

PCE Working Group  
Internet-Draft  
Intended status: Informational  
Expires: September 6, 2018

D. Dhody  
Y. Lee  
Huawei Technologies  
D. Ceccarelli  
Ericsson  
J. Shin  
SK Telecom  
D. King  
Lancaster University  
O. Gonzalez de Dios  
Telefonica I+D  
March 5, 2018

**Hierarchical Stateful Path Computation Element (PCE).  
draft-ietf-pce-stateful-hpce-03**

**Abstract**

A Stateful Path Computation Element (PCE) maintains information on the current network state, including: computed Label Switched Path (LSPs), reserved resources within the network, and pending path computation requests. This information may then be considered when computing new traffic engineered LSPs, and for associated and dependent LSPs, received from Path Computation Clients (PCCs).

The Hierarchical Path Computation Element (H-PCE) architecture, provides an architecture to allow the optimum sequence of inter-connected domains to be selected, and network policy to be applied if applicable, via the use of a hierarchical relationship between PCEs.

Combining the capabilities of Stateful PCE and the Hierarchical PCE would be advantageous. This document describes general considerations and use cases for the deployment of Stateful PCE(s) using the Hierarchical PCE architecture.

**Status of This Memo**

This Internet-Draft is submitted in full conformance with the provisions of [BCP 78](#) and [BCP 79](#).

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at <https://datatracker.ietf.org/drafts/current/>.

Internet-Drafts are draft documents valid for a maximum of six months

and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

## Copyright Notice

Copyright (c) 2018 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to [BCP 78](#) and the IETF Trust's Legal Provisions Relating to IETF Documents (<https://trustee.ietf.org/license-info>) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Simplified BSD License text as described in Section 4.e of the Trust Legal Provisions and are provided without warranty as described in the Simplified BSD License.

## Table of Contents

<a href="#">1.</a>	<a href="#">Introduction</a>	<a href="#">3</a>
<a href="#">1.1.</a>	<a href="#">Requirements Language</a>	<a href="#">3</a>
<a href="#">2.</a>	<a href="#">Terminology</a>	<a href="#">3</a>
<a href="#">3.</a>	<a href="#">Hierarchical Stateful PCE</a>	<a href="#">4</a>
<a href="#">3.1.</a>	<a href="#">Passive Operations</a>	<a href="#">4</a>
<a href="#">3.2.</a>	<a href="#">Active Operations</a>	<a href="#">7</a>
<a href="#">3.3.</a>	<a href="#">PCE Initiation Operation</a>	<a href="#">8</a>
<a href="#">3.3.1.</a>	<a href="#">Per Domain Stitched LSP</a>	<a href="#">8</a>
<a href="#">4.</a>	<a href="#">Other Considerations</a>	<a href="#">10</a>
<a href="#">4.1.</a>	<a href="#">Applicability to Inter-Layer</a>	<a href="#">10</a>
<a href="#">4.2.</a>	<a href="#">Applicability to ACTN</a>	<a href="#">11</a>
<a href="#">5.</a>	<a href="#">Security Considerations</a>	<a href="#">12</a>
<a href="#">6.</a>	<a href="#">Manageability Considerations</a>	<a href="#">12</a>
<a href="#">6.1.</a>	<a href="#">Control of Function and Policy</a>	<a href="#">12</a>
<a href="#">6.2.</a>	<a href="#">Information and Data Models</a>	<a href="#">12</a>
<a href="#">6.3.</a>	<a href="#">Liveness Detection and Monitoring</a>	<a href="#">12</a>
<a href="#">6.4.</a>	<a href="#">Verify Correct Operations</a>	<a href="#">12</a>
<a href="#">6.5.</a>	<a href="#">Requirements On Other Protocols</a>	<a href="#">12</a>
<a href="#">6.6.</a>	<a href="#">Impact On Network Operations</a>	<a href="#">12</a>
<a href="#">7.</a>	<a href="#">IANA Considerations</a>	<a href="#">12</a>
<a href="#">8.</a>	<a href="#">Acknowledgments</a>	<a href="#">12</a>
<a href="#">9.</a>	<a href="#">References</a>	<a href="#">12</a>
<a href="#">9.1.</a>	<a href="#">Normative References</a>	<a href="#">12</a>
<a href="#">9.2.</a>	<a href="#">Informative References</a>	<a href="#">13</a>
<a href="#">Appendix A.</a>	<a href="#">Contributor Addresses</a>	<a href="#">14</a>
	<a href="#">Authors' Addresses</a>	<a href="#">14</a>



## 1. Introduction

The Path Computation Element communication Protocol (PCEP) provides mechanisms for Path Computation Elements (PCEs) to perform path computations in response to Path Computation Clients' (PCCs) requests.

