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Modeling Video Traffic Sources for RMCAT Evaluations
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Abstract

This document describes two reference video traffic source models for evaluating RMCAT candidate algorithms. The first model statistically characterizes the behavior of a live video encoder in response to changing requests on target video rate. The second model is trace-driven, and emulates the encoder output by scaling the pre-encoded video frame sizes from a widely used video test sequence. Both models are designed to strike a balance between simplicity, repeatability, and authenticity in modeling the interactions between a video traffic source and the congestion control module.

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Table of Contents

1.	Introduction	2
2.	Terminology	3
3.	Desired Behavior of A Synthetic Video Traffic Model	3
4.	Interactions Between Synthetic Video Traffic Source and Other Components at the Sender	4
5.	A Statistical Reference Model	6
5.1.	Time-damped response to target rate update	7
5.2.	Temporary burst/oscillation during transient	7
5.3.	Output rate fluctuation at steady state	8
5.4.	Rate range limit imposed by video content	8
6.	A Trace-Driven Model	8
6.1.	Choosing the video sequence and generating the traces	9
6.2.	Using the traces in the syntethic codec	10
6.2.1.	Main algorithm	10
6.2.2.	Notes to the main algorithm	12
6.3.	Varying frame rate and resolution	12
7.	Comparing and Combining The Two Models	13
8.	Implementation Status	14
9.	IANA Considerations	14
10.	References	14
10.1.	Normative References	14
10.2.	Informative References	14
	Authors' Addresses	15

[1.](#) Introduction

When evaluating candidate congestion control algorithms designed for real-time interactive media, it is important to account for the characteristics of traffic patterns generated from a live video encoder. Unlike synthetic traffic sources that can conform perfectly to the rate changing requests from the congestion control module, a live video encoder can be sluggish in reacting to such changes. Output rate of a live video encoder also typically deviates from the target rate due to uncertainties in the encoder rate control process. Consequently, end-to-end delay and loss performance of a real-time media flow can be further impacted by rate variations introduced by the live encoder.

On the other hand, evaluation results of a candidate RMCAT algorithm should mostly reflect performance of the congestion control module, and somewhat decouple from peculiarities of any specific video codec. It is also desirable that evaluation tests are repeatable, and be easily duplicated across different candidate algorithms.

One way to strike a balance between the above considerations is to evaluate RMCAT algorithms using a synthetic video traffic source model that captures key characteristics of the behavior of a live video encoder. To this end, this draft presents two reference models. The first is based on statistical modelling; the second is trace-driven. The draft also discusses the pros and cons of each approach, as well as the possibility to combine both.

2. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described [RFC2119](#) [[RFC2119](#)].

3. Desired Behavior of A Synthetic Video Traffic Model

A live video encoder employs encoder rate control to meet a target rate by varying its encoding parameters, such as quantization step size, frame rate, and picture resolution, based on its estimate of the video content (e.g., motion and scene complexity). In practice, however, several factors prevent the output video rate from perfectly conforming to the input target rate.

Due to uncertainties in the captured video scene, the output rate typically deviates from the specified target. In the presence of a significant change in target rate, it sometimes takes several frames before the encoder output rate converges to the new target. Finally, while most of the frames in a live session are encoded in predictive mode, the encoder can occasionally generate a large intra-coded frame (or a frame partially containing intra-coded blocks) in an attempt to recover from losses, to re-sync with the receiver, or during the transient period of responding to target rate or spatial resolution changes.

Hence, a synthetic video source should have the following capabilities:

- o To change bitrate. This includes ability to change framerate and/or spatial resolution, or to skip frames when required.
- o To fluctuate around the target bitrate specified by the congestion control module.

- o To delay in convergence to the target bitrate.
- o To generate intra-coded or repair frames on demand.

While there exists many different approaches in developing a synthetic video traffic model, it is desirable that the outcome follows a few common characteristics, as outlined below.

- o Low computational complexity: The model should be computationally lightweight, otherwise it defeats the whole purpose of serving as a substitute for a live video encoder.
- o Temporal pattern similarity: The individual traffic trace instances generated by the model should mimic the temporal pattern of those from a real video encoder.
- o Statistical resemblance: The synthetic traffic should match the outcome of the real video encoder in terms of statistical characteristics, such as the mean, variance, peak, and autocorrelation coefficients of the bitrate. It is also important that the statistical resemblance should hold across different time scales, ranging from tens of milliseconds to sub-seconds.
- o Wide range of coverage: The model should be easily configurable to cover a wide range of codec behaviors (e.g., with either fast or slow reaction time in live encoder rate control) and video content variations (e.g, ranging from high-motion to low-motion).

