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Best Current Practices for NAT Traversal for Client-Server SIP draft-ietf-sipping-nat-scenarios-12

Abstract

Traversal of the Session Initiation Protocol (SIP) and the sessions it establishes through Network Address Translators (NATs) is a complex problem. Currently there are many deployment scenarios and traversal mechanisms for media traffic. This document aims to provide concrete recommendations and a unified method for NAT traversal as well as documenting corresponding flows.

Status of this Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

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1. Introduction

NAT (Network Address Translators) traversal has long been identified as a complex problem when considered in the context of the Session Initiation Protocol (SIP)[RFC3261] and it's associated media such as Real Time Protocol (RTP)[RFC3550]. The problem is exacerbated by the variety of NATs that are available in the market place today and the large number of potential deployment scenarios. Details of different NATs behavior can be found in 'NAT Behavioral Requirements for Unicast UDP' [RFC4787].

The IETF has been active on many specifications for the traversal of NATs, including STUN[RFC5389], ICE[RFC5245], symmetric response[RFC3581], symmetric RTP[RFC4961], TURN[RFC5766], SIP Outbound[RFC5626], SDP attribute for RTCP[RFC3605], Multiplexing RTP Data and Control Packets on a Single Port[RFC5761] and others. These each represent a part of the solution, but none of them gives the overall context for how the NATs traversal problem is decomposed and solved through this collection of specifications. This document serves to meet that need.

This document provides a definitive set of 'Best Common Practices' to demonstrate the traversal of SIP and its associated media through NAT devices. The document does not propose any new functionality but does draw on existing solutions for both core SIP signaling and media traversal (as defined in Section 4).

The best practices described in this document are for traditional "client-server"-style SIP. This term refers to the traditional use of the SIP protocol where User Agents talk to a series of intermediaries on a path to connect to a remote User Agent. It seems likely that other groups using SIP, for example "Peer-to-Peer-SIP(P2PSIP), will recommend these same practices between a P2PSIP client and a P2PSIP peer, but will recommend different practices for use between peers in a peer-to-peer network.

The draft is split into distinct sections as follows:

1. A clear definition of the problem statement.

2. Description of proposed solutions for both SIP protocol signaling and media signaling.

3. A set of basic and advanced flow scenarios.

2. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

It should be noted that the use of the term 'Endpoint Independent NAT' in this document refers to a NAT that is both 'Endpoint Independent Filtering NAT' and 'Endpoint Independent Mapping NAT' per RFC 4787 [RFC4787] definition.

3. Problem Statement

The traversal of SIP through NATs can be split into two categories that both require attention - The core SIP signaling and associated media traversal. This document assumes NATs that do not contain SIP-aware Application Layer Gateways(ALG), which makes much of the issues discussed in the document not applicable. ALGs have limitations (as per RFC 4787 [RFC4787]/section 7, RFC 3424 [RFC3424], and [RFC5245]/section 18.6) and experience shows they can have an adverse impact on the functionality of SIP. This includes problems such as requiring the media and signaling to traverse the same device and not working with encrypted signaling and/or payload.

It should also be noted that Session Border Controllers (SBC) doing 'hosted NAT traversal' also make many of the discussions in this document moot. More information can be obtained from [RFC5853] and [I-D.ietf-mmusic-media-path-middleboxes].

The core SIP signaling has a number of issues when traversing through NATs.

Normal SIP response routing over UDP causes the response to be delivered to the source IP address specified in the topmost Via header, or the "received" parameter of the topmost Via header. The port is extracted from the SIP 'Via' header to complete the IP address/port combination for returning the SIP response. While the destination for the response is correct, the port contained in the SIP 'Via' header represents the listening port of the originating client and not the port representing the open pin hole on the NAT. This results in responses being sent back to the NAT but to a port that is likely not open for SIP traffic. The SIP response will then be dropped at the NAT. This is illustrated in Figure 1 which depicts a SIP response being returned to port 5060.

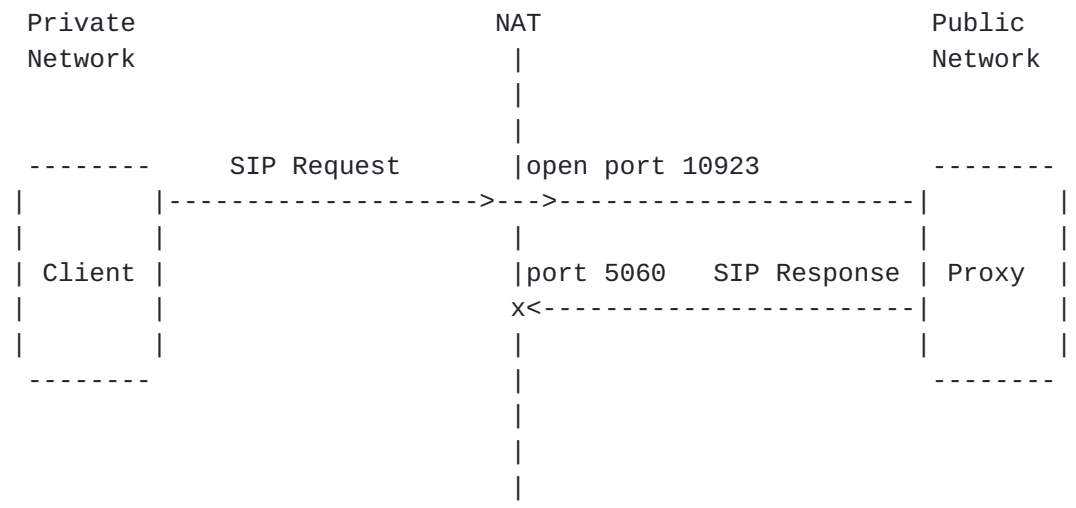


Figure 1: Failed Response

Secondly, there are two cases where new requests re-use existing connections. The first is when using a reliable, connection orientated transport protocol such as TCP, SIP has an inherent mechanism that results in SIP responses reusing the connection that was created/used for the corresponding transactional request. The SIP protocol does not provide a mechanism that allows new requests generated in the reverse direction of the originating client to use, for example, the existing TCP connection created between the client and the server during registration. This results in the registered contact address not being bound to the "connection" in the case of TCP. Requests are then blocked at the NAT, as illustrated in Figure 2. The second case is when unreliable transport protocols such as UDP where external NAT mappings need to be re-used to reach a SIP entity on the private side of the network.

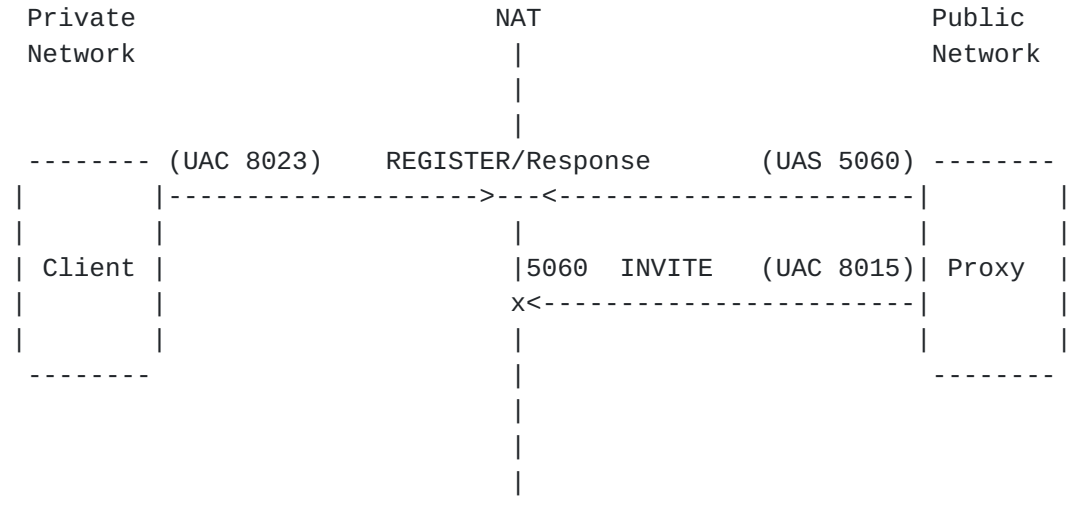


Figure 2: Failed Request

In Figure 2 the original REGISTER request is sent from the client on port 8023 and received by the proxy on port 5060, establishing a connection and opening a pin-hole in the NAT. The generation of a new request from the proxy results in a request destined for the registered entity (Contact IP address) which is not reachable from the public network. This results in the new SIP request attempting to create a connection to a private network address. This problem would be solved if the original connection was re-used. While this problem has been discussed in the context of connection orientated protocols such as TCP, the problem exists for SIP signaling using any transport protocol. The impact of connection reuse of connection orientated transports (TCP, TLS, etc) is discussed in more detail in the connection reuse specification[I-D.ietf-sip-connect-reuse]. The approach proposed for this problem in Section 4 of this document is relevant for all SIP signaling in conjunction with connection reuse, regardless of the transport protocol.

NAT policy can dictate that connections should be closed after a period of inactivity. This period of inactivity may vary from a number seconds to hours. SIP signaling can not be relied upon to keep alive connections for the following two reasons. Firstly, SIP entities can sometimes have no signaling traffic for long periods of time which has the potential to exceed the inactivity timer, and this can lead to problems where endpoints are not available to receive incoming requests as the connection has been closed. Secondly, if a low inactivity timer is specified, SIP signaling is not appropriate as a keep-alive mechanism as it has the potential to add a large amount of traffic to the network which uses up valuable resource and also requires processing at a SIP stack, which is also a waste of processing resources.

Media associated with SIP calls also has problems traversing NAT. RTP [RFC3550] runs over UDP and is one of the most common media transport types used in SIP signaling. Negotiation of RTP occurs with a SIP session establishment using the Session Description Protocol(SDP) [RFC4566] and a SIP offer/answer exchange[RFC3264]. During a SIP offer/answer exchange an IP address and port combination are specified by each client in a session as a means of receiving media such as RTP. The problem arises when a client advertises its address to receive media and it exists in a private network that is not accessible from outside the NAT. Figure 3 illustrates this problem.

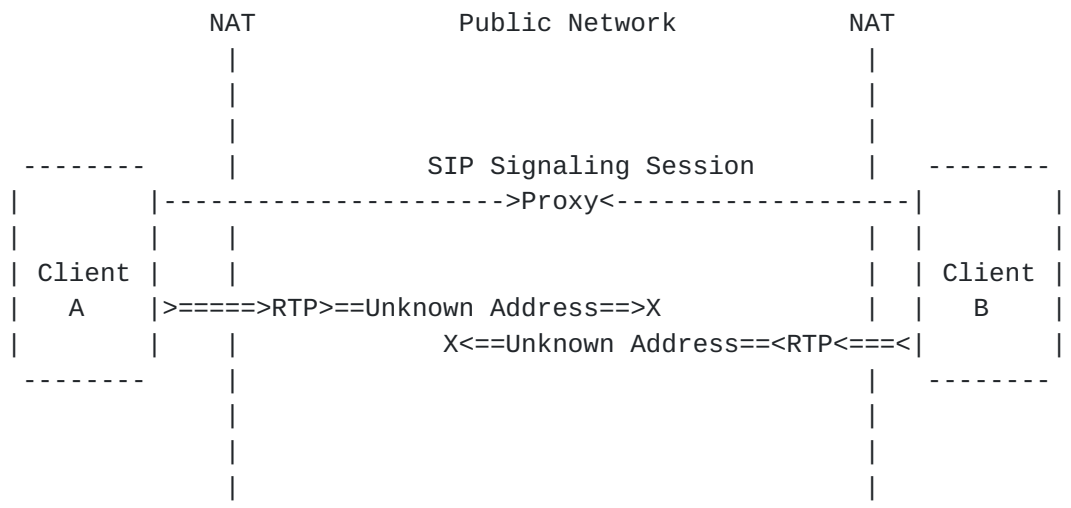


Figure 3: Failed Media

The connection addresses of the clients behind the NATs will nominally contain a private IPv4 address that is not routable across the public Internet. Exacerbating matters even more would be the tendency of Client A to send media to a destination address it received in the signaling confirmation message -- an address that may actually correspond to a host within the private network who is suddenly faced with incoming RTP packets (likewise, Client B may send media to a host within its private network who did not solicit these packets.) And finally, to complicate the problem even further, a number of different NAT topologies with different default behaviors increases the difficulty of arriving at a unified approach. This problem exists for all media transport protocols that might be NATted (e.g., TCP, UDP, SCTP, DCCP).

In general the problems associated with NAT traversal can be categorized as follows.

For signaling:

- Responses do not re-use the NAT mapping and filtering entries created by the request.

- Inbound requests are filtered out by the NAT because there is no long-term connection between the client and the proxy.

For media:

- Each endpoint has a variety of addresses that can be used to reach it (e.g., native interface address, public NATted address). In different situations, a different pair of (local endpoint, remote

endpoint) addresses should be used, and it is not clear when to use which pair.

- Many NATs filter inbound packets if the local endpoint has not recently sent an outbound packet to the sender.

- Classic RTCP usage is to run RTCP on the next highest port. However, NATs do not necessarily preserve port adjacency.

- Classic RTP and RTCP usage is to use different 5-tuples for traffic in each direction. Though not really a problem, doing this through NATs is more work than using the same 5-tuple in both directions.

4. Solution Technology Outline Description

As mentioned previously, the traversal of SIP through existing NATs can be divided into two discrete problem areas: getting the SIP signaling across NATs, and enabling media as specified by SDP in a SIP offer/answer exchange to flow between endpoints.

4.1. SIP Signaling

SIP signaling has two areas that result in transactional failure when traversing through NATs, as described in Section 3 of this document. The remaining sub-sections describe appropriate solutions that result in SIP signaling traversal through NATs, regardless of transport protocol. It is RECOMMENDED that SIP compliant entities follow the guidelines presented in this section to enable traversal of SIP signaling through NATs.

4.1.1. Symmetric Response

As described in Section 3 of this document, when using an unreliable transport protocol such as UDP, SIP responses are sent to the IP address and port combination contained in the SIP 'Via' header field (or default port for the appropriate transport protocol if not present). Figure 4 illustrates the response traversal through the open pin hole using Symmetric techniques defined in RFC 3581 [RFC3581].

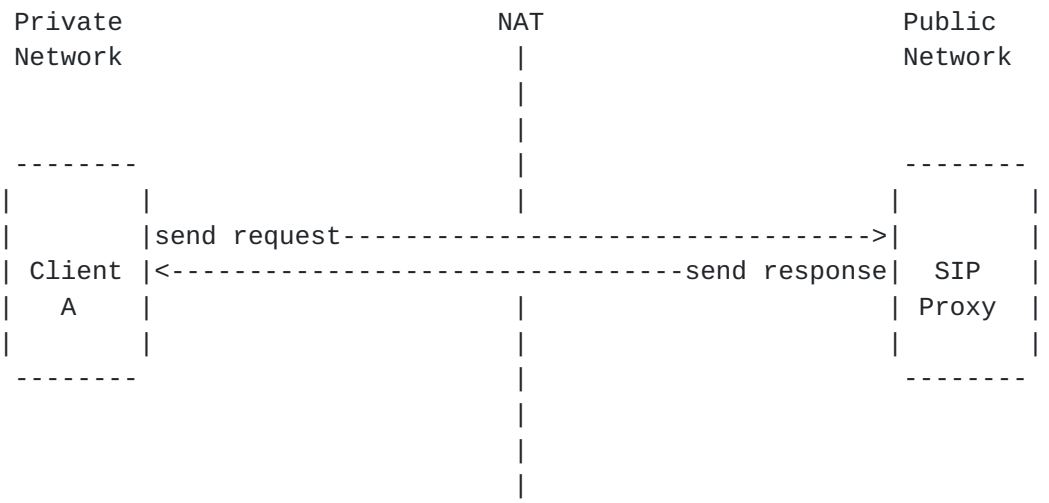


Figure 4: Symmetric Response

The outgoing request from Client A opens a pin hole in the NAT. The SIP Proxy would normally respond to the port available in the SIP Via header, as illustrated in Figure 1. The SIP Proxy honours the

'rport' parameter in the SIP Via header and routes the response to port from which it was sent. The exact functionality for this method of response traversal is called 'Symmetric Response' and the details are documented in RFC 3581 [RFC3581]. Additional requirements are imposed on SIP entities in RFC 3581 [RFC3581] such as listening and sending SIP requests/responses from the same port.

4.1.2. Client Initiated Connections

The second problem with SIP signaling, as defined in Section 3 and illustrated in Figure 2, is to allow incoming requests to be properly routed.

Guidelines for devices such as User Agents that can only generate outbound connections through NATs are documented in 'Managing Client Initiated Connections in the Session Initiation Protocol(SIP)'[RFC5626]. The document provides techniques that use a unique User Agent instance identifier (instance-id) in association with a flow identifier (reg-id). The combination of the two identifiers provides a key to a particular connection (both UDP and TCP) that is stored in association with registration bindings. On receiving an incoming request to a SIP Address-Of-Record (AOR), a proxy/registrar routes to the associated flow created by the registration and thus a route through NATs. It also provides a keepalive mechanism for clients to keep NATs bindings alive. This is achieved by multiplexing a ping/pong mechanism over the SIP signaling connection (STUN for UDP and CRLF/operating system keepalive for reliable transports like TCP). Usage of [RFC5626] is RECOMMENDED. This mechanism is not transport specific and should be used for any transport protocol.

Even if the SIP Outbound draft is not used, clients generating SIP requests SHOULD use the same IP address and port (i.e., socket) for both transmission and receipt of SIP messages. Doing so allows for the vast majority of industry provided solutions to properly function(e.g., SBC hosted NAT traversal). Deployments should also consider the mechanism described in the Connection Reuse[I-D.ietf-sip-connect-reuse] specification for routing bi-directional messages securely between trusted SIP Proxy servers.

4.2. Media Traversal

The issues of media traversal through NATs is not straight forward and requires the combination of a number of traversal methodologies. The technologies outlined in the remainder of this section provide the required solution set.

4.2.1. Symmetric RTP/RTCP

The primary problem identified in Section 3 of this document is that internal IP address/port combinations can not be reached from the public side of NATs. In the case of media such as RTP, this will result in no audio traversing NATs (as illustrated in Figure 3). To overcome this problem, a technique called 'Symmetric RTP/RTCP'[RFC4961] can be used. This involves a SIP endpoint both sending and receiving RTP/RTCP traffic from the same IP address/port combination. When operating behind a NAT and using the 'latching' technique described in [I-D.ietf-mmusic-media-path-middleboxes], SIP user agents MUST implement 'Symmetric RTP/RTCP'. This allows traversal of RTP across the NAT.

4.2.2. RTCP

Normal practice when selecting a port for defining RTP Control Protocol (RTCP) [RFC3550] is for consecutive order numbering (i.e select an incremented port for RTCP from that used for RTP). This assumption causes RTCP traffic to break when traversing certain types of NATs due to various reasons (e.g., already-allocated port, randomized port allocation). To combat this problem a specific address and port need to be specified in the SDP rather than relying on such assumptions. RFC 3605 [RFC3605] defines an SDP attribute that is included to explicitly specify transport connection information for RTCP so a separate, explicit NAT binding can be set up for the purpose. The address details can be obtained using any appropriate method including those detailed in this section (e.g. STUN, TURN, ICE).

A further enhancement to RFC 3605 [RFC3605] is defined in [RFC5761] which specifies 'muxing' both RTP and RTCP on the same IP/PORT combination.

