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PCEP Extensions for LSP scheduling with stateful PCE
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Abstract

This document proposes a set of extensions needed to the stateful Path Computation Element (PCE) communication Protocol (PCEP), so as to enable Labeled Switched Path (LSP) scheduling for path computation and LSP setup/deletion based on the actual network resource usage duration of a traffic service in a centralized network environment as stated in [I-D.ietf-teas-scheduled-resources].

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1. Introduction

The Path Computation Element Protocol (PCEP) defined in [RFC5440] is used between a Path Computation Element (PCE) and a Path Computation Client (PCC) (or other PCE) to enable path computation of Multi-protocol Label Switching (MPLS) Traffic Engineering Label Switched Path (TE LSP).

Further, in order to support use cases described in [RFC8231] specifies a set of extensions to PCEP to enable stateful control of MPLS-TE and GMPLS LSPs via PCEP.

Traditionally, the usage and allocation of network resources, especially bandwidth, can be supported by a Network Management System operation such as path pre-establishment. However, this does not provide efficient network usage since the established paths exclude the possibility of being used by other services even when they are not used for undertaking any service. [I-D.ietf-teas-scheduled-resources] then provides a framework that describes and discusses the problem and proposes an appropriate architecture for the scheduled reservation of TE resources.

With the scheduled reservation of TE resources, it allows network operators to reserve resources in advance according to the agreements with their customers, and allow them to transmit data with scheduling such as specified starting time and duration, for example for a scheduled bulk data replication between data centers. It enables the activation of bandwidth usage at the time the service really being used while letting other services obtain it in spare time. The requirement of scheduled LSP provision is mentioned in [RFC8231] and [RFC7399], so as to provide more efficient network resource usage for traffic engineering, which hasn't been solved yet. Also, for deterministic networks, the scheduled LSP can provide a better network resource usage for guaranteed links. This idea can also be applied in segment routing to schedule the network resources over the whole network in a centralized manner as well.

With this in mind, this document proposes a set of extensions needed to the stateful PCE, so as to enable LSP scheduling for path computation and LSP setup/deletion based on the actual network resource usage duration of a traffic service. A scheduled LSP is characterized by a starting time and a duration. When the end of the LSP life is reached, it is deleted to free up the resources for other LSP (scheduled or not).

2. Conventions used in this document

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

2.1. Terminology

The following terminologies are re-used from existing PCE documents.

- o Active Stateful PCE [RFC8231];
- o Passive Stateful PCE [RFC8231];
- o Delegation [RFC8231];
- o PCE-Initiated LSP [RFC8281];
- o PCC [RFC5440], [RFC8231];
- o PCE [RFC5440], [RFC8231];
- o TE LSP [RFC5440], [RFC8231];
- o TED [RFC5440], [RFC8231];
- o LSP DB [RFC8231];

In addition, this document defines the following terminologies.

Scheduled TE LSP: a LSP with the scheduling attributes,that carries traffic flow demand at a starting time and last for a certain duration. The PCE operates path computation per LSP availability for the required time and duration.

Scheduled LSP DB: a database of scheduled LSPs

Scheduled TED: Traffic engineering database with the awareness of scheduled resources for TE. This database is generated by the PCE from the information in TED and scheduled LSP DB and allows knowing, at any time, the amount of available resources (does not include failures in the future).

Starting time(start-time): This value indicates when the scheduled LSP is used and the corresponding LSP must be setup and active. In other time(i.e., before the starting time or after the starting

time plus Duration), the LSP can be inactive to include the possibility of the resources being used by other services.

Duration: The value indicates the time duration that the LSP is undertaken by a traffic flow and the corresponding LSP must be setup and active. At the end of which, the LSP is teardown and removed from the data base.

3. Motivation and Objectives

A stateful PCE can support better efficiency by using LSP scheduling described in the use case of [RFC8231]. This requires the PCE to maintain the scheduled LSPs and their associated resource usage, e.g. bandwidth for Packet-switched network, as well as have the ability to trigger signaling for the LSP setup/tear-down at the correct time.

