

TCP Maintenance and Minor Extensions (tcpm)  
Internet-Draft  
Intended status: Experimental  
Expires: September 10, 2015

P. Hurtig  
A. Brunstrom  
Karlstad University  
A. Petlund  
Simula Research Laboratory AS  
M. Welzl  
University of Oslo  
March 9, 2015

TCP and SCTP RTO Restart  
draft-ietf-tcpm-rtorestart-05

## Abstract

This document describes a modified algorithm for managing the TCP and SCTP retransmission timers that provides faster loss recovery when there is a small amount of outstanding data for a connection. The modification, RTO Restart (RTOR), allows the transport to restart its retransmission timer more aggressively in situations where fast retransmit cannot be used. This enables faster loss detection and recovery for connections that are short-lived or application-limited.

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## 1. Introduction

TCP uses two mechanisms to detect segment loss. First, if a segment is not acknowledged within a certain amount of time, a retransmission timeout (RTO) occurs, and the segment is retransmitted [[RFC6298](#)]. While the RTO is based on measured round-trip times (RTTs) between the sender and receiver, it also has a conservative lower bound of 1 second to ensure that delayed segments are not mistaken as lost. Second, when a sender receives dupACKs, the fast retransmit algorithm infers segment loss and triggers a retransmission [[RFC5681](#)]. Duplicate acknowledgments are generated by a receiver when out-of-order segments arrive. As both segment loss and segment reordering cause out-of-order arrival, fast retransmit waits for three dupACKs before considering the segment as lost. In some situations, however, the number of outstanding segments is not enough to trigger three dupACKs, and the sender must rely on lengthy RTOs for loss recovery.

The number of outstanding segments can be small for several reasons:

- (1) The connection is limited by the congestion control when the path has a low total capacity (bandwidth-delay product) or the connection's share of the capacity is small. It is also limited by the congestion control in the first few RTTs of a connection or after an RTO when the available capacity is probed using slow-start.
- (2) The connection is limited by the receiver's available buffer space.
- (3) The connection is limited by the application if the available capacity of the path is not fully utilized (e.g. interactive applications), or at the end of a transfer.

While the reasons listed above are valid for any flow, the third reason is most common for applications that transmit short flows, or use a bursty transmission pattern. A typical example of applications

that produce short flows are web-based applications. [RJ10] shows that 70% of all web objects, found at the top 500 sites, are too small for fast retransmit to work. [FDT13] shows that about 77% of all retransmissions sent by a major web service are sent after RTO expiry. Applications with bursty transmission patterns often send

data in response to actions, or as a reaction to real life events. Typical examples of such applications are stock trading systems, remote computer operations, online games, and web-based applications using persistent connections. What is special about this class of applications is that they often are time-dependant, and extra latency can reduce the application service level [P09].

The RTO Restart (RTOR) mechanism described in this document makes the RTO slightly more aggressive when the number of outstanding segments is too small for fast retransmit to work, in an attempt to enable faster loss recovery for all segments while being robust to reordering. While RTOR still conforms to the requirement in [RFC6298] that segments must not be retransmitted earlier than RTO seconds after their original transmission, it could increase the risk of spurious timeout. Spurious timeouts can degrade the performance of flows with multiple bursts of data, as a burst following a spurious timeout might not fit within the reduced congestion window (cwnd). There are, however, several techniques to mitigate the effects of such unnecessary retransmissions (cf. [RFC4015]).

While this document focuses on TCP, the described changes are also valid for the Stream Control Transmission Protocol (SCTP) [RFC4960] which has similar loss recovery and congestion control algorithms.

## 2. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

This document introduces the following variables:

The number of previously unsent segments (prevunsnt): The number of segments that a sender has queued for transmission, but has not yet sent.

