

**Hypertext Transfer Protocol (HTTP) over QUIC**  
**draft-ietf-quic-http-15**

Abstract

The QUIC transport protocol has several features that are desirable in a transport for HTTP, such as stream multiplexing, per-stream flow control, and low-latency connection establishment. This document describes a mapping of HTTP semantics over QUIC. This document also identifies HTTP/2 features that are subsumed by QUIC, and describes how HTTP/2 extensions can be ported to QUIC.

Note to Readers

Discussion of this draft takes place on the QUIC working group mailing list ([quic@ietf.org](mailto:quic@ietf.org)), which is archived at [https://mailarchive.ietf.org/arch/search/?email\\_list=quic](https://mailarchive.ietf.org/arch/search/?email_list=quic) [1].

Working Group information can be found at <https://github.com/quicwg> [2]; source code and issues list for this draft can be found at <https://github.com/quicwg/base-drafts/labels/-http> [3].

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of [BCP 78](#) and [BCP 79](#).

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at <https://datatracker.ietf.org/drafts/current/>.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

This Internet-Draft will expire on April 6, 2019.

## Copyright Notice

Copyright (c) 2018 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to [BCP 78](#) and the IETF Trust's Legal Provisions Relating to IETF Documents (<https://trustee.ietf.org/license-info>) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Simplified BSD License text as described in Section 4.e of the Trust Legal Provisions and are provided without warranty as described in the Simplified BSD License.

## Table of Contents

<a href="#">1.</a>	<a href="#">Introduction</a>	<a href="#">4</a>
<a href="#">1.1.</a>	<a href="#">Notational Conventions</a>	<a href="#">4</a>
<a href="#">2.</a>	<a href="#">Connection Setup and Management</a>	<a href="#">4</a>
<a href="#">2.1.</a>	<a href="#">Draft Version Identification</a>	<a href="#">4</a>
<a href="#">2.2.</a>	<a href="#">Discovering an HTTP/QUIC Endpoint</a>	<a href="#">5</a>
<a href="#">2.2.1.</a>	<a href="#">QUIC Version Hints</a>	<a href="#">5</a>
<a href="#">2.3.</a>	<a href="#">Connection Establishment</a>	<a href="#">6</a>
<a href="#">2.4.</a>	<a href="#">Connection Reuse</a>	<a href="#">7</a>
<a href="#">3.</a>	<a href="#">Stream Mapping and Usage</a>	<a href="#">7</a>
<a href="#">3.1.</a>	<a href="#">HTTP Message Exchanges</a>	<a href="#">8</a>
<a href="#">3.1.1.</a>	<a href="#">Header Formatting and Compression</a>	<a href="#">9</a>
<a href="#">3.1.2.</a>	<a href="#">The CONNECT Method</a>	<a href="#">10</a>
<a href="#">3.1.3.</a>	<a href="#">Request Cancellation</a>	<a href="#">11</a>
<a href="#">3.2.</a>	<a href="#">Request Prioritization</a>	<a href="#">11</a>
<a href="#">3.2.1.</a>	<a href="#">Placeholders</a>	<a href="#">12</a>
<a href="#">3.2.2.</a>	<a href="#">Priority Tree Maintenance</a>	<a href="#">12</a>
<a href="#">3.3.</a>	<a href="#">Unidirectional Streams</a>	<a href="#">13</a>
<a href="#">3.3.1.</a>	<a href="#">Reserved Stream Types</a>	<a href="#">14</a>
<a href="#">3.3.2.</a>	<a href="#">Control Streams</a>	<a href="#">14</a>
<a href="#">3.3.3.</a>	<a href="#">Server Push</a>	<a href="#">15</a>
<a href="#">4.</a>	<a href="#">HTTP Framing Layer</a>	<a href="#">16</a>
<a href="#">4.1.</a>	<a href="#">Frame Layout</a>	<a href="#">16</a>
<a href="#">4.2.</a>	<a href="#">Frame Definitions</a>	<a href="#">17</a>
<a href="#">4.2.1.</a>	<a href="#">Reserved Frame Types</a>	<a href="#">17</a>
<a href="#">4.2.2.</a>	<a href="#">DATA</a>	<a href="#">17</a>
<a href="#">4.2.3.</a>	<a href="#">HEADERS</a>	<a href="#">17</a>
<a href="#">4.2.4.</a>	<a href="#">PRIORITY</a>	<a href="#">18</a>
<a href="#">4.2.5.</a>	<a href="#">CANCEL_PUSH</a>	<a href="#">20</a>
<a href="#">4.2.6.</a>	<a href="#">SETTINGS</a>	<a href="#">21</a>
<a href="#">4.2.7.</a>	<a href="#">PUSH_PROMISE</a>	<a href="#">23</a>
<a href="#">4.2.8.</a>	<a href="#">GOAWAY</a>	<a href="#">24</a>

Bishop

Expires April 6, 2019

[Page 2]

4.2.9.	MAX_PUSH_ID . . . . .	25
5.	Connection Closure . . . . .	25
5.1.	Idle Connections . . . . .	26
5.2.	Connection Shutdown . . . . .	26
5.3.	Immediate Application Closure . . . . .	27
5.4.	Transport Closure . . . . .	28
6.	Error Handling . . . . .	28
6.1.	HTTP/QUIC Error Codes . . . . .	28
7.	Extensions to HTTP/QUIC . . . . .	30
8.	Considerations for Transitioning from HTTP/2 . . . . .	30
8.1.	Streams . . . . .	31
8.2.	HTTP Frame Types . . . . .	31
8.3.	HTTP/2 SETTINGS Parameters . . . . .	33
8.4.	HTTP/2 Error Codes . . . . .	34
9.	Security Considerations . . . . .	35
10.	IANA Considerations . . . . .	35
10.1.	Registration of HTTP/QUIC Identification String . . . . .	35
10.2.	Registration of QUIC Version Hint Alt-Svc Parameter . . . . .	36
10.3.	Frame Types . . . . .	36
10.4.	Settings Parameters . . . . .	37
10.5.	Error Codes . . . . .	38
10.6.	Stream Types . . . . .	41
11.	References . . . . .	42
11.1.	Normative References . . . . .	42
11.2.	Informative References . . . . .	43
11.3.	URIs . . . . .	43
Appendix A.	Change Log . . . . .	43
A.1.	Since <a href="#">draft-ietf-quic-http-14</a> . . . . .	43
A.2.	Since <a href="#">draft-ietf-quic-http-13</a> . . . . .	44
A.3.	Since <a href="#">draft-ietf-quic-http-12</a> . . . . .	44
A.4.	Since <a href="#">draft-ietf-quic-http-11</a> . . . . .	44
A.5.	Since <a href="#">draft-ietf-quic-http-10</a> . . . . .	44
A.6.	Since <a href="#">draft-ietf-quic-http-09</a> . . . . .	45
A.7.	Since <a href="#">draft-ietf-quic-http-08</a> . . . . .	45
A.8.	Since <a href="#">draft-ietf-quic-http-07</a> . . . . .	45
A.9.	Since <a href="#">draft-ietf-quic-http-06</a> . . . . .	45
A.10.	Since <a href="#">draft-ietf-quic-http-05</a> . . . . .	45
A.11.	Since <a href="#">draft-ietf-quic-http-04</a> . . . . .	45
A.12.	Since <a href="#">draft-ietf-quic-http-03</a> . . . . .	46
A.13.	Since <a href="#">draft-ietf-quic-http-02</a> . . . . .	46
A.14.	Since <a href="#">draft-ietf-quic-http-01</a> . . . . .	46
A.15.	Since <a href="#">draft-ietf-quic-http-00</a> . . . . .	46
A.16.	Since <a href="#">draft-shade-quic-http2-mapping-00</a> . . . . .	47
Acknowledgements	. . . . .	47
Author's Address	. . . . .	47

Bishop

Expires April 6, 2019

[Page 3]

## **1. Introduction**

The QUIC transport protocol has several features that are desirable in a transport for HTTP, such as stream multiplexing, per-stream flow control, and low-latency connection establishment. This document describes a mapping of HTTP semantics over QUIC, drawing heavily on the existing TCP mapping, HTTP/2. Specifically, this document identifies HTTP/2 features that are subsumed by QUIC, and describes how the other features can be implemented atop QUIC.

QUIC is described in [[QUIC-TRANSPORT](#)]. For a full description of HTTP/2, see [[RFC7540](#)].

### **1.1. Notational Conventions**

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [BCP 14](#) [[RFC2119](#)] [[RFC8174](#)] when, and only when, they appear in all capitals, as shown here.

Field definitions are given in Augmented Backus-Naur Form (ABNF), as defined in [[RFC5234](#)].

This document uses the variable-length integer encoding from [[QUIC-TRANSPORT](#)].

Protocol elements called "frames" exist in both this document and [[QUIC-TRANSPORT](#)]. Where frames from [[QUIC-TRANSPORT](#)] are referenced, the frame name will be prefaced with "QUIC." For example, "QUIC APPLICATION\_CLOSE frames." References without this preface refer to frames defined in [Section 4.2](#).

## **2. Connection Setup and Management**

### **2.1. Draft Version Identification**

\*RFC Editor's Note:\* Please remove this section prior to publication of a final version of this document.

HTTP/QUIC uses the token "hq" to identify itself in ALPN and Alt-Svc. Only implementations of the final, published RFC can identify themselves as "hq". Until such an RFC exists, implementations MUST NOT identify themselves using this string.

Implementations of draft versions of the protocol MUST add the string "-" and the corresponding draft number to the identifier. For

Bishop

Expires April 6, 2019

[Page 4]

example, [draft-ietf-quic-http-01](#) is identified using the string "hq-01".

Non-compatible experiments that are based on these draft versions MUST append the string "-" and an experiment name to the identifier. For example, an experimental implementation based on [draft-ietf-quic-http-09](#) which reserves an extra stream for unsolicited transmission of 1980s pop music might identify itself as "hq-09-rickroll". Note that any label MUST conform to the "token" syntax defined in [Section 3.2.6 of \[RFC7230\]](#). Experimenters are encouraged to coordinate their experiments on the quic@ietf.org mailing list.

## **[2.2. Discovering an HTTP/QUIC Endpoint](#)**

An HTTP origin advertises the availability of an equivalent HTTP/QUIC endpoint via the Alt-Svc HTTP response header field or the HTTP/2 ALTSVC frame ([[ALTSVC](#)]), using the ALPN token defined in [Section 2.3](#).

For example, an origin could indicate in an HTTP/1.1 or HTTP/2 response that HTTP/QUIC was available on UDP port 50781 at the same hostname by including the following header field in any response:

```
Alt-Svc: hq=":50781"
```

On receipt of an Alt-Svc record indicating HTTP/QUIC support, a client MAY attempt to establish a QUIC connection to the indicated host and port and, if successful, send HTTP requests using the mapping described in this document.

Connectivity problems (e.g. firewall blocking UDP) can result in QUIC connection establishment failure, in which case the client SHOULD continue using the existing connection or try another alternative endpoint offered by the origin.

Servers MAY serve HTTP/QUIC on any UDP port, since an alternative always includes an explicit port.

### **[2.2.1. QUIC Version Hints](#)**

This document defines the "quic" parameter for Alt-Svc, which MAY be used to provide version-negotiation hints to HTTP/QUIC clients. QUIC versions are four-octet sequences with no additional constraints on format. Leading zeros SHOULD be omitted for brevity.

Syntax:

```
quic = DQUOTE version-number [ "," version-number ] * DQUOTE  
version-number = 1*8HEXDIG; hex-encoded QUIC version
```





Where multiple versions are listed, the order of the values reflects the server's preference (with the first value being the most preferred version). Reserved versions MAY be listed, but unreserved versions which are not supported by the alternative SHOULD NOT be present in the list. Origins MAY omit supported versions for any reason.