4.2.3. STUN/TURN/ICE

ICE, STUN and TURN are a suite of 3 inter-related protocols that combine to provide a complete media traversal solution for NATs. The following sections provide details of each component part.

4.2.3.1. STUN

Session Traversal Utilities for NAT or STUN is defined in RFC 5389 [RFC5389]. STUN is a lightweight tool kit and protocol that provides details of the external IP address/port combination used by the NAT device to represent the internal entity on the public facing side of NATs. On learning of such an external representation, a client can use it accordingly as the connection address in SDP to provide NAT

traversal. Using terminology defined in the draft 'NAT Behavioral Requirements for Unicast UDP' [RFC4787], STUN does work with 'Endpoint Independent Mapping' but does not work with either 'Address Dependent Mapping' or 'Address and Port Dependent Mapping' type NATs. Using STUN with either of the previous two NATs mappings to probe for the external IP address/port representation will provide a different result to that required for traversal by an alternative SIP entity. The IP address/port combination deduced for the STUN server would be blocked for RTP packets from the remote SIP user agent.

As mentioned in Section 4.1.2, STUN is also used as a client-to-server keep-alive mechanism to refresh NAT bindings.

4.2.3.2. TURN

As described in the Section 4.2.3.1, the STUN protocol does not work for UDP traversal through certain identified NAT mappings. 'Traversal Using Relays around NAT' is a usage of the STUN protocol for deriving (from a TURN server) an address that will be used to relay packets towards a client. TURN provides an external address (globally routable) at a TURN server that will act as a media relay which attempts to allow traffic to reach the associated internal address. The full details of the TURN specification are defined in [RFC5766]. A TURN service will almost always provide media traffic to a SIP entity but it is RECOMMENDED that this method would only be used as a last resort and not as a general mechanism for NAT traversal. This is because using TURN has high performance costs when relaying media traffic and can lead to unwanted latency.

4.2.3.3. ICE

Interactive Connectivity Establishment (ICE) is the RECOMMENDED method for traversal of existing NATs if Symmetric RTP and media latching is not sufficient. ICE is a methodology for using existing technologies such as STUN, TURN and any other UNSAF[RFC3424] compliant protocol to provide a unified solution. This is achieved by obtaining as many representative IP address/port combinations as possible using technologies such as STUN/TURN (*note - an ICE endpoint can also use other mechanisms (e.g., NAT-PMP[I-D.cheshire-nat-pmp], UPnP IGD[UPnP-IGD]) to learn public IP addresses and ports, and populate a=candidate lines with that information). Once the addresses are accumulated, they are all included in the SDP exchange in a new media attribute called 'candidate'. Each 'candidate' SDP attribute entry has detailed connection information including a media address, priority and transport protocol. The appropriate IP address/port combinations are used in the order specified by the priority. A client compliant to the ICE specification will then locally run STUN servers on all

addresses being advertised using ICE. Each instance will undertake connectivity checks to ensure that a client can successfully receive media on the advertised address. Only connections that pass the relevant connectivity checks are used for media exchange. The full details of the ICE methodology are contained in [RFC5245].

5. NAT Traversal Scenarios

This section of the document includes detailed NAT traversal scenarios for both SIP signaling and the associated media. Signalling NAT traversal is achieved using [RFC5626].

5.1. Basic NAT SIP Signaling Traversal

The following sub-sections concentrate on SIP signaling traversal of NATs. The scenarios include traversal for both reliable and un-reliable transport protocols.

5.1.1. Registration (Registrar/Edge Proxy Co-Located)

The set of scenarios in this section document basic signaling traversal of a SIP REGISTER method through NATs.

5.1.1.1. UDP

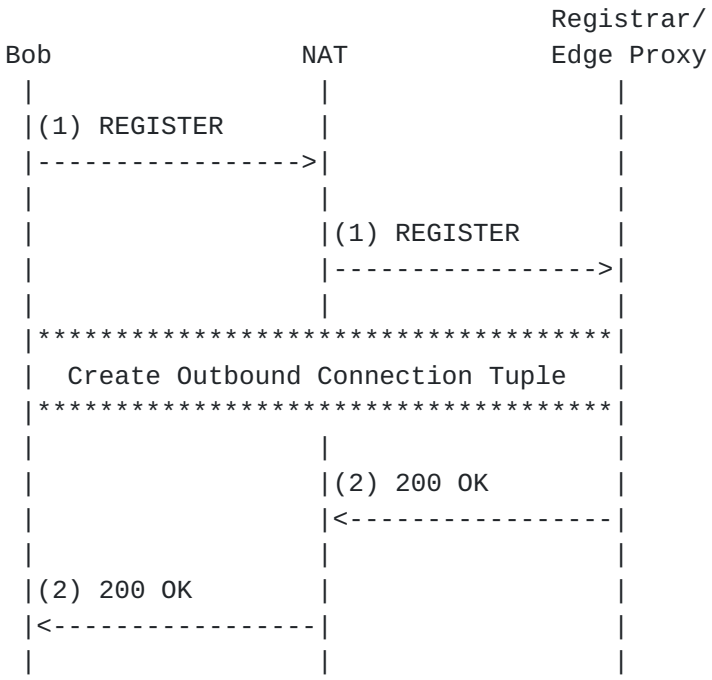


Figure 5: UDP Registration

In this example the client sends a SIP REGISTER request through a NAT. The client will include an 'rport' parameter as described in Section 4.1.1 of this document for allowing traversal of UDP responses. The original request as illustrated in (1) in Figure 5 is a standard SIP REGISTER message:

```
Message 1:

REGISTER sip:example.com SIP/2.0
Via: SIP/2.0/UDP 192.0.2.2;rport;branch=z9hG4bKnashds7
Max-Forwards: 70
From: Bob <sip:bob@example.com>;tag=7F94778B653B
To: Bob <sip:bob@example.com>
Call-ID: 16CB75F21C70
CSeq: 1 REGISTER
Supported: path, outbound
Contact: <sip:bob@192.0.2.2 >;reg-id=1
    ;+sip.instance="<urn:uuid:00000000-0000-1000-8000-AABBCCDDEEFF>"
Content-Length: 0
```

This SIP transaction now generates a SIP 200 OK response, as depicted in (2) from Figure 5:

```
Message 2:

SIP/2.0 200 OK
Via: SIP/2.0/UDP 192.0.2.2;rport=8050;branch=z9hG4bKnashds7;
     received=192.0.2.4
From: Bob <sip:bob@example.com>;tag=7F94778B653B
To: Bob <sip:bob@example.com>;tag=6AF99445E44A
Call-ID: 16CB75F21C70
CSeq: 1 REGISTER
Supported: path, outbound
Require: outbound
Contact: <sip:bob@192.0.2.2 >;reg-id=1;expires=3600
     ;+sip.instance="<urn:uuid:00000000-0000-1000-8000-AABBCCDDEEFF>"
Content-Length: 0
```

The response will be sent to the address appearing in the 'received' parameter of the SIP 'Via' header (address 192.0.2.4). The response will not be sent to the port deduced from the SIP 'Via' header, as per standard SIP operation but will be sent to the value that has been stamped in the 'rport' parameter of the SIP 'Via' header (port 8050). For the response to successfully traverse the NAT, all of the conventions defined in RFC 3581 [RFC3581] MUST be obeyed. Make note of both the 'reg-id' and 'sip.instance' contact header parameters. They are used to establish an Outbound connection tuple as defined in [RFC5626]. The connection tuple creation is clearly shown in Figure 5. This ensures that any inbound request that causes a registration lookup will result in the re-use of the connection path established by the registration. This exonerates the need to manipulate contact header URIs to represent a globally routable

address as perceived on the public side of a NAT.

5.1.1.2. Connection Oriented Transport

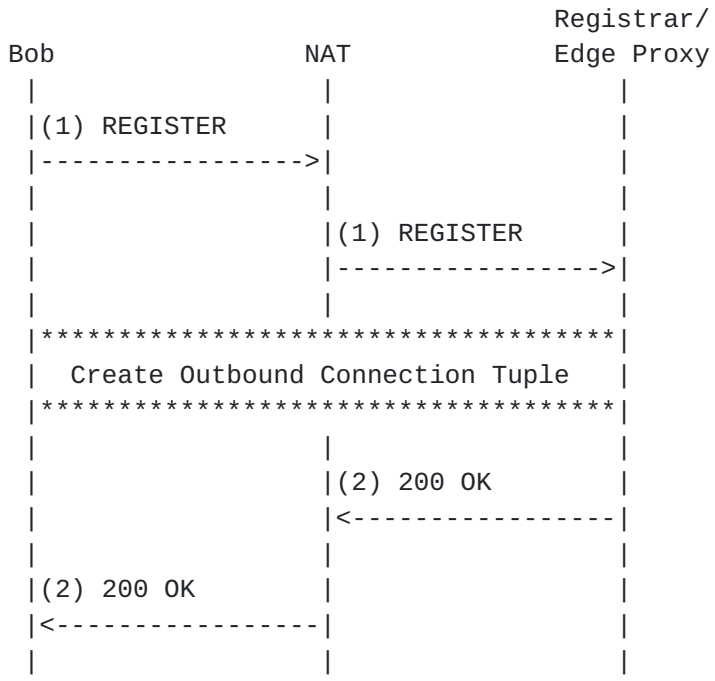


Figure 6

Traversal of SIP REGISTER requests/responses using a reliable, connection orientated protocol such as TCP does not require any additional core SIP signaling extensions, beyond the procedures defined in [RFC5626]. SIP responses will re-use the connection created for the initial REGISTER request, (1) from Figure 6:

Message 1:

```
REGISTER sip:example.com SIP/2.0
Via: SIP/2.0/TCP 192.0.2.2;branch=z9hG4bKnashds7
Max-Forwards: 70
From: Bob <sip:bob@example.com>;tag=7F94778B653B
To: Bob <sip:bob@example.com>
Call-ID: 16CB75F21C70
CSeq: 1 REGISTER
Supported: path, outbound
Contact: <sip:bob@192.0.2.2;transport=tcp>;reg-id=1
     ;+sip.instance="<urn:uuid:00000000-0000-1000-8000-AABBCCDDEEFF>"
```

```
Content-Length: 0

Message 2:


SIP/2.0 200 OK
Via: SIP/2.0/TCP 192.0.2.2;branch=z9hG4bKnashds7
From: Bob <sip:bob@example.com>;tag=7F94778B653B
To: Bob <sip:bob@example.com>;tag=6AF99445E44A
Call-ID: 16CB75F21C70
CSeq: 1 REGISTER
Supported: path, outbound
Require: outbound
Contact: <sip:bob@192.0.2.2;transport=tcp>;reg-id=1;expires=3600
     ;+sip.instance="<urn:uuid:00000000-0000-1000-8000-AABBCCDDEEFF>"
Content-Length: 0
```

This example was included to show the inclusion of the +sip.instance Contact header parameter as defined in the SIP Outbound specification [RFC5626]. This creates an association tuple as described in the previous example for future inbound requests directed at the newly created registration binding with the only difference that the association is with a TCP connection, not a UDP pin hole binding.

5.1.2. Registration(Registrar/Edge Proxy not Co-Located)

This section demonstrates traversal mechanisms when the Registrar component is not co-located with the edge proxy element. The procedures described in this section are identical, regardless of transport protocol and so only one example will be documented in the form of TCP.

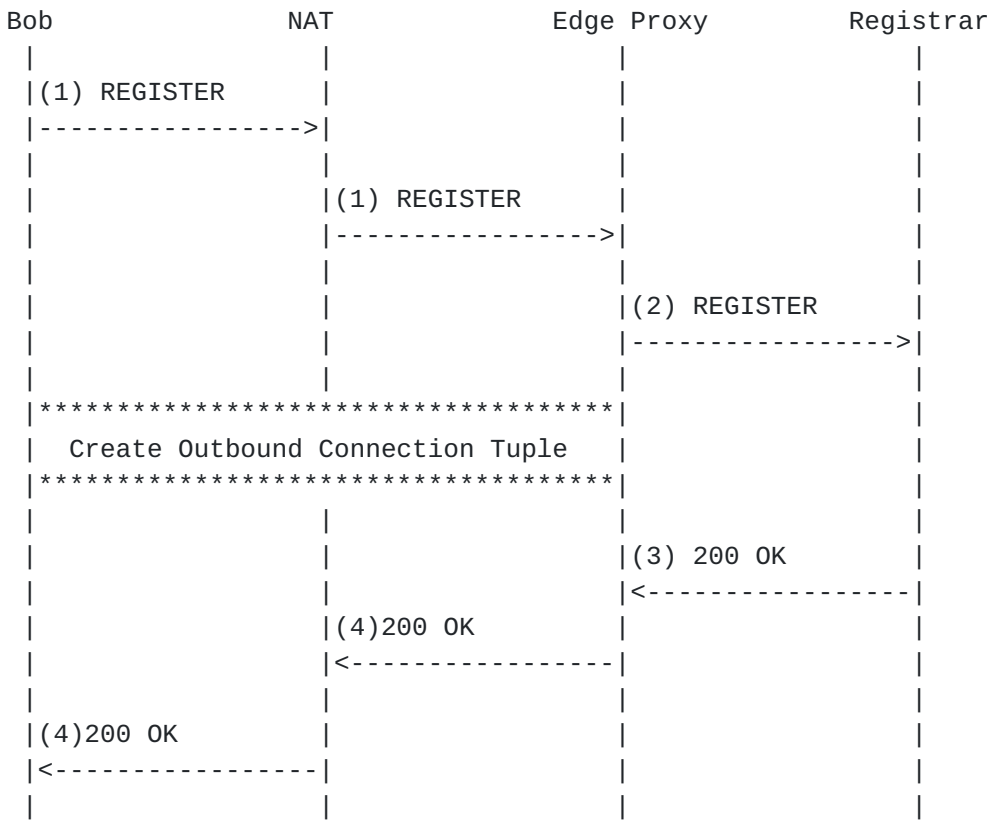


Figure 7: Registration(Registrar/Proxy not Co-Located)

This scenario builds on the previous example contained in Section 5.1.1.2. The primary difference being that the REGISTER request is routed onwards from a Proxy Server to a separated Registrar. The important message to note is (1) in Figure 7. The Edge proxy, on receiving a REGISTER request that contains a 'sip.instance' media feature tag, forms a unique flow identifier token as discussed in [RFC5626]. At this point, the proxy server routes the SIP REGISTER message to the Registrar. The proxy will create the connection tuple as described in SIP Outbound at the same moment as the co-located example, but for subsequent messages to arrive at the Proxy, the proxy needs to indicate its need to remain in the SIP signaling path. To achieve this the proxy inserts to REGISTER message (2) a SIP PATH extension header, as defined in RFC 3327 [RFC3327]. The previously created flow association token is inserted in a position within the Path header where it can easily be retrieved at a later point when receiving messages to be routed to the registration binding (in this case the user part of the SIP URI). The REGISTER message of (1) includes a SIP Route header for the edge proxy.

Message 1:

```
REGISTER sip:example.com SIP/2.0
Via: SIP/2.0/TCP 192.0.2.2;branch=z9hG4bKnashds7
Max-Forwards: 70
From: Bob <sip:bob@example.com>;tag=7F94778B653B
To: Bob <sip:bob@example.com>
Call-ID: 16CB75F21C70
CSeq: 1 REGISTER
Supported: path, outbound
Route: <sip:ep1.example.com;lr>
Contact: <sip:bob@192.0.2.2;transport=tcp>;reg-id=1
     ;+sip.instance="<urn:uuid:00000000-0000-1000-8000-AABBCCDDEEFF>"
Content-Length: 0
```

When proxied in (2) looks as follows:

Message 2:

```
REGISTER sip:example.com SIP/2.0
Via: SIP/2.0/TCP ep1.example.com;branch=z9hG4bKnuiqisi
Via: SIP/2.0/TCP 192.0.2.2;branch=z9hG4bKnashds7
Max-Forwards: 69
From: Bob <sip:bob@example.com>;tag=7F94778B653B
To: Bob <sip:bob@example.com>
Call-ID: 16CB75F21C70
CSeq: 1 REGISTER
Supported: path, outbound
Contact: <sip:bob@192.0.2.2;transport=tcp>;reg-id=1
     ;+sip.instance="<urn:uuid:00000000-0000-1000-8000-AABBCCDDEEFF>"
Path: <sip:VsktcQ/S8p4WPbOnHbuyh5iJvJIW3ib@ep1.example.com;lr;ob>
Content-Length: 0
```

This REGISTER request results in the Path header being stored along with the AOR and it's associated binding at the Registrar. The URI contained in the Path header will be inserted as a pre-loaded SIP 'Route' header into any request that arrives at the Registrar and is directed towards the associated AOR binding. This all but guarantees that all requests for the new registration will be forwarded to the Edge Proxy. In our example, the user part of the SIP 'Path' header URI that was inserted by the Edge Proxy contains the unique token identifying the flow to the client. On receiving subsequent requests, the edge proxy will examine the user part of the pre-loaded SIP 'route' header and extract the unique flow token for use in its connection tuple comparison, as defined in the SIP Outbound specification [RFC5626]. An example which builds on this scenario (showing an inbound request to the AOR) is detailed in Section 5.1.4.2 of this document.

5.1.3. Initiating a Session

This section covers basic SIP signaling when initiating a call from behind a NAT.

5.1.3.1. UDP

Initiating a call using UDP (the Edge Proxy and Authoritative Proxy functionality are co-located).