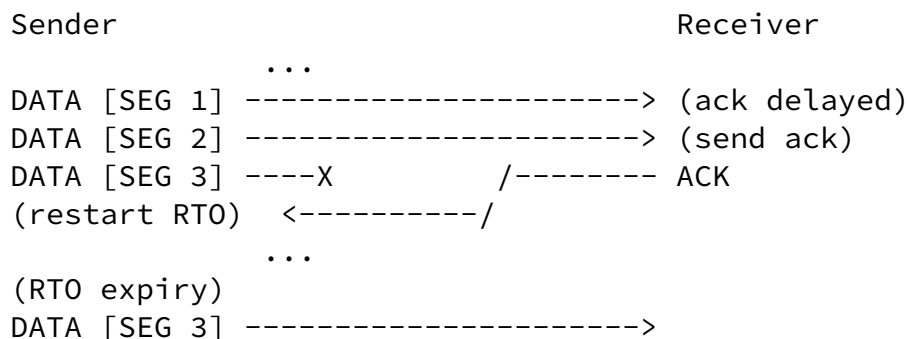
RTO Restart threshold (rrthresh): RTOR is enabled whenever the sum of the number of outstanding and previously unsent segments (prevunsnt) is below this threshold.

### 3. RTO Restart Overview

The RTO management algorithm described in [RFC6298] recommends that the retransmission timer is restarted when an acknowledgment (ACK) that acknowledges new data is received and there is still outstanding data. The restart is conducted to guarantee that unacknowledged segments will be retransmitted after approximately RTO seconds. However, by restarting the timer on each incoming ACK,

retransmissions are not typically triggered RTO seconds after their previous transmission but rather RTO seconds after the last ACK arrived. The duration of this extra delay depends on several factors but is in most cases approximately one RTT. Hence, in most situations, the time before a retransmission is triggered is equal to "RTO + RTT".

The standardized RTO timer management is illustrated in Figure 1 where a TCP sender transmits three segments to a receiver. The arrival of the first and second segment triggers a delayed ACK (delACK) [RFC1122], which restarts the RTO timer at the sender. The RTO is restarted approximately one RTT after the transmission of the third segment. Thus, if the third segment is lost, as indicated in Figure 1, the effective loss detection time is "RTO + RTT" seconds. In some situations, the effective loss detection time becomes even longer. Consider a scenario where only two segments are outstanding. If the second segment is lost, the time to expire the delACK timer will also be included in the effective loss detection time.



## Figure 1: RTO restart example

During normal TCP bulk transfer the current RTO restart approach is not a problem. Actually, as long as enough segments arrive at a receiver to enable fast retransmit, RTO-based loss recovery should be avoided. RTOs should only be used as a last resort, as they drastically lower the congestion window compared to fast retransmit. The current approach can therefore be beneficial -- it is described in [EL04] to act as a "safety margin" that compensates for some of the problems that the authors have identified with the standard RTO calculation. Notably, the authors of [EL04] also state that "this safety margin does not exist for highly interactive applications where often only a single packet is in flight."

Although fast retransmit is preferable there are situations where timeouts are appropriate, or the only choice. For example, if the network is severely congested and no segments arrive RTO-based

recovery should be used. In this situation, the time to recover from the loss(es) will not be the performance bottleneck. However, for connections that do not utilize enough capacity to enable fast retransmit, RTO-based loss detection is the only choice and the time required for this can become a performance bottleneck.

#### 4. RTOR Algorithm

To enable faster loss recovery for connections that are unable to use fast retransmit, RTOR can be used. By resetting the timer to "RTO - T\_earliest", where T\_earliest is the time elapsed since the earliest outstanding segment was transmitted, retransmissions will always occur after exactly RTO seconds. This approach makes the RTO more aggressive than the standardized approach in [RFC6298] but still conforms to the requirement in [RFC6298] that segments must not be retransmitted earlier than RTO seconds after their original transmission.

This document specifies an OPTIONAL sender-only modification to TCP and SCTP which updates step 5.3 in [Section 5 of \[RFC6298\]](#) (and a similar update in [Section 6.3.2 of \[RFC4960\]](#) for SCTP). A sender that implements this method MUST follow the algorithm below:

When an ACK is received that acknowledges new data:

- (1) Set  $T_{\text{earliest}} = 0$ .
- (2) If the sum of the number of outstanding and previously unsent segments ( $\text{prevunsnt}$ ) is less than an RTOR threshold ( $\text{rrthresh}$ ), set  $T_{\text{earliest}}$  to the time elapsed since the earliest outstanding segment was sent.
- (3) Restart the retransmission timer so that it will expire after (for the current value of  $\text{RTO}$ ):
  - (a)  $\text{RTO} - T_{\text{earliest}}$ , if  $\text{RTO} - T_{\text{earliest}} > 0$ .
  - (b)  $\text{RTO}$ , otherwise.

