocated by PCE and encoded in the PCRep, PCUpd, or PCInitiate message (the TE-PATH-BINDING TLV is present in LSP object). Further, a PCE would set this bit to 0 to indicate that the allocation is done by the PCC instead.

The ingress PCC could request the binding label to be allocated by the PCE via a PCRpt message as per [\[RFC8231\]](#). The delegate flag (D-flag) MUST also be set for this LSP. The TE-PATH-BINDING TLV MUST be included with no Binding Value. The PCECC would allocate the binding label and further respond to ingress PCC with PCUpd message as per [\[RFC8231\]](#) and MUST include the TE-PATH-BINDING TLV in an LSP object. The P flag in the LSP object would be set to 1 to indicate that the allocation is made by the PCE.

The PCE could allocate the binding label on its own accord for a PCE-Initiated (or delegated) LSP. The allocated binding label needs to be informed to the PCC. The PCE would use the PCInitiate message [\[RFC8281\]](#) or PCUpd message [\[RFC8231\]](#) towards the PCC and MUST include the TE-PATH-BINDING TLV in the LSP object. The P flag in the LSP object would be set to 1 to indicate that the allocation is made by the PCE.

Before a PCE can allocate a binding label the PCECC capability MUST be exchanged on the PCEP session. Note that the CCI object is not used for binding allocation; this is done to maintain consistency with the rest of the binding label/SID procedures as per [\[I-D.ietf-pce-binding-label-sid\]](#).

6. PCEP Messages

As defined in [\[RFC5440\]](#), a PCEP message consists of a common header followed by a variable-length body made of a set of objects that can be either mandatory or optional. An object is said to be mandatory in a PCEP message when the object must be included for the message to be considered valid. For each PCEP message type, a set of rules is defined that specify the set of objects that the message can carry. An implementation MUST form the PCEP messages using the object ordering specified in this document.

LSP-IDENTIFIERS TLV MUST be included in the LSP object for PCECC LSP.

The message formats in this document are specified using Routing Backus-Naur Form (RBNF) encoding as specified in [\[RFC5511\]](#).

6.1. The PCInitiate Message

The PCInitiate message [[RFC8281](#)] can be used to download or remove the labels, this document extends the message as shown below -

```
<PCInitiate Message> ::= <Common Header>
                           <PCE-initiated-lsp-list>
```

Where:

<Common Header> is defined in [[RFC5440](#)]

```
<PCE-initiated-lsp-list> ::= <PCE-initiated-lsp-request>
                             [<PCE-initiated-lsp-list>]
```

```
<PCE-initiated-lsp-request> ::=
    (<PCE-initiated-lsp-instantiation>|
     <PCE-initiated-lsp-deletion>|
     <PCE-initiated-lsp-central-control>)
```

```
<PCE-initiated-lsp-central-control> ::= <SRP>
                                         <LSP>
                                         <cci-list>
```

```
<cci-list> ::= <CCI>
               [<cci-list>]
```

Where:

<PCE-initiated-lsp-instantiation> and
<PCE-initiated-lsp-deletion> are as per
[\[RFC8281\]](#).

The LSP and SRP object is defined in [[RFC8231](#)].

When PCInitiate message is used for the central controller instructions (labels), the SRP, LSP, and CCI objects MUST be present. The SRP object is defined in [[RFC8231](#)] and if the SRP object is missing, the receiving PCC MUST send a PCErr message with Error-type=6 (Mandatory Object missing) and Error-value=10 (SRP object missing). The LSP object is defined in [[RFC8231](#)] and if the LSP object is missing, the receiving PCC MUST send a PCErr message with Error-type=6 (Mandatory Object missing) and Error-value=8 (LSP object missing). The CCI object is defined in [Section 7.3](#) and if the CCI object is missing, the receiving PCC MUST send a PCErr message with Error-type=6 (Mandatory Object missing) and Error-value=TBD11 (CCI object missing). More than one CCI object MAY be included in the PCInitiate message for a transit LSR.

To clean up entries, the R (remove) bit MUST be set in the SRP object to be encoded along with the LSP and the CCI object.