These distinct behavior features can be characterized via simple statistical models, or a trace-driven approach. We present an example of each in [Section 5](#) and [Section 6](#)

4. Interactions Between Synthetic Video Traffic Source and Other Components at the Sender

Figure 1 depicts the interactions of the synthetic video encoder with other components at the sender, such as the application, the congestion control module, the media packet transport module, etc. Both reference models, as described later in [Section 5](#) and [Section 6](#), follow the same set of interactions.

The synthetic video encoder takes in raw video frames captured by the camera and then dynamically generates a sequence of encoded video frames with varying size and interval. These encoded frames are processed by other modules in order to transmit the video stream over the network. During the lifetime of a video transmission session, the synthetic video encoder will typically be required to adapt its

encoding bitrate, and sometimes the spatial resolution and frame rate.

In our model, the synthetic video encoder module has group of incoming and outgoing interface calls that allow for interaction with other modules. The following are some of the possible incoming interface calls --- marked as (a) in Figure 1 --- that the synthetic video encoder may accept. The list is not exhaustive and can be complemented by other interface calls if deemed necessary.

- o Target rate $R_v(t)$: requested at time t , typically from the congestion control module. Depending on the congestion control algorithm in use, the update requests can either be periodic (e.g., once per second), or on-demand (e.g., only when a drastic bandwidth change over the network is observed).
- o Target frame rate $FPS(t)$: the instantaneous frame rate measured in frames-per-second at time t . This depends on the native camera capture frame rate as well as the target/preferred frame rate configured by the application or user.
- o Frame resolution $XY(t)$: the 2-dimensional vector indicating the preferred frame resolution in pixels at time t . Several factors govern the resolution requested to the synthetic video encoder over time. Examples of such factors are the capturing resolution of the native camera; or the current target rate $R_v(t)$, since very small resolutions do not make sense with very high bitrates, and vice-versa.
- o Instant frame skipping: the request to skip the encoding of one or several captured video frames, for instance when a drastic decrease in available network bandwidth is detected.
- o On-demand generation of intra (I) frame: the request to encode another I frame to avoid further error propagation at the receiver, if severe packet losses are observed. This request typically comes from the error control module.

An example of outgoing interface call --- marked as (b) in Figure 1 --- is the rate range, that is, the dynamic range of the video encoder's output rate for the current video contents: $[R_{min}, R_{max}]$. Here, R_{min} and R_{max} are meant to capture the dynamic rate range the encoder is capable of outputting. This typically depends on the

video content complexity and/or display type (e.g., higher R_{\max} for video contents with higher motion complexity, or for displays of higher resolution). Therefore, these values will not change with R_v , but may change over time if the content is changing.

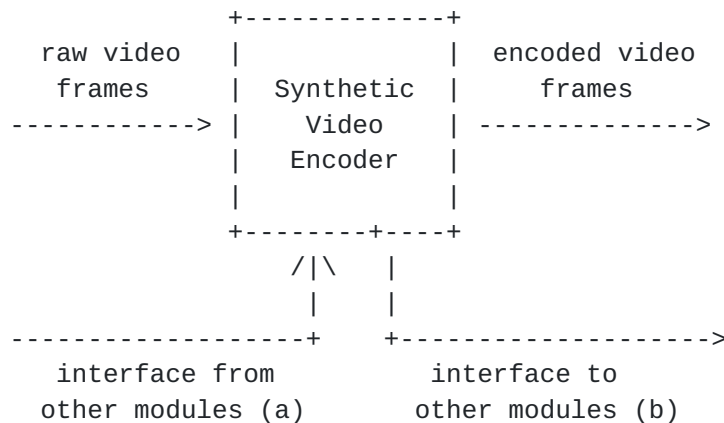


Figure 1: Interaction between synthetic video encoder and other modules at the sender

5. A Statistical Reference Model

In this section, we describe one simple statistical model of the live video encoder traffic source. Figure 2 summarizes the list of tunable parameters in this statistical model. A more comprehensive survey of popular methods for modelling video traffic source behavior can be found in [\[Tanwir2013\]](#).