The RECOMMENDED value of  $\text{rrthresh}$  is four, as it will prevent RTOR from being more aggressive and potentially causing RTOs instead of fast retransmits. This update needs TCP implementations to track the time elapsed since the transmission of the earliest outstanding segment ( $T_{\text{earliest}}$ ). As RTOR is only used when the amount of outstanding and previously unsent data is less than  $\text{rrthresh}$  segments, TCP implementations also need to track whether the amount of outstanding and previously unsent data is more, equal, or less than  $\text{rrthresh}$  segments. Although some packet-based TCP

implementations (e.g. Linux TCP) already track both the transmission times of all segments and also the number of outstanding segments, not all implementations do. [Section 5.3](#) describes how to implement segment tracking for a general TCP implementation. To use RTOR, the calculated expiration time MUST be positive (step 3(a) in the list above); this is required to ensure that RTOR does not trigger retransmissions prematurely when previously retransmitted segments are acknowledged.

## [5.](#) Discussion

In this section, we discuss the applicability and a number of issues surrounding RTOR.

### [5.1.](#) Applicability

The currently standardized algorithm has been shown to add at least one RTT to the loss recovery process in TCP [[LS00](#)] and SCTP [[HB11](#)][PBP09]. For applications that have strict timing requirements (e.g. interactive web) rather than throughput requirements, using RTOR could be beneficial because the RTT and also the delACK timer of receivers are often large components of the effective loss recovery time. Measurements in [[HB11](#)] have shown that the total transfer time of a lost segment (including the original transmission time and the loss recovery time) can be reduced by 35% using RTOR. These results match those presented in [[PGH06](#)][PBP09], where RTOR is shown to significantly reduce retransmission latency.

There are also traffic types that do not benefit from RTOR. One example of such traffic is bulk transmission. The reason why bulk traffic does not benefit from RTOR is that such traffic flows mostly have four or more segments outstanding, allowing loss recovery by fast retransmit. However, there is no harm in using RTOR for such traffic as the algorithm only is active when the amount of outstanding and unsent segments are less than rsthresh (default 4).

Given that RTOR is a mostly conservative algorithm, it is suitable for experimentation as a system-wide default for TCP traffic.

## [5.2.](#) Spurious Timeouts

RTOR can in some situations reduce the loss detection time and thereby increase the risk of spurious timeouts. In theory, the retransmission timer has a lower bound of 1 second [[RFC6298](#)], which limits the risk of having spurious timeouts. However, in practice most implementations use a significantly lower value. Initial measurements, show slight increases in the number of spurious timeouts when such lower values are used [[RHB15](#)]. However, further

experiments, in different environments and with different types of traffic, are encouraged to quantify such increases more reliably.

Does a slightly increased risk matter? Generally, spurious timeouts have a negative effect on the network as segments are transmitted needlessly. However, recent experiments do not show a significant increase in network load for a number of realistic scenarios [[RHB15](#)]. Another problem with spurious retransmissions is related to the

performance of TCP/SCTP, as the congestion window is reduced to one segment when timeouts occur [[RFC5681](#)]. This could be a potential problem for applications transmitting multiple bursts of data within a single flow, e.g. web-based HTTP/1.1 and HTTP/2.0 applications. However, results from recent experiments involving persistent web traffic [[RHB15](#)] only revealed a net gain of using RTOR. Other types of flows, e.g. long-lived bulk flows, are not affected as the algorithm is only applied when the amount of outstanding and unsent segments is less than `rrthresh`. Furthermore, short-lived and application-limited flows are typically not affected as they are too short to experience the effect of congestion control or have a transmission rate that is quickly attainable.

While a slight increase in spurious timeouts has been observed using RTOR, it is not clear whether the effects of this increase mandate any future algorithmic changes or not -- especially since most modern operating systems already include mechanisms to detect [[RFC3522](#)][[RFC3708](#)][[RFC5682](#)] and resolve [[RFC4015](#)] possible problems with spurious retransmissions. Further experimentation is needed to determine this and thereby move this specification from experimental to proposed standard. For instance, RTOR has not been evaluated in the context of mobile networks. Mobile networks often incur highly variable RTTs (delay spikes), due to e.g. handovers, and would therefore be a useful scenario for further experimentation.