A stateful PCE is capable of considering, for the purposes of path computation, not only the network state in terms of links and nodes (referred to as the Traffic Engineering Database or TED) but also the status of active services (previously computed paths, and currently reserved resources, stored in the Label Switched Paths Database (LSP-DB).

[RFC8051] describes general considerations for a stateful PCE deployment and examines its applicability and benefits, as well as its challenges and limitations through a number of use cases.

[RFC8231] describes a set of extensions to PCEP to provide stateful control. A stateful PCE has access to not only the information carried by the network's Interior Gateway Protocol (IGP), but also the set of active paths and their reserved resources for its computations. The additional state allows the PCE to compute constrained paths while considering individual LSPs and their interactions. [RFC8281] describes the setup, maintenance and teardown of PCE-initiated LSPs under the stateful PCE model.

[RFC8231] also describes the active stateful PCE. The active PCE functionality allows a PCE to reroute an existing LSP or make changes to the attributes of an existing LSP, or delegate control of specific LSPs to a new PCE.

The ability to compute shortest constrained TE LSPs in Multiprotocol Label Switching (MPLS) and Generalized MPLS (GMPLS) networks across multiple domains has been identified as a key motivation for PCE development. [RFC6805] describes a Hierarchical PCE (H-PCE) architecture which can be used for computing end-to-end paths for inter-domain MPLS Traffic Engineering (TE) and GMPLS Label Switched Paths (LSPs). Within the Hierarchical PCE (H-PCE) architecture [RFC6805], the Parent PCE (P-PCE) is used to compute a multi-domain path based on the domain connectivity information. A Child PCE (C-PCE) may be responsible for a single domain or multiple domains, it is used to compute the intra-domain path based on its domain topology information.

This document presents general considerations for stateful PCE(s) in hierarchical PCE architecture. In particular, the behavior changes and additions to the existing stateful PCE mechanisms (including PCE-



initiated LSP setup and active PCE usage) in the context of networks using the H-PCE architecture.

The initial section of the document focuses on end to end (E2E) inter-domain TE LSP. [Section 3.3.1](#) describe the operations for the Per Domain LSP that could be stitched.

### **[1.1.](#) Requirements Language**

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [BCP 14](#) [[RFC2119](#)] [[RFC8174](#)] when, and only when, they appear in all capitals, as shown here.

## **[2.](#) Terminology**

The terminology is as per [[RFC4655](#)], [[RFC5440](#)], [[RFC6805](#)], [[RFC8231](#)], and [[RFC8281](#)].

## **[3.](#) Hierarchical Stateful PCE**

As described in [[RFC6805](#)], in the hierarchical PCE architecture, a P-PCE maintains a domain topology map that contains the child domains (seen as vertices in the topology) and their interconnections (links in the topology). The P-PCE has no information about the content of the child domains. Each child domain has at least one PCE capable of computing paths across the domain. These PCEs are known as C-PCEs and have a direct relationship with the P-PCE. The P-PCE builds the domain topology map either via direct configuration (allowing network policy to also be applied) or from learned information received from each C-PCE.

[[RFC8231](#)] specifies new functions to support a stateful PCE. It also specifies that a function can be initiated either from a PCC towards a PCE (C-E) or from a PCE towards a PCC (E-C).

This document extends these functions to support H-PCE Architecture from a C-PCE towards a P-PCE (CE-PE) or from a P-PCE towards a C-PCE (PE-CE). All PCE types herein (i.e., PE or CE) are assumed to be 'stateful PCE'.

A number of interactions are expected in the Hierarchical Stateful PCE architecture, these include:

LSP State Report (CE-PE): a child stateful PCE sends an LSP state report to a Parent Stateful PCE whenever the state of a LSP



changes.

LSP State Synchronization (CE-PE): after the session between the Child and Parent stateful PCEs is initialized, the P-PCE must learn the state of C-PCE's TE LSPs.

