

Network Working Group
Internet-Draft
Intended status: Experimental
Expires: June 14, 2018

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December 11, 2017

TCP Alternative Backoff with ECN (ABE)
draft-ietf-tcpm-alternativebackoff-ecn-05

Abstract

Recent Active Queue Management (AQM) mechanisms allow for burst tolerance while enforcing short queues to minimise the time that packets spend enqueued at a bottleneck. This can cause noticeable performance degradation for TCP connections traversing such a bottleneck, especially if there are only a few flows or their bandwidth-delay-product is large. An Explicit Congestion Notification (ECN) signal indicates that an AQM mechanism is used at the bottleneck, and therefore the bottleneck network queue is likely to be short. This document therefore proposes an update to [RFC3168](#), which changes the TCP sender-side ECN reaction in congestion avoidance to reduce the Congestion Window (cwnd) by a smaller amount than the congestion control algorithm's reaction to inferred packet loss.

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[1.](#) Definitions

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC 2119](#) [[RFC2119](#)].

[2.](#) Introduction

Explicit Congestion Notification (ECN) [[RFC3168](#)] makes it possible for an Active Queue Management (AQM) mechanism to signal the presence of incipient congestion without incurring packet loss. This lets the network deliver some packets to an application that would have been dropped if the application or transport did not support ECN. This

packet loss reduction is the most obvious benefit of ECN, but it is often relatively modest. Other benefits of deploying ECN have been documented in [RFC8087](#) [[RFC8087](#)].

The rules for ECN were originally written to be very conservative, and required the congestion control algorithms of ECN-Capable transport protocols to treat ECN congestion signals exactly the same as they would treat an inferred packet loss [[RFC3168](#)].

Research has demonstrated the benefits of reducing network delays that are caused by interaction of loss-based TCP congestion control and excessive buffering [[BUFFERBLOAT](#)]. This has led to the creation of new AQM mechanisms like PIE [[RFC8033](#)] and CoDel [[CODEL2012](#)][I-D.CoDel], which prevent bloated queues that are common with unmanaged and excessively large buffers deployed across the Internet [[BUFFERBLOAT](#)].

The AQM mechanisms mentioned above aim to keep a sustained queue short while tolerating transient (short-term) packet bursts. However, currently used loss-based congestion control mechanisms cannot always utilise a bottleneck link well where there are short queues. For example, a TCP sender must be able to store at least an end-to-end bandwidth-delay product (BDP) worth of data at the bottleneck buffer if it is to maintain full path utilisation in the face of loss-induced reduction of cwnd [[RFC5681](#)], which effectively doubles the amount of data that can be in flight, the maximum round-trip time (RTT) experience, and the path's effective RTT using the network path.

Modern AQM mechanisms can use ECN to signal the early signs of impending queue buildup long before a tail-drop queue would be forced to resort to dropping packets. It is therefore appropriate for the transport protocol congestion control algorithm to have a more measured response when an early-warning signal of congestion is received in the form of an ECN CE-marked packet. Recognizing these changes in modern AQM practices, more recent rules have relaxed the strict requirement that ECN signals be treated identically to inferred packet loss [[I-D.ECN-exp](#)]. Following these newer, more flexible rules, this document defines a new sender-side-only congestion control response, called "ABE" (Alternative Backoff with ECN). ABE improves TCP's average throughput when routers use AQM controlled buffers that allow for short queues only.

3. Specification

This specification updates the congestion control algorithm of an ECN-Capable TCP transport protocol by changing the TCP sender response to feedback from the TCP receiver that indicates reception

of a CE-marked packet, i.e., receipt of a packet with the ECN-Echo flag (defined in [\[RFC3168\]](#)) set.

It updates the following text in [section 6.1.2](#) of the ECN specification [\[RFC3168\]](#) :

The indication of congestion should be treated just as a congestion loss in non-ECN-Capable TCP. That is, the TCP source halves the congestion window "cwnd" and reduces the slow start threshold "ssthresh".

Replacing this with:

Receipt of a packet with the ECN-Echo flag SHOULD trigger the TCP source to set the slow start threshold (ssthresh) to 0.8 times the FlightSize, with a lower bound of $2 * SMSS$ applied to the result. As in [\[RFC5681\]](#), the TCP sender also reduces the cwnd value to that new ssthresh value.