### [5.3](#). Tracking Outstanding and Previously Unsent Segments

The method of tracking outstanding and previously unsent segments will probably differ depending on the actual TCP implementation. For packet-based TCP implementations, tracking outstanding segments is often straightforward and can be implemented using a simple counter. For byte-based TCP stacks it is a more complex task. [Section 3.2 of \[\[RFC5827\]\(#\)\]](#) outlines a general method of tracking the number of outstanding segments. The same method can be used for RTOR. The implementation will have to track segment boundaries to form an understanding as to how many actual segments have been transmitted, but not acknowledged. This can be done by the sender tracking the boundaries of the `rrthresh` segments on the right side of the current window (which involves tracking `rrthresh + 1` sequence numbers in TCP). This could be done by keeping a circular list of the segment

boundaries, for instance. Cumulative ACKs that do not fall within



this region indicate that at least `rrthresh` segments are outstanding, and therefore RTOR is not enabled. When the outstanding window becomes small enough that RTOR can be invoked, a full understanding of the number of outstanding segments will be available from the `rrthresh + 1` sequence numbers retained. (Note: the implicit sequence number consumed by the TCP FIN bit can also be included in the tracking of segment boundaries.)

Tracking the number of previously unsent segments depends on the segmentation strategy used by the TCP implementation, not whether it is packet-based or byte-based. In the case segments are formed directly on socket writes, the process of determining the number of previously unsent segments should be trivial. In the case that unsent data can be segmented (or re-segmented) as long as it still is unsent, a straightforward strategy could be to divide the amount of unsent data (in bytes) with the SMSS to obtain an estimate. In some cases, such an estimation could be too simplistic, depending on the segmentation strategy of the TCP implementation. However, this estimation is not critical to RTOR. For instance, implementations can use a simplified method by setting `prevunsnt` to `rrthresh` whenever previously unsent data is available, and set `prevunsnt` to zero when no new data is available. This will disable RTOR in the presence of unsent data and only use the number of outstanding segments to enable/disable RTOR. This strategy was used in an earlier version of the algorithm and works well. The addition of tracking `prevunsnt` was only made to optimize a corner case in which RTOR was unnecessarily disabled.

## [6.](#) Related Work

There are several proposals that address the problem of not having enough ACKs for loss recovery. In what follows, we explain why the mechanism described here is complementary to these approaches:

The limited transmit mechanism [[RFC3042](#)] allows a TCP sender to transmit a previously unsent segment for each of the first two dupACKs. By transmitting new segments, the sender attempts to generate additional dupACKs to enable fast retransmit. However, limited transmit does not help if no previously unsent data is ready for transmission. [[RFC5827](#)] specifies an early retransmit algorithm to enable fast loss recovery in such situations. By dynamically lowering the number of dupACKs needed for fast retransmit (`dupthresh`), based on the number of outstanding segments, a smaller number of dupACKs is needed to trigger a retransmission. In some situations, however, the algorithm is of no use or might not work properly. First, if a single segment is outstanding, and lost, it is impossible to use early retransmit. Second, if ACKs are lost, early

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retransmit cannot help. Third, if the network path reorders segments, the algorithm might cause more unnecessary retransmissions than fast retransmit. The recommended value of RTOR's `rrthresh` variable is based on the `dupthresh`, but is possible to adapt to allow tighter integration with other experimental algorithms such as early retransmit.

Tail Loss Probe [TLP] is a proposal to send up to two "probe segments" when a timer fires which is set to a value smaller than the RTO. A "probe segment" is a new segment if new data is available, else a retransmission. The intention is to compensate for sluggish RTO behavior in situations where the RTO greatly exceeds the RTT, which, according to measurements reported in [TLP], is not uncommon. Furthermore, TLP also tries to circumvent the congestion window reset to one segment by instead enabling fast recovery. The Probe timeout (PTO) is normally two RTTs, and a spurious PTO is less risky than a spurious RTO because it would not have the same negative effects (clearing the scoreboard and restarting with slow-start). TLP is a more advanced mechanism than RTOR, requiring e.g. SACK to work, and is often able to reduce loss recovery times more. However, it also increases the amount of spurious retransmissions noticeably, as compared to RTOR [RHB15].