LSP Control Delegation (CE-PE,PE-CE): a C-PCE grants to the P-PCE the right to update LSP attributes on one or more LSPs; the C-PCE may withdraw the delegation or the P-PCE may give up the delegation at any time.

LSP Update Request (PE-CE): a stateful P-PCE requests modification of attributes on a C-PCE's TE LSP.

PCE LSP Initiation Request (PE-CE): a stateful P-PCE requests C-PCE to initiate a TE LSP.

Note that this hierarchy is recursive and thus a Label Switching Router (LSR), as a PCC could delegate the control to a PCE, which may delegate to its parent, which may further delegate it to its parent (if it exist or needed). Similarly update operations could also be applied recursively.

[I-D.ietf-pce-hierarchy-extensions] defines the H-PCE capability TLV that should be used in the OPEN message to advertise the H-PCE capability. [RFC8231] defines the stateful PCE capability TLV. The presence of both TLVs represent the support for stateful H-PCE operations as described in this document.

[I-D.litkowski-pce-state-sync] describes the procedures to allow a stateful communication between PCEs for various use-cases. The procedures and extensions as described in Section 3 of [I-D.litkowski-pce-state-sync] are also applicable to Child and Parent PCE communication.

### **3.1. Passive Operations**

Procedures as described in [RFC6805] are applied, where the ingress C-PCE sends a request to the P-PCE. The P-PCE selects a set of candidate domain paths based on the domain topology and the state of the inter-domain links. It then sends computation requests to the C-PCEs responsible for each of the domains on the candidate domain paths. Each C-PCE computes a set of candidate path segments across its domain and sends the results to the P-PCE. The P-PCE uses this information to select path segments and concatenate them to derive the optimal end-to-end inter-domain path. The end-to-end path is then sent to the C-PCE that received the initial path request, and this C-PCE passes the path on to the PCC that issued the original



request.

As per [[RFC8231](#)], PCC sends an LSP State Report carried on a PCRpt message to the C-PCE, indicating the LSP's status. The C-PCE MAY further propagate the State Report to the P-PCE. A local policy at C-PCE MAY dictate which LSPs to be reported to the P-PCE. The PCRpt message is sent from C-PCE to P-PCE.

State synchronization mechanism as described in [[RFC8231](#)] and [[RFC8233](#)] are applicable to PCEP session between C-PCE and P-PCE as well.

Taking the sample hierarchical domain topology example from [[RFC6805](#)] as the reference topology for the entirety of this document.