Clients MUST ignore any included versions which they do not support. The "quic" parameter MUST NOT occur more than once; clients SHOULD process only the first occurrence.

For example, suppose a server supported both version 0x00000001 and the version rendered in ASCII as "Q034". If it opted to include the reserved versions (from Section 4 of [[QUIC-TRANSPORT](#)]) 0x0 and 0x1abadaba, it could specify the following header field:

```
Alt-Svc: hq=":49288";quic="1,1abadaba,51303334,0"
```

A client acting on this header field would drop the reserved versions (because it does not support them), then attempt to connect to the alternative using the first version in the list which it does support.

### **[2.3.](#) Connection Establishment**

HTTP/QUIC relies on QUIC as the underlying transport. The QUIC version being used MUST use TLS version 1.3 or greater as its handshake protocol. HTTP/QUIC clients MUST indicate the target domain name during the TLS handshake. This may be done using the Server Name Indication (SNI) [[RFC6066](#)] extension to TLS or using some other mechanism.

QUIC connections are established as described in [[QUIC-TRANSPORT](#)]. During connection establishment, HTTP/QUIC support is indicated by selecting the ALPN token "hq" in the TLS handshake. Support for other application-layer protocols MAY be offered in the same handshake.

While connection-level options pertaining to the core QUIC protocol are set in the initial crypto handshake, HTTP/QUIC-specific settings are conveyed in the SETTINGS frame. After the QUIC connection is established, a SETTINGS frame ([Section 4.2.6](#)) MUST be sent by each endpoint as the initial frame of their respective HTTP control stream (see [Section 3.3.2](#)). The server MUST NOT send data on any other stream until the client's SETTINGS frame has been received.



#### **2.4. Connection Reuse**

Once a connection exists to a server endpoint, this connection MAY be reused for requests with multiple different URI authority components. The client MAY send any requests for which the client considers the server authoritative.

An authoritative HTTP/QUIC endpoint is typically discovered because the client has received an Alt-Svc record from the request's origin which nominates the endpoint as a valid HTTP Alternative Service for that origin. As required by [\[RFC7838\]](#), clients MUST check that the nominated server can present a valid certificate for the origin before considering it authoritative. Clients MUST NOT assume that an HTTP/QUIC endpoint is authoritative for other origins without an explicit signal.

A server that does not wish clients to reuse connections for a particular origin can indicate that it is not authoritative for a request by sending a 421 (Misdirected Request) status code in response to the request (see [Section 9.1.2 of \[RFC7540\]](#)).

The considerations discussed in [Section 9.1 of \[RFC7540\]](#) also apply to the management of HTTP/QUIC connections.

### **3. Stream Mapping and Usage**

A QUIC stream provides reliable in-order delivery of bytes, but makes no guarantees about order of delivery with regard to bytes on other streams. On the wire, data is framed into QUIC STREAM frames, but this framing is invisible to the HTTP framing layer. A QUIC receiver buffers and orders received STREAM frames, exposing the data contained within as a reliable byte stream to the application.

When HTTP headers and data are sent over QUIC, the QUIC layer handles most of the stream management.

All client-initiated bidirectional streams are used for HTTP requests and responses. A bidirectional stream ensures that the response can be readily correlated with the request. This means that the client's first request occurs on QUIC stream 0, with subsequent requests on stream 4, 8, and so on. In order to permit these streams to open, an HTTP/QUIC client SHOULD send non-zero values for the QUIC transport parameters "initial\_max\_stream\_data\_bidi\_local". An HTTP/QUIC server SHOULD send non-zero values for the QUIC transport parameters "initial\_max\_stream\_data\_bidi\_remote" and "initial\_max\_bidi\_streams". It is recommended that "initial\_max\_bidi\_streams" be no smaller than 100, so as to not unnecessarily limit parallelism.



These streams carry frames related to the request/response (see [Section 4.2](#)). When a stream terminates cleanly, if the last frame on the stream was truncated, this MUST be treated as a connection error (see HTTP\_MALFORMED\_FRAME in [Section 6.1](#)). Streams which terminate abruptly may be reset at any point in the frame.

HTTP/QUIC does not use server-initiated bidirectional streams. The use of unidirectional streams is discussed in [Section 3.3](#). Both clients and servers SHOULD send a value of three or greater for the QUIC transport parameter "initial\_max\_uni\_streams".

HTTP does not need to do any separate multiplexing when using QUIC - data sent over a QUIC stream always maps to a particular HTTP transaction. Requests and responses are considered complete when the corresponding QUIC stream is closed in the appropriate direction.

### **3.1. HTTP Message Exchanges**

A client sends an HTTP request on a client-initiated bidirectional QUIC stream. A server sends an HTTP response on the same stream as the request.

An HTTP message (request or response) consists of:

1. one header block (see [Section 4.2.3](#)) containing the message header (see [\[RFC7230\]](#), [Section 3.2](#)),
2. the payload body (see [\[RFC7230\]](#), [Section 3.3](#)), sent as a series of DATA frames (see [Section 4.2.2](#)),
3. optionally, one header block containing the trailer-part, if present (see [\[RFC7230\]](#), [Section 4.1.2](#)).

In addition, prior to sending the message header block indicated above, a response may contain zero or more header blocks containing the message headers of informational (1xx) HTTP responses (see [\[RFC7230\]](#), [Section 3.2](#) and [\[RFC7231\]](#), [Section 6.2](#)).

A server MAY interleave one or more PUSH\_PROMISE frames (see [Section 4.2.7](#)) with the frames of a response message. These PUSH\_PROMISE frames are not part of the response; see [Section 3.3.3](#) for more details.

The "chunked" transfer encoding defined in [Section 4.1 of \[RFC7230\]](#) MUST NOT be used.

Trailing header fields are carried in an additional header block following the body. Senders MUST send only one header block in the



trailers section; receivers MUST discard any subsequent header blocks.

An HTTP request/response exchange fully consumes a bidirectional QUIC stream. After sending a request, a client closes the stream for sending; after sending a response, the server closes the stream for sending and the QUIC stream is fully closed.

A server can send a complete response prior to the client sending an entire request if the response does not depend on any portion of the request that has not been sent and received. When this is true, a server MAY request that the client abort transmission of a request without error by triggering a QUIC STOP\_SENDING with error code HTTP\_EARLY\_RESPONSE, sending a complete response, and cleanly closing its streams. Clients MUST NOT discard complete responses as a result of having their request terminated abruptly, though clients can always discard responses at their discretion for other reasons.

Changes to the state of a request stream, including receiving a RST\_STREAM with any error code, do not affect the state of the server's response. Servers do not abort a response in progress solely due to a state change on the request stream. However, if the request stream terminates without containing a usable HTTP request, the server SHOULD abort its response with the error code HTTP\_INCOMPLETE\_REQUEST.

### **3.1.1. Header Formatting and Compression**

HTTP header fields carry information as a series of key-value pairs. For a listing of registered HTTP header fields, see the "Message Header Field" registry maintained at <https://www.iana.org/assignments/message-headers> [4].

Just as in previous versions of HTTP, header field names are strings of ASCII characters that are compared in a case-insensitive fashion. Properties of HTTP header field names and values are discussed in more detail in [Section 3.2 of \[RFC7230\]](#), though the wire rendering in HTTP/QUIC differs. As in HTTP/2, header field names MUST be converted to lowercase prior to their encoding. A request or response containing uppercase header field names MUST be treated as malformed.

As in HTTP/2, HTTP/QUIC uses special pseudo-header fields beginning with ':' character (ASCII 0x3a) to convey the target URI, the method of the request, and the status code for the response. These pseudo-header fields are defined in [Section 8.1.2.3](#) and [8.1.2.4 of \[RFC7540\]](#). Pseudo-header fields are not HTTP header fields. Endpoints MUST NOT generate pseudo-header fields other than those





defined in [\[RFC7540\]](#). The restrictions on the use of pseudo-header fields in [Section 8.1.2.1 of \[RFC7540\]](#) also apply to HTTP/QUIC.

HTTP/QUIC uses QPACK header compression as described in [\[QPACK\]](#), a variation of HPACK which allows the flexibility to avoid header-compression-induced head-of-line blocking. See that document for additional details.

### **[3.1.2.](#) The CONNECT Method**

The pseudo-method CONNECT ([\[RFC7231\]](#), [Section 4.3.6](#)) is primarily used with HTTP proxies to establish a TLS session with an origin server for the purposes of interacting with "https" resources. In HTTP/1.x, CONNECT is used to convert an entire HTTP connection into a tunnel to a remote host. In HTTP/2, the CONNECT method is used to establish a tunnel over a single HTTP/2 stream to a remote host for similar purposes.

A CONNECT request in HTTP/QUIC functions in the same manner as in HTTP/2. The request MUST be formatted as described in [\[RFC7540\]](#), [Section 8.3](#). A CONNECT request that does not conform to these restrictions is malformed. The request stream MUST NOT be half-closed at the end of the request.

A proxy that supports CONNECT establishes a TCP connection ([\[RFC0793\]](#)) to the server identified in the ":authority" pseudo-header field. Once this connection is successfully established, the proxy sends a HEADERS frame containing a 2xx series status code to the client, as defined in [\[RFC7231\]](#), [Section 4.3.6](#).

All DATA frames on the request stream correspond to data sent on the TCP connection. Any DATA frame sent by the client is transmitted by the proxy to the TCP server; data received from the TCP server is packaged into DATA frames by the proxy. Note that the size and number of TCP segments is not guaranteed to map predictably to the size and number of HTTP DATA or QUIC STREAM frames.

The TCP connection can be closed by either peer. When the client ends the request stream (that is, the receive stream at the proxy enters the "Data Recvd" state), the proxy will set the FIN bit on its connection to the TCP server. When the proxy receives a packet with the FIN bit set, it will terminate the send stream that it sends to client. TCP connections which remain half-closed in a single direction are not invalid, but are often handled poorly by servers, so clients SHOULD NOT close a stream for sending while they still expect to receive data from the target of the CONNECT.



A TCP connection error is signaled with RST\_STREAM. A proxy treats any error in the TCP connection, which includes receiving a TCP segment with the RST bit set, as a stream error of type HTTP\_CONNECT\_ERROR ([Section 6.1](#)). Correspondingly, a proxy MUST send a TCP segment with the RST bit set if it detects an error with the stream or the QUIC connection.

### **[3.1.3](#). Request Cancellation**

Either client or server can cancel requests by aborting the stream (QUIC RST\_STREAM or STOP\_SENDING frames, as appropriate) with an error code of HTTP\_REQUEST\_CANCELLED ([Section 6.1](#)). When the client cancels a response, it indicates that this response is no longer of interest. Clients SHOULD cancel requests by aborting both directions of a stream.

When the server cancels its response stream using HTTP\_REQUEST\_CANCELLED, it indicates that no application processing was performed. The client can treat requests cancelled by the server as though they had never been sent at all, thereby allowing them to be retried later on a new connection. Servers MUST NOT use the HTTP\_REQUEST\_CANCELLED status for requests which were partially or fully processed.

Note: In this context, "processed" means that some data from the stream was passed to some higher layer of software that might have taken some action as a result.

If a stream is cancelled after receiving a complete response, the client MAY ignore the cancellation and use the response. However, if a stream is cancelled after receiving a partial response, the response SHOULD NOT be used. Automatically retrying such requests is not possible, unless this is otherwise permitted (e.g., idempotent actions like GET, PUT, or DELETE).