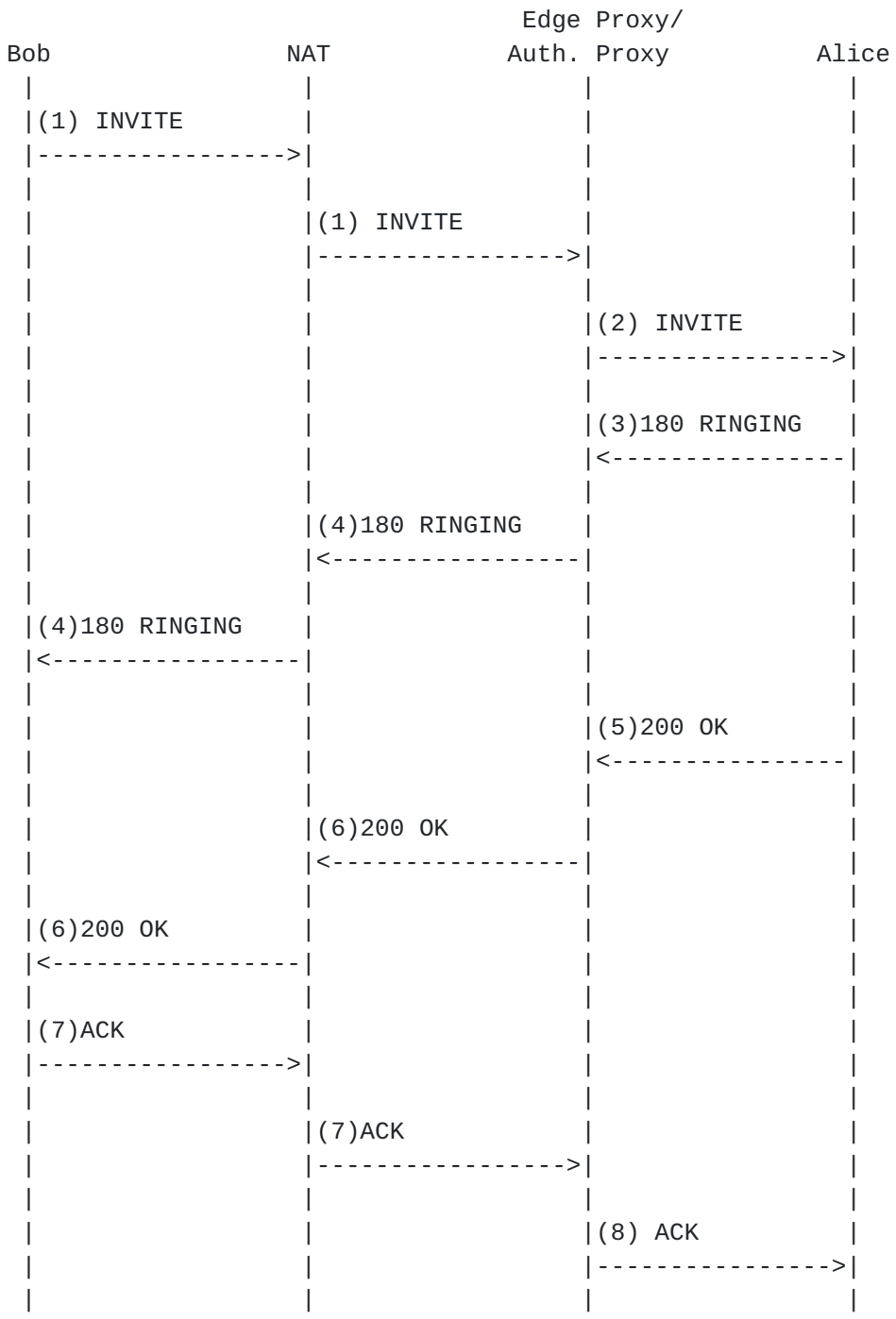


Figure 8: Initiating a Session - UDP

The initiating client generates an INVITE request that is to be sent through the NAT to a Proxy server. The INVITE message is represented in Figure 8 by (1) and is as follows:

Message 1:

```
INVITE sip:alice@a.example SIP/2.0
Via: SIP/2.0/UDP 192.0.2.2;rport;branch=z9hG4bKnashds7
Max-Forwards: 70
From: Bob <sip:bob@example.com>;tag=ldw22z
To: Alice <sip:alice@a.example>
Call-ID: 95KGsk2V/Eis9LcpBYy3
CSeq: 1 INVITE
Supported: outbound
Route: <sip:ep1.example.com;lr>
Contact: <sip:bob@192.0.2.2;ob>
Content-Type: application/sdp
Content-Length: ...

[SDP not shown]
```

There are a number of points to note with this message:

1. Firstly, as with the registration example in Section 5.1.1.1, responses to this request will not automatically pass back through a NAT and so the SIP 'Via' header 'rport' is included as described in the 'Symmetric response' Section 4.1.1 and defined in RFC 3581 [RFC3581].

2. Secondly, the contact inserted contains to ensure that all new requests will be sent to the same flow. Alternatively, a GRUU might have been used. See 4.3/[RFC5626].

In (2), the proxy inserts itself in the Via, adds the rport port number in the previous Via header, adds the received parameter in the previous Via, removes the Route header, and inserts a Record-Route with a token.

Message 2:

```
INVITE sip:alice@192.0.2.4 SIP/2.0
Via: SIP/2.0/UDP ep1.example.com;branch=z9hG4bKnuiqisi
Via: SIP/2.0/UDP 192.0.2.2;rport=8050;branch=z9hG4bKnashds7;
     received=192.0.2.14
Max-Forwards: 69
From: Bob <sip:bob@example.com>;tag=ldw22z
To: Alice <sip:alice@a.example>
Call-ID: 95KGsk2V/Eis9LcpBYy3
CSeq: 1 INVITE
Supported: outbound
Record-Route: <sip:3yJEbr1GYZK9cPYk5Snocez6DzO7w+AX@ep1.example.com;lr>
Contact: <sip:bob@192.0.2.2;ob>
Content-Type: application/sdp
Content-Length: ...

[SDP not shown]
```

5.1.3.2. Connection-oriented Transport

When using a reliable transport such as TCP the call flow and procedures for traversing a NAT are almost identical to those described in Section 5.1.3.1. The primary difference when using reliable transport protocols is that Symmetric response[RFC3581] are not required for SIP responses to traverse a NAT. RFC 3261[RFC3261] defines procedures for SIP response messages to be sent back on the same connection on which the request arrived. See section 9.5/ [RFC5626] for an example call flow of an outgoing call.

5.1.4. Receiving an Invitation to a Session

This section details scenarios where a client behind a NAT receives an inbound request through a NAT. These scenarios build on the previous registration scenario from Section 5.1.1 and Section 5.1.2 in this document.

5.1.4.1. Registrar/Proxy Co-located

The SIP signaling on the interior of the network (behind the user's proxy) is not impacted directly by the transport protocol and so only one example scenario is necessary. The example uses UDP and follows on from the registration installed in the example from Section 5.1.1.1.

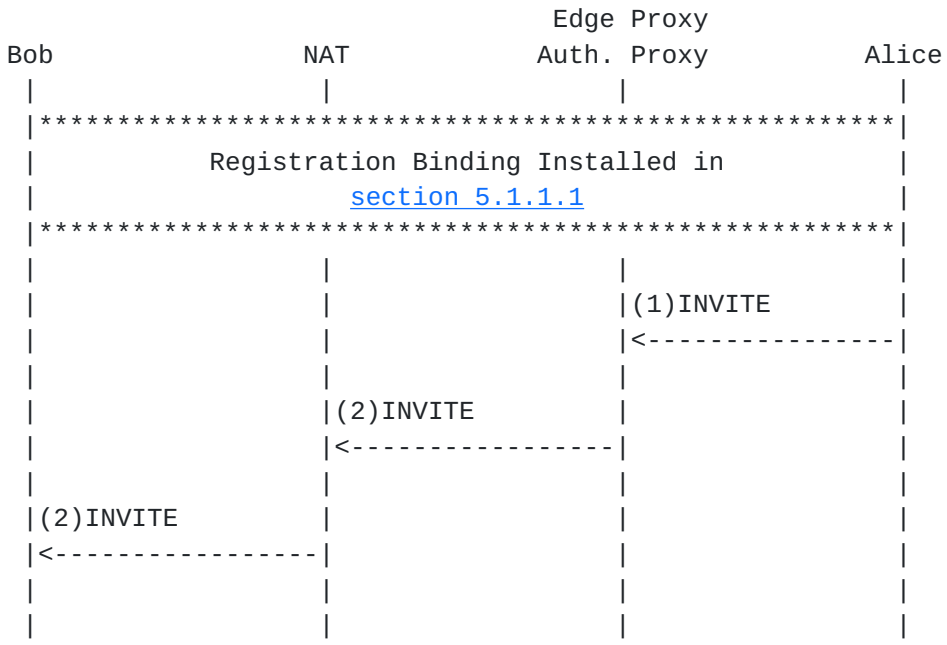


Figure 9: Receiving an Invitation to a Session

An INVITE request arrives at the Authoritative Proxy with a destination pointing to the AOR of that inserted in Section 5.1.1.1. The message is illustrated by (1) in Figure 9 and looks as follows:

```
INVITE sip:bob@example.com SIP/2.0
Via: SIP/2.0/UDP 192.0.2.4;branch=z9hG4bK74huHJ37d
Max-Forwards: 70
From: External Alice <sip:alice@example.com>;tag=02935
To: Bob <sip:bob@example.com>
Call-ID: klmvCxVWGp6MxJp2T2mb
CSeq: 1 INVITE
Contact: <sip:alice@192.0.2.4>
Content-Type: application/sdp
Content-Length: ..

[SDP not shown]
```

The INVITE request matches the registration binding previously installed at the Registrar and the INVITE request-URI is re-written to the selected onward address. The proxy then examines the request URI of the INVITE and compares with its list of connection tuples. It uses the incoming AOR to commence the check for associated open connections/mappings. Once matched, the proxy checks to see if the unique instance identifier (+sip.instance) associated with the binding equals the same instance identifier associated with that

connection tuple. The request is then dispatched on the appropriate binding. This is message (2) from Figure 9 and is as follows:

```
INVITE sip:bob@192.0.2.2 SIP/2.0
Via: SIP/2.0/UDP ep1.example.com;branch=z9hG4kmlds893jhsd
Via: SIP/2.0/UDP 192.0.2.4;branch=z9hG4bK74huHJ37d
Max-Forwards: 69
From: Alice <sip:alice@example.com>;tag=02935
To: client bob <sip:bob@example.com>
Call-ID: klmvCxVWGp6MxJp2T2mb
CSeq: 1 INVITE
Contact: <sip:alice@192.0.2.4>
Content-Type: application/sdp
Content-Length: ..

[SDP not shown]
```

It is a standard SIP INVITE request with no additional functionality. The major difference being that this request will not be forwarded to the address specified in the Request-URI, as standard SIP rules would enforce but will be sent on the flow associated with the registration binding (look-up procedures in RFC 3263 [RFC3263] are overridden). This then allows the original connection/mapping from the initial registration process to be re-used.

5.1.4.2. Edge Proxy/Authoritative Proxy Not Co-located

The core SIP signaling associated with this call flow is not impacted directly by the transport protocol and so only one example scenario is necessary. The example uses UDP and follows on from the registration installed in the example from Section 5.1.2.

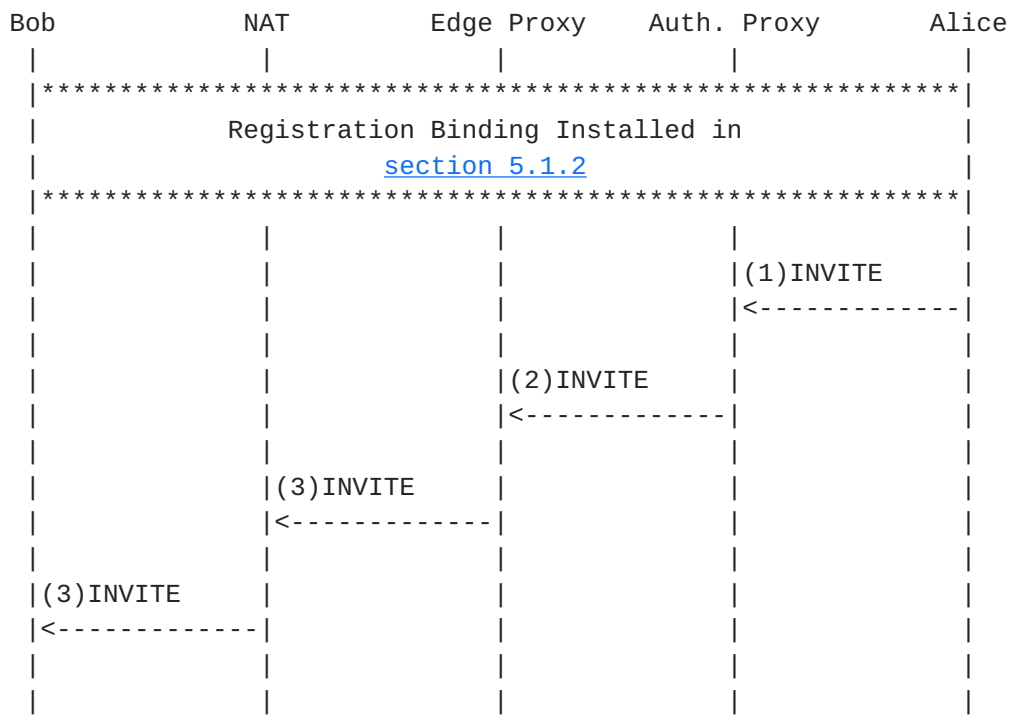


Figure 10: Registrar/Proxy Not Co-located

An INVITE request arrives at the Authoritative Proxy with a destination pointing to the AOR of that inserted in Section 5.1.2. The message is illustrated by (1) in Figure 10 and looks as follows:

```
INVITE sip:bob@example.com SIP/2.0
Via: SIP/2.0/UDP 192.0.2.4;branch=z9hG4bK74huHJ37d
Max-Forwards: 70
From: Alice <sip:alice@example.com>;tag=02935
To: Bob <sip:bob@example.com>
Call-ID: klmvCxVWGp6MxJp2T2mb
CSeq: 1 INVITE
Contact: <sip:external@192.0.2.4>
Content-Type: application/sdp
Content-Length: ..

[SDP not shown]
```

The INVITE request matches the registration binding previously installed at the Registrar and the INVITE request-URI is re-written to the selected onward address. The Registrar also identifies that a SIP PATH header was associated with the registration and pushes it into the INVITE request in the form of a pre-loaded SIP Route header.

It then forwards the request on to the proxy identified in the SIP Route header as shown in (2) from Figure 10:

```
INVITE sip:bob@client.example.com SIP/2.0
Via: SIP/2.0/UDP proxy.example.com;branch=z9hG4bK74fmljnc
Via: SIP/2.0/UDP 192.0.2.4;branch=z9hG4bK74huHJ37d
Route: <sip:VsktcQ/S8p4WPbOnHbuyh5iJvJIW3ib@ep1.example.com;lr;ob>
Max-Forwards: 69
From: Alice <sip:alice@example.net>;tag=02935
To: Bob <sip:Bob@example.com>
Call-ID: klmvCxVWGp6MxJp2T2mb
CSeq: 1 INVITE
Contact: <sip:alice@192.0.2.4>
Content-Type: application/sdp
Content-Length: ..

[SDP not shown]
```

The request then arrives at the outbound proxy for the client. The proxy examines the request URI of the INVITE in conjunction with the flow token that it previously inserted into the user part of the PATH header SIP URI (which now appears in the user part of the Route header in the incoming INVITE). The proxy locates the appropriate flow and sends the message to the client, as shown in (3) from Figure 10:

```
INVITE sip:bob@192.0.2.2 SIP/2.0
Via: SIP/2.0/UDP ep1.example.com;branch=z9hG4nsi30dncmnl
Via: SIP/2.0/UDP proxy.example.com;branch=z9hG4bK74fmljnc
Via: SIP/2.0/UDP 192.0.2.4;branch=z9hG4bK74huHJ37d
Record-Route: <sip:VsktcQ/S8p4WPbOnHbuyh5iJvJIW3ib@ep1.example.com;lr>
Max-Forwards: 68
From: Alice <sip:Alice@example.net>;tag=02935
To: bob <sip:bob@example.com>
Call-ID: klmvCxVWGp6MxJp2T2mb
CSeq: 1 INVITE
Contact: <sip:alice@192.0.2.4>
Content-Type: application/sdp
Content-Length: ..

[SDP not shown]
```

It is a standard SIP INVITE request with no additional functionality at the originator. The major difference being that this request will not follow the address specified in the Request-URI when it reaches the outbound proxy, as standard SIP rules would enforce but will be

sent on the flow associated with the registration binding as indicated in the Route header(look-up procedures in RFC 3263 [RFC3263] are overridden). This then allows the original connection/ mapping from the initial registration to the outbound proxy to be re-used.

5.2. Basic NAT Media Traversal

This section provides example scenarios to demonstrate basic media traversal using the techniques outlined earlier in this document.

In the flow diagrams STUN messages have been annotated for simplicity as follows:

- The "Src" attribute represents the source transport address of the message.
- The "Dest" attribute represents the destination transport address of the message.
- The "Map" attribute represents the server reflexive (XOR-MAPPED-ADDRESS STUN attribute) transport address.
- The "Rel" attribute represents the relayed (RELAY-ADDRESS STUN attribute) transport address.

The meaning of each STUN attribute is extensively explained in the core STUN[RFC5389] and TURN [RFC5766] specifications.

A number of ICE SDP attributes have also been included in some of the examples. Detailed information on individual attributes can be obtained from the core ICE specification[RFC5245].

The examples also contain a mechanism for representing transport addresses. It would be confusing to include representations of network addresses in the call flows and make them hard to follow. For this reason network addresses will be represented using the following annotation. The first component will contain the representation of the client responsible for the address. For example in the majority of the examples "L" (left client), "R" (right client), NAT-PUB" (NAT public), PRIV (Private), and "STUN-PUB" (STUN Public) are used. To allow for multiple addresses from the same network element, each representation can also be followed by a number. These can also be used in combination. For example "L-NAT-PUB-1" would represent a public network address of the left hand side NAT while "R-NAT-PUB-1" would represent a public network address of the right hand side of the NAT. "L-PRIV-1" would represent a private network address of the left hand side of the NAT while "R-PRIV-1"

represents a private address of the right hand side of the NAT.

It should also be noted that during the examples it might be appropriate to signify an explicit part of a transport address. This is achieved by adding either the '.address' or '.port' tag on the end of the representation. For example, 'L-PRIV-1.address' and 'L-PRIV-1.port'.

The use of '$' signifies variable parts in example SIP messages.

5.2.1. Endpoint Independent NAT

This section demonstrates an example of a client both initiating and receiving calls behind an 'Endpoint independent' NAT. An example is included for both STUN and ICE with ICE being the RECOMMENDED mechanism for media traversal.

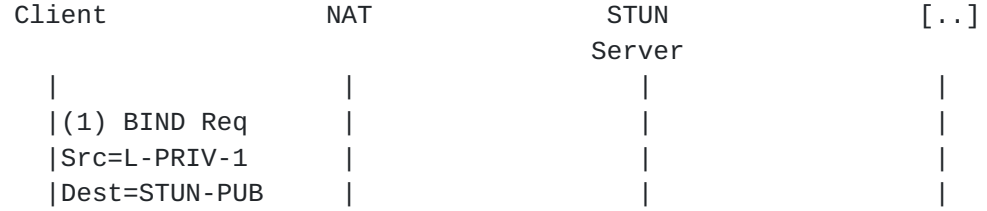
At this time there is no reliable test to determine if a host is behind an 'endpoint independent filtering' NAT or an 'endpoint independent mapping' NAT [RFC5780], and the sort of failure that occurs in this situation is described in Section 5.2.2.1. For this reason, ICE is RECOMMENDED over the mechanism described in this section.