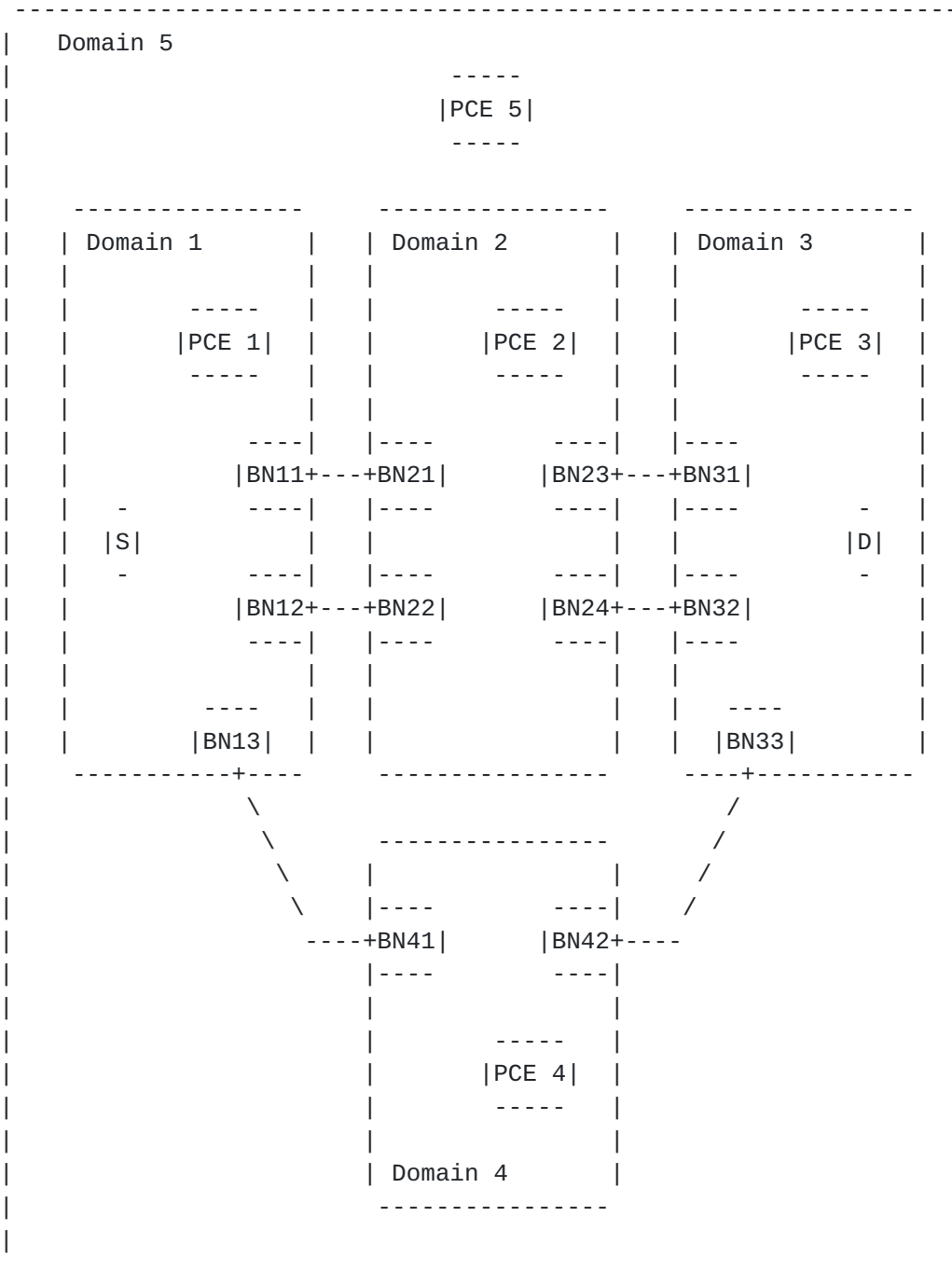


Figure 1: Sample Hierarchical Domain Topology

Steps 1 to 11 are exactly as described in [section 4.6.2](#) (Hierarchical PCE End-to-End Path Computation Procedure) of [\[RFC6805\]](#), the following additional steps are added for stateful PCE:

- (1) The Ingress LSR initiates the setup of the LSP as per the path and reports to the PCE1 the LSP status ("GOING-UP").



- (2) The PCE1 further reports the status of the LSP to the P-PCE (PCE5).
- (3) The Ingress LSR notifies the LSP state to PCE1 when the state is "UP".
- (4) The PCE1 further reports the status of the LSP to the P-PCE (PCE5).

### **3.2. Active Operations**

[RFC8231] describes the case of active stateful PCE. The active PCE functionality uses two specific PCEP messages:

- o Update Request (PCUpd)
- o State Report (PCRpt)

The first is sent by the PCE to a Path Computation Client (PCC) for modifying LSP attributes. The PCC sends back a PCRpt to acknowledge the requested operation or report any change in LSP's state.

As per [RFC8051], Delegation is an operation to grant a PCE, temporary rights to modify a subset of LSP parameters on one or more PCC's LSPs. The C-PCE may further choose to delegate to P-PCE based on a local policy. The PCRpt message with "D" (delegate) flag is sent from C-PCE to P-PCE.

To update an LSP, a PCE send to the PCC, an LSP Update Request using a PCUpd message. For LSP delegated to the P-PCE via the child PCE, the P-PCE can use the same PCUpd message to request change to the C-PCE (the Ingress domain PCE), the PCE further propagates the update request to the PCC.

The P-PCE uses the same mechanism described in [Section 3.1](#) to compute the end to end path using PCReq and PCRep messages.

The following additional steps are also initially performed, for active operations, again using the reference architecture described in Figure 1 (Sample Hierarchical Domain Topology).

- (1) The Ingress LSR delegates the LSP to the PCE1 via PCRpt message with D flag set.
- (2) The PCE1 further delegates the LSP to the P-PCE (PCE5).

Steps 4 to 10 of [section 4.6.2 of \[RFC6805\]](#) are executed to determine the end to end path.