4. Discussion

Much of the technical background to ABE can be found in a research paper [\[ABE2017\]](#). This paper used a mix of experiments, theory and simulations with NewReno [\[RFC5681\]](#) and CUBIC [\[I-D.CUBIC\]](#) to evaluate the technique. The technique was shown to present "...significant performance gains in lightly-multiplexed [few concurrent flows] scenarios, without losing the delay-reduction benefits of deploying CoDel or PIE". The performance improvement is achieved when reacting to ECN-Echo in congestion avoidance by multiplying cwnd and ssthresh with a value in the range $[0.7, 0.85]$.

4.1. Why Use ECN to Vary the Degree of Backoff?

The classic rule-of-thumb dictates that a network path needs to provide a BDP of bottleneck buffering if a TCP connection wishes to optimise path utilisation. A single TCP bulk transfer running through such a bottleneck will have increased its congestion window (cwnd) up to $2 * BDP$ by the time that packet loss occurs. When packet loss is inferred using the retransmission timer and the given packet has not yet been resent by way of the retransmission timer (regarded as a notification of congestion), Standard TCP sets the ssthresh to the maximum of half of the FlightSize and $2 * SMSS$ [\[RFC5681\]](#), which causes the TCP congestion control to go back to allowing only a BDP of packets in flight -- just sufficient to maintain 100% utilisation of the bottleneck on the network path.

AQM mechanisms such as CoDel [\[I-D.CoDel\]](#) and PIE [\[RFC8033\]](#) set a delay target in routers and use congestion notifications to constrain

the queuing delays experienced by packets, rather than in response to impending or actual bottleneck buffer exhaustion. With current default delay targets, CoDel and PIE both effectively emulate a bottleneck with a short queue (section II, [[ABE2017](#)]) while also allowing short traffic bursts into the queue. This provides acceptable performance for TCP connections over a path with a low BDP, or in highly multiplexed scenarios (many concurrent transport flows). However, in a lightly-multiplexed case over a path with a large BDP, conventional TCP backoff leads to gaps in packet transmission and under-utilisation of the path.

Instead of discarding packets, an AQM mechanism is allowed to mark ECN-Capable packets with an ECN CE-mark. The reception of a CE-mark feedback not only indicates congestion on the network path, it also indicates that an AQM mechanism exists at the bottleneck along the path, and hence the CE-mark likely came from a bottleneck with a controlled short queue. Reacting differently to an ECN-signalled congestion than to an inferred packet loss can then yield the benefit of a reduced back-off when queues are short. Using ECN can also be advantageous for several other reasons [[RFC8087](#)].

The idea of reacting differently to inferred packet loss and detection of an ECN-signalled congestion pre-dates this document. For example, previous research proposed using ECN CE-marked feedback to modify TCP congestion control behaviour via a larger multiplicative decrease factor in conjunction with a smaller additive increase factor [[ICC2002](#)]. The goal of this former work was to operate across AQM bottlenecks using Random Early Detection (RED) that were not necessarily configured to emulate a short queue (The current usage of RED as an Internet AQM method is limited [[RFC7567](#)]).

4.2. Focus on ECN as Defined in [RFC3168](#)

Some transport protocol mechanisms rely on ECN semantics that differ from the original ECN definition [[RFC3168](#)] -- for example, Congestion Exposure (ConEx) [[RFC7713](#)] and Datacenter TCP (DCTCP) [[I-D.ietf-tcpm-dctcp](#)] need more accurate ECN information than that offered by the original feedback method. Other mechanisms (e.g., [[I-D.ietf-tcpm-accurate-ecn](#)]) allow the sender to adjust the rate more frequently than once each path RTT. Use of these mechanisms is out of scope for this document.

4.3. Choice of ABE Multiplier

ABE decouples the reaction of a TCP sender to inferred packet loss and ECN-signalled congestion when in the congestion avoidance phase by differentiating the scaling factor used in Equation 4 in [Section 3.1 of \[\[RFC5681\]\(#\)\]](#). The description respectively uses

β_{loss} and β_{ecn} to refer to the multiplicative decrease factors applied in response to inferred packet loss, and in response to a receiver indicating ECN-signalled congestion. For non-ECN-enabled TCP connections, only β_{loss} applies.

