

## AES Ciphersuites for TLS

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### Overview

At present, the symmetric ciphers supported by TLS are RC2, RC4, IDEA, DES and triple DES. The protocol would be enhanced by the addition of AES [[AES](#)] ciphersuites, for the following reasons:

1. RC2, RC4 and IDEA are all subject to intellectual property claims. RSA Security Inc has trademark rights in the names RC2 and RC4, and claims that the RC4 algorithm itself is a trade secret. Ascom Systec Ltd owns a patent on the IDEA algorithm.
2. Triple DES is much less efficient than more modern ciphers.
3. Now the AES process is completed there will be commercial pressure to use the selected cipher. The AES is efficient and has withstood extensive cryptanalytic efforts. The AES is

therefore a desirable choice.

4. Currently the DHE ciphersuites only allow triple DES (along with some ``export'' variants which offer reduced key lengths). At the same time the DHE ciphersuites are the only ones to offer forward secrecy.

This document proposes several new ciphersuites, with the aim of overcoming these problems.

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC 2119](#).

## Cipher Usage

The new ciphersuites proposed here are very similar to the following, defined in [\[TLS\]](#):

```
TLS_RSA_WITH_3DES_EDE_CBC_SHA
TLS_DH_DSS_WITH_3DES_EDE_CBC_SHA
TLS_DH_RSA_WITH_3DES_EDE_CBC_SHA
TLS_DHE_DSS_WITH_3DES_EDE_CBC_SHA
TLS_DHE_RSA_WITH_3DES_EDE_CBC_SHA
TLS_DH_anon_WITH_3DES_EDE_CBC_SHA
```

All the ciphersuites described here use the AES in cipher block chaining (CBC) mode. Furthermore, they use SHA-1 [\[SHA-1\]](#) in an HMAC construction as described in section 5 of [\[TLS\]](#). (Although the TLS ciphersuite names include the text ``SHA'', this actually refers to the modified SHA-1 version of the algorithm.)

The ciphersuites differ in the type of certificate and key exchange method. The ciphersuites defined here use the following options for this part of the protocol:

CipherSuite	Certificate type (if applicable) and key exchange algorithm
TLS_RSA_WITH_AES_128_CBC_SHA	RSA
TLS_DH_DSS_WITH_AES_128_CBC_SHA	DH_DSS
TLS_DH_RSA_WITH_AES_128_CBC_SHA	DH_RSA
TLS_DHE_DSS_WITH_AES_128_CBC_SHA	DHE_DSS
TLS_DHE_RSA_WITH_AES_128_CBC_SHA	DHE_RSA
TLS_DH_anon_WITH_AES_128_CBC_SHA	DH_anon

For the meanings of the terms RSA, DH\_DSS, DH\_RSA, DHE\_DSS, DHE\_RSA and DH\_anon, please refer to sections [7.4.2](#) and [7.4.3](#) of [\[TLS\]](#).



The AES supports key lengths of 128, 192 and 256 bits. At the present time, all of these are believed to be secure against even the best equipped attackers. The overall strength of TLS is such that there is no gain from using a key length longer than 128 bits. Accordingly the AES will use 128 bit keys.

The new ciphersuites will have the following definitions:

```
CipherSuite TLS_RSA_WITH_AES_128_CBC_SHA      = { 0x00, 0x2F };
CipherSuite TLS_DH_DSS_WITH_AES_128_CBC_SHA   = { 0x00, 0x30 };
CipherSuite TLS_DH_RSA_WITH_AES_128_CBC_SHA   = { 0x00, 0x31 };
CipherSuite TLS_DHE_DSS_WITH_AES_128_CBC_SHA  = { 0x00, 0x32 };
CipherSuite TLS_DHE_RSA_WITH_AES_128_CBC_SHA  = { 0x00, 0x33 };
CipherSuite TLS_DH_anon_WITH_AES_128_CBC_SHA  = { 0x00, 0x34 };
```

In the absence of an application profile standard specifying otherwise:

1. Servers MUST provide at least one of  
TLS\_RSA\_WITH\_AES\_128\_CBC\_SHA and  
TLS\_DHE\_RSA\_WITH\_AES\_128\_CBC\_SHA.
2. Clients MUST provide both TLS\_RSA\_WITH\_AES\_128\_CBC\_SHA and  
TLS\_DHE\_RSA\_WITH\_AES\_128\_CBC\_SHA.

(A TLS implementation which does not follow this requirement is non-compliant with this RFC. However, it will still be a valid TLS implementation if it complies with [\[TLS\]](#).)

Implementations MAY provide any of the other ciphersuites described above.

