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# Multicast-only Fast Reroute Based on Topology Independent Loop-free Alternate Fast Reroute draft-ietf-pim-mofrr-tilfa-01

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#### Abstract

Multicast-only Fast Reroute (MoFRR) has been defined in [<u>RFC7431</u>], but the selection of the secondary multicast next hop only according to the loop-free alternate fast reroute, which has restrictions in multicast deployments. This document describes a mechanism for Multicast-only Fast Reroute by using Topology Independent Loop-free Alternate fast reroute, which is independent of network topology and can achieve covering more network environments.

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## **1**. Introduction

As the deployment of video services, operators are paying more and more attention to solutions that minimize the service disruption due to faults in the IP network carrying the packets for these services. Multicast-only Fast Reroute (MoFRR) has been defined in [<u>RFC7431</u>], which can minimize multicast packet loss in a network when node or link failures occur by making simple enhancements to multicast routing protocols such as Protocol Independent Multicast (PIM). But

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the selection of the secondary multicast next hop only according to the loop-free alternate fast reroute in [RFC7431], and there are limitations in multicast deployments for this mechanism. This document describes a new mechanism for Multicast-only Fast Reroute using Topology Independent Loop-free Alternate (TILFA) fast reroute, which is independent of network topology and can achieve covering more network environments.

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in <u>BCP 14</u> [<u>RFC2119</u>] [<u>RFC8174</u>] when, and only when, they appear in all capitals, as shown here.

### 1.2. Terminology

This document use the terms defined in  $[\underline{RFC7431}]$ , and also use the concepts defined in  $[\underline{RFC7490}]$ . The specific content of each term is not described in this document.

## **2**. Problem Statement

#### 2.1. LFA for MoFRR

In [RFC7431] section 3, the secondary Upstream Multicast Hop (UMH) of PIM for MoFRR is a loop-free alternate (LFA). However, the traditional LFA mechanism needs to satisfy at least one neighbor whose next hop to the destination node is an acyclic next hop, existing limitations in network deployments, and can only cover part of the network topology environments. In some network topology, the corresponding secondary UMH cannot be calculated, so PIM cannot establish a standby multicast tree and cannot implement MoFRR protection. Therefore, the current MoFRR of PIM is only available in the network topology applicable to LFA.

The problem of current MoFRR applicability can be illustrated by the example network shown in Figure 1. The metric of R1-R4 link is 20, the metric of R5-R6 link is 100, and the metrics of other links are 10.

Towards multicast source S1, the primary path of the PIM join packet is R3->R2->R1, and the secondary path is R3->R4->R1, which is the same as the LFA path of unicast routing. In this case, the current MoFRR can work well.

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Towards multicast source S2, the primary path of the PIM join packet is R3->R2. However, the LFA does not exist. If R3 sends the packet to R4, R4 will forward it back to R3, as the IGP shortest path from R4 to R1 is R4->R3->R2. In this case, the secondary UMH cannot be calculated by the current MoFRR. Similarly for multicast source S3, the current MoFRR does not work either.

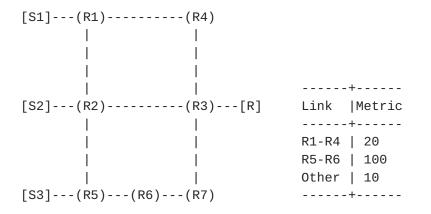


Figure 1: Example Network Topology

### 2.2. RLFA for MoFRR

The remote loop-free alternate (RLFA) defined in [RFC7490] is extended from the LFA and can cover more network deployment scenarios through the tunnel as an alternate path. The RLFA mechanism needs to satisfy at least one node assumed to be N in the network that the fault node is neither on the path from the source node to the N node, nor on the path from the N node to the destination node.

[RFC5496] defines the RPF Vector attribute that can be carried in the PIM Join packet such that the path is selected based on the unicast reachability of the RPF Vector. The secondary multicast tree of MoFRR can be established by using the combination of RLFA mechanism and RPF Vector, which would cover more network topologies than the current MoFRR with LFA.

For example, in the network of Figure 1, the secondary path of PIM join packet towards multicast source S2 cannot be calculated by the current MoFRR, as mentioned above. Based on the RLFA mechanism, R3 sends the packet to R4 along with an RPF Vector containing the IP address of R1, which is the PQ node of R3 with respect to the protected link R2-R3. Then R4 will forward the packet to R1 through the link R1-R4, according to the unicast route for the RPF Vector. R1 continues to forward the packet to R2, and the secondary path, R3->R4->R1->R2, is established by MoFRR with RLFA.

