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Return Routability Check for DTLS 1.2 and DTLS 1.3
draft-ietf-tls-dtls-rrc-01

Abstract

This document specifies a return routability check for use in context of the Connection ID (CID) construct for the Datagram Transport Layer Security (DTLS) protocol versions 1.2 and 1.3.

Discussion Venues

This note is to be removed before publishing as an RFC.

Discussion of this document takes place on the Transport Layer Security Working Group mailing list (tls@ietf.org), which is archived at <https://mailarchive.ietf.org/arch/browse/tls/>.

Source for this draft and an issue tracker can be found at <https://github.com/tlswg/dtls-rrc>.

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Internet-Draft

DTLS Return Routability Check

October 2021

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[1.](#) Introduction

In "classical" DTLS, selecting a security context of an incoming DTLS record is accomplished with the help of the 5-tuple, i.e. source IP address, source port, transport protocol, destination IP address, and destination port. Changes to this 5 tuple can happen for a variety reasons over the lifetime of the DTLS session. In the IoT context,

NAT rebinding is common with sleepy devices. Other examples include end host mobility and multi-homing. Without CID, if the source IP address and/or source port changes during the lifetime of an ongoing DTLS session then the receiver will be unable to locate the correct security context. As a result, the DTLS handshake has to be re-run.

Of course, it is not necessary to re-run the full handshake if session resumption is supported and negotiated.

A CID is an identifier carried in the record layer header of a DTLS datagram that gives the receiver additional information for selecting the appropriate security context. The CID mechanism has been specified in [[I-D.ietf-tls-dtls-connection-id](#)] for DTLS 1.2 and in [[I-D.ietf-tls-dtls13](#)] for DTLS 1.3.

Section 6 of [[I-D.ietf-tls-dtls-connection-id](#)] describes how the use of CID increases the attack surface by providing both on-path and off-path attackers an opportunity for (D)DoS. It then goes on describing the steps a DTLS principal must take when a record with a CID is received that has a source address (and/or port) different from the one currently associated with the DTLS connection. However, the actual mechanism for ensuring that the new peer address is willing to receive and process DTLS records is left open. This document standardizes a return routability check (RRC) as part of the DTLS protocol itself.

The return routability check is performed by the receiving peer before the CID-to-IP address/port binding is updated in that peer's session state database. This is done in order to provide more confidence to the receiving peer that the sending peer is reachable at the indicated address and port.

2. Conventions and Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [BCP 14](#) [[RFC2119](#)] [[RFC8174](#)] when, and only when, they appear in all capitals, as shown here.

This document assumes familiarity with the CID format and protocol defined for DTLS 1.2 [[I-D.ietf-tls-dtls-connection-id](#)] and for DTLS

1.3 [[I-D.ietf-tls-dtls13](#)]. The presentation language used in this document is described in [Section 4 of \[RFC8446\]](#).

3. RRC Extension

This specification uses the `tls_flags` extension defined in [[I-D.ietf-tls-tlsflags](#)] to allow a client and a server to negotiate support for this extension.

The RRC flag is assigned the value (TBD1) and is used in the ClientHello (CH) and the ServerHello (SH).

4. The Return Routability Check Message

When a record with CID is received that has the source address of the enclosing UDP datagram different from the one previously associated with that CID, the receiver MUST NOT update its view of the peer's IP address and port number with the source specified in the UDP datagram before cryptographically validating the enclosed record(s) but instead perform a return routability check.

```
enum {
    invalid(0),
    change_cipher_spec(20),
    alert(21),
    handshake(22),
    application_data(23),
    heartbeat(24), /* RFC 6520 */
    return_routability_check(TBD), /* NEW */
    (255)
} ContentType;
```

```
uint64 Cookie;
```

```
enum {
    path_challenge(0),
    path_response(1),
    reserved(2..255)
} rrc_msg_type;
```

```
struct {
```

```
    rrc_msg_type msg_type;
    select (return_routability_check.msg_type) {
        case path_challenge: Cookie;
        case path_response:  Cookie;
    };
} return_routability_check;
```

The newly introduced `return_routability_check` message contains a cookie. The cookie is a 8-byte field containing arbitrary data.

The `return_routability_check` message **MUST** be authenticated and encrypted using the currently active security context.

The receiver that observes the peer's address and or port update **MUST** stop sending any buffered application data (or limit the data sent to a TBD threshold) and initiate the return routability check that proceeds as follows:

1. A cookie is placed in a `return_routability_check` message of type `path_challenge`;
2. The message is sent to the observed new address and a timeout `T` is started;
3. The peer endpoint, after successfully verifying the received `return_routability_check` message echoes the cookie value in a `return_routability_check` message of type `path_response`;
4. When the initiator receives and verifies the `return_routability_check` message contains the sent cookie, it updates the peer address binding;
5. If `T` expires, or the address confirmation fails, the peer address binding is not updated.

After this point, any pending send operation is resumed to the bound peer address.

[5. Example](#)

The example TLS 1.3 handshake shown in Figure 1 shows a client and a server negotiating the support for CID and for the RRC extension.