## Security Considerations

It is not believed that the new ciphersuites are ever less secure than the corresponding older ones. The AES is believed to be secure, and it has withstood extensive cryptanalytic attack.

The ephemeral Diffie-Hellman ciphersuites provide forward secrecy without any known reduction in security in other areas. To obtain the maximum benefit from these ciphersuites:

1. The ephemeral keys should only be used once. With the TLS protocol as currently defined there is no efficiency gain from reusing ephemeral keys.
2. Ephemeral keys should be destroyed securely when they are no longer required.



3. The random number generator used to create ephemeral keys must not reveal past output even when its internal state is compromised.

[TLS] describes the anonymous Diffie-Hellman (ADH) ciphersuites as deprecated. The ADH ciphersuite defined here is not deprecated. However, when it is used, particular care must be taken:

1. ADH provides confidentiality but not authentication. This means that (if authentication is required) the communicating parties must authenticate to each other by some means other than TLS.
2. ADH is vulnerable to man-in-the-middle attacks, as a consequence of the lack of authentication. The parties must have a way of determining whether they are participating in the same TLS connection. If they are not, they can deduce that they are under attack, and presumably abort the connection.

For example, if the parties share a secret, it is possible to compute a MAC of the TLS Finished message. An attacker would have to negotiate two different TLS connections; one with each communicating party. The Finished messages would be different in each case, because they depend on the master secret. For this reason, the MACs computed by each party would be different.

It is important to note that authentication techniques which do not use the Finished message do not usually provide protection from this attack. For example, the client could authenticate to the server with a password, but it would still be vulnerable to man-in-the-middle attacks.

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During the development of the AES, NIST published the following statement on intellectual property:

### SPECIAL NOTE - Intellectual Property

NIST reminds all interested parties that the adoption of AES is being conducted as an open standards-setting activity. Specifically, NIST has requested that all interested parties identify to NIST any patents or inventions that may be required for the use of AES. NIST hereby gives public notice that it may seek redress under the antitrust laws of the United States against any party in the future who might seek to exercise patent rights





against any user of AES that have not been disclosed to NIST in response to this request for information.

One of the authors of Rijndael signed the following disclaimer when submitting the algorithm to NIST for consideration in the AES process:

I, Joan Daemen, do hereby declare that to the best of my knowledge the practice of the algorithm, reference implementation, and mathematically optimized implementations, I have submitted, known as Rijndael may be covered by the following U.S. and/or foreign patents:

none

I do hereby declare that I am aware of no patent applications which may cover the practice of my submitted algorithm, reference implementation or mathematically optimized implementations.

I do hereby understand that my submitted algorithm may not be selected for inclusion in the Advanced Encryption Standard. I also understand and agree that after the close of the submission period, my submission may not be withdrawn from public consideration for inclusion in the Federal Information Processing Standard (FIPS) for Advanced Encryption Standard (AES). I further understand that I will not receive financial compensation from the government for my submission. I certify that, to the best of my knowledge, I have fully disclosed all patents and patent applications relating to my algorithm. I also understand that the U.S. Government may, during the course of the lifetime of the AES or during the FIPS public review process, modify the algorithm's specifications (e.g., to protect against a newly discovered vulnerability). Should my submission be selected for inclusion in the AES, I hereby agree not to place any restrictions on the use of the algorithm intending it to be available on a worldwide, non-exclusive, royalty-free basis.

I do hereby agree to provide the statements for any patent or patent application identified to cover practice of my algorithm, reference implementation or mathematically optimized implementations and the right to use such implementations for the purposes of the AES evaluation process.



I understand that NIST will announce the selected algorithm(s) and proceed to publish the draft FIPS for public comment. If my algorithm (or the derived algorithm) is not selected for inclusion in the FIPS (including those not selected for second round of public evaluation), I understand that all rights, including use rights of the reference and mathematically optimized implementations, revert back to the submitter (and other owner[s] as appropriate). Additionally, should the U.S. Government not select my algorithm for inclusion in the AES after a period of four years from the close of the submission date for candidate algorithms, all rights revert to the submitter (and other owner[s] as appropriate).

[signed]

Title: Cryptographer  
Dated: 10-6-98  
Place: Brussels

The following disclaimer was signed at the start of the second "round" of the AES process:

Dear Mr Foti [of NIST],

Hereby we confirm that the original patent and patent application information, as provided to NIST with our original submission in June 1998, has not changed. To the best of our knowledge, there are no patents or patent applications covering the practice of the algorithm, reference implementation or the mathematically optimized implementations.

[signed]

Joan Daemen, Vincent Rijmen

#### Acknowledgements

I would like to thank the ietf-tls mailing list contributors who have made helpful suggestions for this document.

#### References

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