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## 2.3. TILFA for MoFRR

However, RLFA is also topology dependent. In the network of Figure 1, towards multicast source S3, the primary path of the PIM join packet is R3->R2->R5, but the RLFA path does not exist. This is because the PQ node of R3 with respect to the protected link R2-R3 does not exist. If R3 sends the packet to R7 along with an RPF Vector containing the IP address of R5, R7 will forward it back to R3, since the IGP shortest path from R7 to R5 is R7->R3->R2->R5. Or, if R3 sends the packet to R7 along with an RPF Vector containing the IP address of R6, but then R6 will forward it back to R7 address of R6, R7 will forward it to R6, but then R6 will forward it back to R7->R3->R2->R5. In this case, the secondary UMH cannot be calculated by MoFRR with RLFA.

RLFA only has enhancement compared to LFA but still has limitations in network deployments. [I-D.ietf-rtgwg-segment-routing-ti-lfa] defined a unicast FRR solution based on the TILFA mechanism. The TILFA mechanism can express the backup path with an explicit path, and has no constraint on the topology, providing a more reliable FRR mechanism. The unicast traffic can be forwarded according to the explicit path list as an alternate path to implement unicast traffic protection, and can achieve full coverage of various networking environments.

The alternate path provided by the TILFA mechanism is actually a Segment List, including one or more Adjacency SIDs of one or more links between the P space and the Q space, and the NodeSID of P space node. PIM can look up the corresponding node IP address in the unicast route according to the NodeSID, and the IP addresses of the two endpoints of the corresponding link in the unicast route according to the Adjacency SIDs, but the multicast protocol packets cannot be directly sent along the path of the Segment List.

PIM join message need to be sent hop-by-hop to establish a standby multicast tree. However, not all of the nodes and links on the unicast alternate path are included in the Segment List. If the PIM protocol packets are transmitted only in unicast mode, then equivalently the PIM packets are transmitted through the unicast tunnel like unicast traffic, and cannot pass through the intermediate nodes of the tunnel. The intermediate nodes of the alternate path cannot forward multicast traffic because there are no PIM state entries on the nodes. PIM needs to create entries on the device hop-by-hop and generate an incoming interface and an outgoing interface list. So it can form an end-to-end complete multicast tree for forwarding multicast traffic according to the Segment List path and can only establish a standby multicast tree.

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## 3. Solution

A secondary Upstream Multicast Hop (UMH) is a candidate next-hop that can be used to reach the root of the tree. In this document the secondary UMH is based on unicast routing to find the Segment List calculated by TILFA.

It is available in principle that the path information of the Segment List is added to the PIM packets to guide the hop-by-hop RPF selection. The IP address of the node corresponding to the Node SID can be used as the segmented root node, and the IP addresses of the interfaces at both endpoints of the link corresponding to the Adjacency SID can be used directly as the local upstream interface and upstream neighbor.

For the PIM protocol, the PIM RPF Vector attribute was defined in [<u>RFC5496</u>], which can carry the node IP address corresponding to the Node SID. The explicit RPF Vector was defined in [<u>RFC7891</u>], which can carry the peer IP address corresponding to the Adjacency SID.

This document can use the above two RPF Vector standards and does not need to extend the PIM protocol, to establish the standby multicast tree according to the Segment List calculated by TILFA, and can achieve full coverage of various networking environments for MoFRR protection of multicast services.

Assume that the Segment List calculated by TILFA is (NodeSID(A), AdjSID(A-B)). Node A belongs to the P Space, and node B belongs to the Q space. The IP address corresponding to NodeSID(A) can be looked up in the local link state database of the IGP protocol, and can be assumed to be IP-a. The IP addresses of the two endpoints of the link corresponding to AdjSID(A-B) can also be looked up in the local link state database of the IGP protocol, and can be assumed to be IP-La and IP-Lb.

In the procedure of PIM, IP-a can be looked as the normal RPF vector attribute and added to the PIM join packet. IP-La can be looked as the local address of the incoming interface, and IP-Lb can be looked as the address of the upstream neighbor. So IP-Lb can be added to the PIM join packet as the explicit RPF Vector attribute.

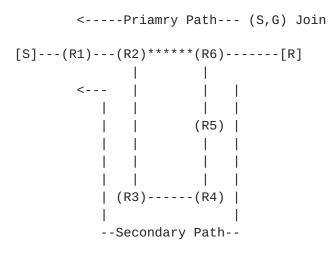
The PIM protocol firstly can select the RPF incoming interface and upstream towards IP-a, and can join hop-by-hop to establish the PIM standby multicast tree until the node A. On the node A, IP-Lb can be looked as the PIM upstream neighbor. The node A can find the incoming interface in the unicast routing table according to the IP-Lb, and IP-Lb is used as the RPF upstream address of the PIM join packet to the node B.