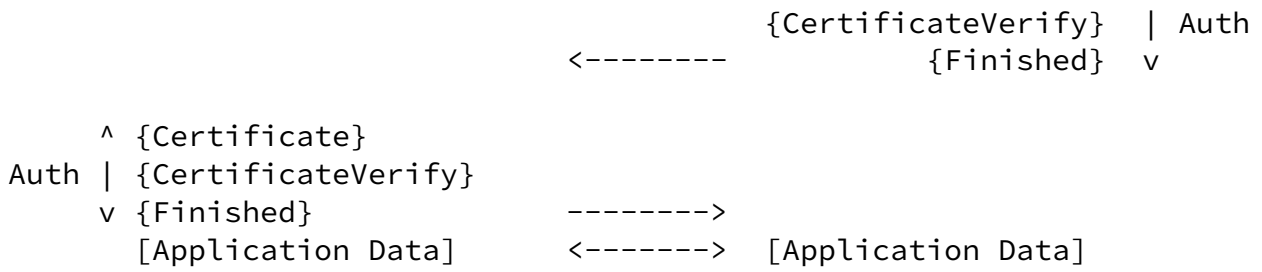
Client

Server

Key ^ ClientHello
Exch | + key_share
| + signature_algorithms
| + tls_flags (RRC)
v + connection_id=empty

----->

ServerHello ^ Key
+ key_share | Exch
+ connection_id=100 |
+ tls_flags (RRC) v
{EncryptedExtensions} ^ Server
{CertificateRequest} v Params
{Certificate} ^



- + Indicates noteworthy extensions sent in the previously noted message.
- * Indicates optional or situation-dependent messages/extensions that are not always sent.
- { } Indicates messages protected using keys derived from a [sender]_handshake_traffic_secret.
- [] Indicates messages protected using keys derived from [sender]_application_traffic_secret_N.

Figure 1: Message Flow for Full TLS Handshake

Once a connection has been established the client and the server exchange application payloads protected by DTLS with an unilaterally used CIDs. In our case, the client is requested to use CID 100 for records sent to the server.

At some point in the communication interaction the IP address used by the client changes and, thanks to the CID usage, the security context to interpret the record is successfully located by the server. However, the server wants to test the reachability of the client at his new IP address.

Client

Server

```

Application Data          =====>
<CID=100>
Src-IP=A
Dst-IP=Z

```

```

<===== Application Data

```

Src-IP=Z
Dst-IP=A

```
<<----->>  
<<  Some  >>  
<<  Time  >>  
<<  Later  >>  
<<----->>
```

```
Application Data      =====>  
<CID=100>  
Src-IP=B  
Dst-IP=Z
```

```
<<< Unverified IP  
Address B >>
```

```
<----- Return Routability Check  
path_challenge(cookie)  
Src-IP=Z  
Dst-IP=B
```

```
Return Routability Check  ----->  
path_response(cookie)  
Src-IP=B  
Dst-IP=Z
```

```
<<< IP Address B  
Verified >>
```

```
<===== Application Data  
Src-IP=Z  
Dst-IP=B
```

Figure 2: Return Routability Example

Note that the return routability checks do not protect against flooding of third-parties if the attacker is on-path, as the attacker can redirect the return routability checks to the real peer (even if those datagrams are cryptographically authenticated). On-path adversaries can, in general, pose a harm to connectivity.

7. IANA Considerations

IANA is requested to allocate an entry to the TLS "ContentType" registry, for the `return_routability_check(TBD)` defined in this document.

IANA is requested to allocate an entry to the TLS Flags registry in the `tls_flags` type:

- * Value: [[IANA please assign a value from the 32-63 value range.]]
- * Flag Name: RRC
- * Message: CH,SH
- * Recommended: Y
- * Reference: [[This document]]

8. Open Issues

Issues against this document are tracked at <https://github.com/tlswg/dtls-rrc/issues>

9. Acknowledgments

We would like to thank Achim Kraus, Hanno Becker, Hanno Boeck, Manuel Pegourie-Gonnard, Mohit Sahni and Rich Salz for their input to this document.

10. Normative References

[I-D.ietf-tls-dtls-connection-id]

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[I-D.ietf-tls-dtls13]

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[RFC8174] Leiba, B., "Ambiguity of Uppercase vs Lowercase in [RFC 2119](#) Key Words", [BCP 14](#), [RFC 8174](#), DOI 10.17487/RFC8174, May 2017, <<https://www.rfc-editor.org/rfc/rfc8174>>.

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[Appendix A](#). History

RFC EDITOR: PLEASE REMOVE THE THIS SECTION

[draft-ietf-tls-dtls-rrc-01](#)

- * Use the TLS flags extension for negotiating RRC
- * Enhanced IANA consideration section
- * Expanded example section
- * Revamp message layout:
 - Use 8-byte fixed size cookies
 - Explicitly separate path challenge from response

[draft-ietf-tls-dtls-rrc-00](#)

- * Draft name changed after WG adoption

[draft-tschofenig-tls-dtls-rrc-01](#)

- * Removed text that overlapped with [draft-ietf-tls-dtls-connection-id](#)

[draft-tschofenig-tls-dtls-rrc-00](#)

- * Initial version

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