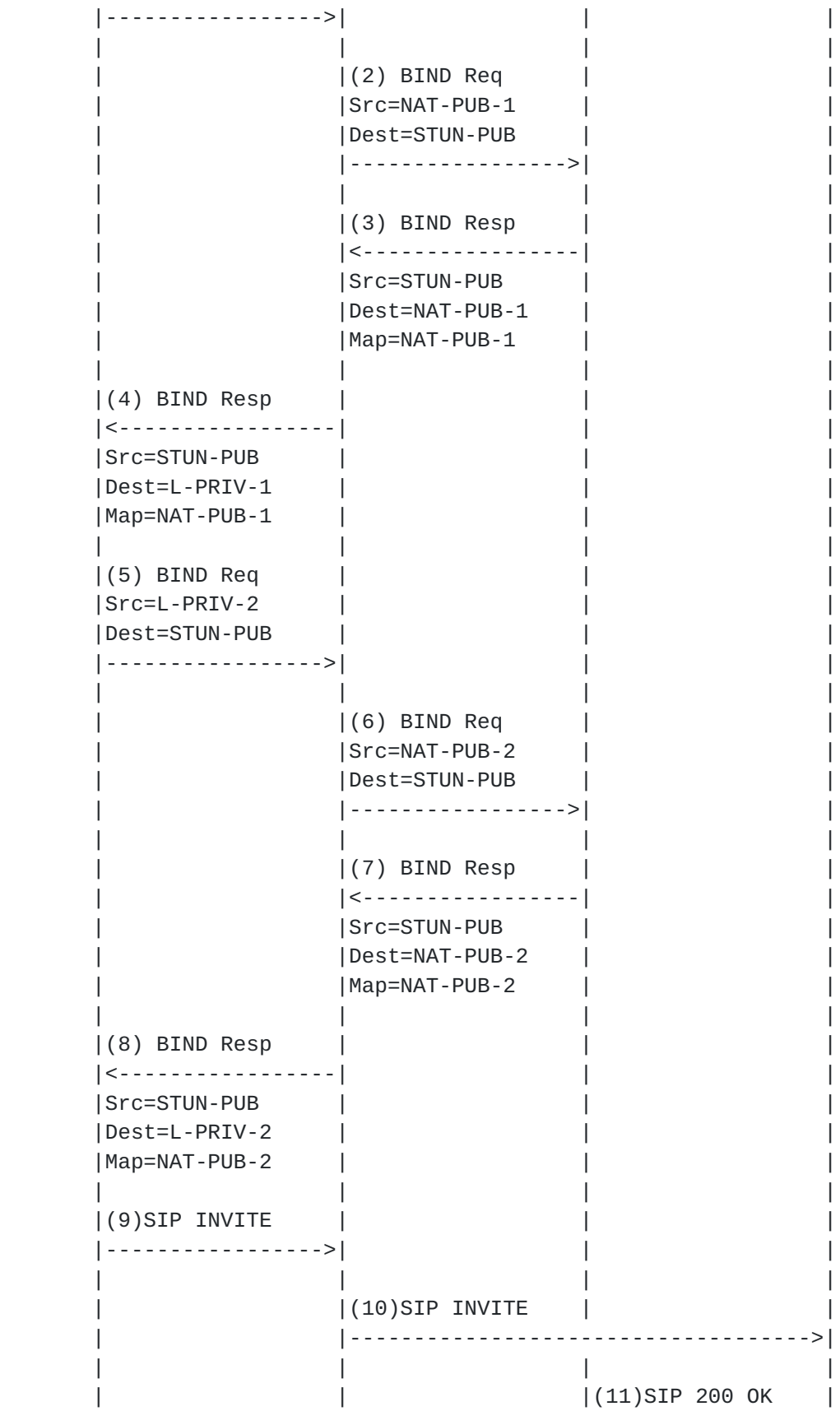
5.2.1.1. STUN Solution

It is possible to traverse media through an 'Endpoint Independent NAT using STUN. The remainder of this section provides simplified examples of the 'Binding Discovery' STUN as defined in [RFC5389]. The STUN messages have been simplified and do not include 'Shared Secret' requests used to obtain the temporary username and password.

5.2.1.1.1. Initiating Session

The following example demonstrates media traversal through a NAT with 'Endpoint-Independent Mapping' properties using the STUN 'Binding Discovery' usage. It is assumed in this example that the STUN client and SIP Client are co-located on the same physical machine. Note that some SIP signaling messages have been left out for simplicity.




(2) BIND Req
Src=NAT-PUB-1
Dest=STUN-PUB
(3) BIND Resp
Src=STUN-PUB
Dest=NAT-PUB-1
Map=NAT-PUB-1
(4) BIND Resp
Src=STUN-PUB
Dest=L-PRIV-1
Map=NAT-PUB-1
(5) BIND Req
Src=L-PRIV-2
Dest=STUN-PUB
(6) BIND Req
Src=NAT-PUB-2
Dest=STUN-PUB
(7) BIND Resp
Src=STUN-PUB
Dest=NAT-PUB-2
Map=NAT-PUB-2
(8) BIND Resp
Src=STUN-PUB
Dest=L-PRIV-2
Map=NAT-PUB-2
(9)SIP INVITE
(10)SIP INVITE
(11)SIP 200 OK

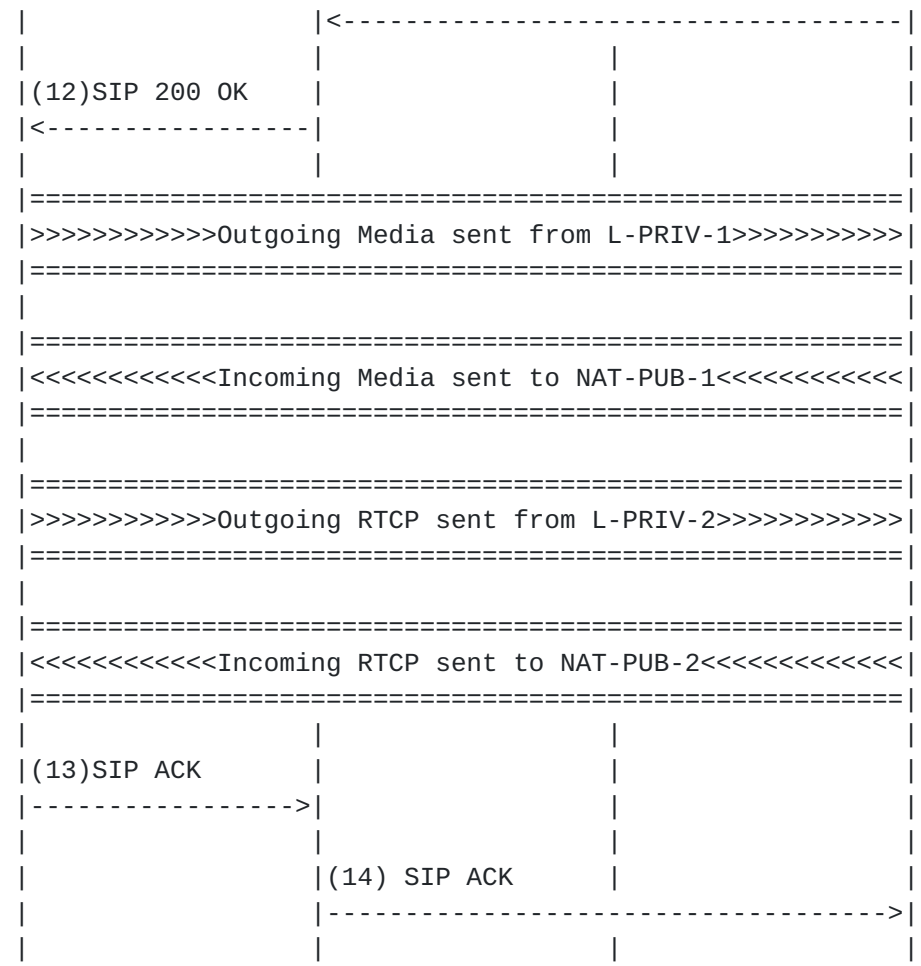


Figure 11: Endpoint Independent NAT - Initiating

- On deciding to initiate a SIP voice session the client starts a local STUN client on the interface and port that is to be used for media (send/receive). The STUN client generates a standard 'Binding Discovery' request as indicated in (1) from Figure 11 which also highlights the source address and port for which the client device wishes to obtain a mapping. The 'Binding Discovery' request is sent through the NAT towards the public internet and STUN server.

- Message (2) traverses the NAT and breaks out onto the public internet towards the public STUN server. Note that the source address of the 'Binding Discovery' request now represents the public address and port from the public side of the NAT.

- The STUN server receives the request and processes it appropriately. This results in a successful 'Binding Discovery' response being generated and returned (3). The message contains

details of the XOR mapped public address (contained in the STUN XOR-MAPPED-ADDRESS attribute) which is to be used by the originating client to receive media (see 'Map=NAT-PUB-1' from (3)).

- The 'Binding Discovery' response traverses back through the NAT using the path created by the 'Binding Discovery' request and presents the new XOR mapped address to the client (4). At this point the process is repeated to obtain a second XOR-mapped address (as shown in (5)-(8)) for a second local address (Address has changed from "L-PRIV-1" to "L-PRIV-2") for an RTCP port.

- The client now constructs a SIP INVITE message(9). Note that traversal of SIP is not covered in this example and is discussed in Section 5.1. The INVITE request will use the addresses it has obtained in the previous STUN transactions to populate the SDP of the SIP INVITE as shown below:

```
v=0
o=test 2890844526 2890842807 IN IP4 $L-PRIV-1.address
c=IN IP4 $NAT-PUB-1.address
t=0 0
m=audio $NAT-PUB-1.port RTP/AVP 0
a=rtcp:$NAT-PUB-2.port
```

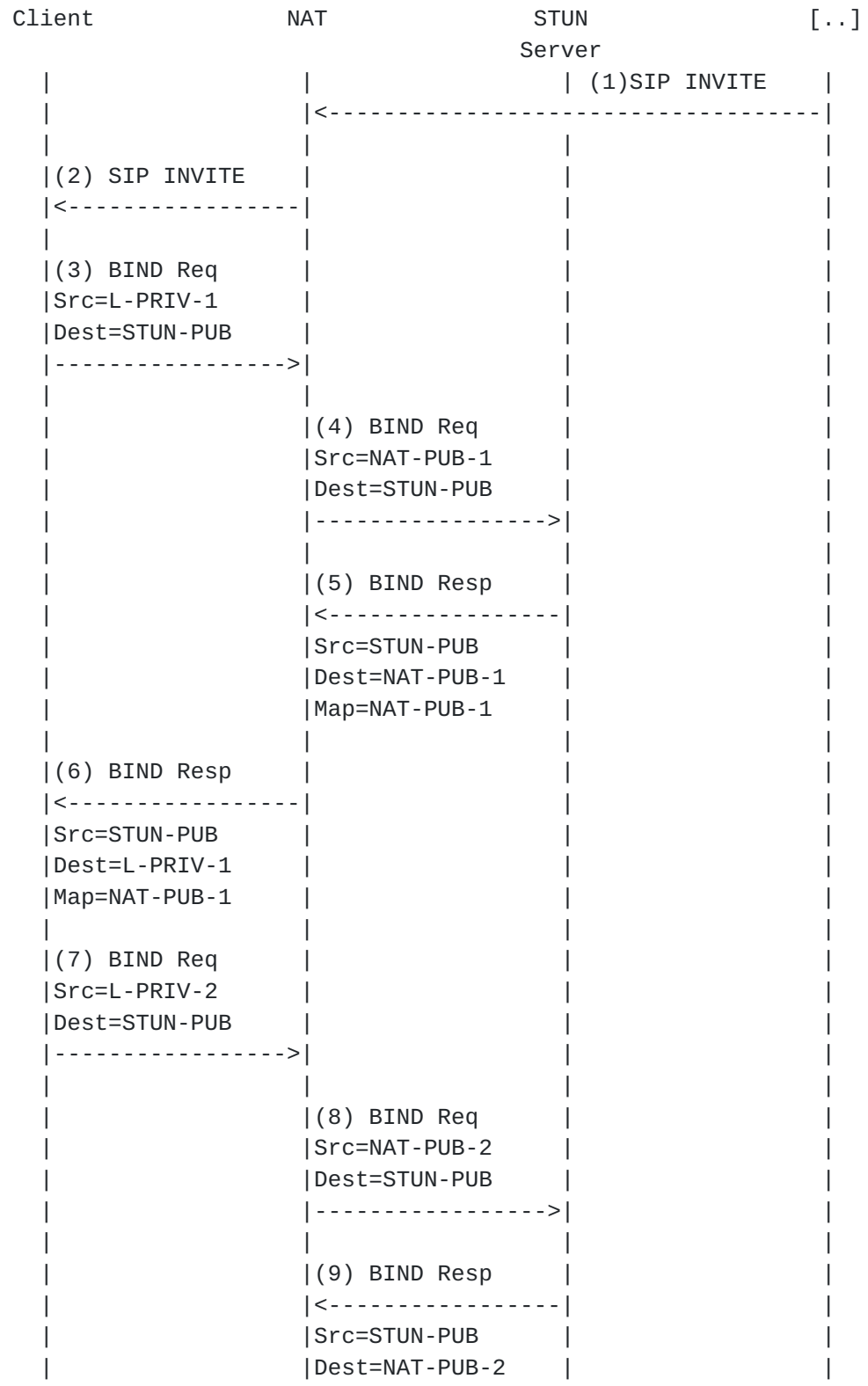
- Note that the XOR-mapped address obtained from the 'Binding Discovery' transactions are inserted as the connection address for the SDP (c=$NAT-PUB-1.address). The Primary port for RTP is also inserted in the SDP (m=audio $NAT-PUB-1.port RTP/AVP 0). Finally, the port gained from the additional 'Binding Discovery' is placed in the RTCP attribute (as discussed in Section 4.2.2) for traversal of RTCP (a=rtcp:$NAT-PUB-2.port).

- The SIP signaling then traverses the NAT and sets up the SIP session (9-12). Note that the left client transmits media as soon as the 200 OK to the INVITE arrives at the client (12). Up until this point the incoming media and RTCP to the left hand client will not pass through the NAT as no outbound association has been created with the far end client. Two way media communication has now been established.

5.2.1.1.2. Receiving Session Invitation

Receiving a session for an 'Endpoint Independent' NAT using the STUN 'Binding Discovery' usage is very similar to the example outlined in Section 5.2.1.1.1. Figure 12 illustrates the associated flow of

messages.


Client
NAT
STUN Server
[..]
(1)SIP INVITE
(2) SIP INVITE
(3) BIND Req
Src=L-PRIV-1
Dest=STUN-PUB
(4) BIND Req
Src=NAT-PUB-1
Dest=STUN-PUB
(5) BIND Resp
Src=STUN-PUB
Dest=NAT-PUB-1
Map=NAT-PUB-1
(6) BIND Resp
Src=STUN-PUB
Dest=L-PRIV-1
Map=NAT-PUB-1
(7) BIND Req
Src=L-PRIV-2
Dest=STUN-PUB
(8) BIND Req
Src=NAT-PUB-2
Dest=STUN-PUB
(9) BIND Resp
Src=STUN-PUB
Dest=NAT-PUB-2

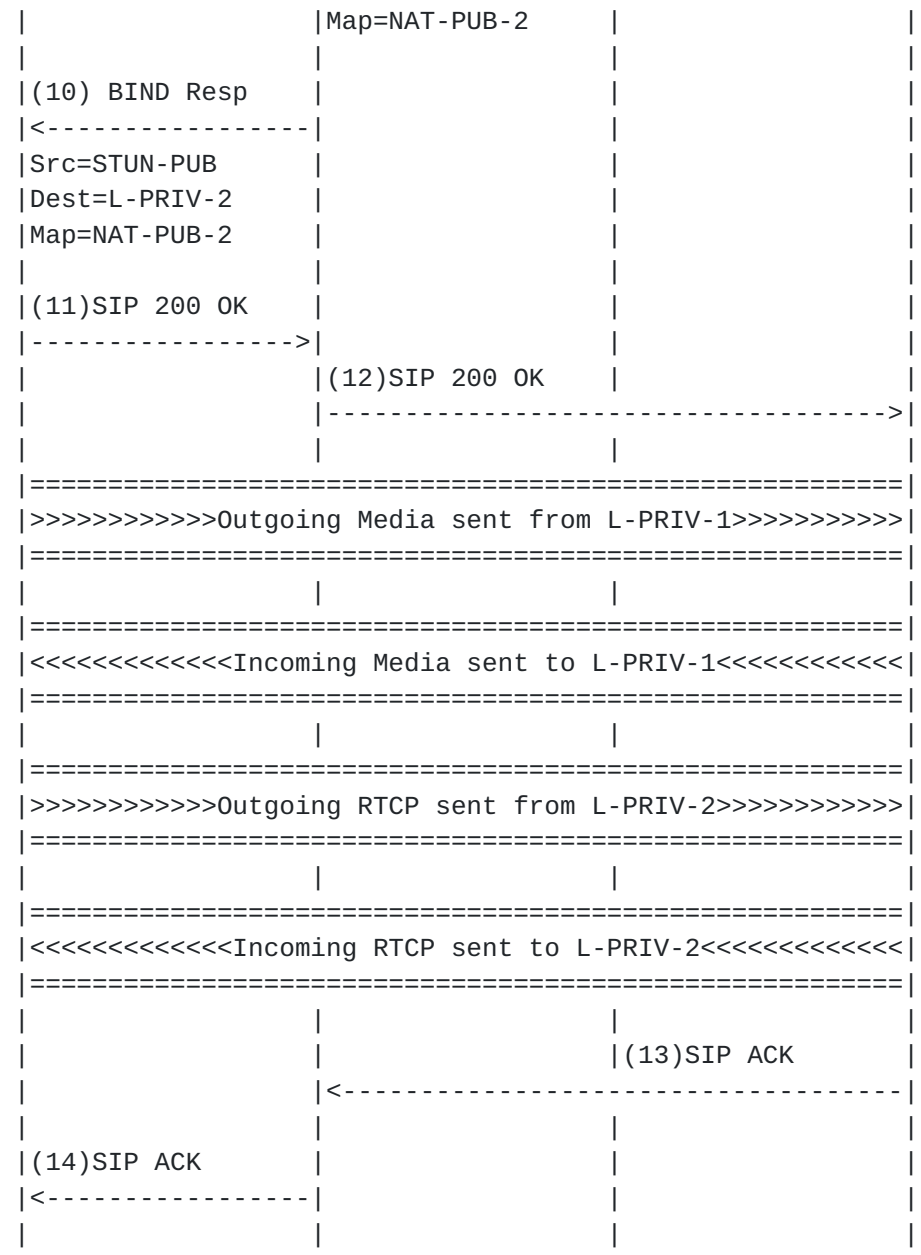


Figure 12: Endpoint Independent NAT - Receiving

- On receiving an invitation to a SIP voice session (SIP INVITE request) the User Agent starts a local STUN client on the appropriate port on which it is to receive media. The STUN client generates a standard 'Binding Discovery' request as indicated in (3) from Figure 12 which also highlights the source address and port for which the client device wishes to obtain a mapping. The 'Binding Discovery' request is sent through the NAT towards the public internet and STUN Server.

- 'Binding Discovery' message (4) traverses the NAT and breaks out onto the public internet towards the public STUN server. Note that the source address of the STUN requests now represents the public address and port from the public side of the NAT.

- The STUN server receives the request and processes it appropriately. This results in a successful 'Binding Discovery' response being generated and returned (5). The message contains details of the mapped public address (contained in the STUN XOR-MAPPED-ADDRESS attribute) which is to be used by the originating client to receive media (see 'Map=NAT-PUB-1' from (5)).

- The 'Binding Discovery' response traverses back through the NAT using the path created by the outgoing 'Binding Discovery' request and presents the new XOR-mapped address to the client (6). At this point the process is repeated to obtain a second XOR-mapped address (as shown in (7)-(10)) for a second local address (local port has now changed and is represented by L-PRIV-2 in (7)) for an RTCP port.

- The client now constructs a SIP 200 OK message (11) in response to the original SIP INVITE requests. Note that traversal of SIP is not covered in this example and is discussed in Section 5.1. SIP Provisional responses are also left out for simplicity. The 200 OK response will use the addresses it has obtained in the previous STUN transactions to populate the SDP of the SIP 200 OK as shown below:

```
v=0
o=test 2890844526 2890842807 IN IP4 $L-PRIV-1.address
c=IN IP4 $NAT-PUB-1.address
t=0 0
m=audio $NAT-PUB-1.port RTP/AVP 0
a=rtcp:$NAT-PUB-2.port
```

- Note that the XOR-mapped address obtained from the initial 'Binding Discovery' transaction is inserted as the connection address for the SDP (c=NAT-PUB-1.address). The Primary port for RTP is also inserted in the SDP (m=audio NAT-PUB-1.port RTP/AVP 0). Finally, the port gained from the second 'Binding Discovery' is placed in the RTCP attribute (as discussed in Section 4.2.2) for traversal of RTCP (a=rtcp:NAT-PUB-2.port).