The CCI object received at the ingress node MUST have the O bit (out-label) set. The CCI Object received at the egress MUST have the O bit unset. If this is not the case, PCC MUST send a PCErr message with Error-Type = TBD5 ("PCECC failure") and Error Value = TBD9 ("Invalid CCI"). Other instances of the CCI object if present, MUST be ignored.

At most two instances of the CCI object can be included, in the case of transit LSR to encode both in-coming and out-going label forwarding instructions. Other instances MUST be ignored for P2P LSP. If the transit LSR did not receive two CCI object with one of them having the O bit set and another with O bit unset, it MUST send a PCErr message with Error-Type = TBD5 ("PCECC failure") and Error Value = TBD9 ("Invalid CCI").

Note that, on receipt of the PCInitiate message with CCI object, the ingress, egress, or transit role of the PCC is identified via the ingress and egress IP address encoded in the LSP-IDENTIFIERS TLV.

6.2. The PCRpt Message

The PCRpt message can be used to report the labels that were allocated by the PCE, to be used during the state synchronization phase or as acknowledgement to PCInitiate message.


```
<PCRpt Message> ::= <Common Header>
                        <state-report-list>
```

Where:

```
<state-report-list> ::= <state-report>[<state-report-list>]
```

```
<state-report> ::= (<lsp-state-report>|
                    <central-control-report>)
```

```
<lsp-state-report> ::= [<SRP>]
                        <LSP>
                        <path>
```

```
<central-control-report> ::= [<SRP>]
                              <LSP>
                              <cci-list>
```

```
<cci-list> ::= <CCI>
               [<cci-list>]
```

Where:

<path> is as per [\[RFC8231\]](#) and the LSP and SRP object are also defined in [\[RFC8231\]](#).

When PCRpt message is used to report the central controller instructions (labels), the LSP and CCI objects MUST be present. The LSP object is defined in [\[RFC8231\]](#) and if the LSP object is missing, the receiving PCE MUST send a PCErr message with Error-type=6 (Mandatory Object missing) and Error-value=8 (LSP object missing). The CCI object is defined in [Section 7.3](#) and if the CCI object is missing, the receiving PCE MUST send a PCErr message with Error-type=6 (Mandatory Object missing) and Error-value=TBD11 (CCI object missing). Two CCI objects can be included in the PCRpt message for a transit LSR.

[7. PCEP Objects](#)

The PCEP objects defined in this document are compliant with the PCEP object format defined in [\[RFC5440\]](#).

[7.1. OPEN Object](#)

This document defines new optional TLVs for use in the OPEN Object.

7.1.1. PCECC Capability sub-TLV

The PCECC-CAPABILITY sub-TLV is an optional TLV for use in the OPEN Object for PCECC capability advertisement in PATH-SETUP-TYPE-CAPABILITY TLV. Advertisement of the PCECC capability implies support of LSPs that are set up through PCECC as per PCEP extensions defined in this document.

Its format is shown in Figure 7.

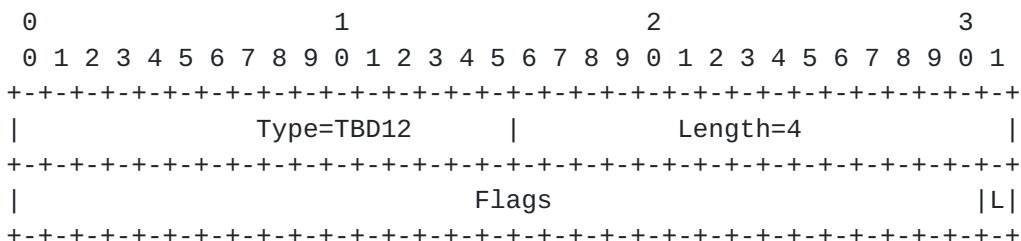


Figure 7: PCECC Capability sub-TLV

The type of the TLV is TBD12 and it has a fixed length of 4 octets.