- (3) The P-PCE (PCE5) sends the update request to the C-PCE (PCE1) via PCUpd message.
- (4) The PCE1 further updates the LSP to the Ingress LSR (PCC).
- (5) The Ingress LSR initiates the setup of the LSP as per the path and reports to the PCE1 the LSP status ("GOING-UP").
- (6) The PCE1 further reports the status of the LSP to the P-PCE (PCE5).
- (7) The Ingress LSR notifies the LSP state to PCE1 when the state is "UP".
- (8) The PCE1 further reports the status of the LSP to the P-PCE (PCE5).

### **3.3. PCE Initiation Operation**

[RFC8281] describes the setup, maintenance and teardown of PCE-initiated LSPs under the stateful PCE model, without the need for local configuration on the PCC, thus allowing for a dynamic network that is centrally controlled and deployed. To instantiate or delete an LSP, the PCE sends the Path Computation LSP Initiate Request (PCInitiate) message to the PCC. In case of inter-domain LSP in Hierarchical PCE architecture, the initiation operations can be carried out at the P-PCE. In which case after P-PCE finishes the E2E path computation, it can send the PCInitiate message to the C-PCE (the Ingress domain PCE), the PCE further propagates the initiate request to the PCC.

The following additional steps are also initially performed, for PCE initiated operations, again using the reference architecture described in Figure 1 (Sample Hierarchical Domain Topology):

- (1) The P-PCE (PCE5) is requested to initiate a LSP.

Steps 4 to 10 of [section 4.6.2 of \[RFC6805\]](#) are executed to determine the end to end path.

- (2) The P-PCE (PCE5) sends the initiate request to the child PCE (PCE1) via PCInitiate message.
- (3) The PCE1 further propagates the initiate message to the Ingress LSR (PCC).
- (4) The Ingress LSR initiates the setup of the LSP as per the path and reports to the PCE1 the LSP status ("GOING-UP").



- (5) The PCE1 further reports the status of the LSP to the P-PCE (PCE5).
- (6) The Ingress LSR notifies the LSP state to PCE1 when the state is "UP".
- (7) The PCE1 further reports the status of the LSP to the P-PCE (PCE5).

### **3.3.1. Per Domain Stitched LSP**

The Hierarchical PCE architecture as per [\[RFC6805\]](#) is primarily used for E2E LSP. With PCE-Initiated capability, another mode of operation is possible, where multiple intra-domain LSPs are initiated in each domain which are further stitched to form an E2E LSP. The P-PCE sends PCInitiate message to each C-PCE separately to initiate individual LSP segments along the domain path. These individual per domain LSP are stitched together by some mechanism, which is out of scope of this document (Refer [\[I-D.dugeon-pce-stateful-interdomain\]](#)).

The following additional steps are also initially performed, for the Per Domain stitched LSP operation, again using the reference architecture described in Figure 1 (Sample Hierarchical Domain Topology):

- (1) The P-PCE (PCE5) is requested to initiate a LSP.

Steps 4 to 10 of [section 4.6.2 of \[RFC6805\]](#) are executed to determine the end to end path, which are broken into per-domain LSPs say -

- o S-BN41
- o BN41-BN33
- o BN33-D

It should be noted that the P-PCE MAY use other mechanisms to determine the suitable per-domain LSPs (apart from [\[RFC6805\]](#)).

For LSP (BN33-D)

- (2) The P-PCE (PCE5) sends the initiate request to the child PCE (PCE3) via PCInitiate message for LSP (BN33-D).
- (3) The PCE3 further propagates the initiate message to BN33.
- (4) BN33 initiates the setup of the LSP as per the path and reports



to the PCE3 the LSP status ("GOING-UP").

- (5) The PCE3 further reports the status of the LSP to the P-PCE (PCE5).
- (6) The node BN33 notifies the LSP state to PCE3 when the state is "UP".
- (7) The PCE3 further reports the status of the LSP to the P-PCE (PCE5).

For LSP (BN41-BN33)

- (8) The P-PCE (PCE5) sends the initiate request to the child PCE (PCE4) via PCInitiate message for LSP (BN41-BN33).
- (9) The PCE4 further propagates the initiate message to BN41.
- (10) BN41 initiates the setup of the LSP as per the path and reports to the PCE4 the LSP status ("GOING-UP").
- (11) The PCE4 further reports the status of the LSP to the P-PCE (PCE5).
- (12) The node BN41 notifies the LSP state to PCE4 when the state is "UP".
- (13) The PCE4 further reports the status of the LSP to the P-PCE (PCE5).