### **[3.2](#). Request Prioritization**

HTTP/QUIC uses a priority scheme similar to that described in [\[RFC7540\]](#), [Section 5.3](#). In this priority scheme, a given stream can be designated as dependent upon another request, which expresses the preference that the latter stream (the "parent" request) be allocated resources before the former stream (the "dependent" request). Taken together, the dependencies across all requests in a connection form a dependency tree. The structure of the dependency tree changes as PRIORITY frames add, remove, or change the dependency links between requests.



The PRIORITY frame [Section 4.2.4](#) identifies a prioritized element. The elements which can be prioritized are:

- o Requests, identified by the ID of the request stream
- o Pushes, identified by the Push ID of the promised resource ([Section 4.2.7](#))
- o Placeholders, identified by a Placeholder ID

An element can depend on another element or on the root of the tree. A reference to an element which is no longer in the tree is treated as a reference to the root of the tree.

Only a client can send PRIORITY frames. A server MUST NOT send a PRIORITY frame.

### **[3.2.1.](#) Placeholders**

In HTTP/2, certain implementations used closed or unused streams as placeholders in describing the relative priority of requests. However, this created confusion as servers could not reliably identify which elements of the priority tree could safely be discarded. Clients could potentially reference closed streams long after the server had discarded state, leading to disparate views of the prioritization the client had attempted to express.

In HTTP/QUIC, a number of placeholders are explicitly permitted by the server using the "SETTINGS\_NUM\_PLACEHOLDERS" setting. Because the server commits to maintain these IDs in the tree, clients can use them with confidence that the server will not have discarded the state.

Placeholders are identified by an ID between zero and one less than the number of placeholders the server has permitted.

### **[3.2.2.](#) Priority Tree Maintenance**

Servers can aggressively prune inactive regions from the priority tree, because placeholders will be used to "root" any persistent structure of the tree which the client cares about retaining. For prioritization purposes, a node in the tree is considered "inactive" when the corresponding stream has been closed for at least two round-trip times (using any reasonable estimate available on the server). This delay helps mitigate race conditions where the server has pruned a node the client believed was still active and used as a Stream Dependency.



Specifically, the server MAY at any time:

- o Identify and discard branches of the tree containing only inactive nodes (i.e. a node with only other inactive nodes as descendants, along with those descendants)
- o Identify and condense interior regions of the tree containing only inactive nodes, allocating weight appropriately

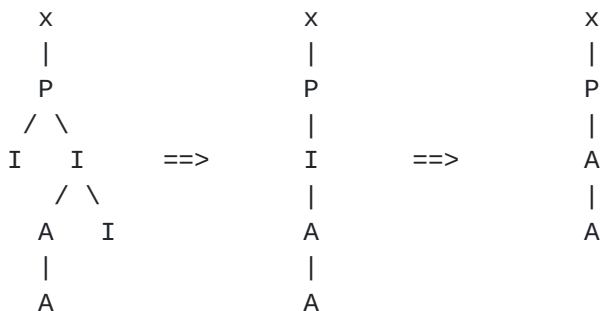


Figure 1: Example of Priority Tree Pruning

In the example in Figure 1, "P" represents a Placeholder, "A" represents an active node, and "I" represents an inactive node. In the first step, the server discards two inactive branches (each a single node). In the second step, the server condenses an interior inactive node. Note that these transformations will result in no change in the resources allocated to a particular active stream.

Clients SHOULD assume the server is actively performing such pruning and SHOULD NOT declare a dependency on a stream it knows to have been closed.

### 3.3. Unidirectional Streams

Unidirectional streams, in either direction, are used for a range of purposes. The purpose is indicated by a stream type, which is sent as a single octet header at the start of the stream. The format and structure of data that follows this header is determined by the stream type.

```

 0 1 2 3 4 5 6 7
+-+--+--+--+--+
|Stream Type (8)|
+-+--+--+--+--+
  
```

Figure 2: Unidirectional Stream Header





Some stream types are reserved ([Section 3.3.1](#)). Two stream types are defined in this document: control streams ([Section 3.3.2](#)) and push streams ([Section 3.3.3](#)). Other stream types can be defined by extensions to HTTP/QUIC.

If the stream header indicates a stream type which is not supported by the recipient, the remainder of the stream cannot be consumed as the semantics are unknown. Recipients of unknown stream types MAY trigger a QUIC STOP\_SENDING frame with an error code of HTTP\_UNKNOWN\_STREAM\_TYPE, but MUST NOT consider such streams to be an error of any kind.

Implementations MAY send stream types before knowing whether the peer supports them. However, stream types which could modify the state or semantics of existing protocol components, including QPACK or other extensions, MUST NOT be sent until the peer is known to support them.

### [3.3.1](#). Reserved Stream Types

Stream types of the format "0x1f \* N" are reserved to exercise the requirement that unknown types be ignored. These streams have no semantic meaning, and can be sent when application-layer padding is desired. They MAY also be sent on connections where no request data is currently being transferred. Endpoints MUST NOT consider these streams to have any meaning upon receipt.

The payload and length of the stream are selected in any manner the implementation chooses.

### [3.3.2](#). Control Streams

The control stream is indicated by a stream type of "0x43" (ASCII 'C'). Data on this stream consists of HTTP/QUIC frames, as defined in [Section 4.2](#).

Each side MUST initiate a single control stream at the beginning of the connection and send its SETTINGS frame as the first frame on this stream. If the first frame of the control stream is any other frame type, this MUST be treated as a connection error of type HTTP\_MISSING\_SETTINGS. Only one control stream per peer is permitted; receipt of a second stream which claims to be a control stream MUST be treated as a connection error of type HTTP\_WRONG\_STREAM\_COUNT. If the control stream is closed at any point, this MUST be treated as a connection error of type HTTP\_CLOSED\_CRITICAL\_STREAM.

A pair of unidirectional streams is used rather than a single bidirectional stream. This allows either peer to send data as soon



they are able. Depending on whether 0-RTT is enabled on the connection, either client or server might be able to send stream data first after the cryptographic handshake completes.

### 3.3.3. Server Push

HTTP/QUIC server push is similar to what is described in HTTP/2 [RFC7540], but uses different mechanisms.

The PUSH\_PROMISE frame ([Section 4.2.7](#)) is sent on the client-initiated bidirectional stream that carried the request that generated the push. This allows the server push to be associated with a request. Ordering of a PUSH\_PROMISE in relation to certain parts of the response is important (see [Section 8.2.1 of \[RFC7540\]](#)).

The PUSH\_PROMISE frame does not reference a stream; it contains a Push ID that uniquely identifies a server push. This allows a server to fulfill promises in the order that best suits its needs. The same Push ID can be used in multiple PUSH\_PROMISE frames (see [Section 4.2.7](#)). When a server later fulfills a promise, the server push response is conveyed on a push stream.

A push stream is indicated by a stream type of "0x50" (ASCII 'P'), followed by the Push ID of the promise that it fulfills, encoded as a variable-length integer. The remaining data on this stream consists of HTTP/QUIC frames, as defined in [Section 4.2](#), and carries the response side of an HTTP message exchange as described in [Section 3.1](#). The header of the request message is carried by a PUSH\_PROMISE frame (see [Section 4.2.7](#)) on the request stream which generated the push. Promised requests MUST conform to the requirements in [Section 8.2 of \[RFC7540\]](#).

Only servers can push; if a server receives a client-initiated push stream, this MUST be treated as a stream error of type HTTP\_WRONG\_STREAM\_DIRECTION.

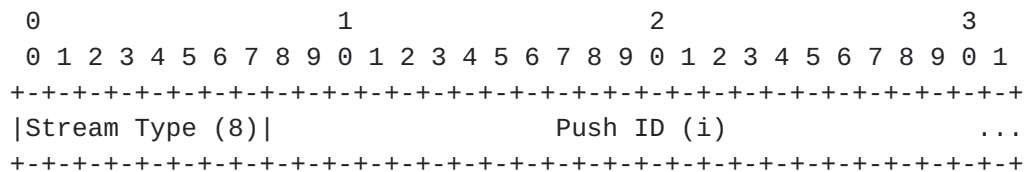


Figure 3: Push Stream Header

Server push is only enabled on a connection when a client sends a MAX\_PUSH\_ID frame (see [Section 4.2.9](#)). A server cannot use server push until it receives a MAX\_PUSH\_ID frame. A client sends additional MAX\_PUSH\_ID frames to control the number of pushes that a



server can promise. A server SHOULD use Push IDs sequentially, starting at 0. A client MUST treat receipt of a push stream with a Push ID that is greater than the maximum Push ID as a connection error of type HTTP\_PUSH\_LIMIT\_EXCEEDED.

Each Push ID MUST only be used once in a push stream header. If a push stream header includes a Push ID that was used in another push stream header, the client MUST treat this as a connection error of type HTTP\_DUPLICATE\_PUSH.

If a promised server push is not needed by the client, the client SHOULD send a CANCEL\_PUSH frame. If the push stream is already open, a QUIC STOP\_SENDING frame with an appropriate error code can be used instead (e.g., HTTP\_PUSH\_REFUSED, HTTP\_PUSH\_ALREADY\_IN\_CACHE; see [Section 6](#)). This asks the server not to transfer the data and indicates that it will be discarded upon receipt.

#### 4. HTTP Framing Layer

Frames are used on the control stream, request streams, and push streams. This section describes HTTP framing in QUIC and highlights some differences from HTTP/2 framing. For more detail on differences from HTTP/2, see [Section 8.2](#).

##### 4.1. Frame Layout

All frames have the following format:

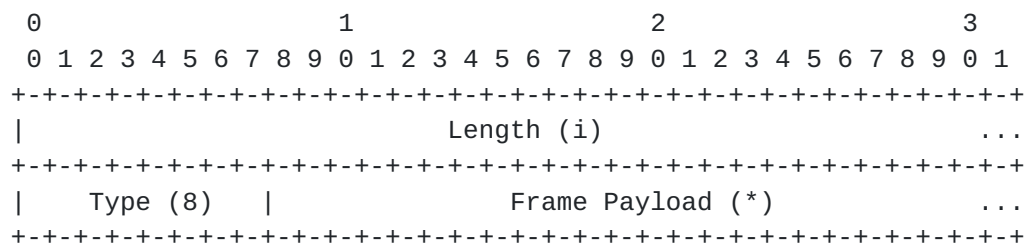


Figure 4: HTTP/QUIC frame format

A frame includes the following fields:

**Length:** A variable-length integer that describes the length of the Frame Payload. This length does not include the Type field.

**Type:** An 8-bit type for the frame.

**Frame Payload:** A payload, the semantics of which are determined by the Type field.



## 4.2. Frame Definitions

### 4.2.1. Reserved Frame Types

Frame types of the format "0xb + (0x1f \* N)" are reserved to exercise the requirement that unknown types be ignored. These frames have no semantic meaning, and can be sent when application-layer padding is desired. They MAY also be sent on connections where no request data is currently being transferred. Endpoints MUST NOT consider these frames to have any meaning upon receipt.

The payload and length of the frames are selected in any manner the implementation chooses.

### 4.2.2. DATA

DATA frames (type=0x0) convey arbitrary, variable-length sequences of octets associated with an HTTP request or response payload.

DATA frames MUST be associated with an HTTP request or response. If a DATA frame is received on either control stream, the recipient MUST respond with a connection error ([Section 6](#)) of type HTTP\_WRONG\_STREAM.