Note that existing configuration tools can be used for LSP scheduling, but as highlighted in section 3.1.3 of [RFC8231] as well as discussions in [I-D.ietf-teas-scheduled-resources], doing this as a part of PCEP in a centralized manner, has obvious advantages.

The objective of this document is to provide a set of extensions to PCEP to enable LSP scheduling for LSPs creation/deletion under the stateful PCE control, according to traffic services from customers, so as to improve the usage of network resources.

4. Architecture Overview

4.1. LSP scheduling Overview

The LSP scheduling allows PCEs and PCCs to provide scheduled LSP for customers' traffic services at its actual usage time, so as to improve the network resource efficient utilization.

For stateful PCE supporting LSP scheduling, there are two types of LSP databases used in this document. One is the LSP-DB defined in PCEP [RFC8231], while the other is the scheduled LSP database (SLSP-DB, see section 6). The SLSP-DB records scheduled LSPs and is used as a complementary to the TED and LSP-DB. Note that the two types of LSP databases can be implemented in one physical database or two different databases. This document does not state any preference here.

Furthermore, a scheduled TED can be generated from the scheduled LSP DB, LSP DB and TED to indicate the network links and nodes with resource availability information for now and future. The scheduled TED should be maintained by all PCEs within the network environment.

In case of implementing PCC-initiated scheduled LSPs, before a PCC delegates a scheduled LSP, it MAY use the PCReq/PCRep messages to learn the path for the scheduled LSP. A PCC MUST delegate a scheduled LSP with information of its scheduling parameters, including the starting time and the duration using PCRpt message. Since the LSP is not yet signaled, at the time of delegation the LSP would be in down state. Upon receiving the delegation of the scheduled LSP, a stateful PCE SHALL check the scheduled TED for the network resource availability on network nodes and computes a path for the LSP with the scheduling information and update to the PCC as per the active stateful PCE techniques [RFC8231].

Note that the active stateful PCE can update to the PCC with the path for the scheduled LSP at any time. However, the PCC should not signal the LSP over the path once receiving these messages since the path is not activated yet until its starting time.

For a multiple PCE environment, in order to synchronize the scheduled LSP DB, the mechanism as described in [I-D.litkowski-pce-state-sync] are used to synchronize between PCEs. The scheduled TED is determined from the synchronized SLSP-DB. The PCE with delegation for the scheduled LSP would report the scheduled LSP to other PCEs, any future update to the scheduled LSP is also updated to other PCEs. This way the state of all scheduled LSPs are synchronized among the PCEs. [RFC7399] discusses some synchronization issues and considerations, that are also applicable to the scheduled databases.

The scheduled LSP can also be initiated by PCE itself. In case of implementing PCE-initiated scheduled LSP, the stateful PCE shall check the network resource availability for the traffic and computes a path for the scheduled LSP and initiate a scheduled LSP at the PCC and synchronize the scheduled LSP to other PCEs. Note that, the PCC could be notified immediately or at the starting time of the scheduled LSP based on the local policy. In case of former SCHED-LSP-ATTRIBUTE MUST be included in the message where as for the latter SCHED-LSP-ATTRIBUTE SHOULD NOT be included.

In both modes, for activation of scheduled LSPs, the PCC could initiate the setup of scheduled LSP at the start time by itself or wait for the PCE to update the PCC to initiate the setup of LSP. Similarly on the scheduling usage expires, the PCC could initiate the removal by itself or wait for the PCE to request the removal of the LSP. This is based on the Flag set in SCHED-LSP-ATTRIBUTE TLV. The state of the scheduled LSP is synchronized to other PCEs using the existing mechanism in [RFC8231] and [I-D.litkowski-pce-state-sync].

4.2. Support of LSP Scheduling

4.2.1. LSP Scheduling

For a scheduled LSP, a user configures it with an arbitrary scheduling duration time from Ta to time Tb, which may be represented as [Ta, Tb].