The value comprises a single field - Flags (32 bits). Currently, the following flag bit is defined:

- o L bit (Label): if set to 1 by a PCEP speaker, the L flag indicates that the PCEP speaker allows PCECC based central controller instructions for label download. The bit MUST be set to 1 by both a PCC and a PCE for the PCECC label download/report on a PCEP session.
- o Unassigned bits MUST be set to 0 on transmission and MUST be ignored on receipt.

7.2. PATH-SETUP-TYPE TLV

The PATH-SETUP-TYPE TLV is defined in [RFC8408]; this document defines a new PST value:

- o PST = TBD1: Path is set up via PCECC mode.

On a PCRpt/PCUpd/PCInitiate message, the PST=TBD1 in PATH-SETUP-TYPE TLV in SRP object indicates that this LSP was set up via a PCECC-based mechanism.

7.3. CCI Object

The Central Controller Instructions (CCI) Object is used by the PCE to specify the forwarding instructions (Label information in the context of this document) to the PCC, and MAY be carried within PCInitiate or PCRpt message for label download/report.

CCI Object-Class is TBD13.

CCI Object-Type is 1 for the MPLS Label.

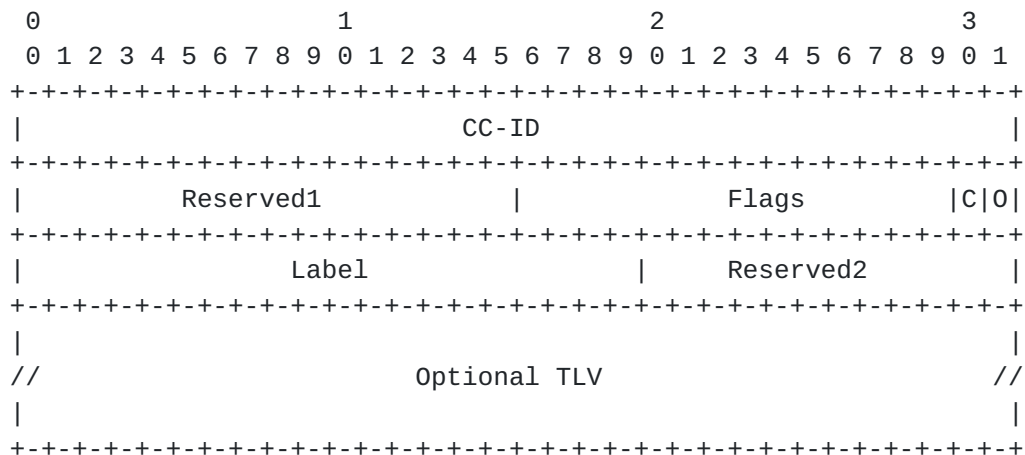


Figure 8: CCI Object

The fields in the CCI object are as follows:

CC-ID: A PCEP-specific identifier for the CCI information. A PCE creates a CC-ID for each instruction, the value is unique within the scope of the PCE and is constant for the lifetime of a PCEP session. The values 0 and 0xFFFFFFFF are reserved and MUST NOT be used.

Reserved1 (16 bit): Set to zero while sending, ignored on receive.

Flags (16 bit): A field used to carry any additional information pertaining to the CCI. Currently, the following flag bits are defined:

- * **0 bit(Out-label)** : If the bit is set to 1, it specifies the label is the OUT label and it is mandatory to encode the next-hop information (via IPV4-ADDRESS TLV or IPV6-ADDRESS TLV or UNNUMBERED-IPV4-ID-ADDRESS TLV in the CCI object). If the bit is not set, it specifies the label is the IN label and it is

optional to encode the local interface information (via IPV4-ADDRESS TLV or IPV6-ADDRESS TLV or UNNUMBERED-IPV4-ID-ADDRESS TLV in the CCI object).

- * C Bit (PCC Allocation): If the bit is set to 1, it indicates that the allocation needs to be done by the PCC for this central controller instruction. A PCE sets this bit to request the PCC to make an allocation from its label space. A PCC would set this bit to indicate that it has allocated the CC-ID and report it to the PCE.
- * All unassigned bits MUST be set to zero at transmission and ignored at receipt.

Label (20-bit): The Label information.

Reserved2 (12 bit): Set to zero while sending, ignored on receive.

