Routing Area Working Group Internet-Draft

Intended status: Standards Track

Expires: June 16, 2016

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# Link State protocols SPF trigger and delay algorithm impact on IGP micro-loops draft-ietf-rtgwg-spf-uloop-pb-statement-02

### Abstract

A micro-loop is a packet forwarding loop that may occur transiently among two or more routers in a hop-by-hop packet forwarding paradigm.

In this document, we are trying to analyze the impact of using different Link State IGP implementations in a single network in regards of micro-loops. The analysis is focused on the SPF triggers and SPF delay algorithm.

### Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

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### 1. Introduction

Link State IGP protocols are based on a topology database on which a SPF (Shortest Path First) algorithm like Dijkstra is implemented to find the optimal routing paths.

Specifications like IS-IS ([RFC1195]) propose some optimization of the route computation (See Appendix C.1) but not all the implementations are following those not mandatory optimizations.

We will call SPF trigger, the events that would lead to a new SPF computation based on the topology.

Link State IGP protocols, like OSPF ([RFC2328]) and IS-IS ([RFC1195]), are using plenty of timers to control the router

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behavior in case of churn : SPF delay, PRC delay, LSP generation delay, LSP flooding delay, LSP retransmission interval ...

Some of those timers are standardized in protocol specification, some are not especially the SPF computation related timers.

For non standardized timers, implementations are free to implement it in any way. For some standardized timer, we can also see that rather than using static configurable values for such timer, implementations may offer dynamically adjusted timers to help controlling the churn.

We will call SPF delay, the timer that exists in most implementations that specifies the required delay before running SPF computation after a SPF trigger is received.

A micro-loop is a packet forwarding loop that may occur transiently among two or more routers in a hop-by-hop packet forwarding paradigm. We can observe that these micro-loops are formed when two routers do not update their Forwarding Information Base (FIB) for a certain prefix at the same time. The micro-loop phenomenon is described in [I-D.ietf-rtgwg-microloop-analysis].

Some micro-loop mitigation techniques have been defined by IETF (e.g. [RFC6976], [I-D.ietf-rtgwg-uloop-delay]) but are not implemented due to complexity or are not providing a complete mitigation.

In multi vendor networks, using different implementations of a link state protocol may favor micro-loops creation during convergence time due to discrepancies of timers. Service Providers are already aware to use similar timers for all the network as best practice, but sometimes it is not possible due to limitation of implementations.

This document will present why it sounds important for service provider to have consistent implementations of Link State protocols across vendors. We are particularly analyzing the impact of using different Link State IGP implementations in a single network in regards of micro-loops. The analysis is focused on the SPF triggers and SPF delay algorithm in a first step.

This document is only stating the problem, and defining some work items but its not intended to provide a solution.

#### 2. Problem statement

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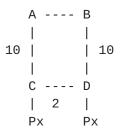


Figure 1

In the figure above, A uses primarily the AC link to reach C. When the AC link fails, IGP convergence occurs. If A converges before B, A will forward traffic to C through B, but as B as not converged yet, B will loop back traffic to A, leading to a micro-loop.

The micro-loop appears due to the asynchronous convergence of nodes in a network when a event occurs.

Multiple factors (and combination of these factors) may increase the probability for a micro-loop to appear :

- o delay of failure notification : the more B is advised of the failure later than A, the more a micro-loop may appear.
- o SPF delay: most of the implementations supports a delay for the SPF computation to try to catch as many events as possible. If A uses a SPF delay timer of x msec and B uses a SPF delay timer of y msec and x < y, B would start converging after A leading to a potential micro-loop.
- o SPF computation time: mostly a matter of CPU power and optimizations like incremental SPF. If A computes SPF faster than B, there is a chance for a micro-loop to appear. CPUs are today faster enough to consider SPF computation time as negligeable (order of msec in a large network).
- o RIB and FIB prefix insertion speed or ordering : highly implementation dependant.

This document will focus on analysis SPF delay (and associated triggers).

# 3. SPF trigger strategies

Depending of the change advertised in LSP/LSA, the topology may be affected or not. An implementation can decide to not run SPF (and only run IP reachability) if the advertised change is not affecting topology.

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Different strategies exists to trigger SPF:

- 1. Always run full SPF whatever the change to process.
- 2. Run only Full SPF when required: e.g. if a link fails, a local node will run an SPF for its local LSP update. If the LSP from the neighbor (describing the same failure) is received after SPF has started, the local node can decide that a new full SPF is not required as the topology has not change.
- 3. If topology does not change, only recompute reachability.

As pointed in <u>Section 1</u>, SPF optimization are not mandatory in specifications, leading to multiple strategies to be implemented.

## 4. SPF delay strategies

Implementations of link state routing protocols use different strategies to delay SPF :

- 1. Two steps.
- 2. Exponential backoff.