When an LSP is configured with arbitrary scheduling duration [Ta, Tb], a path satisfying the constraints for the LSP in the scheduling duration is computed and the LSP along the path is set up to carry traffic from time Ta to time Tb.

4.2.2. Periodical LSP Scheduling

In addition to LSP Scheduling at an arbitrary time period, there are also periodical LSP Scheduling.

A periodical LSP Scheduling represents Scheduling LSP every time interval. It has a scheduling duration such as [Ta, Tb], a number of repeats such as 10 (repeats 10 times), and a repeat cycle/time interval such as a week (repeats every week). The scheduling interval: "[Ta, Tb] repeats n times with repeat cycle C" represents n+1 scheduling intervals as follows:

[Ta, Tb], [Ta+C, Tb+C], [Ta+2C, Tb+2C], ..., [Ta+nC, Tb+nC]

When an LSP is configured with a scheduling interval such as "[Ta, Tb] repeats 10 times with a repeat cycle a week" (representing 11 scheduling intervals), a path satisfying the constraints for the LSP in each of the scheduling intervals represented by the periodical scheduling interval is computed and the LSP along the path is set up to carry traffic in each of the scheduling intervals.

4.2.2.1. Elastic Time LSP Scheduling

In addition to the basic LSP scheduling at an arbitrary time period, another option is elastic time intervals, which is represented as within -P and Q, where P and Q is an amount of time such as 300 seconds. P is called elastic range lower bound and Q is called elastic range upper bound.

For a simple time interval such as [Ta, Tb] with an elastic range, elastic time interval: "[Ta, Tb] within -P and Q" means a time period from (Ta+X) to (Tb+X), where -P <= X <= Q. Note that both Ta and Tb may be shifted the same X.

When an LSP is configured with elastic time interval "[Ta, Tb] within -P and Q", a path is computed such that the path satisfies the constraints for the LSP in the time period from (Ta+X) to (Tb+X) and |X| is the minimum value from 0 to max(P, Q). That is that [Ta+X, Tb+X] is the time interval closest to time interval [Ta, Tb] within the elastic range. The LSP along the path is set up to carry traffic in the time period from (Ta+X) to (Tb+X).

Similarly, for a recurrent time interval with an elastic range, elastic time interval: "[Ta, Tb] repeats n times with repeat cycle C within -P and Q" represents n+1 simple elastic time intervals as follows:

```
[Ta+X0, Tb+X0], [Ta+C+X1, Tb+C+X1], ..., [Ta+nC+Xn, Tb+nC+Xn]
      where -P <= Xi <= Q, i = 0, 1, 2, ..., n.
```

If a user wants to keep the same repeat cycle between any two adjacent time intervals, elastic time interval: "[Ta, Tb] repeats n times with repeat cycle C within -P and Q SYNC" may be used, which represents n+1 simple elastic time intervals as follows:

```
[Ta+X, Tb+X], [Ta+C+X, Tb+C+X], ..., [Ta+nC+X, Tb+nC+X]
     where -P <= X <= Q.
```

4.2.2.2. Graceful Periods

Besides the stated time scheduling, a user may want to have some graceful periods for each or some of the time intervals for the LSP. Two graceful periods may be configured for a time interval. One is the graceful period before the time interval, called grace-before, which extends the lifetime of the LSP for grace-before (such as 30 seconds) before the time interval. The other is the one after the time interval, called grace-after, which extends the lifetime of the LSP for grace-after (such as 60 seconds) after the time interval.

When an LSP is configured with a simple time interval such as [Ta, Tb] with graceful periods such as grace-before GB and grace-after GA, a path is computed such that the path satisfies the constraints for the LSP in the time period from Ta to Tb. The LSP along the path is set up to carry traffic in the time period from (Ta-GB) to (Tb+GA). During graceful periods from (Ta-GB) to Ta and from Tb to (Tb+GA), the LSP is up to carry traffic (maybe in best effort).