For LSP (S-BN41)

- (14) The P-PCE (PCE5) sends the initiate request to the child PCE (PCE1) via PCInitiate message for LSP (S-BN41).
- (15) The PCE1 further propagates the initiate message to node S.
- (16) S initiates the setup of the LSP as per the path and reports to the PCE1 the LSP status ("GOING-UP").
- (17) The PCE1 further reports the status of the LSP to the P-PCE (PCE5).
- (18) The node S notifies the LSP state to PCE1 when the state is "UP".
- (19) The PCE1 further reports the status of the LSP to the P-PCE (PCE5).







path setup as well.

## **4.2. Applicability to ACTN**

[I-D.ietf-teas-actn-framework] describes framework for Abstraction and Control of TE Networks (ACTN), where each Provisioning Network Controller (PNC) is equivalent to C-PCE and P-PCE is the Multi-Domain Service Coordinator (MDSC). The Per domain stitched LSP as per the Hierarchical PCE architecture described in [Section 3.3.1](#) and [Section 4.1](#) is well suited for ACTN.

[I-D.ietf-pce-applicability-actn] examines the applicability of PCE to the ACTN framework. To support the function of multi domain coordination via hierarchy, the stateful hierarchy of PCEs plays a crucial role.

In ACTN framework, Customer Network Controller (CNC) can request the MDSC to check if there is a possibility to meet Virtual Network (VN) requirements (before requesting for VN provision). The H-PCE architecture as described in [\[RFC6805\]](#) can support via the use of PCReq and PCRep messages between the P-PCE and C-PCEs.

## **5. Other Considerations**

### **5.1. Scalability Considerations**

It should be noted that if all the C-PCEs would report all the LSPs in their domain, it could lead to scalability issues for the P-PCE. Thus it is recommended to only report the LSPs which are involved in H-PCE, i.e. the LSPs which are either delegated to the P-PCE or initiated by the P-PCE. Scalability considerations for PCEP as per [\[RFC8231\]](#) continue to apply for the PCEP session between child and parent PCE.

### **5.2. Confidentiality**

As described in [section 4.2 of \[RFC6805\]](#), information about the content of child domains is not shared for both scaling and confidentiality reasons. Along with the confidentiality during path computation, the child PCE could also conceal the path information, a C-PCE may replace a path segment with a path-key [\[RFC5520\]](#), effectively hiding the content of a segment of a path.

## **6. Security Considerations**

The security considerations listed in [\[RFC8231\]](#), [\[RFC6805\]](#) and [\[RFC5440\]](#) apply to this document as well. As per [\[RFC6805\]](#), it is



expected that the parent PCE will require all child PCEs to use full security when communicating with the parent.

Any multi-domain operation necessarily involves the exchange of information across domain boundaries. This is bound to represent a significant security and confidentiality risk especially when the child domains are controlled by different commercial concerns. PCEP allows individual PCEs to maintain confidentiality of their domain path information using path-keys [RFC5520], and the hierarchical PCE architecture is specifically designed to enable as much isolation of domain topology and capabilities information as is possible. The LSP state in the PCRpt message SHOULD continue to use this.

The security consideration for PCE-Initiated LSP as per [RFC8281] is also applicable from P-PCE to C-PCE.

Thus securing the PCEP session (between the P-PCE and the C-PCE) using mechanism like TCP Authentication Option (TCP-AO) [RFC5925] or Transport Layer Security (TLS) [RFC8253] is RECOMMENDED.

## **7. Manageability Considerations**

All manageability requirements and considerations listed in [RFC5440], [RFC6805], [RFC8231], and [RFC8281] apply to Stateful H-PCE defined in this document. In addition, requirements and considerations listed in this section apply.

### **7.1. Control of Function and Policy**

Support of the hierarchical procedure will be controlled by the management organization responsible for each child PCE. The parent PCE must only accept path computation requests from authorized child PCEs. If a parent PCE receives report from an unauthorized child PCE, the report should be dropped. All mechanism as described in [RFC8231] and [RFC8281] continue to apply.

### **7.2. Information and Data Models**

An implementation SHOULD allow the operator to view the stateful and H-PCE capabilities advertised by each peer. The PCEP YANG module [I-D.ietf-pce-pcep-yang] can be extended to include details stateful H-PCE.