[7.3.1.](#) Address TLVs

This document defines the following TLVs for the CCI object to associate the next-hop information in the case of an outgoing label and local interface information in the case of an incoming label.

IPV4-ADDRESS TLV:

```

0                               1                               2                               3
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|                               Type=TBD14                               | Length = 4
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|                               IPv4 address                               |
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+

```

IPV6-ADDRESS TLV:

```

0                               1                               2                               3
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|                               Type=TBD15                               | Length = 16
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|                               |
//                               IPv6 address (16 bytes)                               //
|                               |
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+

```

UNNUMBERED-IPV4-ID-ADDRESS TLV:


```

0          1          2          3
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|          Type=TBD16          | Length = 8          |
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|                               | Node-ID             |
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|                               | Interface ID         |
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+

```

LINKLOCAL-IPV6-ID-ADDRESS TLV:

```

0          1          2          3
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|          Type=TBD17          | Length = 40          |
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
//          Local IPv6 address (16 octets)          //
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|          Local Interface ID          |
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
//          Remote IPv6 address (16 octets)          //
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|          Remote Interface ID          |
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+

```

Figure 9: Address TLVs

The address TLVs are as follows:

IPV4-ADDRESS TLV: an IPv4 address.

IPV6-ADDRESS TLV: an IPv6 address.

UNNUMBERED-IPV4-ID-ADDRESS TLV: a Node ID / Interface ID tuple.

LINKLOCAL-IPV6-ID-ADDRESS TLV: a pair of (global IPv6 address,
interface ID) tuples.

8. Implementation Status

[Note to the RFC Editor - remove this section before publication, as well as remove the reference to [RFC 7942](#).]

This section records the status of known implementations of the protocol defined by this specification at the time of posting of this Internet-Draft, and is based on a proposal described in [[RFC7942](#)].

The description of implementations in this section is intended to assist the IETF in its decision processes in progressing drafts to RFCs. Please note that the listing of any individual implementation here does not imply endorsement by the IETF. Furthermore, no effort has been spent to verify the information presented here that was supplied by IETF contributors. This is not intended as, and must not be construed to be, a catalog of available implementations or their features. Readers are advised to note that other implementations may exist.

According to [RFC7942], "this will allow reviewers and working groups to assign due consideration to documents that have the benefit of running code, which may serve as evidence of valuable experimentation and feedback that have made the implemented protocols more mature. It is up to the individual working groups to use this information as they see fit".

8.1. Huawei's Proof of Concept based on ONOS

The PCE function was developed in the ONOS open source platform. This extension was implemented on a private version as a proof of concept for PCECC.

- o Organization: Huawei
- o Implementation: Huawei's PoC based on ONOS
- o Description: PCEP as a southbound plugin was added to ONOS. To support PCECC, an earlier version of this I-D was implemented. Refer <https://wiki.onosproject.org/display/ONOS/PCEP+Protocol>
- o Maturity Level: Prototype
- o Coverage: Partial
- o Contact: satishk@huawei.com

9. Security Considerations

The security considerations described in [RFC8231] and [RFC8281] apply to the extensions described in this document. Additional considerations related to a malicious PCE are introduced.

9.1. Malicious PCE

PCE has complete control over PCC to update the labels and can cause the LSP's to behave inappropriately and cause major impact to the network. As a general precaution, it is RECOMMENDED that this PCEP

extension be activated on authenticated and encrypted sessions across PCEs and PCCs belonging to the same administrative authority, using Transport Layer Security (TLS) [[RFC8253](#)], as per the recommendations and best current practices in [[RFC7525](#)].

[10.](#) Manageability Considerations

[10.1.](#) Control of Function and Policy

A PCE or PCC implementation SHOULD allow to configure to enable/disable PCECC capability as a global configuration.

[10.2.](#) Information and Data Models

[RFC7420] describes the PCEP MIB, this MIB can be extended to get the PCECC capability status.

The PCEP YANG module [[I-D.ietf-pce-pcep-yang](#)] could be extended to enable/disable PCECC capability.

[10.3.](#) Liveness Detection and Monitoring

Mechanisms defined in this document do not imply any new liveness detection and monitoring requirements in addition to those already listed in [[RFC5440](#)].