In other words, in response to inferred packet loss:

$$\text{ssthresh} = \max(\text{FlightSize} * \beta_{\text{loss}}, 2 * \text{SMSS})$$

and in response to an indication of an ECN-signalled congestion:

$$\text{ssthresh} = \max(\text{FlightSize} * \beta_{\text{ecn}}, 2 * \text{SMSS})$$

and

$$\text{cwnd} = \text{ssthresh}$$

where FlightSize is the amount of outstanding data in the network, upper-bounded by the smaller of the sender's cwnd and the receiver's advertised window (rwnd) [RFC5681]. The higher the values of β_{loss} and β_{ecn} , the less aggressive the response of any individual backoff event.

The appropriate choice for β_{loss} and β_{ecn} values is a balancing act between path utilisation and draining the bottleneck queue. More aggressive backoff (smaller $\beta_{\text{*}}$) risks underutilising the path, while less aggressive backoff (larger $\beta_{\text{*}}$) can result in slower draining of the bottleneck queue.

The Internet has already been running with at least two different β_{loss} values for several years: the standard value is 0.5 [RFC5681], and the Linux implementation of CUBIC [I-D.CUBIC] has used a multiplier of 0.7 since kernel version 2.6.25 released in 2008. ABE proposes no change to β_{loss} used by current TCP implementations.

β_{ecn} depends on how the response of a TCP connection to shallow AQM marking thresholds is optimised. β_{loss} reflects the preferred response of each congestion control algorithm when faced with exhaustion of buffers (of unknown depth) signalled by packet loss. Consequently, for any given TCP congestion control algorithm the choice of β_{ecn} is likely to be algorithm-specific, rather than a constant multiple of the algorithm's existing β_{loss} . The recommended β_{ecn} value in this document is only applicable for Standard TCP congestion control.

A range of tests (section IV, [ABE2017]) with NewReno and CUBIC over CoDel and PIE in lightly-multiplexed scenarios have explored this

choice of parameter. The results of these tests indicate that CUBIC connections benefit from β_{ecn} of 0.85 (cf. $\beta_{\text{loss}} = 0.7$), and NewReno connections see improvements with β_{ecn} in the range 0.7 to 0.85 (cf. $\beta_{\text{loss}} = 0.5$).

5. ABE Requirements

This update is a sender-side only change. Like other changes to congestion control algorithms, it does not require any change to the TCP receiver or to network devices. It does not require any ABE-specific changes in routers or the use of Accurate ECN feedback [[I-D.ietf-tcpm-accurate-ecn](#)] by a receiver.

[RFC3168](#) states that the congestion control response to an ECN-signalled congestion is the same as the response to a dropped packet [[RFC3168](#)]. [[I-D.ECN-exp](#)] updates this specification to allow systems to provide a different behaviour when they experience ECN-signalled congestion rather than packet loss. The present specification defines such an experiment and has thus been assigned an Experimental status before being proposed as a Standards-Track update.

The purpose of the Internet experiment is to collect experience with deployment of ABE, and confirm the safety in deployed networks using this update to TCP congestion control.

When used with bottlenecks that do not support ECN-marking the specification does not modify the transport protocol.

To evaluate the benefit, this experiment therefore requires support in AQM routers for ECN-marking of packets carrying the ECN-Capable Transport, ECT(0), codepoint [[RFC3168](#)].

If the method is only deployed by some senders, and not by others, the senders that use this method can gain some advantage, possibly at the expense of other flows that do not use this updated method. Because this advantage applies only to ECN-marked packets and not to packet loss indications, in the worst case (e.g., an ABE-compliant TCP sender using $\beta_{\text{ecn}} = 1.0$) the ECN-Capable bottleneck will still fall back to dropping packets, and the result is no different than if the TCP sender was using traditional loss-based congestion control.

A TCP sender reacts to loss or ECN marks only once per round-trip time. Hence, if a sender would first be notified of an ECN mark and then learn about loss in the same round-trip, it would only react to the first notification (ECN) but not to the second (loss). [RFC3168](#) specified a reaction to ECN that was equal to the reaction to loss [[RFC3168](#)].