- The SIP signaling then traverses the NAT and sets up the SIP session (11-14). Note that the left hand client transmits media as soon as the 200 OK to the INVITE is sent to the UAC(11). Up

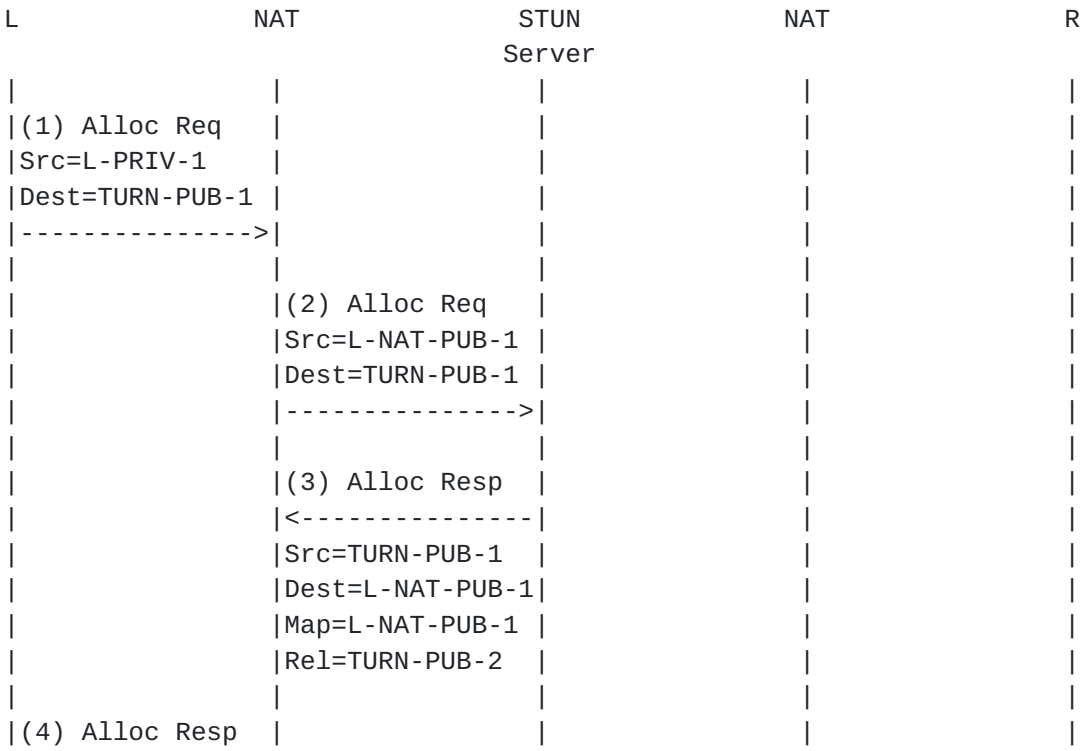
until this point the incoming media from the right hand client will not pass through the NAT as no outbound association has been created with the far end client. Two way media communication has now been established.

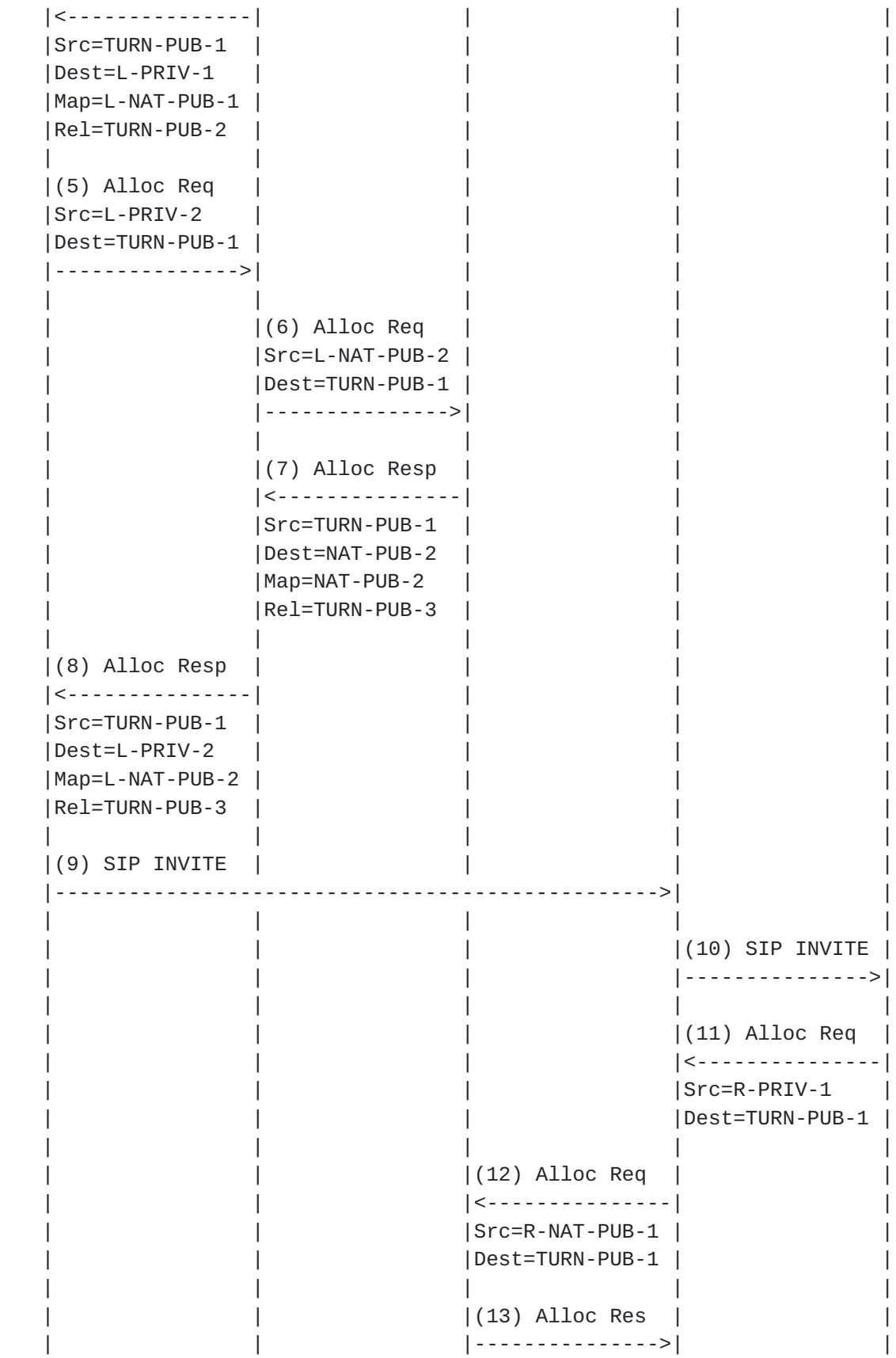
5.2.1.2. ICE Solution

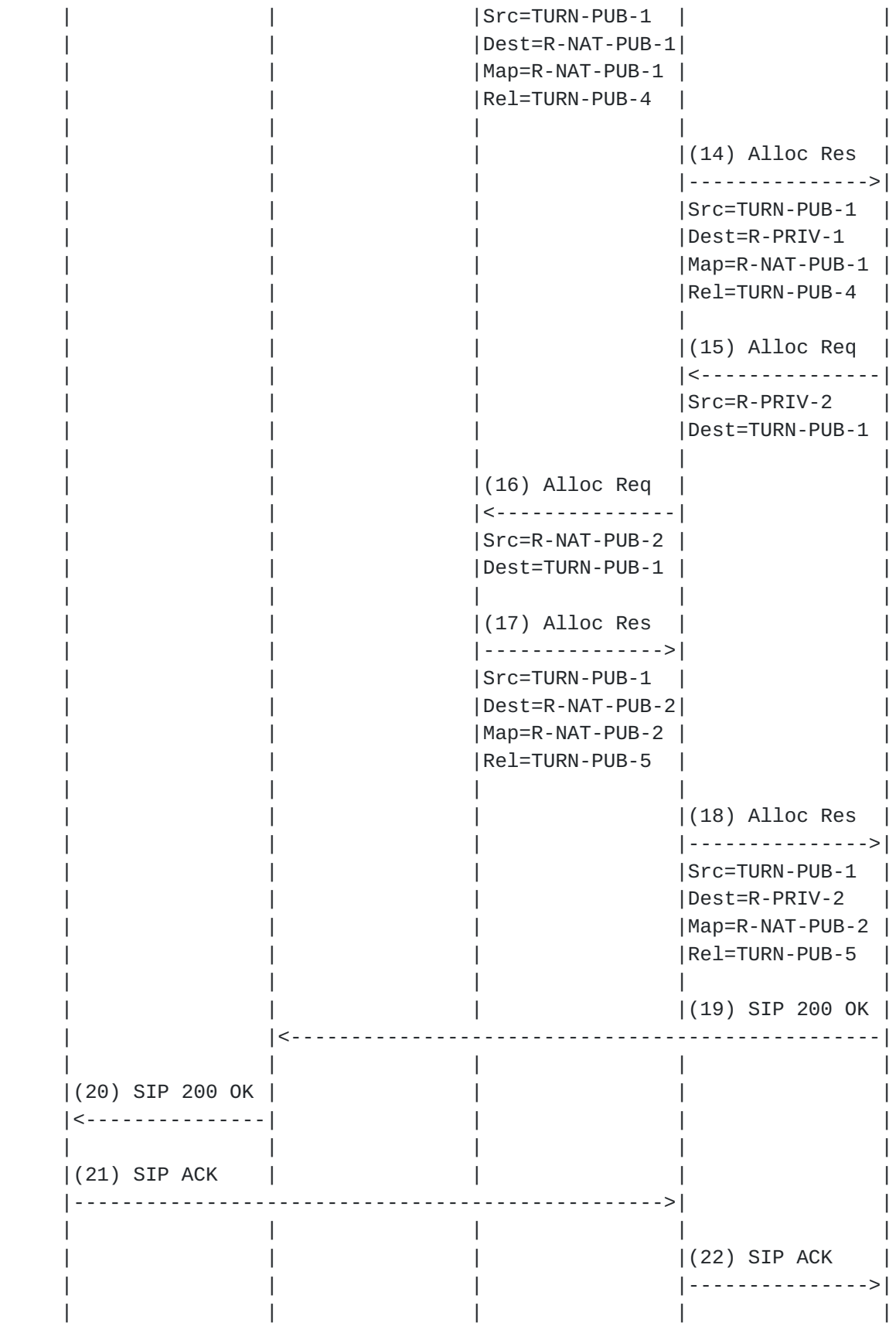
The preferred solution for media traversal of NAT is using ICE, as described in Section 4.2.3.3, regardless of the NAT type. The following examples illustrate the traversal of an 'Endpoint Independent' NAT when initiating the session. The example only covers ICE in association with the 'Binding Discovery' and TURN. It is worth noting that the TURN server provides both STUN functions(to learn your public mapping) and TURN functions (media relaying). It is also worth noting that in example described in Section 5.2.1.2.1, that both SIP clients 'L' and 'R' are contacting the same TURN server. This is not necessary for ICE, STUN, TURN to function; all that is necessary is that the STUN and TURN server(s) be in the same addressing domain - that is accessible on the Internet.

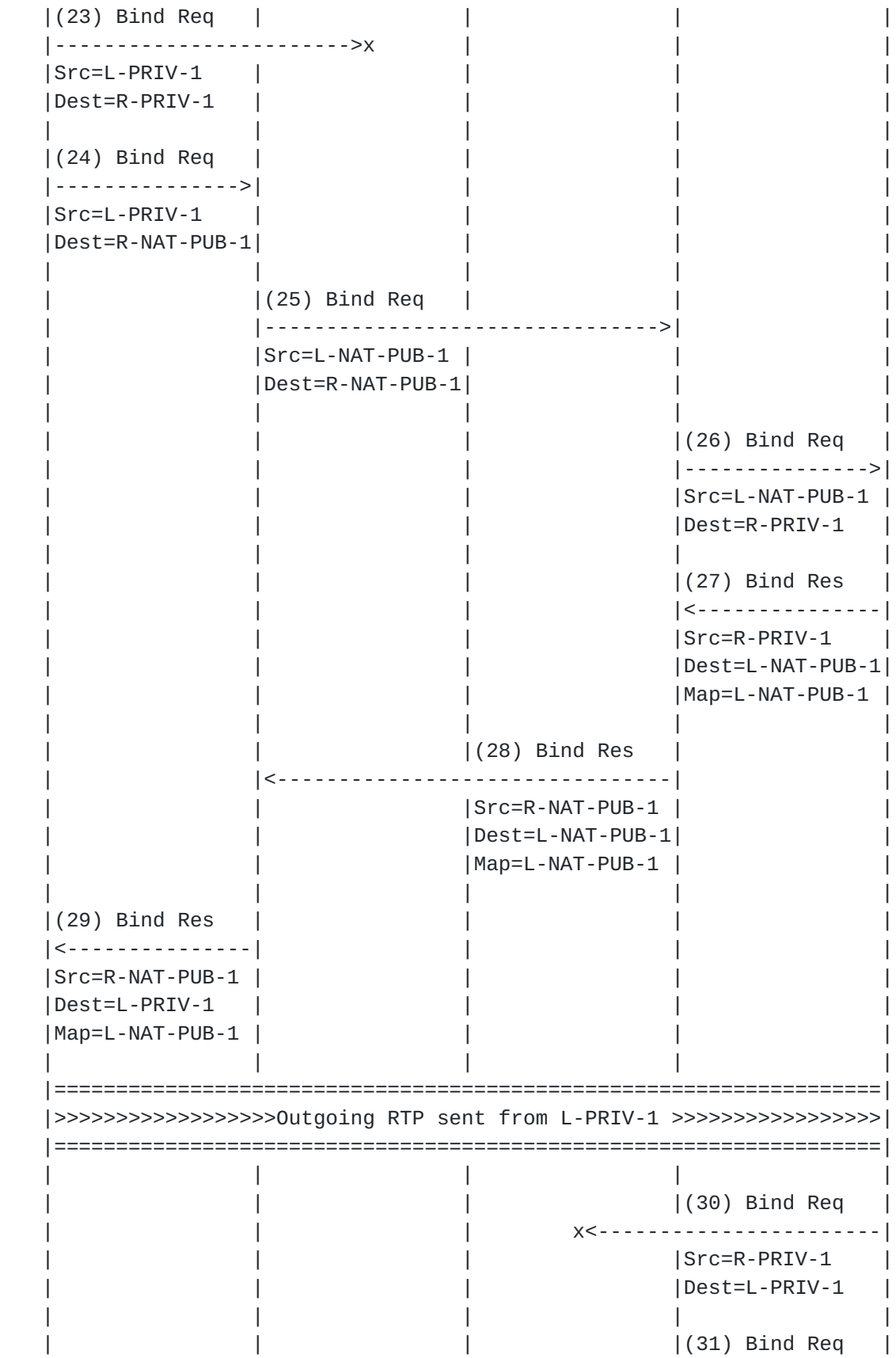
5.2.1.2.1. Initiating Session

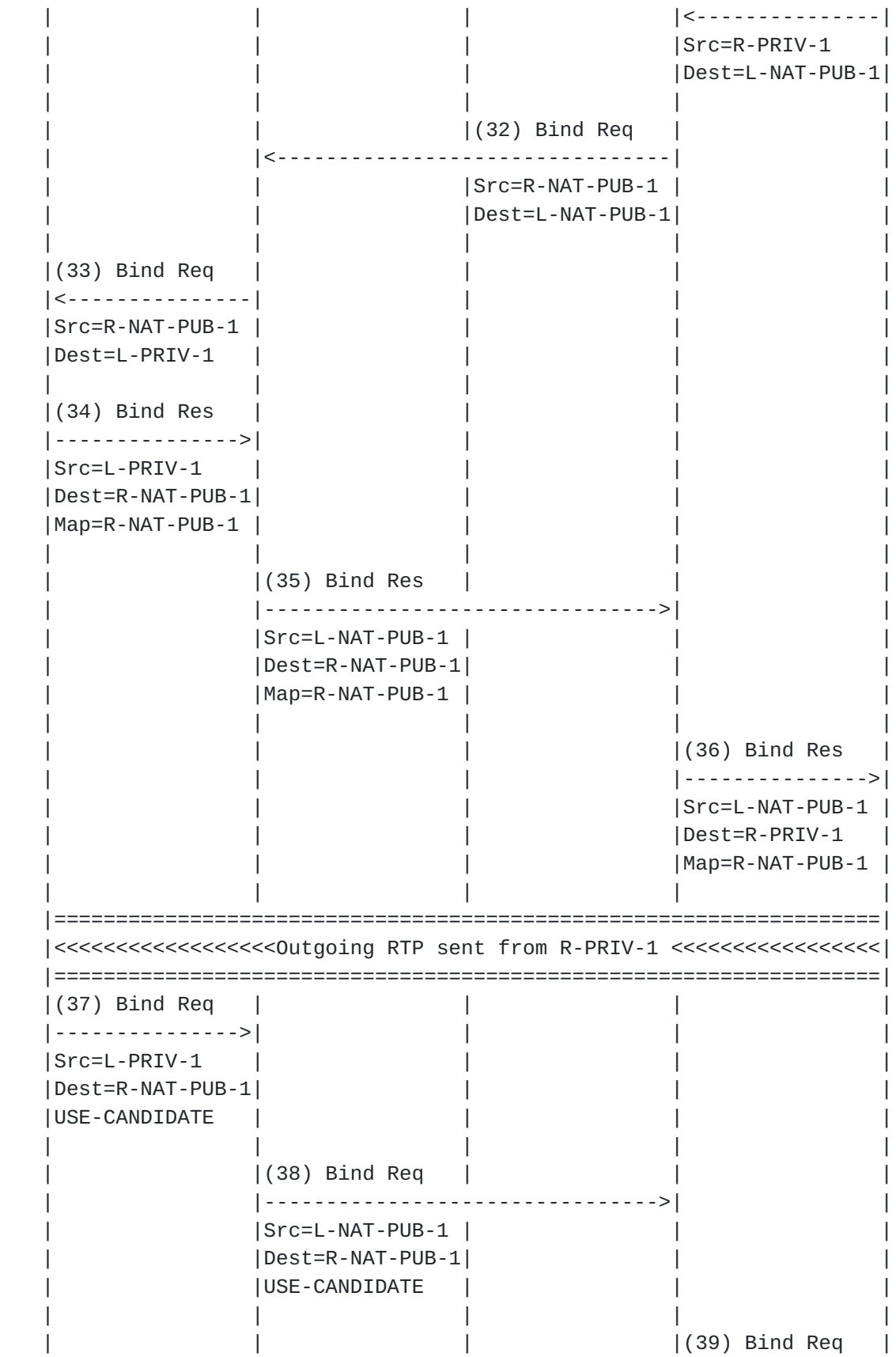
The following example demonstrates an initiating traversal through an 'Endpoint independent' NAT using ICE.