Notation	Parameter Name	Example Value
$R_v(t)$	Target rate request at time t	1 Mbps
$R_o(t)$	Output rate at time t	1.2 Mbps
τ_v	Encoder reaction latency	0.2 s
K_d	Burst duration during transient	5 frames
K_r	Burst size during transient	5:1
$R_e(t)$	Error in output rate at time t	0.2 Mbps
SIGMA	standard deviation of normally distributed relative rate error	0.1
DELTA	upper and lower bound (+/-) of uniformly distributed relative rate error	0.1
R_{\min}	minimum rate supported by video encoder or content activity	150 Kbps
R_{\max}	maximum rate supported by video encoder or content activity	1.5Mbps

Figure 2: List of tunable parameters in a statistical video traffic source model.

5.1. Time-damped response to target rate update

While the congestion control module can update its target rate request $R_v(t)$ at any time, our model dictates that the encoder will only react to such changes after τ_v seconds from a previous rate transition. In other words, when the encoder has reacted to a rate change request at time t , it will simply ignore all subsequent rate change requests until time $t+\tau_v$.

5.2. Temporary burst/oscillation during transient

The output rate R_o during the period $[t, t+\tau_v]$ is considered to be in transient. Based on observations from video encoder output data, we model the transient behavior of an encoder upon reacting to a new target rate request in the form of largely varying output sizes. It is assumed that the overall average output rate R_o during this period matches the target rate R_v . Consequently, the occasional burst of large frames are followed by smaller-than average encoded frames.

This temporary burst is characterized by two parameters:

- o burst duration K_d : number frames in the burst event; and

- o burst size K_r : ratio of a burst frame and average frame size at steady state.

It can be noted that these burst parameters can also be used to mimic the inserion of a large on-demand I frame in the presence of severe packet losses. The values of K_d and K_r are fitted to reflect the typical ratio between I and P frames for a given video content.

5.3. Output rate fluctuation at steady state

We model output rate R_o as randomly fluctuating around the target rate R_v after convergence. There are two variants in modeling the random fluctuation $R_e = R_o - R_v$:

- o As normal distribution: with a mean of zero and a standard deviation $SIGMA$ specified in terms of percentage of the target rate. A typical value of $SIGMA$ is 10 percent of target rate.
- o As uniform distribution bounded between $-DELTA$ and $DELTA$. A typical value of $DELTA$ is 10 percent of target rate.

The distribution type (normal or uniform) and model parameters ($SIGMA$ or $DELTA$) can be learned from data samples gathered from a live encoder output.

5.4. Rate range limit imposed by video content

The output rate R_o is further clipped within the dynamic range $[R_{min}, R_{max}]$, which in reality are dictated by scene and motion complexity of the captured video content. In our model, these parameters are specified by the application.

6. A Trace-Driven Model

We now present the second approach to model a video traffic source. This approach is based on running an actual live video encoder offline on a set of chosen raw video sequences and using the encoder's output traces for constructing a synthetic live encoder. With this approach, the recorded video traces naturally exhibit temporal fluctuations around a given target rate request $R_v(t)$ from the congestion control module.

The following list summarizes this approach's main steps:

- 1) Choose one or more representative raw video sequences.

- 2) Using an actual live video encoder, encode the sequences at various bitrates. Keep just the sequences of frame sizes for each bitrate.
- 3) Construct a data structure that contains the output of the previous step. The data structure should allow for easy bitrate lookup.
- 4) Upon a target bitrate request $R_v(t)$ from the controller, look up the closest bitrates among those previously stored. Use the frame size sequences stored for those bitrates to approximate the frame sizes to output.
- 5) The output of the synthetic encoder contains "encoded" frames with zeros as contents but with realistic sizes.

[Section 6.1](#) explains steps 1), 2), and 3), [Section 6.2](#) elaborates on steps 4) and 5). Finally, [Section 6.3](#) briefly discusses the possibility to extend the model for supporting variable frame rate and/or variable frame resolution.

[6.1](#). Choosing the video sequence and generating the traces

The first step we need to perform is a careful choice of a set of video sequences that are representative of the use cases we want to model. Our use case here is video conferencing, so we must choose a low-motion sequence that resembles a "talking head", for instance a news broadcast or a video capture of an actual conference call.

The length of the chosen video sequence is a tradeoff. If it is too long, it will be difficult to manage the data structures containing the traces we will produce in the next steps. If it is too short, there will be an obvious periodic pattern in the output frame sizes, leading to biased results when evaluating congestion controller performance. In our experience, a one-minute-long sequence is a fair tradeoff.