### **7.3. Liveness Detection and Monitoring**

Mechanisms defined in this document do not imply any new liveness



detection and monitoring requirements in addition to those already listed in [[RFC5440](#)].

#### **7.4. Verify Correct Operations**

Mechanisms defined in this document do not imply any new operation verification requirements in addition to those already listed in [[RFC5440](#)] and [[RFC8231](#)].

#### **7.5. Requirements On Other Protocols**

Mechanisms defined in this document do not imply any new requirements on other protocols.

#### **7.6. Impact On Network Operations**

Mechanisms defined in [[RFC5440](#)] and [[RFC8231](#)] also apply to PCEP extensions defined in this document.

The stateful H-PCE technique brings the applicability of stateful PCE as described in [[RFC8051](#)], for the LSP traversing multiple domains.

### **8. IANA Considerations**

There are no IANA considerations.

### **9. Acknowledgments**

Thanks to Manuela Scarella, Haomian Zheng, Sergio Marmo, Stefano Parodi, Giacomo Agostini, Jeff Tantsura and Rajan Rao for suggestions.

### **10. References**

#### **10.1. Normative References**

[RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", [BCP 14](#), [RFC 2119](#), DOI 10.17487/RFC2119, March 1997, <<http://www.rfc-editor.org/info/rfc2119>>.

[RFC5440] Vasseur, JP., Ed. and JL. Le Roux, Ed., "Path Computation Element (PCE) Communication Protocol (PCEP)", [RFC 5440](#), DOI 10.17487/RFC5440, March 2009, <<http://www.rfc-editor.org/info/rfc5440>>.

[RFC6805] King, D., Ed. and A. Farrel, Ed., "The Application of the



Path Computation Element Architecture to the Determination of a Sequence of Domains in MPLS and GMPLS", [RFC 6805](#), DOI 10.17487/RFC6805, November 2012, <<http://www.rfc-editor.org/info/rfc6805>>.

[RFC8174] Leiba, B., "Ambiguity of Uppercase vs Lowercase in [RFC 2119](#) Key Words", [BCP 14](#), [RFC 8174](#), DOI 10.17487/RFC8174, May 2017, <<https://www.rfc-editor.org/info/rfc8174>>.

[RFC8231] Crabbe, E., Minei, I., Medved, J., and R. Varga, "Path Computation Element Communication Protocol (PCEP) Extensions for Stateful PCE", [RFC 8231](#), DOI 10.17487/RFC8231, September 2017, <<https://www.rfc-editor.org/info/rfc8231>>.

[RFC8281] Crabbe, E., Minei, I., Sivabalan, S., and R. Varga, "Path Computation Element Communication Protocol (PCEP) Extensions for PCE-Initiated LSP Setup in a Stateful PCE Model", [RFC 8281](#), DOI 10.17487/RFC8281, December 2017, <<https://www.rfc-editor.org/info/rfc8281>>.

## **[10.2](#). Informative References**

[RFC4655] Farrel, A., Vasseur, J., and J. Ash, "A Path Computation Element (PCE)-Based Architecture", [RFC 4655](#), DOI 10.17487/RFC4655, August 2006, <<http://www.rfc-editor.org/info/rfc4655>>.

[RFC5520] Bradford, R., Ed., Vasseur, JP., and A. Farrel, "Preserving Topology Confidentiality in Inter-Domain Path Computation Using a Path-Key-Based Mechanism", [RFC 5520](#), DOI 10.17487/RFC5520, April 2009, <<http://www.rfc-editor.org/info/rfc5520>>.

[RFC5623] Oki, E., Takeda, T., Le Roux, JL., and A. Farrel, "Framework for PCE-Based Inter-Layer MPLS and GMPLS Traffic Engineering", [RFC 5623](#), DOI 10.17487/RFC5623, September 2009, <<http://www.rfc-editor.org/info/rfc5623>>.

[RFC5925] Touch, J., Mankin, A., and R. Bonica, "The TCP Authentication Option", [RFC 5925](#), DOI 10.17487/RFC5925, June 2010, <<http://www.rfc-editor.org/info/rfc5925>>.

