TLS S. Santesson
Internet-Draft 3xA Security AB
Intended status: Standards Track H. Tschofeniq

Expires: September 29, 2013

Nokia Siemens Networks

March 28, 2013

# Transport Layer Security (TLS) Cached Information Extension draft-ietf-tls-cached-info-14.txt

#### Abstract

Transport Layer Security (TLS) handshakes often include fairly static information, such as the server certificate and a list of trusted Certification Authorities (CAs). This information can be of considerable size, particularly if the server certificate is bundled with a complete certificate path (including all intermediary certificates up to the trust anchor public key).

This document defines an extension that omits the exchange of already available information. The TLS client informs a server of cached information, for example from a previous TLS handshake, allowing the server to omit the already available information.

#### Status of this Memo

This Internet-Draft is submitted in full conformance with the provisions of  $\underline{\mathsf{BCP}}$  78 and  $\underline{\mathsf{BCP}}$  79.

Internet-Drafts are working documents of the Internet Engineering Task Force (IETF). Note that other groups may also distribute working documents as Internet-Drafts. The list of current Internet-Drafts is at <a href="http://datatracker.ietf.org/drafts/current/">http://datatracker.ietf.org/drafts/current/</a>.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

This Internet-Draft will expire on September 29, 2013.

## Copyright Notice

Copyright (c) 2013 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to  $\underline{\mathsf{BCP}\ 78}$  and the IETF Trust's Legal Provisions Relating to IETF Documents

(<a href="http://trustee.ietf.org/license-info">http://trustee.ietf.org/license-info</a>) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Simplified BSD License text as described in Section 4.e of the Trust Legal Provisions and are provided without warranty as described in the Simplified BSD License.

#### Table of Contents

<u>1</u> .	Introduction	. 3
<u>2</u> .	Terminology	. 4
<u>3</u> .	Cached Information Extension	. <u>5</u>
<u>4</u> .	Exchange Specification	. 7
4	1. Omitting the Certificate Chain	. 7
4	2. Omitting the Trusted CAs	. 8
<u>5</u> .	Example	. <u>9</u>
<u>6</u> .	Security Considerations	. 11
<u>7</u> .	IANA Considerations	. 12
7	1. New Entry to the TLS ExtensionType Registry	. 12
7	2. New Registry for CachedInformationType	. 12
<u>8</u> .	Acknowledgments	. 13
<u>9</u> .	References	. 14
9	1. Normative References	. 14
9	2. Informative References	. 14
Aut	ors' Addresses	. 15

#### 1. Introduction

Transport Layer Security (TLS) handshakes often include fairly static information, such as the server certificate and a list of trusted Certification Authorities (CAs). This information can be of considerable size, particularly if the server certificate is bundled with a complete certificate path (including all intermediary certificates up to the trust anchor public key).

Optimizing the exchange of information to a minimum helps to improve performance in environments where devices are connected to a network with characteristics like low bandwidth, high latency and high loss rate. These types of networks exist, for example, when smart objects are connected using a low power IEEE 802.15.4 radio. For more information about the challenges with smart object deployments please see [RFC6574].

This specification defines a TLS extension that allows a client and a server to exclude transmission of cached information from the TLS handshake.

A typical example exchange may therefore look as follows. First, the client and the server executes the usual TLS handshake. The client may, for example, decide to cache the certificate provided by the server. When the TLS client connects to the TLS server some time in the future, without using session resumption, it then attaches the cached information extension defined in this document to the client hello message to indicate that it had cached the certificate, and it provides the fingerprint of it. If the server's certificate had not changed then the TLS server does not need to send the full certificate to the client again. In case the information had changed, the certificate payload is transmitted to the client to allow the client to update it's state information.

## 2. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "MUST", "MUST NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

#### 3. Cached Information Extension

This document defines a new extension type (cached information(TBD)), which is used in client hello and server hello messages. The extension type is specified as follows.

```
enum {
     cached_information(TBD), (65535)
} ExtensionType;
```

The extension\_data field of this extension, when included in the client hello, MUST contain the CachedInformation structure.

```
enum {
     certificate_chain(1), trusted_cas(2) (255)
} CachedInformationType;
struct {
     CachedInformationType type;
     HashAlgorithm hash;
     opaque hash_value<1..255>;
} CachedObject;
struct {
     CachedObject cached_info<1..2^16-1>;
} CachedInformation;
```

When the CachedInformationType identifies a certificate\_chain, then the hash value field MUST include the hash calculated over the certificate\_list element of the Certificate payload provided by the TLS server in an earlier exchange, excluding the three length bytes of the certificate\_list vector.