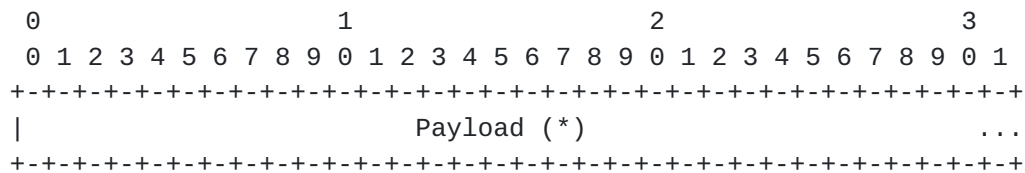


Figure 5: DATA frame payload

DATA frames MUST contain a non-zero-length payload. If a DATA frame is received with a payload length of zero, the recipient MUST respond with a stream error ([Section 6](#)) of type HTTP\_MALFORMED\_FRAME.

### 4.2.3. HEADERS

The HEADERS frame (type=0x1) is used to carry a header block, compressed using QPACK. See [[QPACK](#)] for more details.

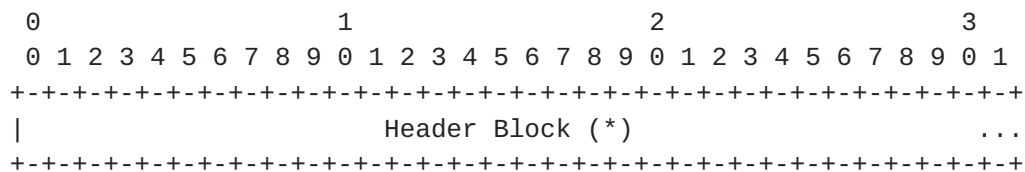


Figure 6: HEADERS frame payload





HEADERS frames can only be sent on request / push streams.

#### 4.2.4. PRIORITY

The PRIORITY (type=0x02) frame specifies the sender-advised priority of a stream and is substantially different in format from [\[RFC7540\]](#). In order to ensure that prioritization is processed in a consistent order, PRIORITY frames MUST be sent on the control stream. A PRIORITY frame sent on any other stream MUST be treated as a HTTP\_WRONG\_STREAM error.

The format has been modified to accommodate not being sent on a request stream, to allow for identification of server pushes, and the larger stream ID space of QUIC. The semantics of the Stream Dependency, Weight, and E flag are otherwise the same as in HTTP/2.

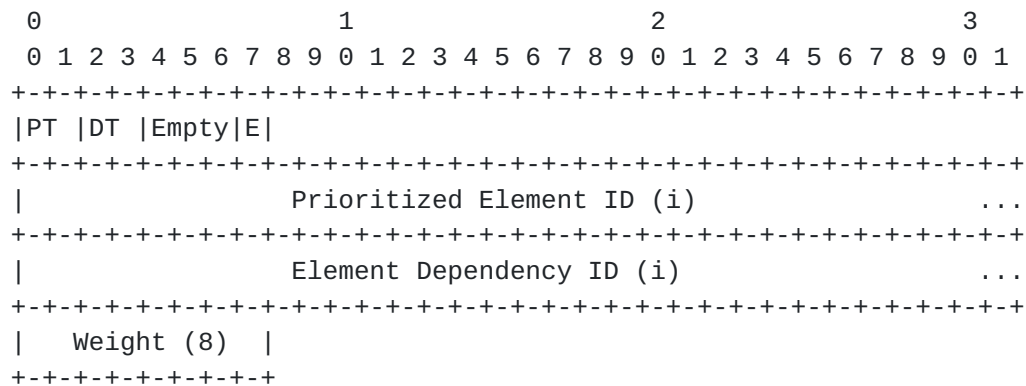


Figure 7: PRIORITY frame payload

The PRIORITY frame payload has the following fields:

**Prioritized Type:** A two-bit field indicating the type of element being prioritized.

**Dependency Type:** A two-bit field indicating the type of element being depended on.

**Empty:** A three-bit field which MUST be zero when sent and MUST be ignored on receipt.

**Exclusive:** A flag which indicates that the stream dependency is exclusive (see [\[RFC7540\]](#), [Section 5.3](#)).

**Prioritized Element ID:** A variable-length integer that identifies the element being prioritized. Depending on the value of Prioritized Type, this contains the Stream ID of a request stream,



the Push ID of a promised resource, or a Placeholder ID of a placeholder.

**Element Dependency ID:** A variable-length integer that identifies the element on which a dependency is being expressed. Depending on the value of Dependency Type, this contains the Stream ID of a request stream, the Push ID of a promised resource, or a Placeholder ID of a placeholder. For details of dependencies, see [Section 3.2](#) and [\[RFC7540\], Section 5.3](#).

**Weight:** An unsigned 8-bit integer representing a priority weight for the stream (see [\[RFC7540\], Section 5.3](#)). Add one to the value to obtain a weight between 1 and 256.

A PRIORITY frame identifies an element to prioritize, and an element upon which it depends. A Prioritized ID or Dependency ID identifies a client-initiated request using the corresponding stream ID, a server push using a Push ID (see [Section 4.2.7](#)), or a placeholder using a Placeholder ID (see [Section 3.2.1](#)).

The values for the Prioritized Element Type and Element Dependency Type imply the interpretation of the associated Element ID fields.

Type Bits	Type Description	Element ID Contents
00	Request stream	Stream ID
01	Push stream	Push ID
10	Placeholder	Placeholder ID
11	Root of the tree	Ignored

Note that the root of the tree cannot be referenced using a Stream ID of 0, as in [\[RFC7540\]](#); QUIC stream 0 carries a valid HTTP request. The root of the tree cannot be reprioritized. A PRIORITY frame that prioritizes the root of the tree MUST be treated as a connection error of type HTTP\_MALFORMED\_FRAME.

When a PRIORITY frame claims to reference a request, the associated ID MUST identify a client-initiated bidirectional stream. A server MUST treat receipt of PRIORITY frame with a Stream ID of any other type as a connection error of type HTTP\_MALFORMED\_FRAME.



A PRIORITY frame that references a non-existent Push ID or a Placeholder ID greater than the server's limit MUST be treated as a HTTP\_MALFORMED\_FRAME error.

A PRIORITY frame MUST contain only the identified fields. A PRIORITY frame that contains more or fewer fields, or a PRIORITY frame that includes a truncated integer encoding MUST be treated as a connection error of type HTTP\_MALFORMED\_FRAME.

#### 4.2.5. CANCEL\_PUSH

The CANCEL\_PUSH frame (type=0x3) is used to request cancellation of server push prior to the push stream being created. The CANCEL\_PUSH frame identifies a server push request by Push ID (see [Section 4.2.7](#)) using a variable-length integer.

When a server receives this frame, it aborts sending the response for the identified server push. If the server has not yet started to send the server push, it can use the receipt of a CANCEL\_PUSH frame to avoid opening a stream. If the push stream has been opened by the server, the server SHOULD send a QUIC RST\_STREAM frame on those streams and cease transmission of the response.

A server can send this frame to indicate that it won't be sending a response prior to creation of a push stream. Once the push stream has been created, sending CANCEL\_PUSH has no effect on the state of the push stream. A QUIC RST\_STREAM frame SHOULD be used instead to cancel transmission of the server push response.

A CANCEL\_PUSH frame is sent on the control stream. Sending a CANCEL\_PUSH frame on a stream other than the control stream MUST be treated as a stream error of type HTTP\_WRONG\_STREAM.

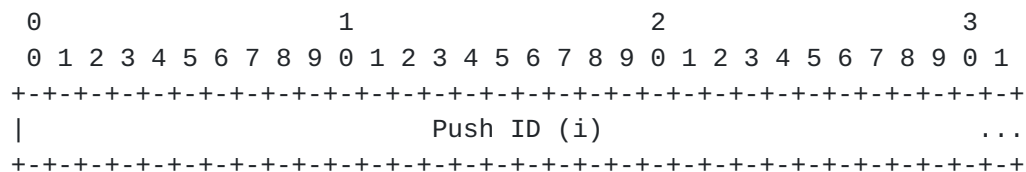


Figure 8: CANCEL\_PUSH frame payload

The CANCEL\_PUSH frame carries a Push ID encoded as a variable-length integer. The Push ID identifies the server push that is being cancelled (see [Section 4.2.7](#)).

If the client receives a CANCEL\_PUSH frame, that frame might identify a Push ID that has not yet been mentioned by a PUSH\_PROMISE frame.



An endpoint MUST treat a CANCEL\_PUSH frame which does not contain exactly one properly-formatted variable-length integer as a connection error of type HTTP\_MALFORMED\_FRAME.

#### 4.2.6. SETTINGS

The SETTINGS frame (type=0x4) conveys configuration parameters that affect how endpoints communicate, such as preferences and constraints on peer behavior, and is different from [RFC7540]. Individually, a SETTINGS parameter can also be referred to as a "setting".

SETTINGS parameters are not negotiated; they describe characteristics of the sending peer, which can be used by the receiving peer. However, a negotiation can be implied by the use of SETTINGS - a peer uses SETTINGS to advertise a set of supported values. The recipient can then choose which entries from this list are also acceptable and proceed with the value it has chosen. (This choice could be announced in a field of an extension frame, or in its own value in SETTINGS.)

Different values for the same parameter can be advertised by each peer. For example, a client might be willing to consume a very large response header, while servers are more cautious about request size.

Parameters MUST NOT occur more than once. A receiver MAY treat the presence of the same parameter more than once as a connection error of type HTTP\_MALFORMED\_FRAME.

The payload of a SETTINGS frame consists of zero or more parameters, each consisting of an unsigned 16-bit setting identifier and a value which uses the QUIC variable-length integer encoding.

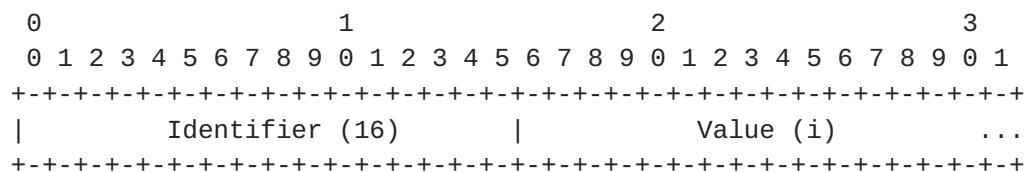


Figure 9: SETTINGS value format

Each value MUST be compared against the remaining length of the SETTINGS frame. Any value which purports to cross the end of the frame MUST cause the SETTINGS frame to be considered malformed and trigger a connection error of type HTTP\_MALFORMED\_FRAME.

An implementation MUST ignore the contents for any SETTINGS identifier it does not understand.





SETTINGS frames always apply to a connection, never a single stream. A SETTINGS frame MUST be sent as the first frame of either control stream (see [Section 3](#)) by each peer, and MUST NOT be sent subsequently or on any other stream. If an endpoint receives a SETTINGS frame on a different stream, the endpoint MUST respond with a connection error of type HTTP\_WRONG\_STREAM. If an endpoint receives a second SETTINGS frame, the endpoint MUST respond with a connection error of type HTTP\_MALFORMED\_FRAME.

The SETTINGS frame affects connection state. A badly formed or incomplete SETTINGS frame MUST be treated as a connection error ([Section 6](#)) of type HTTP\_MALFORMED\_FRAME.

#### **[4.2.6.1](#). Defined SETTINGS Parameters**

The following settings are defined in HTTP/QUIC:

SETTINGS\_NUM\_PLACEHOLDERS (0x3): This value SHOULD be non-zero. The default value is 16.

SETTINGS\_MAX\_HEADER\_LIST\_SIZE (0x6): The default value is unlimited.

Settings values of the format "0x?a?a" are reserved to exercise the requirement that unknown parameters be ignored. Such settings have no defined meaning. Endpoints SHOULD include at least one such setting in their SETTINGS frame. Endpoints MUST NOT consider such settings to have any meaning upon receipt.