4.3. Scheduled LSP creation

In order to realize PCC-Initiated scheduled LSP in a centralized network environment, a PCC has to separate the setup of a LSP into two steps. The first step is to request/delegate and get a LSP but

not signal it over the network. The second step is to signal the scheduled LSP over the LSRs (Labeled switched Router) at its starting time.

For PCC-Initiated scheduled LSPs, a PCC can delegate the scheduled LSP by sending a path computation report (PCRpt) message by including its demanded resources with the scheduling information to a stateful PCE. Note the PCC MAY use the PCReq/PCRep with scheduling information before delegating.

Upon receiving the delegation via PCRpt message, the stateful PCE computes the path for the scheduled LSP per its starting time and duration based on the network resource availability stored in scheduled TED (see section 4.1).

The stateful PCE will send a PCUpd message with the scheduled path information as well as the scheduled resource information for the scheduled LSP to the PCC. The PCE SHOULD add the scheduled LSP into its scheduled LSP DB and update its scheduled TED. The PCE SHOULD also synchronize to other PCEs within the network if there is any, so as to keep their scheduling information synchronized as per [I-D.litkowski-pce-state-sync].

For PCE-Initiated Scheduled LSP, the stateful PCE can compute a path for the scheduled LSP per requests from network management systems automatically based on the network resource availability in the scheduled TED, send a PCInitiate message with the path information back to the PCC. Based on the local policy, the PCInitiate message could be sent immediately to ask PCC to create a scheduled LSP (as per this document) or the PCInitiate message could be sent at the start time to the PCC to create a normal LSP (as per [RFC8281]).

In both modes:

- The stateful PCE is required to update its local scheduled LSP DB and scheduled TED with the scheduled LSP. Besides, it shall send a PCRpt message with the scheduled LSP to other PCEs within the network, so as to achieve the scheduling traffic engineering information synchronization as per [I-D.litkowski-pce-state-sync].

- Upon receiving the PCUpd message or PCInitiate message for the scheduled LSP from PCEs with a found path, the PCC knows that it is a scheduled path for the LSP and does not trigger signaling for the LSP setup on LSRs.

- The stateful PCE can update the Scheduled LSP parameters on any network events using the PCUpd message to PCC. These changes are

also synchronized to other PCEs as per [I-D.litkowski-pce-state-sync].

- o Based on the configuration (and the C flag in scheduled TLVs), when it is time (i.e., at the start time) for the LSP to be set up, either the PCC triggers the LSP to be signaled or the delegated PCE sends a PCUpd message to the head end LSR providing the updated path to be signaled (with A flag set to indicate LSP activation).

4.4. Scheduled LSP Modifications

After a scheduled LSP is configured, a user may change its parameters including the requested time as well as the bandwidth.

In PCC-Initiated case, the PCC can send a PCRpt message for the scheduled LSP with updated parameters as well as scheduled information included in the SCHED-LSP-ATTRIBUTE TLV (see Section 5.2.1) or SCHED-PD-LSP-ATTRIBUTE TLV (see Section 5.2.2) carried in the LSP Object. The PCE would take the updated resources and schedule into considerations and update the new path for the scheduled LSP to the PCC as well as synchronize to other PCEs in the network. In case path cannot be set based on new requirements the same should be conveyed by the use of empty ERO in the PCEP messages.

In PCE-Initiated case, the Stateful PCE would recompute the path based on updated parameters as well as scheduled information. In case it has already conveyed to the PCC this information by sending a PCInitiate message, it should update the path and other scheduling and resource information by sending a PCUpd message.

In any case, the scheduled databases SHALL be updated and the PCE MUST synchronize this information to other PCEs as per [I-D.litkowski-pce-state-sync].