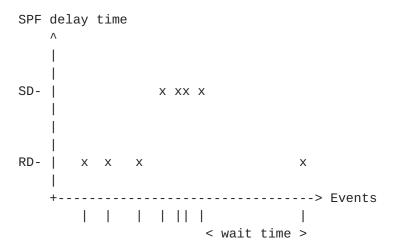
## 4.1. Two step SPF delay

The SPF delay is managed by four parameters :

- o Rapid delay : amount of time to wait before running SPF.
- o Rapid runs: amount of consecutive SPF runs that can run using rapid delay. When amount is exceeded router moves to slow delay.
- o Slow delay: amount of time to wait before running SPF.
- o Wait time: amount of time to wait without events before going back to rapid delay.

Example : Rapid delay = 50msec, Rapid runs = 3, Slow delay = 1sec, Wait time = 2sec

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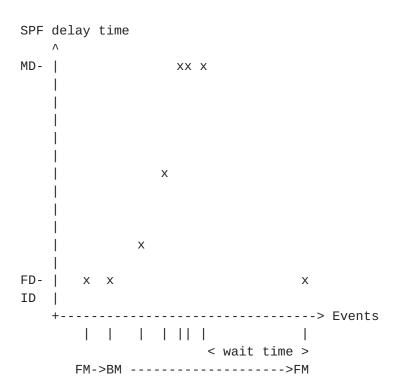


# 4.2. Exponential backoff

The algorithm has two mode: fast mode and backoff mode. In backoff mode, the SPF delay is increasing exponentially at each run. The SPF delay is managed by four parameters:

- o First delay: amount of time to wait before running SPF. This delay is used only when SPF is in fast mode.
- o Incremental delay: amount of time to wait before running SPF. This delay is used only when SPF is in backoff mode and increments exponentially at each SPF run.
- o Maximum delay : maximum amount of time to wait before running SPF.
- o Wait time: amount of time to wait without events before going back to fast mode.

Example : First delay = 50msec, Incremental delay = 50msec, Maximum delay = 1sec, Wait time = 2sec



# **5**. Mixing strategies

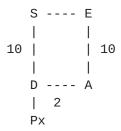


Figure 2

In the diagram above, we consider a flow of packet from S to D. We consider that S is using optimized SPF triggering (Full SPF is triggered only when necessary), and two steps SPF delay (rapid=150ms,rapid-runs=3, slow=1s). As implementation of S is optimized, Partial Reachability Computation (PRC) is available. We consider the same timers as SPF for delaying PRC. We consider that E is using a SPF trigger strategy that always compute Full SPF and exponential backoff strategy for SPF delay (start=150ms, inc=150ms, max=1s)

We also consider the following sequence of events (note: the time scale does not intend to represent a real router time scale where jitters are introduced to all timers):

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- o t0=0 ms : a prefix is declared down in the network. We consider this event to happen at time=0.
- o 200ms: the prefix is declared as up.
- o 400ms: a prefix is declared down in the network.
- o 1000ms : S-D link fails.

++		-+	+
Time	Network Event	Router S events	Router E events
t0=0	Prefix DOWN		+ 
10ms		Schedule PRC (in	Schedule SPF (in
İ		150ms)	150ms)
160ms		PRC starts	I   SPF starts
161ms		PRC ends	
162ms		RIB/FIB starts	
163ms			SPF ends
164ms			RIB/FIB starts
175ms		RIB/FIB ends	
178ms			RIB/FIB ends
	Prefix UP		 
212ms		Schedule PRC (in	
i i		150ms)	
214ms			Schedule SPF (in
1			150ms)
		   PRC starts	 
372ms		PRC ends	! 
373ms		i	SPF starts
373ms		RIB/FIB starts	
375ms			SPF ends
376ms			RIB/FIB starts
383ms		RIB/FIB ends	
385ms			RIB/FIB ends
	Prefix DOWN	 	 
410ms		Schedule PRC (in	Schedule SPF (in
		300ms)	300ms)
i i		1	
l İ		1	
		1	
1			

	710ms   711ms   712ms		PRC starts   PRC ends   RIB/FIB starts	SPF starts       
İ	713ms			SPF ends
	714ms			RIB/FIB starts
	716ms		RIB/FIB ends 	RIB/FIB ends   
	1000ms	S-D link DOWN		
	1010ms		Schedule SPF (in	Schedule SPF (in
			150ms)	600ms)
	1160ms		SPF starts	
	1161ms		SPF ends	
	1162ms	Micro-loop may	RIB/FIB starts	
		start from here		
	1175ms		RIB/FIB ends	
!				
-				
-	1610mc		 	CDE otorto
1	1612ms   1615ms		 	SPF starts     SPF ends
1	1616ms		 	SPF ends     RIB/FIB starts
1	1626ms	Micro-loop ends	I 	RIB/FIB starts
  -		•	 <del> </del>	L KID/IID CHGS

# Route computation event time scale

In the table above, we can see that due to discrepancies in SPF management, after multiple events (different types of event), SPF delays are completely misaligned between nodes leading to long microloop creation.