Because the setting has no defined meaning, the value of the setting can be any value the implementation selects.

Additional settings MAY be defined by extensions to HTTP/QUIC.

#### **[4.2.6.2](#). Initial SETTINGS Values**

When a 0-RTT QUIC connection is being used, the client's initial requests will be sent before the arrival of the server's SETTINGS frame. Clients MUST store the settings the server provided in the session being resumed and MUST comply with stored settings until the server's current settings are received. Remembered settings apply to the new connection until the server's SETTINGS frame is received.

A server can remember the settings that it advertised, or store an integrity-protected copy of the values in the ticket and recover the information when accepting 0-RTT data. A server uses the HTTP/QUIC settings values in determining whether to accept 0-RTT data.



A server MAY accept 0-RTT and subsequently provide different settings in its SETTINGS frame. If 0-RTT data is accepted by the server, its SETTINGS frame MUST NOT reduce any limits or alter any values that might be violated by the client with its 0-RTT data.

When a 1-RTT QUIC connection is being used, the client MUST NOT send requests prior to receiving and processing the server's SETTINGS frame.

#### 4.2.7. PUSH\_PROMISE

The PUSH\_PROMISE frame (type=0x05) is used to carry a request header set from server to client, as in HTTP/2.

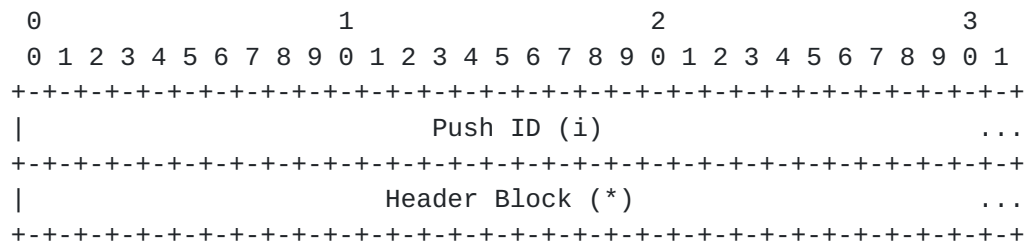


Figure 10: PUSH\_PROMISE frame payload

The payload consists of:

**Push ID:** A variable-length integer that identifies the server push request. A push ID is used in push stream header ([Section 3.3.3](#)), CANCEL\_PUSH frames ([Section 4.2.5](#)), and PRIORITY frames ([Section 4.2.4](#)).

**Header Block:** QPACK-compressed request header fields for the promised response. See [[QPACK](#)] for more details.

A server MUST NOT use a Push ID that is larger than the client has provided in a MAX\_PUSH\_ID frame ([Section 4.2.9](#)). A client MUST treat receipt of a PUSH\_PROMISE that contains a larger Push ID than the client has advertised as a connection error of type HTTP\_MALFORMED\_FRAME.

A server MAY use the same Push ID in multiple PUSH\_PROMISE frames. This allows the server to use the same server push in response to multiple concurrent requests. Referencing the same server push ensures that a PUSH\_PROMISE can be made in relation to every response in which server push might be needed without duplicating pushes.

A server that uses the same Push ID in multiple PUSH\_PROMISE frames MUST include the same header fields each time. The octets of the



header block MAY be different due to differing encoding, but the header fields and their values MUST be identical. Note that ordering of header fields is significant. A client MUST treat receipt of a PUSH\_PROMISE with conflicting header field values for the same Push ID as a connection error of type HTTP\_MALFORMED\_FRAME.

Allowing duplicate references to the same Push ID is primarily to reduce duplication caused by concurrent requests. A server SHOULD avoid reusing a Push ID over a long period. Clients are likely to consume server push responses and not retain them for reuse over time. Clients that see a PUSH\_PROMISE that uses a Push ID that they have since consumed and discarded are forced to ignore the PUSH\_PROMISE.

#### 4.2.8. GOAWAY

The GOAWAY frame (type=0x7) is used to initiate graceful shutdown of a connection by a server. GOAWAY allows a server to stop accepting new requests while still finishing processing of previously received requests. This enables administrative actions, like server maintenance. GOAWAY by itself does not close a connection.

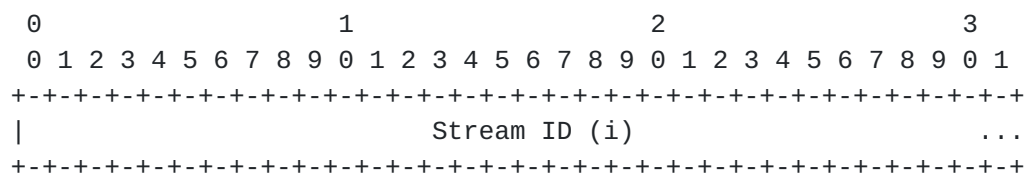


Figure 11: GOAWAY frame payload

The GOAWAY frame carries a QUIC Stream ID for a client-initiated bidirectional stream encoded as a variable-length integer. A client MUST treat receipt of a GOAWAY frame containing a Stream ID of any other type as a connection error of type HTTP\_MALFORMED\_FRAME.

Clients do not need to send GOAWAY to initiate a graceful shutdown; they simply stop making new requests. A server MUST treat receipt of a GOAWAY frame as a connection error ([Section 6](#)) of type HTTP\_UNEXPECTED\_GOAWAY.

The GOAWAY frame applies to the connection, not a specific stream. An endpoint MUST treat a GOAWAY frame on a stream other than the control stream as a connection error ([Section 6](#)) of type HTTP\_WRONG\_STREAM.

See [Section 5.2](#) for more information on the use of the GOAWAY frame.



#### 4.2.9. MAX\_PUSH\_ID

The MAX\_PUSH\_ID frame (type=0xD) is used by clients to control the number of server pushes that the server can initiate. This sets the maximum value for a Push ID that the server can use in a PUSH\_PROMISE frame. Consequently, this also limits the number of push streams that the server can initiate in addition to the limit set by the QUIC MAX\_STREAM\_ID frame.

The MAX\_PUSH\_ID frame is always sent on a control stream. Receipt of a MAX\_PUSH\_ID frame on any other stream MUST be treated as a connection error of type HTTP\_WRONG\_STREAM.

A server MUST NOT send a MAX\_PUSH\_ID frame. A client MUST treat the receipt of a MAX\_PUSH\_ID frame as a connection error of type HTTP\_MALFORMED\_FRAME.

The maximum Push ID is unset when a connection is created, meaning that a server cannot push until it receives a MAX\_PUSH\_ID frame. A client that wishes to manage the number of promised server pushes can increase the maximum Push ID by sending a MAX\_PUSH\_ID frame as the server fulfills or cancels server pushes.

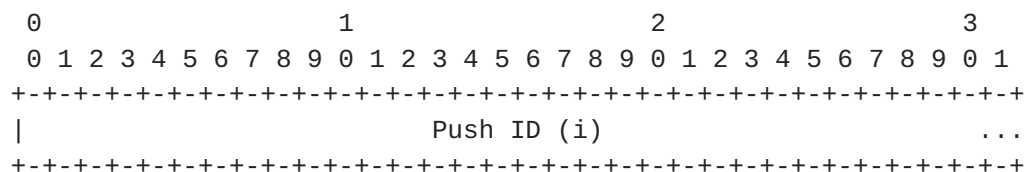


Figure 12: MAX\_PUSH\_ID frame payload

The MAX\_PUSH\_ID frame carries a single variable-length integer that identifies the maximum value for a Push ID that the server can use (see [Section 4.2.7](#)). A MAX\_PUSH\_ID frame cannot reduce the maximum Push ID; receipt of a MAX\_PUSH\_ID that contains a smaller value than previously received MUST be treated as a connection error of type HTTP\_MALFORMED\_FRAME.

A server MUST treat a MAX\_PUSH\_ID frame payload that does not contain a single variable-length integer as a connection error of type HTTP\_MALFORMED\_FRAME.

## 5. Connection Closure

Once established, an HTTP/QUIC connection can be used for many requests and responses over time until the connection is closed. Connection closure can happen in any of several different ways.





### **5.1. Idle Connections**

Each QUIC endpoint declares an idle timeout during the handshake. If the connection remains idle (no packets received) for longer than this duration, the peer will assume that the connection has been closed. HTTP/QUIC implementations will need to open a new connection for new requests if the existing connection has been idle for longer than the server's advertised idle timeout, and SHOULD do so if approaching the idle timeout.

HTTP clients are expected to use QUIC PING frames to keep connections open while there are responses outstanding for requests or server pushes. If the client is not expecting a response from the server, allowing an idle connection to time out is preferred over expending effort maintaining a connection that might not be needed. A gateway MAY use PING to maintain connections in anticipation of need rather than incur the latency cost of connection establishment to servers. Servers SHOULD NOT use PING frames to keep a connection open.

### **5.2. Connection Shutdown**

Even when a connection is not idle, either endpoint can decide to stop using the connection and let the connection close gracefully. Since clients drive request generation, clients perform a connection shutdown by not sending additional requests on the connection; responses and pushed responses associated to previous requests will continue to completion. Servers perform the same function by communicating with clients.

Servers initiate the shutdown of a connection by sending a GOAWAY frame ([Section 4.2.8](#)). The GOAWAY frame indicates that client-initiated requests on lower stream IDs were or might be processed in this connection, while requests on the indicated stream ID and greater were not accepted. This enables client and server to agree on which requests were accepted prior to the connection shutdown. This identifier MAY be lower than the stream limit identified by a QUIC MAX\_STREAM\_ID frame, and MAY be zero if no requests were processed. Servers SHOULD NOT increase the QUIC MAX\_STREAM\_ID limit after sending a GOAWAY frame.

Once sent, the server MUST cancel requests sent on streams with an identifier higher than the indicated last Stream ID. Clients MUST NOT send new requests on the connection after receiving GOAWAY, although requests might already be in transit. A new connection can be established for new requests.

If the client has sent requests on streams with a higher Stream ID than indicated in the GOAWAY frame, those requests are considered



cancelled ([Section 3.1.3](#)). Clients SHOULD reset any streams above this ID with the error code HTTP\_REQUEST\_CANCELLED. Servers MAY also cancel requests on streams below the indicated ID if these requests were not processed.

Requests on Stream IDs less than the Stream ID in the GOAWAY frame might have been processed; their status cannot be known until they are completed successfully, reset individually, or the connection terminates.

Servers SHOULD send a GOAWAY frame when the closing of a connection is known in advance, even if the advance notice is small, so that the remote peer can know whether a stream has been partially processed or not. For example, if an HTTP client sends a POST at the same time that a server closes a QUIC connection, the client cannot know if the server started to process that POST request if the server does not send a GOAWAY frame to indicate what streams it might have acted on.

A client that is unable to retry requests loses all requests that are in flight when the server closes the connection. A server MAY send multiple GOAWAY frames indicating different stream IDs, but MUST NOT increase the value they send in the last Stream ID, since clients might already have retried unprocessed requests on another connection. A server that is attempting to gracefully shut down a connection SHOULD send an initial GOAWAY frame with the last Stream ID set to the current value of QUIC's MAX\_STREAM\_ID and SHOULD NOT increase the MAX\_STREAM\_ID thereafter. This signals to the client that a shutdown is imminent and that initiating further requests is prohibited. After allowing time for any in-flight requests (at least one round-trip time), the server MAY send another GOAWAY frame with an updated last Stream ID. This ensures that a connection can be cleanly shut down without losing requests.

Once all accepted requests have been processed, the server can permit the connection to become idle, or MAY initiate an immediate closure of the connection. An endpoint that completes a graceful shutdown SHOULD use the HTTP\_NO\_ERROR code when closing the connection.