Once we have chosen the raw video sequence, denoted S , we use a live encoder, e.g. [\[H264\]](#) or [\[HEVC\]](#) to produce a set of encoded sequences. As discussed in [Section 3](#), a live encoder's output bitrate can be tuned by varying three input parameters, namely, quantization step size, frame rate, and picture resolution. In order to simplify the choice of these parameters for a given target rate, we assume a fixed frame rate (e.g. 25 fps) and a fixed resolution (e.g., 480p). See [section 6.3](#) for a discussion on how to relax these assumptions.

Following these simplifications, we run the chosen encoder by setting a constant target bitrate at the beginning, then letting the encoder vary the quantization step size internally while encoding the input video sequence. Besides, we assume that the first frame is encoded as an I-frame and the rest are P-frames. We further assume that the encoder algorithm does not use knowledge of frames in the future so as to encode a given frame.

We define R_{\min} and R_{\max} as the minimum and maximum bitrate at which the synthetic codec is to operate. We divide the bitrate range between R_{\min} and R_{\max} in $n_s + 1$ bitrate steps of length $l = (R_{\max} - R_{\min}) / n_s$. We then use the following simple algorithm to encode the raw video sequence.

```

r = R_min
while r <= R_max do
    Traces[r] = encode_sequence(S, r, e)
    r = r + l

```

where function `encode_sequence` takes as parameters, respectively, a raw video sequence, a constant target rate, and an encoder algorithm; it returns a vector with the sizes of frames in the order they were encoded. The output vector is stored in a map structure called `Traces`, whose keys are bitrates and values are frame size vectors.

The choice of a value for n_s is important, as it determines the number of frame size vectors stored in map `Traces`. The minimum value one can choose for n_s is 1, and its maximum value depends on the amount of memory available for holding the map `Traces`. A reasonable value for n_s is one that makes the steps' length $l = 200$ kbps. We will further discuss step length l in the next section.

6.2. Using the traces in the syntethic codec

The main idea behind the trace-based synthetic codec is that it mimics a real live codec's rate adaptation when the congestion controller updates the target rate $R_v(t)$. It does so by switching to a different frame size vector stored in the map `Traces` when needed.

6.2.1. Main algorithm

We maintain two variables r_{current} and t_{current} :

* r_{current} points to one of the keys of the map `Traces`. Upon a change in the value of $R_v(t)$, typically because the congestion controller detects that the network conditions have changed, r_{current} is updated to the greatest key in `Traces` that is less than

or equal to the new value of $R_v(t)$. For the moment, we assume the value of $R_v(t)$ to be clipped in the range $[R_{\min}, R_{\max}]$.

```

r_current = r
such that
  ( r in keys(Traces) and
    r <= R_v(t) and
    (not(exists) r' in keys(Traces) such that r < r' <= R_v(t)) )

```

* t_{current} is an index to the frame size vector stored in $\text{Traces}[r_{\text{current}}]$. It is updated every time a new frame is due. We assume all vectors stored in Traces to have the same size, denoted size_traces . The following equation governs the update of t_{current} :

```

if t_current < SkipFrames then
  t_current = t_current + 1
else
  t_current = ((t_current+1-SkipFrames) % (size_traces-
SkipFrames))
               + SkipFrames

```

where operator % denotes modulo, and SkipFrames is a predefined constant that denotes the number of frames to be skipped at the beginning of frame size vectors after t_{current} has wrapped around. The point of constant SkipFrames is avoiding the effect of periodically sending a (big) I-frame followed by several smaller-than-normal P-frames. We typically set SkipFrames to 20, although it could be set to 0 if we are interested in studying the effect of sending I-frames periodically.

We initialize r_{current} to R_{\min} , and t_{current} to 0.

When a new frame is due, we need to calculate its size. There are three cases:

- a) $R_{\min} \leq R_v(t) < R_{\max}$: In this case we use linear interpolation of the frame sizes appearing in $\text{Traces}[r_{\text{current}}]$ and $\text{Traces}[r_{\text{current}} + 1]$. The interpolation is done as follows:

```

size_lo = Traces[r_current][t_current]
size_hi = Traces[r_current + 1][t_current]
distance_lo = ( R_v(t) - r_current ) / 1
framesize = size_hi * distance_lo + size_lo * (1 - distance_lo)

```

- b) $R_v(t) < R_{\min}$: In this case, we scale the trace sequence with the lowest bitrate, in the following way:


```

factor = R_v(t) / R_min
framesize = max(1, factor * Traces[R_min][t_current])

```

c) $R_v(t) \geq R_{\max}$: We also use scaling for this case. We use the trace sequence with the greatest bitrate:

```

factor = R_v(t) / R_max
framesize = factor * Traces[R_max][t_current]

```

In case b), we set the minimum to 1 byte, since the value of factor can be arbitrarily close to 0.