The same issue can also appear with only single type of events as displayed below :

+ -		-+-		.+	++
	Time			Router S events	'
İ	t0=0	İ	Link DOWN	Ī	l I
	10ms			Schedule SPF (in	Schedule SPF (in
				150ms)	150ms)
	160ms			SPF starts	SPF starts
	161ms			SPF ends	I I
	162ms			RIB/FIB starts	

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163ms     164ms			SPF ends   RIB/FIB starts
175ms     178ms   		RIB/FIB ends     	   RIB/FIB ends   
200ms     212ms           214ms   	Link DOWN	   Schedule SPF (in     150ms)   	 
370ms   372ms   373ms   373ms   375ms   376ms   383ms   385ms		SPF starts SPF ends RIB/FIB starts RIB/FIB ends RIB/FIB ends	SPF starts   SPF ends   RIB/FIB starts   RIB/FIB ends
400ms     410ms   	Link DOWN	   Schedule SPF (in     150ms)	   Schedule SPF (in     300ms)   
560ms     561ms     562ms	Micro-loop may start from here	   SPF starts     SPF ends     RIB/FIB starts       RIB/FIB ends	
710ms   713ms   714ms   716ms     1000ms   1010ms	Micro-loop ends Link DOWN	           Schedule SPF (in     1s)	SPF starts   SPF ends   RIB/FIB starts   RIB/FIB ends     Schedule SPF (in   600ms)
	Micro-loop may	       	   SPF starts     SPF ends     RIB/FIB starts

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1	start from here			
1626ms				RIB/FIB ends
2012ms		SPF starts		
2014ms		SPF ends		
2015ms		RIB/FIB starts		
2025ms	Micro-loop ends	RIB/FIB ends		
+		-+	-+-	

Route computation event time scale

### 6. Proposed work items

In order to enhance the current LinkState IGP behaviour, authors would encourage working on standardization of some behaviours.

Authors are proposing the following work items :

- o Standardize SPF trigger strategy.
- o Standardize computation timer scope : single timer for all computation operations, separated timers ...
- o Standardize "slowdown" timer algorithm including its association to a particular timer: authors of this document does not presume that the same algorithm must be used for all timers.

Using the same event sequence as in figure 2, we may expect fewer and/or shorter micro-loops using standardized implementations.

+	+		+	++
Time	 	Network Event	Router S events	
t0=0	   	Prefix DOWN	l	
	į		150ms)	150ms)
160ms	İ		PRC starts	PRC starts
161ms			PRC ends	I I
162ms			RIB/FIB starts	PRC ends
163ms				RIB/FIB starts
175ms	- 1		RIB/FIB ends	

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1	.76ms			RIB/FIB ends
	00ms	Prefix UP	 	
i ¯	212ms      213ms   		Schedule PRC (in     150ms)   	   Schedule PRC (in     150ms)
3   3   3   3	370ms     372ms     373ms     374ms     383ms     884ms		   PRC starts     PRC ends   RIB/FIB starts     RIB/FIB ends   	PRC starts   PRC starts   PRC ends   RIB/FIB starts   RIB/FIB ends
•	00ms   10ms         	Prefix DOWN     	   Schedule PRC (in     300ms)	
7   7   7	/10ms   /11ms   /12ms   /13ms   /16ms		PRC starts PRC ends RIB/FIB starts RIB/FIB ends	PRC starts   PRC ends   RIB/FIB starts   RIB/FIB ends
•	   000ms.   010ms.   	S-D link DOWN   	 	   Schedule SPF (in     150ms)   
1   1   1	.160ms   .161ms   .162ms   .163ms   .175ms	   Micro-loop may   start from here	SPF starts SPF ends RIB/FIB starts RIB/FIB ends	SPF starts   SPF ends   RIB/FIB starts
1	.177ms	Micro-loop ends 	 +	RIB/FIB ends

Route computation event time scale

As displayed above, there could be some other parameters like router computation power, flooding timers that may also influence microloops. In the figure 5, we consider E to be a bit slower than S,

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leading to micro-loop creation. Despite of this, we expect that by aligning implementations at least on SPF trigger and SPF delay, service provider may reduce number or duration of micro-loops.

## 7. Security Considerations

This document does not introduce any security consideration.

# Acknowledgements

Authors would like to thank Mike Shand for his useful comments.

### 9. IANA Considerations

This document has no action for IANA.

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