### **5.3. Immediate Application Closure**

An HTTP/QUIC implementation can immediately close the QUIC connection at any time. This results in sending a QUIC APPLICATION\_CLOSE frame to the peer; the error code in this frame indicates to the peer why the connection is being closed. See [Section 6](#) for error codes which can be used when closing a connection.

Before closing the connection, a GOAWAY MAY be sent to allow the client to retry some requests. Including the GOAWAY frame in the



same packet as the QUIC APPLICATION\_CLOSE frame improves the chances of the frame being received by clients.

#### **5.4. Transport Closure**

For various reasons, the QUIC transport could indicate to the application layer that the connection has terminated. This might be due to an explicit closure by the peer, a transport-level error, or a change in network topology which interrupts connectivity.

If a connection terminates without a GOAWAY frame, clients **MUST** assume that any request which was sent, whether in whole or in part, might have been processed.

### **6. Error Handling**

QUIC allows the application to abruptly terminate (reset) individual streams or the entire connection when an error is encountered. These are referred to as "stream errors" or "connection errors" and are described in more detail in [[QUIC-TRANSPORT](#)].

This section describes HTTP/QUIC-specific error codes which can be used to express the cause of a connection or stream error.

#### **6.1. HTTP/QUIC Error Codes**

The following error codes are defined for use in QUIC RST\_STREAM, STOP\_SENDING, and APPLICATION\_CLOSE frames when using HTTP/QUIC.

STOPPING (0x00): This value is reserved by the transport to be used in response to QUIC STOP\_SENDING frames.

HTTP\_NO\_ERROR (0x01): No error. This is used when the connection or stream needs to be closed, but there is no error to signal.

HTTP\_PUSH\_REFUSED (0x02): The server has attempted to push content which the client will not accept on this connection.

HTTP\_INTERNAL\_ERROR (0x03): An internal error has occurred in the HTTP stack.

HTTP\_PUSH\_ALREADY\_IN\_CACHE (0x04): The server has attempted to push content which the client has cached.

HTTP\_REQUEST\_CANCELLED (0x05): The client no longer needs the requested data.



HTTP\_INCOMPLETE\_REQUEST (0x06): The client's stream terminated without containing a fully-formed request.

HTTP\_CONNECT\_ERROR (0x07): The connection established in response to a CONNECT request was reset or abnormally closed.

HTTP\_EXCESSIVE\_LOAD (0x08): The endpoint detected that its peer is exhibiting a behavior that might be generating excessive load.

HTTP\_VERSION\_FALLBACK (0x09): The requested operation cannot be served over HTTP/QUIC. The peer should retry over HTTP/1.1.

HTTP\_WRONG\_STREAM (0x0A): A frame was received on a stream where it is not permitted.

HTTP\_PUSH\_LIMIT\_EXCEEDED (0x0B): A Push ID greater than the current maximum Push ID was referenced.

HTTP\_DUPLICATE\_PUSH (0x0C): A Push ID was referenced in two different stream headers.

HTTP\_UNKNOWN\_STREAM\_TYPE (0x0D): A unidirectional stream header contained an unknown stream type.

HTTP\_WRONG\_STREAM\_COUNT (0x0E): A unidirectional stream type was used more times than is permitted by that type.

HTTP\_CLOSED\_CRITICAL\_STREAM (0x0F): A stream required by the connection was closed or reset.

HTTP\_WRONG\_STREAM\_DIRECTION (0x0010): A unidirectional stream type was used by a peer which is not permitted to do so.

HTTP\_EARLY\_RESPONSE (0x0011): The remainder of the client's request is not needed to produce a response. For use in STOP\_SENDING only.

HTTP\_MISSING\_SETTINGS (0x0012): No SETTINGS frame was received at the beginning of the control stream.

HTTP\_GENERAL\_PROTOCOL\_ERROR (0x00FF): Peer violated protocol requirements in a way which doesn't match a more specific error code, or endpoint declines to use the more specific error code.

HTTP\_MALFORMED\_FRAME (0x01XX): An error in a specific frame type. The frame type is included as the last octet of the error code. For example, an error in a MAX\_PUSH\_ID frame would be indicated with the code (0x10D).





## 7. Extensions to HTTP/QUIC

HTTP/QUIC permits extension of the protocol. Within the limitations described in this section, protocol extensions can be used to provide additional services or alter any aspect of the protocol. Extensions are effective only within the scope of a single HTTP/QUIC connection.

This applies to the protocol elements defined in this document. This does not affect the existing options for extending HTTP, such as defining new methods, status codes, or header fields.

Extensions are permitted to use new frame types ([Section 4.2](#)), new settings ([Section 4.2.6.1](#)), new error codes ([Section 6](#)), or new unidirectional stream types ([Section 3.3](#)). Registries are established for managing these extension points: frame types ([Section 10.3](#)), settings ([Section 10.4](#)), error codes ([Section 10.5](#)), and stream types ([Section 10.6](#)).

Implementations MUST ignore unknown or unsupported values in all extensible protocol elements. Implementations MUST discard frames and unidirectional streams that have unknown or unsupported types. This means that any of these extension points can be safely used by extensions without prior arrangement or negotiation.

Extensions that could change the semantics of existing protocol components MUST be negotiated before being used. For example, an extension that changes the layout of the HEADERS frame cannot be used until the peer has given a positive signal that this is acceptable. In this case, it could also be necessary to coordinate when the revised layout comes into effect.

This document doesn't mandate a specific method for negotiating the use of an extension but notes that a setting ([Section 4.2.6.1](#)) could be used for that purpose. If both peers set a value that indicates willingness to use the extension, then the extension can be used. If a setting is used for extension negotiation, the default value MUST be defined in such a fashion that the extension is disabled if the setting is omitted.

## 8. Considerations for Transitioning from HTTP/2

HTTP/QUIC is strongly informed by HTTP/2, and bears many similarities. This section describes the approach taken to design HTTP/QUIC, points out important differences from HTTP/2, and describes how to map HTTP/2 extensions into HTTP/QUIC.

HTTP/QUIC begins from the premise that HTTP/2 code reuse is a useful feature, but not a hard requirement. HTTP/QUIC departs from HTTP/2



primarily where necessary to accommodate the differences in behavior between QUIC and TCP (lack of ordering, support for streams). We intend to avoid gratuitous changes which make it difficult or impossible to build extensions with the same semantics applicable to both protocols at once.

These departures are noted in this section.

### **8.1. Streams**

HTTP/QUIC permits use of a larger number of streams ( $2^{62}-1$ ) than HTTP/2. The considerations about exhaustion of stream identifier space apply, though the space is significantly larger such that it is likely that other limits in QUIC are reached first, such as the limit on the connection flow control window.

### **8.2. HTTP Frame Types**

Many framing concepts from HTTP/2 can be elided away on QUIC, because the transport deals with them. Because frames are already on a stream, they can omit the stream number. Because frames do not block multiplexing (QUIC's multiplexing occurs below this layer), the support for variable-maximum-length packets can be removed. Because stream termination is handled by QUIC, an `END_STREAM` flag is not required. This permits the removal of the `Flags` field from the generic frame layout.

Frame payloads are largely drawn from [\[RFC7540\]](#). However, QUIC includes many features (e.g. flow control) which are also present in HTTP/2. In these cases, the HTTP mapping does not re-implement them. As a result, several HTTP/2 frame types are not required in HTTP/QUIC. Where an HTTP/2-defined frame is no longer used, the frame ID has been reserved in order to maximize portability between HTTP/2 and HTTP/QUIC implementations. However, even equivalent frames between the two mappings are not identical.

Many of the differences arise from the fact that HTTP/2 provides an absolute ordering between frames across all streams, while QUIC provides this guarantee on each stream only. As a result, if a frame type makes assumptions that frames from different streams will still be received in the order sent, HTTP/QUIC will break them.

For example, implicit in the HTTP/2 prioritization scheme is the notion of in-order delivery of priority changes (i.e., dependency tree mutations): since operations on the dependency tree such as reparenting a subtree are not commutative, both sender and receiver must apply them in the same order to ensure that both sides have a consistent view of the stream dependency tree. HTTP/2 specifies



priority assignments in PRIORITY frames and (optionally) in HEADERS frames. To achieve in-order delivery of priority changes in HTTP/QUIC, PRIORITY frames are sent on the control stream and the PRIORITY section is removed from the HEADERS frame.

Likewise, HPACK was designed with the assumption of in-order delivery. A sequence of encoded header blocks must arrive (and be decoded) at an endpoint in the same order in which they were encoded. This ensures that the dynamic state at the two endpoints remains in sync. As a result, HTTP/QUIC uses a modified version of HPACK, described in [\[QPACK\]](#).

Frame type definitions in HTTP/QUIC often use the QUIC variable-length integer encoding. In particular, Stream IDs use this encoding, which allow for a larger range of possible values than the encoding used in HTTP/2. Some frames in HTTP/QUIC use an identifier rather than a Stream ID (e.g. Push IDs in PRIORITY frames). Redefinition of the encoding of extension frame types might be necessary if the encoding includes a Stream ID.

Because the Flags field is not present in generic HTTP/QUIC frames, those frames which depend on the presence of flags need to allocate space for flags as part of their frame payload.

Other than this issue, frame type HTTP/2 extensions are typically portable to QUIC simply by replacing Stream 0 in HTTP/2 with a control stream in HTTP/QUIC. HTTP/QUIC extensions will not assume ordering, but would not be harmed by ordering, and would be portable to HTTP/2 in the same manner.

Below is a listing of how each HTTP/2 frame type is mapped:

DATA (0x0): Padding is not defined in HTTP/QUIC frames. See [Section 4.2.2](#).

HEADERS (0x1): As described above, the PRIORITY region of HEADERS is not supported. A separate PRIORITY frame MUST be used. Padding is not defined in HTTP/QUIC frames. See [Section 4.2.3](#).

PRIORITY (0x2): As described above, the PRIORITY frame is sent on the control stream and can reference either a Stream ID or a Push ID. See [Section 4.2.4](#).

RST\_STREAM (0x3): RST\_STREAM frames do not exist, since QUIC provides stream lifecycle management. The same code point is used for the CANCEL\_PUSH frame ([Section 4.2.5](#)).



SETTINGS (0x4): SETTINGS frames are sent only at the beginning of the connection. See [Section 4.2.6](#) and [Section 8.3](#).

PUSH\_PROMISE (0x5): The PUSH\_PROMISE does not reference a stream; instead the push stream references the PUSH\_PROMISE frame using a Push ID. See [Section 4.2.7](#).

PING (0x6): PING frames do not exist, since QUIC provides equivalent functionality.

GOAWAY (0x7): GOAWAY is sent only from server to client and does not contain an error code. See [Section 4.2.8](#).

WINDOW\_UPDATE (0x8): WINDOW\_UPDATE frames do not exist, since QUIC provides flow control.

CONTINUATION (0x9): CONTINUATION frames do not exist; instead, larger HEADERS/PUSH\_PROMISE frames than HTTP/2 are permitted, and HEADERS frames can be used in series.

Frame types defined by extensions to HTTP/2 need to be separately registered for HTTP/QUIC if still applicable. The IDs of frames defined in [[RFC7540](#)] have been reserved for simplicity. See [Section 10.3](#).

### **[8.3](#). HTTP/2 SETTINGS Parameters**

An important difference from HTTP/2 is that settings are sent once, at the beginning of the connection, and thereafter cannot change. This eliminates many corner cases around synchronization of changes.