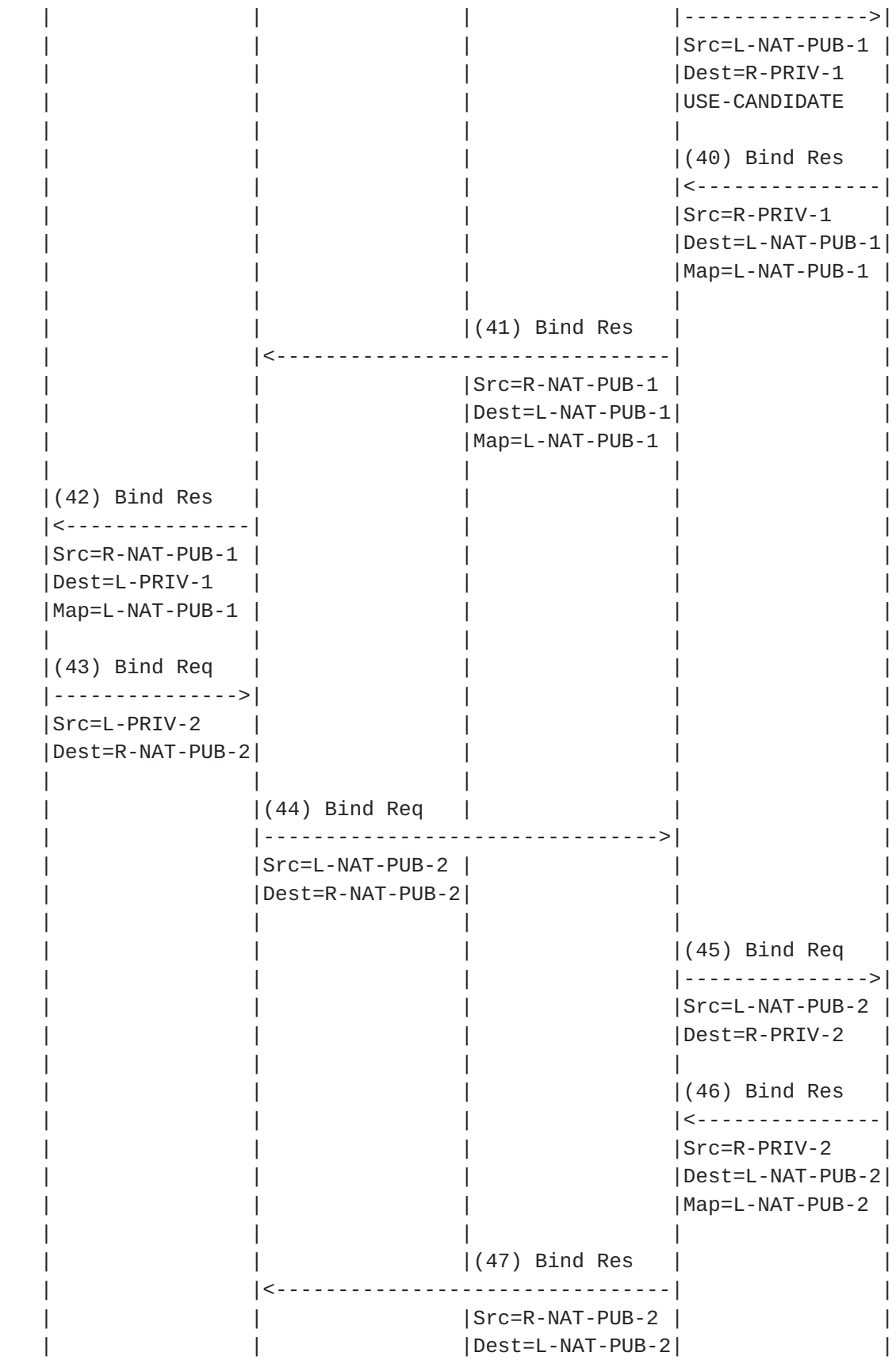


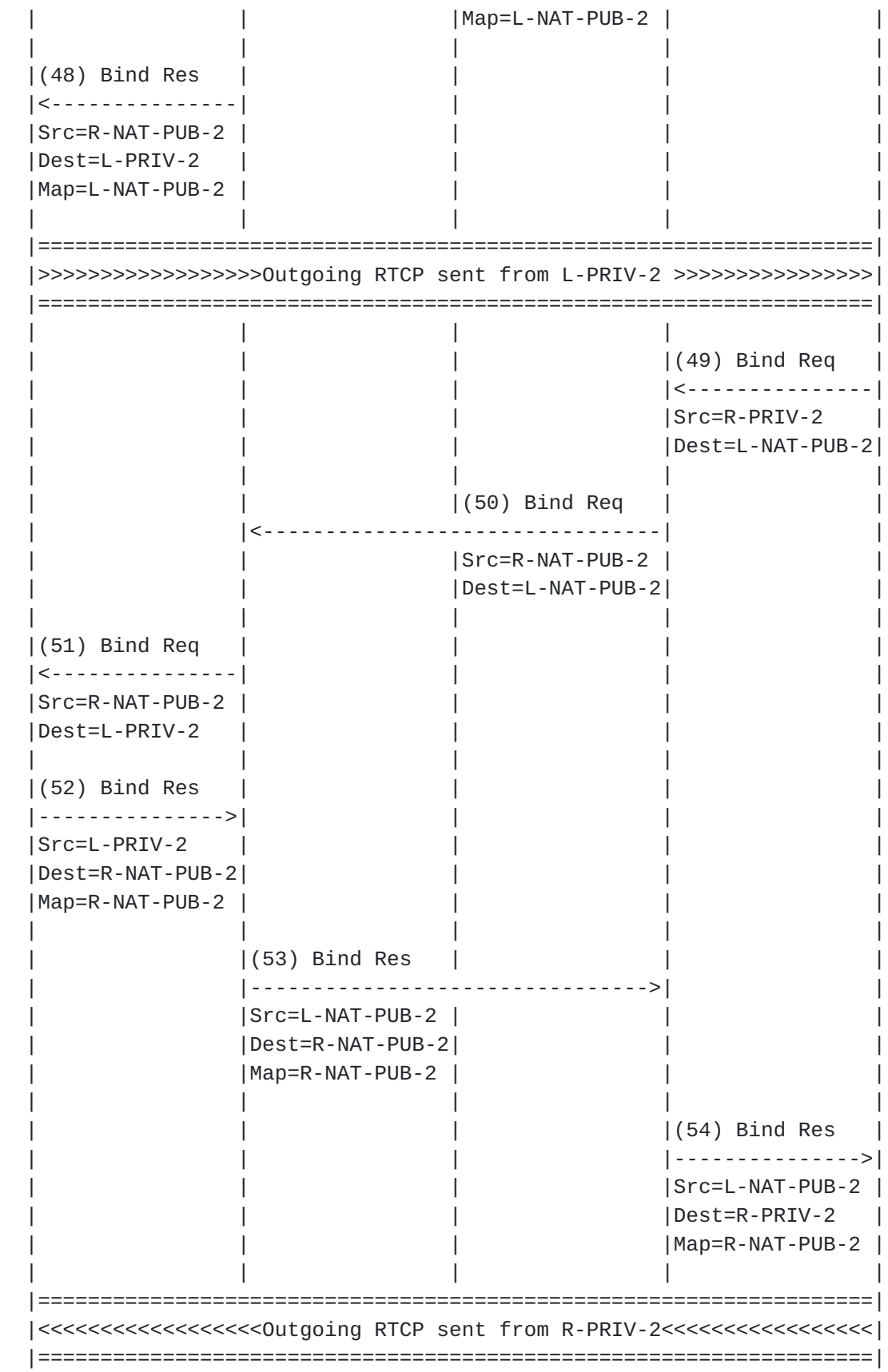

|<----------------|
|Src=TURN-PUB-1 |
|Dest=L-PRIV-1 |
|Map=L-NAT-PUB-1 |
|Rel=TURN-PUB-2 |
|(5) Alloc Req |
|Src=L-PRIV-2 |
|Dest=TURN-PUB-1 |
|---------------->|
|(6) Alloc Req |
|Src=L-NAT-PUB-2 |
|Dest=TURN-PUB-1 |
|---------------->|
|(7) Alloc Resp |
|<----------------|
|Src=TURN-PUB-1 |
|Dest=NAT-PUB-2 |
|Map=NAT-PUB-2 |
|Rel=TURN-PUB-3 |
|(8) Alloc Resp |
|<----------------|
|Src=TURN-PUB-1 |
|Dest=L-PRIV-2 |
|Map=L-NAT-PUB-2 |
|Rel=TURN-PUB-3 |
|(9) SIP INVITE |
|--->|
|(10) SIP INVITE |
|---------------->|
|(11) Alloc Req |
|<----------------|
|Src=R-PRIV-1 |
|Dest=TURN-PUB-1 |
|(12) Alloc Req |
|<----------------|
|Src=R-NAT-PUB-1 |
|Dest=TURN-PUB-1 |
|(13) Alloc Res |
|---------------->|

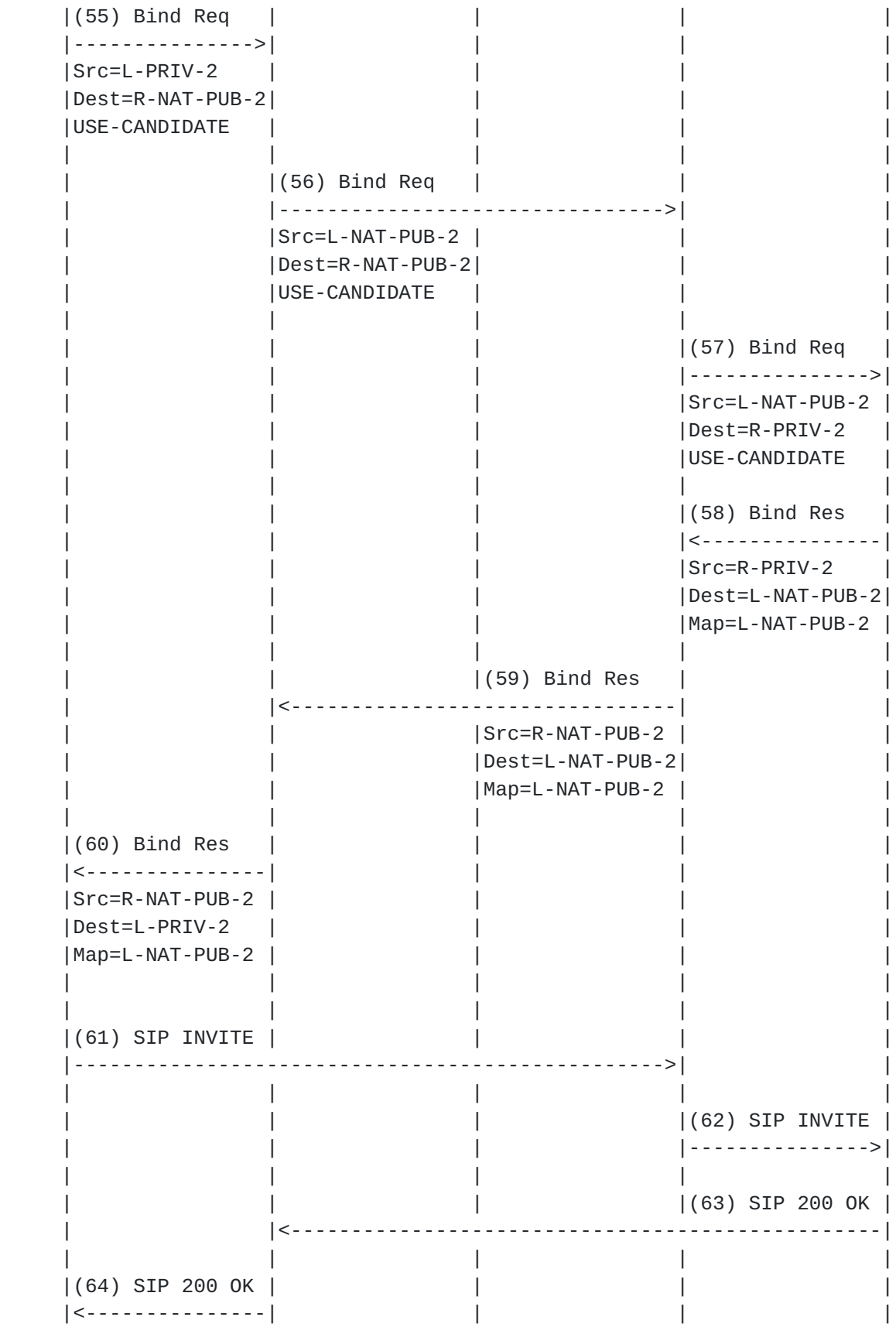

Src=TURN-PUB-1
Dest=R-NAT-PUB-1
Map=R-NAT-PUB-1
Rel=TURN-PUB-4
(14) Alloc Res
Src=TURN-PUB-1
Dest=R-PRIV-1
Map=R-NAT-PUB-1
Rel=TURN-PUB-4
(15) Alloc Req
Src=R-PRIV-2
Dest=TURN-PUB-1
(16) Alloc Req
Src=R-NAT-PUB-2
Dest=TURN-PUB-1
(17) Alloc Res
Src=TURN-PUB-1
Dest=R-NAT-PUB-2
Map=R-NAT-PUB-2
Rel=TURN-PUB-5
(18) Alloc Res
Src=TURN-PUB-1
Dest=R-PRIV-2
Map=R-NAT-PUB-2
Rel=TURN-PUB-5
(19) SIP 200 OK
(20) SIP 200 OK
(21) SIP ACK
(22) SIP ACK


(23) Bind Req
Src=L-PRIV-1
Dest=R-PRIV-1
(24) Bind Req
Src=L-PRIV-1
Dest=R-NAT-PUB-1
(25) Bind Req
Src=L-NAT-PUB-1
Dest=R-NAT-PUB-1
(26) Bind Req
Src=L-NAT-PUB-1
Dest=R-PRIV-1
(27) Bind Res
Src=R-PRIV-1
Dest=L-NAT-PUB-1
Map=L-NAT-PUB-1
(28) Bind Res
Src=R-NAT-PUB-1
Dest=L-NAT-PUB-1
Map=L-NAT-PUB-1
(29) Bind Res
Src=R-NAT-PUB-1
Dest=L-PRIV-1
Map=L-NAT-PUB-1
Outgoing RTP sent from L-PRIV-1
(30) Bind Req
Src=R-PRIV-1
Dest=L-PRIV-1
(31) Bind Req


|<----------------|
|Src=R-PRIV-1 |
|Dest=L-NAT-PUB-1|
|(32) Bind Req |
|<--------------------------------|
|Src=R-NAT-PUB-1 |
|Dest=L-NAT-PUB-1|
|(33) Bind Req |
|<----------------|
|Src=R-NAT-PUB-1 |
|Dest=L-PRIV-1 |
|(34) Bind Res |
|---------------->|
|Src=L-PRIV-1 |
|Dest=R-NAT-PUB-1|
|Map=R-NAT-PUB-1 |
|(35) Bind Res |
|-------------------------------->|
|Src=L-NAT-PUB-1 |
|Dest=R-NAT-PUB-1|
|Map=R-NAT-PUB-1 |
|(36) Bind Res |
|---------------->|
|Src=L-NAT-PUB-1 |
|Dest=R-PRIV-1 |
|Map=R-NAT-PUB-1 |
<<<<<<<<<<<<<<<<Outgoing RTP sent from R-PRIV-1 <<<<<<<<<<<<<<<<
|(37) Bind Req |
|---------------->|
|Src=L-PRIV-1 |
|Dest=R-NAT-PUB-1|
|USE-CANDIDATE |
|(38) Bind Req |
|-------------------------------->|
|Src=L-NAT-PUB-1 |
|Dest=R-NAT-PUB-1|
|USE-CANDIDATE |
|(39) Bind Req |


|---------->|
|Src=L-NAT-PUB-1 |
|Dest=R-PRIV-1 |
|USE-CANDIDATE |
|(40) Bind Res |
|<----------|
|Src=R-PRIV-1 |
|Dest=L-NAT-PUB-1|
|Map=L-NAT-PUB-1 |
|(41) Bind Res |
|<----------|
|Src=R-NAT-PUB-1 |
|Dest=L-NAT-PUB-1|
|Map=L-NAT-PUB-1 |
|(42) Bind Res |
|<----------|
|Src=R-NAT-PUB-1 |
|Dest=L-PRIV-1 |
|Map=L-NAT-PUB-1 |
|(43) Bind Req |
|---------->|
|Src=L-PRIV-2 |
|Dest=R-NAT-PUB-2|
|(44) Bind Req |
|---------->|
|Src=L-NAT-PUB-2 |
|Dest=R-NAT-PUB-2|
|(45) Bind Req |
|---------->|
|Src=L-NAT-PUB-2 |
|Dest=R-PRIV-2 |
|(46) Bind Res |
|<----------|
|Src=R-PRIV-2 |
|Dest=L-NAT-PUB-2|
|Map=L-NAT-PUB-2 |
|(47) Bind Res |
|<----------|
|Src=R-NAT-PUB-2 |
|Dest=L-NAT-PUB-2|


Map=L-NAT-PUB-2
(48) Bind Res
Src=R-NAT-PUB-2
Dest=L-PRIV-2
Map=L-NAT-PUB-2
Outgoing RTCP sent from L-PRIV-2
(49) Bind Req
Src=R-PRIV-2
Dest=L-NAT-PUB-2
(50) Bind Req
Src=R-NAT-PUB-2
Dest=L-NAT-PUB-2
(51) Bind Req
Src=R-NAT-PUB-2
Dest=L-PRIV-2
(52) Bind Res
Src=L-PRIV-2
Dest=R-NAT-PUB-2
Map=R-NAT-PUB-2
(53) Bind Res
Src=L-NAT-PUB-2
Dest=R-NAT-PUB-2
Map=R-NAT-PUB-2
(54) Bind Res
Src=L-NAT-PUB-2
Dest=R-PRIV-2
Map=R-NAT-PUB-2
Outgoing RTCP sent from R-PRIV-2


(55) Bind Req
Src=L-PRIV-2
Dest=R-NAT-PUB-2
USE-CANDIDATE
(56) Bind Req
Src=L-NAT-PUB-2
Dest=R-NAT-PUB-2
USE-CANDIDATE
(57) Bind Req
Src=L-NAT-PUB-2
Dest=R-PRIV-2
USE-CANDIDATE
(58) Bind Res
Src=R-PRIV-2
Dest=L-NAT-PUB-2
Map=L-NAT-PUB-2
(59) Bind Res
Src=R-NAT-PUB-2
Dest=L-NAT-PUB-2
Map=L-NAT-PUB-2
(60) Bind Res
Src=R-NAT-PUB-2
Dest=L-PRIV-2
Map=L-NAT-PUB-2
(61) SIP INVITE
(62) SIP INVITE
(63) SIP 200 OK
(64) SIP 200 OK

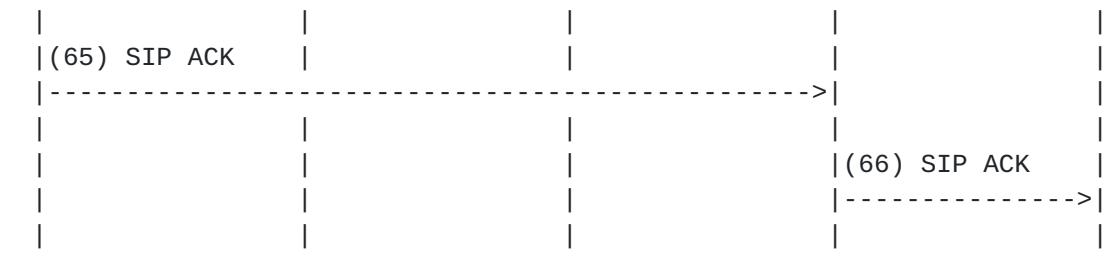


Figure 13: Endpoint Independent NAT with ICE

- On deciding to initiate a SIP voice session the SIP client 'L' starts a local STUN client. The STUN client generates a TURN Allocate request as indicated in (1) from Figure 13 which also highlights the source address and port combination for which the client device wishes to obtain a mapping. The Allocate request is sent through the NAT towards the public internet.