[10.4.](#) Verify Correct Operations

Mechanisms defined in this document do not imply any new operation verification requirements in addition to those already listed in [[RFC5440](#)] and [[RFC8231](#)].

[10.5.](#) Requirements On Other Protocols

PCEP extensions defined in this document do not put new requirements on other protocols.

[10.6.](#) Impact On Network Operations

PCEP extensions defined in this document do not put new requirements on network operations.

[11.](#) IANA Considerations

11.1. PCEP TLV Type Indicators

IANA is requested to allocate the following TLV Type Indicator values within the "PCEP TLV Type Indicators" sub-registry of the PCEP Numbers registry:

Value	Meaning	Reference
TBD14	IPV4-ADDRESS TLV	This document
TBD15	IPV6-ADDRESS TLV	This document
TBD16	UNNUMBERED-IPV4-ID-ADDRESS TLV	This document
TBD17	LINKLOCAL-IPV6-ID-ADDRESS TLV	This document

11.2. PATH-SETUP-TYPE-CAPABILITY Sub-TLV Type Indicators

[RFC8408] requested the creation of "PATH-SETUP-TYPE-CAPABILITY Sub-TLV Type Indicators" sub-registry. Further IANA is requested to allocate the following code-point:

Value	Meaning	Reference
TBD12	PCECC-CAPABILITY	This document

11.3. PCECC-CAPABILITY sub-TLV's Flag field

This document defines the PCECC-CAPABILITY sub-TLV and requests that IANA to create a new sub-registry to manage the value of the PCECC-CAPABILITY sub-TLV's 32-bits Flag field. New values are to be assigned by Standards Action [[RFC8126](#)]. Each bit should be tracked with the following qualities:

- o Bit number (counting from bit 0 as the most significant bit)
- o Capability description
- o Defining RFC

Currently, there is one allocation in this registry.

Bit	Name	Reference
31	Label	This document
0-30	Unassigned	This document

11.4. Path Setup Type Registry

[RFC8408] created a sub-registry within the "Path Computation Element Protocol (PCEP) Numbers" registry called "PCEP Path Setup Types". IANA is requested to allocate a new code point within this registry, as follows:

Value	Description	Reference
TBD1	Traffic engineering path is set up using PCECC mode	This document

11.5. PCEP Object

IANA is requested to allocate new code-point in the "PCEP Objects" sub-registry for the CCI object as follows:

Object-Class	Value	Name	Reference
TBD13		CCI Object-Type	This document
	0		Reserved
	1		MPLS Label

11.6. CCI Object Flag Field

IANA is requested to create a new sub-registry to manage the Flag field of the CCI object called "CCI Object 16-bits Flag Field". New values are to be assigned by Standards Action [[RFC8126](#)]. Each bit should be tracked with the following qualities:

- o Bit number (counting from bit 0 as the most significant bit)
- o Capability description
- o Defining RFC

Two bits to be defined for the CCI Object flag field in this document as follows:

Bit	Description	Reference
0-13	Unassigned	This document
14	C Bit - PCC allocation	This document
15	O Bit - Specifies label is out-label	This document

11.7. PCEP-Error Object

IANA is requested to allocate new error types and error values within the "PCEP-ERROR Object Error Types and Values" sub-registry of the PCEP Numbers registry for the following errors:

Error-Type	Meaning
-----	-----
6	Mandatory Object missing.

Error-value = TBD11 : CCI object missing

10	Reception of an invalid object.	
	Error-value = TBD2 :	Missing PCECC Capability sub-TLV
19	Invalid operation.	
	Error-value = TBD3 :	Attempted PCECC operations when PCECC capability was not advertised
	Error-value = TBD4 :	Stateful PCE capability was not advertised
	Error-value = TBD8 :	Unknown Label
TBD5	PCECC failure.	
	Error-value = TBD6 :	Label out of range.
	Error-value = TBD7 :	Instruction failed.
	Error-value = TBD9 :	Invalid CCI.
	Error-value = TBD10 :	Unable to allocate the specified CCI.

12. Acknowledgments

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