Some transport-level options that HTTP/2 specifies via the SETTINGS frame are superseded by QUIC transport parameters in HTTP/QUIC. The HTTP-level options that are retained in HTTP/QUIC have the same value as in HTTP/2.

Below is a listing of how each HTTP/2 SETTINGS parameter is mapped:

SETTINGS\_HEADER\_TABLE\_SIZE: See [Section 4.2.6.1](#).

SETTINGS\_ENABLE\_PUSH: This is removed in favor of the MAX\_PUSH\_ID which provides a more granular control over server push.

SETTINGS\_MAX\_CONCURRENT\_STREAMS: QUIC controls the largest open Stream ID as part of its flow control logic. Specifying SETTINGS\_MAX\_CONCURRENT\_STREAMS in the SETTINGS frame is an error.





SETTINGS\_INITIAL\_WINDOW\_SIZE: QUIC requires both stream and connection flow control window sizes to be specified in the initial transport handshake. Specifying SETTINGS\_INITIAL\_WINDOW\_SIZE in the SETTINGS frame is an error.

SETTINGS\_MAX\_FRAME\_SIZE: This setting has no equivalent in HTTP/QUIC. Specifying it in the SETTINGS frame is an error.

SETTINGS\_MAX\_HEADER\_LIST\_SIZE: See [Section 4.2.6.1](#).

In HTTP/QUIC, setting values are variable-length integers (6, 14, 30, or 62 bits long) rather than fixed-length 32-bit fields as in HTTP/2. This will often produce a shorter encoding, but can produce a longer encoding for settings which use the full 32-bit space. Settings ported from HTTP/2 might choose to redefine the format of their settings to avoid using the 62-bit encoding.

Settings need to be defined separately for HTTP/2 and HTTP/QUIC. The IDs of settings defined in [\[RFC7540\]](#) have been reserved for simplicity. See [Section 10.4](#).

#### **[8.4](#). HTTP/2 Error Codes**

QUIC has the same concepts of "stream" and "connection" errors that HTTP/2 provides. However, because the error code space is shared between multiple components, there is no direct portability of HTTP/2 error codes.

The HTTP/2 error codes defined in [Section 7 of \[RFC7540\]](#) map to the HTTP/QUIC error codes as follows:

NO\_ERROR (0x0): HTTP\_NO\_ERROR in [Section 6.1](#).

PROTOCOL\_ERROR (0x1): No single mapping. See new HTTP\_MALFORMED\_FRAME error codes defined in [Section 6.1](#).

INTERNAL\_ERROR (0x2): HTTP\_INTERNAL\_ERROR in [Section 6.1](#).

FLOW\_CONTROL\_ERROR (0x3): Not applicable, since QUIC handles flow control. Would provoke a QUIC\_FLOW\_CONTROL\_RECEIVED\_TOO\_MUCH\_DATA from the QUIC layer.

SETTINGS\_TIMEOUT (0x4): Not applicable, since no acknowledgement of SETTINGS is defined.

STREAM\_CLOSED (0x5): Not applicable, since QUIC handles stream management. Would provoke a QUIC\_STREAM\_DATA\_AFTER\_TERMINATION from the QUIC layer.



FRAME\_SIZE\_ERROR (0x6): No single mapping. See new error codes defined in [Section 6.1](#).

REFUSED\_STREAM (0x7): Not applicable, since QUIC handles stream management. Would provoke a QUIC\_TOO\_MANY\_OPEN\_STREAMS from the QUIC layer.

CANCEL (0x8): HTTP\_REQUEST\_CANCELLED in [Section 6.1](#).

COMPRESSION\_ERROR (0x9): HTTP\_QPACK\_DECOMPRESSION\_FAILED in [\[QPACK\]](#).

CONNECT\_ERROR (0xa): HTTP\_CONNECT\_ERROR in [Section 6.1](#).

ENHANCE\_YOUR\_CALM (0xb): HTTP\_EXCESSIVE\_LOAD in [Section 6.1](#).

INADEQUATE\_SECURITY (0xc): Not applicable, since QUIC is assumed to provide sufficient security on all connections.

HTTP\_1\_1\_REQUIRED (0xd): HTTP\_VERSION\_FALLBACK in [Section 6.1](#).

Error codes need to be defined for HTTP/2 and HTTP/QUIC separately. See [Section 10.5](#).

## **9. Security Considerations**

The security considerations of HTTP/QUIC should be comparable to those of HTTP/2 with TLS. Note that where HTTP/2 employs PADDING frames to make a connection more resistant to traffic analysis, HTTP/QUIC can rely on QUIC's own PADDING frames or employ the reserved frame and stream types discussed in [Section 4.2.1](#) and [Section 3.3.1](#).

When HTTP Alternative Services is used for discovery for HTTP/QUIC endpoints, the security considerations of [\[ALTSVC\]](#) also apply.

The modified SETTINGS format contains nested length elements, which could pose a security risk to an incautious implementer. A SETTINGS frame parser MUST ensure that the length of the frame exactly matches the length of the settings it contains.

## **10. IANA Considerations**

### **10.1. Registration of HTTP/QUIC Identification String**

This document creates a new registration for the identification of HTTP/QUIC in the "Application Layer Protocol Negotiation (ALPN) Protocol IDs" registry established in [\[RFC7301\]](#).

The "hq" string identifies HTTP/QUIC:



Protocol: HTTP/QUIC

Identification Sequence: 0x68 0x71 ("hq")

Specification: This document

### **10.2. Registration of QUIC Version Hint Alt-Svc Parameter**

This document creates a new registration for version-negotiation hints in the "Hypertext Transfer Protocol (HTTP) Alt-Svc Parameter" registry established in [[RFC7838](#)].

Parameter: "quic"

Specification: This document, [Section 2.2.1](#)

### **10.3. Frame Types**

This document establishes a registry for HTTP/QUIC frame type codes. The "HTTP/QUIC Frame Type" registry manages an 8-bit space. The "HTTP/QUIC Frame Type" registry operates under either of the "IETF Review" or "IESG Approval" policies [[RFC8126](#)] for values from 0x00 up to and including 0xef, with values from 0xf0 up to and including 0xff being reserved for Experimental Use.

While this registry is separate from the "HTTP/2 Frame Type" registry defined in [[RFC7540](#)], it is preferable that the assignments parallel each other. If an entry is present in only one registry, every effort SHOULD be made to avoid assigning the corresponding value to an unrelated operation.

New entries in this registry require the following information:

Frame Type: A name or label for the frame type.

Code: The 8-bit code assigned to the frame type.

Specification: A reference to a specification that includes a description of the frame layout and its semantics, including any parts of the frame that are conditionally present.

The entries in the following table are registered by this document.



Frame Type	Code	Specification
DATA	0x0	<a href="#">Section 4.2.2</a>
HEADERS	0x1	<a href="#">Section 4.2.3</a>
PRIORITY	0x2	<a href="#">Section 4.2.4</a>
CANCEL_PUSH	0x3	<a href="#">Section 4.2.5</a>
SETTINGS	0x4	<a href="#">Section 4.2.6</a>
PUSH_PROMISE	0x5	<a href="#">Section 4.2.7</a>
Reserved	0x6	N/A
GOAWAY	0x7	<a href="#">Section 4.2.8</a>
Reserved	0x8	N/A
Reserved	0x9	N/A
MAX_PUSH_ID	0xD	<a href="#">Section 4.2.9</a>

Additionally, each code of the format "0xb + (0x1f \* N)" for values of N in the range (0..7) (that is, "0xb", "0x2a", "0x49", "0x68", "0x87", "0xa6", "0xc5", and "0xe4"), the following values should be registered:

Frame Type: Reserved - GREASE

Specification: [Section 4.2.1](#)

#### 10.4. Settings Parameters

This document establishes a registry for HTTP/QUIC settings. The "HTTP/QUIC Settings" registry manages a 16-bit space. The "HTTP/QUIC Settings" registry operates under the "Expert Review" policy [[RFC8126](#)] for values in the range from 0x0000 to 0xffff, with values between 0xf000 and 0xffff being reserved for Experimental Use. The designated experts are the same as those for the "HTTP/2 Settings" registry defined in [[RFC7540](#)].

While this registry is separate from the "HTTP/2 Settings" registry defined in [[RFC7540](#)], it is preferable that the assignments parallel each other. If an entry is present in only one registry, every



Bishop

Expires April 6, 2019

[Page 37]

effort SHOULD be made to avoid assigning the corresponding value to an unrelated operation.

New registrations are advised to provide the following information:

Name: A symbolic name for the setting. Specifying a setting name is optional.

Code: The 16-bit code assigned to the setting.

Specification: An optional reference to a specification that describes the use of the setting.

The entries in the following table are registered by this document.

Setting Name	Code	Specification
Reserved	0x2	N/A
NUM_PLACEHOLDERS	0x3	<a href="#">Section 4.2.6.1</a>
Reserved	0x4	N/A
Reserved	0x5	N/A
MAX_HEADER_LIST_SIZE	0x6	<a href="#">Section 4.2.6.1</a>

Additionally, each code of the format "0x?a?a" where each "?" is any four bits (that is, "0x0a0a", "0x0a1a", etc. through "0xfafa"), the following values should be registered:

Name: Reserved - GREASE

Specification: [Section 4.2.6.1](#)

## 10.5. Error Codes

This document establishes a registry for HTTP/QUIC error codes. The "HTTP/QUIC Error Code" registry manages a 16-bit space. The "HTTP/QUIC Error Code" registry operates under the "Expert Review" policy [[RFC8126](#)].

Registrations for error codes are required to include a description of the error code. An expert reviewer is advised to examine new registrations for possible duplication with existing error codes. Use of existing registrations is to be encouraged, but not mandated.



New registrations are advised to provide the following information:

Name: A name for the error code. Specifying an error code name is optional.

Code: The 16-bit error code value.

Description: A brief description of the error code semantics, longer if no detailed specification is provided.

Specification: An optional reference for a specification that defines the error code.

The entries in the following table are registered by this document.

Name	Code	Description	Specification
STOPPING	0x000 0	Reserved by QUIC	[QUIC-TRANSPORT ]
HTTP_NO_ERROR	0x000 1	No error	<a href="#">Section 6.1</a>
HTTP_PUSH_REFUSED	0x000 2	Client refused pushed content	<a href="#">Section 6.1</a>
HTTP_INTERNAL_ERROR	0x000 3	Internal error	<a href="#">Section 6.1</a>
HTTP_PUSH_ALREADY_IN_CACHE	0x000 4	Pushed content already cached	<a href="#">Section 6.1</a>
HTTP_REQUEST_CANCELLED	0x000 5	Data no longer needed	<a href="#">Section 6.1</a>
HTTP_INCOMPLETE_REQUEST	0x000 6	Stream terminated early	<a href="#">Section 6.1</a>
HTTP_CONNECT_ERROR	0x000 7	TCP reset or error on CONNECT request	<a href="#">Section 6.1</a>

Bishop

Expires April 6, 2019

[Page 39]

HTTP_EXCESSIVE_LOAD	0x0008	Peer generating excessive load	<a href="#">Section 6.1</a>
HTTP_VERSION_FALLBACK	0x0009	Retry over HTTP/1.1	<a href="#">Section 6.1</a>
HTTP_WRONG_STREAM	0x000A	A frame was sent on the wrong stream	<a href="#">Section 6.1</a>
HTTP_PUSH_LIMIT_EXCEEDED	0x000B	Maximum Push ID exceeded	<a href="#">Section 6.1</a>
HTTP_DUPLICATE_PUSH	0x000C	Push ID was fulfilled multiple times	<a href="#">Section 6.1</a>
HTTP_UNKNOWN_STREAM_TYPE	0x000D	Unknown unidirectional stream type	<a href="#">Section 6.1</a>
HTTP_WRONG_STREAM_COUNT	0x000E	Too many unidirectional streams	<a href="#">Section 6.1</a>
HTTP_CLOSED_CRITICAL_STREAM	0x000F	Critical stream was closed	<a href="#">Section 6.1</a>
HTTP_WRONG_STREAM_DIRECTION	0x0010	Unidirectional stream in wrong direction	<a href="#">Section 6.1</a>
HTTP_EARLY_RESPONSE	0x0011	Remainder of request not needed	<a href="#">Section 6.1</a>
HTTP_MISSING_SETTINGS	0x0012	No SETTINGS frame received	<a href="#">Section 6.1</a>
HTTP_MALFORMED_FRAME	0x01XX	Error in frame	<a href="#">Section 6.1</a>

Bishop

Expires April 6, 2019

[Page 40]

		formatting or	
		use	

### 10.6. Stream Types

This document establishes a registry for HTTP/QUIC unidirectional stream types. The "HTTP/QUIC Stream Type" registry manages an 8-bit space. The "HTTP/QUIC Stream Type" registry operates under either of the "IETF Review" or "IESG Approval" policies [[RFC8126](#)] for values from 0x00 up to and including 0xef, with values from 0xf0 up to and including 0xff being reserved for Experimental Use.