4.5. Scheduled LSP activation and deletion

In PCC-Initiated case, based on the configuration (and the C flag in scheduled TLVs), when it is time (i.e., at the start time) for the LSP to be set up, either the PCC triggers the LSP to be signaled or the delegated PCE sends a PCUpd message to the head end LSR providing the updated path to be signaled (with A flag set to indicate LSP activation). The PCC would report the status of the active LSP as per the procedures in [RFC8231] and at this time the LSP MUST be considered as part of the LSP-DB. The A flag MUST be set in the scheduled TLVs to indicate that the LSP is active now. After the scheduled duration expires, based on the C flag, the PCC triggers the

LSP deletion on it self or the delegated PCE sends a PCUpd message to the PCC to delete the LSP as per the procedures in [RFC8231].

In PCE-Initiated case, based on the local policy, if the scheduled LSP is already conveyed to the PCC at the time of creation, the handling of LSP activation and deletion is handled in the same way as PCC-Initiated case as per the setting of C flag. In other case, the PCE would send the PCInitiate message at the start time to the PCC to create a normal LSP without the scheduled TLVs and remove the LSP after the duration expires as per [RFC8281].

Additionally, the scheduled databases SHALL be updated and synchronization to other PCEs MUST be done as per [I-D.litkowski-pce-state-sync].

5. PCEP Objects and TLVs

5.1. Stateful PCE Capability TLV

After a TCP connection for a PCEP session has been established, a PCC and a PCE indicates its ability to support LSP scheduling during the PCEP session establishment phase. For a multiple-PCE environment, the PCEs should also establish PCEP session and indicate its ability to support LSP scheduling among PCEP peers. The Open Object in the Open message contains the STATEFUL-PCE-CAPABILITY TLV defined in [RFC8231]. Note that the STATEFUL- PCE-CAPABILITY TLV is defined in [RFC8231] and updated in [RFC8281] and [RFC8232]". In this document, we define a new flag bit B (SCHED-LSP-CAPABLITY) flag for the STATEFUL- PCE-CAPABILITY TLV to indicate the support of LSP scheduling and another flag bit PD (PD-LSP-CAPABLITY) to indicate the support of LSP periodical scheduling.

B (LSP-SCHEDULING-CAPABILITY - 1 bit) [Bit Position - TBD3]: If set to 1 by a PCC, the B Flag indicates that the PCC allows LSP scheduling; if set to 1 by a PCE, the B Flag indicates that the PCE is capable of LSP scheduling. The B bit MUST be set by both PCEP peers in order to support LSP scheduling for path computation.

PD (PD-LSP-CAPABLITY - 1 bit): [Bit Position - TBD4] If set to 1 by a PCC, the PD Flag indicates that the PCC allows LSP scheduling periodically; if set to 1 by a PCE, the PD Flag indicates that the PCE is capable of periodical LSP scheduling. The PD bit MUST be set by both PCEP peers in order to support periodical LSP scheduling for path computation.

5.2. LSP Object

The LSP object is defined in [RFC8231]. This document add an optional SCHED-LSP-ATTRIBUTE TLV for normal LSP scheduling and an optional SCHED-PD-LSP-ATTRIBUTE TLV for periodical LSP scheduling.

The presence of SCHED-LSP-ATTRIBUTE TLV in the LSP object indicates that this LSP is requesting scheduled parameters while the SCHED-PD-LSP-ATTRIBUTE TLV indicates that this scheduled LSP is periodical. The scheduled LSP attribute TLV MUST be present in LSP Object for each scheduled LSP carried in the PCEP messages. For periodical LSPs, the SCHED-PD-LSP-ATTRIBUTE TLV can be used in LSP Object for each periodic scheduled LSP carried in the PCEP messages.

5.2.1. SCHED-LSP-ATTRIBUTE TLV

The SCHED-LSP-ATTRIBUTE TLV can be included as an optional TLV within the LSP object for LSP scheduling for the requesting traffic service.

This TLV SHOULD be included only if both PCEP peers have set the B (LSP-SCHEDULING-CAPABILITY bit) in STATEFUL-PCE-CAPABILITY TLV carried in open message.