- The Allocate message (2) traverses the NAT to the public internet towards the public TURN server. Note that the source address of the Allocate request now represents the public address and port from the public side of the NAT (L-NAT-PUB-1).

- The TURN server receives the Allocate request and processes it appropriately. This results in a successful Allocate response being generated and returned (3). The message contains details of the server reflexive address which is to be used by the originating client to receive media (see 'Map=L-NAT-PUB-1') from (3)). It also contains an appropriate TURN-relayed address that can be used at the STUN server (see 'Rel=TURN-PUB-2').

- The Allocate response traverses back through the NAT using the binding created by the initial Allocate request and presents the new mapped address to the client (4). The process is repeated and a second STUN derived set of address' are obtained, as illustrated in (5)-(8) in Figure 13. At this point the User Agent behind the NAT has pairs of derived external server reflexive and relayed representations. The client can also gather IP addresses and ports via other mechanisms (e.g., NAT-PMP[I-D.cheshire-nat-pmp], UPnP IGD[UPnP-IGD]) or similar.

- The client now constructs a SIP INVITE message (9). The INVITE request will use the addresses it has obtained in the previous STUN/TURN interactions to populate the SDP of the SIP INVITE. This should be carried out in accordance with the semantics defined in the ICE specification[RFC5245], as shown below in Figure 14:

```
v=0
o=test 2890844526 2890842807 IN IP4 $L-PRIV-1
c=IN IP4 $L-PRIV-1.address
t=0 0
a=ice-pwd:$LPASS
a=ice-ufrag:$LUNAME
m=audio $L-PRIV-1.port RTP/AVP 0
a=rtpmap:0 PCMU/8000
a=rtcp:$L-PRIV-2.port
a=candidate:$L1 1 UDP 2130706431 $L-PRIV-1.address $L-PRIV-1.port
            typ host
a=candidate:$L1 2 UDP 2130706430 $L-PRIV-2.address $L-PRIV-2.port
            typ host
a=candidate:$L2 1 UDP 1694498815 $L-NAT-PUB-1.address $L-NAT-PUB-1.port
            typ srflx raddr $L-PRIV-1.address rport $L-PRIV-1.port
a=candidate:$L2 2 UDP 1694498814 $L-NAT-PUB-2.address $L-NAT-PUB-2.port
            typ srflx raddr $L-PRIV-1.address rport $L-PRIV-2.port
a=candidate:$L3 1 UDP 16777215 $STUN-PUB-2.address $STUN-PUB-2.port
            typ relay raddr $L-PRIV-1.address rport $L-PRIV-1.port
a=candidate:$L3 2 UDP 16777214 $STUN-PUB-3.address $STUN-PUB-3.port
            typ relay raddr $L-PRIV-1.address rport $L-PRIV-2.port
```

Figure 14: ICE SDP Offer

- The SDP has been constructed to include all the available candidates that have been assembled. The first set of candidates (as identified by Foundation $L1) contain two local addresses that have the highest priority. They are also encoded into the connection (c=) and media (m=) lines of the SDP. The second set of candidates, as identified by Foundation $L2, contains the two server reflexive addresses obtained from the STUN server for both RTP and RTCP traffic (identified by candidate-id $L2). This entry has been given a priority lower than the pair $L1 by the client. The third and final set of candidates represents the relayed addresses (as identified by $L3) obtained from the STUN server. This pair has the lowest priority and will be used as a last resort if both $L1 or $L2 fail.

- The SIP signaling then traverses the NAT and sets up the SIP session (9)-(10). On advertising a candidate address, the client should have a local STUN server running on each advertised candidate address. This is for the purpose of responding to incoming STUN connectivity checks.

- On receiving the SIP INVITE request (10) client 'R' also starts local STUN servers on appropriate address/port combinations and gathers potential candidate addresses to be encoded into the SDP

(as the originating client did). Steps (11-18) involve client 'R' carrying out the same steps as client 'L'. This involves obtaining local, server reflexive and relayed addresses. Client 'R' is now ready to generate an appropriate answer in the SIP 200 OK message (19). The example answer follows in Figure 14:

```
v=0
o=test 3890844516 3890842803 IN IP4 $R-PRIV-1
c=IN IP4 $R-PRIV-1.address
t=0 0
a=ice-pwd:$RPASS
m=audio $R-PRIV-1.port RTP/AVP 0
a=rtpmap:0 PCMU/8000
a=rtcp:$R-PRIV-2.port
a=candidate:$L1 1 UDP 2130706431 $R-PRIV-1.address $R-PRIV-1.port
            typ host
a=candidate:$L1 2 UDP 2130706430 $R-PRIV-2.address $R-PRIV-2.port
            typ host
a=candidate:$L2 1 UDP 1694498815 $R-NAT-PUB-1.address $R-NAT-PUB-1.port
            typ srflx raddr $R-PRIV-1.address rport $R-PRIV-1.port
a=candidate:$L2 2 UDP 1694498814 $R-NAT-PUB-2.address $R-NAT-PUB-2.port
            typ srflx raddr $R-PRIV-1.address rport $R-PRIV-1.port
a=candidate:$L3 1 UDP 16777215 $STUN-PUB-2.address $STUN-PUB-4.port
            typ relay raddr $R-PRIV-1.address rport $R-PRIV-1.port
a=candidate:$L3 2 UDP 16777214 $STUN-PUB-3.address $STUN-PUB-5.port
            typ relay raddr $R-PRIV-1.address rport $R-PRIV-1.port
```

Figure 15: ICE SDP Answer

- The two clients have now exchanged SDP using offer/answer and can now continue with the ICE processing - User Agent 'L' assuming the role controlling agent, as specified by ICE. The clients are now required to form their Candidate check lists to determine which will be used for the media streams. In this example User Agent 'L's 'Foundation 1' is paired with User Agent 'R's 'Foundation 1', User Agent 'L's 'Foundation 2' is paired with User Agent 'R's 'Foundation 2', and finally User Agent 'L's 'Foundation 3' is paired with User Agent 'R's 'Foundation 3'. User Agents 'L' and 'R' now have a complete candidate check list. Both clients now use the algorithm provided in ICE to determine candidate pair priorities and sort into a list of decreasing priorities. In this example, both User Agent 'L' and 'R' will have lists that firstly specifies the host address (Foundation $L1), then the server reflexive address (Foundation $L2) and lastly the relayed address (Foundation $L3). All candidate pairs have an associate state as specified in ICE. At this stage, all of the candidate pairs for

User Agents 'L' and 'R' are initialized to the 'Frozen' state. The User Agents then scan the list and move the candidates to the 'Waiting' state. At this point both clients will periodically, starting with the highest candidate pair priority, work their way down the list issuing STUN checks from the local candidate to the remote candidate (of the candidate pair). As a STUN Check is attempted from each local candidate in the list, the candidate pair state transitions to 'In-Progress'. As illustrated in (23), client 'L' constructs a STUN connectivity check in an attempt to validate the remote candidate address received in the SDP of the 200 OK (20) for the highest priority in the check list. As a private address was specified in the active address in the SDP, the STUN connectivity check fails to reach its destination causing a STUN failure. Client 'L' transitions the state for this candidate pair to 'Failed'. In the mean time, Client 'L' is attempting a STUN connectivity check for the second candidate pair in the returned SDP with the second highest priority (24). As can be seen from messages (24) to (29), the STUN Bind request is successful and returns a positive outcome for the connectivity check. Client 'L' is now free to send media to the peer using the candidate pair. Immediately after sending its 200 Okay, Client 'R' also carries out the same set of binding requests. It firstly (in parallel) tries to contact the active address contained in the SDP (30) which results in failure.

- In the mean time, a successful response to a STUN connectivity check by User Agent 'R' (27) results in a tentative check in the reverse direction - this is illustrated by messages (31) to (36). Once this check has succeeded, User Agent 'R' can transition the state of the appropriate candidate to 'Succeeded', and media can be sent (RTP). The previously (31-36) described check confirm on both sides (User Agent 'L' and 'R') that connectivity can be achieved using the appropriate candidate pair. User Agent 'L', as the controlling client now sends another connectivity check for the candidate pair, this time including the 'USE-CANDIDATE' attribute as specified in ICE to signal the chosen candidate. This exchange is illustrated in messages (37) to (42).

- As part of the process in this example, both 'L' and 'R' will now complete the same connectivity checks for part 2 of the component named for the favored 'Foundation' selected for use with RTCP. The connectivity checks for part '2' of the candidate component are shown in 'L'(43-48) and 'R'(49-54). Once this has succeeded, User Agent 'L' as the controlling client sends another connectivity check for the candidate pair. This time the 'USE-CANDIDATE' attribute is again specified to signal the chosen candidate for component '2'.

- The candidates have now been fully verified (and selected) and as they are the highest priority, an updated offer (61-62) is now sent from the offerer (client 'L') to the answerer (client 'R') representing the new active candidates. The new offer would look as follows:

```
v=0
o=test 2890844526 2890842808 IN IP4 $L-PRIV-1
c=IN IP4 $L-NAT-PUB-1.address
t=0 0
a=ice-pwd:$LPASS
a=ice-ufrag:$LUNAME
m=audio $L-NAT-PUB-1.port RTP/AVP 0
a=rtpmap:0 PCMU/8000
a=rtcp:$L-NAT-PUB-2.port
a=candidate:$L2 1 UDP 2203948363 $L-NAT-PUB-1.address $L-NAT-PUB-1.port
            typ srflx raddr $L-PRIV-1.address rport $L-PRIV-1.port
a=candidate:$L2 2 UDP 2172635342 $L-NAT-PUB-2.address $L-NAT-PUB-2.port
            typ srflx raddr $L-PRIV-1.address rport $L-PRIV-2.port
```

Figure 16: ICE SDP Updated Offer

- The resulting answer (63-64) for 'R' would look as follows:

```
v=0
o=test 3890844516 3890842804 IN IP4 $R-PRIV-1
c=IN IP4 $R-PRIV-1.address
t=0 0
a=ice-pwd:$RPASS
a=ice-ufrag:$RUNAME
m=audio $R-PRIV-1.port RTP/AVP 0
a=rtpmap:0 PCMU/8000
a=rtcp:$R-PRIV-2.port
a=candidate:$L2 1 UDP 2984756463 $R-NAT-PUB-1.address $R-NAT-PUB-1.port
            typ srflx raddr $R-PRIV-1.address rport $R-PRIV-1.port
a=candidate:$L2 2 UDP 2605968473 $R-NAT-PUB-2.address $R-NAT-PUB-2.port
            typ srflx raddr $R-PRIV-1.address rport $R-PRIV-2.port
```

Figure 17: ICE SDP Updated Answer

5.2.2. Address/Port-Dependent NAT

5.2.2.1. STUN Failure

This section highlights that while using STUN techniques is the preferred mechanism for traversal of NAT, it does not solve every case. The use of basic STUN on its own will not guarantee traversal through every NAT type, hence the recommendation that ICE is the preferred option.

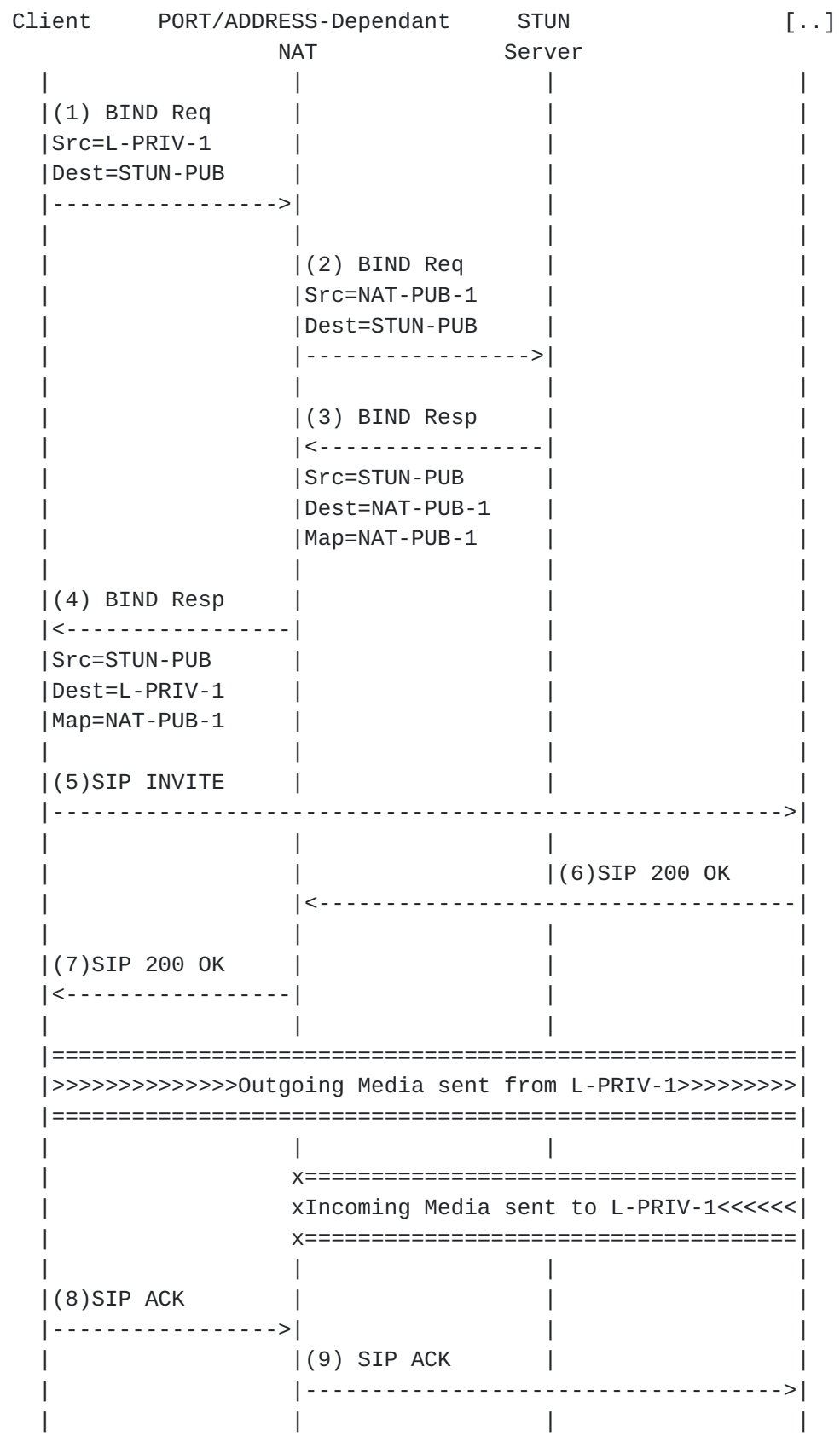

Client
PORT/ADDRESS-Dependant NAT
STUN Server
[..]
(1) BIND Req
Src=L-PRIV-1
Dest=STUN-PUB
(2) BIND Req
Src=NAT-PUB-1
Dest=STUN-PUB
(3) BIND Resp
Src=STUN-PUB
Dest=NAT-PUB-1
Map=NAT-PUB-1
(4) BIND Resp
Src=STUN-PUB
Dest=L-PRIV-1
Map=NAT-PUB-1
(5)SIP INVITE
(6)SIP 200 OK
(7)SIP 200 OK
>>>>>>>>>>>>>Outgoing Media sent from L-PRIV-1>>>>>>>>>
xIncoming Media sent to L-PRIV-1<<<<<<
(8)SIP ACK
(9) SIP ACK

Figure 18: Port/Address-Dependant NAT with STUN - Failure

The example in Figure 18 is conveyed in the context of a client behind the 'Port/Address-Dependant' NAT initiating a call. It should be noted that the same problem applies when a client receives a SIP invitation and is behind a Port/Address-Dependant NAT.

- In Figure 18 the client behind the NAT obtains a server reflexive representation using standard STUN mechanisms (1)-(4) that have been used in previous examples in this document (e.g Section 5.2.1.1.1).

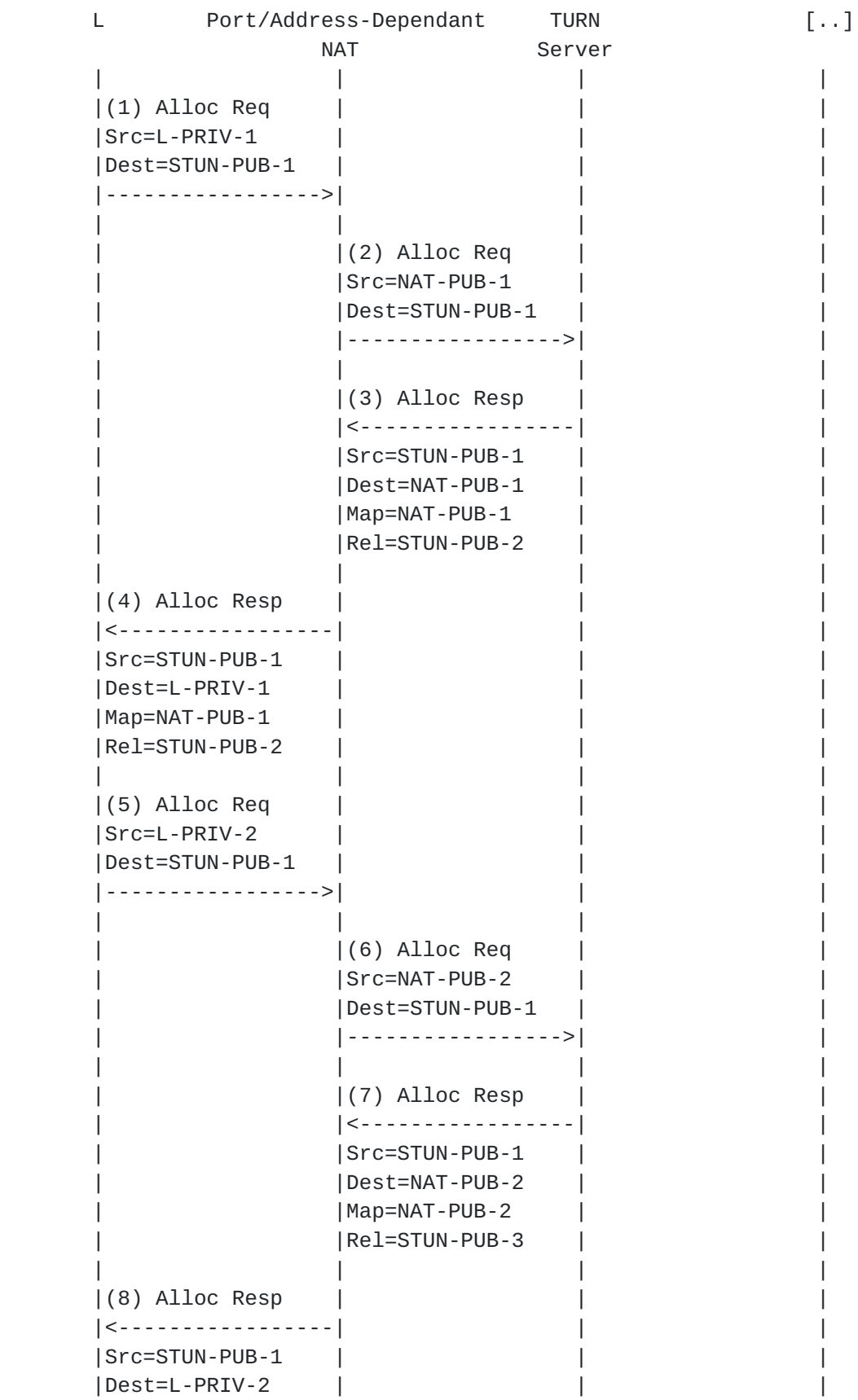
- The external mapped address (server reflexive) obtained is also used in the outgoing SDP contained in the SIP INVITE request(5).