ABE also makes one congestion-response each RTT that congestion is signalled, and therefore there is no response to loss within the same round-trip time, since ABE has already made a reduction of the congestion window. ABE will however respond for each round-trip time that congestion continues to be signaled. This consecutive reduction can protect the network against long-standing unfairness in the case of AQM algorithms that do not keep a small average queue length.

The result of this Internet experiment will include an investigation of cases such as the ones listed above, and be reported by presentation to the TCPM WG (or IESG) or an implementation report at the end of the experiment.

6. Acknowledgements

Authors N. Khademi, M. Welzl and G. Fairhurst were part-funded by the European Community under its Seventh Framework Programme through the Reducing Internet Transport Latency (RITE) project (ICT-317700). The views expressed are solely those of the authors.

The authors would like to thank Stuart Cheshire for many suggestions when revising the draft, and the following people for their contributions to [ABE2017]: Chamil Kulatunga, David Ros, Stein Gjessing, Sebastian Zander. Thanks also to (in alphabetical order) Roland Bless, Bob Briscoe, David Black, Markku Kojo, John Leslie, Lawrence Stewart, Dave Taht and the TCPM working group for providing valuable feedback on this document.

The authors would finally like to thank everyone who provided feedback on the congestion control behaviour specified in this update received from the IRTF Internet Congestion Control Research Group (ICCRG).

7. IANA Considerations

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This document includes no request to IANA.

8. Implementation Status

ABE is implemented as a patch for Linux and FreeBSD. It is meant for research and available for download from <http://heim.ifi.uio.no/naeemk/research/ABE/>. This code was used to produce the test results that are reported in [ABE2017]. An evolved version of the patch for FreeBSD is currently under review for potential inclusion in the mainline kernel [ABE-FreeBSD].

9. Security Considerations

The described method is a sender-side only transport change, and does not change the protocol messages exchanged. The security considerations for ECN [[RFC3168](#)] therefore still apply.

This is a change to TCP congestion control with ECN that will typically lead to a change in the capacity achieved when flows share a network bottleneck. This could result in some flows receiving more than their fair share of capacity. Similar unfairness in the way that capacity is shared is also exhibited by other congestion control mechanisms that have been in use in the Internet for many years (e.g., CUBIC [[I-D.CUBIC](#)]). Unfairness may also be a result of other factors, including the round trip time experienced by a flow. ABE applies only when ECN-marked packets are received, not when packets are lost, hence use of ABE cannot lead to congestion collapse.

10. Revision Information

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-05. Refined the description of the experiment based on feedback at IETF-100. Incorporated comments from David Black.

-04. Incorporates review comments from Lawrence Stewart and the remaining comments from Roland Bless. References are updated.

-03. Several review comments from Roland Bless are addressed. Consistent terminology and equations. Clarification on the scope of recommended β_{ecn} value.

-02. Corrected the equations in [Section 4.3](#). Updated the affiliations. Lower bound for $cwnd$ is defined. A recommendation for window-based transport protocols is changed to cover all transport protocols that implement a congestion control reduction to an ECN congestion signal. Added text about ABE's FreeBSD mainline kernel status including a reference to the FreeBSD code review page. References are updated.

-01. Text improved, mainly incorporating comments from Stuart Cheshire. The reference to a technical report has been updated to a published version of the tests [[ABE2017](#)]. Used "AQM Mechanism" throughout in place of other alternatives, and more consistent use of technical language and clarification on the intended purpose of the experiments required by EXP status. There was no change to the technical content.

-00. [draft-ietf-tcpm-alternativebackoff-ecn-00](#) replaces [draft-khademi-tcpm-alternativebackoff-ecn-01](#). Text describing the nature of the experiment was added.

Individual draft -01. This I-D now refers to [draft-black-tsvwg-ecn-experimentation-02](#), which replaces [draft-khademi-tsvwg-ecn-response-00](#) to make a broader update to [RFC3168](#) for the sake of allowing experiments. As a result, some of the motivating and discussing text that was moved from [draft-khademi-alternativebackoff-ecn-03](#) to [draft-khademi-tsvwg-ecn-response-00](#) has now been re-inserted here.

Individual draft -00. [draft-khademi-tsvwg-ecn-response-00](#) and [draft-khademi-tcpm-alternativebackoff-ecn-00](#) replace [draft-khademi-alternativebackoff-ecn-03](#), following discussion in the TSVWG and TCPM working groups.

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