When the CachedInformationType identifies a trusted\_cas, then the hash\_value MUST include a hash calculated over the certificate\_authorities element of the CertificateRequest payload provided by the TLS server in an earlier exchange, excluding the two length bytes of the certificate\_authorities vector.

The hash algorithm used to calculate hash values is conveyed in the 'hash' field of the CachedObject element. The list of registered hash algorithms can be found in the TLS HashAlgorithm Registry, which was created by RFC 5246 [RFC5246]. The value zero (0) for 'none' is not an allowed choice for a hash algorithm and MUST NOT be used.

This document establishes a registry for CachedInformationType types and additional values can be added following the policy described in Section 7.

#### 4. Exchange Specification

Clients supporting this extension MAY include the "cached\_information" extension in the (extended) client hello, which MAY contain zero or more CachedObject attributes.

A server supporting this extension MAY include the "cached\_information" extension in the (extended) server hello, which MAY contain one or more CachedObject attributes it supports. By returning the "cached information" extension the server indicates that it supports caching of each present CachedObject that matches the specified hash value. The server MAY support other cached objects that are not present in the extension.

Note: If clients make use of the Server Name Indication [RFC6066] then clients may need to cache multiple data items for a single server since servers may host multiple 'virtual' servers at a single underlying network address.

Following a successful exchange of the "cached\_information" extensions in the client and server hello, the server omits sending the corresponding handshake message. How information is omitted from the handshake message is defined per cached info type. Section 4.1 and Section 4.2 defines the syntax of the fingerprinted information.

The handshake protocol MUST proceed using the information as if it was provided in the handshake protocol. The Finished message MUST be calculated over the actual data exchanged in the handshake protocol. That is, the Finished message will be calculated over the information that was omitted from transmission by means of its present hash in the client hello and not through its presence in the handshake exchange.

The server MUST NOT include more than one fingerprint for a single information element, i.e., at maximum only one CachedObject structure per replaced information is provided.

## 4.1. Omitting the Certificate Chain

When an object of type 'certificate\_chain' is provided in the client hello, the server MAY replace the sequence of certificates with an empty sequence with an actual length field of zero (=empty vector).

The original handshake message syntax is defined in RFC 5246 [RFC5246] and has the following structure:

```
opaque ASN.1Cert<1..2^24-1>;
struct {
    ASN.1Cert certificate_list<0..2^24-1>;
} Certificate;
```

Note that [I-D.ietf-tls-oob-pubkey] allows the certificate payload to contain only the SubjectPublicKeyInfo instead of the full information typically found in a certificate. Hence, when this specification is used in combination with [I-D.ietf-tls-oob-pubkey] and the negotiated certificate type is a raw public key then the TLS server omits sending a Certificate payload that contains an ASN.1Cert structure of the SubjectPublicKeyInfo.

## 4.2. Omitting the Trusted CAs

When a fingerprint for an object of type 'trusted\_cas' is provided in the client hello, the server MAY send a DistinguishedName in the Certificate Request message with an actual length field of zero (=empty vector).

The original handshake message syntax is defined in <a href="RFC 5246">RFC 5246</a>] and has the following structure:

```
opaque DistinguishedName<1..2^16-1>;
struct {
    ClientCertificateType certificate_types<1..2^8-1>;
    SignatureAndHashAlgorithm
        supported_signature_algorithms<2^16-1>;
    DistinguishedName certificate_authorities<0..2^16-1>;
} CertificateRequest;
```

## 5. Example

Figure 1 illustrates an example exchange using the TLS cached info extension. In the normal TLS handshake exchange shown in flow (A) the TLS server provides its certificate in the Certificate payload to the client, see step [1]. This allows the client to store the certificate for future use. After some time the TLS client again interacts with the same TLS server and makes use of the TLS cached info extension, as shown in flow (B). The TLS client indicates support for this specification via the cached\_information extension, see [2], and indicates that it has stored the certificate\_chain from the earlier exchange. With [3] the TLS server indicates that it also supports this specification and informs the client that it also supports caching of other objects beyond the 'certificate\_chain', namely 'trusted\_cas' (also defined in this document), and the 'foobar' extension (i.e., an imaginary extension that yet needs to be defined). With [4] the TLS server omits sending the certificate chain, as described in <u>Section 4.1</u>.

```
(A) Initial (full) Exchange
client_hello ->
                      <- server_hello,
                          certificate, // [1]
                          server_key_exchange,
                          server_hello_done
client_key_exchange,
change_cipher_spec,
finished
                         ->
                      <- change_cipher_spec,
                         finished
Application Data
                      <----> Application Data
(B) TLS Cached Extension Usage
client_hello,
cached_information=(certificate_chain) -> // [2]
                      <- server_hello,
                          cached_information= // [3]
                          (certificate_chain, trusted_cas, foo-bar)
                          certificate, // [4]
                          server_key_exchange,
                          server_hello_done
client_key_exchange,
change_cipher_spec,
finished
                         ->
                      <- change_cipher_spec,
                         finished
Application Data
                      <--->
                                     Application Data
```