The format of the SCHED-LSP-ATTRIBUTE TLV is shown in the following figure:

```
 0                   1                   2                   3
 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|             Type (TBD1)       |          Length               |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|R|C|A|  Flags  |               Reserved (0)                    |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|                       Start-Time                              |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|                         Duration                              |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|                GrB            |             GrA               |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|      Elastic-Lower-Bound      |      Elastic-Upper-Bound      |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
```

The type of the TLV is [TBD1] and the TLV has a fixed length of 20 octets.

The fields in the format are:

Flags (8 bits): Flowing flags are defined in this document

R (1 bit): Set to 1 to indicate the Start-Time is a relative time, which is relative to the current time; set to 0 to indicate the the Start-Time is an absolute time.

C (1 bit): Set to 1 to indicate the PCC is responsible to setup and remove the scheduled LSP based on the Start-Time and duration.

A (1 bit): Set to 1 to indicate the scheduled LSP has been activated and should be considered as part of LSP-DB (instead of Scheduled LSP-DB).

Reserved (24 bits): This field MUST be set to zero on transmission and MUST be ignored on receipt.

Start-Time (32 bits): This value in seconds, indicates when the scheduled LSP is used to carry traffic and the corresponding LSP must be setup and activated.

Duration (32 bits): The value in seconds, indicates the duration that the LSP is undertaken by a traffic flow and the corresponding LSP must be up to carry traffic. At the expiry of this duration, the LSP is tear down and deleted.

The Start-Time indicates a calendar time (e.g.,2018/12/13 8:29:58), at or before which the scheduled LSP must be set up. The value of the Start-Time represents the number of seconds since 00:00 hours,Jan 1, 1970 UTC when R bit is set to 0. When R bit is set to 1, it represents the number of seconds from the current time.

In addition, it contains an non zero grace-before and grace-after if graceful periods are configured. It includes an non zero elastic range lower bound and upper bound if there is an elastic range configured.

- o GrB (Grace-Before -16 bits): The graceful period time length in seconds before the starting time.

- o GrA (Grace-After -16 bits): The graceful period time length in seconds after time interval [starting time, starting time + duration].

- o Elastic-Lower-Bound (16 bits): The maximum amount of time in seconds that time interval can shift to lower/left.

- o Elastic-Upper-Bound (16 bits): The maximum amount of time in seconds that time interval can shift to upper/right.

5.2.2. SCHED-PD-LSP-ATTRIBUTE TLV

The periodical LSP is a special case of LSP scheduling. The traffic service happens in a series of repeated time intervals. The SCHED-PD-LSP-ATTRIBUTE TLV can be included as an optional TLV within the LSP object for this periodical LSP scheduling.

This TLV SHOULD be included only if both PCEP peers have set the B (LSP-SCHEDULING-CAPABILITY bit) and PD (PD-LSP-CAPABLITY bit) in STATEFUL-PCE-CAPABILITY TLV carried in open message.

The format of the SCHED-PD-LSP-ATTRIBUTE TLV is shown in the following figure:

```
 0                   1                   2                   3
 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|            Type (TBD2)        |         Length                |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|R|C|A|  Flags  | Opt   |  NR   |        Reserved (0)           |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|                           Start-Time                          |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|                              Duration                         |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|                      Repeat-time-length                       |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|                GrB            |             GrA               |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
|      Elastic-Lower-Bound      |      Elastic-Upper-Bound      |
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+
```

The type of the TLV is [TBD2] and the TLV has a fixed length of 24 octets. The description, format and meaning of the Flags (R, C and A bit), Start-Time, Duration, GrB, GrA, Elastic-Lower-Bound and Elastic-Upper-Bound fields remains same as SCHED-LSP-ATTRIBUTE TLV.

The following fields are new :

Opt: (4 bits) Indicates options to repeat.