6.2.2. Notes to the main algorithm

* Reacting to changes in target bitrate. Similarly to the statistical model presented in [Section 5](#), the trace-based synthetic codec can have a time bound, τ_v , to reacting to target bitrate changes. If the codec has reacted to an update in $R_v(t)$ at time t , it will delay any further update to $R_v(t)$ to time $t + \tau_v$. Note that, in any case, the value of τ_v cannot be chosen shorter than the time between frames, i.e. the inverse of the frame rate.

* I-frames on demand. The synthetic codec could be extended to simulate the sending of I-frames on demand, e.g., as a reaction to losses. To implement this extension, the codec's API is augmented with a new function to request a new I-frame. Upon calling such function, t_{current} is reset to 0.

* Variable length l of steps defined between R_{\min} and R_{\max} . In the main algorithm's description, the step length l is fixed. However, if the range $[R_{\min}, R_{\max}]$ is very wide, it is also possible to define a set of steps with a non-constant length. The idea behind this modification is that the difference between 400 kbps and 600 kbps as bitrate is much more important than the difference between 4400 kbps and 4600 kbps. For example, one could define steps of length 200 Kbps under 1 Mbps, then length 300 kbps between 1 Mbps and 2 Mbps, 400 kbps between 2 Mbps and 3 Mbps, and so on.

6.3. Varying frame rate and resolution

The trace-based synthetic codec model explained in this section is relatively simple because we have fixed the frame rate and the frame resolution. The model could be extended to have variable frame rate, variable spatial resolution, or both.

When the encoded picture quality at a given bitrate is low, one can potentially decrease the frame rate (if the video sequence is currently in low motion) or the spatial resolution in order to

improve quality-of-experience (QoE) in the overall encoded video. On the other hand, if target bitrate increases to a point where there is no longer a perceptible improvement in the picture quality of individual frames, then one might afford to increase the spatial resolution or the frame rate (useful if the video is currently in high motion).

Many techniques have been proposed to choose over time the best combination of encoder quantization step size, frame rate, and spatial resolution in order to maximize the quality of live video codecs [[Ozer2011](#)][[Hu2010](#)]. Future work may consider extending the trace-based codec to accommodate variable frame rate and/or resolution.

From the perspective of congestion control, varying the spatial resolution typically requires a new intra-coded frame to be generated, thereby incurring a temporary burst in the output traffic pattern. The impact of frame rate change tends to be more subtle: reducing frame rate from high to low leads to sparsely spaced larger encoded packets instead of many densely spaced smaller packets. Such difference in traffic profiles may still affect the performance of congestion control, especially when outgoing packets are not paced at the transport module. We leave the investigation of varying frame rate to future work.

7. Comparing and Combining The Two Models

It is worthwhile noting that the statistical and trace-based models each has its own advantages and drawbacks. Both models are fairly simple to implement. However, it takes significantly more effort to fit the parameters of a statistical model to actual encoder output data whereas a trace-based model does not require such fitting. On the other hand, once validated, the statistical model is more flexible in mimicking a wide range of encoder/content behavior by simply varying the corresponding parameters in the model. In contrast, a trace-driven model relies, by definition, on additional data collection efforts for accommodating new codecs or video contents.

In general, trace-based model is more realistic for mimicking ongoing, steady-state behavior of a video traffic source whereas statistical model is more versatile for simulating transient events (e.g., when target rate changes from A to B with temporary bursts during the transition). Therefore, it may be desirable to combine both approaches into a hybrid model, using traces for steady-state and statistical model for transients.

8. Implementation Status

The statistical model has been implemented as a traffic generator module within the [[ns-2](#)] network simulation platform.

More recently, both the statistical and trace-driven models have been implemented as a stand-alone traffic source module. This can be easily integrated into network simulation platforms such as [[ns-2](#)] and [[ns-3](#)], as well as testbeds using a real network. The stand-alone traffic source module is available as an open source implementation at [[Syncodecs](#)].

9. IANA Considerations

There are no IANA impacts in this memo.

10. References

10.1. Normative References

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