Figure 1: Example Message Exchange

## **6**. Security Considerations

This specification defines a mechanism to reference stored state using a fingerprint. Sending a fingerprint of cached information in an unencrypted handshake, as the client and server hello is, may allow an attacker or observer to correlate independent TLS exchanges. While some information elements used in this specification, such as server certificates, are public objects and usually not sensitive in this regard, others may be. Those who implement and deploy this specification should therefore make an informed decision whether the cached information is inline with their security and privacy goals. In case of concerns, it is advised to avoid sending the fingerprint of the data objects in clear.

The hash algorithm used in this specification is required to have reasonable random properties in order to provide reasonably unique identifiers. There is no requirement that this hash algorithm must have strong collision resistance.

#### 7. IANA Considerations

#### 7.1. New Entry to the TLS ExtensionType Registry

IANA is requested to add an entry to the existing TLS ExtensionType registry, defined in  $\frac{RFC}{5246}$  [ $\frac{RFC5246}{6}$ ], for cached\_information(TBD) defined in this document.

## 7.2. New Registry for CachedInformationType

IANA is requested to establish a registry for TLS CachedInformationType values. The first entries in the registry are

- o certificate\_chain(1)
- o trusted\_cas(2)

The policy for adding new values to this registry, following the terminology defined in  $\underline{\mathsf{RFC}}\ 5226\ [\underline{\mathsf{RFC5226}}]$ , is as follows:

- o 0-63 (decimal): Standards Action
- o 64-223 (decimal): Specification Required
- o 224-255 (decimal): reserved for Private Use

## 8. Acknowledgments

We would like to thank the following persons for your detailed document reviews:

- o Paul Wouters and Nikos Mavrogiannopoulos (December 2011)
- o Rob Stradling (February 2012)
- o Ondrej Mikle (in March 2012)

Additionally, we would like to thank the TLS working group chairs, Eric Rescorla and Joe Salowey, as well as the security area directors, Sean Turner and Stephen Farrell, for their feedback and support.

## 9. References

#### 9.1. Normative References

- [RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", <u>BCP 14</u>, <u>RFC 2119</u>, March 1997.
- [RFC3874] Housley, R., "A 224-bit One-way Hash Function: SHA-224", RFC 3874, September 2004.
- [RFC5246] Dierks, T. and E. Rescorla, "The Transport Layer Security (TLS) Protocol Version 1.2", <u>RFC 5246</u>, August 2008.
- [RFC6066] Eastlake, D., "Transport Layer Security (TLS) Extensions: Extension Definitions", <u>RFC 6066</u>, January 2011.

#### 9.2. Informative References

- [I-D.ietf-tls-oob-pubkey]
  - Wouters, P., Tschofenig, H., Gilmore, J., Weiler, S., and T. Kivinen, "Out-of-Band Public Key Validation for Transport Layer Security (TLS)", <a href="mailto:draft-ietf-tls-oob-pubkey-07">draft-ietf-tls-oob-pubkey-07</a> (work in progress), February 2013.
- [RFC5226] Narten, T. and H. Alvestrand, "Guidelines for Writing an IANA Considerations Section in RFCs", <u>BCP 26</u>, <u>RFC 5226</u>, May 2008.
- [RFC6574] Tschofenig, H. and J. Arkko, "Report from the Smart Object Workshop", <u>RFC 6574</u>, April 2012.

# Authors' Addresses

Stefan Santesson 3xA Security AB Scheelev. 17 Lund 223 70 Sweden

Email: sts@aaa-sec.com

Hannes Tschofenig Nokia Siemens Networks Linnoitustie 6 Espoo 02600 Finland

Phone: +358 (50) 4871445

Email: Hannes.Tschofenig@gmx.net

URI: http://www.tschofenig.priv.at