Options = 1: repeat every day;

Options = 2: repeat every week;

Options = 3: repeat every month;

Options = 4: repeat every year;

Options = 5: repeat every Repeat-time-length.

NR:(4 bits) The number of repeats. In each of repeats, LSP carries traffic. If value is set to 0xFFFF, it indicates forever.

Reserved (16 bits): This field MUST be set to zero on transmission and MUST be ignored on receipt.

Repeat-time-length:(32 bits) The time length in seconds after which LSP starts to carry traffic again for the Duration.

6. The PCEP Messages

6.1. The PCRpt Message

Path Computation State Report (PCRpt) is a PCEP message sent by a PCC to a PCE to report the status of one or more LSPs as per [RFC8231]. Each LSP State Report in a PCRpt message MAY contain the actual LSP's path, bandwidth, operational and administrative status, etc. An LSP Status Report carried on a PCRpt message is also used in delegation or revocation of control of an LSP to/from a PCE. In case of scheduled LSP, the scheduled TLVs MUST be carried in the LSP object and the ERO conveys the intended path for the scheduled LSP. The scheduled LSP MUST be delegated to a PCE. This message is also used to synchronize the scheduled LSPs to other PCE as described in [I-D.litkowski-pce-state-sync].

6.2. The PCUpd Message

Path Computation Update Request (PCUpd) is a PCEP message sent by a PCE to a PCC to update LSP parameters, on one or more LSPs as per [RFC8231]. Each LSP Update Request on a PCUpd message MUST contain all LSP parameters that a PCE wishes to be set for a given LSP. In case of scheduled LSP, the scheduled TLVs MUST be carried in the LSP object and the ERO conveys the intended path for the scheduled LSP. In case no path can be found, an empty ERO is used. The A bit is used in PCUpd message to indicate the activation of the scheduled LSP in case the PCE is responsible for the activation (as per the C bit). This message is also used to synchronize the scheduled LSPs to other PCE as described in [I-D.litkowski-pce-state-sync].

6.3. The PCInitiate Message

An LSP Initiate Request (PCInitiate) message is a PCEP message sent by a PCE to a PCC to trigger LSP instantiation or deletion as per [RFC8281]. In case of scheduled LSP, based on the local policy, PCE MAY convey the scheduled LSP to the PCC by including the scheduled TLVs in the LSP object. Or the PCE would initiate the LSP only at the start time of the scheduled LSP as per the [RFC8281] without the use of scheduled TLVs.

6.4. The PCReq message

The Path Computation Request (PCReq) message is a PCEP message sent by a PCC to a PCE to request a path computation [RFC5440] and it MAY contain the LSP object to identify the LSP for which the path computation is requested. In case of scheduled LSP, the scheduled TLVs MUST be carried in the LSP object in PCReq message to request the path computation based on scheduled TED and LSP-DB. A PCC MAY use PCReq message to obtain the scheduled path before delegating the LSP.

6.5. The PCRep Message

The Path Computation Reply (PCRep) message is a PCEP message sent by a PCE to a PCC in reply to a path computation request [RFC5440]. A PCRep message can contain either a set of computed paths if the request can be satisfied, or a negative reply if not. The negative reply may indicate the reason why no path could be found. In case of scheduled LSP, the scheduled TLVs MUST be carried in the LSP object in PCRep message to indicate the path computation based on scheduled TED and LSP-DB. A PCC and PCE MAY use PCReq and PCRep message to obtain the scheduled path before delegating the LSP.

6.6. The PCErr Message

[Editor's Note - Error Handling will be taken up in the next revision of the draft]

7. Security Considerations

This document defines LSP-SCHEDULING-CAPABILITY TLV and SCHED- LSP-ATTRIBUTE TLV which does not add any new security concerns beyond those discussed in [RFC5440] and [RFC8231]. But in some deployments the scheduling information could provide details about the network operations that could be deemed as extra sensitive. Additionally, snooping of PCEP messages with such data or using PCEP messages for network reconnaissance may give an attacker sensitive information about the operations of the network. Thus, such deployment should

employ suitable PCEP security mechanisms like TCP Authentication Option (TCP-AO) [RFC5925] or [RFC8253]. The procedure based on Transport Layer Security (TLS) in [RFC8253] is considered a security enhancement and thus is much better suited for the sensitive service-aware information.