TLP is applicable in situations where RTOR does not apply, and it could overrule (yielding a similar general behavior, but with a lower timeout) RTOR in cases where the number of outstanding segments is smaller than four and no new segments are available for transmission. The PTO has the same inherent problem of restarting the timer on an incoming ACK, and could be combined with a strategy similar to RTOR's to offer more consistent timeouts.

## 7. Acknowledgements

The authors wish to thank Godred Fairhurst, Yuchung Cheng, Mark Allman, Anantha Ramaiah, Richard Scheffenegger, Nicolas Kuhn, Alexander Zimmermann, and Michael Scharf for commenting on the draft and the ideas behind it.

All the authors are supported by RITE (<http://riteproject.eu/>), a research project (ICT-317700) funded by the European Community under its Seventh Framework Program. The views expressed here are those of the author(s) only. The European Commission is not liable for any use that may be made of the information in this document.

## [8.](#) IANA Considerations

This memo includes no request to IANA.

## [9.](#) Security Considerations

This document discusses a change in how to set the retransmission timer's value when restarted. Therefore, the security considerations found in [\[RFC6298\]](#) apply to this document. No additional security problems have been identified with RT0 Restart at this time.

## [10.](#) Changes from Previous Versions

RFC-Editor note: please remove this section prior to publication.

### [10.1.](#) Changes from [draft-ietf-...-04](#) to -05

- o Introduced variable to track the number of previously unsent segments.
- o Clarified many concepts, e.g. extended the description on how to track outstanding and previously unsent segments.
- o Added a reference to initial measurements on the effects of using RTOR.
- o Improved wording throughout the document.

### [10.2.](#) Changes from [draft-ietf-...-03](#) to -04

- o Changed the algorithm to allow RTOR when there is unsent data available, but the cwnd does not allow transmission.
- o Changed the algorithm to not trigger if  $RTOR \leq 0$ .
- o Made minor adjustments throughout the document to adjust for the algorithmic change.
- o Improved the wording throughout the document.

[10.3.](#) Changes from [draft-ietf-...-02](#) to -03

- o Updated the document to use "RTOR" instead of "RTO Restart" when referring to the modified algorithm.
- o Moved document terminology to a section of its own.
- o Introduced the rrthresh variable in the terminology section.

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- o Added a section to generalize the tracking of outstanding segments.
- o Updated the algorithm to work when the number of outstanding segments is less than four and one segment is ready for transmission, by restarting the timer when new data has been sent.
- o Clarified the relationship between fast retransmit and RTOR.
- o Improved the wording throughout the document.

[10.4.](#) Changes from [draft-ietf-...-01](#) to -02

- o Changed the algorithm description in [Section 3](#) to use formal [RFC 2119](#) language.
- o Changed last paragraph of [Section 3](#) to clarify why the RTO restart algorithm is active when less than four segments are outstanding.
- o Added two paragraphs in [Section 4.1](#) to clarify why the algorithm can be turned on for all TCP traffic without having any negative effects on traffic patterns that do not benefit from a modified timer restart.
- o Improved the wording throughout the document.
- o Replaced and updated some references.

[10.5.](#) Changes from [draft-ietf-...-00](#) to -01

- o Improved the wording throughout the document.

- o Removed the possibility for a connection limited by the receiver's advertised window to use RT0 restart, decreasing the risk of spurious retransmission timeouts.
- o Added a section that discusses the applicability of and problems related to the RT0 restart mechanism.
- o Updated the text describing the relationship to TLP to reflect updates made in this draft.
- o Added acknowledgments.

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#### Authors' Addresses

Per Hurtig  
Karlstad University  
Universitetsgatan 2  
Karlstad 651 88  
Sweden

Phone: +46 54 700 23 35  
Email: per.hurtig@kau.se

Anna Brunstrom  
Karlstad University  
Universitetsgatan 2  
Karlstad 651 88  
Sweden

Phone: +46 54 700 17 95  
Email: anna.brunstrom@kau.se

Andreas Petlund  
Simula Research Laboratory AS  
P.O. Box 134  
Lysaker 1325  
Norway

Phone: +47 67 82 82 00  
Email: apetlund@simula.no

Michael Welzl  
University of Oslo  
PO Box 1080 Blindern  
Oslo N-0316  
Norway

Phone: +47 22 85 24 20  
Email: michawe@ifi.uio.no