New entries in this registry require the following information:

Stream Type: A name or label for the stream type.

Code: The 8-bit code assigned to the stream type.

Specification: A reference to a specification that includes a description of the stream type, including the layout semantics of its payload.

Sender: Which endpoint on a connection may initiate a stream of this type. Values are "Client", "Server", or "Both".

The entries in the following table are registered by this document.

Stream Type	Code	Specification	Sender
Control Stream	0x43	<a href="#">Section 3.3.2</a>	Both
Push Stream	0x50	<a href="#">Section 3.3.3</a>	Server

Additionally, for each code of the format "0x1f \* N" for values of N in the range (0..8) (that is, "0x00", "0x1f", "0x3e", "0x5d", "0x7c", "0x9b", "0xba", "0xd9", "0xf8"), the following values should be registered:

Stream Type: Reserved - GREASE

Specification: [Section 3.3.1](#)

Sender: Both





## **11.** References

### **11.1.** Normative References

- [ALTSVC] Nottingham, M., McManus, P., and J. Reschke, "HTTP Alternative Services", [RFC 7838](#), DOI 10.17487/RFC7838, April 2016, <<https://www.rfc-editor.org/info/rfc7838>>.
- [QPACK] Krasic, C., Bishop, M., and A. Frindell, Ed., "QPACK: Header Compression for HTTP over QUIC", [draft-ietf-quic-qpack-03](#) (work in progress), October 2018.
- [QUIC-TRANSPORT] Iyengar, J., Ed. and M. Thomson, Ed., "QUIC: A UDP-Based Multiplexed and Secure Transport", [draft-ietf-quic-transport-14](#) (work in progress), October 2018.
- [RFC0793] Postel, J., "Transmission Control Protocol", STD 7, [RFC 793](#), DOI 10.17487/RFC0793, September 1981, <<https://www.rfc-editor.org/info/rfc793>>.
- [RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", [BCP 14](#), [RFC 2119](#), DOI 10.17487/RFC2119, March 1997, <<https://www.rfc-editor.org/info/rfc2119>>.
- [RFC5234] Crocker, D., Ed. and P. Overell, "Augmented BNF for Syntax Specifications: ABNF", STD 68, [RFC 5234](#), DOI 10.17487/RFC5234, January 2008, <<https://www.rfc-editor.org/info/rfc5234>>.
- [RFC6066] Eastlake 3rd, D., "Transport Layer Security (TLS) Extensions: Extension Definitions", [RFC 6066](#), DOI 10.17487/RFC6066, January 2011, <<https://www.rfc-editor.org/info/rfc6066>>.
- [RFC7230] Fielding, R., Ed. and J. Reschke, Ed., "Hypertext Transfer Protocol (HTTP/1.1): Message Syntax and Routing", [RFC 7230](#), DOI 10.17487/RFC7230, June 2014, <<https://www.rfc-editor.org/info/rfc7230>>.
- [RFC7231] Fielding, R., Ed. and J. Reschke, Ed., "Hypertext Transfer Protocol (HTTP/1.1): Semantics and Content", [RFC 7231](#), DOI 10.17487/RFC7231, June 2014, <<https://www.rfc-editor.org/info/rfc7231>>.



- [RFC7540] Belshe, M., Peon, R., and M. Thomson, Ed., "Hypertext Transfer Protocol Version 2 (HTTP/2)", [RFC 7540](#), DOI 10.17487/RFC7540, May 2015, <<https://www.rfc-editor.org/info/rfc7540>>.
- [RFC7838] Nottingham, M., McManus, P., and J. Reschke, "HTTP Alternative Services", [RFC 7838](#), DOI 10.17487/RFC7838, April 2016, <<https://www.rfc-editor.org/info/rfc7838>>.
- [RFC8174] Leiba, B., "Ambiguity of Uppercase vs Lowercase in [RFC 2119](#) Key Words", [BCP 14](#), [RFC 8174](#), DOI 10.17487/RFC8174, May 2017, <<https://www.rfc-editor.org/info/rfc8174>>.

### **[11.2.](#) Informative References**

- [RFC7301] Friedl, S., Popov, A., Langley, A., and E. Stephan, "Transport Layer Security (TLS) Application-Layer Protocol Negotiation Extension", [RFC 7301](#), DOI 10.17487/RFC7301, July 2014, <<https://www.rfc-editor.org/info/rfc7301>>.
- [RFC8126] Cotton, M., Leiba, B., and T. Narten, "Guidelines for Writing an IANA Considerations Section in RFCs", [BCP 26](#), [RFC 8126](#), DOI 10.17487/RFC8126, June 2017, <<https://www.rfc-editor.org/info/rfc8126>>.

### **[11.3.](#) URIs**

- [1] [https://mailarchive.ietf.org/arch/search/?email\\_list=quic](https://mailarchive.ietf.org/arch/search/?email_list=quic)
- [2] <https://github.com/quicwg>
- [3] <https://github.com/quicwg/base-drafts/labels/-http>
- [4] <https://www.iana.org/assignments/message-headers>

## **[Appendix A.](#) Change Log**

\*RFC Editor's Note:\* Please remove this section prior to publication of a final version of this document.

### **[A.1.](#) Since [draft-ietf-quic-http-14](#)**

- o Recommend sensible values for QUIC transport parameters (#1720,#1806)
- o Define error for missing SETTINGS frame (#1697,#1808)



- o Setting values are variable-length integers (#1556,#1807) and do not have separate maximum values (#1820)
- o Expanded discussion of connection closure (#1599,#1717,#1712)
- o HTTP\_VERSION\_FALLBACK falls back to HTTP/1.1 (#1677,#1685)

#### **A.2. Since [draft-ietf-quic-http-13](#)**

- o Reserved some frame types for grease (#1333, #1446)
- o Unknown unidirectional stream types are tolerated, not errors; some reserved for grease (#1490, #1525)
- o Require settings to be remembered for 0-RTT, prohibit reductions (#1541, #1641)
- o Specify behavior for truncated requests (#1596, #1643)

#### **A.3. Since [draft-ietf-quic-http-12](#)**

- o TLS SNI extension isn't mandatory if an alternative method is used (#1459, #1462, #1466)
- o Removed flags from HTTP/QUIC frames (#1388, #1398)
- o Reserved frame types and settings for use in preserving extensibility (#1333, #1446)
- o Added general error code (#1391, #1397)
- o Unidirectional streams carry a type byte and are extensible (#910,#1359)
- o Priority mechanism now uses explicit placeholders to enable persistent structure in the tree (#441,#1421,#1422)

#### **A.4. Since [draft-ietf-quic-http-11](#)**

- o Moved QPACK table updates and acknowledgments to dedicated streams (#1121, #1122, #1238)

#### **A.5. Since [draft-ietf-quic-http-10](#)**

- o Settings need to be remembered when attempting and accepting 0-RTT (#1157, #1207)



**A.6. Since [draft-ietf-quic-http-09](#)**

- o Selected QCRAM for header compression (#228, #1117)
- o The server\_name TLS extension is now mandatory (#296, #495)
- o Specified handling of unsupported versions in Alt-Svc (#1093, #1097)

**A.7. Since [draft-ietf-quic-http-08](#)**

- o Clarified connection coalescing rules (#940, #1024)

**A.8. Since [draft-ietf-quic-http-07](#)**

- o Changes for integer encodings in QUIC (#595, #905)
- o Use unidirectional streams as appropriate (#515, #240, #281, #886)
- o Improvement to the description of GOAWAY (#604, #898)
- o Improve description of server push usage (#947, #950, #957)

**A.9. Since [draft-ietf-quic-http-06](#)**

- o Track changes in QUIC error code usage (#485)

**A.10. Since [draft-ietf-quic-http-05](#)**

- o Made push ID sequential, add MAX\_PUSH\_ID, remove SETTINGS\_ENABLE\_PUSH (#709)
- o Guidance about keep-alive and QUIC PINGs (#729)
- o Expanded text on GOAWAY and cancellation (#757)

**A.11. Since [draft-ietf-quic-http-04](#)**

- o Cite [RFC 5234](#) (#404)
- o Return to a single stream per request (#245, #557)
- o Use separate frame type and settings registries from HTTP/2 (#81)
- o SETTINGS\_ENABLE\_PUSH instead of SETTINGS\_DISABLE\_PUSH (#477)
- o Restored GOAWAY (#696)





- o Identify server push using Push ID rather than a stream ID (#702,#281)
- o DATA frames cannot be empty (#700)

**A.12.** Since [draft-ietf-quic-http-03](#)

None.

**A.13.** Since [draft-ietf-quic-http-02](#)

- o Track changes in transport draft

**A.14.** Since [draft-ietf-quic-http-01](#)

- o SETTINGS changes (#181):
  - \* SETTINGS can be sent only once at the start of a connection; no changes thereafter
  - \* SETTINGS\_ACK removed
  - \* Settings can only occur in the SETTINGS frame a single time
  - \* Boolean format updated
- o Alt-Svc parameter changed from "v" to "quic"; format updated (#229)
- o Closing the connection control stream or any message control stream is a fatal error (#176)
- o HPACK Sequence counter can wrap (#173)
- o 0-RTT guidance added
- o Guide to differences from HTTP/2 and porting HTTP/2 extensions added (#127,#242)

**A.15.** Since [draft-ietf-quic-http-00](#)

- o Changed "HTTP/2-over-QUIC" to "HTTP/QUIC" throughout (#11,#29)
- o Changed from using HTTP/2 framing within Stream 3 to new framing format and two-stream-per-request model (#71,#72,#73)
- o Adopted SETTINGS format from [draft-bishop-httpbis-extended-settings-01](#)



- o Reworked SETTINGS\_ACK to account for indeterminate inter-stream order (#75)
- o Described CONNECT pseudo-method (#95)
- o Updated ALPN token and Alt-Svc guidance (#13,#87)
- o Application-layer-defined error codes (#19,#74)

**A.16. Since [draft-shade-quic-http2-mapping-00](#)**

- o Adopted as base for [draft-ietf-quic-http](#)
- o Updated authors/editors list

#### Acknowledgements

The original authors of this specification were Robbie Shade and Mike Warres.

A substantial portion of Mike's contribution was supported by Microsoft during his employment there.

#### Author's Address

Mike Bishop (editor)  
Akamai

Email: [mbishop@evequefou.be](mailto:mbishop@evequefou.be)