8. Manageability Consideration

8.1. Control of Function and Policy

The LSP-Scheduling feature MUST BE controlled per tunnel by the active stateful PCE, the values for parameters like starting time, duration SHOULD BE configurable by customer applications and based on the local policy at PCE.

8.2. Information and Data Models

An implementation SHOULD allow the operator to view the capability defined in this document. To serve this purpose, the PCEP YANG module [I-D.ietf-pce-pcep-yang] could be extended.

8.3. Liveness Detection and Monitoring

Mechanisms defined in this document do not imply any new liveness detection and monitoring requirements in addition to those already listed in [RFC5440].

8.4. Verify Correct Operations

Mechanisms defined in this document do not imply any new operation verification requirements in addition to those already listed in [RFC5440].

8.5. Requirements On Other Protocols

Mechanisms defined in this document do not imply any new requirements on other protocols.

8.6. Impact On Network Operations

Mechanisms defined in this document do not have any impact on network operations in addition to those already listed in [RFC5440].

9. IANA Considerations

9.1. PCEP TLV Type Indicators

This document defines the following new PCEP TLV. IANA maintains a sub-registry "PCEP TLV Type Indicators" in the "Path Computation Element Protocol (PCEP) Numbers" registry. IANA is requested to make the following allocations from this sub-registry.

Value	Meaning	Reference
TBD1	SCHED-LSP-ATTRIBUTE	This document
TBD2	SCHED-PD-LSP-ATTRIBUTE	This document

9.2. STATEFUL-PCE-CAPABILITY TLV Flag field

This document defines new bits in the Flags field in the STATEFUL-PCE-CAPABILITY TLV in the OPEN object. IANA maintains a sub-registry "STATEFUL-PCE-CAPABILITY TLV Flag Field" in the "Path Computation Element Protocol (PCEP) Numbers" registry. IANA is requested to make the following allocations from this sub-registry.

The following values are defined in this document:

Bit	Description	Reference
TBD3	LSP-SCHEDULING-CAPABILITY (B-bit)	This document
TBD4	PD-LSP-CAPABLITY (PD-bit)	This document

9.3. Schedule TLVs Flag Field

IANA is requested to create a new sub-registry, named "Schedule TLVs Flag Field", within the "Path Computation Element Protocol (PCEP) Numbers" registry to manage the Flag field in the SCHED-LSP-ATTRIBUTE and SCHED-PD-LSP-ATTRIBUTE TLVs. New values are assigned by Standards Action [RFC8126]. Each bit should be tracked with the following qualities:

- o Bit number (counting from bit 0 as the most significant bit)

- o Capability description

- o Defining RFC

The following values are defined in this document:

Bit	Description	Reference
0	R-bit	This document
1	C-bit	This document
2	A-bit	This document

10. Acknowledgments

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Appendix A. Scheduled LSP information synchronization

As for a stateful PCE, it maintains a database of LSPs (LSP-DB) that are active in the network, so as to reveal the available network resources and place new LSPs more cleverly.

With the scheduled LSPs, they are not activated while creation, but should be considered when operating future path computation. Hence, a scheduled LSP Database (SLSP-DB) is suggested to maintain all scheduled LSP information.

The information of SLSP-DB MUST be shared and synchronized among all PCEs within the centralized network by using PCRpt and PCUpd message

with scheduled LSP information as per the mechanism described in [I-D.litkowski-pce-state-sync].

The PCE should generate and maintain a scheduled TED based on LSP DB, scheduled LSP DB and TED, which is used to indicate the network resource availability on network nodes for LSP path computation.

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