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After the PIM join packet is received on the node B, the PIM protocol can find no more RPF Vector attributes and select the RPF incoming interface and upstream towards the multicast source directly, and then can continue to join hop-by-hop to establish the PIM standby multicast tree until the router directly connected the source.

## **4**. Illustration

This section provides an illustration of MoFRR based on TI-LFA. The example topology is shown in Figure 2. The metric of R3-R4 link is 100, and the metrics of other links are 10. All the link metrics are bidirectional.





The IP addresses and SIDs, which may be involved in the MoFRR calculation, are configured as following:

IPv4 Data Plane (MPLS-SR)

| Node   | IP Address                               | Node SID |
|--------|--|----------|
| R4     | 4.4.4.4/32                               | 16004    |
| R3->R4 | IP Address<br>14.0.0.1/24<br>14.0.0.2/24 |          |

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IPv6 Data Plane (SRv6)

Node IP Address Node SID (End) R4 4::4/128 A4::100 Link IP Address Adjacency SID (End.X) R3->R4 BB::1/96 A3::1 R4->R3 BB::2/96 A4::1

The primary path of the PIM join packet is R6->R2->R1, and the secondary path is R6->R5->R4->R3->R2->R1. However, the traditional LFA does not work properly for the secondary path, because the shortest path to R2 from R5 (or from R4) still goes through R6-R2 link. In this case, R6 needs to calculate the secondary UMH using the proposed MoFRR method based on TI-LFA.

According to the TI-LFA algorithm, P-Space and Q-Space are shown in Figure 3. The TI-LFA repair path consists of the Node SID of R4 and the Adjacency SID of R4->R3. On MPLS-SR data plane, the repair segment list is (16004, 24002). On SRv6 data plane, the repair segment list is (A4::100, A4::1).

. . . . . . . .

Figure 3: P-Space and Q-Space

In the procedure of PIM, the IP addresses associated with the repair segment list are looked up in the IGP link state database.

On IPv4 data plane, the Node SID 16004 corresponds to 4.4.4.4, which will be carried in the RPF Vector Attribute. The Adjacency SID 24002 corresponds to local address 14.0.0.2 and remote peer address 14.0.0.1, and 14.0.0.1 will be carried in the Explicit RPF Vector Attribute.

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On IPv6 data plane, the End SID A4::100 corresponds to 4::4, which will be carried in the RPF Vector Attribute. The End.X SID A4::1 corresponds to local address BB::2 and remote peer address BB::1, and BB::1 will be carried in the Explicit RPF Vector Attribute.

Then, R6 installs the secondary UMH with these RPF Vectors.

```
+----+

|Type = 0|

|4.4.4.4 |

+----+

|Type = 4| |Type = 4|

|14.0.0.1| |14.0.0.1|

+----+ No RPF Vector
```

R6---->R5---->R4---->R3---->R2--->R1

Figure 4: Forwarding PIM Join Packet along Secondary Path (IPv4)

On IPv4 data plane, the forwarding of PIM join packet along the secondary path is shown in Figure 4.

R6 inserts two RPF Vector Attributes in the PIM join packet, which are 4.4.4.4 of Type 0 (RPF Vector Attribute) and 14.0.0.1 of Type 4 (Explicit RPF Vector Attribute). Then R6 forwards the packet along the secondary path.

When R5 receives the packet, R5 performs a unicast route lookup of the first RPF Vector 4.4.4.4 and sends the packet to R4.

R4 is the owner of 4.4.4.4, so it removes the first RPF Vector from the packet and forwards according to the following RPF Vector. R4 sends the packet to R3 according to the next RPF Vector 14.0.0.1, since its PIM neighbor R3 corresponds to 14.0.0.1.

When R3 receives the packet, as the owner of 14.0.0.1, it removes the RPF Vector. Then the packet has no RPF Vector, and will be forwarded to the source through R3->R2->R1 according to unicast routes.

After the PIM join packet reaches R1, a secondary multicast tree, R1->R2->R3->R4->R5->R6, is established hop-by-hop for (S, G) by MoFRR based on TI-LFA.

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+----+ |Type = 0| |4::4 | +----+ +---+ |Type = 4| |Type = 4| |BB::1 | |BB::1 | +---+ No RPF Vector

R6---->R5---->R4---->R3---->R2--->R1

Figure 5: Forwarding PIM Join Packet along Secondary Path (IPv6)

On IPv6 data plane, the forwarding of PIM join packet along the secondary path is shown in Figure 5. The procedure is similar with IPv4 data plane.

## **<u>5</u>**. Security Considerations

This document does not change the security properties of PIM. For general PIM-SM protocol Security Considerations, see [<u>RFC7761</u>].

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TBD

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