[RFC8051] Zhang, X., Ed. and I. Minei, Ed., "Applicability of a Stateful Path Computation Element (PCE)", [RFC 8051](#), DOI 10.17487/RFC8051, January 2017, <<http://www.rfc-editor.org/info/rfc8051>>.



- [RFC8233] Dhody, D., Wu, Q., Manral, V., Ali, Z., and K. Kumaki, "Extensions to the Path Computation Element Communication Protocol (PCEP) to Compute Service-Aware Label Switched Paths (LSPs)", [RFC 8233](#), DOI 10.17487/RFC8233, September 2017, <<https://www.rfc-editor.org/info/rfc8233>>.
- [RFC8253] Lopez, D., Gonzalez de Dios, O., Wu, Q., and D. Dhody, "PCEPS: Usage of TLS to Provide a Secure Transport for the Path Computation Element Communication Protocol (PCEP)", [RFC 8253](#), DOI 10.17487/RFC8253, October 2017, <<https://www.rfc-editor.org/info/rfc8253>>.
- [I-D.ietf-teas-actn-framework]  
Ceccarelli D. and Y. Lee, "Framework for Abstraction and Control of Transport Networks", [draft-ietf-teas-actn-framework-11](#) (work in progress), October 2017.
- [I-D.ietf-pce-applicability-actn]  
Dhody, D., Lee, Y., and D. Ceccarelli, "Applicability of Path Computation Element (PCE) for Abstraction and Control of TE Networks (ACTN)", [draft-ietf-pce-applicability-actn-03](#) (work in progress), March 2018.
- [I-D.litkowski-pce-state-sync]  
Litkowski, S., Sivabalan, S., and D. Dhody, "Inter Stateful Path Computation Element communication procedures", [draft-litkowski-pce-state-sync-02](#) (work in progress), August 2017.
- [I-D.ietf-pce-hierarchy-extensions]  
Zhang, F., Zhao, Q., Dios, O., Casellas, R., and D. King, "Extensions to Path Computation Element Communication Protocol (PCEP) for Hierarchical Path Computation Elements (PCE)", [draft-ietf-pce-hierarchy-extensions-03](#) (work in progress), July 2016.
- [I-D.ietf-pce-pcep-yang]  
Dhody, D., Hardwick, J., Beeram, V., and j. jeffrant@gmail.com, "A YANG Data Model for Path Computation Element Communications Protocol (PCEP)", [draft-ietf-pce-pcep-yang-06](#) (work in progress), January 2018.
- [I-D.dugeon-pce-stateful-interdomain]  
Dugeon, O. and J. Meuric, "PCEP Extension for Stateful Inter-Domain Tunnels", [draft-dugeon-pce-stateful-interdomain-00](#) (work in progress), October 2017.



**Appendix A. Contributor Addresses**

Avantika  
Huawei Technologies  
Divyashree Techno Park, Whitefield  
Bangalore, Karnataka 560066  
India

EMail: s.avantika.avantika@gmail.com

Xian Zhang  
Huawei Technologies  
Bantian, Longgang District  
Shenzhen, Guangdong 518129  
P.R.China

EMail: zhang.xian@huawei.com

Udayasree Palle  
Huawei Technologies  
Divyashree Techno Park, Whitefield  
Bangalore, Karnataka 560066  
India

EMail: udayasreereddy@gmail.com

**Authors' Addresses**

Dhruv Dhody  
Huawei Technologies  
Divyashree Techno Park, Whitefield  
Bangalore, Karnataka 560066  
India

EMail: dhruv.ietf@gmail.com

Young Lee  
Huawei Technologies  
5340 Legacy Drive, Building 3  
Plano, TX 75023  
USA

EMail: leeyoung@huawei.com

Daniele Ceccarelli  
Ericsson



Torshamnsgatan, 48  
Stockholm  
Sweden

EMail: daniele.ceccarelli@ericsson.com

Jongyoon Shin  
SK Telecom  
6 Hwangsaeul-ro, 258 beon-gil, Bundang-gu, Seongnam-si,  
Gyeonggi-do 463-784  
Republic of Korea

EMail: jongyoon.shin@sk.com

Daniel King  
Lancaster University  
UK

EMail: d.king@lancaster.ac.uk

Oscar Gonzalez de Dios  
Telefonica I+D  
Don Ramon de la Cruz 82-84  
Madrid, 28045  
Spain

Phone: +34913128832  
Email: ogondio@tid.es