- In this example the client is still able to send media to the external client. The problem occurs when the client outside the NAT tries to use the reflexive address supplied in the outgoing INVITE request to traverse media back through the 'Port/Address Dependent' NAT.

- A 'Port/Address-Dependant' NAT has differing rules from the 'Endpoint Independent' type of NAT (as defined in RFC4787 [RFC4787]). For any internal IP address and port combination, data sent to a different external destination does not provide the same public mapping at the NAT. In Figure 18 the STUN query produced a valid external mapping for receiving media. This mapping, however, can only be used in the context of the original STUN request that was sent to the STUN server. Any packets that attempt to use the mapped address, that do not originate from the STUN server IP address and optionally port, will be dropped at the NAT. Figure 18 shows the media being dropped at the NAT after (7) and before (8). This then leads to one way audio.

5.2.2.2. TURN Solution

As identified in Section 5.2.2.1, STUN provides a useful tool for the traversal of the majority of NATs but fails with Port/Address Dependent NAT. The TURN extensions [RFC5766] address this scenario. TURN extends STUN to allow a client to request a relayed address at the TURN server rather than a reflexive representation. This then introduces a media relay in the path for NAT traversal (as described in Section 4.2.3.2). The following example explains how TURN solves the previous failure when using STUN to traverse a 'Port/Address Dependent' type NAT. It should be noted that TURN works most effectively when used in conjunction with ICE. Using TURN on its own results in all media being relayed through a TURN server which is not efficient.


L
Port/Address-Dependant NAT
TURN Server
[..]
(1) Alloc Req
Src=L-PRIV-1
Dest=STUN-PUB-1
(2) Alloc Req
Src=NAT-PUB-1
Dest=STUN-PUB-1
(3) Alloc Resp
Src=STUN-PUB-1
Dest=NAT-PUB-1
Map=NAT-PUB-1
Rel=STUN-PUB-2
(4) Alloc Resp
Src=STUN-PUB-1
Dest=L-PRIV-1
Map=NAT-PUB-1
Rel=STUN-PUB-2
(5) Alloc Req
Src=L-PRIV-2
Dest=STUN-PUB-1
(6) Alloc Req
Src=NAT-PUB-2
Dest=STUN-PUB-1
(7) Alloc Resp
Src=STUN-PUB-1
Dest=NAT-PUB-2
Map=NAT-PUB-2
Rel=STUN-PUB-3
(8) Alloc Resp
Src=STUN-PUB-1
Dest=L-PRIV-2

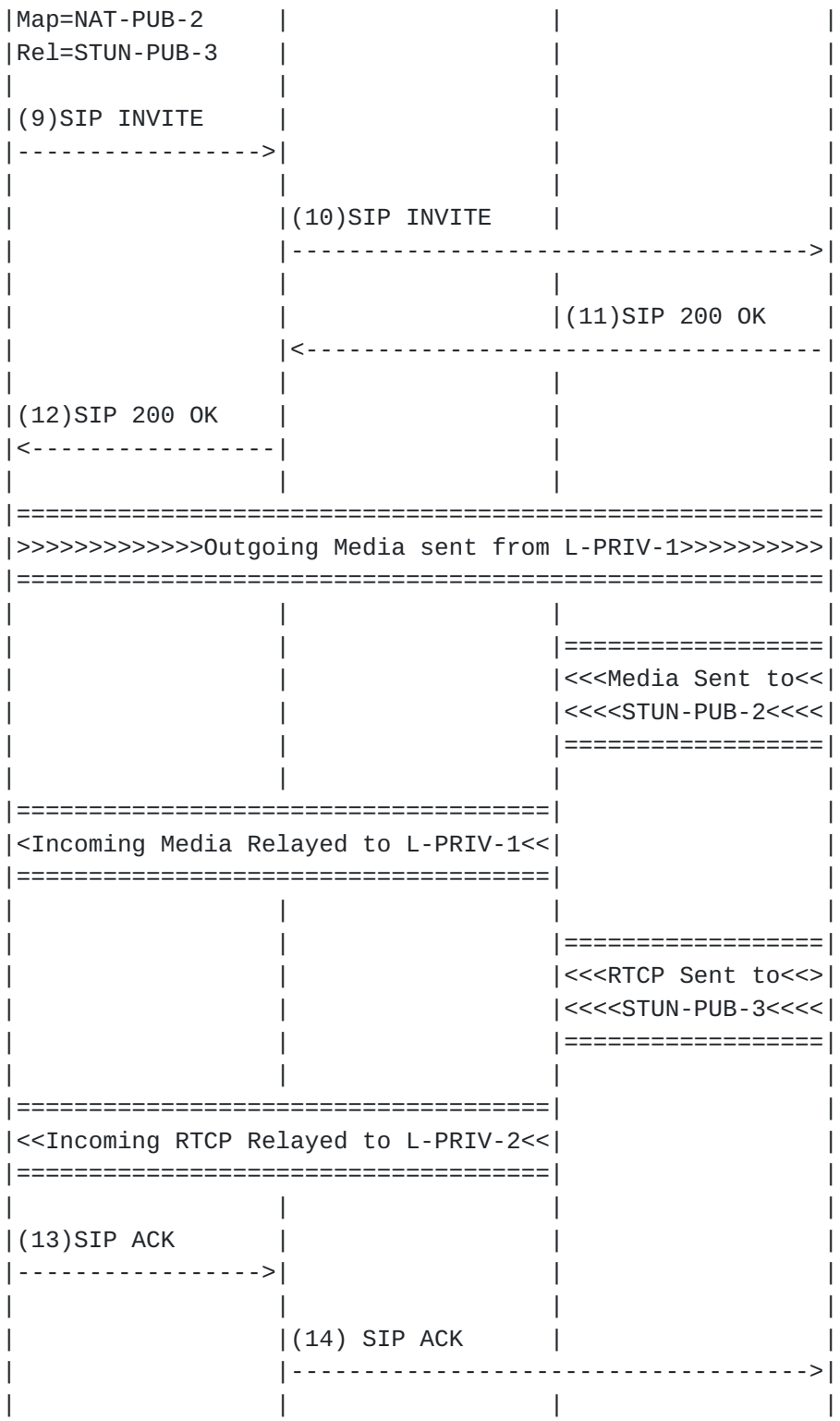


Figure 19: Port/Address-Dependant NAT with TURN - Success

- In this example, client 'L' issues a TURN allocate request(1) to obtained a relay address at the STUN server. The request traverses through the 'Port/Address-Dependant' NAT and reaches the STUN server (2). The STUN server generates an Allocate response (3) that contains both a server reflexive address (Map=NAT-PUB-1) of the client and also a relayed address (Rel=STUN-PUB-2). The relayed address maps to an address mapping on the STUN server which is bound to the public pin hole that has been opened on the NAT by the Allocate request. This results in any traffic sent to the TURN server relayed address (Rel=STUN-PUB-2) being forwarded to the external representation of the pin hole created by the Allocate request(NAT-PUB-1).

- The TURN derived address (STUN-PUB-2) arrives back at the originating client (4) in an Allocate response. This address can then be used in the SDP for the outgoing SIP INVITE request as shown in the following example (note that the example also includes client 'L' obtaining a second relay address for use in the RTCP attribute (5-8)):

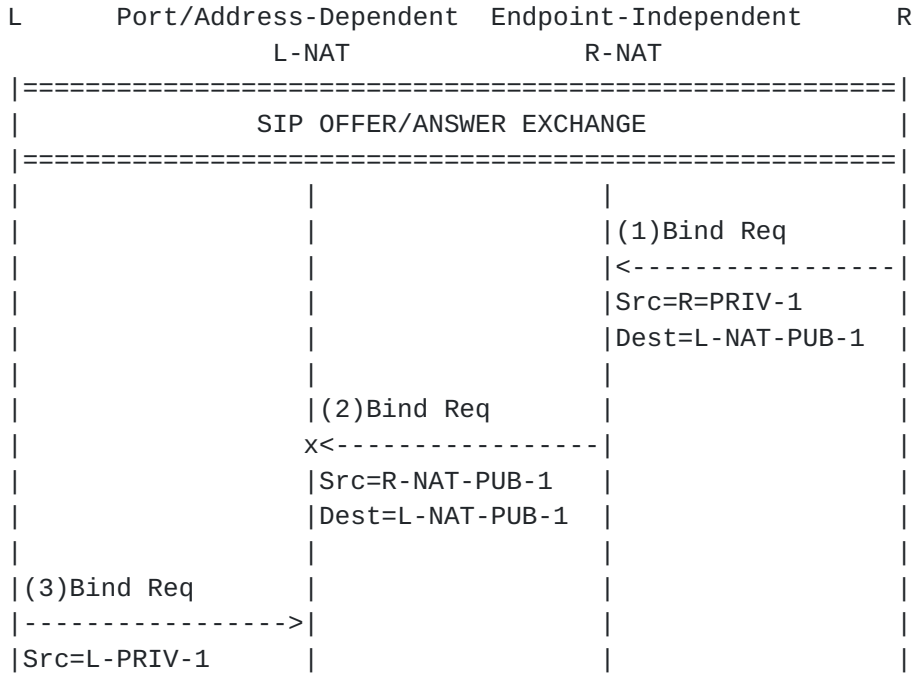
```
v=0
o=test 2890844342 2890842164 IN IP4 $L-PRIV-1
c=IN IP4 $STUN-PUB-2.address
t=0 0
m=audio $STUN-PUB-2.port RTP/AVP 0
a=rtcp:$STUN-PUB-3.port
```

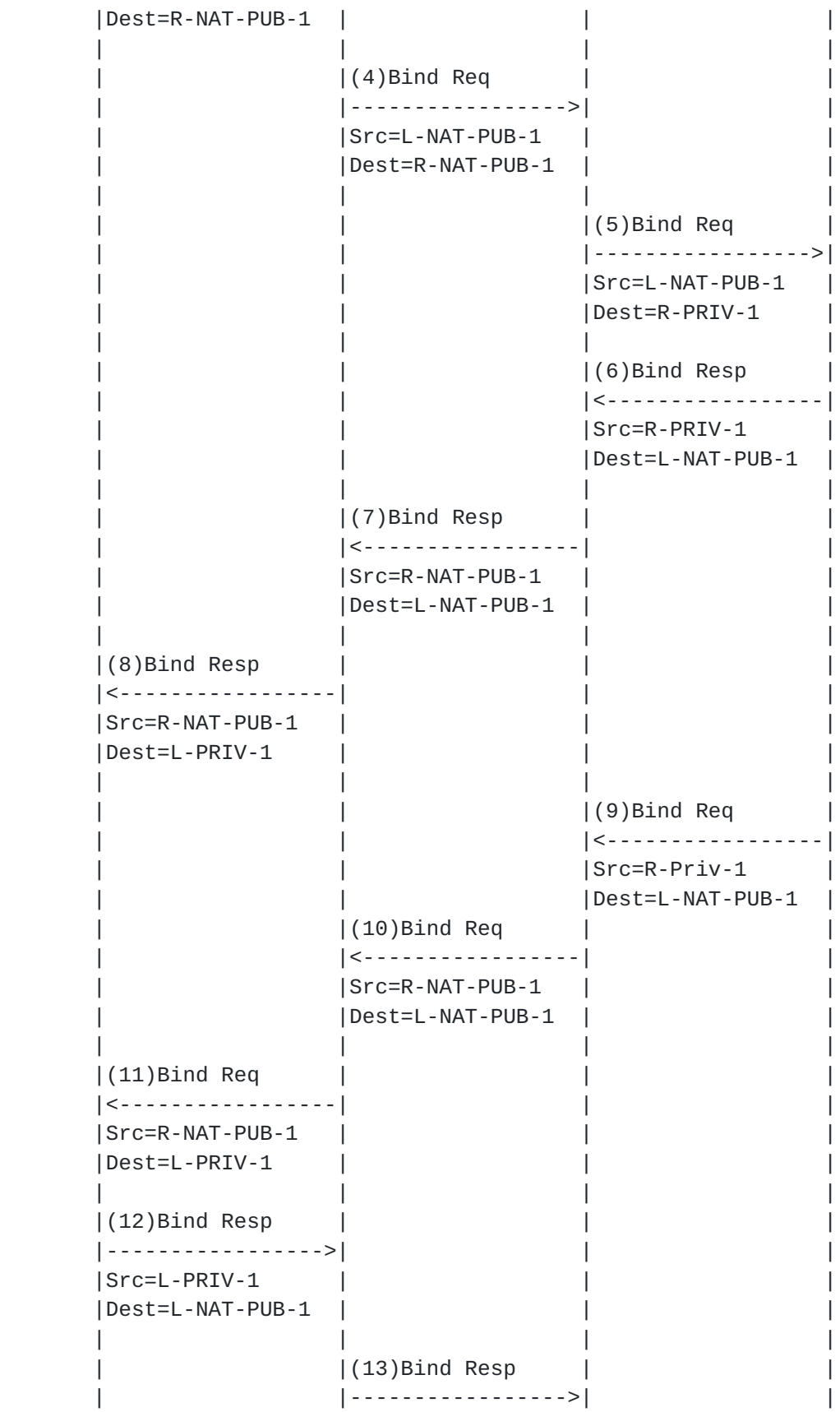
- On receiving the INVITE request, the UAS is able to stream media and RTCP to the relay address (STUN-PUB-2 and STUN-PUB-3) at the STUN server. As shown in Figure 19 (between messages (12) and (13), the media from the UAS is directed to the relayed address at the STUN server. The STUN server then forwards the traffic to the open pin holes in the Port/Address-Dependant NAT (NAT-PUB-1 and NAT-PUB-2). The media traffic is then able to traverse the 'Port/Address-Dependant' NAT and arrives back at client 'L'.

- TURN on its own will work for 'Port/Address-Dependent' and other types of NAT mentioned in this specification but should only be used as a last resort. The relaying of media through an external entity is not an efficient mechanism for NAT traversal and comes at a high processing cost.

5.2.2.3. ICE Solution

The previous two examples have highlighted the problem with using core STUN for all forms of NAT traversal and a solution using TURN for the Address/Port-Dependent NAT case. The RECOMMENDED mechanism for traversing all varieties of NAT is using ICE, as detailed in Section 4.2.3.3. ICE makes use of core STUN, TURN and any other mechanism (e.g., NAT-PMP[I-D.cheshire-nat-pmp], UPnP IGD[UPnP-IGD]) to provide a list of prioritized addresses that can be used for media traffic. Detailed examples of ICE can be found in Section 5.2.1.2.1. These examples are associated with an 'Endpoint-Independent' type NAT but can be applied to any NAT type variation, including 'Address/Port-Dependant' type NAT. The ICE procedures carried out are the same. For a list of candidate addresses, a client will choose where to send media dependant on the results of the STUN connectivity checks and associated priority (highest priority wins). It should be noted that the inclusion of a NAT displaying Address/Port-Dependent properties does not automatically result in relayed media. In fact, ICE processing will avoid use of media relay with the exception of two clients which both happen to be behind a NAT using Address/Port-Dependent characteristics. The connectivity checks and associated selection algorithm enable traversal in this case. Figure 20 and following description provide a guide as to how this is achieved using the ICE connectivity checks. This is an abbreviated example that assumes successful SIP offer/answer exchange and illustrates the connectivity check flow.




Dest=R-NAT-PUB-1
(4)Bind Req
Src=L-NAT-PUB-1
Dest=R-NAT-PUB-1
(5)Bind Req
Src=L-NAT-PUB-1
Dest=R-PRIV-1
(6)Bind Resp
Src=R-PRIV-1
Dest=L-NAT-PUB-1
(7)Bind Resp
Src=R-NAT-PUB-1
Dest=L-NAT-PUB-1
(8)Bind Resp
Src=R-NAT-PUB-1
Dest=L-PRIV-1
(9)Bind Req
Src=R-Priv-1
Dest=L-NAT-PUB-1
(10)Bind Req
Src=R-NAT-PUB-1
Dest=L-NAT-PUB-1
(11)Bind Req
Src=R-NAT-PUB-1
Dest=L-PRIV-1
(12)Bind Resp
Src=L-PRIV-1
Dest=L-NAT-PUB-1
(13)Bind Resp

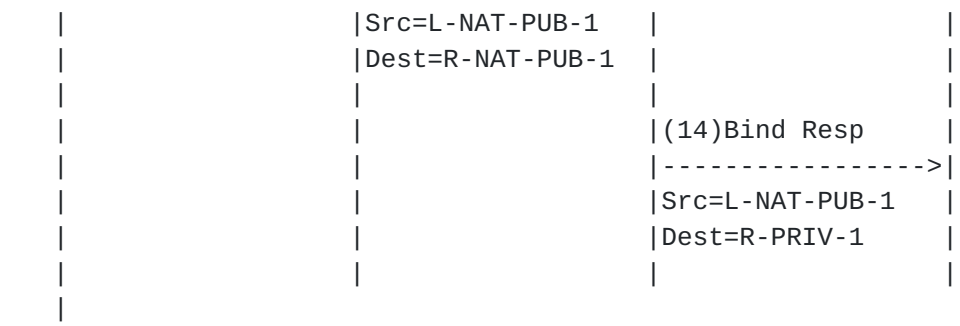


Figure 20: Single Port/Address-Dependant NAT - Success

In this abbreviated example, Client R has already received a SIP INVITE request and is starting its connectivity checks with Client L. Client R generates a connectivity check (1) and sends to client L's information as presented in the SDP offer. The request arrives at client L's Port/Address dependent NAT and fails to traverse as there is no NAT binding. This would then move the connectivity check to a failed state. In the mean time client L has received the SDP answer in the SIP request and will also commence connectivity checks. A check is dispatched (3) to Client R. The check is able to traverse the NAT due to the association set up in the previously failed check(1). The full Bind request/response is shown in steps (3)-(8). As part of a candidate pair, Client R will now successfully be able to complete the checks, as illustrated in steps (9)-(14). The result is a successful pair of candidates that can be used without the need to relay any media.

In conclusion, the only time media needs to be relayed is a result of clients both behind Address/Port Dependant NAT type. As you can see from the example in this section, neither side would be able to complete connectivity checks with the exception of the Relayed candidates.

6. IPv4-IPv6 Transition

This section describes how IPv6-only SIP user agents can communicate with IPv4-only SIP user agents. While the techniques discussed in this draft primarily contain examples of traversing NATs to allow communications between hosts in private and public networks, they are by no means limited to such scenarios. The same NAT traversal techniques can also be used to establish communication in a heterogeneous network environment -- e.g., communication between an IPv4 host and an IPv6 host.

6.1. IPv4-IPv6 Transition for SIP Signaling

IPv4-IPv6 translations at the SIP level usually take place at dual-stack proxies that have both IPv4 and IPv6 DNS entries. Since this translations do not involve NATs that are placed in the middle of two SIP entities, they fall outside the scope of this document. A detailed description of this type of translation can be found in [I-D.ietf-sipping-v6-transition]

7. Security Considerations

There are no Security Considerations beyond the ones inherited by reference.

8. IANA Considerations

There are no IANA Considerations.

9. IAB Considerations

There are no IAB considerations.

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